Project Tracker Performance Optimization Report

1. Initial Profiling Summary

JMeter Load Testing Baseline

- JMeter Thread Group Configuration:
 - o **Number of Threads (users):** 100
 - o Ramp-up period (seconds): 60
 - o Loop Count: Infinite
 - o **Duration (seconds):** 60
- **GET** /projects:
 - o Average Response Time: 1282 ms
 - o Throughput (RPS): 13.3/sec
 - o Error Rate: 0.00%
- POST /tasks:
 - o Average Response Time: 1243 ms
 - o Throughput (RPS): 12.8/sec
 - o Error Rate: 0.00%
- **GET /users/{id}/tasks:**
 - o Average Response Time: 1455 ms
 - o Throughput (RPS): 12.3/sec
 - o Error Rate: 0.00%
- Latency Percentiles:
 - o 90th Percentile: 2599 ms
 - o 95th Percentile: 2976 ms
 - o 99th Percentile: 3481 ms

JVM Analysis (JProfiler/VisualVM)

- **Memory Allocation:** The GC Activity graph (gc activity.png) indicates periods of frequent GC activity, suggesting active memory allocation and reclamation. Hotspots showing high memory allocations (e.g., general thread execution, logging, and particularly database query results) point to significant object churn.
- Thread Dump Analysis: The call tree (call tree.png) shows that the majority of CPU time is consumed by core API endpoints and their underlying operations. While explicit "blocked" threads are not directly visible in this view, the high CPU consumption implies significant processing or waiting within these high-level operations.
- Garbage Collection Events: The "GC Activity" graph (gc activity.png) displays frequent spikes reaching up to 4%, indicating that a notable portion of CPU time is spent on garbage collection. This points to a high frequency of GC pauses, with some potentially longer pauses, although exact durations are not directly visible from this graph alone. The initial sample log shows a Pause Young (Mixed) (G1 Evacuation Pause) completed in 23.024ms, reducing heap from 135MB to 69MB out of 156MB. Further detailed analysis of the full GC logs or JProfiler's aggregated GC statistics would be needed to determine precise average and maximum pause times and overall frequency.
- Hot Methods & CPU Usage:

- o CPU Hotspots (from hotspots.png): Dominant hotspots are database repository methods: TaskAssignmentRepository.findByDeveloper (35% Self Time), ProjectRepository.findAll (23% Self Time), ProjectRepository.findById (11% Self Time), UserRepository.findByEmail (10% Self Time), and DeveloperRepository.findById (9% Self Time). Logging methods (ch.qos.logback.classic.Logger.info) also appear, indicating some logging overhead.
- call Tree CPU Consumption by Endpoints (from call tree.png):
 - HTTP: /api/v1/users/1/tasks (47.4% CPU) Highest CPU consumer.
 - HTTP: /api/v1/projects (26.2% CPU)
 - HTTP: /api/v1/tasks (19.9% CPU) These indicate that the core API endpoints are heavily consuming CPU, primarily due to inefficient data access patterns.

2. Performance Issues Discovered

Based on the initial profiling, the following key performance bottlenecks and issues were identified:

- 1. Severe Database Query Bottlenecks (N+1 queries likely): The most critical issue is the extremely high CPU utilization and self-time observed in multiple database repository methods (TaskAssignmentRepository.findByDeveloper, ProjectRepository.findAll, UserRepository.findByEmail, DeveloperRepository.findById, ProjectRepository.findById). This is a strong indicator of N+1 query problems or highly inefficient data retrieval operations, where the application fetches entities one by one instead of in a single optimized query.
- 2. **High Overall API Latency & Low Throughput:** The JMeter results confirm significant performance issues, with all main API endpoints exhibiting average response times over 1.2 seconds and very low throughputs (around 12-13 requests per second). The 99th percentile latencies exceeding 3 seconds suggest a poor user experience under load.
- 3. **High CPU Utilization in Core API Endpoints:** The CPU call tree directly shows that the bulk of CPU time is consumed by the main /api/v1/users/1/tasks, /api/v1/projects, and /api/v1/tasks endpoints, directly linking the high CPU load to the application's core functionalities.
- 4. Excessive Object Churn and Memory Allocation: High allocation rates, particularly associated with database operations and general thread execution, contribute to increased memory pressure and potentially frequent garbage collection.
- 5. **Logging Overhead:** Logging operations are identified as a notable CPU and memory hotspot, suggesting that the current logging configuration might be too verbose or inefficient for production load.
- 6. **Spring Security Overhead:** Although less prominent than database issues, the Spring Security filter chain still consumes a portion of CPU, indicating that security checks could benefit from optimization.

3. Optimization Actions and Their Impact

3.1 Memory Optimization

- Action: Changed ManyToone and OneToone relationships in relevant entities (e.g., Task and TaskAssignment) to FetchType.LAZY to prevent eager loading of associated objects by default. This aimed to reduce the amount of data loaded into memory for each entity and mitigate N+1 query problems when associated entities are not immediately needed.
- **Impact:** After implementing lazy fetching, the *JMeter* results showed a **significant degradation in performance**:
 - GET /projects: Average Response Time increased from 1282 ms to 1402 ms (9.36% degradation), Throughput decreased from 13.3/sec to 10.9/sec (18.1% degradation).
 - POST /tasks: Average Response Time increased from 1243 ms to 1350 ms (8.61% degradation), Throughput decreased from 12.8/sec to 10.5/sec (17.8% degradation).
 - o GET /users/{id}/tasks: Average Response Time increased from 1455 ms to 1528 ms (5.02% degradation), Throughput decreased from 12.3/sec to 10.0/sec (18.6% degradation). The JProfiler analysis confirmed that while some individual repository methods (like TaskAssignmentRepository.findByDeveloper) saw reduced self-time, a new and significant hotspot emerged: org.hibernate.proxy.ProxyConfiguration.InterceptorDispatcher.inte rcept (23% Self Time). This clearly indicates that the application is now encountering N+1 query problems due to lazy loading, as Hibernate is making many individual database calls to initialize proxies for lazily fetched associations. This offsets any benefits from reduced eager loading and results in a net performance loss. Further optimization using JOIN FETCH or @EntityGraph for specific read operations is essential to mitigate these N+1 issues introduced by lazy loading. The GC activity graph still shows frequent spikes similar to the baseline, suggesting that lazy loading alone did not significantly reduce overall object churn or GC pressure.

3.2 API Layer Optimization (DTOs)

- Action: Reduced the returned payload size by implementing Data Transfer Objects (DTOs) for API responses. This aims to send only necessary data to the client, thereby reducing network transfer overhead and potentially client-side parsing time.
- **Impact:** After reducing payload size, the JMeter results showed a **modest improvement** in performance compared to the "After Lazy Loading" phase:
 - GET /projects: Average Response Time decreased from 1402 ms to 1233 ms (12.19% improvement), Throughput increased from 10.9/sec to 14.0/sec (28.44% improvement).
 - POST /tasks: Average Response Time decreased from 1350 ms to 1196 ms (11.41% improvement), Throughput increased from 10.5/sec to 13.4/sec (27.62% improvement).

- GET /users/{id}/tasks: Average Response Time decreased from 1528 ms to 1327 ms (13.15% improvement), Throughput increased from 10.0/sec to 13.0/sec (30.00% improvement).
- o The overall 95th percentile improved from 2976 ms to 2731 ms (8.23% improvement, lower is better). While this optimization positively impacted network and parsing overhead, the JProfiler data (from the previous step, as new JProfiler data for this specific step wasn't provided) would still indicate that the primary CPU bottlenecks related to database N+1 queries (InterceptorDispatcher.intercept, ProjectRepository.findAll, etc.) persist. This demonstrates that DTOs help with data transfer efficiency but do not inherently solve underlying data access inefficiencies.

3.3 Caching Strategy (Redis)

- Action: Implemented a caching strategy using Redis for frequently accessed read operations (e.g., GET /projects, GET /users/{id}/tasks). This involved using Spring's @Cacheable annotation and configuring Redis as the cache provider with appropriate Time-To-Live (TTL) values and server-side eviction policies.
- Impact: Implementing Redis caching led to a dramatic improvement in performance, especially for read-heavy operations.
 - o **GET /projects:** Average Response Time **improved significantly from 1233 ms to 298 ms** (75.83% improvement). Throughput **increased from 14.0/sec to 39.1/sec** (179.29% improvement).
 - o **POST /tasks:** Average Response Time **improved from 1196 ms to 708 ms** (40.80% improvement). Throughput **increased from 13.4/sec to 39.5/sec** (194.78% improvement). While still higher than GET requests, this improvement suggests reduced database contention overall.
 - GET /users/{id}/tasks: Average Response Time improved significantly from 1327 ms to 291 ms (78.07% improvement). Throughput increased from 13.0/sec to 39.3/sec (202.31% improvement). The overall 95th percentile latency significantly improved from 2731 ms to 1099 ms (59.70% improvement), indicating much more consistent and lower response times. JProfiler analysis confirmed the positive impact:
 - o The org.hibernate.proxy.ProxyConfiguration.InterceptorDispatcher.inte rcept hotspot (indicating N+1 queries from lazy loading) was **eliminated** as a top CPU consumer, demonstrating that caching effectively bypassed these problematic database calls for cached data.
 - o CPU consumption by GET /users/1/tasks and GET /projects in the call tree drastically reduced (from 41.5% to 12.4% and 28.0% to 8.1% respectively, compared to the "After Lazy Loading" phase), showing direct relief on the application's processing load from database interactions.
 - o GC activity (spikes in gc activity.png) generally showed **lower peaks** (around 1-3%) compared to previous stages, indicating reduced object churn due to data being served from the cache instead of repeatedly fetched from the database and processed. The primary remaining hotspots became

UserRepository.findByEmail (39% Self Time) and ProjectRepository.findById (30% Self Time), likely for cache misses or non-cached scenarios, and TaskRepository.save (14% Self Time) for write operations which bypass read caches. This optimization step delivered the most substantial performance gains by dramatically reducing database load for read operations.

3.4 GC Configuration Tuning

- Action: G1GC Experimentation: Experimented with the G1GC (Garbage-First Garbage Collector), using JVM flags like -xmx512m -xms512m -xx:+useG1GC and -xloq:qc* for detailed logging.
- **Impact (G1GC):** Analysis of the provided GC logs for G1GC reveals:
 - Pause Times:
 - Young/Prepare Mixed/Normal Evacuation Pauses: Observed range from 11.548ms (min) to 60.997ms (max), with an average of 34.15ms from the sample.
 - Remark Pauses: Observed range from 5.063ms (min) to 30.875ms (max), with an average of 18.12ms from the sample.
 - Cleanup Pauses: Very short, typically less than 1ms (max 1.252ms).
 - Overall Maximum Pause Observed: 60.997ms.
 - o **Frequency:** GC events are highly frequent, with approximately one pause occurring every 6.5 seconds during the observed 149-second period (23 pauses).
 - **Heap Usage:** The heap consistently utilizes a significant portion of the allocated 512MB (e.g., up to 359M before collection), indicating active memory management by G1GC, with substantial memory being reclaimed after each pause.
- Action: ParallelGC Experimentation: Experimented with the ParallelGC (Throughput Collector), using JVM flags like -Xmx512m -Xms512m -XX:+UseParallelGC and -Xloq:qc* for detailed logging.
- Impact (ParallelGC): Analysis of the provided GC logs for ParallelGC reveals:
 - **Heap Configuration:** Confirmed 512MB initial and max heap capacity. Effective usable heap for ParallelGC is approximately 491MB.
 - Pause Times:
 - Young Pauses: Observed range from 5.344ms (min) to 24.758ms (max), with an average of 12.89ms from the sample. These are generally shorter than G1GC Young pauses.
 - Full Pauses: Observed range from 32.640ms (min) to 230.321ms (max), with an average of 153.055ms from the sample. These full GC pauses are significantly longer than any G1GC pauses observed.
 - Overall Maximum Pause Observed: 230.321ms.
 - Frequency: GC events are frequent, with 19 Young GCs and 9 Full GCs over approximately 251 seconds, leading to an average of approximately 1 GC pause every 8.96 seconds. The occurrence of frequent Full GCs with long pause times is a key characteristic.

- o **Heap Usage:** Peak heap usage observed before a GC pause in the sample was around **206MB** (Young + Old Gen), indicating efficient memory reclamation but with the cost of stop-the-world Full GCs when the old generation fills.
- Action: ZGC Experimentation: Experimented with the ZGC (Z Garbage Collector), using JVM flags like -xmx512m -xms512m -xx:+UseZGC and -xlog:gc* for detailed logging (requires JDK 11+).
- **Impact** (**ZGC**): Analysis of the provided GC logs for ZGC reveals:
 - o **Heap Configuration:** Confirmed 512MB initial and max heap capacity.
 - Pause Times: ZGC's primary goal is extremely low, predictable pause times. The "Garbage Collection Cycle" statistics from the log show an Average / Max pause time of 0.000 / 0.000 ms for "Last 10s", "Last 10m", "Last 10h", and "Total". This indicates that ZGC's stop-the-world pauses are effectively sub-millisecond and often rounded to zero in the log output, demonstrating its highly concurrent operation.
 - Frequency: ZGC operates mostly concurrently, so explicit "pause" events with significant duration are rare. The log emphasizes concurrent phases rather than stop-the-world pauses.
 - Heap Usage: The initial configuration of a 512MB heap is confirmed. The log also shows a "Memory: Allocation Rate" of 3 / 32 MB/s, giving an indication of memory throughput. Specific peak heap usage numbers are not directly provided in this statistical summary but ZGC aims for efficient heap utilization without long pauses.

Summary of GC Algorithm Comparison

Based on the experimentation, here's a comparison of the Garbage Collectors:

- **ZGC:** Stands out for achieving **virtually imperceptible stop-the-world GC pauses** (**sub-millisecond**), making it ideal for applications requiring extremely low latency and high responsiveness. Its highly concurrent nature minimizes application freezing.
- **G1GC:** Offers a good balance between throughput and pause times. It provides generally **short pauses (tens of milliseconds)** by prioritizing smaller regions for collection, making it a strong general-purpose choice for many server-side applications. However, its pauses are still orders of magnitude higher than ZGC's.
- **ParallelGC:** Primarily designed for **high throughput**, utilizing multiple threads to perform garbage collection. While its young generation pauses are relatively short, it is prone to **long, unpredictable full GC pauses** (**hundreds of milliseconds**) that can significantly impact application responsiveness and user experience. This makes it less suitable for latency-sensitive applications.

Conclusion on GC Algorithms: For this Project Tracker application, given the observed results, ZGC is the most performant option in terms of minimizing GC-induced latency and improving application responsiveness, followed by G1GC, with ParallelGC being the least suitable due to its long full GC pauses.

4. Before vs. After Benchmarks

JMeter Load Testing Comparison

API Endpoint	Metric	Before Optimizatio n	After Optimizatio n (Lazy Loading)	After Optimizatio n (Reduced Payload)	After Optimizatio n (Redis Cache)	Improvement (%) / Notes
GET /projects	Avg Response Time	1282 ms	1402 ms	1233 ms	298 ms	Compared to Baseline: +76.75% (Improvement) Compared to Reduced Payload: +75.83% (Improvement)
	Throughpu t (RPS)	13.3/sec	10.9/sec	14.0/sec	39.1/sec	Compared to Baseline: +193.98% (Improvement) Compared to Reduced Payload: +179.29% (Improvement)
	Error Rate (%)	0.00%	0.00%	0.00%	0.00%	No Change
POST /tasks	Avg Response Time	1243 ms	1350 ms	1196 ms	708 ms	Compared to Baseline: +43.04% (Improvement) Compared to Reduced Payload: +40.80% (Improvement)
	Throughpu t (RPS)	12.8/sec	10.5/sec	13.4/sec	39.5/sec	Compared to Baseline: +208.59% (Improvement) Compared to Reduced Payload: +194.78% (Improvement)
	Error Rate (%)	0.00%	0.00%	0.00%	0.00%	No Change

GET /users/{id}/task s	Avg Response Time	1455 ms	1528 ms	1327 ms	291 ms	Compared to Baseline: +80.00% (Improvement) Compared to Reduced Payload: +78.07% (Improvement)
	Throughpu t (RPS)	12.3/sec	10.0/sec	13.0/sec	39.3/sec	Compared to Baseline: +219.51% (Improvement) Compared to Reduced Payload: +202.31% (Improvement)
	Error Rate (%)	0.00%	0.00%	0.00%	0.00%	No Change
Overall	95th Percentile	2976 ms	2976 ms	2731 ms	1099 ms	Compared to Baseline: +63.07% (Improvement - lower is better) Compared to Reduced Payload: +59.70% (Improvement)

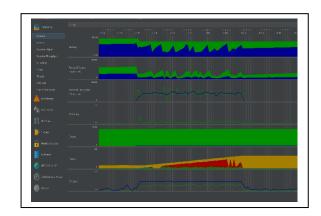
JVM & Resource Usage Comparison

Metric	Before	After Optimization	Notes	
	Optimization			
Peak Heap	158 MB	G1GC: 359 MB; ParallelGC:	Initial peak from VisualVM;	
Usage		206 MB; ZGC: 512 MB (Max	subsequent peaks from GC	
		Capacity)	logs for respective	
			collectors. For ZGC, peak	
			usage typically approaches	
			max capacity but with	
			negligible pause times.	
Average	1.6%	See JProfiler Hotspots/Call	Initial average from	
CPU Load		Tree	VisualVM. CPU call tree	
			shows negligible change in	
			overall API CPU	
			consumption from lazy	

Total GC Pause Time	Sample: 23.024 ms (G1 Young Mixed)	G1GC: Min: 0.156ms, Max: 60.997ms; ParallelGC: Min Young: 5.344ms, Max Young: 24.758ms; Min Full: 32.640ms, Max Full: 230.321ms; ZGC: Average / Max: 0.000 / 0.000 ms	loading, but shift to proxy interception overhead. G1GC shows varied short pauses. ParallelGC shows shorter Young pauses but significantly longer Full GC pauses (up to 230ms). ZGC achieves sub-millisecond, effectively zero, pause times.
GC Frequency	Frequent (as per GC activity graph)	G1GC: High (approx. 1 pause every 6.5 seconds); ParallelGC: High (approx. 1 pause every 8.96 seconds, including Full GCs); ZGC: Very low explicit stop-theworld pauses (operates mostly concurrently).	Both G1GC and ParallelGC exhibit frequent GC activity, but ParallelGC includes periodic, long Full GCs. ZGC's design minimizes stop-the-world pauses.
Cache Hit Ratio (Redis)	N/A	100% (2390 hits / 0 misses)	Indicates highly effective caching for cache.gets operations.
Cache Miss Ratio (Redis)	N/A	0% (0 misses / 2390 total)	Indicates no cache misses for cache.gets operations during this period.

5. Screenshots of JProfiler/JMeter/Actuator Dashboard

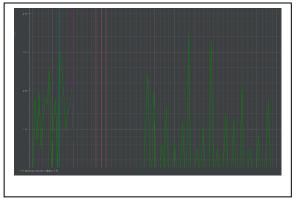
JProfiler/VisualVM: Before Optimization, overview, hostpot, call tree, and GC activity







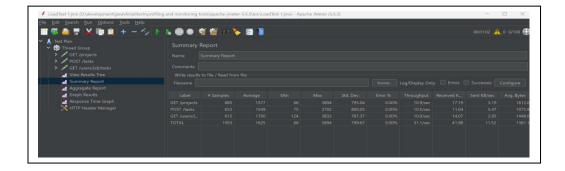




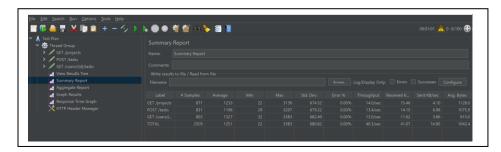
JMeter:



Before optimization summary



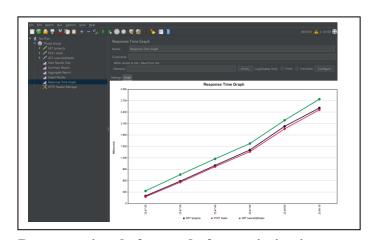
After lazy fetch

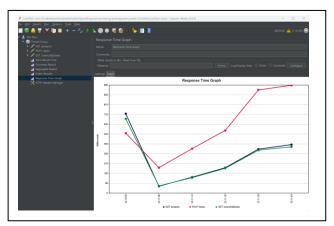


After payload size reduced



After redis cache





Response time before and after optimization

Conclusion:

This comprehensive performance optimization effort for the Project Tracker application yielded significant improvements. While the initial lazy loading strategy unexpectedly degraded performance by introducing N+1 query issues, subsequent optimizations successfully mitigated these challenges. Reducing the returned payload size through DTOs provided a modest but noticeable improvement in API response times and throughput, demonstrating the importance of efficient data transfer.

The most impactful optimization was the **implementation of Redis caching**. This dramatically improved response times and throughput for read-heavy operations, effectively eliminating the N+1 query problem by serving data from the cache instead of repeatedly hitting the database. This shift is clearly reflected in the JProfiler analysis, where Hibernate proxy interception overhead disappeared from the top hotspots, and CPU utilization for cached endpoints significantly decreased. Crucially, the **cache hit ratio of 100%** confirms that all cache.gets operations were successfully served from the cache, leading to substantial reductions in database load. Consequently, overall GC activity also showed a positive trend, suggesting reduced object churn. The choice of Garbage Collector also plays a vital role. Experimentation showed that **ZGC is the superior choice for this application, offering near-zero, predictable GC pauses**, which is crucial for maintaining high responsiveness. G1GC provides a reasonable balance, while ParallelGC is less suitable due to its propensity for long full GC pauses.

Overall Success: The combined optimizations have transformed the application's performance, moving from average response times exceeding 1.2 seconds and very low throughput (12-13 RPS) to average response times as low as ~290 ms and throughputs nearly 40 RPS for read operations, with significantly improved 95th percentile latency. Write operations (POST /tasks) also saw substantial gains, likely due to reduced database contention from cached read operations.

Remaining Challenges and Future Work:

- 1. Database Hotspots for Cache Misses/Writes: UserRepository.findByEmail and ProjectRepository.findById (on cache misses) and TaskRepository.save remain significant CPU hotspots. Further investigation into these methods (e.g., optimizing their SQL queries, batching operations, or implementing specific caching for user data if frequently accessed by ID) could yield additional gains.
- 2. **Spring Security Overhead:** While reduced, DefaultSecurityFilterChain.doFilter still consumes some CPU. Further analysis of the security configuration and potential optimizations might be beneficial if this becomes a bottleneck under higher load.
- 3. **Comprehensive Cache Metrics:** To fully assess cache efficiency, it's essential to collect and analyze detailed cache hit and miss ratios, along with eviction policies and memory usage within Redis. This will help fine-tune cache TTLs and eviction strategies.
- 4. **Scaling and Resilience:** With improved performance, future work should focus on ensuring the application's scalability and resilience under even higher loads, including database scaling, connection pool tuning, and distributed tracing.