

Dynamic Programming II

Coin Change?

How many ways to make **N** cents with pennies, nickels, dimes, and quarters?

Order matters

E.g. $6 = \{(1,1,1,1,1,1), (1,5), (5,1)\}$

Coin Change?

What is the optimal subproblem?

- Make change for smaller **N**?
- Make change with fewer coins?

Coin Change?

```
ans[0] = 1
for (i = 1 to N)
    for(c in coins)
        if(i-c >= 0)
            ans[i] += ans[i-c];
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What is time complexity?

Coin Change II

How many ways to make **N** cents with pennies, nickels, dimes, and quarters?

Order **doesn't** matter

E.g. $6 = \{(1,1,1,1,1,1), (1,5)\}$

Coin Change II

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Hmm...

Coin Change II

Simpler case: only pennies and nickels

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Make $5M$ cents using only nickels, then
 $N-5M$ cents using only pennies

Coin Change II

What is the optimal subproblem?

- Make change for smaller N *and* with fewer coins
- $\text{ans}[i,j]$ = ways to make j cents using first i coins

Coin Change II

simple test case: coins are {1¢, 2¢}

i \ j	0	1	2	3	4	5
0						
1						
2						

ans[i,j] = ways to make j ¢ w/ first i coins

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Coin Change II

```
ans[0, 0] = 1
for (c = 0 to num_coins)
    for (i = 0 to N)
        if(c >= 1)
            ans[c, i] += ans[c-1, i]
        if(i - vals[c] >= 0)
            ans[c, i] += ans[c, i -
vals[c]]
```

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```

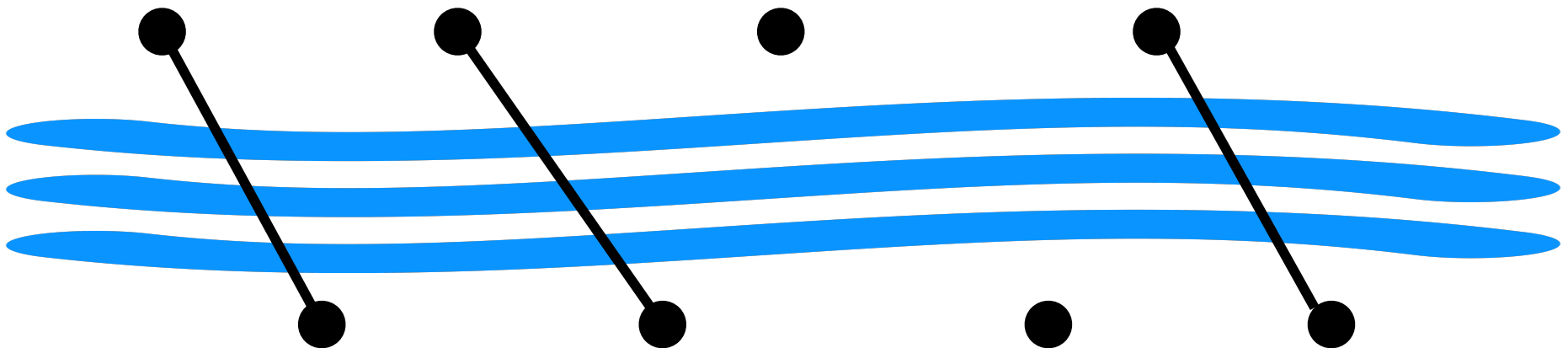
What is time complexity?

Bridge Building

Given points $(p_i, -1)$ and $(q_i, 1)$ on opposite banks of river, build some bridges

Rules:

- no bridges can cross
- every $(p_i, -1)$ must have a bridge

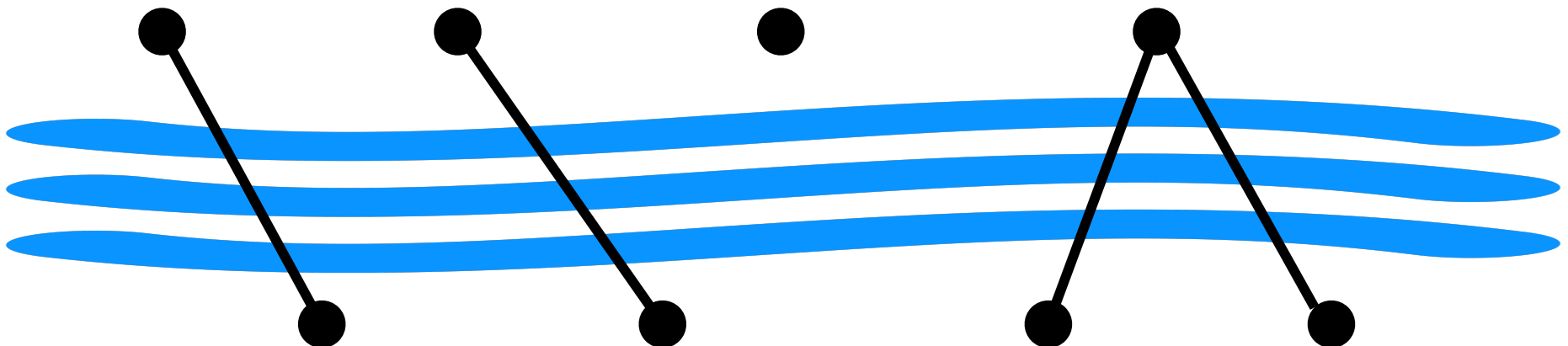


Bridge Building

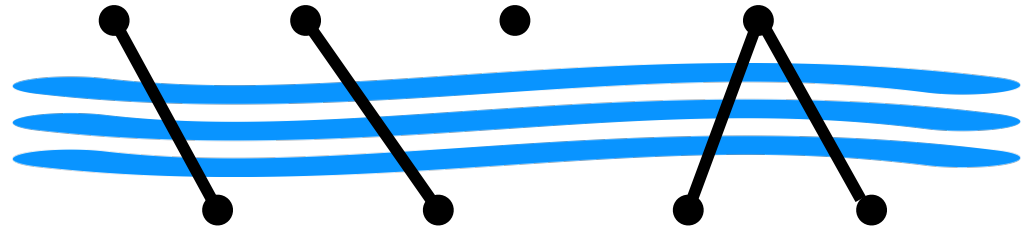
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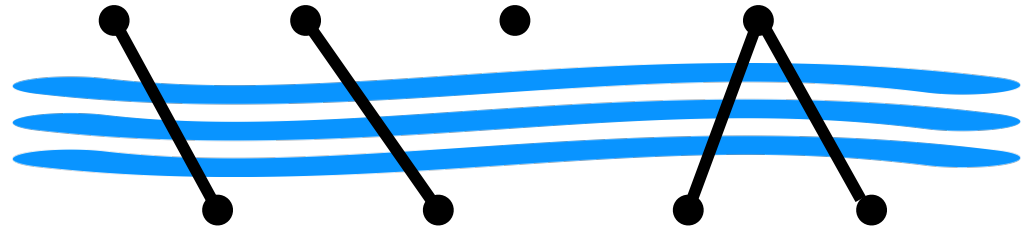
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What is valid bridge network with shortest total length?

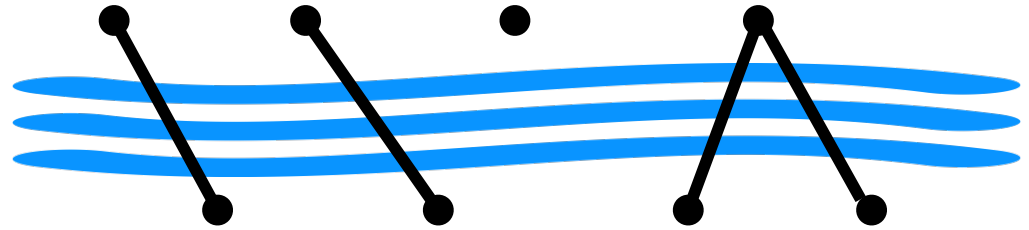
Bridge Building



Greedy algorithm

For each bottom point, connect it to
nearest point on top

Bridge Building



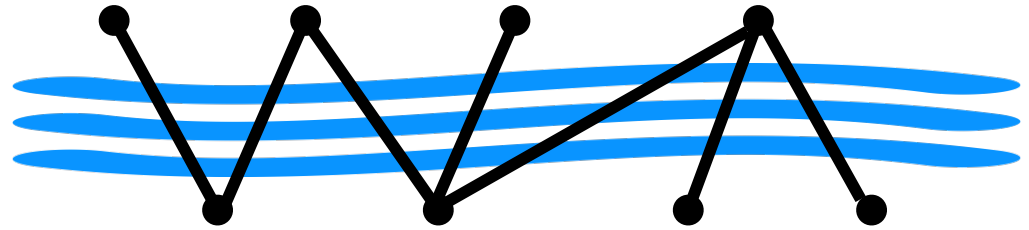
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Rules:

- no bridges can cross
- every $(p_i, -1)$ must have a bridge
- must be **maximally-connected**

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Bridge Building



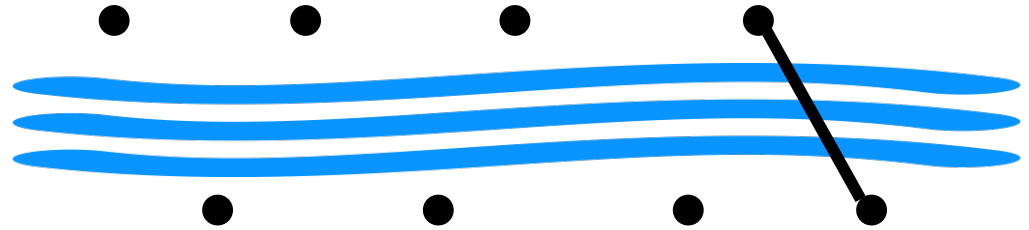
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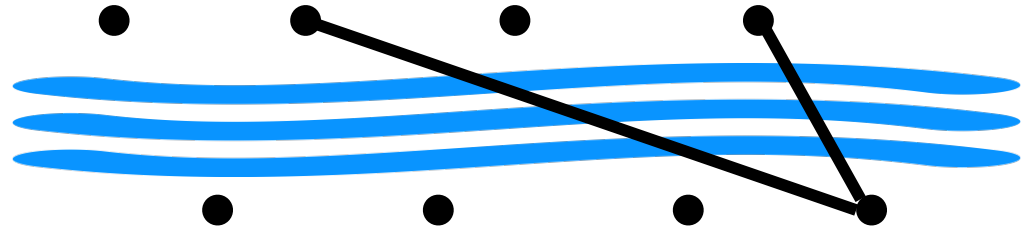
Bridge Building



Some observations:

- rightmost bottom point **must** connect rightmost top point

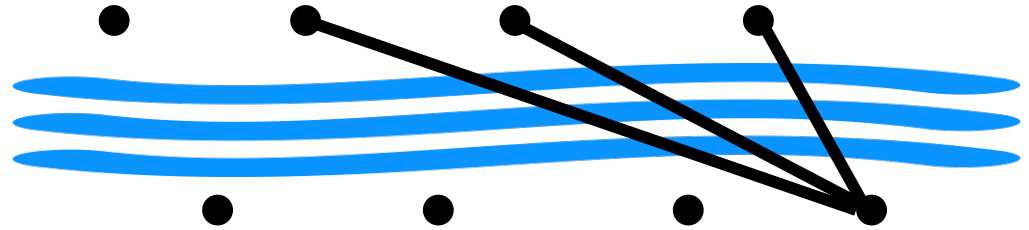
Bridge Building



Some observations:

- rightmost bottom point **must** connect rightmost top point
- could also connect to other top points

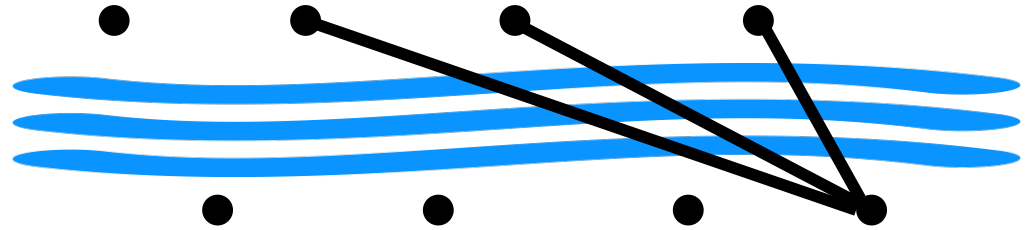
Bridge Building



Some observations:

- rightmost bottom point **must** connect rightmost top point
- could also connect to other top points
- if it connects to a top point, must connect to all other points to the right

Bridge Building



$\text{best}(i,j)$ = shortest network using only up to
bottom point i bottom and top point j

```
best(i,j):  
    ans = infinity  
    for (k = 0 to j)  
        len = best(i-1,k)  
        for (l = k to j)  
            len += sqrt( ( $\mathbf{p}_i - \mathbf{q}_l$ ) * ( $\mathbf{p}_i - \mathbf{q}_l$ ) + 4 )  
        ans = min(ans, len)  
    return ans
```

Traveling Salesman

Given 16 cities and distances d_{ij} between them, what is cheapest itinerary that starts at city 0 and visits all cities exactly once?

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Traveling Salesman

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Find best paths visiting fewer cities?

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Find best paths visiting fewer cities?

- what if we **knew** that best path ended in city **c**?
- now can recurse

Traveling Salesman

$\text{best}[\mathbf{S}, \mathbf{c}]$ = cheapest path visiting all cities in \mathbf{S}
and ending at city \mathbf{c}

```
ans = infinity
for ( $\mathbf{x}$  in  $\mathbf{S} - \{\mathbf{c}, 0\}$ )
    len =  $\text{best}[\mathbf{S} - \{\mathbf{c}\}, \mathbf{x}]$ 
    len +=  $\mathbf{d}_{\mathbf{x}\mathbf{c}}$ 
    ans = min(ans, len)
return ans
```

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What is time complexity?

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What is time complexity?

- 2^{15} subsets containing city 0
- 15 possible ending cities
- $2^{15} * 15 \approx 500,000$

Traveling Salesman

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and ending at city \mathbf{c}

What is time complexity?

- 2^{15} subsets containing city 0
- 15 possible ending cities
- $2^{15} * 15 \approx 500,000$
- NP-hard problem... no free lunch