

# PROBLEM SET 8: RSA AND PRIME TESTING (DUE NOVEMBER 15)

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**Problem 1.** Do 31.7-1 in the book.

*Solution:* We have  $\phi(319) = \phi(29 \cdot 11) = 28 \cdot 10 = 280$ . We have to find  $d$  with  $3d \equiv 1 \pmod{280}$ . Now  $280 = 3 \cdot 93 + 1$ , so  $3 \cdot 93 \equiv -1 \pmod{280}$  and  $3 \cdot (-93) \equiv 1 \pmod{280}$ . Take  $d = (-93) + 280 = 197$ . To encrypt  $M = 100$ , we compute  $100^3 \bmod 319$ . Now  $100^2 = 31 \cdot 319 + 111$  and  $100 \cdot 111 = 11100 = 34 \cdot 319 + 254$ . So finally  $100^3 \equiv 100 \cdot 111 \equiv 254 \pmod{319}$ .

**Problem 2.** Do 31.8-3 in the book.

*Solution:* If  $\gcd(x-1, n) = n$ , then  $x \equiv 1 \pmod{n}$ . If  $\gcd(x-1, n) = 1$ , then from  $n \mid (x-1)(x+1)$  follows that  $n$  divides  $x+1$  and  $x \equiv -1 \pmod{n}$ . This shows that  $\gcd(x-1, n)$  is a nontrivial divisor of  $n$ . The prove that  $\gcd(x+1, n)$  is a nontrivial divisor goes similarly.

**Problem 3.** Do 31.9-1 in the book.

*Solution:* Consider the sequence modulo 73 as in figure 31.7(c) in the book. The first time, the value of  $y$  lies within the loop is when  $y$  is set to  $x_8 = 814$ . Then we have  $y \equiv 11 \pmod{73}$ . The loop modulo 73 has length four. We get  $x_{12} = 84 \equiv 11 \pmod{73}$  again. The computation of  $\gcd(y - x_1, 1387)$  (where  $y$  is set equal to  $x_8$ ) yields  $\gcd(814 - 84, 1387) = 73$ . This is the first time that the divisor 73 will be printed. (the divisor 19 will be printed earlier).

**Problem 4.** \* Do 31.8-2 in the book.

*Solution:* From the formula for  $\phi(n)$  it is clear that

$$\lambda(n) = \text{lcm}(\phi(p_1^{e_1}), \dots, \phi(p_r^{e_r})) \text{ divides } \phi(n) = \phi(p_1^{e_1}) \cdots \phi(p_r^{e_r}).$$

Suppose that  $a$  is relatively prime to  $n$ . Then  $a$  is relatively prime to  $p_i^{e_i}$  for all  $i$  and

$$a^{\lambda(n)} \equiv 1 \pmod{p_i^{e_i}}$$

because  $\lambda(n)$  is divisible by  $\phi(p_i^{e_i})$ . It follows that  $a^{\lambda(n)} - 1$  is divisible by  $p_1^{e_1} p_2^{e_2} \cdots p_r^{e_r} = n$  and  $a^{\lambda(n)} \equiv 1 \pmod{n}$ . Suppose that  $n = p_1^{e_1} p_2^{e_2} \cdots p_r^{e_r}$  is Carmichael. Now  $p_i^{e_i-1}$  divides  $\phi(p_i^{e_i})$ ,  $\lambda(n)$  and  $n - 1$ . But  $p_i^{e_i-1}$  also divides  $n$ , hence  $p_i^{e_i-1}$  divides  $n - (n - 1) = 1$ . This can only happen when  $e_i = 1$  for all  $i$ , which means that  $n$  is squarefree. Suppose that  $n = pq$  with  $p < q$

primes. If  $n$  is Carmichael then  $\lambda(n)$  divides  $n - 1$ , so  $p - 1$  and  $q - 1$  divide  $n - 1$ . Write  $n - 1 = a(q - 1)$ . Clearly  $a > p$  because  $p(q - 1) = n - p < n$ . So  $n - 1 \geq (p + 1)(q - 1)$  which implies that  $n - 1 \geq pq + q - p - 1 = n - 1 + (q - p)$ . We conclude that  $p \geq q$  which contradicts our assumption that  $p < q$ .