K8s Service Load Balancing with BPF & XDP Daniel Borkmann & Martynas Pumputis, Cilium.io

Linux Plumbers 2020

Agenda

- → Part 1: Kubernetes networking 101
- → Part 2: Cilium's service LB & lessons learned
- → Part 3: New BPF kernel extensions





Part 1: K8s networking 101



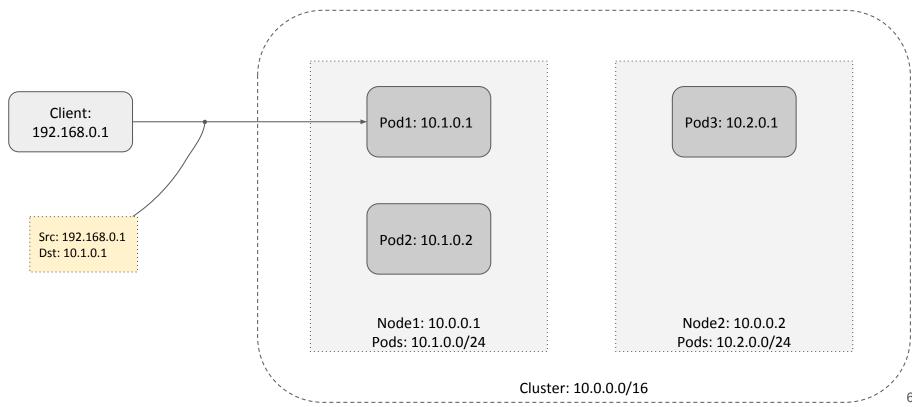
Each Pod must be reachable by its IP addr within a cluster:

- Handled by K8s CNI (e.g. Cilium)
- IP allocation/management and networking
- Doesn't say anything about access from outside



1) Pod IP





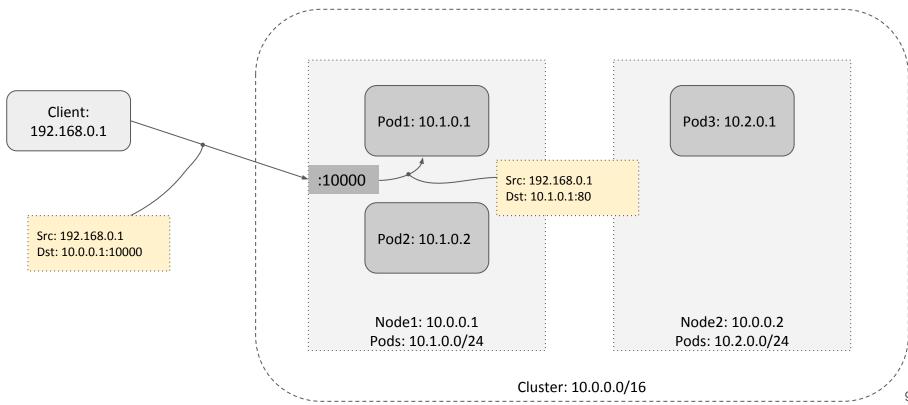


Downside: Pods come and go, no guarantee a Pod IP will ever be preserved.



2) "HostPort"





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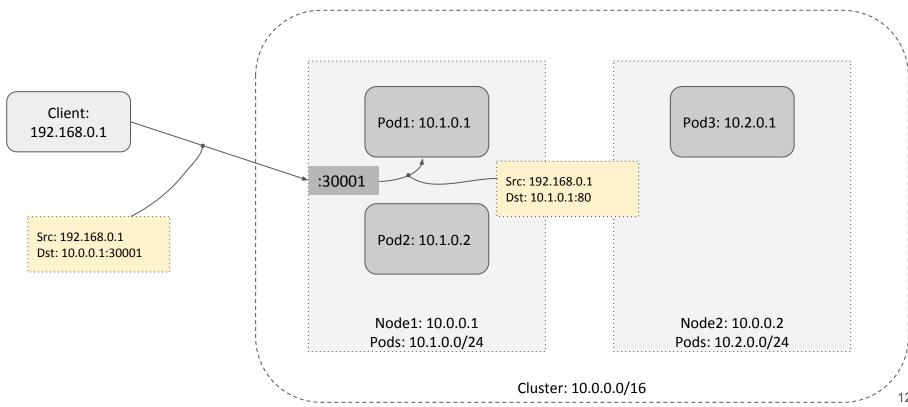


Downside: HostPort to local Pod is 1:1 mapping, that is, only 1 Pod can back-up the HostPort on a node. Use disadviced.

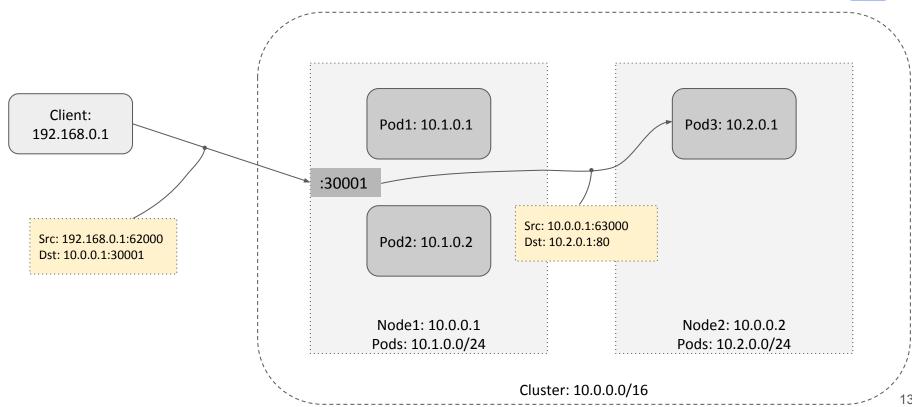


3) "NodePort service"











Advantage: Multiple Pods can back-up NodePort service. Pods can be local or remote to the node.



Advantage: Each node in the cluster reserves the same NodePort port and has same view of backends. Every node becomes a LB.



Advantage: Connectivity from host ns on every node w/o DNS through any local address, e.g. 127.0.0.1:NodePort.

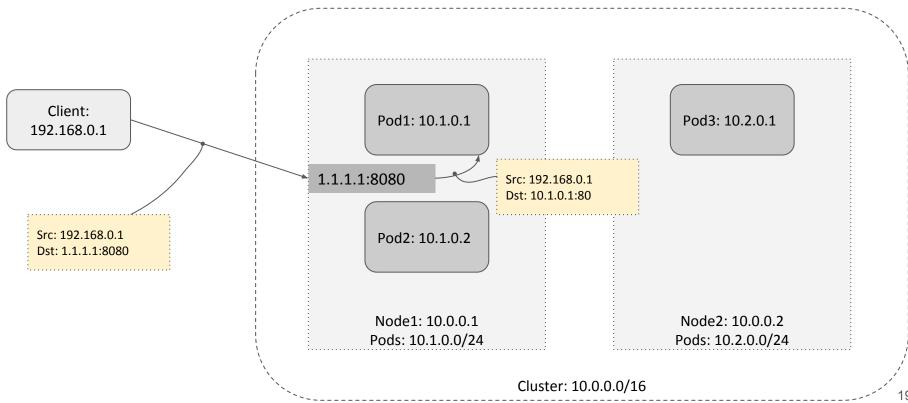


Disadvantage: SNAT-based implementations (common) hide client IP addr and introduce extra hop for replies if backend is remote.

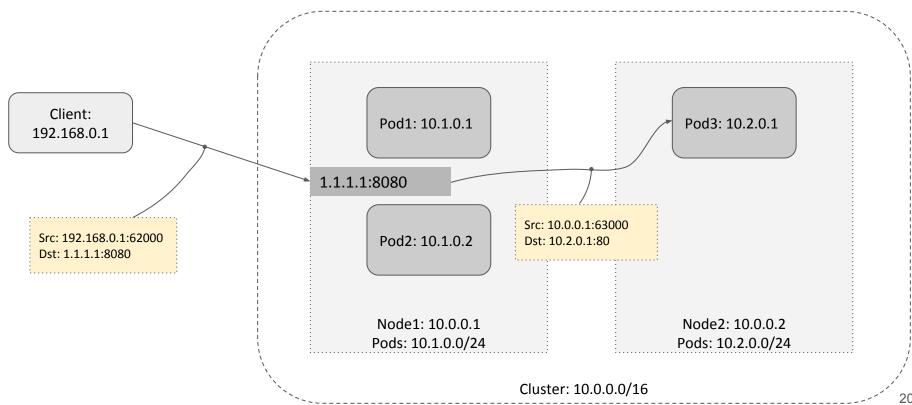


4) "Services with external IPs"











Advantage: Can impersonate any public IP inside the cluster as long as network routes to these nodes.



Downside: External IPs are not managed by K8s. Need to be announced (e.g. via BGP) to route traffic to node.

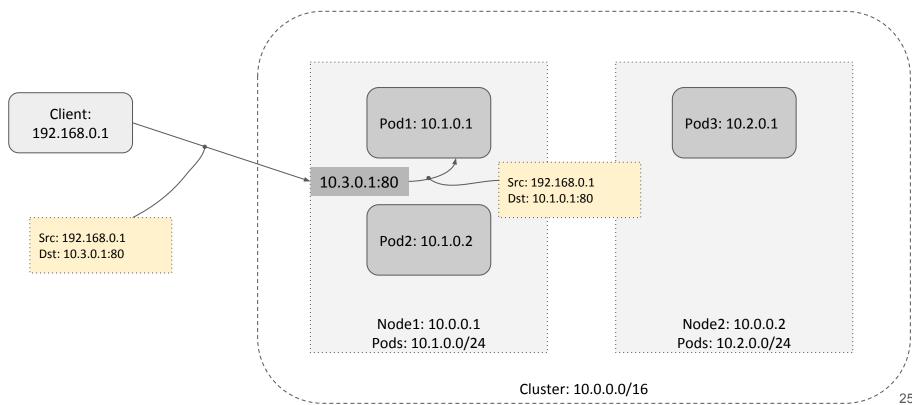


Downside: No in-cluster access to service due to potential of traffic spoofing.

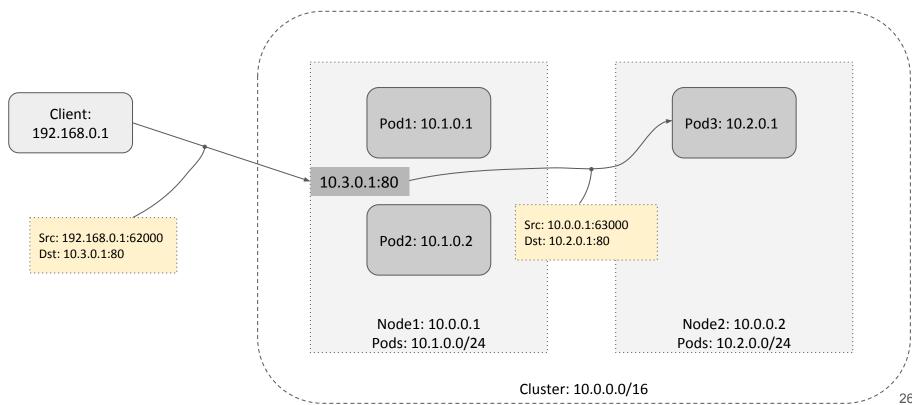


5.1) "LoadBalancer service" (on-prem)











Advantage: LoadBalancer IPs managed via K8s. Not via CNI plugin, but LoadBalancer implementation.



LoadBalancer implementation done by Cloud providers or MetalLB for on-prem. MetalLB can announce via ARP/NDP or BGP.

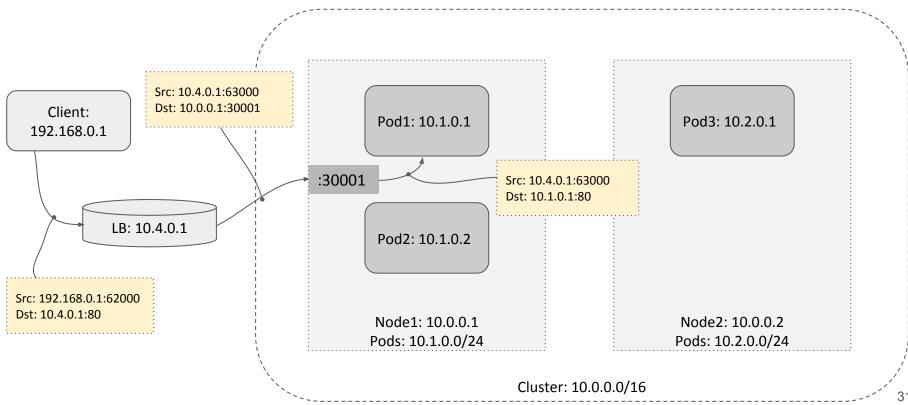


MetalLB does IP address allocation and external announcement, but does not sit in critical fast path (hence works with XDP).

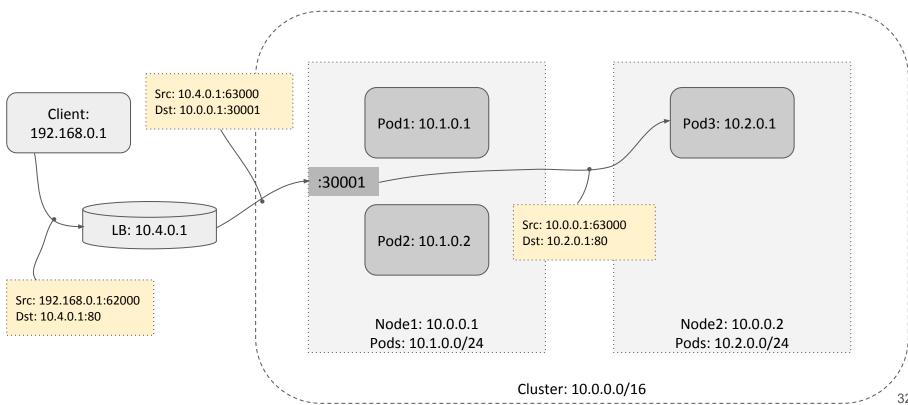


5.2) "LoadBalancer service" (cloud)











Advantage: No additional user setup wrt BGP etc. All major cloud providers offer this for their managed K8s (EKS, GKE, AKS).



Cloud LB performs health checks to probe individual backend nodes from its LB whether they respond.

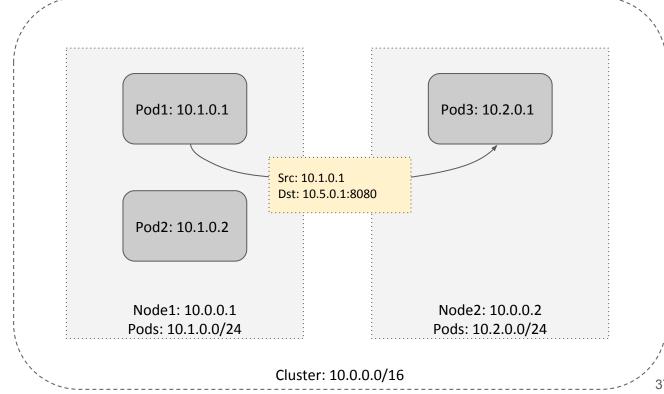


Downside: Two layers of LB, vendor specific LB annotations which are not yet standardised in K8s. Cloud LB programming time can be slow.



6) "ClusterIP service"







Dedicated IP range for ClusterIP, non-routable (always translated locally to backend). For in-cluster access only.



For one LoadBalancer service K8s creates: LoadBalancer, NodePort, ClusterIP services with same set of backends.



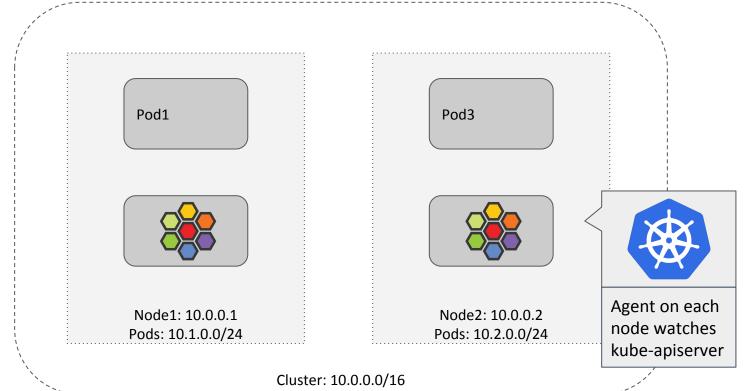
There are also various K8s features for services like sessionAffinity or externalTrafficPolicy (out of scope here).



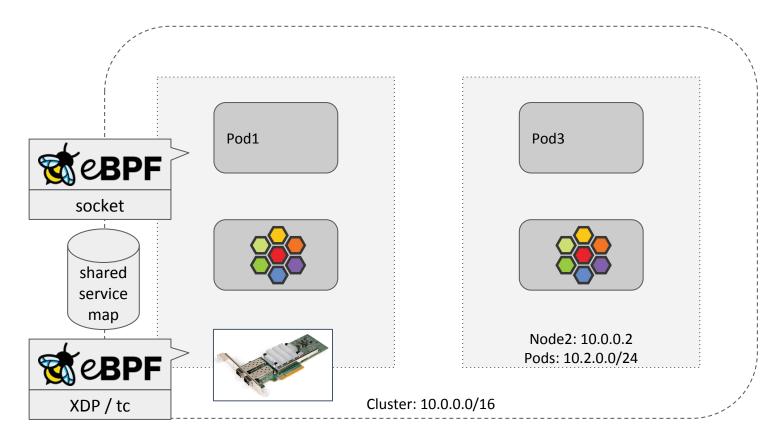


Cilium service LB implements data path for all K8s service types via BPF.





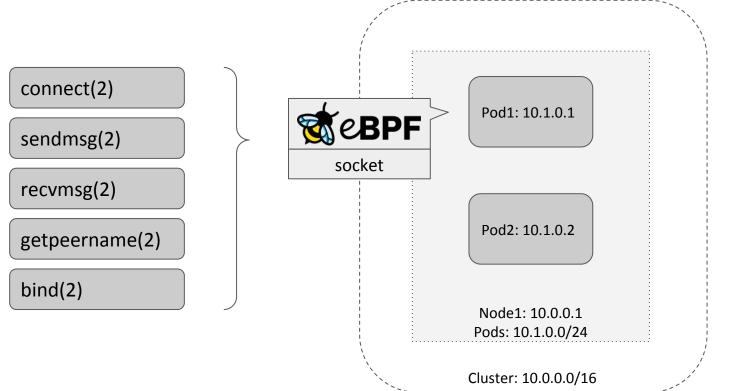






1) E-W (in-cluster): BPF socket LB







K8s Pods are still cgroup v1. Cilium mounts cgroup v2, attaches BPF to root cgroup. Hybrid use works well for root v2.



connect + sendmsg BPF progs do fwd xlation of struct sockaddr. recvmsg + getpeername BPF progs rev xlation. No packet-based NAT! Done for TCP + UDP on v4, v6, v4-in-v6.



Scoped lookup in service map for sock LB vs tc / XDP in order to permit different backends depending on node internal/external traffic.



Wildcarded lookup in service map for sock LB in order to expose service on local or loopback addresses.



Approach to xlate all services on every cluster node at socket layer also faster than kube-proxy: no additional hops in network.



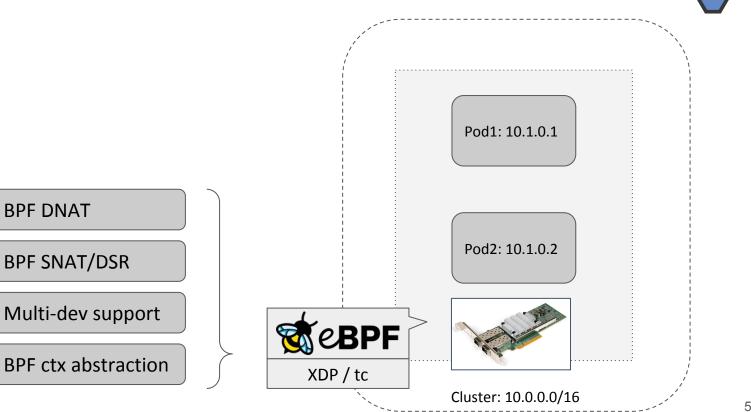
bind BPF prog rejects application requests from binding to NodePort.



Globally unique bpf_get_socket_cookie() and bpf_get_netns_cookie() crucial helpers.



2) N-S: BPF via tc or XDP





D/SNAT engine, DSR, conntrack, etc all implemented in tc BPF. For XDP support Q was whether to abstract context or reimplement ...



Ended up refactoring almost all parts of our BPF code base to make it ctx generic.



Rationale: avoids bit rot, optimizations & fixes in generic code apply to both XDP + tc.



Example, generic code:

```
if (ctx_adjust_room(ctx, 0x8, BPF_ADJ_ROOM_NET,
                             ctx_adjust_room_dsr_flags()))
                return DROP_INVALID;
        if (ctx_store_bytes(ctx, ETH_HLEN + sizeof(*ip4),
                             &opt, sizeof(opt), 0) < 0)
                return DROP_INVALID;
skb specific implementation
                                             xdp specific implementation
```



Example, context specifics:

```
#define __ctx_buff
                                        xdp_md
#define __ctx_is
                                        __ctx_xdp
#include "common.h"
#include "../helpers_xdp.h"
#include "../builtins.h"
#include "../section.h"
#include "../loader.h"
#include "../csum.h"
#define CTX ACT OK
                                        XDP PASS
#define CTX_ACT_DROP
                                        XDP_DROP
                                        XDP_TX /* hairpin only */
#define CTX ACT_TX
```



Most helpers in skb context need inline equivalents for XDP. LLVM tends to optimise which then fails verifier. Inline asm as rescue.





```
static <u>__always_inline</u> <u>__maybe_unused int</u>
xdp_load_bytes(const struct xdp_md *ctx, __u64 off, void *to, const __u64 len)
        void *from;
        int ret;
        asm volatile("r1 = *(u32 *)(%[ctx] +0)\n\t"
                     "r2 = *(u32 *)(%[ctx] +4)\n\t"
                      "%[off] &= %[offmax]\n\t"
                     "r1 += %[off]\n\t"
                     "%[from] = r1\n\t"
                     "r1 += %[len]\n\t"
                     "if r1 > r2 goto +2\n\t"
                      "%[ret] = 0 \ln t"
                      "goto +1\n\t"
                      "%[ret] = %[errno]\n\t"
                      : [ret]"=r"(ret), [from]"=r"(from)
                      : [ctx]"r"(ctx), [off]"r"(off), [len]"ri"(len),
                        [offmax]"i"( CTX OFF MAX), [errno]"i"(-EINVAL)
                      : "r1", "r2");
        if (!ret)
                memcpy(to, from, len);
        return ret;
```



v5.6 kernel was a milestone on XDP side: 'XDP for the masses' on public cloud via ena & hv_netvsc driver.



For max support on variety of drivers though, only bare minimum features must be assumed: XDP_PASS/DROP/TX



Cilium only supports native XDP on user side. Generic XDP only utilized for CI purpose.



Reasons to avoid generic XDP two-fold: given this runs on every end node in the cluster, we cannot linearize every skb & bypass GRO.



Own optimised mem{cpy,zero,cmp,move}(). Compile error for LLVM builtin functions.



LLVM builtins end up as byte-wise ops for non-stack data. No context on efficient unaligned access.



bpf_ktime_get_ns() overhead noticeable
under XDP. Made clock source selectable,
switched to bpf_jiffies64(). Approx +1.1Mpps.



No cb[] for passing data between tail calls in XDP. Initially used xdp_adjust_meta().



Bad for 2 reasons: missing driver support, high rate of cache-misses. Switched to per-CPU scratch map, approx +1.2Mpps.



bpf_map_update_elem() in fast-path hitting bucket spinlock. If assumptions allow, can be converted to lock-free lookup first.



bpf_fib_lookup() expensive, can be avoided and compiled out e.g. for hairpin LB. Approx +1.5Mpps in our test env.



Also don't gamble with LLVM & enforce BPF's tail call patching via text_poke for static slots.





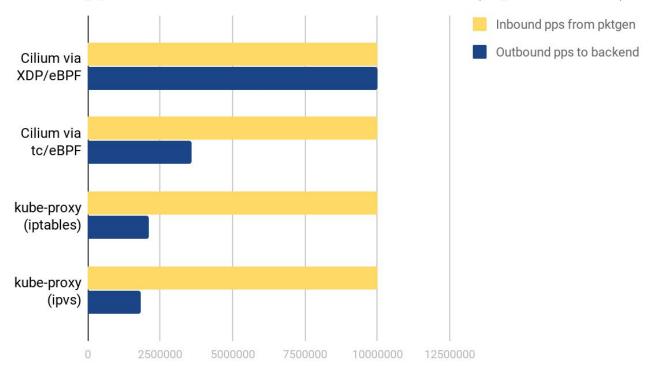
```
static __always_inline __maybe_unused void
tail_call_static(const struct __ctx_buff *ctx, const void *map,
                const __u32 slot)
        if (!__builtin_constant_p(slot))
               throw build bug();
        asm volatile("r1 = %[ctx]\n\t"
                    "r2 = %[map]\n\t"
                     "r3 = %[slot]\n\t"
                     "call 12\n\t"
                     :: [ctx]"r"(ctx), [map]"r"(map), [slot]"i"(slot)
                     : "r0", "r1", "r2", "r3", "r4", "r5");
```



pktgen hairpin test on XDP layer for remote K8s service backend. 10Mpps inbound ...



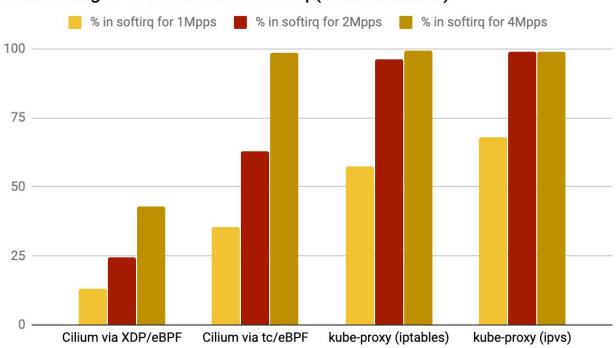
Forwarding performance of tested Kubernetes node (higher is better)







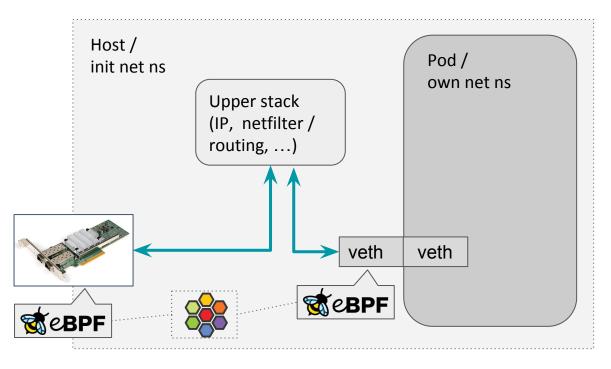
Forwarding CPU overhead in softirq (lower is better)



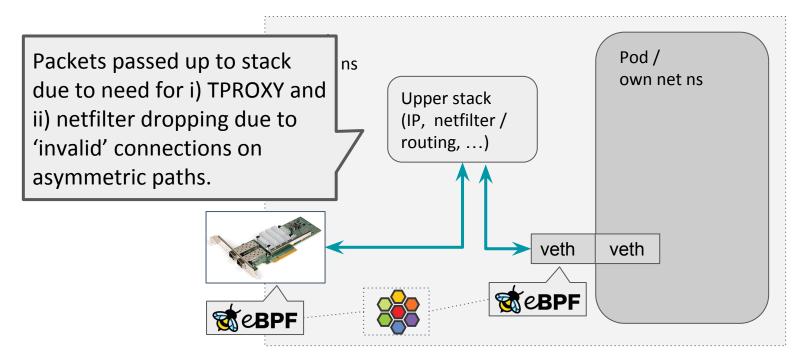


Part 3: New BPF kernel extensions (for Cilium use case)

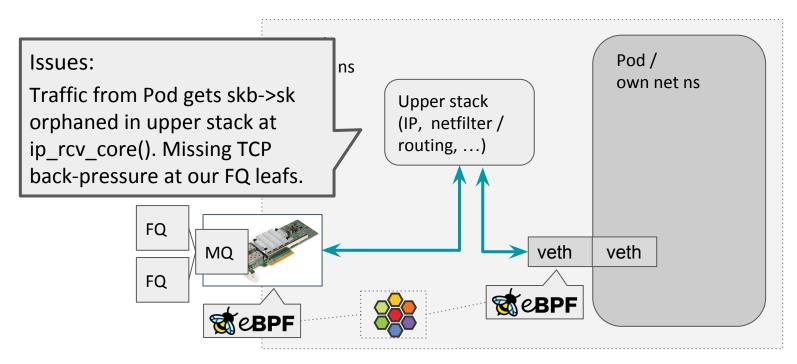




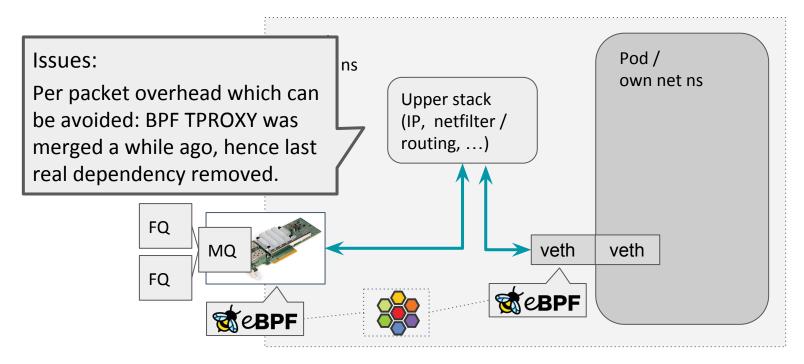




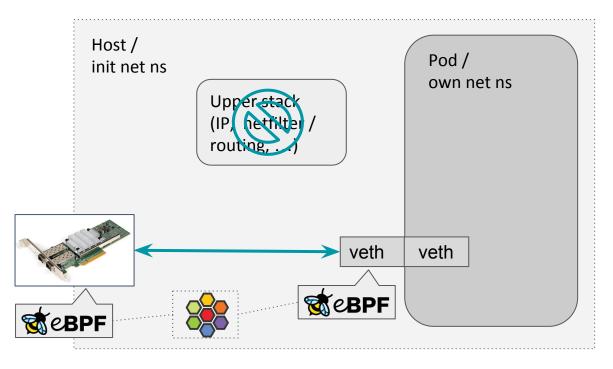








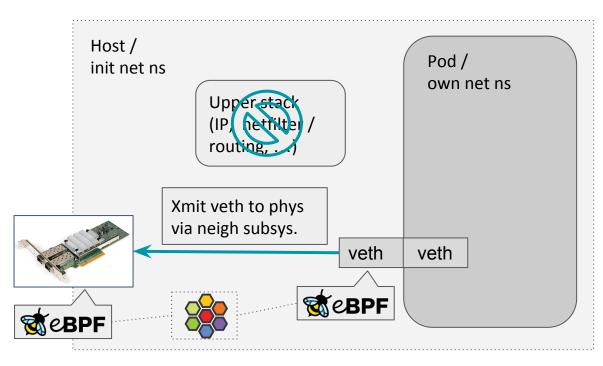




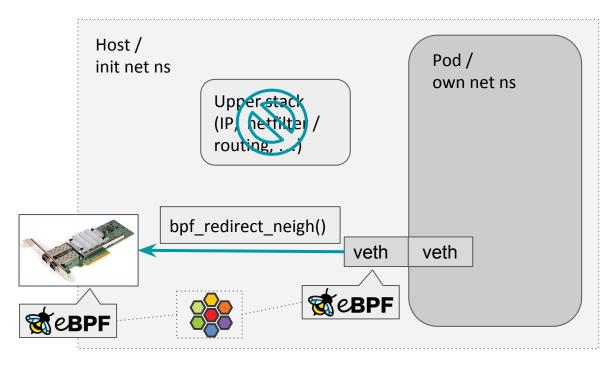


2 new BPF helpers for tc:
 bpf_redirect_neigh() and
 bpf_redirect_peer()...
borrowing ideas from ipvlan.

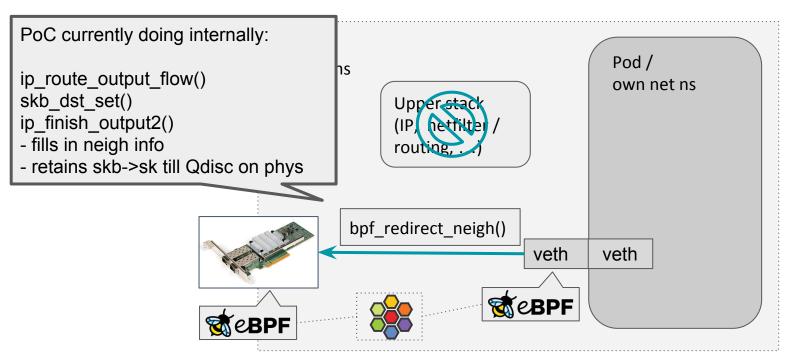




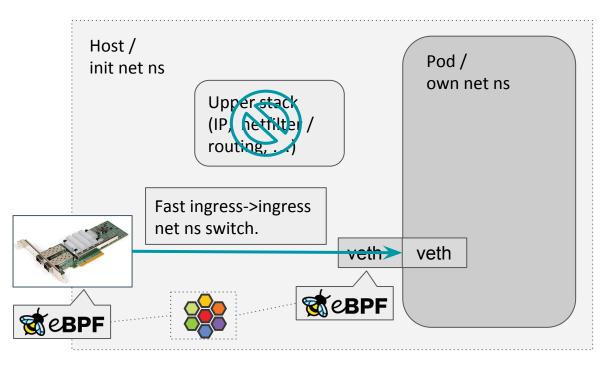




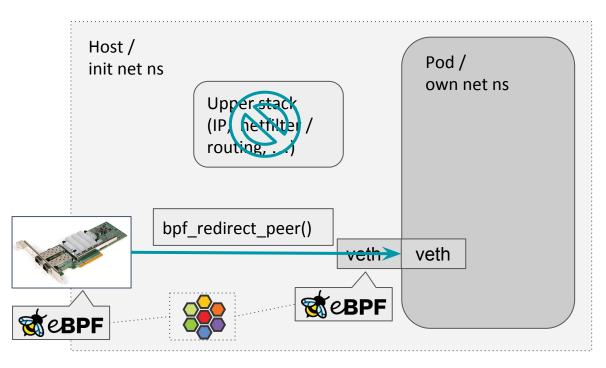




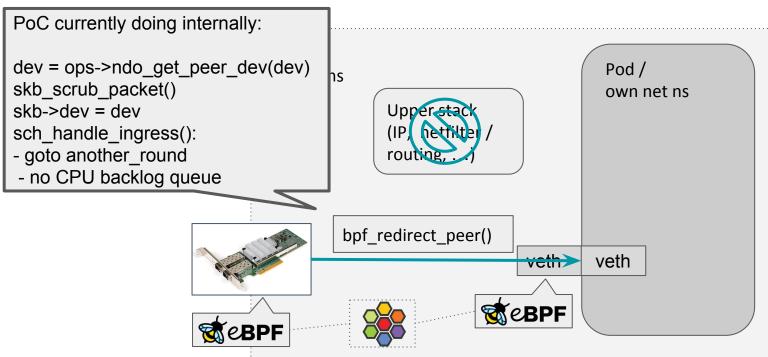




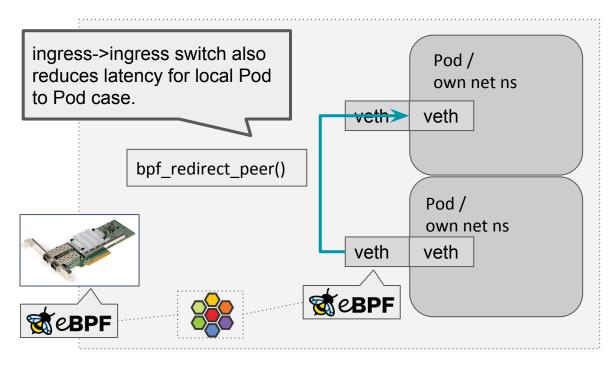




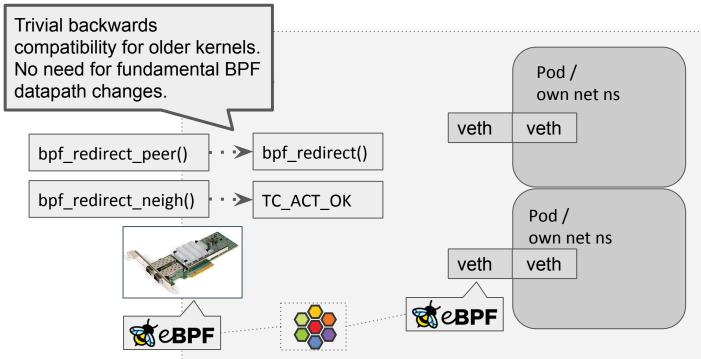






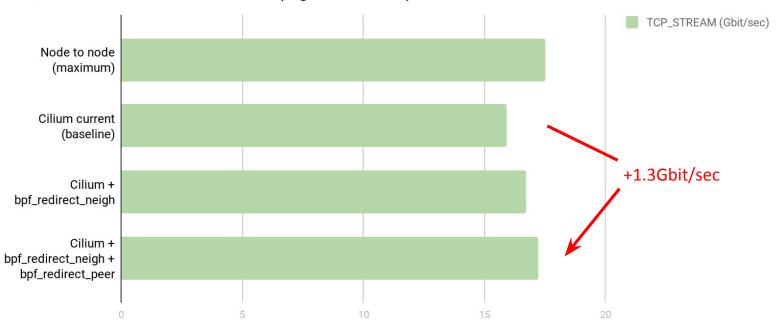






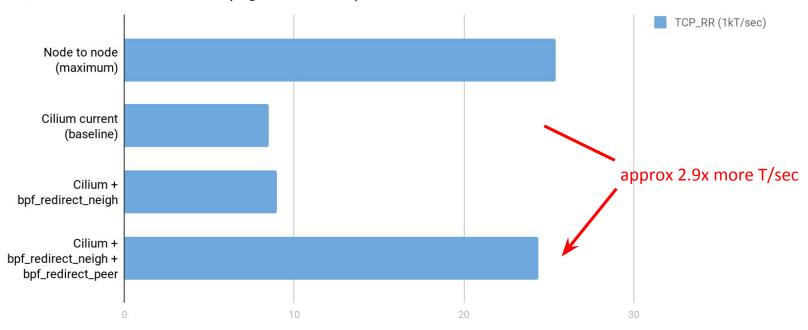


TCP_STREAM remote node to Pod (higher is better)





TCP_RR remote node to Pod (higher is better)



Thanks! Questions?

→ Try it out: https://cilium.link/kubeproxy-free

→ Cilium: https://github.com/cilium/cilium

→ PoC code: https://git.kernel.org/[...]/dborkman/bpf.git



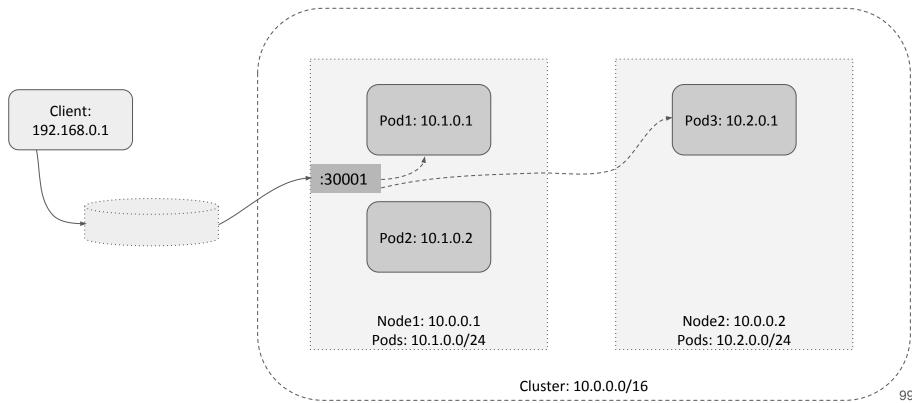


Appendix: K8s networking 101



3/4/5.a) "externalTrafficPolicy=Cluster"







Backends can be local or remote.

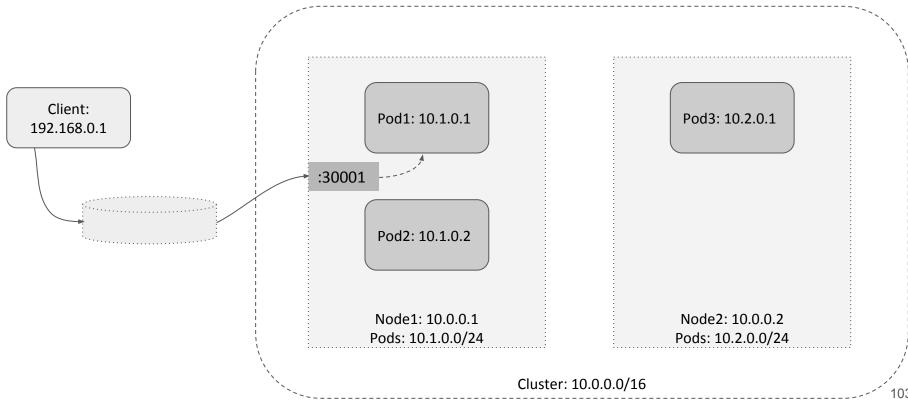


Traffic can be spread evenly across cluster, but for remote backends client source IP is lost unless DSR is available & used (Cilium).



3/4/5.b) "externalTrafficPolicy=Local"







Backends must only ever be local.



Client source IP is preserved. Potential of traffic imbalance though. Needs Pod anti-affinity against hostname.



Nodes w/o service backends drop requests. This is probed via cloud LB health check to update its backends.