```
DGA 03
5) M= (ab)no = 10a+6 N= (cd)no = 10c+d
      A=ac B=6d C=(a-6)(d-c)
    => MN = (ab) no - (cd) no = 100 A + 10 A + 10 B + B + 10 C
 a) ZZ: (ab), (cd), = 100 A+10 A+10B+B+10C
      100 A + 10 A + 10 B + B + 10C = 100 ac + 10ac + 10bd+bd+10 (a-6)(de)
     = 100ac + 10ac +10bd+bd+10ad-10ac-10bd+10bc
     = 10 a (10c+d) + b (10c+d) = (10c+d) (10a+b) = (ab), (cd),
 6) ges: ahnlicher Algorithmus für n-stellige Zahlen
      M=10 a+6 N=10 2 C+d
       A=ac B=bd C=(a-b)(d-c)
      => MN= 102L= A+10L= A+10L= B+B+10L= C
           = 102 (ad-orc-bd+6c)
           = 102(2) ac + bd + 10 (2) ad +10 (2) bc
           = 10 (10 (10 (2) c + d) + b (10 (2) c + d)
            = (10 (2) a+ b) (10 (2) c+d) = MIV
      Algorithm (M, N, n):
        if n == 1: return M * N;
        a= M // 10 + + (n/2); 6= M % 10 + + (n/2);
        C=N//10 da (4/2); d=N% 10 x a (4/2).
        A = Algorithm (a, c); B = Algorithm (b, d);
        C= Algorithm (a-b, d-c);
         return 10 ** n * A+ 10 ** (n/2) * A+10 ** (n/2) * B+B+10 ** (n/2) *C;
  C) Angenommen n=2 KEN T(n)... An zahl einstelligen Multiphikantionen von zwei
      n-stelligen Jahlen, T(1)=1 T(n)=3T(2)+4 a=3 b=2 p(n)=4
      \log_{10}(a-E) = \log_{2}(3-1) = n f(n)=4=O(n) = T(n)=O(n\log_{2}3)
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