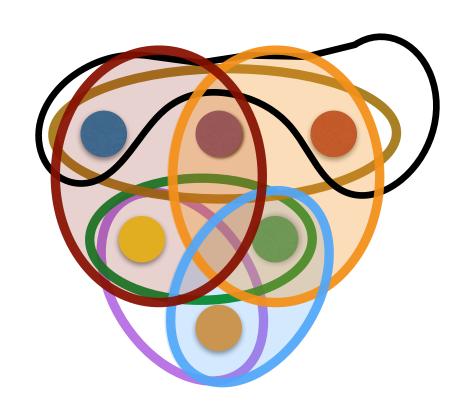
Set cover, linear programming and randomized rounding



Is the output a cover? Maybe, maybe not

Number of elements covered:

\(\sum_e 1 \) (e covered)

On average:

\(\sum_{\mathbf{e}} \text{Pr[e covered]} \)

Consider an element e. With what probability is it covered by output?

 $\Pr[\mathbf{e} \ \mathbf{covered}] =$ $\Pr[\mathbf{there} \ \mathbf{exists} \ S \ \mathbf{in} \ \mathbf{output:} \ \mathbf{e} \in \mathbf{S}]$

 $\Pr[ext{there exists } S ext{ in output:} \\ e \in S] = \\ 1 - \Pr[ext{for all } S ext{ containing } e : \\ S ext{ not in output]}$

Independence

If independence: $Pr[A \text{ and } B] = Pr[A] \times Pr[B]$

 $\Pr[\text{for all } S \text{ containing } e: \\ S \text{ not in output}] = \\ \prod_{S:e \in S} \Pr[S \text{ not in output}]$

 $Pr[S \text{ not in output}] = 1 - x_S$

Together

$$Pr[e covered] = 1 - \prod_{S:e \in S} (1 - x_S)$$

Algebra

$$X = e^{\ln X}$$

$$\ln(\mathbf{XY}) = \ln \mathbf{X} + \ln \mathbf{Y}$$

$$\prod_{\mathbf{S}: \mathbf{e} \in \mathbf{S}} (\mathbf{1} - \mathbf{x}_{\mathbf{S}}) = \mathbf{e}^{\sum_{\mathbf{S}: \mathbf{e} \in \mathbf{S}} \ln(\mathbf{1} - \mathbf{x}_{\mathbf{S}})}$$

Algebra

$$\ln(1 - X) \le -X$$

$$e^{\sum_{\mathbf{S}:\mathbf{e}\in\mathbf{S}}\ln(\mathbf{1}-\mathbf{x}_{\mathbf{S}})}\leq e^{-\sum_{\mathbf{S}:\mathbf{e}\in\mathbf{S}}\mathbf{x}_{\mathbf{S}}$$

Use LP constraint

$$\sum_{s:e\in s} x_s \ge 1$$

$$e^{-\sum_{s:e\in s} x_s} \le e^{-1}$$

Combining

$$Pr[e covered] \ge 1 - 1/e$$

Average number of elements covered:

$$\#(elements)(1-1/e)$$

$$1 - 1/e = 0.63...$$

Recap

Randomized rounding gives collection of sets with average cost at most OPT and covering on average 63% of the elements.

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