

Let a continuous family of symmetric matrices be given:

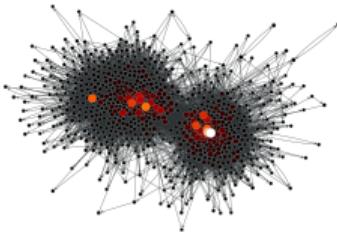
$$\begin{pmatrix} a_{11}(t) & \cdots & a_{1n}(t) \\ \vdots & & \vdots \\ a_{n1}(t) & \cdots & a_{nn}(t) \end{pmatrix}$$

Then for every parameter value t , classically there is

- ▶ a full list of eigenvalues $\lambda_1(t), \dots, \lambda_n(t)$ and
- ▶ an eigenvector basis $(v_1(t), \dots, v_n(t))$.



Can locally the functions λ_i be chosen to be continuous?
How about the v_i ?



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Can locally the functions λ_i be chosen to be continuous? **Yes.**
How about the v_i ? **No.**

– *an invitation* –

New modal operators for constructive mathematics

Type Theory, Constructive Mathematics and Geometric Logic

CIRM
May 2nd, 2023

Ingo Blechschmidt
j.w.w. Alexander Oldenziel

Questions

- 1 Why has the inductive revolution been so powerful?
- 2 Why do proofs using Zorn's maximal ideals work so well in constructive algebra?
- 3 Why are elements of $\bigcap_{\mathfrak{p}} \mathfrak{p}$ not necessarily nilpotent?
- 4 How can we extract computational content from classical proofs?

Infinite data

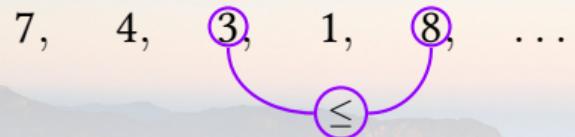
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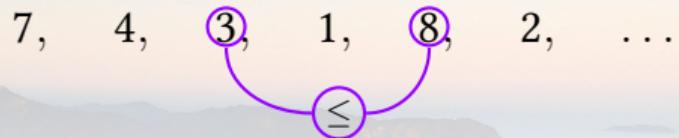
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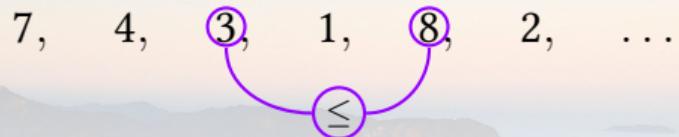


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Thm. Every sequence $\alpha : \mathbb{N} \rightarrow \mathbb{N}$ is **good** in that there exist $i < j$ with $\alpha(i) \leq \alpha(j)$.

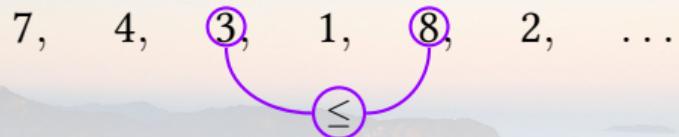
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Proof. (offensive?) By LEM, there is a minimum $\alpha(i)$. Set $j := i + 1$. □

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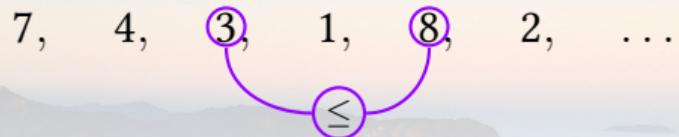
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Def. A preorder X is **well** iff every sequence $\mathbb{N} \rightarrow X$ is good.

Examples. (\mathbb{N}, \leq) , $X \times Y$, X^* , $\text{Tree}(X)$.

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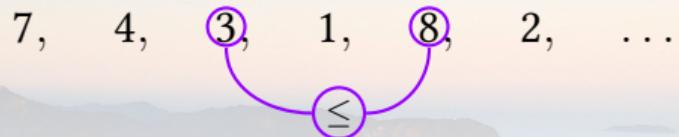
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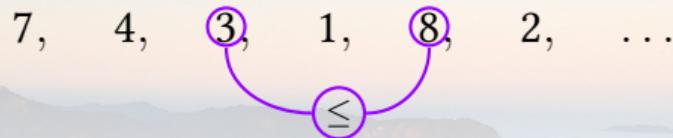
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which might not have enough.*

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Def. For a predicate P on finite lists over a set X , inductively define:

$$\frac{P\sigma}{P \mid \sigma} \quad \frac{\forall(x \in X). \ P \mid \sigma x}{P \mid \sigma}$$

Def. A preorder is **well** iff $\text{Good} \mid []$, where $\text{Good } \sigma \equiv (\exists(i < j). \sigma[i] \leq \sigma[j])$.

Computational content from classical proofs

Def. A transitive relation ($<$) on a set X is ...

- 1 **well-founded*** iff there is no **infinite chain** $x_0 > x_1 > \dots$,
- 2 **well-founded** iff for every $x \in X$, $\text{Acc}(x)$,

where Acc is inductively defined by:

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Prop. Let (X, \leq) be preorder. Let “ $x < y$ ” mean $x \leq y \wedge \neg(y \leq x)$.
Then: If X is well*, then $(<)$ is well-founded*.

Proof. An infinite strictly descending chain would also be good. □

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Can we extract a constructive proof that well preorders are well-founded?

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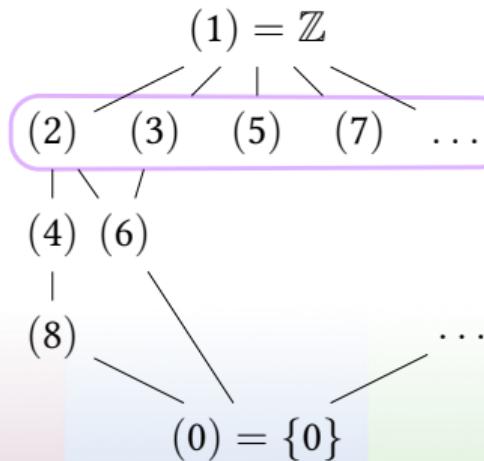
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maximal among the proper ideals

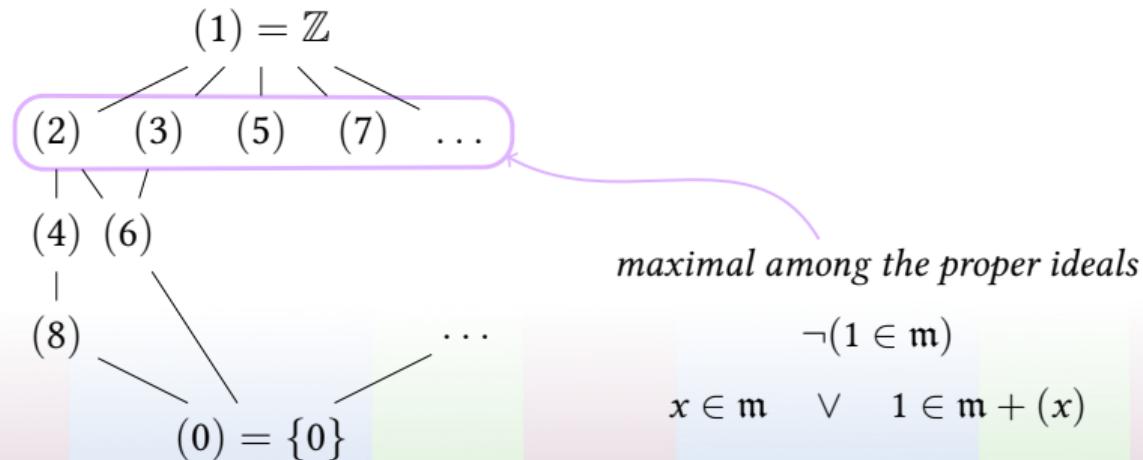
$$\neg(1 \in \mathfrak{m})$$

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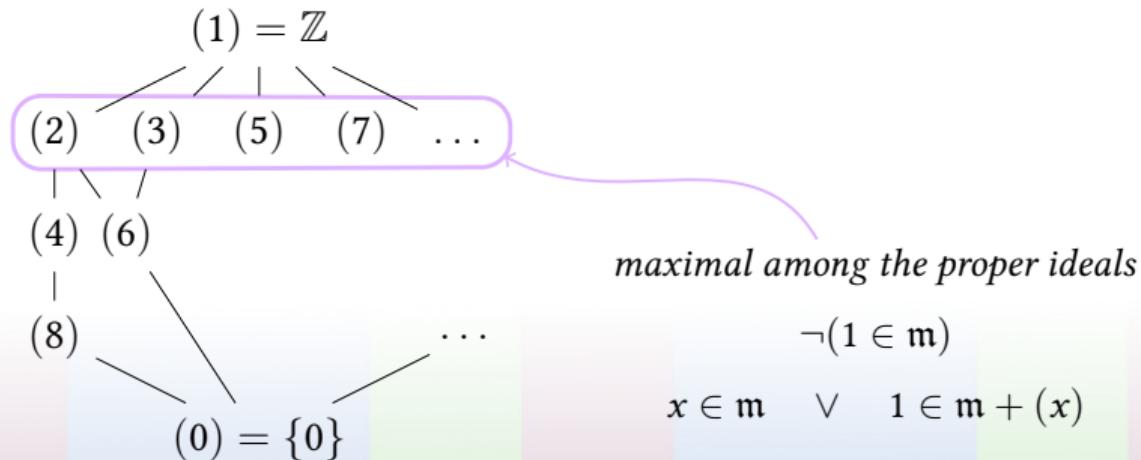
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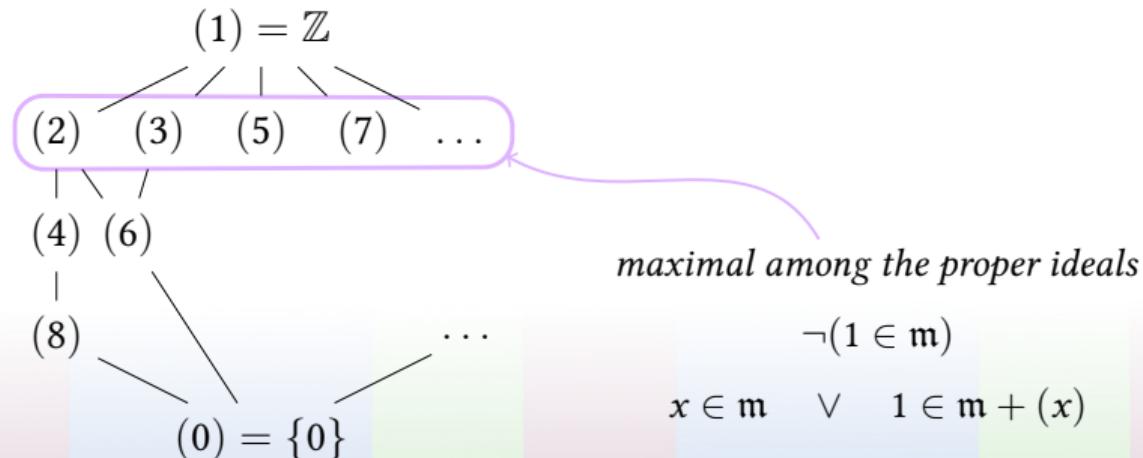


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- 4 In the general case: No,
but *first-order consequences* of the existence of a maximal ideal do hold.

Questions

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Constructive forcing (= Grothendieck toposes)

Let L be a **forcing notion**, a preorder equipped with a **covering system**.¹ **Filters** $F \subseteq L$ are subsets which are upward-closed, downward-directed and split the covering system.²

¹A covering system consists of a set $\text{Cov}(\sigma) \subseteq P(\downarrow\sigma)$ of *coverings* for each element $\sigma \in L$ subject only to the following simulation condition: If $\tau \preccurlyeq \sigma$ and $R \in \text{Cov}(\sigma)$, there should be a covering $S \in \text{Cov}(\tau)$ such that $S \subseteq \downarrow R$.

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3	f.g. ideals	—	ideals
4	f.g. ideals	$\{\sigma + (a), \sigma + (b)\}$ for each $ab \in \sigma$, $\{\}$ if $1 \in \sigma$	prime ideals
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6	$\{\star\}$	$\{\star \mid \varphi\} \cup \{\star \mid \neg\varphi\}$	witnesses of LEM

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Def. Given a monotone predicate P on L , inductively define:

$$\frac{P_\sigma}{P \mid \sigma} \quad \frac{\forall(\tau \in R). P \mid \tau}{P \mid \sigma} \quad (R \in \text{Cov}(\sigma))$$

We use quantifier-like notation: “ $\nabla\sigma. P\sigma$ ” means $P \mid \sigma$.

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A forcing notion is a template for a **forcing extension** V^∇ of the base universe V :

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 - ▶ $(<)$ is well-founded iff the generic strictly descending chain validates \perp .

The modal multiverse

In general, “ φ holds in V^∇ ” and “ φ holds in V ” are *not* equivalent.

- ▶ For **positive** extensions, they are equivalent for coherent implications.
 - e.g. the “Barr cover”.
- ▶ For **positive overt** extensions, they are equivalent for bounded first-order formulas.
 - e.g. V^∇ containing the generic surjection $\mathbb{N} \twoheadrightarrow X$, if X is inhabited.

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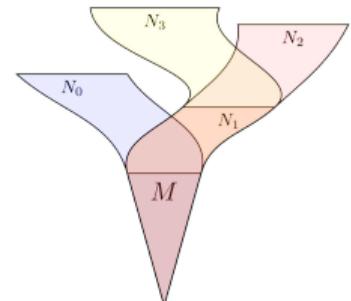
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Def. A statement φ holds ...

- ▶ **everywhere** ($\Box\varphi$) iff it holds in every extension.
- ▶ **somewhere** ($\Diamond\varphi$) iff it holds in some positive extension.
- ▶ **proximally** ($\lozenge\varphi$) iff it holds in some positive overt extension.

Foreshadowed by:

- 1984 André Joyal, Miles Tierney. *An extension of the Galois theory of Grothendieck*.
- 1987 Andreas Blass. *Well-ordering and induction in intuitionistic logic and topoi*.
- 2010s Milly Maietti, Steve Vickers. Ongoing work on arithmetic universes.
- 2011 Joel David Hamkins. *The set-theoretic multiverse*.
- 2013 Shawn Henry. *Classifying topoi and preservation of higher order logic by geometric morphisms*.



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For every inhabited set X , *proximally* there is an enumeration $\mathbb{N} \twoheadrightarrow X$.

A preorder is well iff *everywhere*, every sequence is good.

A ring element is nilpotent iff all prime ideals *everywhere* contain it.

For every ring, *proximally* there is a maximal ideal.

A relation is well-founded iff *everywhere*, there is no descending chain.

Somewhere,
the law of excluded middle holds.

The multiverse perspective

- 1 *Why has the inductive revolution been so powerful?*

Because the inductive conditions are equivalent to truth in *all* forcing extensions.

- 2 *Why do proofs using Zorn's maximal ideals work so well in constructive algebra?*

Because every ring proximally has a maximal ideal.

- 3 *Why are elements of $\bigcap_{\mathfrak{p}} \mathfrak{p}$ not necessarily nilpotent?*

Because we forgot the prime ideals in forcing extensions.

- 4 *How can we extract computational content from classical proofs?*

By traveling the multiverse (upwards, keeping ties to the base), exploiting that

- LEM holds *somewhere* and
- DC holds *proximally*.

```

module _ (A : Set) where

open import Data.List
open import Data.List.Membership.Propositional
open import Data.Product

data Eventually (P : List A → Set) : List A → Set where
  now
    : {σ : List A}
    → P σ
    → Eventually P σ
  later
    : {σ : List A} {a : A}
    → ((τ : List A) → a ∈ (σ ++ τ) → Eventually P (σ ++ τ))
    → Eventually P σ

```

U:**- Countable.agda All L1 <N> (Agda:Checked +5 Undo-Tree)



U:%*- *All Done* All L1 <M> (AgdaInfo Undo-Tree)

Partial Agda formalization available.