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Theorem 4.4

$A \in \mathbb{R}^{N \times n}$ with independent isotropic rows,

$m := E \max_{i \leq N} \|A_i\|_2^2$. Then

$$E \max_{j \leq n} |S_j(A) - \sqrt{N}| \leq C \sqrt{m \log \min(m, N)}$$

Lemma 4.5

Let X_i be a finite sequence of independent random vector in some Banach space and (ε_i) be i.i.d. symm Bernoulli, Then

$$E \left\| \sum (X_i - EX_i) \right\| \leq 2 E \left\| \sum \varepsilon_i X_i \right\|$$

Lemma 4.6 Let z be a non-negative r.v.

Then $E|z^2 - 1| \geq \max(E|z - 1|, (E|z - 1|)^2)$

Pf of Thm 4.4

Step 1: we need to control

$$E := E \left\| \frac{1}{N} A^* A - I \right\|$$

$$= E \left\| \frac{1}{N} \sum_{i=1}^N A_i \otimes A_i - I \right\|$$

$$\stackrel{\substack{\uparrow \\ \text{lemma 4.5}}}{\leq} \frac{2}{N} E \left\| \sum_{i=1}^N \varepsilon_i A_i \otimes A_i \right\|$$

Rudelson
conditioning on A

$$\leq \frac{2}{N} E \left(C \sqrt{\log \min(N, n)} \cdot \max_{i \leq N} \|A_i\|_2 \left\| \sum_{i=1}^N A_i \otimes A_i \right\|^{\frac{1}{2}} \right)$$

$$\stackrel{\text{Cauchy ineq}}{\leq} \tilde{C} \sqrt{\frac{m \log(\min(N, n))}{N}} (E+1)^{\frac{1}{2}}$$

$:= 2$

$$E(\dots) \leq m$$

$$\Rightarrow E \leq \sqrt{2^2 + \frac{2^4}{4}} + \frac{2^2}{2} \leq \max(3 \cdot 2, 9 \cdot 2^2).$$

the constants are not the best.

Step 2: Diagonalization

$$\left\| \frac{1}{N} A^* A - I \right\| \stackrel{\substack{\uparrow \\ \text{Diagonalization}}}{=} \left\| \frac{1}{N} U^* \begin{pmatrix} S_1^2(A) & & \\ & \ddots & \\ & & S_{\min(N,n)}^2(A) \end{pmatrix} U - I \right\| \overset{U^*U}{\leftarrow}$$

$$\stackrel{\text{unitary invariance}}{=} \left\| \frac{1}{N} \begin{pmatrix} S_1^2(A) - N & & \\ & \ddots & \\ & & S_{\min(N,n)}^2(A) - N \end{pmatrix} \right\|$$

$$= \max \left(\left| \frac{S_{\min}^2(A)}{N} - 1 \right|, \left| \frac{S_{\max}^2(A)}{N} - 1 \right| \right)$$

$$\max(32, 92^2) \geq E \left\| \frac{1}{N} A^* A - I \right\|$$

$$= E \max \left(\left| \frac{S_{\min}^2(A)}{N} - 1 \right|, \left| \frac{S_{\max}^2(A)}{N} - 1 \right| \right)$$

$$\geq \max \left\{ E \left(\left| \frac{S_{\min}^2(A)}{N} - 1 \right| \right), E \left(\left| \frac{S_{\max}^2(A)}{N} - 1 \right| \right) \right\}$$

lemma 4.6

$$\geq \max \left\{ E \left| \frac{S_{\min}(A)}{\sqrt{N}} - 1 \right|, \left(E \left| \frac{S_{\min}(A)}{\sqrt{N}} - 1 \right| \right)^2, \right.$$

$$\left. E \left| \frac{S_{\max}(A)}{\sqrt{N}} - 1 \right|, \left(E \left| \frac{S_{\max}(A)}{\sqrt{N}} - 1 \right| \right)^2 \right\}$$

$$\Rightarrow 32 \geq \max \left(E \left| \frac{S_{\min}(A)}{\sqrt{N}} - 1 \right|, E \left| \frac{S_{\max}(A)}{\sqrt{N}} - 1 \right| \right)$$

Thus

$$E \max_{j \leq n} \left| \frac{S_j(A)}{\sqrt{N}} - 1 \right| = E \max \left(\left| \frac{S_{\min}(A)}{\sqrt{N}} - 1 \right|, \left| \frac{S_{\max}(A)}{\sqrt{N}} - 1 \right| \right)$$

$$\leq E \left(\left| \frac{S_{\min}(A)}{\sqrt{N}} - 1 \right| + \left| \frac{S_{\max}(A)}{\sqrt{N}} - 1 \right| \right)$$

$$\leq 2 \max (\dots)$$

$$\leq 62$$

Multiplying by \sqrt{N} complete the proof.

(We will omit heavy-tail columns, do sub-gau/exp columns)

5. Random matrices with independent columns

Results of section 4 can be applied to matrices with independent columns.
(consider the transpose)

One obtains that with high probability

$$\sqrt{n} - C\sqrt{N} \leq S_{\min}(A) \leq S_{\max}(A) \leq \sqrt{n} + C\sqrt{N}$$

Lower bound is no longer meaningful for $n < N$.

We seek similar inequality with
— roles of n and N .

Why does the previous approach fail?

The previous approach amounts, to analyzing the spectrum of $AA^* \in \mathbb{R}^{N \times N}$.

If $N > n$, this matrix must be singular, so no sharp bounds can be expected.

Instead, one should again studying

$$\mathbb{R}^{n \times n} \ni A^*A = (\langle A_j, A_k \rangle)_{j,k=1}^n$$

$G \nearrow \quad \nwarrow$
Gram Mat

It follows from Lemma 2.33 that it suffices to show that G is an approximate isometry.

For G , neither the row, nor the column are independent. However, in each of the entries two of the column enter in a (bi)linear fashion pattern.

This can be seen an analogy to the chaos studied in Thm 2.11

↓ 11: 0_ , 1.27 watch video

→ Try a variant of decoupling (Lemma 2.11) different for subgaussian and heavy-tailed.

5.1 Subgaussian columns

Theorem 5.1 (Subgaussian columns)

Let A be an $N \times n$ matrix ($N \geq n$) whose column A_i are independent subgaussian isotropic random vectors in \mathbb{R}^N with $\|A_j\|_2 = \sqrt{N}$ a.s.

Then for every $t \geq 0$, one has $\sqrt{N} - C\sqrt{n} - t \leq S_{\min}(A) \leq S_{\max}(A) \leq \sqrt{N} + C\sqrt{n} + t$ we need a more strict assumption here, in formal one we only need a bound

With probability $\geq 1 - 2\exp(-ct^2)$, where

$C = C'_k$, $C' = C'_k > 0$ depend only on the subgaussian norm bound $K := \max_j \|A_j\|_{\psi_2}$ of the columns.

Difference to Thm 4.1

Normalization of the columns (necessarity as we will see)

Remark: By Lemma 2.33, the conclusion of the Theorem

is equiv to $\|\frac{1}{N} A^* A - I\| \leq C \sqrt{\frac{n}{N}} + \frac{t}{\sqrt{N}}$.

Hence $\|\frac{1}{N} G - I\| = O(\sqrt{\frac{n}{N}})$ w. high prob,

the normalized Gram matrix is an approximate isometry.

For the proof, we decoupling for random subsets.

Definition 5.2 (Random Subsets)

Let $n \in \mathbb{N}$, $m \in \{0, \dots, n\} := [n]$.
← can be arbitrary

Consider independent $\{0, 1\}$ -valued random variables

$\delta_1, \dots, \delta_n$ with $E(\delta_i) = \frac{m}{n}$, we call that these are independent selection.

Then we define a random subset of $[n]$ of average size m to the set-valued r.v. given by $T = \{i \in [n] : \delta_i = 1\}$.

Note that $E|T| = E \sum_{j=1}^n \delta_j = m$.

Lemma 5.3 (Decoupling with random subsets)

Consider a double array of real numbers $(a_{ij})_{i=1}^n$ s.t. $a_{ii} = 0$ for all i . Then

$$\sum_{i,j \in [n]} a_{ij} = 4 E \sum_{i \in T, j \in T^c} a_{ij},$$

where T is a random subset of $\{1, \dots, n\}$ with average size $\frac{n}{2}$.

In particular,

$$4 \min_{T \subseteq [n]} \sum_{\substack{i \in T, \\ j \in T^c}} a_{ij} \leq \sum_{i, j \in [n]} a_{ij} \leq 4 \max_{T \subseteq [n]} \sum_{i \in T} \sum_{j \in T^c} a_{ij}$$

where the min and max are taken over all $T \subseteq [n]$.

Pf: Using the definition of a random subset, we write

$T = \{i \in [n], \delta_i = 1\}$ where the δ_i are independent, sections with

$$\mathbb{E} \delta_i = \frac{1}{2}.$$

We obtain

$$\mathbb{E} \sum_{i \in T} \sum_{j \in T^c} a_{ij} = \mathbb{E} \sum_{i, j \in \{1, \dots, n\}} \delta_i (1 - \delta_j) a_{ij} = \frac{1}{4} \sum_{i, j \in [n]} a_{ij}$$

The second part follows by estimating \mathbb{E} by max and min.

Proof of Thm 5.1:

Step 1: w.l.o.g. assume that the columns A_i have zero mean.

Indeed, multiplying all of the columns A_i by independent symmetric Bernoulli random variables ε_i keeps zero mean while

not changing the singular of the matrix.