

تجزئة لغة

$$1) L = \{aaa, aaab, aaabb, aaabbb, aaabbbb, \dots\} \quad (1)$$

$$L_1 = aaaa^*(a + b + bb + bbb + bbbb)$$

$$2) L(ab+b)^* L(b) L(a+ab)^*$$

$$(L(ab) \cup L(b))^* L(b) (L(a) \cup L(ab))^*$$

$$\{b, ab\}^* \{b\} \{a, ab\}^* = \{b, ba, bab, baba, babab, \dots\}$$

3)

$$(n+m) \text{ is odd} \rightarrow \begin{matrix} \underline{n} \text{ odd} & \underline{m} \text{ even} \end{matrix} \rightarrow a(aa)^*(bb)^*$$

$$\hookrightarrow \begin{matrix} \underline{n} \text{ even} & \underline{m} \text{ odd} \end{matrix} \rightarrow (aa)^* b(bb)^*$$

$$r = a(aa)^*(bb)^* + (aa)^* b(bb)^*$$

$$4) L((a+b+bb)aa^+) = L(a+b+bb)L(aa^+)$$

$$= (L(a) \cup L(b) \cup L(bb))L(a)L(a^+)$$

$$(\{a\} \cup \{b\} \cup \{bb\})\{a\}\{a\}^+$$

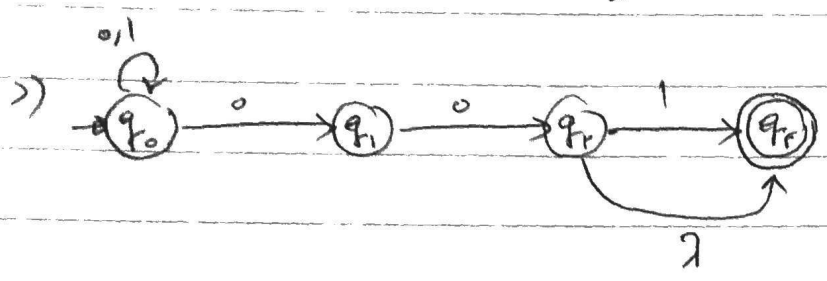
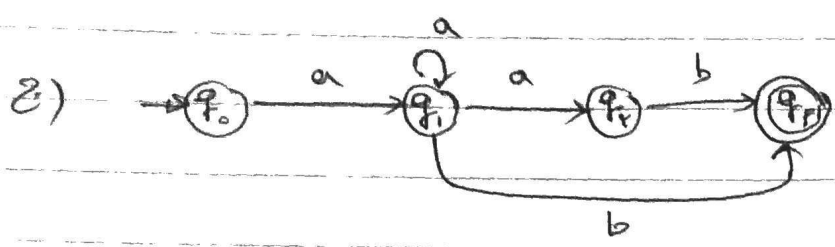
$$\{a, b, bb\}\{a\}\{a, aa, aaa, \dots\}$$

$$= \{aa, aaa, aaaa, \dots, baa, baaa, \dots, bb aa, bb aaa, \dots\}$$

$$L(n) = \{b^m a^n : n \geq 1, m \geq 0\}$$

۱) $r = 0^*1(00^*1 + (10^*1)^*)^*$

۲) $r = a^*b$



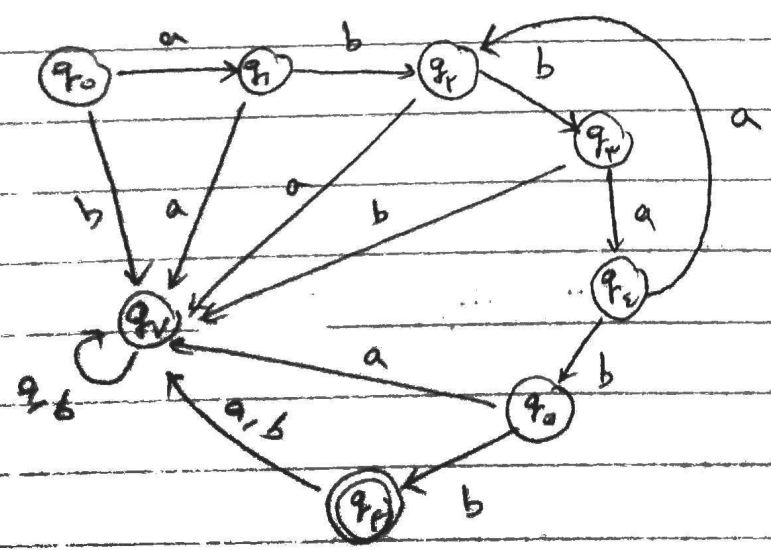
۱) راه اول: نوشتن اشتقاق تدریجی زیر

$S \rightarrow abA \rightarrow abbaB \rightarrow abbabb$

$S \rightarrow abA \rightarrow abbaB \rightarrow abbaaA \rightarrow abbaababB \rightarrow abbaabbabb$

$S \rightarrow abA \rightarrow \dots \rightarrow abbaabaababb$

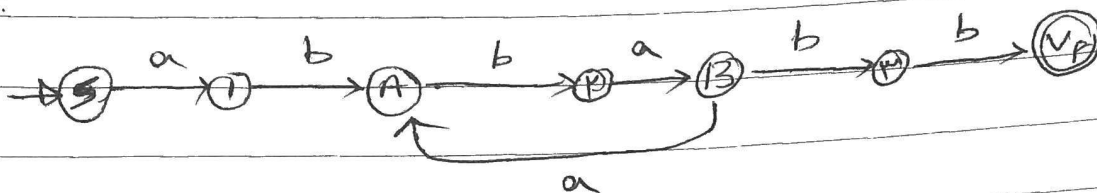
$L(M) = \{w: (abb)(aab)^n(abb); n \geq 0\}$



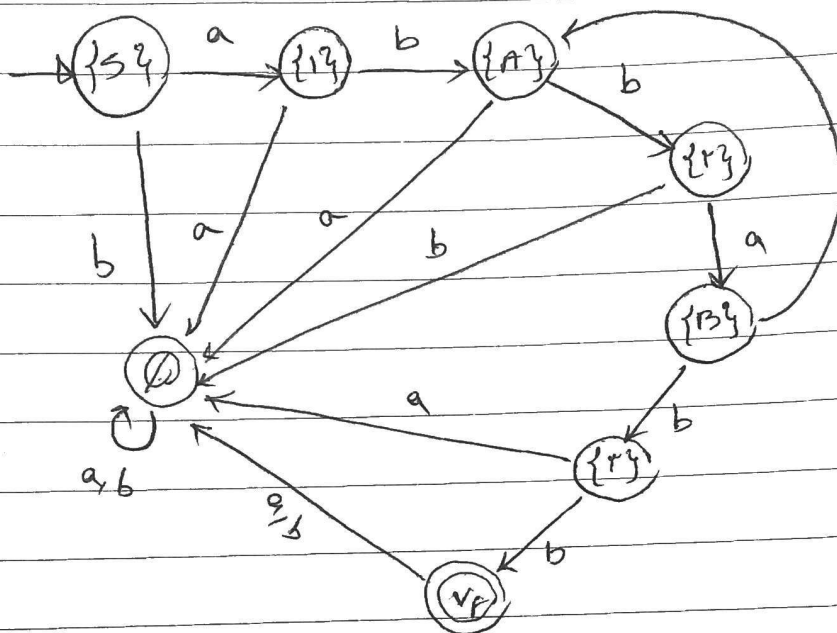
الف)

3) به رسم نوشتن NFA از یک گرامر وابسته به DFA

NFA:



DFA:



جواب:

ب) $S \rightarrow abS \rightarrow abA \rightarrow abbaB \rightarrow abbaabb$

$S \rightarrow abS \rightarrow abA \rightarrow abbaB \rightarrow abbaaA \rightarrow abbaababb$

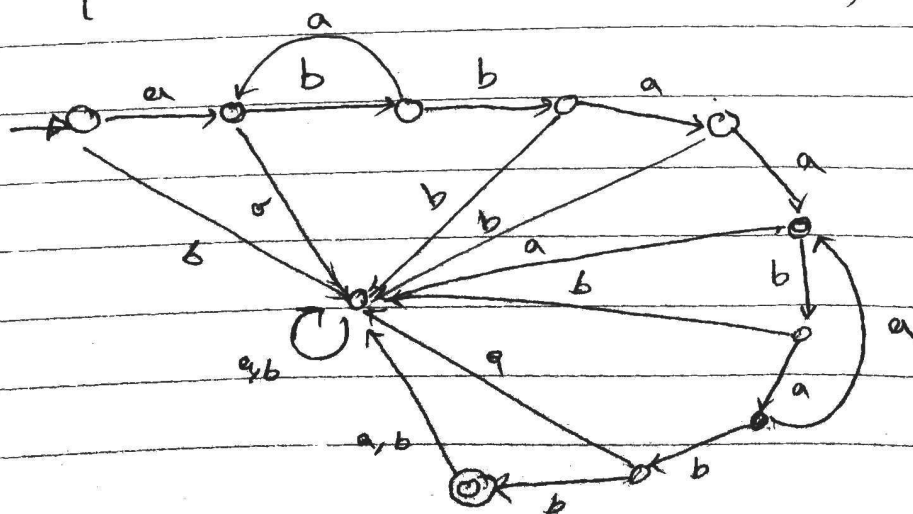
$S \rightarrow abS \rightarrow ababS \rightarrow abababS$

$\rightarrow abababA$

$\rightarrow abababba$

$\rightarrow abababbaabababb$

$L = \{ (ab)^n (abb) (aab)^m (abb) : n, m \geq 0 \}$



DFA به رسم نوشتن DFA
با استفاده از DFA

2) $L = \{a^n b^m : n \leq 2, m \leq 2\}$

(4)

$R = (\epsilon + a + aa + aaa)(\epsilon + b + bb + bbb + bbbb)$

مشتق: $S \rightarrow aA | bB | \epsilon$

$A \rightarrow aB | aaB | abB | aabB | \epsilon$

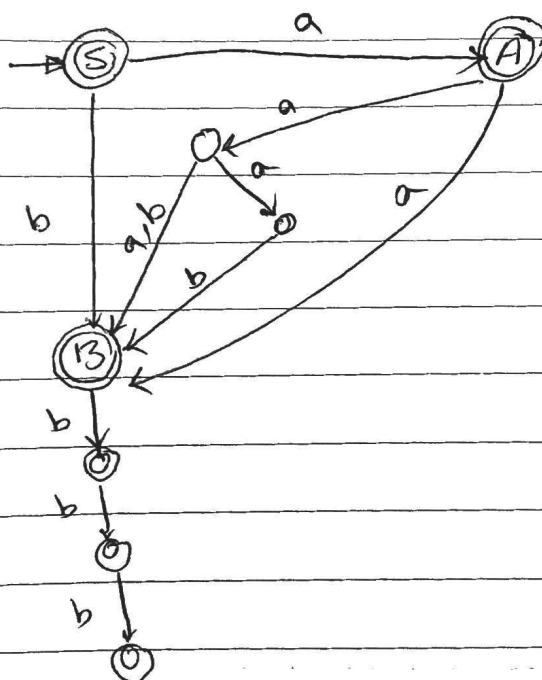
$B \rightarrow b | bb | bbb | \epsilon$

مشتق: $S \rightarrow Bb | Aa | \epsilon$

$B \rightarrow Ab | Abb | Abbb | Aab | Aabb | Aabbb | \epsilon$

$A \rightarrow a | aa | \epsilon$

NFA:



این NFA دیگر ادیتوری
 که این زبان را تولید
 کند قابل قبول است

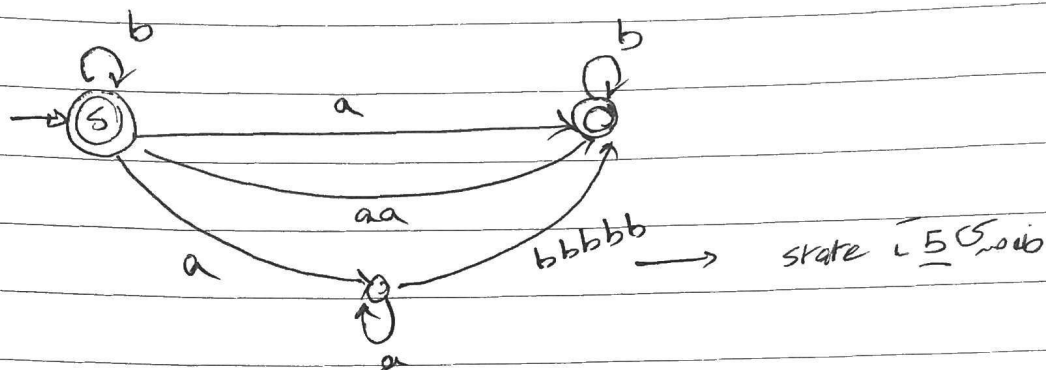
1) $L_2 = \{a^n b^m, n \leq m\} \cup \{a^n b^m, n > m\}$

$r_1 = b^* + ab^* + aab^*$

$r_2 = a^* b b b b b b^*$

$r = r_1 \cup r_2 = b^* + ab^* + aab^* + a^* b b b b b b^*$

NFA:



Deriv:

$S \rightarrow bS \mid \lambda$

$S \rightarrow aE \mid aA \mid aaF$

$A \rightarrow aA \mid b b b b b F$

Deriv: $S \rightarrow Sb \mid \lambda$

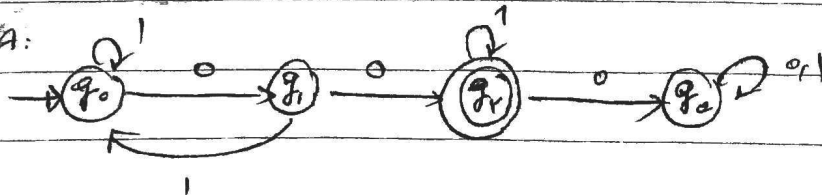
$S \rightarrow Fa \mid Faa \mid Abbbb$

$A \rightarrow Aa \mid Fa$

$F \rightarrow Fb \mid \lambda$

2) ~~$r = (1^* 0 1)^* 0 0 (1^*)$~~ $r = (1^* 0 1)^* 0 0 (1^*)$

NFA:



Deriv: $S \rightarrow 1S \mid A$

$A \rightarrow 0F \mid 1S$

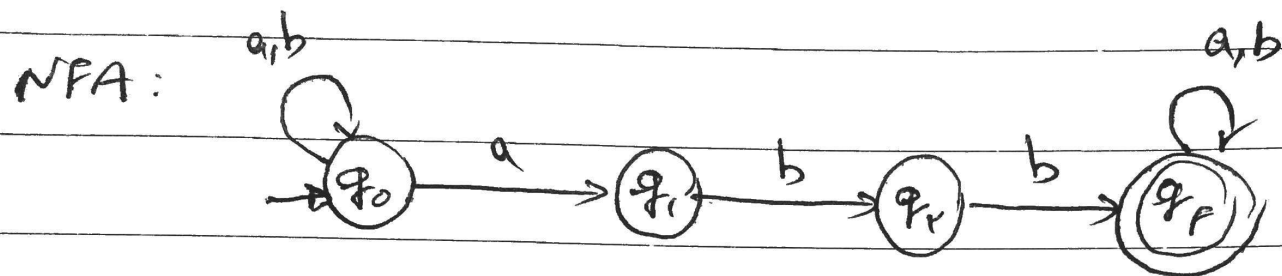
$F \rightarrow 1F \mid \lambda$

Deriv: $S \rightarrow S1 \mid A0$

$A \rightarrow F0$

$F \rightarrow F1 \mid A1 \mid \lambda$

$$> r = (a+b)^* abb(a+b)^*$$



$$\text{grammar: } S \rightarrow aS \mid bS \mid aB$$

$$B \rightarrow bbF$$

$$F \rightarrow aF \mid bF \mid \lambda$$

$$\text{grammar: } S \rightarrow Sa \mid Sb \mid Abb$$

$$A \rightarrow Fa$$

$$F \rightarrow Fa \mid Fb \mid \lambda$$