

# 학습 목표

트리(Tree) 구조와 용어에 대해 학습하고 간단하게 구현할 수 있다



#### Data Structures in Python Chapter 7 - 1

- Tree Introduction
- Tree Traversals
- Tree Algorithms

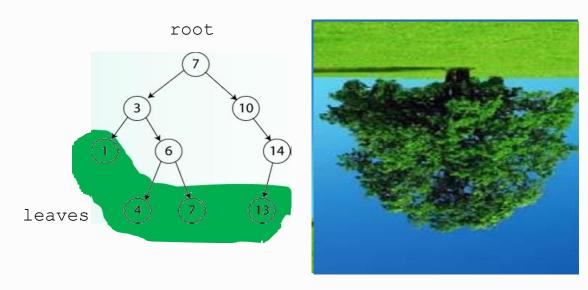
### Agenda & Readings

- Agenda
  - Tree Terminology
  - Binary Tree Properties
  - Binary Tree and Node Representation
- Reference:
  - Problem Solving with Algorithms and Data Structures
  - Chapter 6 Tree

#### What is a Tree?

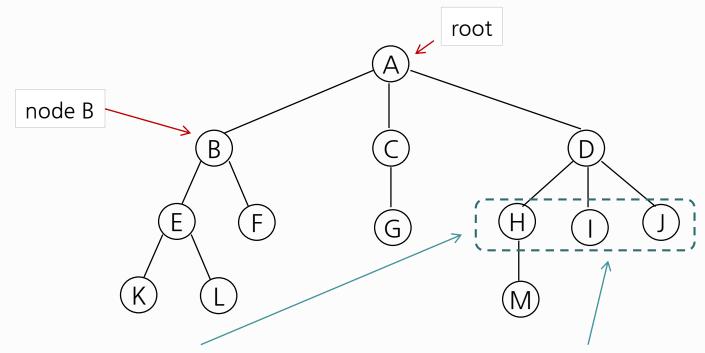
- A non-linear data structure
- An abstraction for a hierarchical structure

 It is defined as a set of points called nodes and a set of lines called edges where an edge connects two distinct nodes.



#### Introduction - Terminology

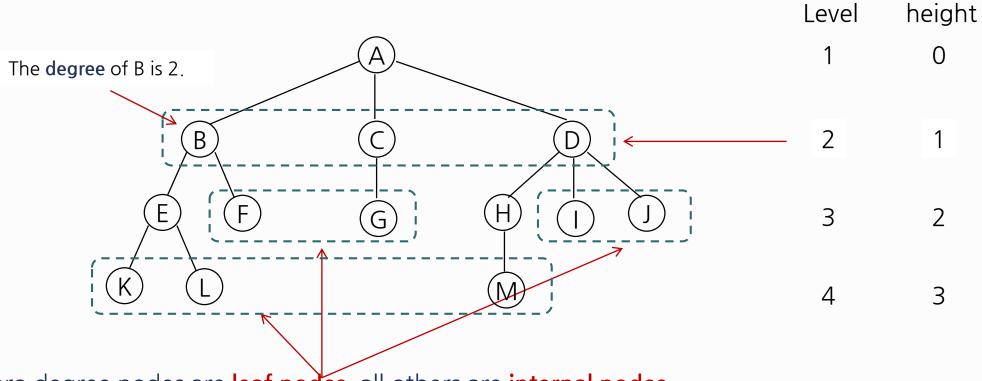
- A tree data structure: it is like a linked list that has a first node, this node is called as the root of the tree.
- Example. A tree with a root storing the value 'A'



- The children of D are H, I, and J; H, I, and J are siblings.
- The parent of D is A.

#### Introduction - Terminology

Definition. child, parent, sibling, degree, leaf nodes, level, and internal node



- Zero degree nodes are leaf nodes, all others are internal nodes.
  - An internal node is any node that has at least one non-empty child.
- The degree of a node is the number of children.
- The degree of a tree is the maximum of the degree of the nodes in the tree.

#### Introduction - Representation of trees

**Exercise.** The tree representing the HTML document below:

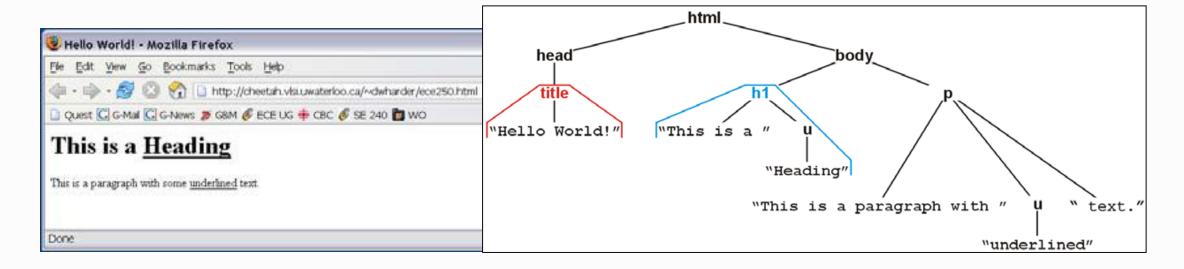
#### Introduction - Representation of trees

Exercise. The tree representing the HTML document below:



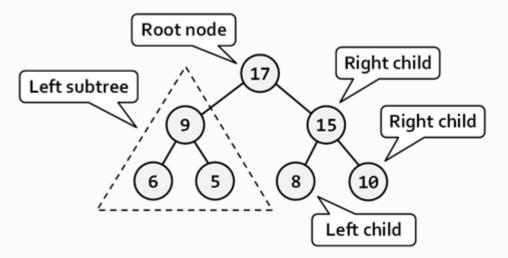
#### Introduction - Representation of trees

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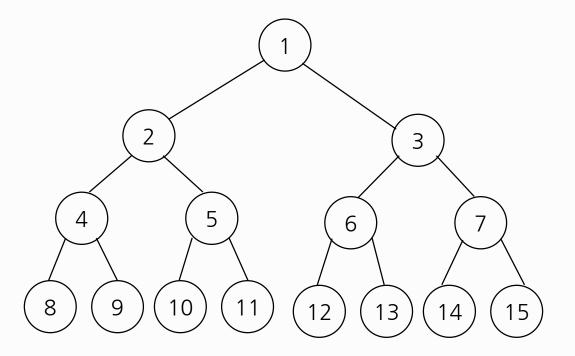


#### Binary trees

- Definition: A tree such that each node has exactly two children.
  - Notice, exactly two children not up to two children! Because exactly two children means a left child and/or right child, no middle child.
  - Each child is either empty or another binary tree.
  - Given this constraint, we can label the two children as left and right nodes or subtrees.

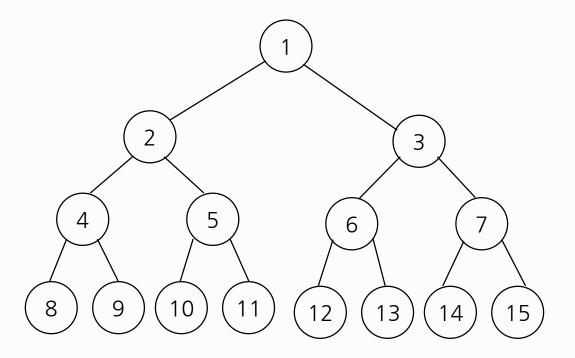


- Q: Maximum number of nodes in binary trees in each level and all levels?
- Q: What is the max level k if there are n nodes? k(n) = ?



A full binary tree

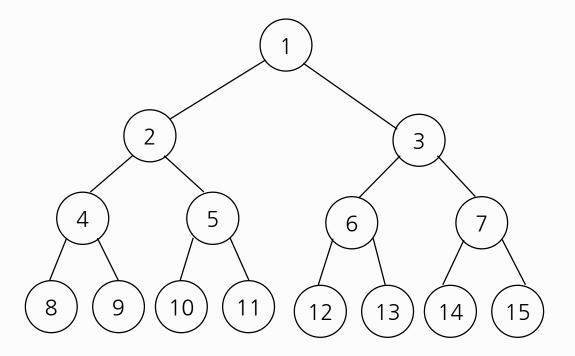
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A full binary tree

Level	Node Numbers at Each Level	Total Numbers of Nodes
1	$1 = 2^0$	
2	$2 = 2^1$	
3	$4 = 2^2$	
4	$8 = 2^3$	
•	·	
11	$1024 = 2^{10}$	
k		

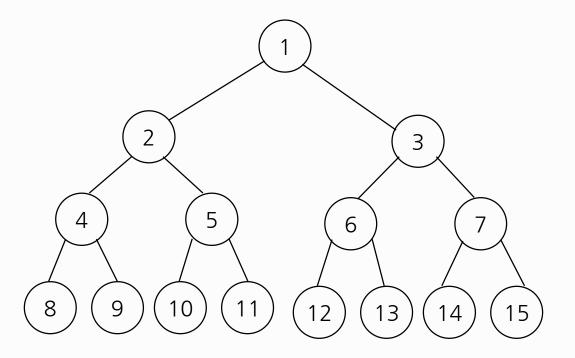
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A full binary tree

Level	Node Numbers at Each Level	Total Numbers of Nodes
1	$1 = 2^0$	$1 = 2^1 - 1$
2	$2 = 2^1$	$3 = 2^2 - 1$
3	$4 = 2^2$	$7 = 2^3 - 1$
4	$8 = 2^3$	$15 = 2^4 - 1$
•	·	
11	$1024 = 2^{10}$	$2047 = 2^{11} - 1$
•		
k		

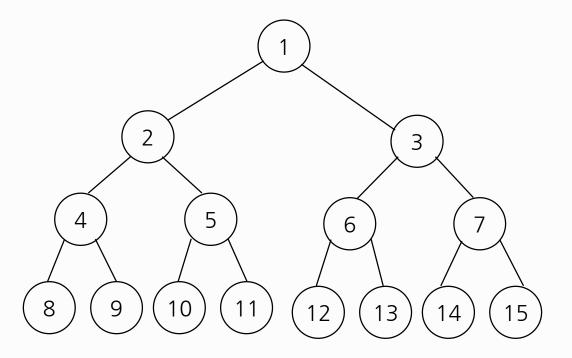
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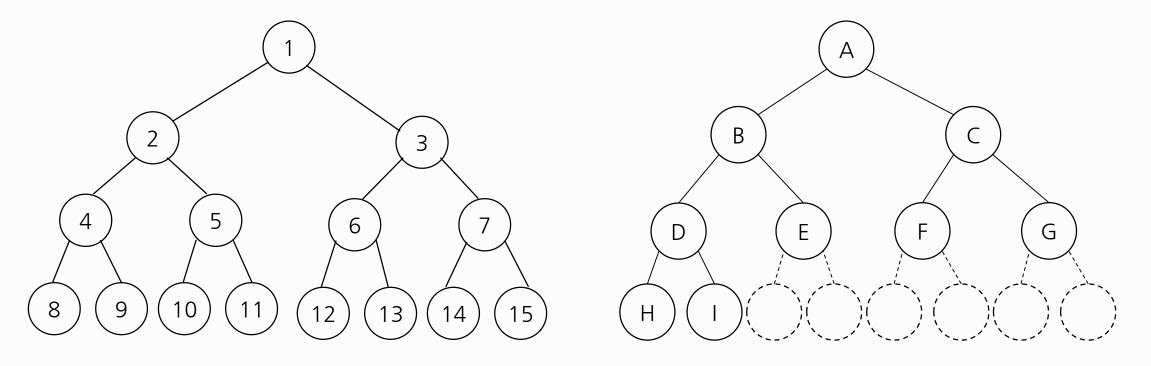
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4	$8 = 2^3$	$15 = 2^4 - 1$
11	$1024 = 2^{10}$	$2047 = 2^{11} - 1$
k	$2^{k-1}$	$2^k - 1$
h	2 <sup>h</sup>	2 <sup>h+1</sup> - 1

- **Definition:** A full binary tree of level k is a binary tree having  $2^k$  1 nodes,  $k \ge 0$ .
- **Definition**: A binary tree with n nodes and level k is **complete** iff its nodes correspond to the nodes numbered from 1 to n in the full binary tree of level k.



A full binary tree

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- **Definition**: A binary tree with n nodes and level k is **complete** iff its nodes correspond to the nodes numbered from 1 to n in the full binary tree of level k.

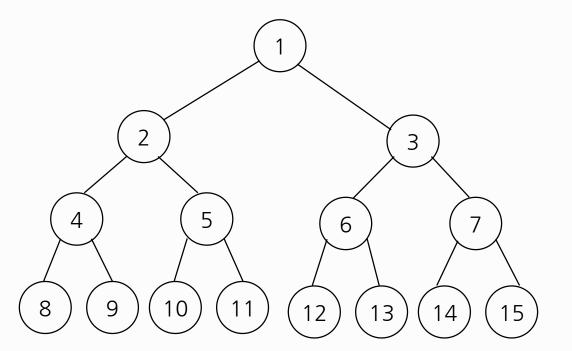


A full binary tree

A complete binary tree

#### Binary trees - Array representation

- Q: Let's suppose that you have a **complete binary tree** in an array, how can we locate node x's parent or child?
- A complete binary tree with n nodes, any node index i,  $1 \le i \le n$ , we have
  - parent(i) is at  $\lfloor i/2 \rfloor$  If i = 1, i is at the root and has no parent
  - leftChild(i) is at 2i if 2i <= n. If 2i > n, then i has no left child.
  - rightChild(i) is at 2i+1 if 2i+1 <= n. If 2i+1 > n, then i has no right child.



Wow! Can we use this to all binary trees? Why not?

#### Problem remains:

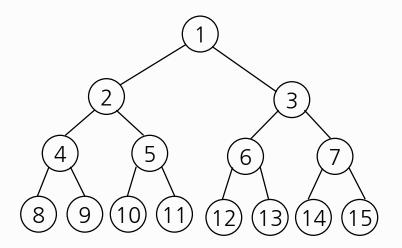
The problem with storing an arbitrary binary tree using an array is the inefficiency *in memory usage*.

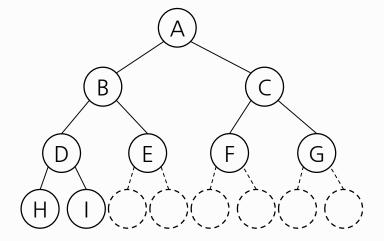
- (1) The maximum number of **nodes on level k** of a binary tree is  $k \ge 1$
- (2) The maximum number of nodes in a level k binary tree of is  $k \ge 1$
- (3) The maximum level of a **complete binary tree** with **n** nodes is [x] is the smallest integer  $\geq x$ .

- (1) The maximum number of **nodes on level k** of a binary tree is  $2^{k-1}$ ,  $k \ge 1$
- (2) The maximum number of nodes in a binary tree of level k is  $2^k 1$ ,  $k \ge 1$
- (3) The maximum level of a **complete binary tree** with **n** nodes is  $k(n) = [log_2(n+1)], [x]$  is the smallest integer  $\geq x$ .

$$n=2^k-1$$
  $n+1=2^k$   $\log(n+1)=\log 2^k$   $\log(n+1)=k$   $\log(n+1)=\log(n+1)$  since k is an integer, and includes  $\log(n)=\log(n)+1$  the max level of complete binary tree.

- Observation: The max level of a full binary tree of n nodes is k = floor(log(n)) + 1:
  - Many operations with trees have a run time that goes with the max level of some path within the tree;
  - If we have a full binary tree (or something *close* to it), we know that those operations run in  $O(\log n)$ .





A full binary tree

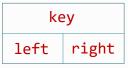
A complete binary tree

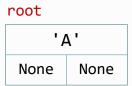
#### Node and Tree representations:

```
class Node:
    def __init__(self, key):
        self.key = key
        self.left = None
        self.right = None

    def insertLeft(self, key):
    ...
```

```
root = Node('A')
print(root)
```





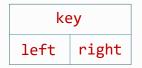
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    def insertLeft(self, key):
    ...
```

```
root = Node('A')
print(root)
```

```
<__main__.Node object at
0x0000015985006A00>
```



root
'A'
None None

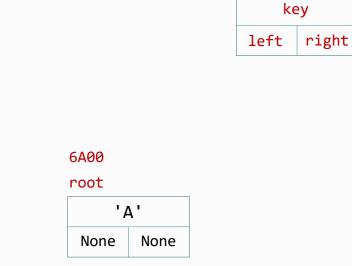
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```

Create 'B' and link with left of 'A'

```
root = Node('A')
node = Node('B')
```



7A00

node

None

'B'

None

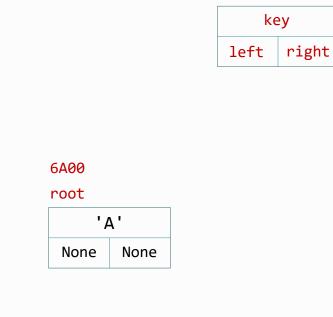
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Create 'B' and link with left of 'A'.

```
root = Node('A')
node = Node('B')
root.left = node
```



7A00

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None

'B'

None

Node and Tree representations:

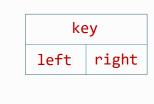
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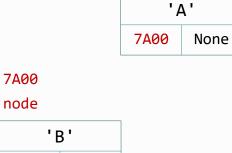
    def insertLeft(self, key):
    ...
```

Create 'B' and link with left of 'A'.

```
root = Node('A')
node = Node('B')
root.left = node
```

Simplify the code and diagram.





None

None

6A00 root

Node and Tree representations:

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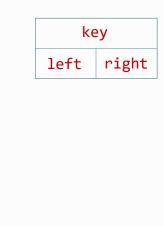
def insertLeft(self, key):
    ...
```

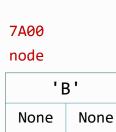
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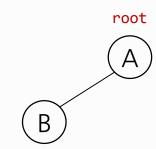
```
root = Node('A')
node = Node('B')
root.left = node
```

Simplify the code and diagram.

```
root = Node('A')
root.left = Node('B')
```







6A00 root

7A00

'Α'

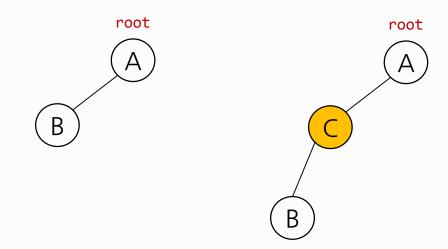
None

Node and Tree representations:

```
class Node:
    def __init__(self, key):
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        self.left = None
        self.right = None

def insertLeft(self, key):
    ...
```

```
root = Node('A')
root.insertLeft('B')
root.insertLeft('C')
```



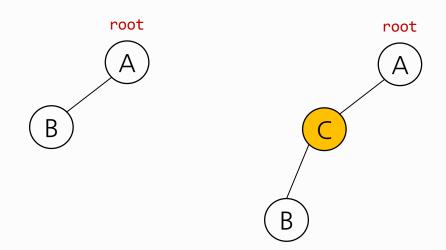
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```
root = Node('A')
root.insertLeft('B')
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```

```
def insertLeft(self, key):
    if self.left == None:
        self.left = Node(key)
    else:
        ...
```



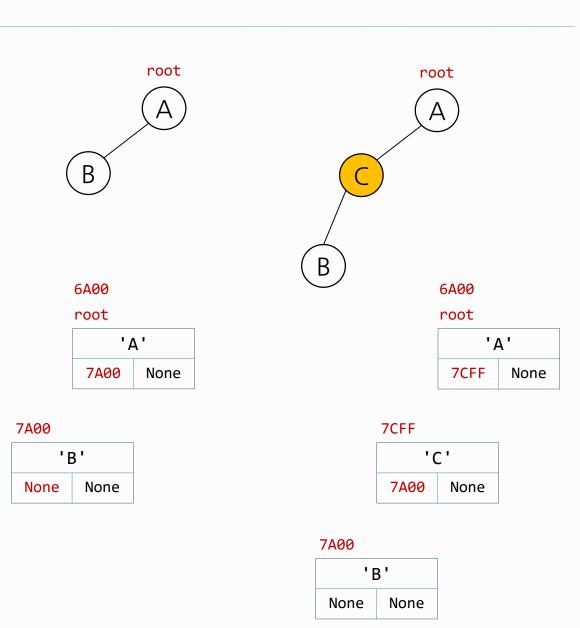
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```



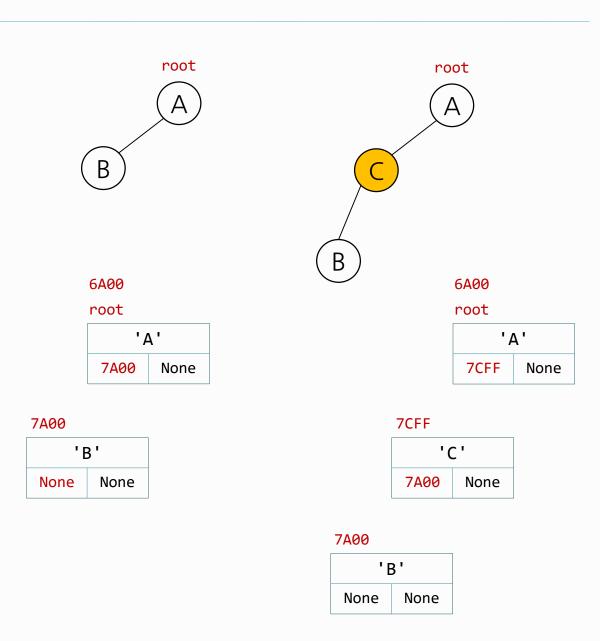
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    def insertLeft(self, key):
    ...
```

```
root = Node('A')
root.insertLeft('B')
root.insertLeft('C')
```

```
def insertLeft(self, key):
    if self.left == None:
        self.left = Node(key)
    else:
        node = Node(key)
        node.left = self.left
        self.left = node
```



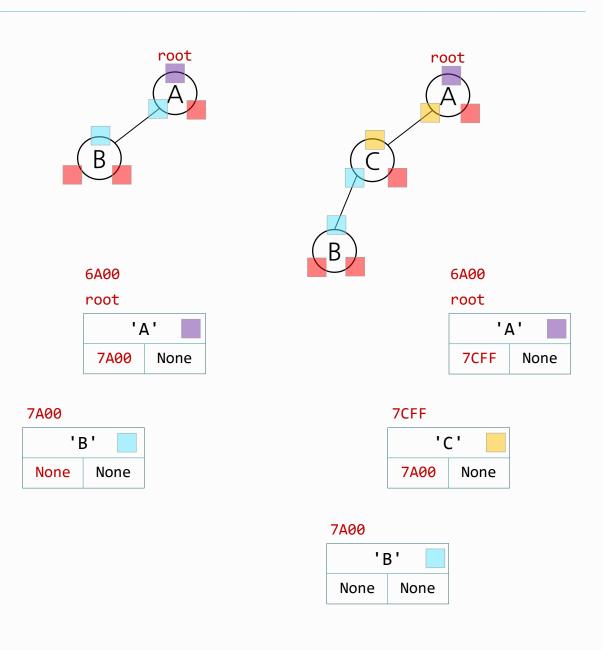
Node and Tree representations:

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    def insertLeft(self, key):
    ...
```

```
root = Node('A')
root.insertLeft('B')
root.insertLeft('C')

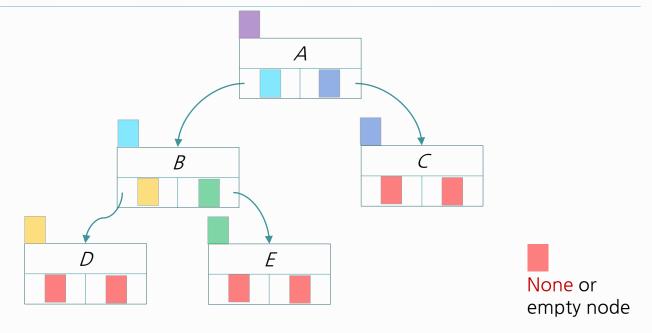
def insertLeft(self, key):
    if self.left == None:
        self.left = Node(key)
    else:
        node = Node(key)
        node.left = self.left
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```



Node and Tree representations:

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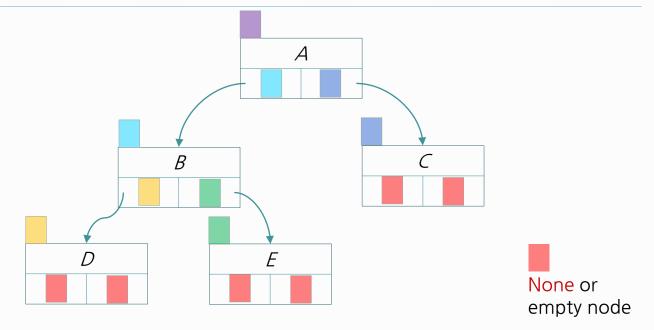


- Q. Is this node structure good enough?
  - Not easy to find its parent node. A parent field could be added if necessary.

Node and Tree representations:

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class Node:
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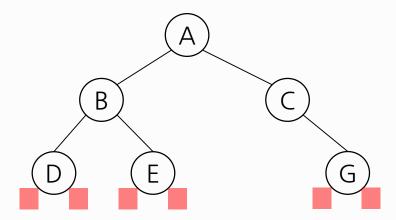


- Q. Is this node structure good enough?
  - Not easy to find its parent node. A parent field could be added if necessary.
- Q. It is similar to a doubly-linked list(DLL). What is different?
  - One head, but many tails. None points empty node conceptually.

#### **Binary trees - Exercise 1**

Build a tree shown below using insertLeft() and insertRight().

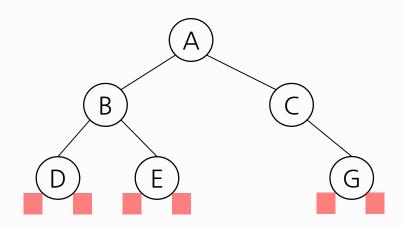
```
class Node:
    def __init__(self, key):
        self.key = key
        self.left = None
        self.right = None
    def insertLeft(self, key):
    def insertRight(self, key):
        . . .
if __name__ == '__main__':
    root = Node('A')
    root.insertLeft('B')
    root.insertRight('C')
   # your code here
```



#### Binary trees - Exercise 2

Extend Exercise 1 such that it exactly reproduces the output using the as shown below.

```
class Node:
    def __init__(self, key):
        self.key = key
        self.left = None
        self.right = None
    def insertLeft(self, key):
    def insertRight(self, key):
        . . .
if __name__ == '__main__':
    for node in [a, b, c, d, e, g]
        if node.key:
            print(...
        . . .
```



```
A:(B, C)
B:(D, E)
C:(None, G)
D:(None, None)
E:(None, None)
G:(None, None)
```

## 학습 정리

- 1) 트리(Tree)는 비선형 자료 구조이다
- 2) 완전이진트리(Complete binary tree)에 대한 시간복잡도는  $O(\log n)$ 으로 빠른 알고리즘이다

