

# Cool Numbers

This problem can be easily solved using dynamic programming.

We maintain a 2-Dimensional  $dp$  matrix  $dp[n+1][10]$ , where each entry  $dp[i][j]$  represents the number of *cool* numbers starting at index  $i$  from the left and with first digit as  $j$  (here we allow  $j$  to be 0, and we deal with the case of first digit to be non-zero while computing the final answer).

Since difference between two consecutive digits is atmost  $c$ , the  $dp$  relations are as follows:

$$dp[i][j] = \sum_{k=\max(0, j-c)}^{\min(9, j+c)} dp[i+1][k].$$

For the base case, we set  $dp[n][j] = 1$  for each  $9 \geq j \geq 0$ . The proof of this can be argued inductively, with the induction claim being that for each first digit  $j$ ,  $dp[i][j]$  gives the number of *cool* numbers (where 0 can also be the first digit) beginning at index  $i$ . Now, for any number beginning at index  $i-1$ , if it is *cool*, then the next digit will differ by atmost  $c$ , and the the number from index  $i$  till  $n$  must also be *cool*, hence the above relation covers all such *cool* numbers and counts them exactly once. Hence our transition is valid.

To compute the final answer  $ans$ , simply do :  $ans = \sum_{k=1}^{k=9} dp[1][k]$  (here we exclude 0 as the starting digit).

Time Complexity:  $O(bnc)$

Space Complexity:  $O(bn)$

Here  $b = 10$  is the base of the number system chosen.