Computational model of lambda calculus

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Contents

1 General Definitions

Definition 1.1. *An enumeration operator (or e-operator) Ψ^A is a r.e. set. For any $A \subseteq \mathbb{N}$*

$$x \in \Psi^A \iff \exists u \, (finite\ D_u \subseteq A) \,\&\, ((x,u) \in \Psi) \tag{1}$$

Definition 1.2. *If A is a r.e. set then Ψ_A is the enumeration operator defined by it, namely*

$$x \in \Psi_A^B \iff \exists u \, (D_u is finite) \,\&\, ((x,u) \in A \wedge D_u \subseteq B) \tag{2}$$

Definition 1.3. *If θ is an enumeration operator then G_θ is a well-defined r.e. set defining it, namely*

$$(x,u) \in G_\theta \iff x \in \theta^{D_u} \tag{3}$$

Lemma 1.1. *If Ψ is an enumeration operator, then $\Psi_{G_\Psi} = \Psi$*

Proof. TO BE DONE @anton □

Definition 1.4. *Let η be an assignment of r.e. sets to the variables of lambda calculus. With every λ-term E we inductively associate a r.e. set $[\![E]\!]_\eta$:*

1. $[\![x]\!]_\eta = \eta(x)$
2. $[\![E_1 E_2]\!]_\eta = \Psi_{[\![E_1]\!]_\eta}([\![E_2]\!]_\eta)$
3. $[\![\lambda x.E]\!]_\eta = G_{\lambda X.[\![E]\!]_{\eta[x:=X]}}$

Where $\lambda X.[\![E]\!]_{\eta[x:=X]}$ is a function

$$A \in W \mapsto [\![E]\!]_{\eta[x:=A]} \tag{4}$$

Lemma 1.2. *For any environment η and term t, $[\![t]\!]_\eta$ is an c.e. set.*

Proof. By structural induction on the definition of $[\![\mathrm{t}]\!]_\eta$.

1. $[\![x]\!]_\eta = \eta(x)$ by definition
2. To show that $\Psi_{[\![E_1]\!]_\eta}([\![E_2]\!]_\eta)$ is a c.e. set we prove that Ψ_A^B is an enumeration operator which follows from

$$n \in \Psi_A^B \iff \exists u (D_u \subseteq B \,\&\, < n, u > \in A) \tag{5}$$

3. To be done.

□

Lemma 1.3. *For the following theorem we will need one lemma beforehand for better readability, namely*

$$[\![u]\!]_{\eta[x:=[\![v]\!]_\eta]} = [\![u[x \mapsto v]]\!]_\eta \tag{6}$$

Proof. Structural induction on u.

1. u = x, then:

$$[\![x]\!]_{\eta[x:=[\![v]\!]_\eta]} = \eta(x) = [\![v]\!]_\eta = [\![x[x \mapsto v]]\!]_\eta \tag{7}$$

2. u = y $\neq$ x, then:

$$[\![y]\!]_{\eta[x:=[\![v]\!]_\eta]} = \eta(y) = [\![y]\!]_\eta = [\![y[x \mapsto v]]\!]_\eta \tag{8}$$

3. u = pq, then:

$$\begin{aligned}
[\![pq]\!]_{\eta[x:=[\![v]\!]_\eta]} &\overset{def1.4.2}{=} \Psi_{[\![p]\!]_{\eta[x:=[\![v]\!]_\eta]}}([\![q]\!]_{\eta[x:=[\![v]\!]_\eta]}) \\
&\overset{ind.hyp.}{=} \Psi_{[\![p[x\mapsto v]]\!]_\eta([\![q[x\mapsto v]]\!]_\eta])} \\
&\overset{def1.4.2}{=} [\![p[x \mapsto v]q[x \mapsto v]]\!]_\eta \\
&\overset{defApp}{=} [\![pq[x \mapsto v]]\!]_\eta
\end{aligned}$$

4. u = λ_y p, then:

$$\begin{aligned}
[\![\lambda_y p]\!]_{\eta[x:=[\![v]\!]_\eta]} &\overset{def1.4.3}{=} G_{\Lambda Y.[\![p]\!]_{\eta[x:][\![v]\!]_\eta;y:=Y]}} \\
&\overset{ind.hyp}{=} G_{\Lambda Y.[\![p[x\mapsto v]]\!]_{\eta[y:=Y]}} \\
&\overset{def1.4.3}{=} [\![\lambda_y p[x \mapsto v]]\!]_\eta
\end{aligned}$$

□

Theorem 1.4. *If* $E_1 \overset{\beta}{\twoheadrightarrow} E_2$ *then* $[\![E_1]\!]_\eta = [\![E_2]\!]_\eta$ *for any* η.

Proof. We will prove one step of the β reduction and then by induction the rest will follow

We have that $(\lambda x E_1)E_2 = E_1[x \mapsto E_2]$ and we will prove that $[\![(\lambda x E_1)E_2]\!]_\eta = [\![E_1[x \mapsto E_2]]\!]_\eta$

$$\begin{aligned}
[\![(\lambda x E_1)E_2]\!]_\eta &\overset{def.1.4.2}{=} \Psi_{[\![\lambda x E_1]\!]_\eta}([\![E_2]\!]_\eta) \\
&\overset{def.1.4.3}{=} \Psi_{G_{\Lambda X[\![E_1]\!]_{\eta[x:=X]}}}([\![E_2]\!]_\eta) \\
&\overset{Lemma1.1}{=} \Lambda X[\![E_1]\!]_{\eta[x:=X]}([\![E_2]\!]_\eta) \\
&\overset{def\Lambda}{=} [\![E_1]\!]_{\eta[x:=[\![E_2]\!]_\eta]} \\
&\overset{Lemma1.3}{=} [\![E_1[x \mapsto E_2]]\!]_\eta
\end{aligned}$$

□

Definition 1.5. *We define the fixed point combinator Y to be:*

$$Y \stackrel{def}{=} \lambda f.((\lambda x.f(x\,x))\,((\lambda x.f(x\,x)) \tag{9}$$

Exercise 1.1. *Show that any λ-term F has fixed point YF, such that*

$$YF \;=\; F(YF). \tag{10}$$

Proof.

$$\begin{aligned} YF &= ((\lambda x.F(x\,x))\,((\lambda x.F(x\,x)) \\ &= F((\lambda x.F(x\,x))\,((\lambda x.F(x\,x)) \\ &= F(\lambda f.((\lambda x.f(x\,x))\,((\lambda x.f(x\,x))F) \\ &= F(YF) \end{aligned}$$

□

Exercise 1.2. *Let Ψ be any e-operator. Show that Ψ can be expressed as:*

$$\Psi \;=\; \Psi_{[\![x]\!]_\eta} \tag{11}$$

for suitable choise of assigment η.

Proof. Since Ψ is an e-operator, then $\forall B\,\exists e:\ \Psi = \Phi_e$, therefore:

$$\begin{aligned} n \in \Psi^B &\iff \exists u\ <n,u>\in W_e\ \&\ D_u \subseteq B \\ &\iff \exists u\ <n,u>\in [\![x]\!]_\eta\ \&\ D_u \subseteq B,\ \ where\ \ \eta(x) = W_e \\ &\iff n \in \Psi^B_{[\![x]\!]_\eta} \end{aligned}$$

□

Exercise 1.3. *Let Ψ be any e-operator. Show that Ψ has a r.e. fixed point W, such that:*

$$\Psi^W \;=\; W \tag{12}$$

Proof. Let $W = [\![Yx]\!]_\eta$, then:

$$[\![Yx]\!]_\eta = [\![x(Yx)]\!]_\eta = \Phi_{[\![x]\!]_\eta}([\![Yx]\!]_\eta) \tag{13}$$

$$\begin{aligned} n \in \Phi_{[\![x]\!]_\eta}([\![Yx]\!]_\eta) &\iff \exists u\ <n,u>\in [\![x]\!]_\eta\ \&\ D_u \subseteq [\![Yx]\!]_\eta,\ let\,\eta(x) = W_e \\ &\iff \exists u\ <n,u>\in W_e\ \&\ D_u \subseteq [\![Yx]\!]_\eta \\ &\iff \eta \in \Psi^{[\![Yx]\!]_\eta} \end{aligned}$$

□

Exercise 1.4. *Show that for any assigment η:*

$$[\![Yx]\!]_{\eta[x:=\emptyset]} = \emptyset \tag{14}$$

Proof. We know that:

$$[\![Yx]\!]_{\eta[x:=\emptyset]} = [\![x(Yx)]\!]_{\eta[x:=\emptyset]} = \Phi_{[\![x]\!]_{\eta[x:=\emptyset]}}([\![Yx]\!]_{\eta[x:=\emptyset]}) \tag{15}$$

$$\begin{aligned} n \in \Phi_{[\![x]\!]_{\eta[x:=\emptyset]}}([\![Yx]\!]_{\eta[x:=\emptyset]}) &\iff \exists u \; < n, u > \in [\![x]\!]_{\eta[x:=\emptyset]} \;\&\; D_u \subseteq [\![Yx]\!]_{\eta[x:=\emptyset]} \\ &\iff \exists u \; < n, u > \in \emptyset \;\&\; D_u \subseteq [\![Yx]\!]_{\eta[x:=\emptyset]} \\ &\Rightarrow \neg \exists n \; : \; n \in \Phi_{[\![x]\!]_{\eta[x:=\emptyset]}} \end{aligned}$$

□

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References

[1] S. B. Cooper, *Computability Theory*. 2003.

[2] P. Odifreddi, *Classical recursion theory*. 1989.