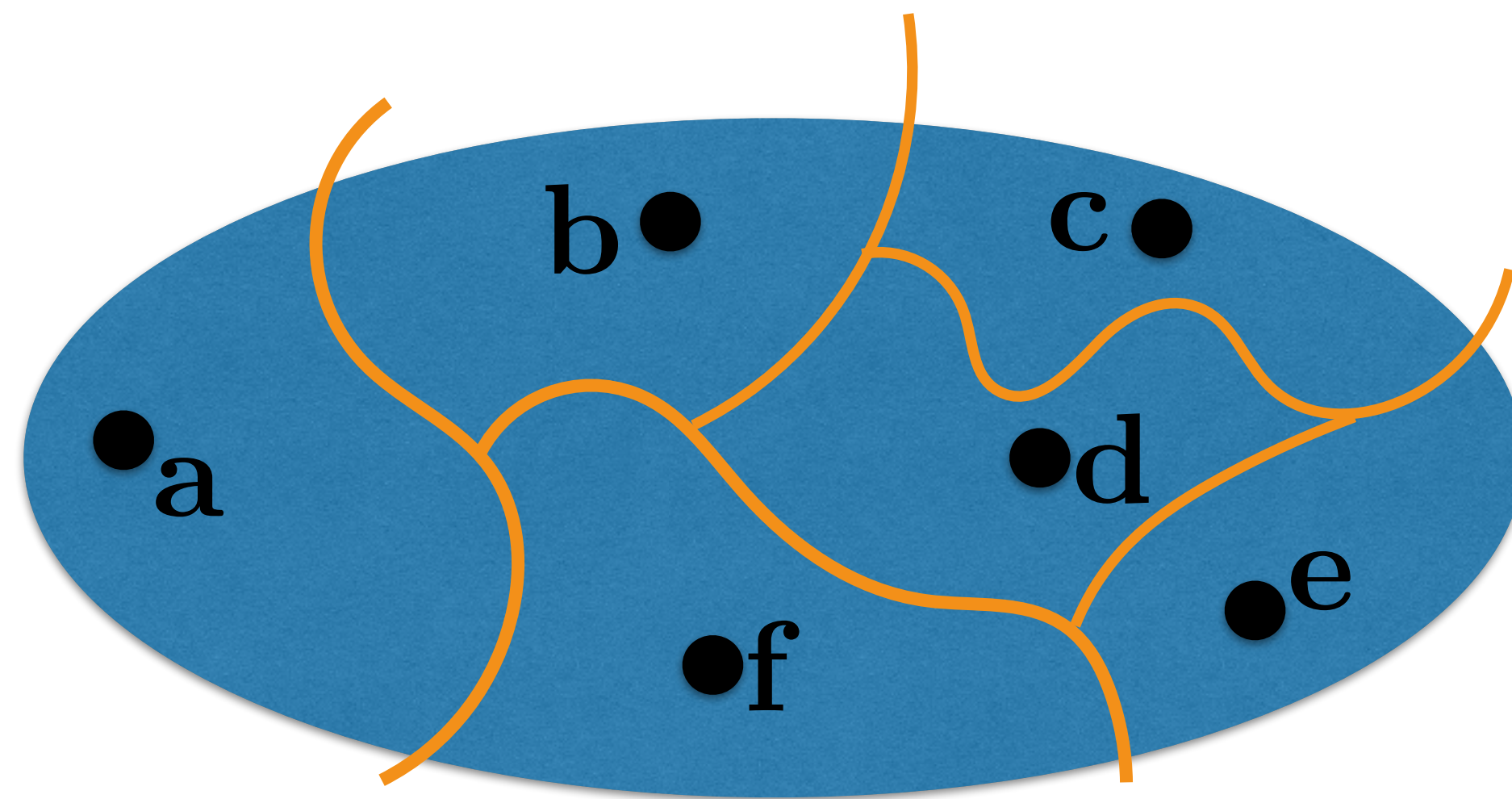


# Multiway cut, linear programming and randomized rounding



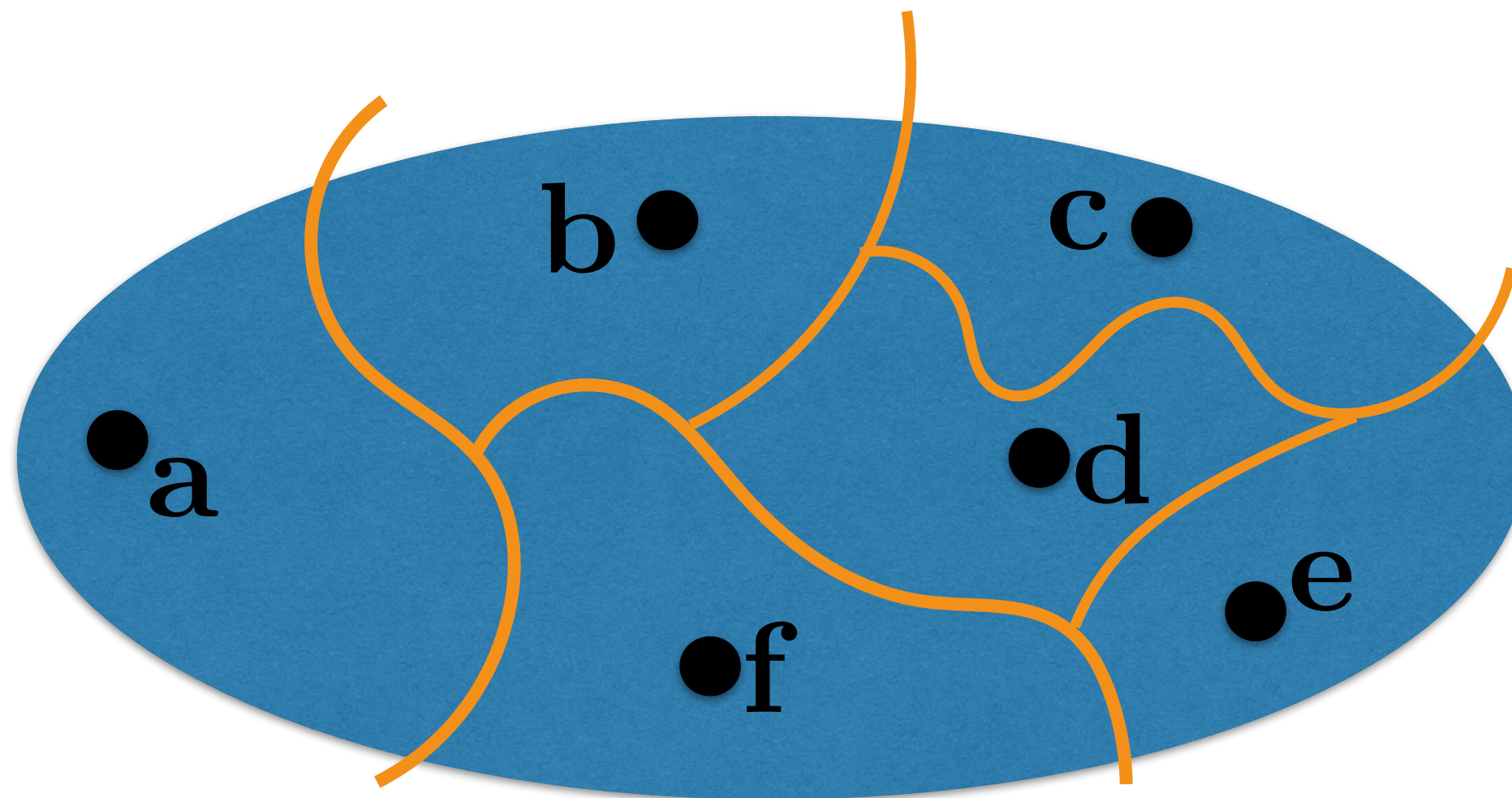
# Problem

Let  $k$  be a fixed integer.

**Given:** graph  $G$  with edge weights,  
 $k$  vertices called “terminals”

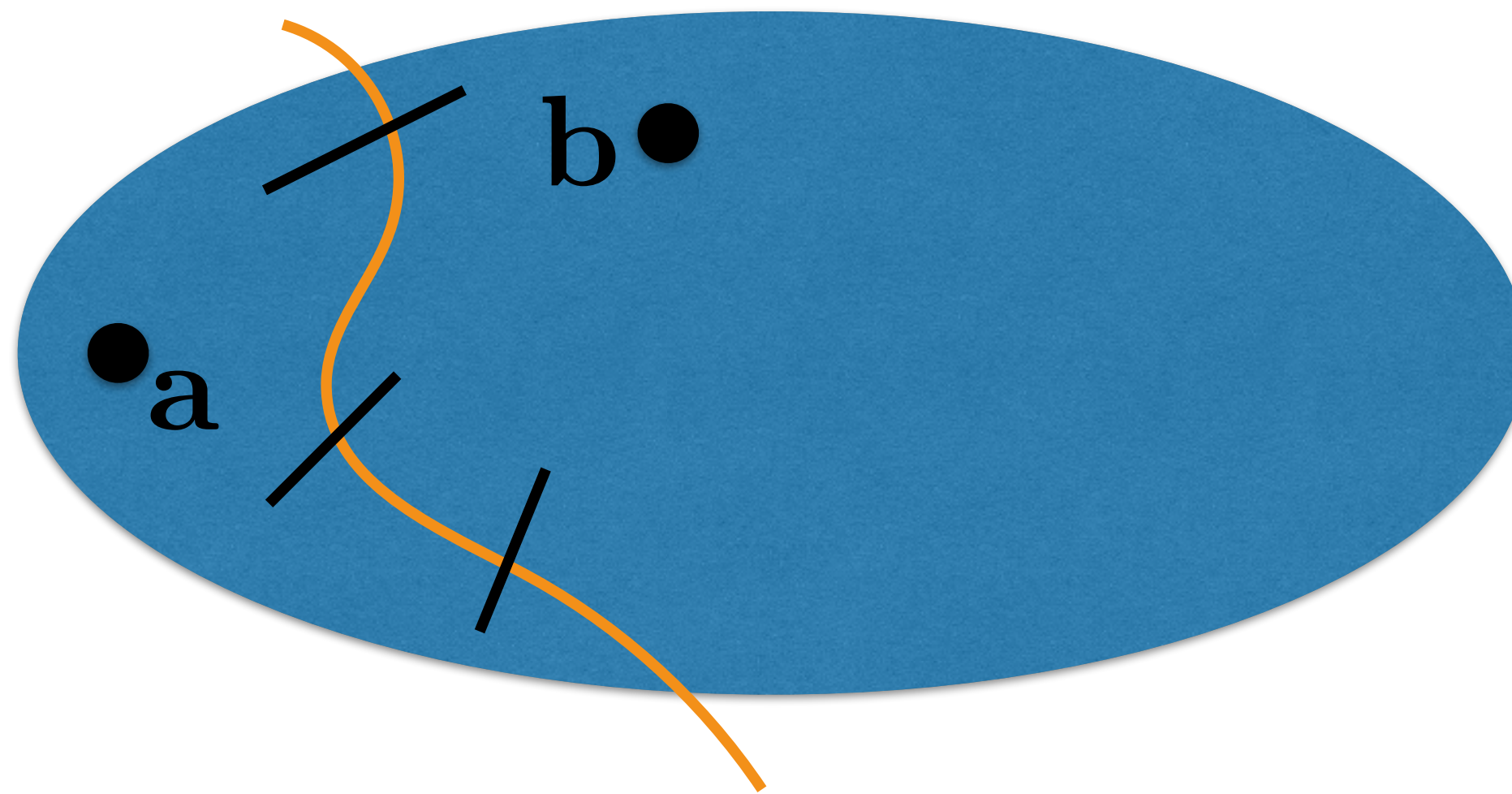
**Find:** subset  $F$  of edges such that  
Removing  $F$  disconnects the  
terminals from one another

**Goal:** minimize weight of  $F$

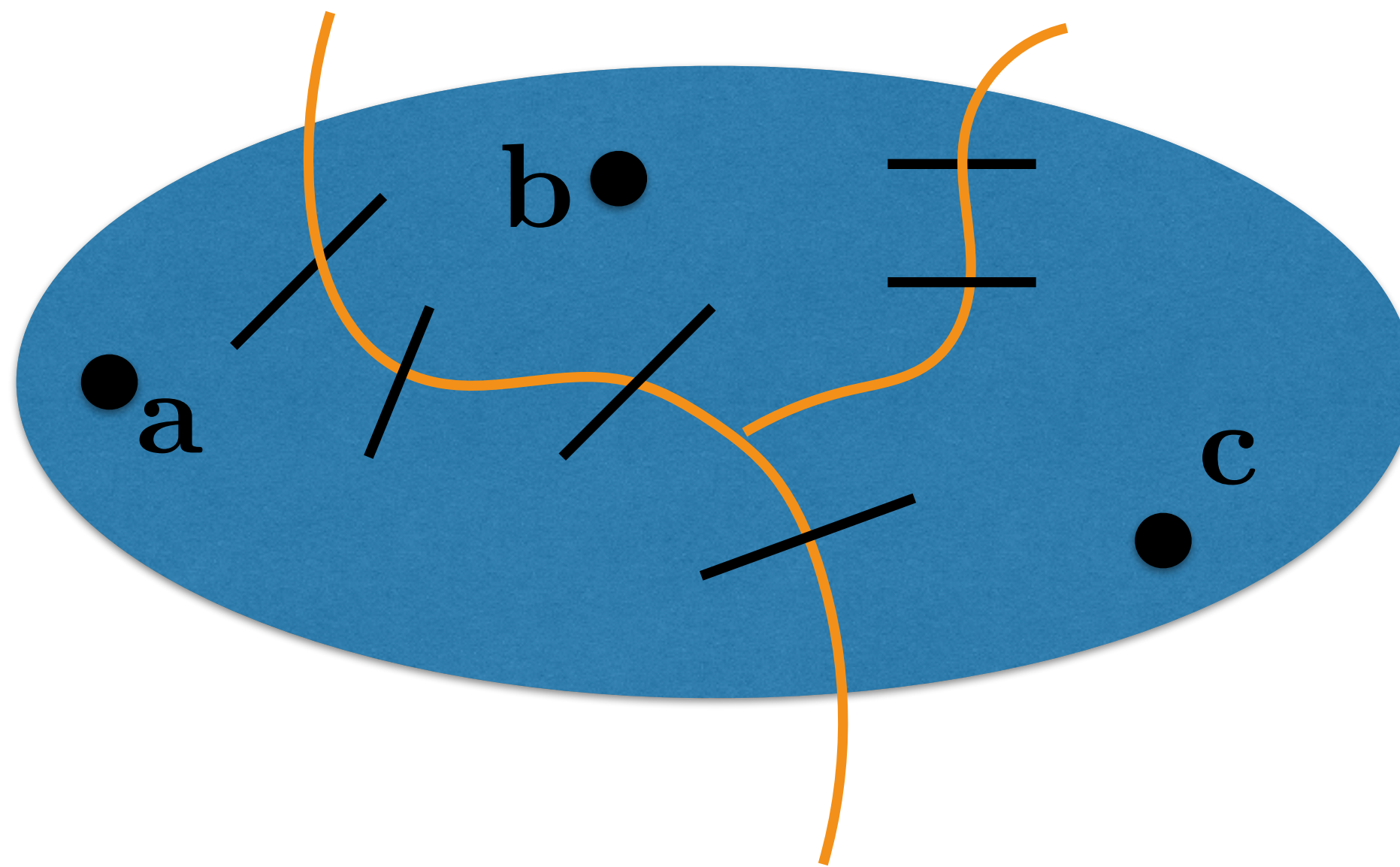




**$k=2$**   
**min cut**  
**in  $P$**



**$k=3$**   
**NP-hard**



Simple algorithm for  $k=3$

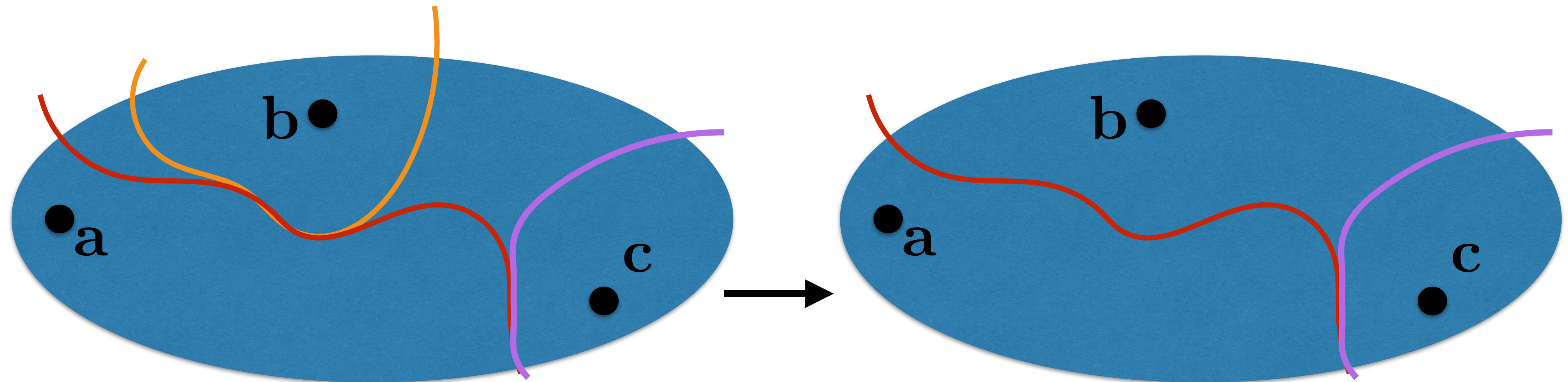
Terminals:  $a, b, c$

Output the two smallest of

$\{\text{Mincut}(a, \{b, c\}),$

$\text{Mincut}(b, \{c, a\}),$

$\text{Mincut}(c, \{a, b\})\}$



# Analysis

- It takes polynomial time...
- It's a correct multiway cut:  
a,b,c are separated from one another
- But how good is it?



## Cost of output

Output is at most

**$(2/3) (\text{Mincut}(a,bc) + \text{Mincut}(b,ac) + \text{Mincut}(c,ab))$**

- OPT is at least  $\text{Mincut}(a,bc)$
- OPT is at least  $\text{Mincut}(b,ac)$
- OPT is at least  $\text{Mincut}(c,ab)$

So OPT is at least

**$(\text{Mincut}(a,bc) + \text{Mincut}(b,ac) + \text{Mincut}(c,ab))/3$**

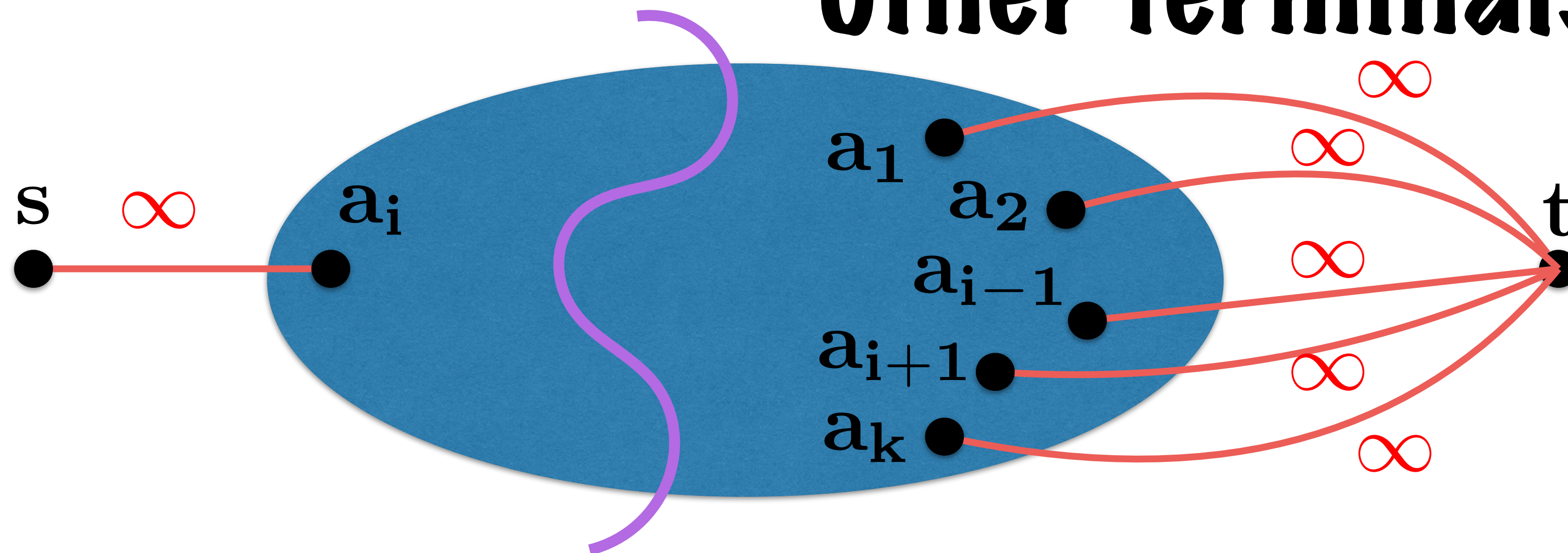
**Alg is a 2 approximation**

**Extension: algorithm for k**

**Terminals:  $a_1, a_2, \dots, a_k$**

**Output the  $k-1$  smallest of  
 $\{\text{Mincut}(a_i, \{a_1, \dots, a_{i-1}, a_{i+1}, \dots, a_k\}): i=1, 2, \dots, k\}$**

**To compute the min cut separation of  $a_i$  from the  
other terminals:**



**Mincut( $s, t$ )**

# Analysis of output

**The  $k-1$  smallest cuts cost less than a random choice of  $k-1$ , so**

$$\text{Cost}(\text{Output}) \leq \frac{k-1}{k} \sum_i \text{Mincut}(a_i, \{a_1, \dots, a_{i-1}, a_{i+1}, \dots, a_k\})$$



# Analysis of OPT

**Let  $F(i) = \{\text{edges of OPT that separate } a_i \text{ from some } a_j\}$**

**Consider  $e$  in OPT**

**$e$  separates some  $a_i$  from some  $a_j$   
 $e$  belongs to  $F(i)$  and to  $F(j)$**

**so**

$$\sum_i \text{Cost}(F_i) = 2\text{OPT}$$

**$F(i) = \{\text{edges of OPT that separate } a_i \text{ from some } a_j\}$**

**$F(i)$  separates  $a_i$   
from all other terminals  
so**

$$\text{Cost}(F_i) \geq \text{Mincut}(a_i, \{a_1, \dots, a_{i-1}, a_{i+1}, \dots, a_k\})$$

# Together

$$\begin{aligned} & \text{Cost}(\text{Output}) \\ & \leq \frac{k-1}{k} \sum_i \text{Cost}(\mathbf{F}_i) \\ & \leq 2\left(1 - \frac{1}{k}\right) \text{OPT} \end{aligned}$$

**k=2: Optimal**

**k=3: it's a 4/3 approximation**



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