# Fibered categories

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Fibered and indexed categories formalize objects and morphisms which can be pulled back along morphisms of a base category. One reason they are important is that stacks are fibered categories with certain extra properties.

These are informal notes prepared for the October 2014 meeting of the *Kleine Bayerische AG* at the University of Regensburg. The notes summarize part of chapter 3 of Angelo Vistoli's *Notes on Grothendieck topologies, fibered categories and descent theory*.

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### 1 Indexed categories

**Definition 1.1.** Let  $\mathcal{C}$  be a category. A  $\mathcal{C}$ -indexed category is a pseudofunctor  $F: \mathcal{C}^{op} \to Cat$ .

The notion of a *pseudofunctor* is a slight weakening of the notion of an ordinary functor. It is best explained by looking at a example. By Cat we mean the (2-)category of categories, functors, and natural transformations. We ignore all set-theoretical issues.

**Example 1.2.** Let C be a category of spaces (topological spaces, manifolds, schemes, schemes over a fixed scheme, ...). Then there is a C-indexed category of vector bundles:

$$X \longmapsto \operatorname{VB}(X) = \text{category of vector bundles on } X$$
 
$$(f:X\to Y) \longmapsto (f^*:\operatorname{VB}(Y)\to\operatorname{VB}(X)) = \text{pullback along } f$$

Pullback of vector bundles is not associative on the nose; for maps  $X \xrightarrow{f} Y \xrightarrow{g} Z$  it is not the case that  $(g \circ f)^* = f^* \circ g^*$  as functors  $VB(Z) \to VB(X)$ . In fact, recall from general category theory that it does not make sense to compare functors on equality, since this involves comparing objects on equality. A pseudofunctor does not need to satisfy the functor axioms on the nose, it suffices that they are satisfied up to coherent natural isomorphisms. We give the precise definition below.

**Example 1.3.** Let  $\mathcal{C}$  be a category of schemes. Then there is a  $\mathcal{C}$ -indexed category of quasicoherent sheaves of modules, defined by  $X \mapsto \mathrm{QCoh}(X)$ .

**Example 1.4.** There is a Ring-indexed category of modules, given by  $R \mapsto \text{Mod}(R)$  (the category of R-modules) and the restriction of scalars functors.

**Example 1.5.** Let C be a category with pullbacks. There is a canonical C-indexed category  $\mathbb{C}$ , the *self-indexing of* C:

$$\begin{array}{ccc} X & \longmapsto & \mathcal{C}/X \\ (f:X\to Y) & \longmapsto & (f^*:\mathcal{C}/Y\to\mathcal{C}/X) \end{array}$$

As usual,  $\mathcal{C}/X$  is the slice category of morphisms to X; its objects are morphisms  $T \to X$ , where  $T \in \mathcal{C}$  is arbitrary, and its morphisms are commutative triangles. The functor  $f^*$  is defined by sending an object  $(T \to Y)$  to some chosen pullback  $(T \times_Y X \to X)$ .

Remark 1.6. A morphism  $T \xrightarrow{\pi} X$  in the category of sets can be thought of as an X-indexed family of sets, namely the fibers  $\pi^{-1}[\{x\}]$ , where x ranges over X. By analogy, we visualize a morphism  $T \to X$  in an arbitrary category  $\mathcal{C}$  as an X-indexed family of objects in  $\mathcal{C}$ . Thus  $\mathcal{C}/X$  can be thought of as the category of X-indexed families of objects in  $\mathcal{C}$ , and the functors  $f^*$  reindex families.

**Example 1.7.** Any presheaf  $F: \mathcal{C}^{\text{op}} \to \text{Set}$  (that is, an ordinary functor) gives rise to a  $\mathcal{C}$ -indexed category by postcomposing with the embedding  $\text{Set} \to \text{Cat}$ , which associates to any set its induced discrete category. In particular, any object  $A \in \mathcal{C}$  gives rise to a representable  $\mathcal{C}$ -indexed category A given by  $T \mapsto \text{Hom}_{\mathcal{C}}(T, A)$ .

Stacks are special kinds of indexed categories, so indexed categories should be geometric in some sense. Where is the geometry? Where are the topological spaces, the points?

Let  $\mathcal{C}$  be a category of spaces and F a  $\mathcal{C}$ -indexed category. We can infuse geometric content by imagining, for any space  $A \in \mathcal{C}$ , the category F(A) to be the category of maps  $A \to F$ . So even though F does not define topological structure directly, we can *probe* it by test spaces. This is the general idea of Grothendieck's functor

of points philosophy (see http://ncatlab.org/nlab/show/motivation+for+sheaves, +cohomology+and+higher+stacks for a beautiful exposition).

How good is this idea? We expect to be able to precompose morphisms  $f: A \to F$  by morphisms  $g: B \to A$  in  $\mathcal{C}$  to obtain morphisms  $f \circ g: B \to F$ . This is indeed possible:  $f \circ g$  can be realized as  $g^*(f)$ .

But the resulting notion fails in general to be local: If  $A = \bigcup_i U_i$  is covered by open subsets, we expect to be able to define maps  $A \to F$  by glueing compatible maps  $U_i \to F$ . This is not possible with an arbitrary indexed category – the notion of a topology does not enter the definition in any way. A stack is an indexed category where this kind of pathology does not arise.

**Example 1.8.** The *points* of the indexed category VB of vector bundles (example 1.2) are by definition (and by analogy with the classical situation) the maps  $pt \to VB$ . So for any natural number n, there is a point of VB (corresponding to a vector space of dimension n). These points have lots of automorphisms, namely the invertible  $(n \times n)$ -matrices.

Remark 1.9. In speaking of morphisms from a test space A to an indexed category F we commit a type error. This can be fixed by replacing A with its induced indexed category  $\underline{A}$ . A 2-categorical Yoneda lemma then gives a canonical equivalence  $\operatorname{Hom}(\underline{A}, F) \simeq F(A)$ .

We are now ready to appreciate a precise definition of a pseudofunctor.

**Definition 1.10.** Let  $\mathcal{C}$  be a category. A pseudofunctor  $F: \mathcal{C}^{\mathrm{op}} \to \mathrm{Cat}$  consists of

- a category F(X) for each object X of C,
- a functor  $F(f) =: f^* : F(Y) \to F(X)$  for each morphism  $X \xrightarrow{f} Y$  in  $\mathcal{C}$ ,
- a natural isomorphism  $\eta_X : (\mathrm{id}_X)^* \Rightarrow \mathrm{Id}_{F(X)}$  for each object X in  $\mathcal{C}$ , and
- a natural isomorphism  $\alpha_{f,g}: f^* \circ g^* \Rightarrow (g \circ f)^*$  for each diagram  $X \xrightarrow{f} Y \xrightarrow{g} Z$  in C such that the following *coherence conditions* hold:
  - For any  $X \xrightarrow{f} Y$  in C,  $\alpha_{\mathrm{id}_X,f} = \eta_X f^* : \mathrm{id}_X^* \circ f^* \Rightarrow f^*$ .
  - For any  $X \xrightarrow{f} Y$  in C,  $\alpha_{f,id_Y} = f^* \eta_Y : f^* \circ id_Y^* \Rightarrow f^*$ .
  - For any  $X \xrightarrow{f} Y \xrightarrow{g} Z \xrightarrow{h} W$  in C, the diagram commutes.

#### XXX diagram

The coherence conditions are needed for the following reason: When dealing with vector bundles or similar objects which can be pulled back, we are used to identify iterated pullbacks with pullbacks along compositions (" $f^* \circ g^* = (g \circ f)^*$ "). This is justified only because between any two pullback expressions, e. g.

$$(f^* \circ (h \circ g)^* \circ p^* \circ q^*)(E)$$
 and  $((g \circ f)^* \circ (q \circ p \circ h)^*)(E)$ ,

we have exactly one agreed-upon canonical isomorphism (of course, there are lots of other non-canonical isomorphisms as well). In a general indexed category, we can construct these canonical isomorphisms out of the given  $\eta_X$  and  $\alpha_{f,g}$ ; the coherence conditions ensure that – no matter in which way we do this – we always obtain the same canonical isomorphism.

Remark 1.11. Indexed categories are also important in categorical logic and topos theory. For instance, any geometric morphism  $\mathcal{F} \xrightarrow{f} \mathcal{E}$  of toposes induces an  $\mathcal{E}$ -indexed category  $\mathbb{F}$ , given by  $X \mapsto \mathcal{F}/f^*X$ . From the point of view of the internal language of  $\mathcal{E}$ , an  $\mathcal{E}$ -indexed category looks like an ordinary (large, but locally small) category. This is useful for doing relative category theory.

### 2 Fibered categories

Recall that an I-indexed family  $(X_i)_{i\in I}$  of sets can equivalently be given by

- a map  $I \to \text{Ob Set}, i \mapsto X_i$ , or
- a map  $\coprod_{i \in I} X_i \xrightarrow{\pi} I$ . (Recover  $X_i$  as the fiber of i under  $\pi$ .)

Indexed categories are like the first approach, while fibered categories are like the second approach.

**Definition 2.1.** Let  $\mathcal{C}$  be a category. A category fibered over  $\mathcal{C}$  is a category  $\mathcal{F}$  together with a functor  $\mathcal{F} \xrightarrow{p} \mathcal{C}$  such that "pullbacks exist": For any morphism  $Y \xrightarrow{f} X$  in  $\mathcal{C}$  and any object E in the fiber category  $\mathcal{F}(X)$  (the subcategory of  $\mathcal{F}$  consisting of those objects which map to X and those morphisms which map to X and X under X and exist an object X of X and a cartesian morphism X with X with X with X with X and X with X and X and X and X with X with X with X with X with X and X with X with X and X and X with X

$$E' \xrightarrow{\phi} E$$

$$\downarrow \qquad \qquad \downarrow$$

$$Y \xrightarrow{f} X$$

The notion of a cartesian morphism is modeled on the universal property of usual pullbacks: A morphisms  $\phi: E' \to E$  with  $p(\phi) = f$  as in the definition is cartesian if and only if, for any morphism  $g: Z \to Y$ , for any object E'' in  $\mathcal{F}(Z)$ , and any morphism  $\psi: E'' \to E$  such that  $p(\psi) = f \circ g$ , there exists a unique morphism  $\alpha: E'' \to E'$  such that  $p(\alpha) = g$  and  $\phi \circ \alpha = \psi$ .

Remark 2.2. The definition of a fibered category as stated is evil in the technical sense that it requires certain equalities between objects (namely, p(E') = Y). Since

important examples of categories which ought to be fibered categories actually fulfill these evil equalities, this defect is tolerable in practice. But for some purposes, a non-evil variant is necessary; see for instance http://ncatlab.org/nlab/show/Grothendieck+fibration#StreetFibration.

**Example 2.3.** The incarnation of example 1.2 as a fibered category is the category VB of vector bundles on arbitrary base spaces. The projection functor VB  $\rightarrow \mathcal{C}$  associates to any vector bundle its base space.

**Definition 2.4.** An indexed functor  $\mathcal{F} \to \mathcal{G}$  between fibered categories over  $\mathcal{C}$  is an ordinary functor  $\mathcal{F} \to \mathcal{G}$  which commutes with the projection functors to  $\mathcal{C}$  and which carries cartesian morphisms to cartesian morphisms.

### 3 Equivalence between indexed and fibered categories

**Proposition 3.1.** There is a pseudoequivalence between the 2-category of indexed categories, indexed functors, and indexed natural transformations; and the 2-category of fibered categories, fibered functors, and fibered natural transformations.

From fibered categories to indexed categories. If suitable choice principles are available, any fibered category  $\mathcal{F} \to \mathcal{C}$  defines a  $\mathcal{C}$ -indexed category  $F: \mathcal{C}^{\mathrm{op}} \to \mathrm{Cat}$  by setting  $F(X) := \mathcal{F}(X)$ . For morphisms  $f: Y \to X$  in  $\mathcal{C}$ , the pullback functors  $f^*: F(X) \to F(Y)$  are defined by associating to any object E a chosen pullback E' as in the definition of a fibered category. The action of  $f^*$  on morphisms is given by the universal property of cartesian morphisms.

One can check that this construction gives indeed rise to a pseudofunctor. In fact, the proof is very similar to the proof that the self-indexing  $X \mapsto \mathcal{C}/X$  is a pseudofunctor.

From indexed categories to fibered categories. In the opposite direction, any  $\mathcal{C}$ -indexed category  $F:\mathcal{C}^{\mathrm{op}}\to\mathrm{Cat}$  induces a fibered category  $f\in\mathcal{C}$  via the Grothendieck construction. Mimicking the case of families of sets, the objects of  $f\in\mathcal{C}$  are pairs  $f\in\mathcal{C}$ ,  $f\in\mathcal{C}$ ,  $f\in\mathcal{C}$ . A morphism  $f\in\mathcal{C}$  is a pair  $f\in\mathcal{C}$  and a morphism  $f\in\mathcal{C}$  sends an object  $f\in\mathcal{C}$  to  $f\in\mathcal{C}$  and a morphism  $f\in\mathcal{C}$  one can check that this category over  $f\in\mathcal{C}$  is indeed a fibered category. The cartesian morphisms  $f\in\mathcal{C}$  is indeed a fibered category. The cartesian morphisms  $f\in\mathcal{C}$  is indeed a fibered category. The cartesian morphisms  $f\in\mathcal{C}$  is indeed a fibered category. The cartesian morphisms  $f\in\mathcal{C}$  is indeed a fibered category. The cartesian morphisms  $f\in\mathcal{C}$  is indeed a fibered category.

**Example 3.2.** Let  $\mathcal{C}$  be a category with pullbacks. Then  $\int \mathbb{C}$  (see example 1.5) is equivalent to the *category of morphisms in*  $\mathcal{C}$ . This category has arbitrary morphisms  $Y \to X$  as objects and commutative squares as morphisms. The equivalence is given by sending an object  $(X \in \mathcal{C}, (Y \xrightarrow{f} X))$  to f.

**Example 3.3.** Let  $\mathcal{C}$  be an arbitrary category and A a fixed object of  $\mathcal{C}$ . Then  $\int \underline{A}$  (see example 1.7) is equivalent to the slice category  $\mathcal{C}/A$ , by sending an object  $(X \in \mathcal{C}, X \xrightarrow{f} A)$  to f.

### 4 Special types of indexed and fibered categories

**Definition 4.1.** A category fibered in groupoids (sets) [equivalence relations] is a fibered category  $\mathcal{F}$  such that the fiber categories  $\mathcal{F}(X)$  are all groupoids (discrete categories) [categories equivalent to discrete categories].

Equivalently, the associated pseudofunctor  $\mathcal{C}^{op} \to \operatorname{Cat}$  factors over Grpd  $\hookrightarrow \operatorname{Cat}$  (Set  $\hookrightarrow \operatorname{Cat}$ ) [EqRel  $\hookrightarrow \operatorname{Cat}$ ].

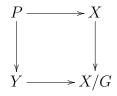
**Proposition 4.2.** A functor  $\mathcal{F} \xrightarrow{p} \mathcal{C}$  is ...

- a category fibered in groupoids if and only if any morphism in  $\mathcal{F}$  is cartesian and for any morphism  $Y \xrightarrow{f} X$  in  $\mathcal{C}$  and any object E in  $\mathcal{F}(X)$ , there exists a morphism  $E' \xrightarrow{\phi} E$  in  $\mathcal{F}$  such that  $p(\phi) = f$ .
- a category fibered in sets if and only if for any morphism  $Y \xrightarrow{f} X$  in C and any object E in  $\mathcal{F}(X)$ , there exists a unique morphism  $E' \xrightarrow{\phi} E$  in  $\mathcal{F}$  such that  $p(\phi) = f$ .

## 5 An important example: quotient stacks

For this example, we fix a category  $\mathcal{C}$  of spaces. We will be so sketchy that it does not matter which category  $\mathcal{C}$  actually is. The source for this section is the article *Algebraic stacks* of Tomás Gómez.

Let G be a group acting on a space X. Assume that the action is free, i.e. for any  $g \in G$  and any  $x \in X$ ,  $g \cdot x = x$  implies g = 1. Then recall that there is a well-defined projection  $X \to X/G$ ; that this map is a G-principal bundle; and that this map is the universal G-principal bundle equipped with a G-equivariant map to X: For any G-principal bundle  $P \to Y$  equipped with a G-equivariant map  $P \to X$ , there is a unique map  $Y \to X/G$  such that the diagram



is a pullback diagram. Thus the space X/G has the pleasant property that maps  $Y \to X/G$  correspond to G-principal bundles  $P \to Y$  equipped with G-equivariant maps  $Y \to X$ . (If such a bundle is given, construct  $Y \to X/G$  by selecting, for any  $y \in Y$ , a preimage under  $P \to Y$ , and applying  $P \to X \to X/G$  to this preimage. One can check that the resulting element in X/G is independent of the choice.)

Unfortunately, if the action of G on X is not free, X/G fails to exist as an object of C or fails to have this universal property. But there is always a quotient stack [X/G] (at least an indexed category) with this property, tautologically defined by sending a test

object Y of C to the category of G-principal bundles  $P \to Y$  equipped with G-equivariant maps  $Y \to X$ .

Indeed, one can canonically define a map  $\underline{X} \to [X/G]$  and check that is the universal bundle.

**Proposition 5.1.** The stack [X/G] is representable, i. e. of the form  $\underline{A}$  for some space A in C, if and only if the action of G on X is free.

*Proof* (sketch). For the "only if" direction, check that the morphism  $X \to A$  corresponding to the universal bundle  $X \to A$  is a G-principal bundle isomorphic to  $X \to X/G$ .  $\square$ 

**Example 5.2.** Let X = pt. Then the quotient stack [X/G] is commonly named "BG".

Remark 5.3. The category of points of [X/G], i. e. [X/G](pt), is equivalent to the action groupoid (weak quotient) X//G. This groupoid has as objects the elements of X, and the set of morphisms  $x \to x'$  is the set of group elements g such that  $g \cdot x = x'$ . The action groupoid is the "right replacement" for the ill-behaved space of orbits. A very accessible exposition is in Higher-Dimensional Algebra VII: Groupoidification by John Baez, Alexander Hoffnung and Christopher Walker: http://arxiv.org/abs/0908.4305.

Remark 5.4. I believe that in some cases, the category of coherent sheaves on [X/G] can simply be described as the category of G-equivariant sheaves on X.