

USING THE INTERNAL LANGUAGE OF TOPOSES IN ALGEBRAIC GEOMETRY

INGO BLECHSCHMIDT

ABSTRACT. There are several important topoi associated to a scheme, for instance the petit and gros Zariski topoi. These come with an internal mathematical language which closely resembles the usual formal language of mathematics, but is “local on the base scheme”:

For example, from the internal perspective, the structure sheaf looks like an ordinary local ring (instead of a sheaf of rings with local stalks) and vector bundles look like ordinary free modules (instead of sheaves of modules satisfying a certain condition). The translation of internal statements and proofs is facilitated by an easy mechanical procedure.

These expository notes give an introduction to this topic and show how the internal point of view can be exploited to give simpler definitions and more conceptual proofs of the basic notions and observations in algebraic geometry. No prior knowledge about topos theory and formal logic is assumed.

CONTENTS

1.	Introduction	1
2.	Kripke–Joyal semantics	1
3.	Sheaves of rings	3
3.1.	Reducedness	3
3.2.	Locality	3
3.3.	Field properties	3
4.	Sheaves of modules	4
5.	Rational functions and Cartier divisors	4
6.	Relative spectrum	4
7.	Modalities	4
8.	Unsorted	4

1. INTRODUCTION

2. KRIPKE–JOYAL SEMANTICS

Let X be a topological space. Later, X will be the underlying space of a scheme.

Definition 2.1 (Kripke–Joyal semantics of a sheaf topos). The meaning of

$$U \models \varphi \quad (\text{“}\varphi \text{ holds on } U\text{”})$$

for open subsets $U \subseteq X$ and formulas φ is given by the following rules, recursively in the structure of φ :

$$\begin{array}{ll}
U \models f = g : \mathcal{F} & :\iff f|_U = g|_U \in \Gamma(U, \mathcal{F}) \\
U \models \varphi \wedge \psi & :\iff U \models \varphi \text{ and } U \models \psi \\
U \models \varphi \vee \psi & :\iff \text{there exists a covering } U = \bigcup_i U_i \text{ s. th. for all } i: \\
& \quad U_i \models \varphi \text{ or } U_i \models \psi \\
U \models \varphi \Rightarrow \psi & :\iff \text{for all open } V \subseteq U: V \models \varphi \text{ implies } V \models \psi \\
U \models \forall f : \mathcal{F}. \varphi(f) & :\iff \text{for all sections } f \in \Gamma(V, \mathcal{F}), V \subseteq U: V \models \varphi(f) \\
U \models \exists f : \mathcal{F}. \varphi(f) & :\iff \text{there exists a section } f \in \Gamma(U, \mathcal{F}) \text{ s. th. } U \models \varphi(f) \\
& \quad \text{there exists a covering } U = \bigcup_i U_i \text{ s. th. for all } i: \\
& \quad \text{there exists } f_i \in \Gamma(U_i, \mathcal{F}) \text{ s. th. } U_i \models \varphi(f_i) \\
U \models \forall \mathcal{F}. \varphi(\mathcal{F}) & :\iff \text{for all sheaves } \mathcal{F} \text{ on } V, V \subseteq U: V \models \varphi(\mathcal{F}) \\
U \models \exists \mathcal{F}. \varphi(\mathcal{F}) & :\iff \text{there exists a covering } U = \bigcup_i U_i \text{ s. th. for all } i: \\
& \quad \text{there exists a sheaf } \mathcal{F}_i \text{ on } U_i \text{ s. th. } U_i \models \varphi(\mathcal{F}_i)
\end{array}$$

Remark 2.2. The last two rules, concerning *unbounded quantification*, are not part of the classical Kripke–Joyal semantics, but instead of Mike Shulman’s stack semantics [?], a slight extension. They are needed so that we can formulate universal properties in the internal language.

The rules are not all arbitrary. They are finely concerted to make the following propositions true, which are crucial for a proper appreciation of the internal language.

Proposition 2.3 (Locality of the internal language). *Let $U = \bigcup_i U_i$ be covered by open subsets. Let φ be a formula. Then*

$$U \models \varphi \quad \text{iff} \quad U_i \models \varphi \text{ for each } i.$$

Proof. Induction on the structure of φ . Note that the canceled rules would make this proposition false. \square

Proposition 2.4 (Soundness of the internal language). *If a formula φ implies a further formula ψ in intuitionistic logic, then*

$$U \models \varphi \quad \text{implies} \quad U \models \psi.$$

Proof. Proof by induction on the structure of formal intuitionistic proofs; we are to show that any inference rule of intuitionistic logic is satisfied by the Kripke–Joyal semantics. For instance, there is the following rule governing disjunction:

If $\varphi \vee \psi$ holds, and both φ and ψ imply a further formula χ , then χ holds.

So we are to prove that if $U \models \varphi \vee \psi$, $U \models (\varphi \Rightarrow \chi)$, and $U \models (\psi \Rightarrow \chi)$, then $U \models \chi$. This is done as follows: By assumption, there exists a covering $U = \bigcup_i U_i$ such that on each U_i , $U_i \models \varphi$ or $U_i \models \psi$. Again by assumption, we may conclude that $U_i \models \chi$ for each i . The statement follows because of the locality of the internal language.

A complete list of which rules are to prove is in [?, D1.3.1]. \square

- geometric formulas

- geometric constructions
- simplification rules
- first steps: invertibility, nilpotency (needed later)

3. SHEAVES OF RINGS

3.1. Reducedness. Recall that a scheme X is *reduced* if and only if all stalks $\mathcal{O}_{X,x}$ are reduced rings. Since the condition on a ring R to be reduced is a geometric implication,

$$\forall s : R. s^2 = 0 \implies s = 0,$$

we immediately obtain the following characterization of reducedness in the internal language:

Proposition 3.1. *A scheme X is reduced iff, from the internal point of view, the ring \mathcal{O}_X is reduced.*

3.2. Locality. Recall the usual definition of a local ring: a ring possessing exactly one maximal ideal. This is a higher-order condition and in particular not of a geometric form. Therefore, for our purposes, it's better to adopt the following elementary definition of a local ring.

Definition 3.2. A *local ring* is a ring R such that $1 \neq 0$ in R and for all $x, y \in R$

$$x + y \text{ invertible} \implies x \text{ invertible} \vee y \text{ invertible}.$$

In classical logic, it's an easy exercise to show the equivalence of this definition with the usual one. In intuitionistic logic, we would need to be more precise in order to even ask the question of equivalence, since intuitionistically, the notion of a maximal ideal bifurcates into several non-equivalent notions.

Proposition 3.3. *In the internal language of a scheme X (or a locally ringed space), the ring \mathcal{O}_X is a local ring.*

Proof. The stated locality condition is a conjunction of two geometric implications (the first one being $1 = 0 \Rightarrow \perp$, the second being the displayed one) and holds on each stalk. \square

3.3. Field properties. From the internal point of view, the structure sheaf \mathcal{O}_X of a scheme X is *almost* a field, in the sense that any element which is not invertible is nilpotent. This is a genuine property of schemes, not shared with general locally ringed spaces.

Proposition 3.4. *Let X be a scheme. Then*

$$\text{Sh}(X) \models \forall s : \mathcal{O}_X. \neg(\ulcorner s \text{ invertible} \urcorner) \Rightarrow \ulcorner s \text{ nilpotent} \urcorner.$$

Proof. By the locality of the internal language and since X can be covered by open affine subsets, it's enough to show that for any affine scheme $X = \text{Spec } A$ and global function $s \in \Gamma(X, \mathcal{O}_X) = A$ it holds that

$$X \models \neg(\ulcorner s \text{ invertible} \urcorner) \text{ implies } X \models \ulcorner s \text{ nilpotent} \urcorner.$$

The meaning of the antecedent is that any open subset on which s is invertible is empty. So in particular, the standard open subset $D(s)$ is empty. Therefore s is an element of any prime ideal of A and thus nilpotent. This implies the a priori weaker statement $X \models \ulcorner s \text{ nilpotent} \urcorner$ (which would allow s to have different indices of nilpotency on an open covering). \square

Corollary 3.5. *Let X be a scheme. If X is reduced, the ring \mathcal{O}_X is a field from the internal point of view, in the sense that*

$$\mathrm{Sh}(X) \models \forall s : \mathcal{O}_X. \neg(\ulcorner s \text{ invertible} \urcorner) \Rightarrow s = 0.$$

The converse holds as well.

Proof. We can prove this purely in the internal language: It suffices to give an intuitionistic proof of the fact that a local ring which satisfies the condition of the previous proposition fulfills the stated field condition if and only if it is reduced. This is straightforward. \square

This field property is very useful. We will put it to good use when giving a simple proof of the fact that \mathcal{O}_X -modules of finite type on a reduced scheme are locally free on a dense open subset (proposition ??).

- Remark that intuitionistically, the notion of a field bifurcates into several inequivalent notions
- discreteness

4. SHEAVES OF MODULES

- of finite type, of finite presentation, coherent
- basic lemmas
- flatness
- important hard exercise

5. RATIONAL FUNCTIONS AND CARTIER DIVISORS

- internal definition of K_X
- internal definition of Cartier divisors
- correspondence between Cartier divisors and sub- \mathcal{O}_X -modules of K_X

6. RELATIVE SPECTRUM

- ...

7. MODALITIES

- negneg
- spreading of properties from stalk to neighbourhood
- internal sheafification

8. UNSORTED

- “functoriality”
- Kähler differentials
- completion of the natural numbers, rank function
- closed and open subschemes
- reduced closed subscheme
- Koszul resolution
- meta properties, uses (e.g. nilpotent on stalks iff globally nilpotent, some lemmas about limits of modules)
- locally small categories
- big Zariski topos

- open/closed immersions
 - morphisms of schemes...
 - proper maps...
 - limits and colimits...
 - related work: Mulvey/Burden, Wraith, Vickers, the Bohr topos crew, Awodey, ...
- E-mail address:* iblech@web.de