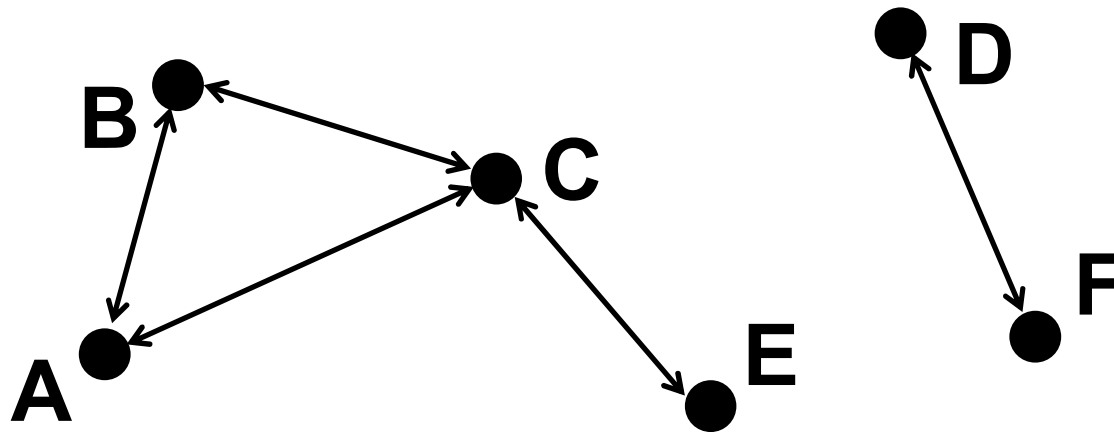


Graphs and Search Algorithms

Social Network Backend

Want to build social network app tracking:

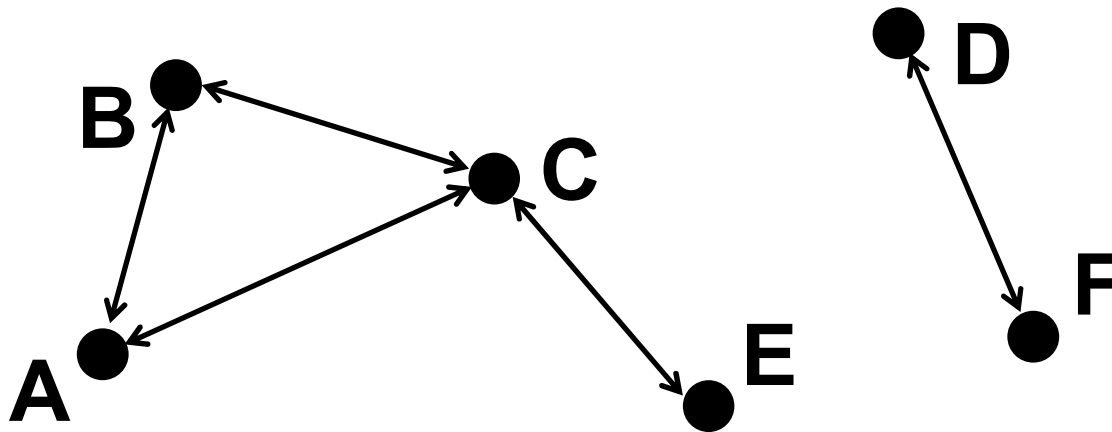
- users
- *friend* relationship between some pairs of users



Graph Basics

Network is a **graph**:

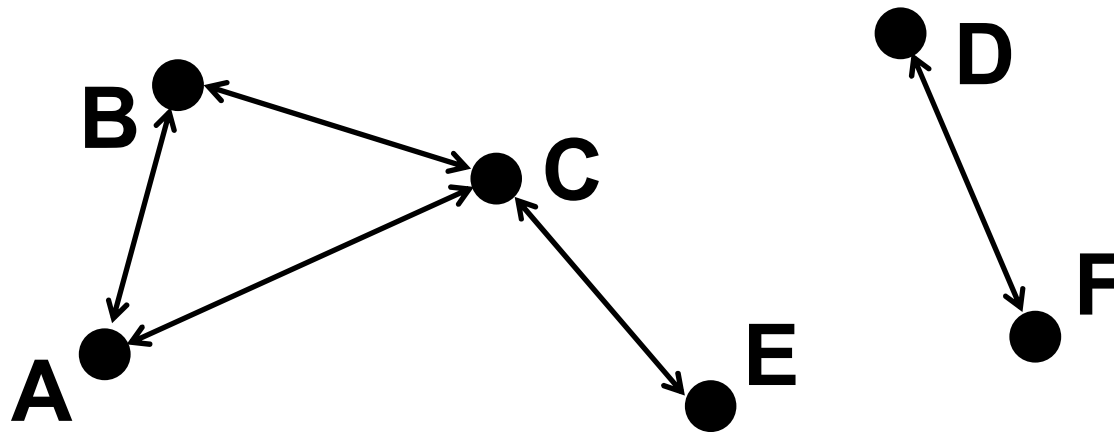
- **nodes or vertices** $V = \{A, B, C, D, E, F\}$



Graph Basics

Network is a **graph**:

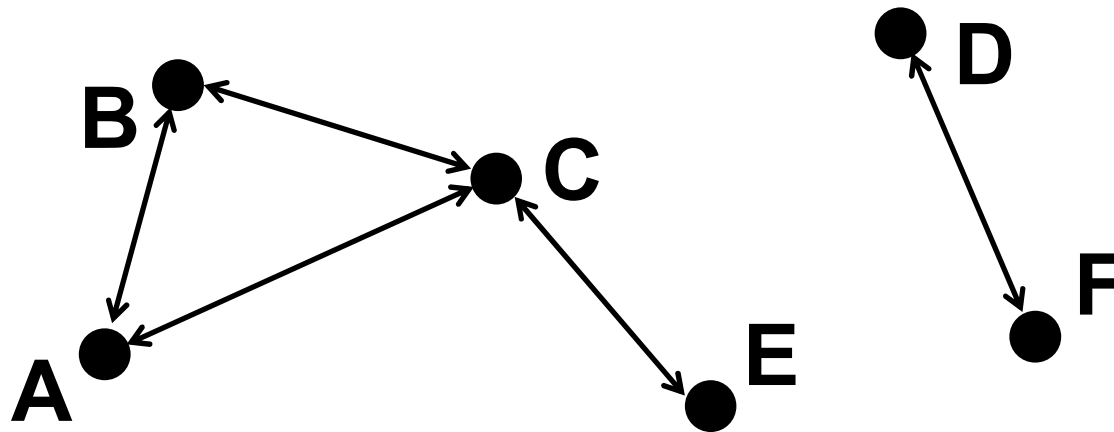
- **nodes** or **vertices** $V = \{A, B, C, D, E, F\}$
- **edges** $E = \{\{A, B\}, \{A, C\}, \{B, C\}, \dots\}$



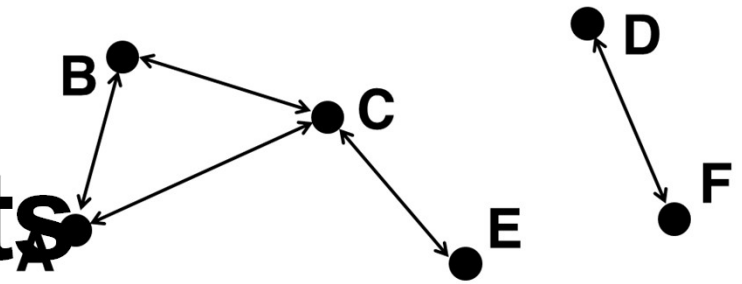
Graph Basics

Network is a **graph**:

- **nodes** or **vertices** $V = \{A, B, C, D, E, F\}$
- **edges** $E = \{\{A, B\}, \{A, C\}, \{B, C\}, \dots\}$
 - can be **directed** (one-way) or **undirected**



Graph Data Structs

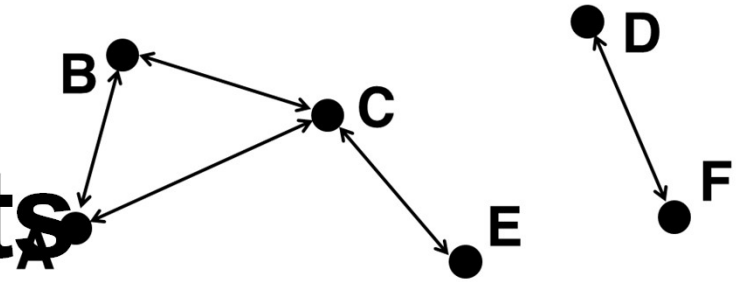


1. Raw vertex/edge lists

$$V = \{A, B, C, D, E, F\}$$

$$E = \{\{A, B\}, \{A, C\}, \{B, C\}, \dots\}$$

Graph Data Structs



1. Raw vertex/edge lists

$$V = \{A, B, C, D, E, F\}$$

$$E = \{\{A, B\}, \{A, C\}, \{B, C\}, \dots\}$$

2. Adjacency list

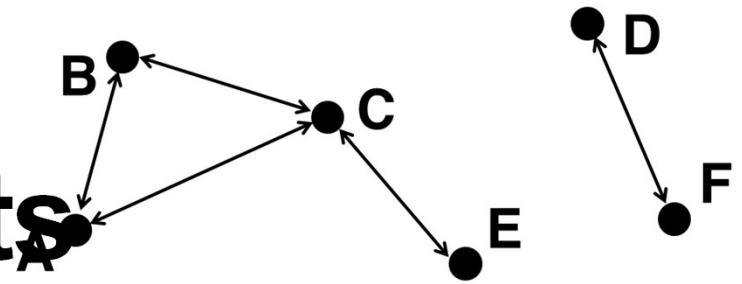
$$A : \{B, C\}$$

$$B : \{A, C\}$$

$$C : \{A, B, E\}$$

...

Graph Data Structs



1. Raw vertex/edge lists

$$V = \{A, B, C, D, E, F\}$$

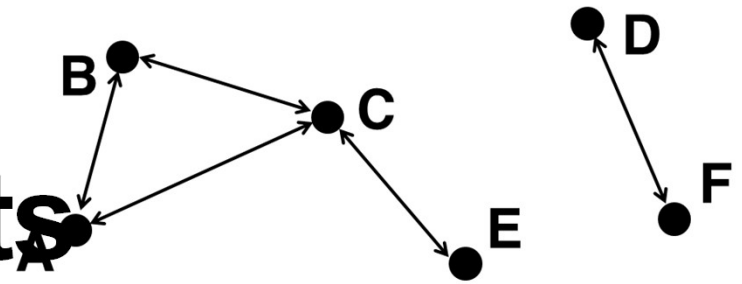
$$E = \{\{A, B\}, \{A, C\}, \{B, C\}, \dots\}$$

2. Adjacency list

3. Adjacency matrix

	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>F</i>
<i>A</i>	0	1	1	0	0	0
<i>B</i>						
<i>C</i>						
<i>D</i>						
<i>E</i>						
<i>F</i>						

Graph Data Structs



1. Raw vertex/edge lists

$$V = \{A, B, C, D, E, F\}$$

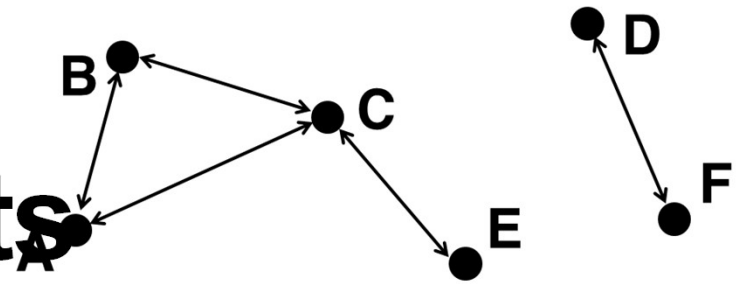
$$E = \{\{A, B\}, \{A, C\}, \{B, C\}, \dots\}$$

2. Adjacency list

3. Adjacency matrix

	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>F</i>
<i>A</i>	0	1	1	0	0	0
<i>B</i>	1	0	1	0	0	0
<i>C</i>	1	1	0	0	1	0
<i>D</i>	0	0	0	0	0	1
<i>E</i>	0	0	1	0	0	0
<i>F</i>	0	0	0	1	0	0

Graph Data Structs



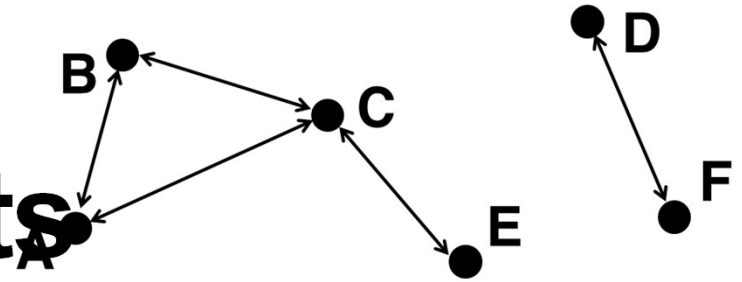
What is the space cost of each option?

Raw lists:

Adjacency list:

Adjacency matrix:

Graph Data Structs



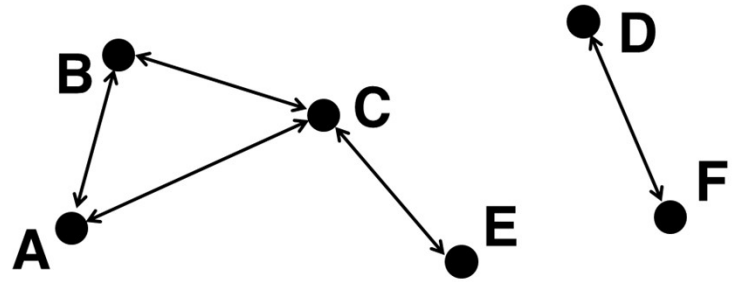
What is the space cost of each option?

Raw lists: $O(|V| + |E|)$

Adjacency list: $O(|V| + |E|)$

Adjacency matrix: $O(|V|^2)$

Graph Operations



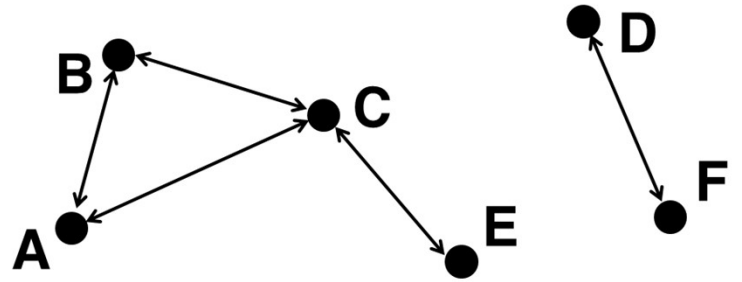
Given two vertices, are they neighbors?

Raw lists:

Adjacency list:

Adjacency matrix:

Graph Operations



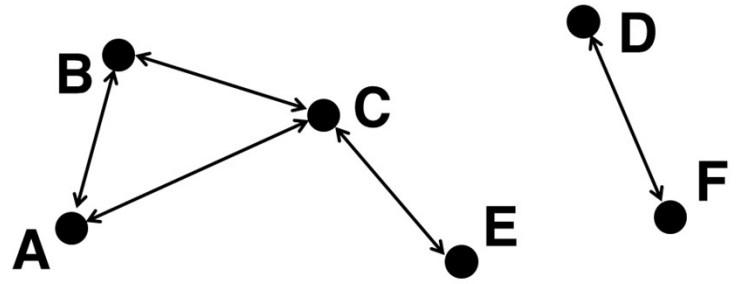
Given two vertices, are they neighbors?

Raw lists: search entire edge list $O(|E|)$

Adjacency list:

Adjacency matrix:

Graph Operations



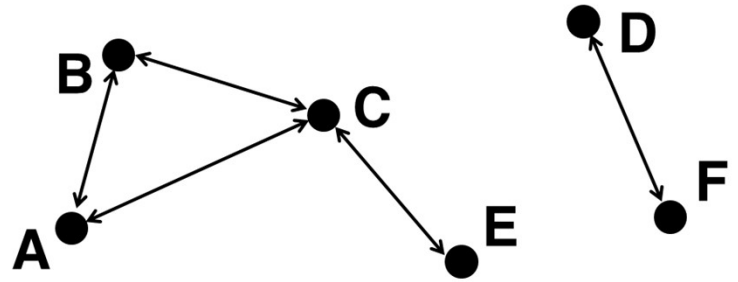
Given two vertices, are they neighbors?

Raw lists: search entire edge list $O(|E|)$

Adjacency list: search one adjacency list
(technically $O(|V|)$)

Adjacency matrix:

Graph Operations



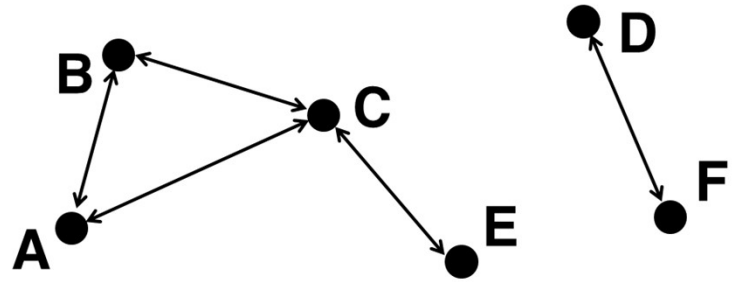
Given two vertices, are they neighbors?

Raw lists: search entire edge list $O(|E|)$

Adjacency list: search one adjacency list
(technically $O(|V|)$)

Adjacency matrix: look up entry $O(1)$

Graph Operations



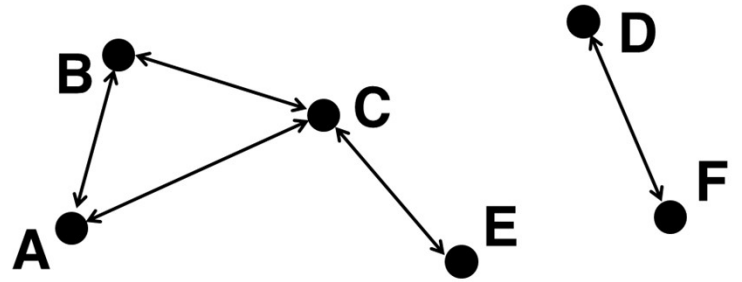
Given a vertex, who are the neighbors?

Raw lists: search entire edge list $O(|E|)$

Adjacency list:

Adjacency matrix:

Graph Operations



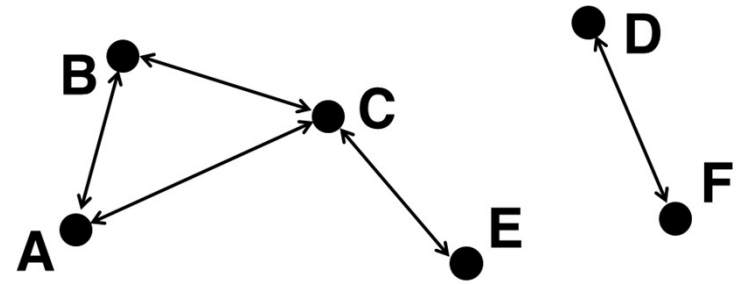
Given a vertex, who are the neighbors?

Raw lists: search entire edge list $O(|E|)$

Adjacency list: nothing to do... $O(1)$

Adjacency matrix: search row of matrix
 $O(|V|)$

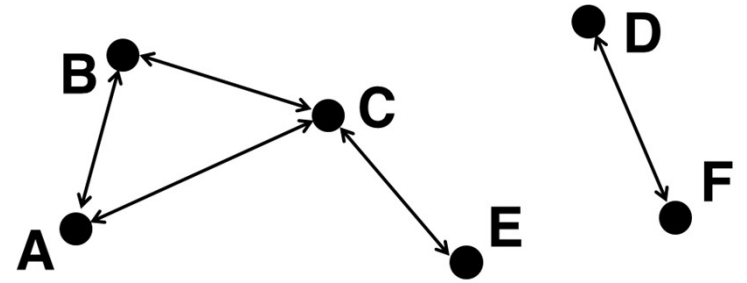
Friend Network



Given a social network containing people (vertices) and friend relationships (edges), **A** is in the same friend network as **B** if

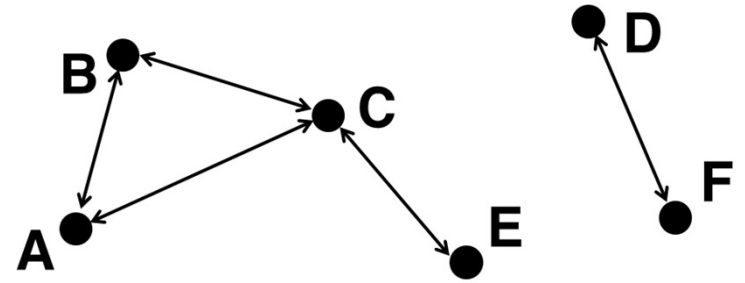
- they are the same person
- **A** is friends with someone that is in the same friend network as **B**

Friend Network



Are **A** and **E** in the same friend network?

Friend Network

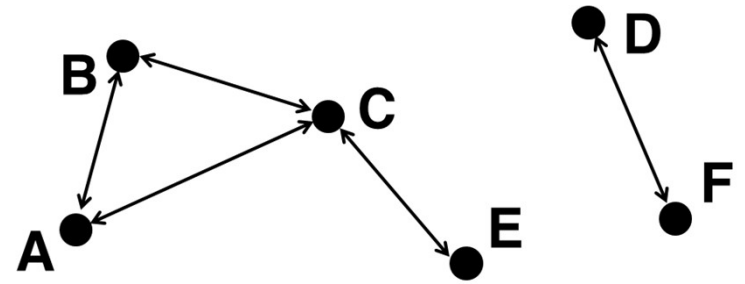


Are **A** and **E** in the same friend network?

Basic idea: start at **A** and “flood fill” along edges, and see if we ever hit **E**

(We will need to create a “visited” flag for vertices)

Friend Network



friendNetwork(**A**,**B**)

for each vertex **v**:

v.visited = false;

return search(**A**);

search(**v**)

if(**v** == **B**) return true;

if(**v**.visited) return false;

v.visited = true;

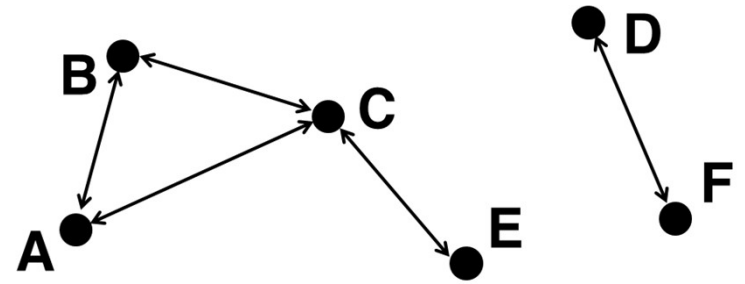
for each neighbor **w**:

 if(search(**w**))

 return true;

return false;

Friend Network



friendNetwork(**A**,**B**)

for each vertex **v**:

v.visited = false;

return search(**A**);

(are there potential issues?)

search(**v**)

if(**v** == **B**) return true;

if(**v**.visited) return false;

v.visited = true;

for each neighbor **w**:

 if(search(**w**))

 return true;

return false;

Iterative Version

friendNetwork(**A**,**B**)

for each vertex **v**:

v.visited = false;

stack **S** = {**A**};

while(!**S**.empty())

v = **S**.pop();

 if(**v** == **B**) return
 true;

 if(**v**.visited) continue;

v.visited = true;

 for each neighbor **w**:

S.push(**w**);

return false;

Iterative Version

friendNetwork(A,B)

for each vertex **v**:

v.visited = false;

stack **S** = {**A**};

depth-first search (DFS)



while(!**S**.empty())

v = **S**.pop();

 if(**v** == **B**) return
 true;

 if(**v**.visited) continue;

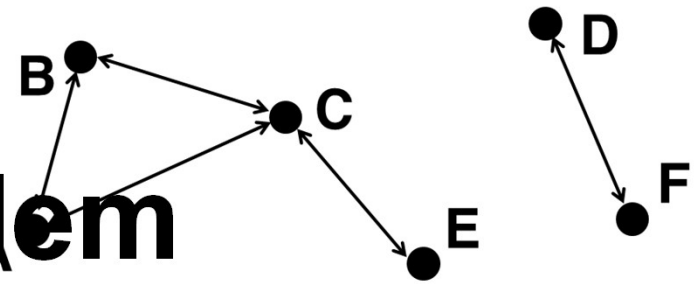
v.visited = true;

 for each neighbor **w**:

S.push(**w**);

return false;

Kevin Bacon Problem



Given a social network and two people **A**, **B**, what is the shortest chain of friends from **A** to **B**?

Ex:

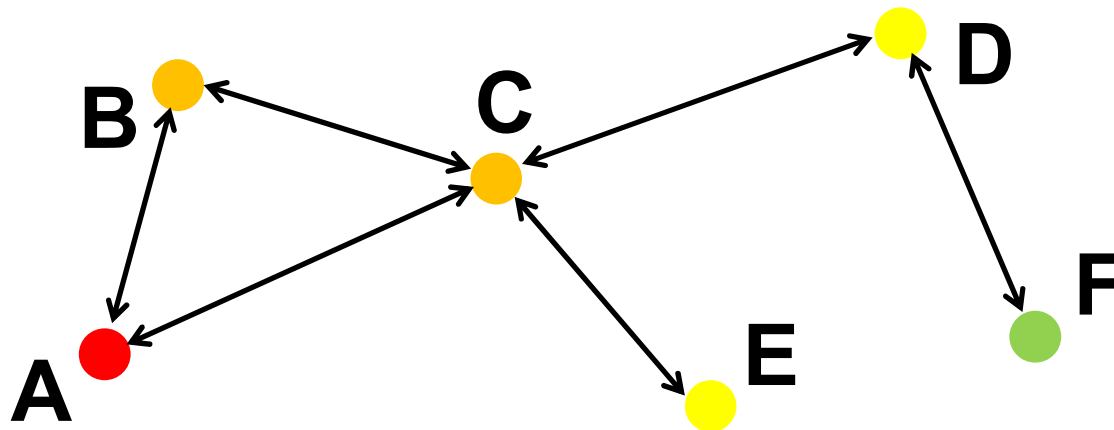
$$\text{bacon}(\mathbf{A}, \mathbf{A}) = 0$$

$$\text{bacon}(\mathbf{A}, \mathbf{E}) = 2$$

$$\text{bacon}(\mathbf{C}, \mathbf{D}) = \text{infinity}$$

Kevin Bacon Problem

Intuition: when calculating $\text{bacon}(\mathbf{A}, *)$ we still want to flood-fill, but we need to guarantee we search friends **before** friends-of-friends



Breadth-First Search

friendNetwork(**A**,**B**)

for each vertex **v**:

v.visited = false;

queue **Q** = {**A**};

while(!**Q**.empty())

v = **Q**.pop();

 if(**v** == **B**) return
 true;

 if(**v**.visited) continue;

v.visited = true;

 for each neighbor **w**:

Q.push(**w**);

return false;

Kevin Bacon Problem

bacon(A,B)

for each vertex **v**:

v.visited = false;

v.dist = infinity;

queue Q = {**A**};

A.dist = 0;

while(!Q.empty())

v = Q.pop();

 if(**v** == **B**) return **v**.dist;

 if(**v**.visited) continue;

v.visited = true;

 for each neighbor **w**:

 Q.push(**w**);

w.dist = **v**.dist + 1;

return infinity;