A Proof-Theoretic View of Basic Intuitionistic Conditional Logic

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Conditional logic

- Account of conditionals 'if φ , then ψ ' different from material implication
- Reasoning with conditional sentences:

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\varphi \mapsto \psi 'If \varphi were the case, then \psi would be the case' \varphi \Leftrightarrow \psi 'If \varphi were the case, then \psi might be the case'
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- ► Tipically, extension of classical logic: CPL + would/might-principles
- Many and diverse applications: counterfactuals, conditional obligations, preferences, ceteris paribus, comparative similarity
- Several alternative semantics: selection functions, preferential models, sphere models
- ► Big family of logics



and many more

Basic conditional logic

- CK introduced by Chellas 1975
- Conditional logic core: few fundamental principles shared by most conditional logics
- Congruence on the left

$$\mathsf{RA}_{\square} \frac{\varphi \leftrightarrow \rho}{(\varphi \bowtie \psi) \leftrightarrow (\rho \bowtie \psi)}$$

Normality on the right

$$\begin{array}{ll} \mathsf{RC}_{\square} & \dfrac{\psi \leftrightarrow \chi}{\left(\varphi \boxminus \psi\right) \leftrightarrow \left(\varphi \boxminus \chi\right)} & \text{(congruence)} \\ \mathsf{CM}_{\square} & \left(\varphi \boxminus \psi \land \chi\right) \rightarrow \left(\varphi \boxminus \psi\right) \land \left(\varphi \boxminus \chi\right) & \text{(monotonicity)} \\ \mathsf{CC}_{\square} & \left(\varphi \boxminus \psi\right) \land \left(\varphi \boxminus \chi\right) \rightarrow \left(\varphi \boxminus \psi \land \chi\right) & \text{(agglomeration)} \\ \mathsf{CN}_{\square} & \varphi \boxminus \top & \text{(necessitation)} \end{array}$$

- ► Smallest normal conditional logic
- ► CK is to conditional logics as K is to modal logics
- Basic system to be extended with further conditional principles

Intuitionistic conditional logic

- Can intuitionistic logic support reasoning on conditionals?
- Which classical systems have intuitionistic counterparts?
- Analogy with intuitionistic modal logic

Basic intuitionistic conditional logic

Initiated by Weiss 2019

- Intuitionistic variant of Chellas' CK
- ▶ Defined in the language with only □→
- Axiomatically, IPL + □→-axioms of CK
- Kripke semantics ('Weiss models' in our paper) combining intutionistic order and Chellas 'selection function' (in Bozic-Dosen style)
- Called ConstCK^{□→} in our paper

Continued by Ciardelli & Liu 2020

- Alternative semantics with general frames
- Further extensions

Dufty & de Groot 2025

Very recently, further analysis and extensions

Two problems raised by Weiss 2019

1. Extensions of CK without intuitionistic counterparts

- ▶ Stalnaker's logic C2, classical system stronger than CK
- ► Characteristic axiom: conditional excluded middle CEM_□ $(\varphi \mapsto \psi) \lor (\varphi \mapsto \neg \psi)$
- ► CEM_□ + conditional modus ponens $(\varphi \mapsto \psi) \to (\varphi \to \psi) \Rightarrow_{\mathsf{IPL}} \varphi \lor \neg \varphi$
- Weiss' claim: Intuitionistic variant not possible

2. Addition of the might conditional ⋄→

- ▶ ⋄→ behaviour not determined by □→
- Many different relations bewteen □→ and ⋄→ are possible
- ▶ Analogous to □, ◊ of intuitionistic modal logics
- ▶ Left open by Weiss 2019

Tackling the problems

1. Ciardelli & Liu 2020: Never give up

Intuitionistic variant of Stalnaker's C2 possible considering the classically but not intuitionistically equivalent axiom of conditional determinacy $(\varphi \bowtie \psi \lor \chi) \to (\varphi \bowtie \psi) \lor (\varphi \bowtie \chi)$

2. Olkhovikov's 2024 proposal: IntCK

- ▶ Intuitionistic variant of CK with both □→ and ⋄→
- ► Fischer Servi style semantics
- But: extension not conservative: validates

 → principles that are not valid in ConstCK

 →
- ▶ Analogus to intuitionistic modal logics: □-fragment of IK vs iK (Das & Marin 2023)

Our contribution

Basic intuitionistic conditional logic through the lens of proof-theory

- ► Make order: systematic view of the existing systems
- ightharpoonup Conservative extension of ConstCK $^{\square \rightarrow}$ with \Leftrightarrow
- ▶ Byproduct: Aternative intuitionistic variant of CEM□

Why proof theory...

... instead of axioms?

- Axioms equivalent over CPL but not over IPL. How to select?
- Example of CEM_□: Weiss 2019 vs. Ciardelli & Liu 2020 vs. alternative formulations
- ► Which axioms for ⇔?
- Which interactions between □→ and ⋄→?
- ... lost in the intuitionistic conditional jungle



Why proof theory...

instead of semantics?

- A possible way: follow the heredity property. But:
- Not necessarily conservative over ConstCK (e.g. Olkhovikov's IntCK)
- Many different ways to satisfy it (think about intu. variants of modal logic K)
- ... still easy to get lost



Proof-theoretic compass to orientate in the intuitionistic conditional jungle

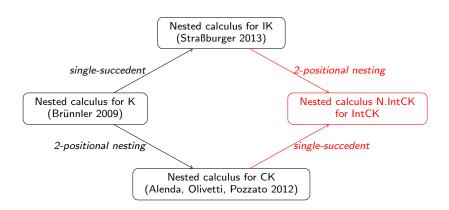


The direction we follow: Multi- vs. single-succedent sequents relation

Multi- vs. single-succedent relation

- ▶ Understand the relations bewteen CK and existing intu. counterparts of it
- ▶ Define further intuitionistic counterparts and extensions
- ▶ Known to hold between CPL and IPL since Gentzen 1935
- Applied to intuitionistic modal logics (Straßburger 2013, Das & Marin 2023, ...)

Nested calculus for IntCK



Nested calculus for IntCK

- From Straßburger 2013: treatment of intuitionistic base, input and output polarities
- ► From Alenda, Olivetti, Pozzato 2012: indexing nested components with formulas to evaluate conditional operators
- Conditional rules:

- Sound and complete calculus for IntCK
- ► Admissibility of structural rules and cut elimination

Sequent calculus for ConstCK^{□→}

Sequent calculus for CK (Pattinson & Schröder 2011)

single-succedent

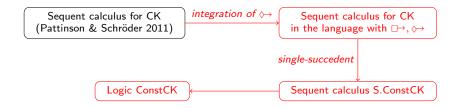
Sequent calculus S.ConstCK $^{\square \rightarrow}$ for ConstCK $^{\square \rightarrow}$

Conditional rule:

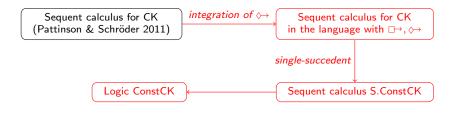
$$\square \rightarrow \frac{\{\varphi \Leftrightarrow \rho_i\}_{i \leq n} \qquad \sigma_1, ..., \sigma_n \Rightarrow \psi}{\Gamma, \rho_1 \square \rightarrow \sigma_1, ..., \rho_n \square \rightarrow \sigma_n \Rightarrow \varphi \square \rightarrow \psi}$$

- ▶ Sound and complete calculus for ConstCK
- ► Admissibility of structural rules and cut elimination

Adding \leftrightarrow to ConstCK \longrightarrow : the logic ConstCK



Adding \leftrightarrow to ConstCK \longrightarrow : the logic ConstCK



Conditional rules:

Adding \leftrightarrow to ConstCK \longrightarrow : the logic ConstCK

Properties of S.ConstCK and ConstCK

- Admissibility of structural rules and cut elimination
- ► Equivalent axiomatic system: ConstCK:

$$\begin{split} & \mathsf{IPL} + \mathsf{congruence} \ \mathsf{rules} \ \mathsf{of} \ \mathsf{left} \ \mathsf{and} \ \mathsf{right} \ + \\ & \mathsf{CM}_{\square} \ (\varphi \boxminus \psi \land \chi) \to (\varphi \boxminus \psi) \land (\varphi \boxminus \chi) \\ & \mathsf{CC}_{\square} \ (\varphi \boxminus \psi) \land (\varphi \boxminus \chi) \to (\varphi \boxminus \psi \land \chi) \\ & \mathsf{CN}_{\square} \ \varphi \boxminus \top \\ & \mathsf{CN}_{\Diamond} \ \neg (\varphi \diamondsuit \bot) \\ & \mathsf{CK}_{\Diamond} \ (\varphi \boxminus (\psi \to \chi)) \to ((\varphi \diamondsuit \psi) \to (\varphi \diamondsuit \chi)) \end{split}$$

- ► Soundness and completeness w.r.t. Constructive Chellas models
- ▶ \diamondsuit -free fragment of ConstCK amounts to ConstCK ::

For
$$\varphi \in \mathcal{L}^{\square \rightarrow}$$
, ConstCK $\vdash \varphi$ iff ConstCK $^{\square \rightarrow} \vdash \varphi$.

Why these names for intuitionistic conditional logics

Int/const modal logics embeddable into int conditional logics via the translation

$$(\Box \varphi)^t := \top \Longrightarrow \varphi^t \quad (\Diamond \varphi)^t := \top \Longleftrightarrow \varphi^t$$

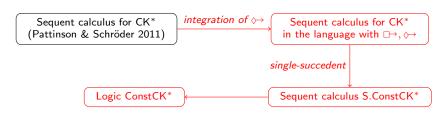
In particular:

Extensions of ConstCK

We study intuitionistic counterparts of the extensions of CK with:

- ▶ Identity: $ID_{\square} \varphi \longrightarrow \varphi$
- ▶ Conditional modus ponens: $MP_{\square} (\varphi \square \!\!\!\! \rightarrow \psi) \rightarrow (\varphi \rightarrow \psi)$
- ► Conditional excluded middle: CEM_□ $(\varphi \mapsto \psi) \lor (\varphi \mapsto \neg \psi)$

That is, the conditional logics covered in Pattinson & Schröder 2011



Extensions of ConstCK

Conditional rules:

$$\Box^{\mathrm{id}} \frac{\{\varphi \Leftrightarrow \rho_{i}\}_{i \leq n} \quad \sigma_{1}, ..., \sigma_{n}, \varphi \Rightarrow \psi}{\Gamma, \rho_{1} \boxminus \sigma_{1}, ..., \rho_{n} \boxminus \sigma_{n} \Rightarrow \varphi \boxminus \psi}$$

$$\Diamond^{\mathrm{id}} \frac{\{\varphi \Leftrightarrow \rho_{i}\}_{i \leq n} \quad \varphi \Leftrightarrow \eta \quad \sigma_{1}, ..., \sigma_{n}, \varphi, \psi \Rightarrow \vartheta}{\Gamma, \rho_{1} \boxminus \sigma_{1}, ..., \rho_{n} \boxminus \sigma_{n}, \varphi \Leftrightarrow \psi \Rightarrow \eta \Leftrightarrow \vartheta}$$

$$\Diamond^{\mathrm{id}} \frac{\{\varphi \Leftrightarrow \rho_{i}\}_{i \leq n} \quad \varphi \Leftrightarrow \eta \quad \sigma_{1}, ..., \sigma_{n}, \varphi, \psi \Rightarrow \varphi}{\Gamma, \rho_{1} \boxminus \sigma_{1}, ..., \rho_{n} \boxminus \sigma_{n}, \varphi \Leftrightarrow \psi \Rightarrow \Delta}$$

$$\mathsf{mp}_{\square} \frac{\Gamma, \varphi \boxminus \psi \Rightarrow \varphi \quad \Gamma, \varphi \boxminus \psi, \psi \Rightarrow \Delta}{\Gamma, \varphi \boxminus \psi \Rightarrow \Delta} \quad \mathsf{mp}_{\Diamond} \frac{\Gamma \Rightarrow \varphi \quad \Gamma \Rightarrow \psi}{\Gamma \Rightarrow \varphi \Leftrightarrow \psi}$$

$$\Diamond^{\mathrm{cem}} \frac{\{\varphi \Leftrightarrow \rho_{i}\}_{i \leq n} \quad \{\varphi \Leftrightarrow \xi_{j}\}_{j \leq k} \quad \varphi \Leftrightarrow \eta \quad \sigma_{1}, ..., \sigma_{n}, \chi_{1}, ..., \chi_{k}, \psi \Rightarrow \vartheta}{\Gamma, \rho_{1} \boxminus \sigma_{1}, ..., \rho_{n} \boxminus \sigma_{n}, \xi_{1} \Leftrightarrow \chi_{1}, ..., \xi_{k} \Leftrightarrow \chi_{k}, \varphi \Leftrightarrow \psi \Rightarrow \eta \Leftrightarrow \vartheta}$$

$$\Diamond^{\mathrm{cem}} \frac{\{\varphi \Leftrightarrow \rho_{i}\}_{i \leq n} \quad \{\varphi \Leftrightarrow \xi_{j}\}_{j \leq k} \quad \sigma_{1}, ..., \sigma_{n}, \chi_{1}, ..., \chi_{k}, \psi \Rightarrow \vartheta}{\Gamma, \rho_{1} \boxminus \sigma_{1}, ..., \rho_{n} \boxminus \sigma_{n}, \xi_{1} \Leftrightarrow \chi_{1}, ..., \xi_{k} \Leftrightarrow \chi_{k}, \varphi \Leftrightarrow \psi \Rightarrow \Delta}$$

Extensions of ConstCK

Equivalent axiomatics systems characterised by axioms

- Constructive Chellas models for S.ConstCKID, S.ConstCKMP, S.ConstCKMPID: soundness and completeness
- \$\rightarrow\$-free fragments of S.ConstCKID, S.ConstCKMPID coincide with the logics by Weiss 2019
- ▶ A different counterpart for logics with CEM: no \Box → axiom, but CEM \Diamond

Let's stop here for today

Where we arrived using the proof-theoretic compass

- Clear proof-theoretic relations between CK and its intuitionistic counterparts
- ▶ New intuitionistic counterpart of CK conservative over ConstCK →
- Definition of some extensions
- Semantical characterisation of most of them

A cardinal direction is not the same as a GPS track

There still are details to be checked and choices to be made

- Schröder, Pattinson, Hausmann 2010: calculi for logics with cautious monotonicity. Is the same strategy still applicable?
- Different ways of being single-succedent: e.g. modal logics CK vs. WK
- Extensions of nested calculi

Just an example of a possible fork

CEM rule by Alenda, Olivetti, Pozzato 2012:

$$_{cem}\frac{\varphi^{\bullet},\eta^{\circ}\quad\eta^{\bullet},\varphi^{\circ}\quad\Gamma\{[\varphi:\Sigma,\Delta],[\eta:\Delta]\}}{\Gamma\{[\varphi:\Sigma],[\eta:\Delta]\}}$$

- ► Single-succedent intuitionistic base (Straßburger 2013 style):
 - ► CEM $_{\Diamond}$ ($\varphi \diamondsuit \psi$) \wedge ($\varphi \diamondsuit \chi$) \rightarrow ($\varphi \diamondsuit \psi \wedge \chi$) derivable
 - ► CD $(\varphi \mapsto \psi \lor \chi) \to (\varphi \mapsto \psi) \lor (\varphi \mapsto \chi)$ not derivable
- Multi-succedent intuitionistic base (Kuznets & Straßburger 2019 style):
 - ► CEM $_{\Diamond}$ ($\varphi \diamondsuit \psi$) \wedge ($\varphi \diamondsuit \chi$) \rightarrow ($\varphi \diamondsuit \psi \wedge \chi$) derivable
 - ► CD $(\varphi \mapsto \psi \lor \chi) \to (\varphi \mapsto \psi) \lor (\varphi \mapsto \chi)$ derivable
- Different intuitionistionizations give rise to different counterparts
- Many questions open
 - Cut admissibility?
 - Semantics?
 - Relations with classical counterparts?
 - **.**..



The exploration of the intuitionistic conditional jungle is only just beginning

Thank you!