

학습 목표

해시(Hash) 테이블과 해시 함수를 이해하고 특징들을 학습한다



Data Structures in Python Chapter 6

- Hash Table
- Collision Resolution
- Double Hashing & Rehashing
- HashMap Coding

Agenda & Readings

- Agenda
 - Hashing
 - Hash Table
 - Hash Function
- Reference:
 - Problem Solving with Algorithms and Data Structures
 - Chapter 5 Hashing

Overview

Hashing or Hash Table Data Structure:

Data structures so far

Array of size n	unsorted list	sorted array	Trees BST – average AVL – worst	Heap, Priority Queue	Hashing
insert	find+0(1)	0(n)	O(log n)	O(log n)	
find	0(n)	O(log n)	O(log n)	O(log n)	
remove	find+0(1)	0(n)	O(log n)	O(log n)	

Overview

 Hashing or Hash Table Data Structure: supports insertion, deletion and search in average case constant time O(1).

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find	0(n)	O(log n)	O(log n)	O(log n)	O(1)
remove	find+0(1)	0(n)	O(log n)	O(log n)	O(1)

Overview

 Hashing or Hash Table Data Structure: supports insertion, deletion and search in average case constant time O(1).

Hash table

- It is data structure that stores key-value pairs.
- The key is sent to a hash function that performs arithmetic operations on it.
- The result is called hash value that is the index of the key-value pair in the hash table.

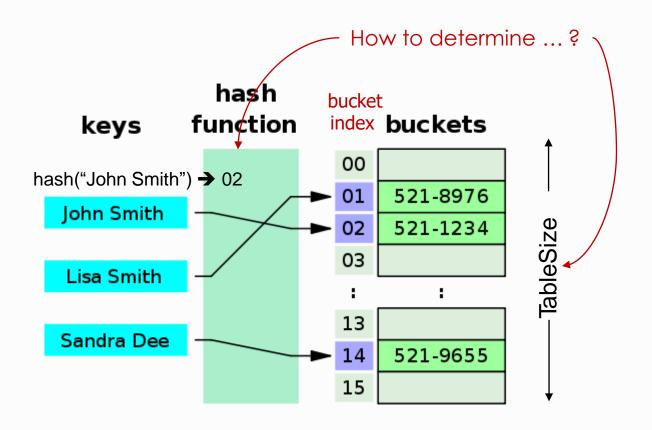
Hash function

- hash(key) → integer value
- hash("string key") → integer value

Hash table is an array of fixed size elements

Let us suppose that there are one billion of names and numbers. Find, insert, and remove a number by a given name in **O(1)**. Time Complexity keys buckets hash("John Smith") 521-8976 John Smith **TableSize** 521-1234 Lisa Smith Sandra Dee 521-9655

- Hash table is an array of fixed size elements
- Array elements indexed by a key mapped to a bucket index[0 .. TableSize-1]
- Mapping from key to index using hash(), hash function
 - e.g., hash("John Smith") → 02



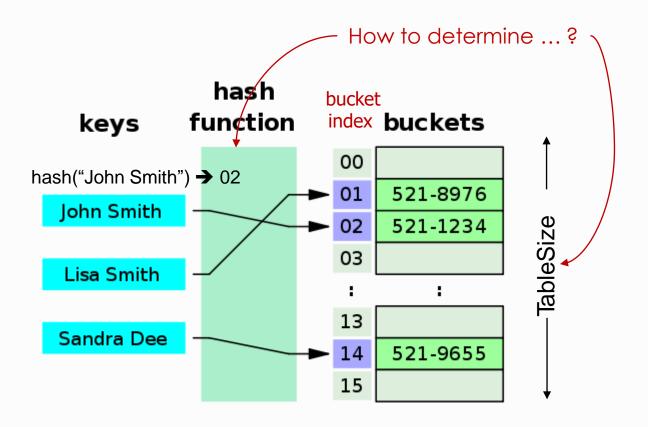
- insert
 - HashTable[hash("John Smith")] = <"John Smith", 521-1234>
- remove
 - HashTable[hash("John Smith")] = None
- find
 - HashTable[hash("John Smith")] returns the element hashed for "John Smith"

hash bucket fun¢tion index buckets keys 00 hash("John Smith") → 02 521-8976 01 John Smith 02 521-1234 **TableSize** 03 Lisa Smith Sandra Dee 14 521-9655 15

How to determine ...?

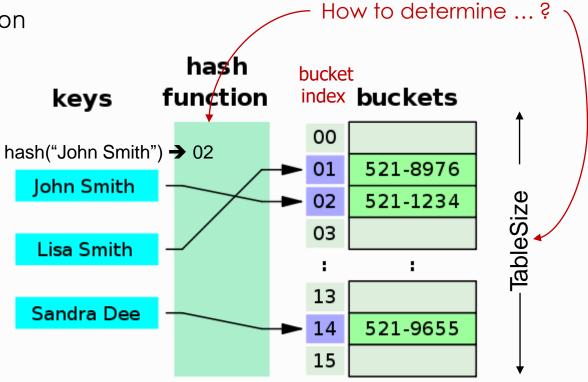
What happens if hash("John Smith") == hash("Joe Blow")? "Collision"

- Factors affecting Hash Table Design
 - Hash function
 - Table size Usually fixed at the start
 - Collision handling schemes Array or Linked List



Hash Function

- It maps an element's key into a valid hash table index
 - hash(key) → hash table index
- Note that this is (slightly) different from saying:
 - hash(string) → int
 - Because the key can be of any type
 - e.g., "hash(int) → int" is also a hash function



Hash Function Properties

- It maps an element's key into a valid hash table index
 - hash(key) → hash table index
- It maps key to integer
 - Constraint: Integer should be between [0, TableSize-1]
- A hash function can result in a many-to-one mapping (causing collision)
 - Collision occurs when hash function maps two or more keys to same array index
- Collisions cannot be avoided but its chances can be reduced using a "good" hash function

Hash Function - Effective use of table size

- Simple hash function (assume integer keys)
 - hash(Key) = Key % TableSize
- For random keys, hash() distributes keys evenly over table
 - What if TableSize = 100 and keys are ALL multiples of 10?
 - Better if TableSize is a prime number

Hash Function Example: String Keys

- Using a very simple function to map strings to integers:
 - Add up character ASCII values (0-255) to produce integer keys
 - e.g., "abcd" = 97 + 98 + 99 + 100 = 394
 - hash("abcd") = 394 % TableSize
- Potential problems:
 - Anagrams will map to the same index
 - hash("abcd") = hash("dbac")
 - Small strings may not use all of table
 - strlen(s) * 255 < TableSize
 - Time proportional to length of the string

Hash Function Example: String Keys

- Another approach:
 - Treat first 3 characters of string as base-27 integer(26 letters plus space)
 - e.g., Key = $s[0] + (27^1 * s[1]) + (27^2 * s[2])$
 - Better than previous approach because ...
- But, potential problems:
- Apple Apply collisionAppointment Apricot

Hash Function Example: String Keys

- Last approach:
 - Use all N characters of string as an N-digit and base-K number
 - Choose K to be prime number larger than number of different digits (characters)
 - i.e., K = 29, 31, 37
 - If L = Length of string S, then

$$hash(S) = \sum_{i=0}^{L-i} S[L-1-i] * 37^{i} \% TableSize$$
 (1)

- Use Horner's rule to compute hash(S).
- Limit L for long strings
- Potential problems
 - Overflow
 - Larger runtime

```
# a hash function for strings
def hash(key, tablesize)
  code = 0
  for x in key: code = code * 37 + x
  code %= tablesize
  if code < 0: code += tablesize
  return code</pre>
```

Summary

- Using a hash table we can, on average (if table large enough and hash function suitable), insert, delete and search for items in constant time - O(1).
- The hash function is the mapping between an item and the slot where the item is stored.
- A collision occurs when an item is mapped to an occupied slot.
- A perfect hash function is able to map m items into a table of size m with no collisions. Perfect hash functions are hard to come by.
- Handling collisions systematically is required collision resolution.

학습 정리

1) 해시 테이블을 이용하면, 삽입/삭제/검색 작업을 O(1)으로 해결할 수 있다

2) 해시 함수가 반환하는 해시 값(Value)은 항상 정수이다

3) 해시 테이블을 구현할 때, 충돌을 최소화할 수 있는 해시 함수를 사용해야 한다

