

## “A New Ordering for Efficient Sphere Decoding” Constrained ILS Reduction

Consider the ILS problem,  $\min_{x \in X} \|Ax - y\|$ . Define  $G = (A^{-1})^T$  and  $G_i$  references the  $i^{th}$  column of  $G$ . The algorithm starts with all elements of  $x$  unfixed and an ‘index’ set  $Index = 1 \dots m$  which will be used to reference columns of  $G$ .

In the first iteration, we find the column of  $G$  which maximizes the distance to the second nearest integer point. Meaning that for each  $i \in Index$  we compute  $a = \arg \min_{x \in X} |y^T G_i - x|$  which gives the  $a = x_m$  that minimizes the partial residual assuming that we choose column  $i$  to be the  $m^{th}$  column in the final matrix  $A$ . We then compute  $b = \arg \min_{x \in X \setminus a} |y^T G_i - x|$  which gives the  $b = x_m$  that has the second smallest partial residual. We compute the distance between the target vector  $y$  and the affine set defined by fixing  $x_m = b$  as  $dist = \|(1/(G_i^T G_i)(G_i G_i^T))(y - A_i * b)\|$  where the term  $1/(G_i^T G_i)(G_i G_i^T)$  is just the orthogonal projector that projects onto the  $i^{th}$  column of  $G$ . We want to keep track of which value of  $i$  gives the maximum value for  $dist$ .

After completing the above process we should have values for  $a$ ,  $b$  and  $i$ , where  $i$  was the index giving the maximum value for  $dist$ . We then set the  $m^{th}$  entry in a permutation vector to move the  $i^{th}$  column to the  $m^{th}$ , record the value of  $a$  because it will be the  $m^{th}$  entry of the babai point and remove  $i$  from the Index set  $Index_i = []$ .

Since we have now fixed  $x_m$  we must project and shift the target vector as follows,  $y = (y - A_i * a) - (1/(G_i^T G_i)(G_i G_i^T))(y - A_i * a)$ . This is equivalent to applying the constraint  $x_m = a$ .

We must also project all of the remaining columns of  $G$  as follows,  $\forall j \in Index \quad G_j = G_j - (1/(G_i^T G_i)(G_i G_i^T)) * G_j$ .

The algorithm proceeds to repeat the above process  $m$  times. So that in each iteration there is one less column of  $G$  in the Index set. We are finished when the index set is empty.