

Canonical

Simplify, Monomorphize, Destruct

```
theorem add_assoc (a b c : Nat) : (a + b) + c = a + (b + c) :=  
  by canonical
```

```
theorem add_comm (a b : Nat) : a + b = b + a :=  
  by canonical 60
```

```
theorem pow_add (a m n : Nat) : a^(m + n) = a^m * a^n  
  by canonical [mul_one, mul_assoc]
```

"Why can't Canonical output tactics?"

Tactics do not commute with metavariable assignments.

?n : Nat

0

?0 + x = x

simp

Tactics do not produce constraints on metavariables.

Tactics require interaction with the Lean server.

Tactics do not have soundness and completeness guarantees.

s i m p

Essential for reasoning about equality and bi-implication.

*“all it does is repeatedly replace (or rewrite) subterms of the form A by B ,
for all applicable facts of the form $A = B$ or $A \leftrightarrow B$ ”*

Type theorists have a tool for this too: **definitional reduction rules**.

$\text{Nat.rec } m \ a \ b \ \text{zero} \mapsto a$
 $\text{Nat.rec } m \ a \ b \ n.\text{succ} \mapsto b \ n \ (\text{Nat.rec } m \ a \ b \ n)$

$\text{Prod.fst } \langle a, b \rangle \mapsto a$
 $\text{Prod.snd } \langle a, b \rangle \mapsto b$

$(\text{let } t := M; t) \mapsto M$

General Reduction Rule System

Canonical now accepts custom definitional reduction rules: $0 + n \mapsto n$

Reduction rules are applied **pervasively**, as though `simp` were always applied.

Instead of having Canonical generate `simp` invocations (impossibly difficult) we now only need to reason about terms in the quotient of `simp` lemmas!

We forbid Canonical from writing reducible expressions, like $0 + n$

canonical

We register Lean's definitional equalities as reduction rules.

We register equation compiler lemmas as reduction rules.

We register relevant **simp** lemmas as reduction rules.

```
Nat.add : (a.277 : Nat) → (a.278 : Nat) → Nat := Canonical.Pi (Nat
  (λ x.350 ↪ Nat) (PProd.fst (Canonical.Pi Nat (λ a.293 ↪ Nat))
    (Nat.rec (λ t.294 ↪ Sort) PUnit (λ n.295 n_ih.296 ↪ PProd
      (Canonical.Pi Nat (λ a.300 ↪ Nat)) n_ih.296) a.278) (Nat.rec (λ
        t.301 ↪ PProd (Canonical.Pi Nat (λ a.303 ↪ Nat)) (Nat.rec (λ
          t.304 ↪ Sort) PUnit (λ n.305 n_ih.306 ↪ PProd (Canonical.Pi Nat
            (λ a.308 ↪ Nat)) n_ih.306) t.301)) (PProd.mk (Canonical.Pi Nat (λ
              a.310 ↪ Nat)) PUnit (Canonical.Pi.mk Nat (λ a.317 ↪ Nat) (λ a.318
                ↪ a.318)) PUnit.unit) (λ n.326 n_ih.327 ↪ PProd.mk (Canonical.Pi
                  Nat (λ a.329 ↪ Nat)) (PProd (Canonical.Pi Nat (λ a.331 ↪ Nat))
                    (Nat.rec (λ t.332 ↪ Sort) PUnit (λ n.333 n_ih.334 ↪ PProd
                      (Canonical.Pi Nat (λ a.336 ↪ Nat)) n_ih.334) n.326))
                      (Canonical.Pi.mk Nat (λ a.338 ↪ Nat) (λ a.339 ↪ Nat.succ
                        (Canonical.Pi.f Nat (λ x.349 ↪ Nat) (PProd.fst (Canonical.Pi Nat
                          (λ a.343 ↪ Nat)) (Nat.rec (λ t.344 ↪ Sort) PUnit (λ n.345
                            n_ih.346 ↪ PProd (Canonical.Pi Nat (λ a.348 ↪ Nat)) n_ih.346)
                              n.326) n_ih.327) a.339))) n_ih.327) a.278)) a.277
```

We proactively identify premises that can be registered as reduction rules.

Canonical carefully tracks the usage of these reduction rules.

Non-definitional reduction rules are reconstructed as **simp** invocations.

Nat.add : (a.412 : Nat) → (a.413 : Nat) → Nat := [
 Nat.add x.409 Nat.zero ↪ x.409,
 Nat.add x.414 (Nat.succ b.415) ↪ Nat.succ (Nat.add x.414 b.415)

Unsoundness

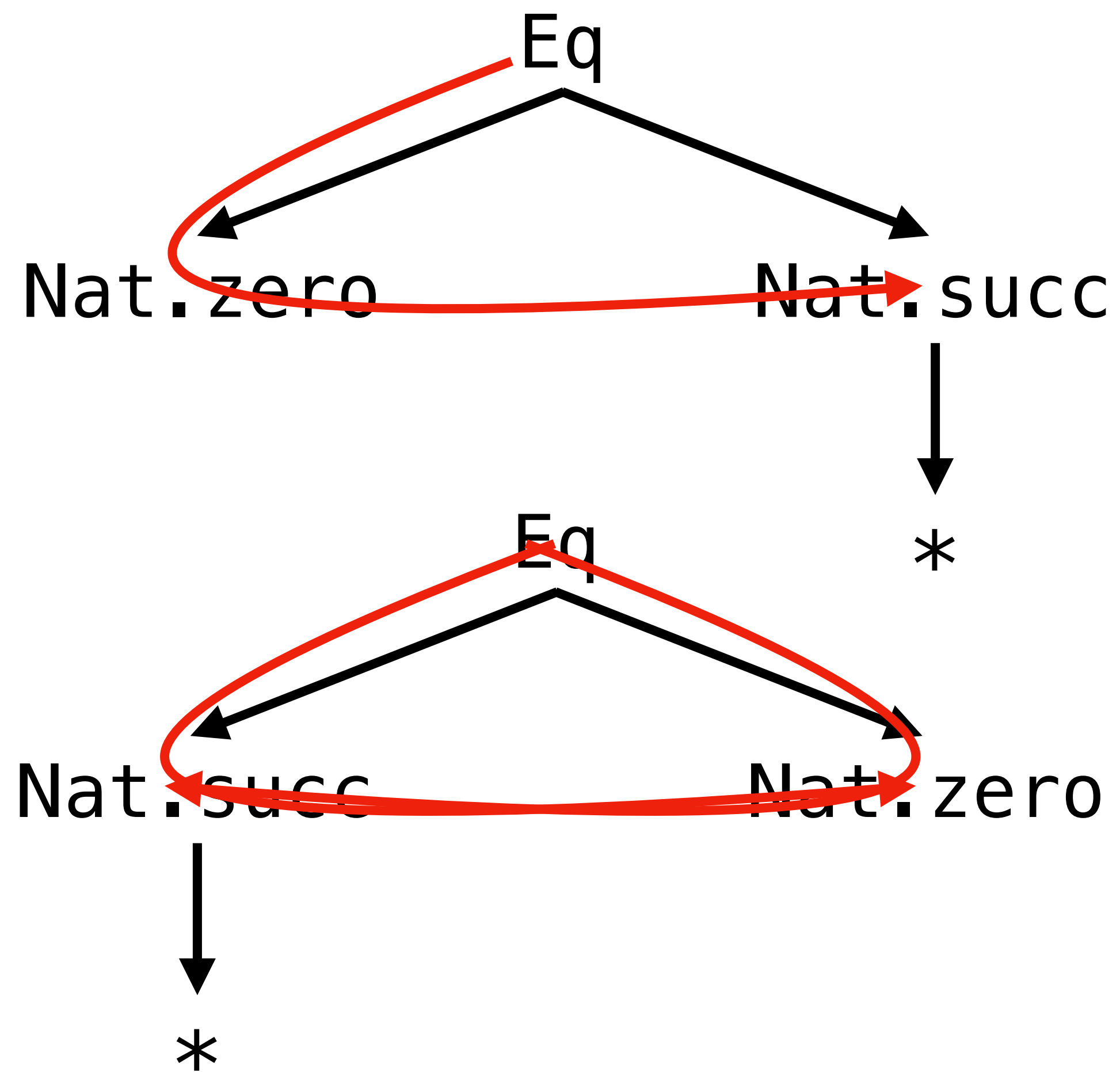
Definitional equality is stronger than `simp`.

```
variable (n : Nat) (i : Fin n) (P : {n : Nat} → Fin n → Prop)

example (h : P i) : ∃ (j : Fin (0 + n)), P j :=
  by simp only [Nat.zero_add] <;> exact Exists.intro i h
```

Uncommon, but you can always `–simp`.

Pattern Matching



Implementation

Lexicographic path ordering

Forbidding redexes syntactically

Should ?X be a match?

$$\text{Proj.fst } ?X \equiv a$$

Enrich reasoning with type-theoretic structure, and it becomes free.

Monomorphize

`example (a b : \mathbb{N}) : a + b = b + a := by canonical [add_comm]`

```
Nat.add : (a.496 : Nat) → (a.497 : Nat) → Nat := [Nat.add x.493 Nat.zero ↪ x.493, Nat.add x.498 (Nat.succ b.499) ↪  
  Nat.succ (Nat.add x.498 b.499)]
```

```
instAddNat : Add Nat := [instAddNat ↪ Add.mk Nat (λ a.376 a.377 ↪ Nat.add a.376 a.377)]
```

```
Add : (α.348 : Sort) → Sort := []
```

```
Add.mk : (α.361 : Sort) → (add.362 : (a.365 : α.361) → (a.366 : α.361) → α.361) → Add α.361 := []
```

```
Add.add : (α.371 : Sort) → (self.372 : Add α.371) → (a.373 : α.371) → (a.374 : α.371) → α.371 := [Add.add * (Add.mk *  
  field) arg0 arg1 ↪ field arg0 arg1]
```

```
HAdd : (α.313 : Sort) → (β.314 : Sort) → (γ.315 : Sort) → Sort := []
```

```
HAdd.mk : (α.302 : Sort) → (β.303 : Sort) → (γ.304 : Sort) → (hAdd.305 : (a.308 : α.302) → (a.309 : β.303) → γ.304) →  
  HAdd α.302 β.303 γ.304 := []
```

```
HAdd.hAdd : (α.322 : Sort) → (β.323 : Sort) → (γ.324 : Sort) → (self.325 : HAdd α.322 β.323 γ.324) → (a.326 : α.322)  
  → (a.327 : β.323) → γ.324 := [HAdd.hAdd * * * (HAdd.mk * * * field) arg0 arg1 ↪ field arg0 arg1]
```

```
instHAdd : (α.345 : Sort) → (inst.346 : Add α.345) → HAdd α.345 α.345 α.345 := [instHAdd α.341 inst.342 ↪ HAdd.mk  
  α.341 α.341 α.341 (λ a.349 a.350 ↪ Add.add α.341 inst.342 a.349 a.350)]
```

```
AddCommMagma : (G.845 : Sort) → Sort := []
```

```
AddCommMagma.mk : (G.863 : Sort) → (toAdd.864 : Add G.863) → (add_comm.865 : (a.868 : G.863) → (b.869 : G.863) → Eq  
  G.863 (HAdd.hAdd G.863 G.863 G.863 (instHAdd G.863 toAdd.864) a.868 b.869) (HAdd.hAdd G.863 G.863 G.863 (instHAdd  
  G.863 toAdd.864) b.869 a.868)) → AddCommMagma G.863 := []
```

```
AddCommMagma.toAdd : (G.888 : Sort) → (self.889 : AddCommMagma G.888) → Add G.888 := [AddCommMagma.toAdd *  
  (AddCommMagma.mk * field *) ↪ field]
```

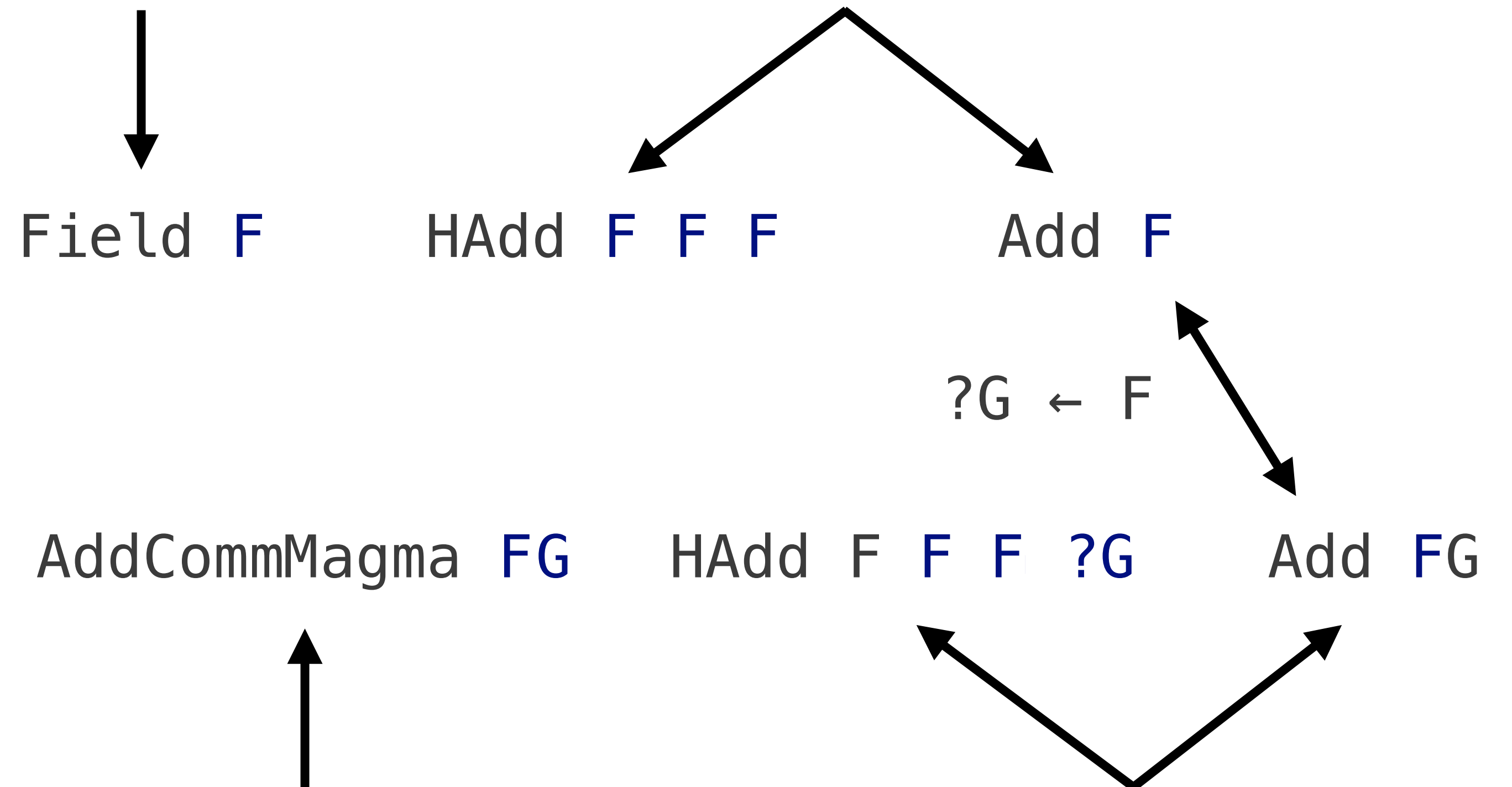
```
add_comm : (G.840 : Sort) → (inst.841 : AddCommMagma G.840) → (a.842 : G.840) → (b.843 : G.840) → Eq G.840 (HAdd.hAdd  
  G.840 G.840 G.840 (instHAdd G.840 (AddCommMagma.toAdd G.840 inst.841)) a.842 b.843) (HAdd.hAdd G.840 G.840 G.840  
  (instHAdd G.840 (AddCommMagma.toAdd G.840 inst.841) b.843 a.842) := []
```

Monomorphize

```
HAdd_hAdd_0 : (a.331 : Nat) → (a.332 : Nat) → Nat := [  
  HAdd_hAdd_0 n.796 Nat.zero ↪ n.796,  
  HAdd_hAdd_0 n.822 (Nat.succ m.823) ↪ Nat.succ (HAdd_hAdd_0 n.822 m.823)  
]  
add_comm_0 : (a.457 : Nat) → (b.458 : Nat) →  
  Eq Nat (HAdd_hAdd_0 a.457 b.458) (HAdd_hAdd_0 b.458 a.457) := []
```

Implementation

example {F : Type} [Field F] (x y : F) : x + y = x + y



theorem add_comm (a b : F) : a + b = b + a

Implementation

```
@HAdd.hAdd Nat Nat Nat (@instHAdd Nat instAddNat) 1 2
```



```
@HAdd.hAdd _ _ _ (@instHAdd _ instAddNat) _ _
```



```
@HAdd.hAdd Nat Nat Nat (@instHAdd Nat instAddNat) _ _
```



```
HAdd_hAdd_0 : Nat → Nat → Nat := ...
```

Destruct

$$\exists x > 0, \forall y, y < x$$

```
Canonical.Pi : (A.2214 : Sort) → (B.2215 : (a.2217 : A.2214) → Sort) → Sort := []
Canonical.Pi.f : (A.2244 : Sort) → (B.2245 : (a.2249 : A.2244) → Sort) → (self.2246 :
  Canonical.Pi A.2244 (λ a.2252 ↦ B.2245 a.2252)) → (a.2247 : A.2244) → B.2245 a.2247 :=
  [Canonical.Pi.f * * (Canonical.Pi.mk * * field) arg0 ↦ field arg0]
Canonical.Pi.mk : (A.2223 : Sort) → (B.2224 : (a.2228 : A.2223) → Sort) → (f.2225 : (a.2229 :
  A.2223) → B.2224 a.2229) → Canonical.Pi A.2223 (λ a.2233 ↦ B.2224 a.2233) := []
```

```
And : (a.2142 : Sort) → (b.2143 : Sort) → Sort := []
And.right : (a.2171 : Sort) → (b.2172 : Sort) → (self.2173 : And a.2171 b.2172) → b.2172 :=
  [And.right * * (And.intro * * * field) ↦ field]
And.left : (a.2160 : Sort) → (b.2161 : Sort) → (self.2162 : And a.2160 b.2161) → a.2160 :=
  [And.left * * (And.intro * * * field) ↦ field]
And.intro : (a.2148 : Sort) → (b.2149 : Sort) → (left.2150 : a.2148) → (right.2151 : b.2149)
  → And a.2148 b.2149 := []
```

```
Exists : (α.1988 : Sort) → (p.1989 : (a.1991 : α.1988) → Sort) → Sort := []
Exists.rec : (α.2044 : Sort) → (p.2045 : (a.2053 : α.2044) → Sort) → (motive.2046 : (t.2054 :
  Exists α.2044 (λ a.2057 ↦ p.2045 a.2057)) → Sort) → (intro.2047 : (w.2059 : α.2044) →
  (h.2060 : p.2045 w.2059) → motive.2046 (Exists.intro α.2044 (λ a.2067 ↦ p.2045 a.2067)
  w.2059 h.2060)) → (t.2048 : Exists α.2044 (λ a.2071 ↦ p.2045 a.2071)) → motive.2046
  t.2048 := [Exists.rec α.2036 p.2037 motive.2038 intro.2039 (Exists.intro * * w.2040 h.2041)
  ↦ intro.2039 w.2040 h.2041]
Exists.intro : (α.1997 : Sort) → (p.1998 : (a.2002 : α.1997) → Sort) → (w.1999 : α.1997) →
  (h.2000 : p.1998 w.1999) → Exists α.1997 (λ a.2006 ↦ p.1998 a.2006) := []
```

```
LT_lt_0 : (a.2203 : Nat) → (a.2204 : Nat) → Sort := []
```

```
Exists Nat (λ a.2137 ↦ And (LT_lt_0 Nat.zero a.2137) (Canonical.Pi Nat (λ a.2255 ↦ LT_lt_0
  a.2255 a.2137)))
```

→

```
a_w.2089 : Nat
a_h_left.2090 : LT_lt_0 Nat.zero
  a_w.2089
a_h_right.2091 : (y.2197 : Nat) →
  LT_lt_0 y.2197 a_w.2089
```

Implementation

Destruct constructs a **bijection** from the input type to the unpacked types.

$$\text{destruct}(A \times B) \leftrightarrow \text{destruct}(A) ++ \text{destruct}(B)$$
$$\begin{aligned} \text{destruct}(A \rightarrow B) &\leftrightarrow \\ &(\text{destruct } B).map \lambda b \mapsto (\text{destruct } A \rightarrow b) \end{aligned}$$

Dependence makes this tricky...

**When a problem is too hard,
change the problem.**