

Correctness:

Claim:

For all two subsequences $(a_n)_{n \in \mathbb{N}}, (b_m)_{m \in \mathbb{N}}$ of characters. s.t $1 \leq n, m$

$f(n, m)$ returns the LCS of let (a_n) and (b_m)

Proof:(by induction on n and m)

Base case:

(3) Assume $n = 0$ or $m = 0$. Then f returns 0 , as required.

(4) Assume $n = 1$ and $m = 1$.

if $a_1 = b_1$ then f returns $1 + f(0,0) =_{(3)} 1 + 0 = 1$, as required.

Let $(a_n)_{n \in \mathbb{N}}, (b_m)_{m \in \mathbb{N}}$ be any two subsequences of characters.

(*) Assume that the claim holds for any $k < n, l < m$.

Assume that $a_n = b_m$.

(5) Therefore, we count 1 and add to that the return value of $f(n - 1, m - 1)$.

From **(*)**, $f(n - 1, m - 1)$ returns the LCS of

$(a_1, a_2, \dots, a_{n-1})$ and $(b_1, b_2, \dots, b_{m-1})$, say that this value is r , then of course that the LCS of (a_n) and (b_m) is $1 + r$. From here and **(5)** we get what is required.

Assume that $a_n \neq b_m$.

Therefore, the LCS of (a_n) and (b_m) can only be at

(6) $(a_1, a_2, \dots, a_{n-1})$ and (b_1, b_2, \dots, b_m) or at

(7) (a_1, a_2, \dots, a_n) and $(b_1, b_2, \dots, b_{m-1})$.

From **(*)**, $f(n - 1, m)$ return the LCS of the subsequence at **(6)**, say this value is r_1

and $f(n, m - 1)$ returns the LCS of the subsequence at **(7)**, say this value is r_2

Of course, the result is the maximum value between r_1 and r_2 .

That means, the result is $\max(f(n - 1, m), f(n, m - 1))$. As required.

We assumed that the claim holds for any $k < n, l < m$ and shows that the claim holds for n and m . Therefore, we completed the inductive step.

Conclusion: By mathematical induction, it's proved that the claim holds for all $1 \leq n, m$.

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