

1. Rekursives Flood Filling ${\tt FloodFill}(I,u,v,label)$ ▷ Recursive Version if coordinate (u,v) is within image boundaries and I(u,v)=1 then Set $I(u,v)\leftarrow label$ ${\tt FloodFill}(I,u\!+\!1,v,label)$ 11: FLOODFILL(I, u, v+1, label)12:

4er Nachbarschaft

Vorteil: Einfach zu implementieren

 ${\tt FloodFill}(I,u,v\!-\!1,label)$

FLOODFILL(I, u-1, v, label)

- Nachteil: Stack-Speicher schnell erschöpft
 - Nur für kleine Bilder geeignet

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13:

14:

```
2. Depth-First Flood Filling
 16:
        {\tt FloodFill}(I,u,v,label)
                                                                                    \quad \triangleright \  \, \mathbf{Depth\text{-}First} \  \, \mathbf{Version}
             Create an empty \mathbf{stack}\ S
             Put the seed coordinate \langle u,v\rangle onto the stack: Push(S, \langle u,v\rangle) while S is not empty do
19:
                   Get the next coordinate from the top of the stack: \langle x,y \rangle \leftarrow \operatorname{Pop}(S)
20:
21:
                    if coordinate (x,y) is within image boundaries and I(x,y)=1
                          then
                          then

Set I(x,y) \leftarrow label

PUSH(S, \langle x+1, y \rangle)

PUSH(S, \langle x, y+1 \rangle)

PUSH(S, \langle x, y-1 \rangle)
22:
23:
24.
26:
27:
                          \mathrm{Push}(S,\langle x\!-\!1,y\rangle)
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```

2. Depth-First Flood Filling

- Verwendung einer Datenstruktur
 - Stack von Koordinatenpunkten (u,v)
 - Z.B. Java Klasse Stack
 - Dynamische Speicheralloziierung vom Heap
- Vorteil:
 - Keine Probleme mit dem Stack-Speicher
- Nachteil:
 - Stack kann lang werden, wenn die Objekte sehr groß sind

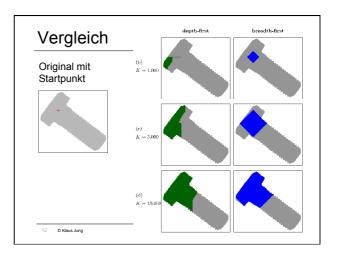
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3. Breadth-First Flood Filling

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FloodFill(I, u, v, label)
                                                                                \quad \triangleright \ Breadth\text{-}First \ Version
29:
             Create an empty queue Q
              Insert the seed coordinate \langle u,v\rangle into the queue: \mathtt{Enqueue}(Q,\langle u,v\rangle)
31
              while Q is not empty \mathbf{do}
                    Get the next coordinate from the front of the queue:
32:
                          \langle x,y \rangle \leftarrow \mathsf{Dequeue}(Q)
33:
                    if coordinate \langle x,y\rangle is within image boundaries and I(x,y)=1
                          _{
m then}
                          then Set I(x,y) \leftarrow label Enqueue(Q,\langle x+1,y\rangle) Enqueue(Q,\langle x,y+1\rangle) Enqueue(Q,\langle x,y+1\rangle) Enqueue(Q,\langle x,y-1\rangle) Enqueue(Q,\langle x-1,y\rangle)
34.
35:
37:
39:
             return
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```

3. Breadth-First Flood Filling

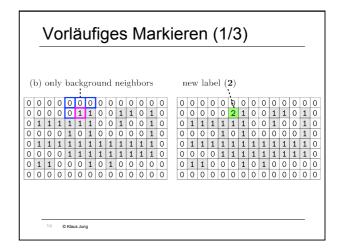
- Verwendung einer Datenstruktur
 - Queue von Koordinatenpunkten (u,v)
 - Z.B. mit Java Klasse java.util.LinkedList
 - DEQUEUE mit removeLast()
 - ENQUEUE mit addFirst()
 - Dynamische Speicheralloziierung vom Heap
- Vorteile:
 - · Keine Probleme mit dem Stack-Speicher
 - Geringerer Speicherbedarf als bei Depth-First

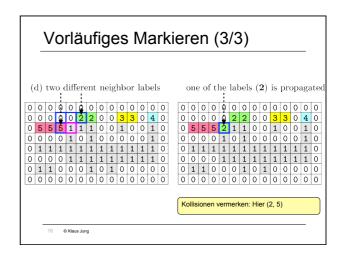


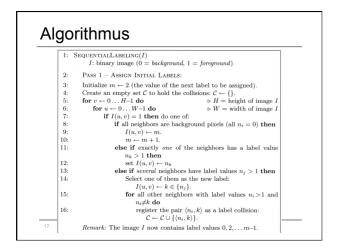
Sequentielle Regionenmarkierung

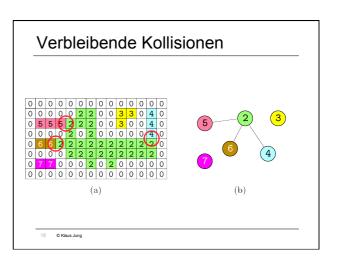
- Auch als "region labeling" bezeichnet
- "Gleichzeitiges" Markieren aller Regionen
- Zwei Schritte
 - 1. Vorläufiges Markieren
 - 2. Auflösung von Kollisionen

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Kollisionen Auflösen

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17: PASS 2 - RESOLVE LABEL COLLISIONS:

18: Let \mathcal{L} = \{2,3,\dots m-1\} be the set of preliminary region labels.

19: Create a partitioning of \mathcal{L} as a vector of sets, one set for each label value: \mathcal{R} \leftarrow [\mathcal{R}_2, \mathcal{R}_3, \dots, \mathcal{R}_{m-1}] = [\{2\}, \{3\}, \{4\}, \dots, \{m-1\}], so \mathcal{R}_i = \{i\} for all i \in \mathcal{L}.

20: for all collisions (a, b) \in \mathcal{C} do

21: Find in \mathcal{R} the sets \mathcal{R}_a, \mathcal{R}_b containing the labels a, b, resp.: \mathcal{R}_a \leftarrow the set which currently contains label a
\mathcal{R}_b \leftarrow the set which currently contains label b

22: if \mathcal{R}_a \neq \mathcal{R}_b (a and b are contained in different sets) then

Merge sets \mathcal{R}_a and \mathcal{R}_b by moving all elements of \mathcal{R}_b to \mathcal{R}_a: \mathcal{R}_a \leftarrow \mathcal{R}_a \cup \mathcal{R}_b
\mathcal{R}_b \leftarrow \{\}

Remark: All equivalent label values (i.e., all labels of pixels in the same region) are now contained in the same sets within \mathcal{R}.
```

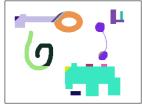
Bild neu Indizieren

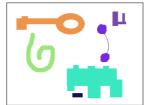
```
24: Pass 3: Relabel the Image:
25: Iterate through all image pixels (u,v):
26: if I(u,v) > 1 then
27: Find the set \mathcal{R}_i in \mathcal{R} which contains label I(u,v).
28: Choose one unique, representative element k from the set \mathcal{R}_i
(e.g., the minimum value, k \leftarrow \min(\mathcal{S})).
29: Replace the image label: I(u,v) \leftarrow k.
30: return the labeled image I.
```

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Beispiel

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einzelne Regionen

verbundene Regionen

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