

Database Models: Beyond RDBMS

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The Battle of the Data Models

- ❖ Data models are perhaps the most important part of developing software
- ❖ They have a profound effect on:
 - how the software is written
 - how we think about the problem that we are solving.
- ❖ There are many different kinds of data model
 - Each data model embodies assumptions about how it is going to be used.
- ❖ We will now look at a range of general-purpose data models for data storage and querying

When we have some data...

❖ Relational Databases solve most data problems

Why?

❖ Persistence

- We can store data, and it will remain stored!

❖ Integration

- We can integrate lots of different apps through a central DB

❖ SQL

- Standard, well understood, very expressive

❖ Transactions

- ACID transactions, strong consistency

Transactions – ACID Properties

❖ Atomic

- All of the work in a transaction completes (commit) or none of it completes

❖ Consistent

- A transaction transforms the database from one consistent state to another consistent state. Consistency is defined in terms of constraints.

❖ Isolated

- The results of any changes made during a transaction are not visible until the transaction has committed. Concurrent interactions behave as though they occurred serially

❖ Durable

- The results of a committed transaction survive failures

The Relational Model

- ❖ The relational model, proposed by Edgar Codd in 1970, is still the best-known data model today.
 - data is organized into relations (in SQL: tables), where each relation is an unordered collection of tuples (rows).
- ❖ The dominance of relational databases has been around for +30 years.
 - An “eternity” in computing history.
- ❖ Other databases at that time forced application developers to think a lot about the internal representation of the data in the database.
 - The goal of the relational model was to hide that implementation detail behind a cleaner interface.

Rivals of the Relational Model

- ❖ Over the years, there have been many competing approaches to data storage and querying.
 - Object databases came and went again in the late 1980s and early 1990s.
 - XML databases appeared in the early 2000s, but have only seen niche adoption.
- ❖ Much of what you see on the web today is still powered by relational databases
 - Online publishing, discussion, social networking, e-commerce, games, software-as-a-service productivity applications, or much more.
- ❖ Now, NoSQL is the most recent attempt to overthrow the relational model's dominance.

Current trends and Issues

- ❖ A few key trends and issues have motivated change in relational data storage technologies
 - ...In use cases
 - ...In technology
- ❖ Key trends include:
 - Increasing **volume of data and traffic**
 - More complex data connectedness
- ❖ Key Issues include:
 - The **impedance mismatch** problem

Impedance Mismatch

- ❖ **Object** Orientation

- based on software engineering principals

- ❖ **Relational** Paradigms

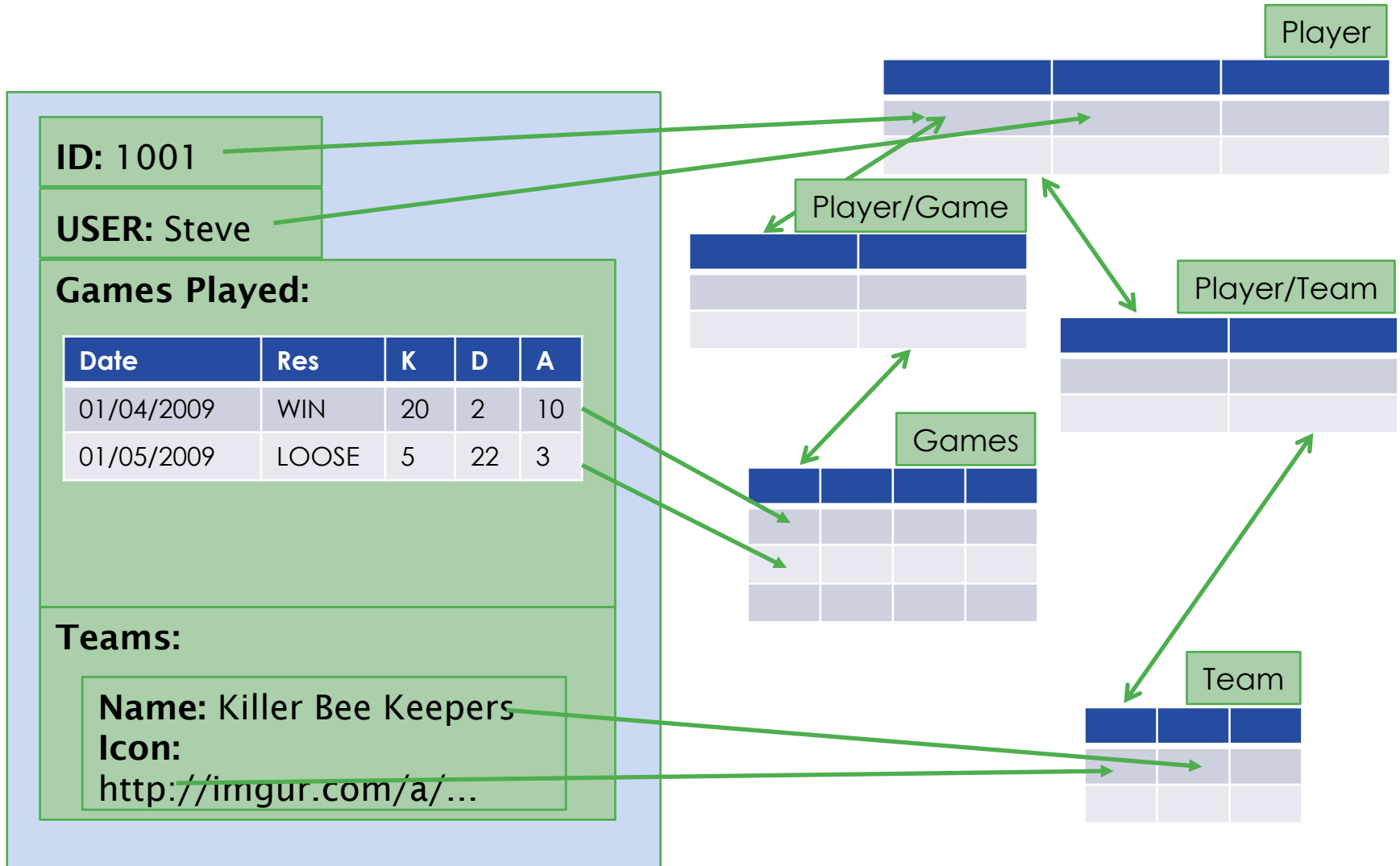
- based on mathematics and set theory

- ❖ Mapping from one world to the other has problems

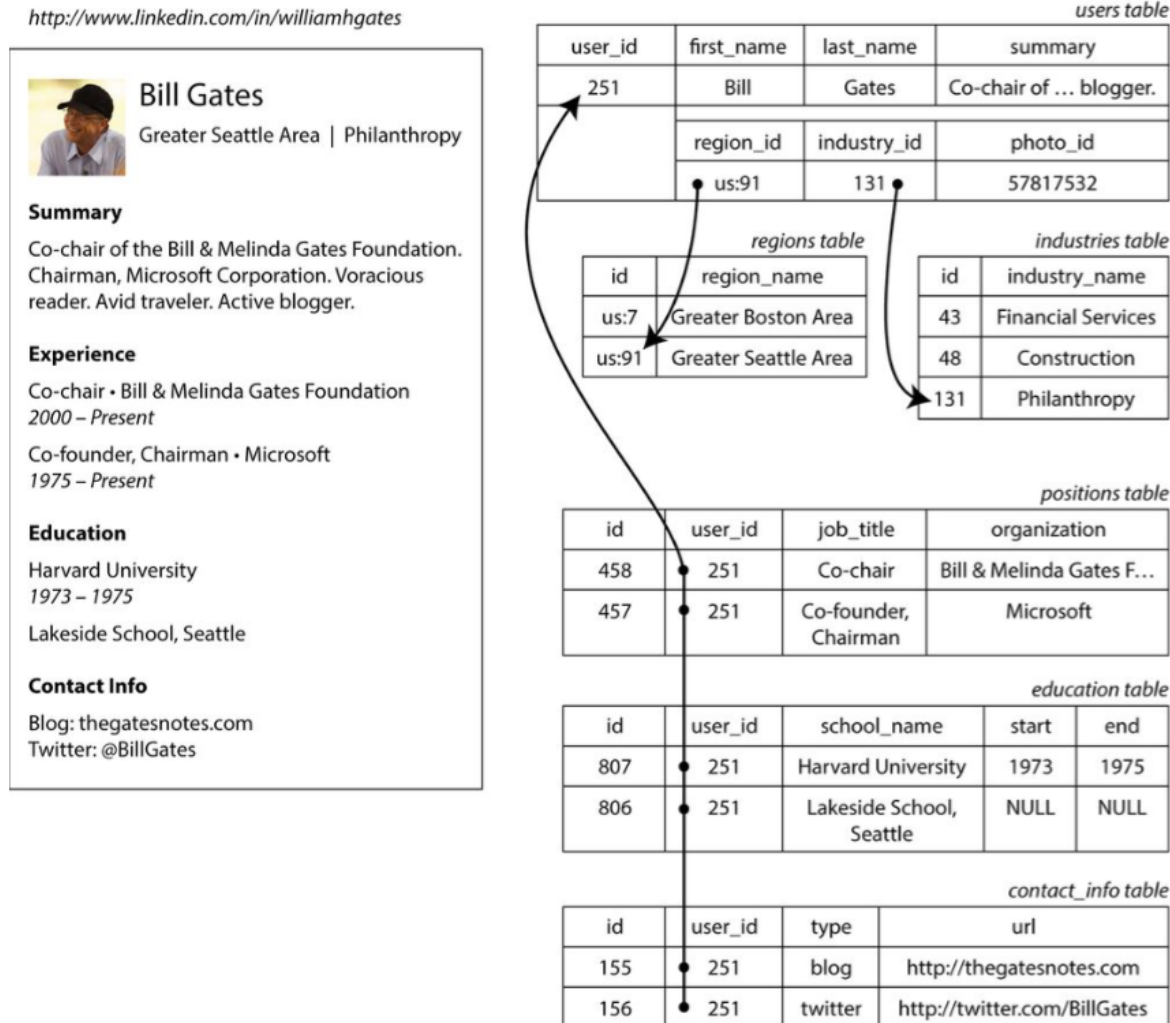
- ❖ To store data persistently in modern programs a single logical structure must be split up

- The nice word is **normalised**

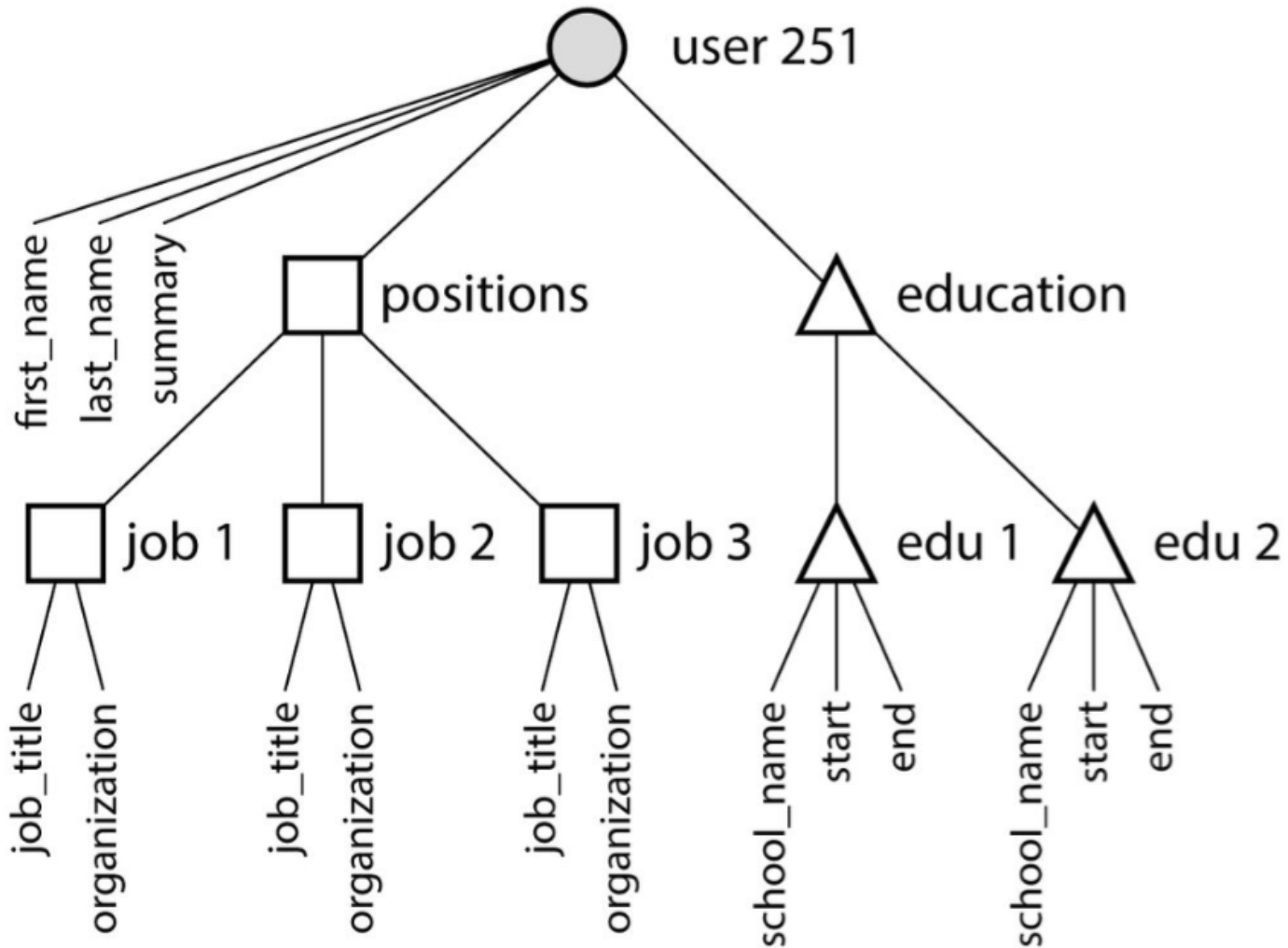
Impedance Mismatch – example



Impedance Mismatch – example



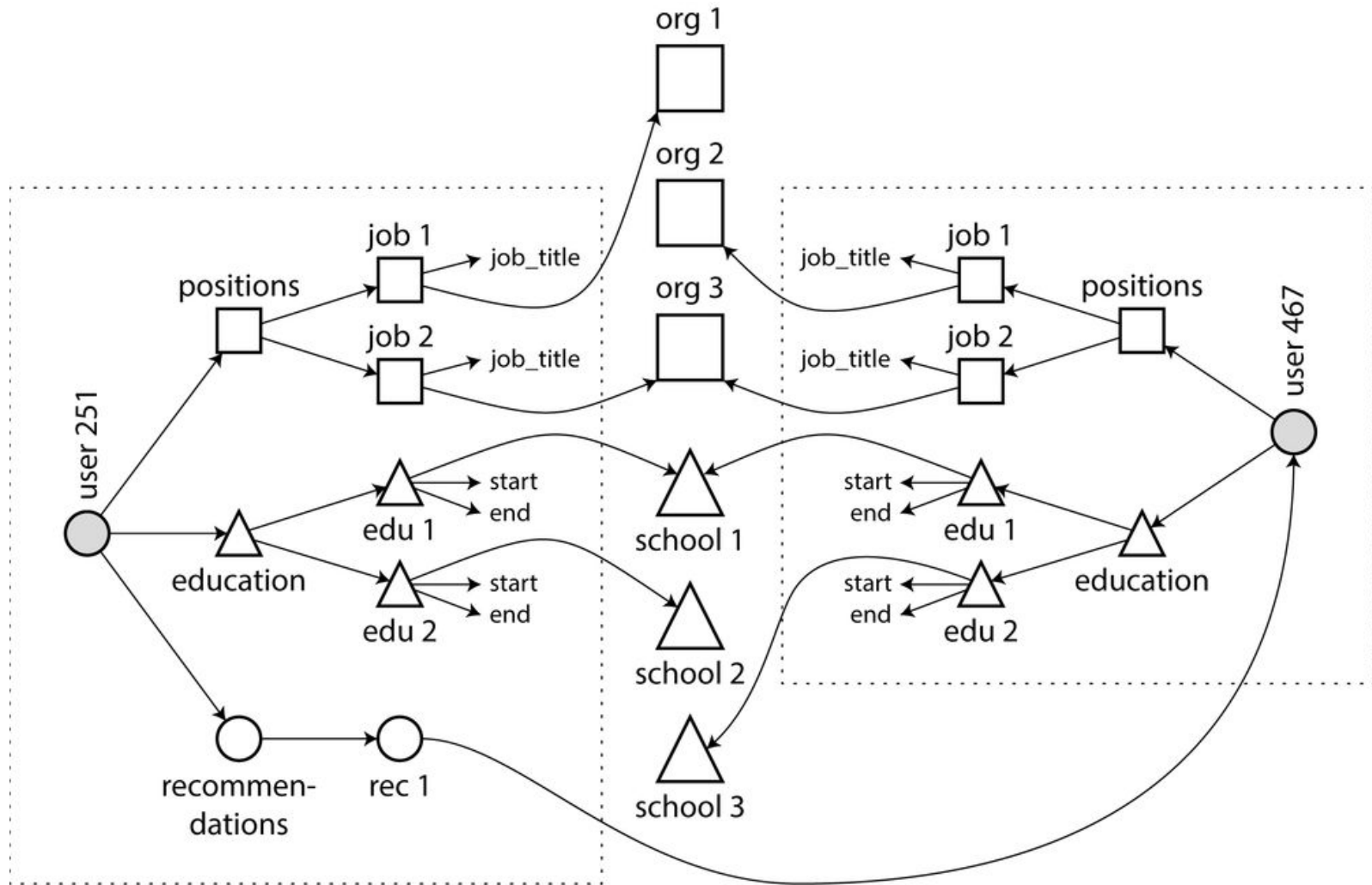
One-to-Many relations



Normalization

- ❖ Why IDs (region_id, industry_id, ..) and not plain-text?
 - Consistent style and spelling across profiles,
 - Avoiding ambiguity, e.g. if there are several cities with the same name,
 - The name is stored only in one place, so it is easy to update,
 - Simplify translation into other languages,
- ❖ A database in which entities like region and industry are referred to by ID is called **normalized**.
- ❖ A database that duplicates the names and properties of entities on each document is **denormalized**.

Many-to-Many relationships



Increased Data Volume

- ❖ We are creating, storing, processing more data than ever before!
 - *“From 2005 to 2020, the digital universe will grow by a factor of 300, from 130 exabytes to 40,000 exabytes, or 40 trillion gigabytes (more than 5,200 gigabytes for every man, woman, and child in 2020).
From now until 2020, the digital universe will about double every two years.”,*

IDC – The Digital Universe in 2020.

<http://www.emc.com/leadership/digital-universe/index.htm>

Increased Data Connectivity

- ❖ The data we're producing has fundamentally changed
 - from isolated Text Documents (early 1990s)
 - ... to html pages with links (early web)
 - ... to blogs with pingback, RSS feeds (web 2.0)
 - ... to social networks (... add links between people)
 - ... to massive linked open data sets (web 3.0... one of them anyway)

Dealing with data size Trends

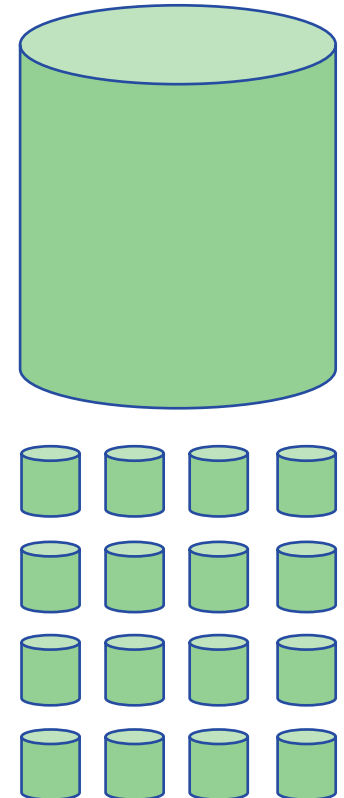
❖ Two options when dealing with these trends:

❖ Build Bigger Database machines

- This can be expensive
- Fundamental limits to machine size

❖ Build Clusters of smaller machines

- Lots of small machines (commodity machines)
- Each machine is cheap, potentially unreliable
- Needs a DBMS which understands clusters



RDBMS have fundamental issues

- ❖ In dealing with (horizontal) scale
 - Designed to work on single, large machines
 - Difficult to distribute effectively
- ❖ More subtle: An Impedance Mismatch
 - We create logical structures in memory
 - and then rip them apart to stick it in an RDBMS
 - The RDBMS data model often disjoint from its intended use
 - (Normalisation sucks sometimes)
 - Uncomfortable to program with (joins and ORM etc.)

The NoSQL Movement

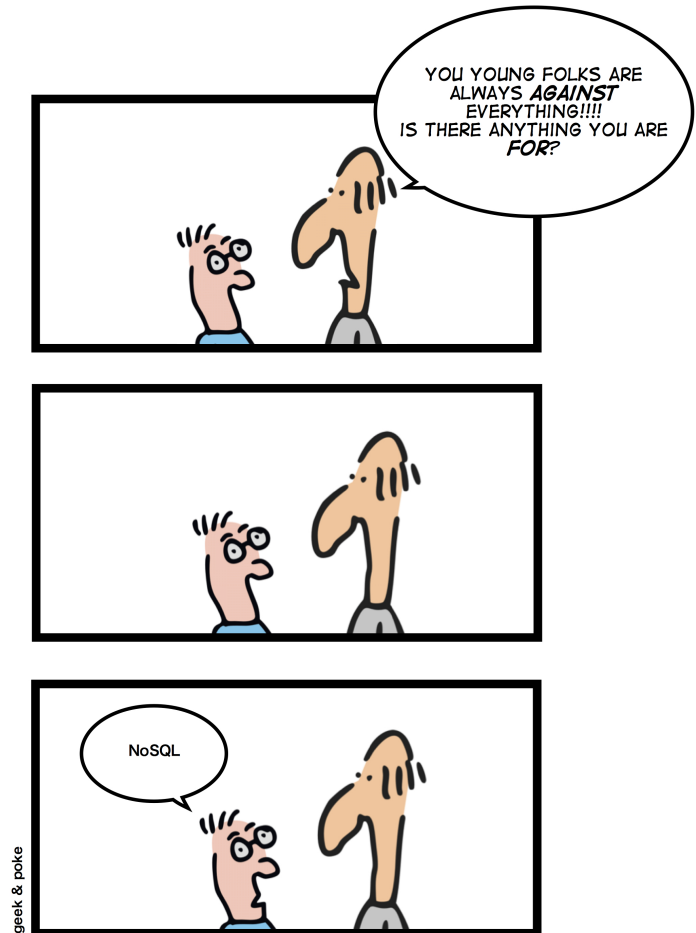
NoSQL

- ❖ The term NoSQL is unfortunate, since it doesn't refer to any technology
 - “Not only SQL”
- ❖ Nevertheless, the term struck a nerve, and quickly spread through the web startup community and beyond.
- ❖ Several interesting database systems are now associated with the #NoSQL hashtag.

The NoSQL movement

❖ Key attributes include:

- **Non-Relational**
 - They can be, but aren't good at it
- **Simple API**
 - No Join
- **BASE & CAP Theorem**
 - No ACID requirements
- **Schema-free**
 - Implicit schema, application side
- **Inherently Distributed**
 - Some more so than others
- **Open Source**
 - mostly



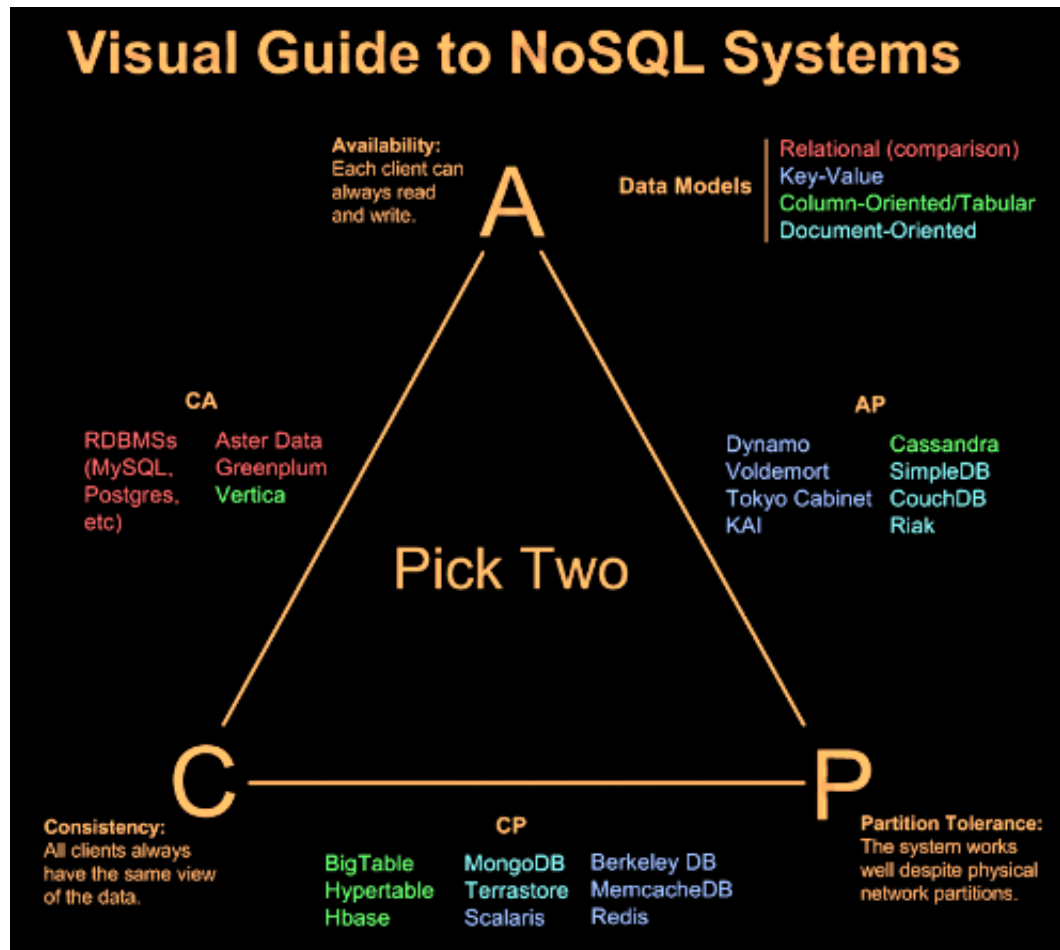
BASE Transactions

- ❖ Acronym contrived to be the opposite of ACID
 - Basically Available,
 - Soft state,
 - Eventually Consistent
- ❖ Characteristics
 - Weak consistency – stale data OK
 - Availability first
 - Best effort
 - Approximate answers OK
 - Aggressive (optimistic)
 - Simpler and faster

Brewer's CAP Theorem

- ❖ A distributed system can support only two of the following characteristics:
 - **Consistent**
 - writes are atomic, all subsequent requests retrieve the new value
 - **Available**
 - The database will always return a value so long as the server is running
 - **Partition Tolerant**
 - The system will still function even if the cluster network is partitioned (i.e. the cluster loses contact with parts of itself)
- ❖ The overly stated well cited issue is:
 - We can only ever build an algorithm which satisfies 2 of 3.
 - Formal Proof by Gilbert and Lynch:
<http://portal.acm.org/citation.cfm?doid=564585.564601>

CAP Theorem



Types of NoSQL Databases

❖ Core types

- **Key-value** stores
- **Document** stores
- **Column** stores
- **Graph** databases

❖ Non-core types

- **Object** databases
- Native **XML** databases
- **RDF** stores
- ...

Key-Value Databases – Basics

❖ Data model

- The most simple NoSQL database type
- Works as a simple hash table (mapping)

❖ Key-value pairs

- **Key** (id, identifier, primary key) – usually a string.
- **Value**: can be anything (text, structure, image, etc.) – a black box for the database system.

❖ Query patterns

- Create, update or remove value for a given key
- Get value for a given key

❖ Characteristics

- great performance, easily scaled, ...
- not for complex queries nor complex data

Key-Value Databases – Basics

❖ Suitable use cases

- Session data, user profiles, user preferences, shopping carts, ...
 - I.e. when values are only accessed via keys

❖ When not to use

- Relationships among entities
- Queries requiring access to the content of the value part
- Set operations involving multiple key-value pairs

❖ Examples

- Redis, MemcachedDB, Riak KV, Amazon SimpleDB, Berkeley DB, Oracle NoSQL, LevelDB, Project Voldemort
- *Multi-model*: OrientDB, ArangoDB

Document Databases – Basics

- ❖ Data model - Documents
 - Self-describing complex data structure
 - **Hierarchical tree structures** (JSON, XML, ...)
 - Scalar values, maps, lists, sets, nested documents, ...
 - Identified by a unique identifier (key, ...)
- ❖ Document data stores understand their documents
 - Queries can run against values of document fields
 - Indexes can be constructed for document fields
- ❖ Query patterns
 - Create, update or remove a document
 - Retrieve documents according to complex queries
- ❖ Difference from Key-Value stores
 - Extended key-value stores. The value part is examinable!

Document Databases – Basics

```
{
  "_id": "1",
  "name": "steve",
  "games_owned": [
    {"name": "Super Meat Boy"},
    {"name": "FTL"},
  ],
}
```

```
{
  "_id": "2",
  "name": "darren",
  "handle": "zerocool",
  "games_owned": [
    {"name": "FTL"},
    {"name": "Assassin's Creed 3", "dev": "ubisoft"},
  ],
}
```

Document Databases – Basics

❖ Suitable use cases

- Event logging, content management systems, blogs, web analytics, e-commerce applications, ...
- I.e. for **structured documents with similar schema**

❖ When not to use

- Set operations involving multiple documents
- Design of document structure is constantly changing
 - I.e. when the required level of granularity would outbalance the advantages of aggregates

❖ Examples

- MongoDB, Couchbase, Amazon DynamoDB, CouchDB, RethinkDB, RavenDB, Terrastore
- *Multi-model*: MarkLogic, OrientDB, OpenLink Virtuoso, ArangoDB

Column Databases – Basics

❖ Data model

- Column family (table)
 - Table is a collection of similar rows (not necessarily identical)
- Row
 - Row is a collection of columns - should encompass a group of data that is accessed together
 - Associated with a unique row key
- Column
 - Column consists of a column name and column value (and possibly other metadata records)
 - Scalar values, but also flat sets, lists or maps may be allowed

❖ Query patterns

- Create, update or remove a row within a given column family
- Select rows according to a row key or simple conditions

Column Databases – Basics

❖ Suitable use cases

- Event logging, content management systems, blogs, ...
 - I.e. for structured flat data with similar schema
- Batch processing via mapreduce

❖ When not to use

- ACID transactions are required
- Complex queries: aggregation (SUM, AVG, ...), joining, ...
- Early prototypes: i.e. when database design may change

❖ Examples

- Apache Cassandra, Apache HBase, Apache Accumulo, Hypertable, Google Bigtable

Graph Databases – Basics

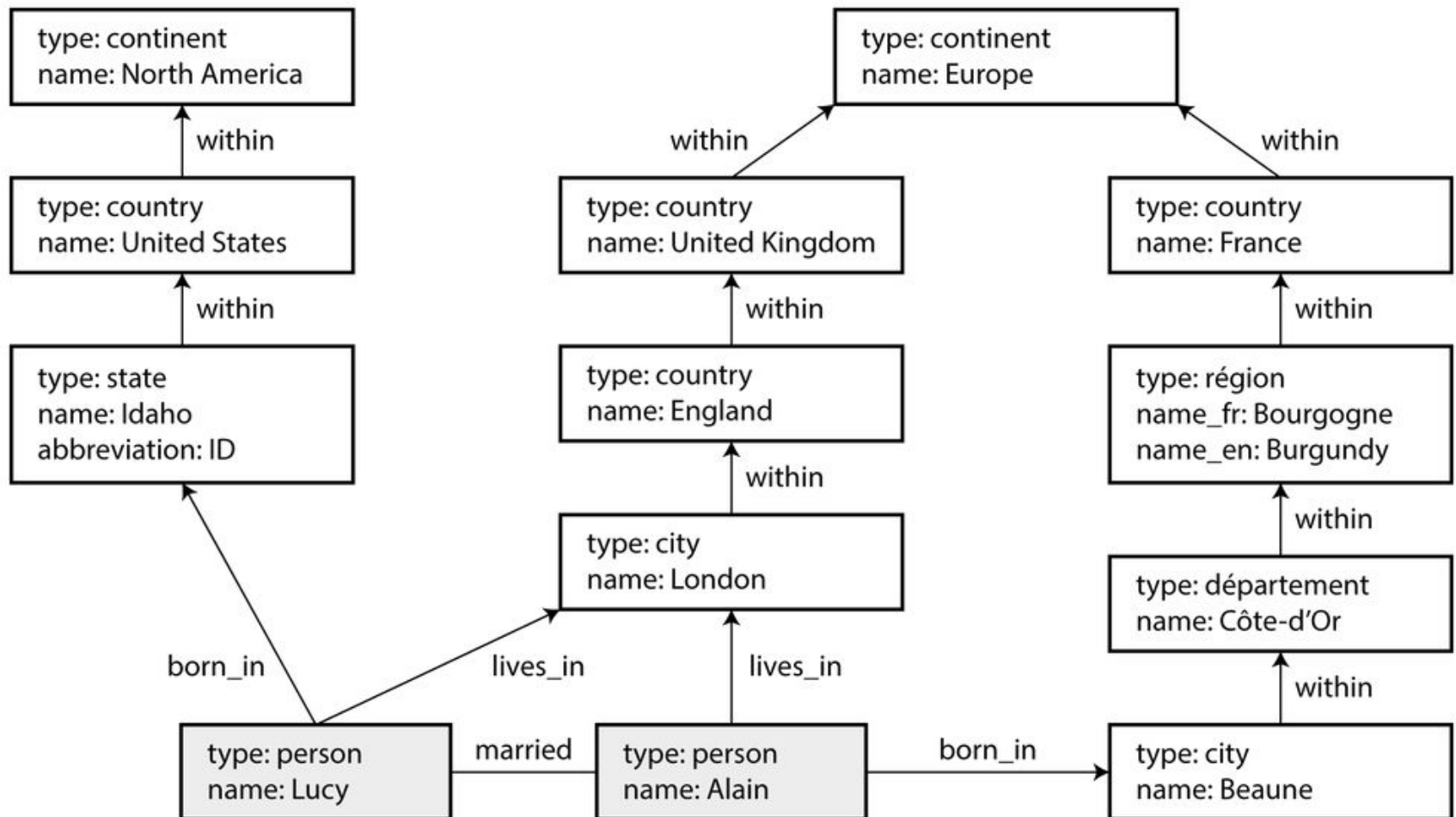
❖ Data Model

- Focus on modelling graphs' structure and properties
- Directed / undirected graphs, i.e. collections of ...
 - nodes (vertices) for real-world entities, and
 - relationships (edges) between these nodes
- Both the nodes and relationships can have properties

❖ Query patterns

- Create, update or remove a node / relationship in a graph
 - Graph algorithms (shortest paths, spanning trees, ...)
 - General graph traversals
 - Sub-graph queries or super-graph queries
 - Similarity based queries (approximate matching)

Graph Databases – Basics



Graph Databases – Basics

❖ Suitable use cases

- Social networks, routing, dispatch, and location-based services, recommendation engines, chemical compounds, biological pathways, linguistic trees, ...

❖ When not to use

- Extensive batch operations are required
 - Multiple nodes / relationships are to be affected
- Too large graphs to be stored
 - Graph distribution is difficult or impossible at all

❖ Examples

- Neo4j, Titan, Apache Giraph, InfiniteGraph, FlockDB
- *Multi-model*: OrientDB, OpenLink Virtuoso, ArangoDB

Native XML Databases – Basics

❖ Data model

- XML documents
- **Tree structure** with nested elements, attributes, and text values (beside other less important constructs)
- Documents are organized into collections

❖ Query languages

- **XPath**: XML Path Language (navigation)
- **XQuery**: XML Query Language (querying)
- **XSLT**: XSL Transformations (transformation)

❖ Examples

- Sedna, Tamino, BaseX, eXist-db
- *Multi-model*: MarkLogic, OpenLink Virtuoso

RDF Databases – Basics

❖ Data model

– RDF **triples**

- Components: **subject**, **predicate**, and **object**
- Each triple represents a statement about a real-world entity

– Triples can be viewed as graphs

- Vertices for subjects and objects
- Edges directly correspond to individual statements

❖ Query language

– **SPARQL**: SPARQL Protocol and RDF Query Language

❖ Examples

- Apache Jena, rdf4j (Sesame), Algebraix
- *Multi-model*: MarkLogic, OpenLink Virtuoso

NoSQL Databases

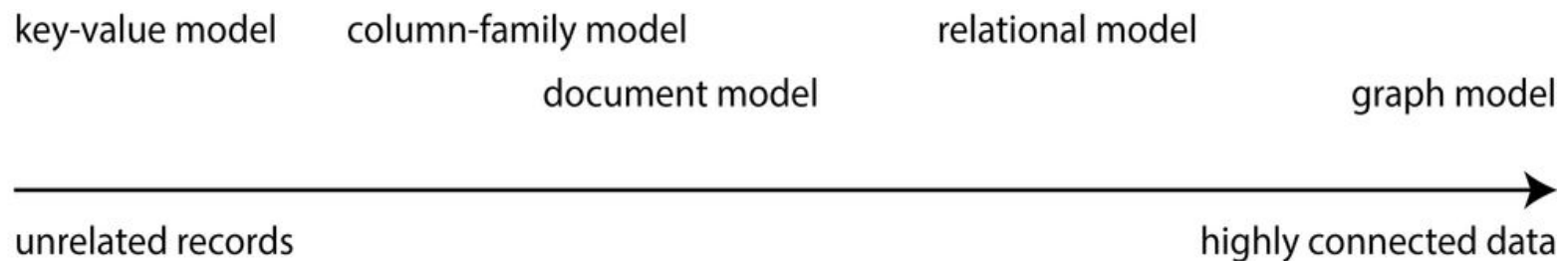
- ❖ The end of relational databases?
- ❖ **Certainly no!**
 - They are still suitable for most projects (90%)
 - Familiarity, stability, feature set, available support, ...
- ❖ However, we should also consider different database models and systems
 - **Polyglot persistence** = usage of different data stores in different circumstances

Databases and data connectivity

❖ Relational model

❖ NoSQL models

- Key-value stores
- Document stores
- Column stores
- Graph databases



What next?

❖ Basic principles

- Data formats: JSON, YAML, XML, RDF, ...
- Distribution, scaling, sharding, replication, consistency
- Parallelism, transactions, visualization, processing of graphs

❖ NoSQL technologies: principles, models, interfaces, languages, ...

- Core databases: Redis, MongoDB, Cassandra, Neo4j
- MapReduce: Apache Hadoop

Resources

- ❖ Martin Kleppmann, ***Designing Data-Intensive Applications***, O'Reilly Media, Inc., 2017.
- ❖ Pramod J Sadalage and Martin Fowler, ***NoSQL Distilled*** Addison-Wesley, 2012.
- ❖ Eric Redmond, Jim R. Wilson. ***Seven databases in seven weeks***, Pragmatic Bookshelf, 2012.
- ❖ Hector Garcia-Molina, Jeffrey D. Ullman, Jennifer Widom, ***Database systems: the complete book*** (2nd Ed.), Pearson Education, 2009.