

# When GNNs Met a Word Equations Solver: Learning to Rank Equations

*FroCoS 2025*

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University of Lorraine  
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October 1, 2025



# THE CONFERENCE MORNING SESSION

DAY 1  
7:00am

Welcome, everyone!



DAY 2  
7:00am

Sorry, I haven't had  
my coffee yet...



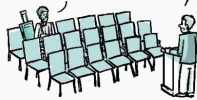
DAY 3  
7:00am

(Awkward silence)



LAST DAY  
7:00am

Thanks for attending.  
I couldn't find an  
earlier flight.



## Elements

- Empty string:  $\epsilon$
- Letters:  $a, b, \dots$
- Variables:  $X, Y, Z, \dots$

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## Easy, right...?

$$\begin{aligned} &AaAkjhohhhkokookkkokkookkhkoohohCAYZjhBkjhkkkokhkkEkkoWQB \\ &= aABBBncYECaa \end{aligned}$$

- Variables:  $A, C, Y, Z, B, E, W, Q$
- Letters:  $a, k, j, h, o, n, c$

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Z3, cvc5, Ostrich, Z3-noodler, and Woorpje cannot solve it in 300 seconds

# Challenges

- Difficult to solve
  - ▷ Decidable
  - ▷ Complexity: NP-Hard and in PSPACE
  - ▷ Practical implementation are incomplete
- Important in modeling string constraints in verification tasks
  - ▷ User inputs, string manipulations, ...

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Can machine learning help to solve conjunctive word equation problem?



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- Is it another “we used ML for [insert problem] and...”?

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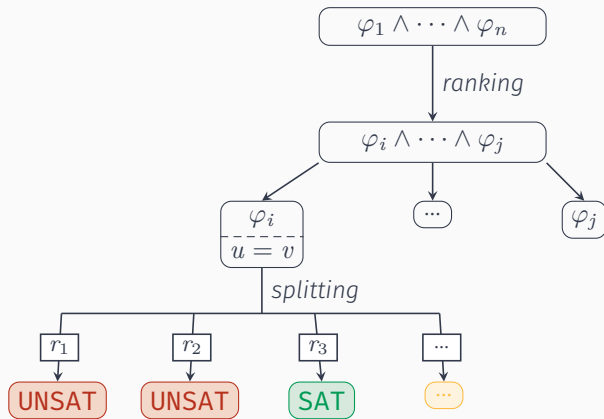
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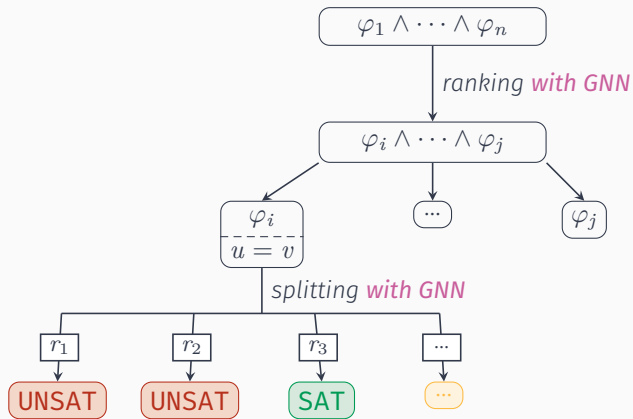
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- Is human intelligence over now?
  - ▷ ~~It depends on the human~~
  - ▷ Obviously not!

# Big Picture



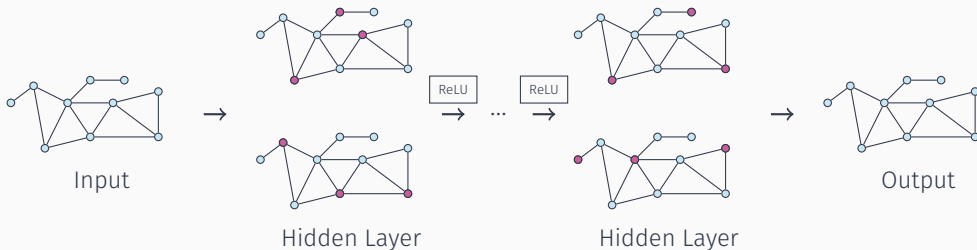
# Big Picture



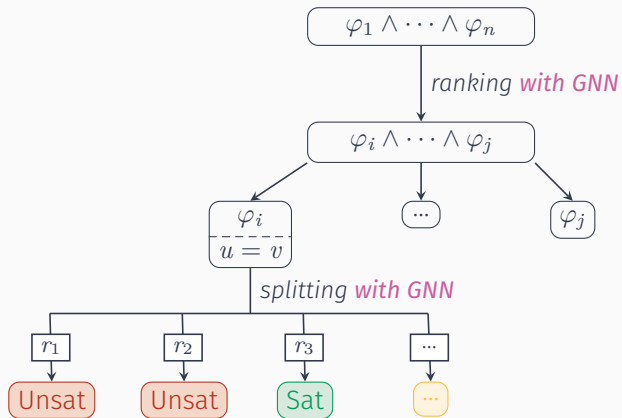


## Graph Neural Networks (GNN)

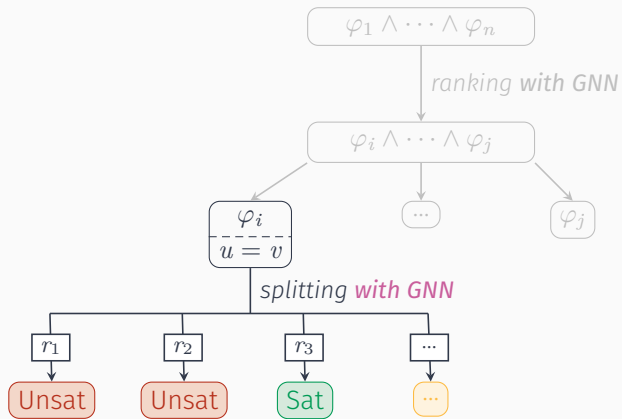
- Designed to work directly with graph-structured data
- Take graph as input, output node and graph representations
- Capture the structural features of graph



# Big Picture



# Big Picture



## Calculus

- Nielsen transformation
- Splitting rules, iterative deepening & backtracking
- Continuously simplify the first terms of both sides
- DragonLi



$$Xabb = YaZ$$

where  $a$  is a letter,  $u$  and  $v$  are terms, and  $X$  and  $Y$  are variables.

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<sup>1</sup>Abdulla, Atig, Cailler, Liang and Rüemmer. *Guiding word equation solving using graph neural networks*. International Symposium on Automated Technology for Verification and Analysis (ATVA 2024).

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$$\begin{array}{ccc} Xabb = YaZ \\ \downarrow r_1 \quad \downarrow r_2 \quad \downarrow r_3 \end{array}$$

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$$\begin{array}{c} Xabb = YaZ \\ \downarrow r_2 X \mapsto Y \end{array}$$

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$$\boxed{l} \boxed{u} = \boxed{l'} \boxed{v}$$

where  $l$ ,  $l'$ ,  $a$  and  $b$  are a letters,  $u$  and  $v$  are terms,  $X$  and  $Y$  are variables, and  $X'$  and  $Y'$  are fresh variables.



# Rules

$$\boxed{u} = \boxed{v} \quad \begin{array}{c} \swarrow \\ l = l' \\ \nwarrow \end{array} \quad \boxed{a} \boxed{u} = \boxed{a} \boxed{v}$$

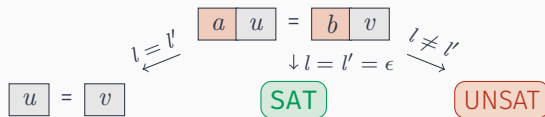
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$$\begin{array}{ccc} \boxed{u} = \boxed{v} & \begin{array}{c} l = l' \\ \swarrow \quad \searrow \end{array} & \begin{array}{c} \boxed{\epsilon} = \boxed{\epsilon} \\ \downarrow l = l' = \epsilon \\ \text{SAT} \end{array} \end{array}$$

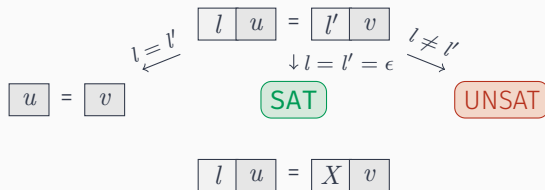
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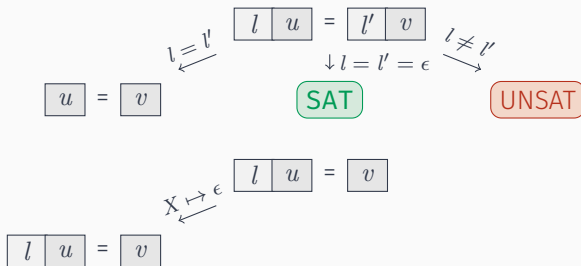
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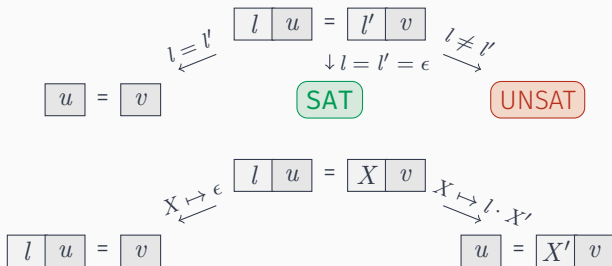
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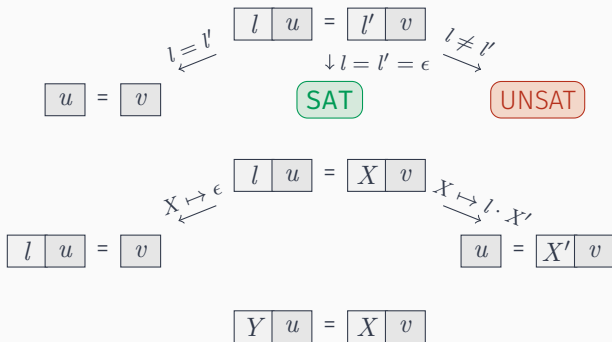
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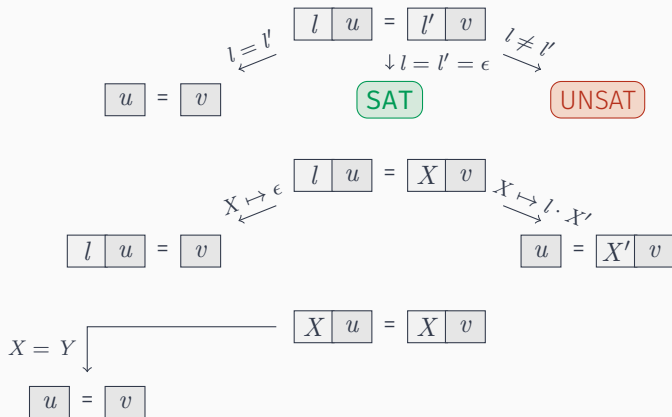
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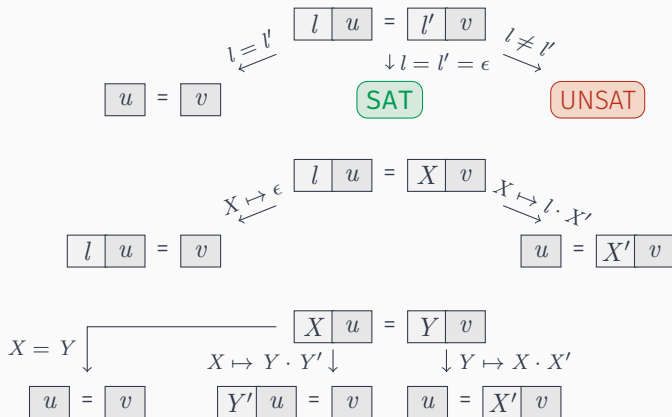
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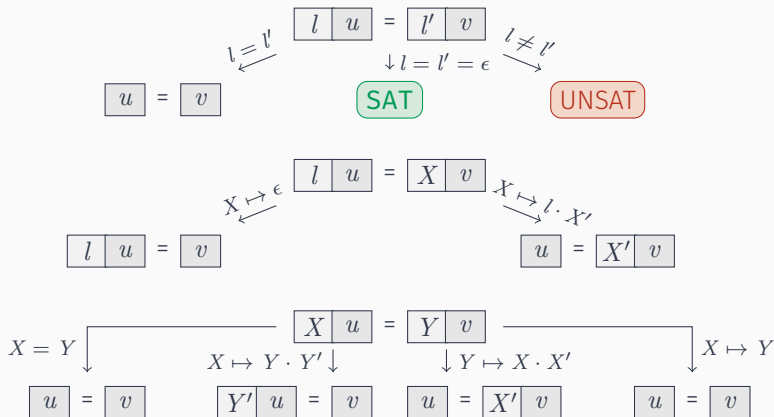


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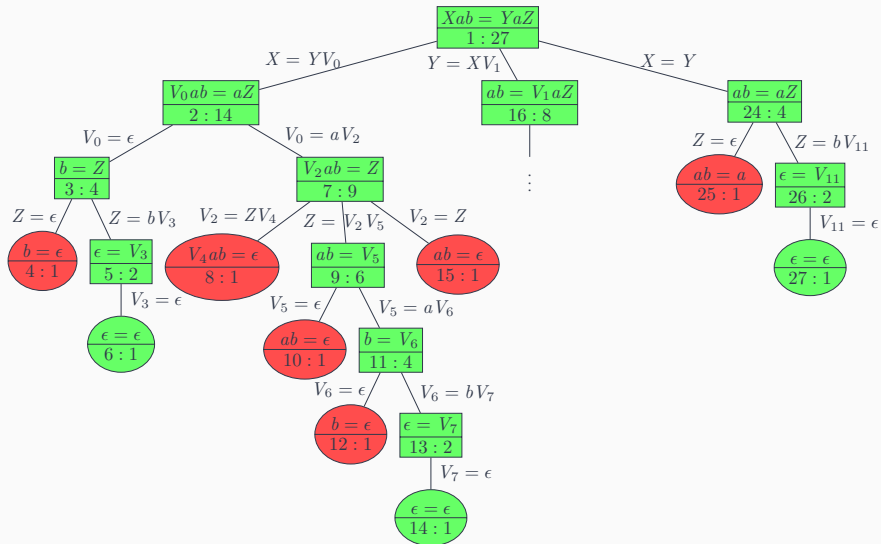
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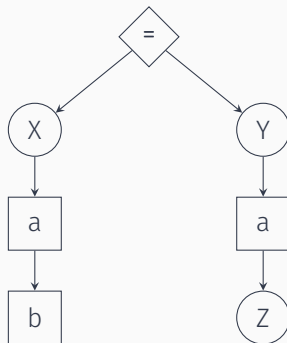
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# Here Is a Proof...



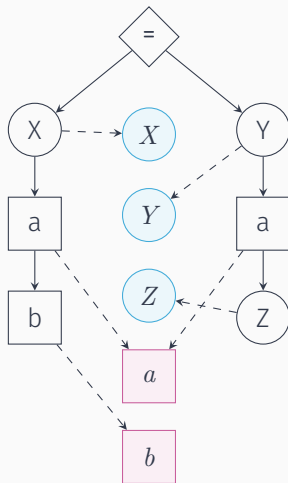
## ...and Here Is an Equation!

$$Xab = YaZ$$

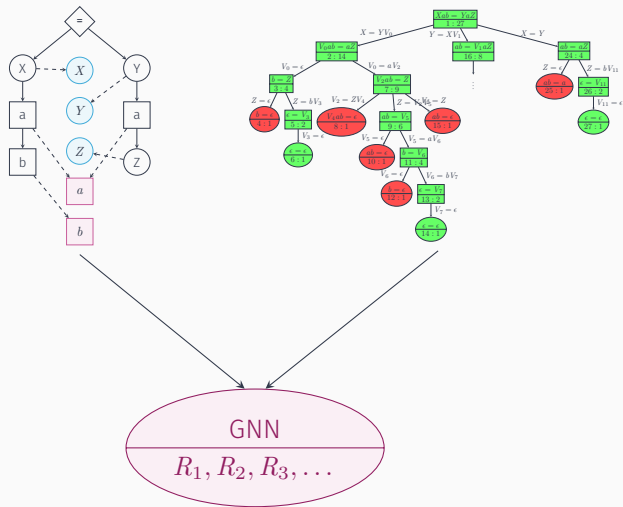


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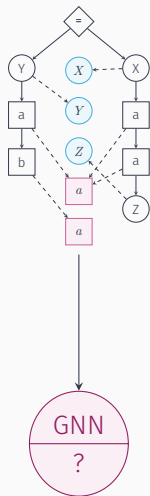
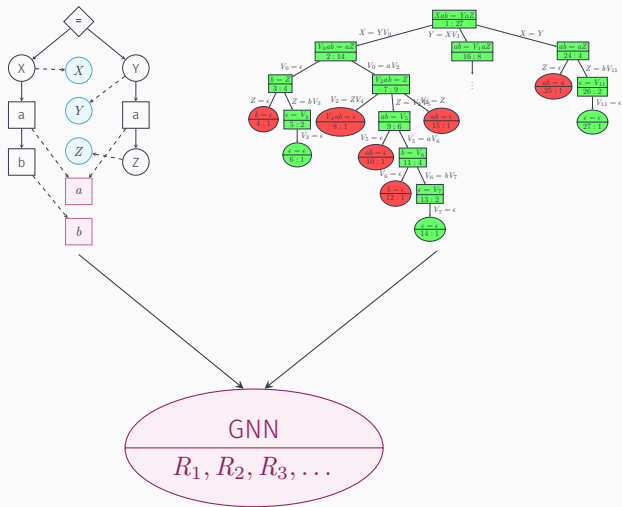
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# Better Together!



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# DragonLi, I Choose You!

- Python
- Focus on SAT
- Results comparable to state-of-the-art solvers





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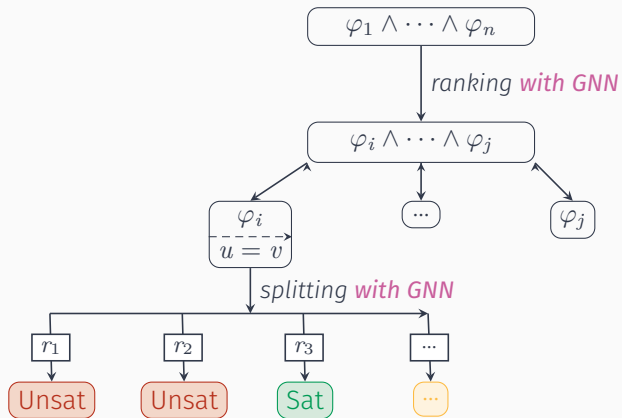
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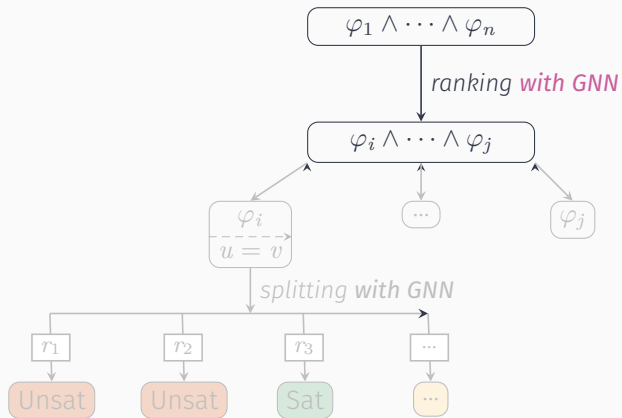
How to deal with more than one equation?

How to improve on UNSAT problems?

# Big Picture



# Big Picture



*AaAkjhhooohohCAYZjhBkjhkkkokhkkEk = aABBBncYECaa ∧ a = b*

$$AaAkjhhooohohCAYZjhBkjhkkkokhkkEk = aABBBncYECaa \wedge a = b$$

- Left first: Timeout
- Right first: UNSAT

$$AaAkjhhooohohCA YZjhBkjhkkkokhkkEk = aABBBncYECaa \wedge a = b$$

- Left first: **Timeout**
- Right first: **UNSAT**
- Basic heuristic (+ length):
  1.  $\epsilon = \epsilon$
  2.  $\epsilon = u \cdot v$  or  $u \cdot v = \epsilon$
  3.  $a \cdot u = b \cdot v$  or  $u \cdot a = v \cdot b$
  4.  $a \cdot u = a \cdot v$
  5. Otherwise

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Can a GNN do better?

## Minimal Unsatisfiable Set (MUS)

A set of equations  $U$  of a conjunctive word equation is a MUS iff the conjunction of  $U$  is unsatisfiable, and for all equations  $e \in U$ , the conjunction of  $U \setminus \{e\}$  is satisfiable.

$$a = X \wedge b = Y \wedge aa = Xa \wedge aab = XaY \wedge X = b$$



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$$\{a = X, X = b\}$$

# How to (m)Use Them?

## Training Data

- Problems not solved by **DragonLi** are given to **Z3**, **Z3-Noodler**, **cvc5**, **Ostrich** and **Woorpje**
- Extraction of MUS

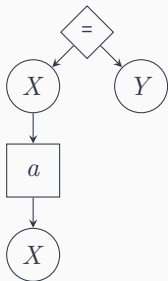
## Labelling

For a conjunctive word equation  $\phi = e_1 \wedge \dots \wedge e_n$  and  $U$  its MUS:

$$\text{label}(e_i) = \begin{cases} 1 & \text{if } e_i \in U \text{ and } |e_i| = \min(\{|e| \mid e \in U\}) \\ 0 & \text{otherwise.} \end{cases}$$

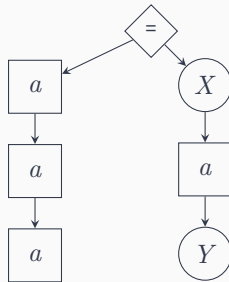
# I Wouldn't Count on It...

$$XaX = Y$$



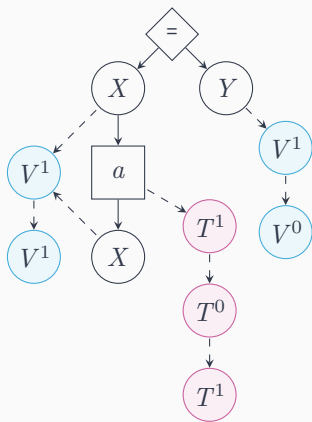
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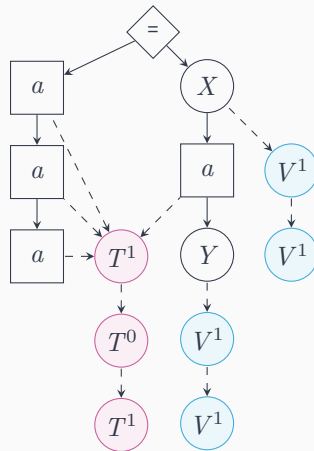
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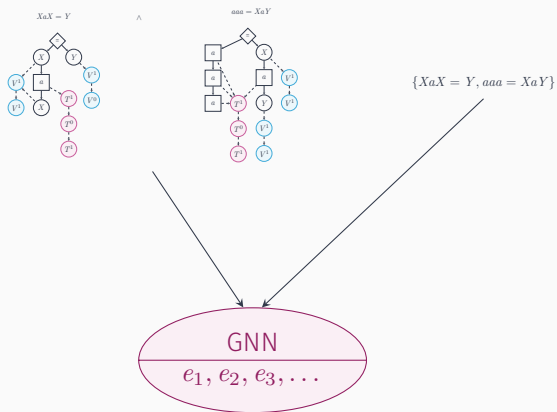
$\wedge$

$$\begin{aligned} Occ(X) &= 3 \\ Enc(X) &= 11 \\ Occ(Y) &= 2 \\ Enc(Y) &= 10 \\ Occ(a) &= 5 \\ Enc(a) &= 101 \end{aligned}$$

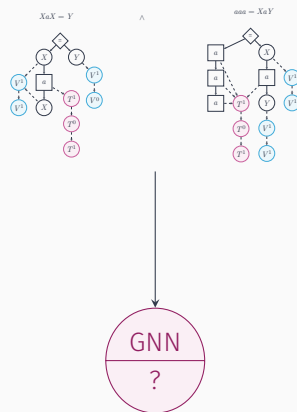
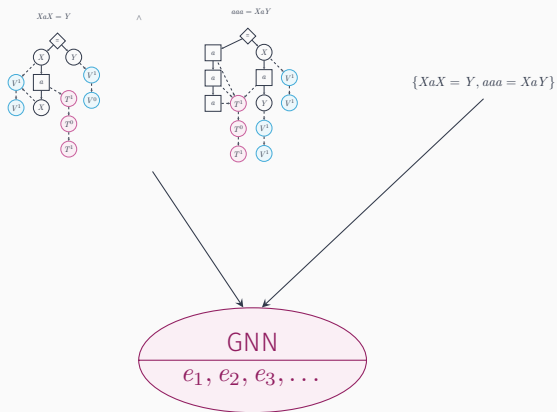
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# Better Together (bis repetita)



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  - ▷ Different task
  - ▷ Varying size of the input
  - ▷ Focus on **UNSAT**



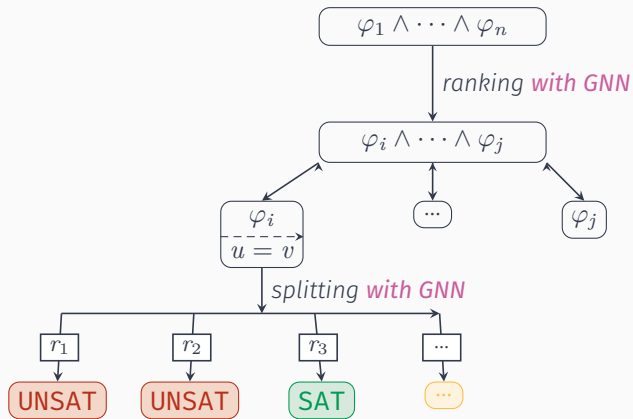
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- Why this encoding?
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- When do you call the GNN?
  - ▷ Once at the beginning with 5 equation of priority level 5
  - ▷ Limit overhead

# Big Picture



# I Need Problems...

- Real-world benchmarks:
  - ▷ SMT-LIB and Zalgivinder
  - ▷ Remove length constraint, boolean operator, and regex
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- Real-world benchmarks:
  - ▷ SMT-LIB and Zalgivinder
  - ▷ Remove length constraint, boolean operator, and regex
  - ▷ Too easy
- Artificial benchmarks:
  - ▷  $s = s$
  - ▷ At most 60 letters (among a set  $T$ )
  - ▷ Replace  $n$  times substrings on both side by  $m$  variables
  - ▷ Between 0 and 100 equations

## Linearity

A conjunctive word equation is called *linear* iff each variable appear **at most once** in each equation.

$$X = a \wedge Y = b \quad \checkmark$$

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$$XaX = aaa \wedge YaX = YbZ \quad \times$$



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- **A1** : fresh variable,  $T = 6$ ,  $n \in [0, 5]$ ,  $m \in [0, 5]$  (*linear*)
- **A2** :  $A1 + T = 26$ ,  $n \in [0, 16]$ , and  $m = 1$  (*linear*)
- **B** :  $A1 +$  No fresh variable (*non-linear*)
- **C** :  $X_n a X_n b X_{n-1} \cdots b X_1 = a X_n X_{n-1} X_{n-1} b \cdots X_1 X_1 b a a + b \mapsto A_1$  (*highly non-linear*)

# Experimental Setup

- **DragonLi**: basic, random and GNN-based
- **Z3**, **Z3-Noodler**, **cvc5**, **Ostrich**, and **Woopje**
- 300 seconds timeout
- 2 Intel Xeon E5 2630 v4 at 2.20 GHz/core and 128GB memory
- GNNs trained on NVIDIA A100 GPUs
- 60 000 problems per benchmark for training
- 1000 problems per benchmark for experiments

# Results — A1

Solver	Number of solved problems					Average solving time (split number)			
	SAT	UNSAT	UNI	CS	CU	SAT	UNSAT	CS	CU
DragonLi	24	955	0	13	678	5.6 (244.8)	6.5 (1085.3)	5.0 (94.4)	5.7 (126.3)
Random-DragonLi	22	944	0			5.6 (198.8)	6.3 (932.6)	5.6 (137.6)	5.7 (180.5)
GNN-DragonLi	24	961	0			6.1 (164.7)	7.5 (1974.8)	6.1 (96.4)	6.3 (60.5)
cvc5	24	952	1			0.5	0.6	0.1	0.3
Z3	17	960	0			8.7	0.4	1.1	0.1
Z3-Noodler	22	939	2			5.7	0.3	4.8	0.1
Ostrich	17	931	0			15.0	5.5	8.0	4.7
Woorpje	23	744	0			3.0	12.5	0.1	12.2

## Results — A2

Solver	Number of solved problems					Average solving time (split number)			
	SAT	UNSAT	UNI	CS	CU	SAT	UNSAT	CS	CU
DragonLi	59	824	0	3	0	8.5 (4233.4)	11.8 (1231.3)	4.7 (27.3)	-
Random-DragonLi	44	806	1			24.7 (29779.6)	6.2 (210.9)	4.6 (27.3)	-
GNN-DragonLi	59	836	4			8.4 (1330.6)	11.6 (1074.1)	5.9 (27.3)	-
cvc5	67	142	15			0.6	56.0	0.1	-
Z3	8	870	10			1.1	0.6	0.1	-
Z3-Noodler	22	7	1			15.4	3.8	0.4	-
Ostrich	13	18	2	-	-	24.8	38.8	8.6	-
Woorpje	0	0	0			-	-	-	-

## Results — B

Solver	Number of solved problems					Average solving time (split number)			
	SAT	UNSAT	UNI	CS	CU	SAT	UNSAT	CS	CU
DragonLi	11	805	0	4	294	4.9 (62.5)	5.2 (81.5)	4.9 (29.2)	5.3 (82.4)
Random-DragonLi	10	894	0			5.0 (58.7)	5.8 (295.2)	5.0 (27.25)	5.2 (73.1)
GNN-DragonLi	11	821	0			6.5 (65.1)	6.8 (70.0)	6.5 (28.25)	6.8 (60.2)
cvc5	12	915	0			0.1	0.6	0.1	0.7
Z3	11	859	3			0.1	0.2	0.1	0.1
Z3-Noodler	24	911	1			4.9	0.4	1.3	0.4
Ostrich	12	917	2			6.9	3.7	3.3	4.2
Woorpje	19	330	1	1		29.5	6.0	0.2	5.0
	3	62	0			65.0	28.4	0.2	223.1

## Results — C

Solver	Number of solved problems					Average solving time (split number)			
	SAT	UNSAT	UNI	CS	CU	SAT	UNSAT	CS	CU
DragonLi	2	0	0	-	-	5.1 (85.5)	-	-	-
Random-DragonLi	2	0	0	-	-	5.0 (85.5)	-	-	-
GNN-DragonLi	-	-	-	-	-	-	-	-	-
cvc5	0	909	17	-	1	-	46.9	-	17.3
Z3	1	821	12	1		0.8	1.7	0.8	0.1
Z3-Noodler	7	657	4	1		0.2	94.1	0.1	1.0
Ostrich	0	61	0	-		-	77.2	-	27.1
Woorpje	3	62	0	1		65.0	28.4	0.2	223.1

# “Why Is Your Work Good?”

- **DragonLi** performs well on linear problems
- Fewer splits
- No improvement on non-linear problems
- Need for dedicated strategies (equation length, letter count, ...)
- Re-implementation can offer improvements
- Increasing training dataset increases the number of problems solved

# This is the End!

## Contributions

- Extension of **DragonLi** to handle conjunctive word equations
- Multi-classification task to handle variable number of inputs
- MUS & new graph representation
- Good results on linear problems

## What's Next?

- Optimize GNN overhead
- Call timing
- Length constraint and regular expressions



Thank you! 😊

<https://github.com/ChenchengLiang/DragonLi>

