

# 파이썬으로 배우는 데이터 구조



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# 학습 목표

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힙(Heap)과 BST와의 차이점을 이해하고

힙(Heap)의 시간복잡도를 알 수 있다

## **Data Structures in Python**

### **Chapter 8**

- **Heap and Priority Queue**
- **Heap Coding**
- **Min/MaxHeap and Heap sort**

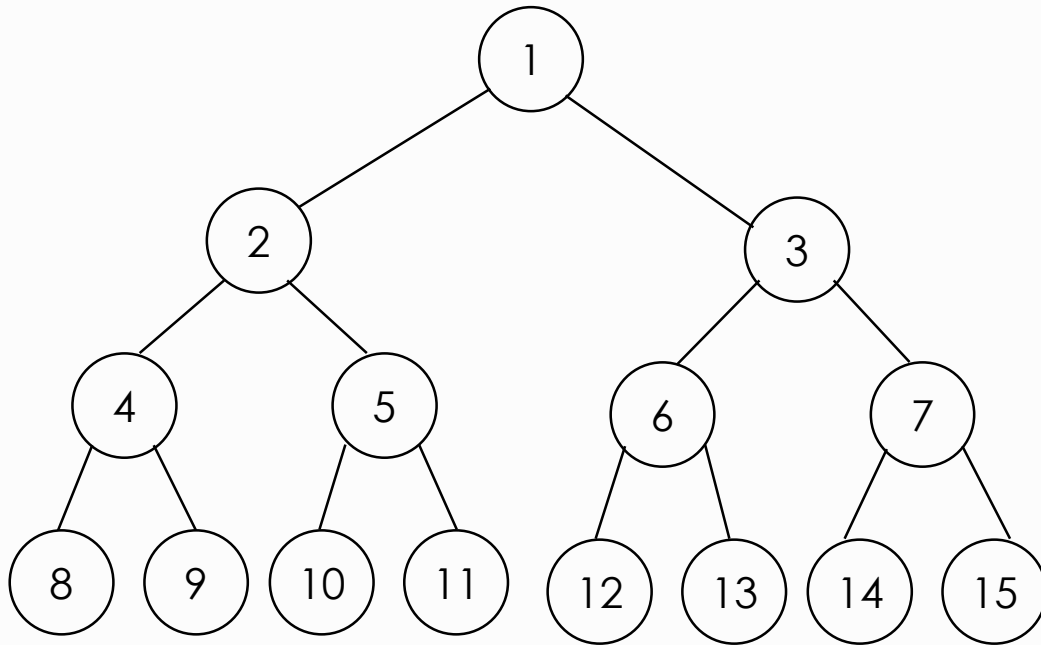
# Agenda & Readings

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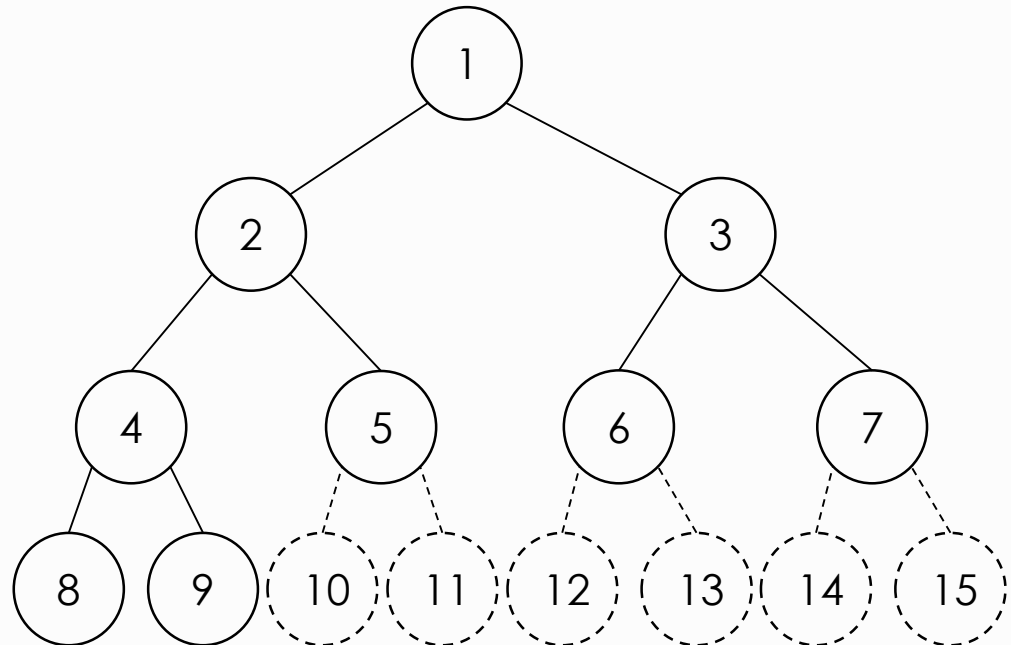
- Heap and Priority Queue
  - Complete Binary Tree (Review)
  - Heap and Priority Queue
  - Heap ADT
  - Time Complexity
- Reference:
  - Problem Solving with Algorithms and Data Structures

# Binary trees - Properties

- **Definition:** A **full** binary tree of level  $k$  is a binary tree having  $2^k - 1$  nodes,  $k \geq 0$ .
- **Definition:** A binary tree with  $n$  nodes and level  $k$  is **complete** if and only if its nodes correspond to the nodes numbered from 1 to  $n$  in the full binary tree of level  $k$ .



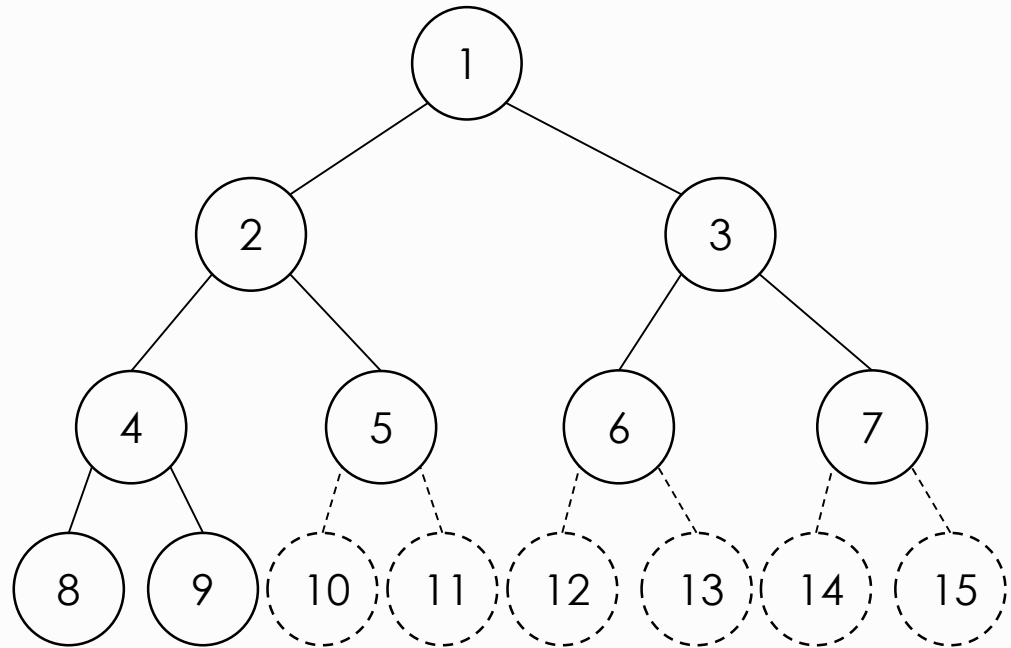
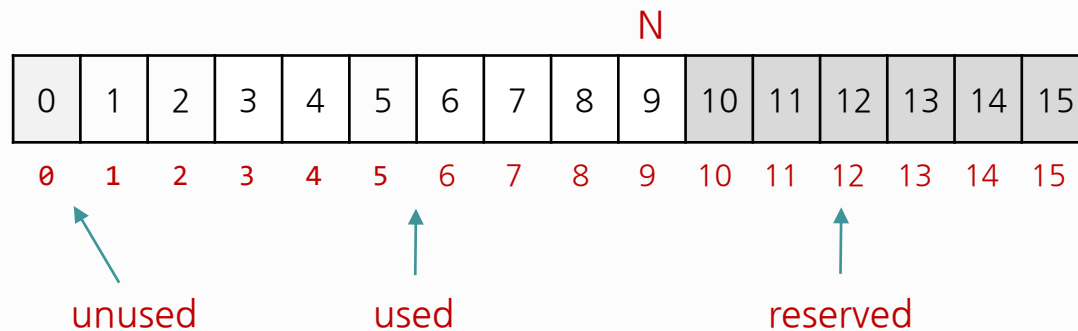
*A **full** binary tree*



*A **complete** binary tree*

# Binary trees - Array representation

- A **complete** binary tree with  $n$  nodes, any node index  $i$ ,  $1 \leq i \leq n$ , we have
  - $\text{parent}(i)$  is at  $\lfloor i/2 \rfloor$  If  $i = 1$ ,  $i$  is at the root and has no parent
  - $\text{leftChild}(i)$  is at  $2i$  if  $2i \leq n$ . If  $2i > n$ , then  $i$  has no left child.
  - $\text{rightChild}(i)$  is at  $2i+1$  if  $2i+1 \leq n$ . If  $2i+1 > n$ , then  $i$  has no right child.



*A **complete** binary tree*

# Heaps & Priority Queues

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  - Because it provides an efficient implementation for **priority queues**.

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- **Priority queues**.
  - Queues with priorities associated to.
  - **Example:** A line waiting to be served at a bank and served FIFO except if a senior or a disabled person arrives in the line. They are served first. Seniors and disabled persons have higher priority than others.



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  - **Example:** A line waiting to be served at a bank and served FIFO except if a senior or a disabled person arrives in the line. They are served first. Seniors and disabled persons have higher priority than others.
- **A typical ADT for Priority Queue**
  - Get the top priority element (min or max)
  - Insert an element
  - Delete the top priority element
  - Decrease the priority of an element

- $O(1)$
- $O(\log n)$
- $O(\log n)$
- $O(\log n)$

# Heaps & Priority Queues

- Challenge: Find the largest M items in a stream of N items.
  - Constraints: Not enough memory to store N items.
- N huge, M large*

Order of insert of finding the largest M **in a stream of N items**

implementation	insert	delete	min/max
unordered array	1	N	N
ordered array	N	1	1
goal	<b>log N</b>	<b>log N</b>	1 or log N

**Mission Impossible?**

min/max

priority queue

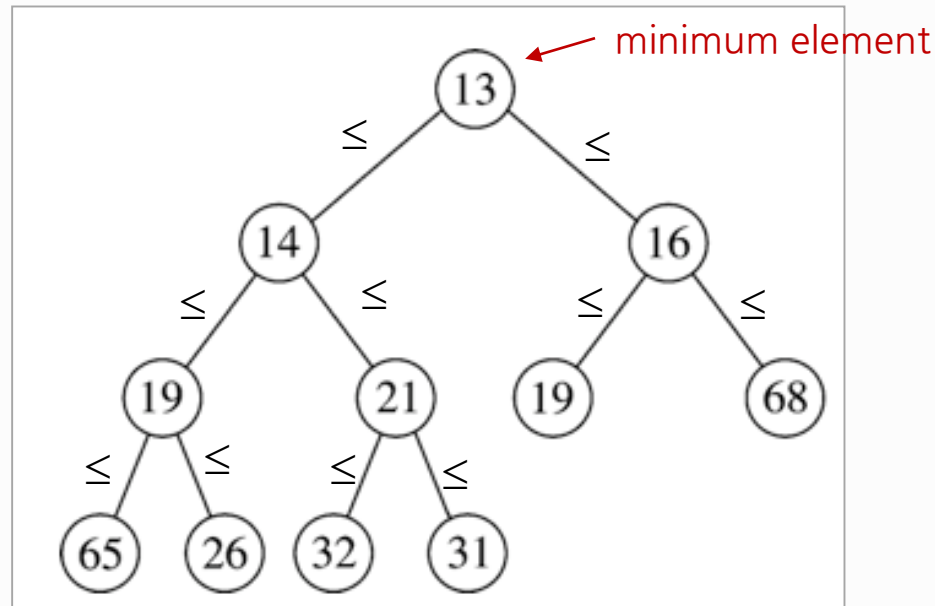
# Binary heap

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- **Binary heap**: array representation of a **heap-ordered** complete binary tree
- **Properties:**
  - **Heap-ordered:**  
Parent's key **no smaller** than children's keys for max-heap. (**no greater** for min-heap)
  - **Heap-structure:**  
A complete binary tree

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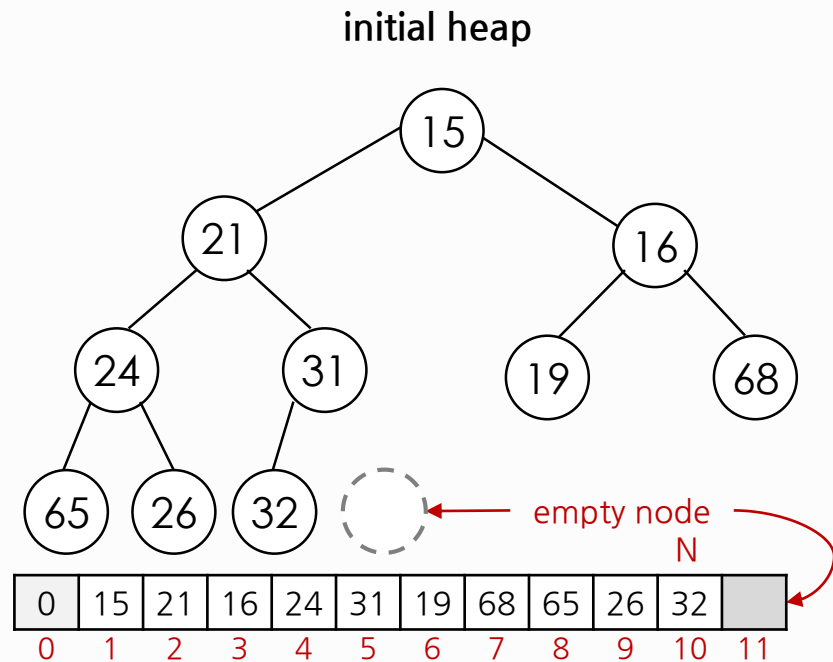


- Duplicates are allowed
- No order implied for elements which do not share ancestor-descendant relationship

## min-heap: insert(heap, 14)

Algorithm:

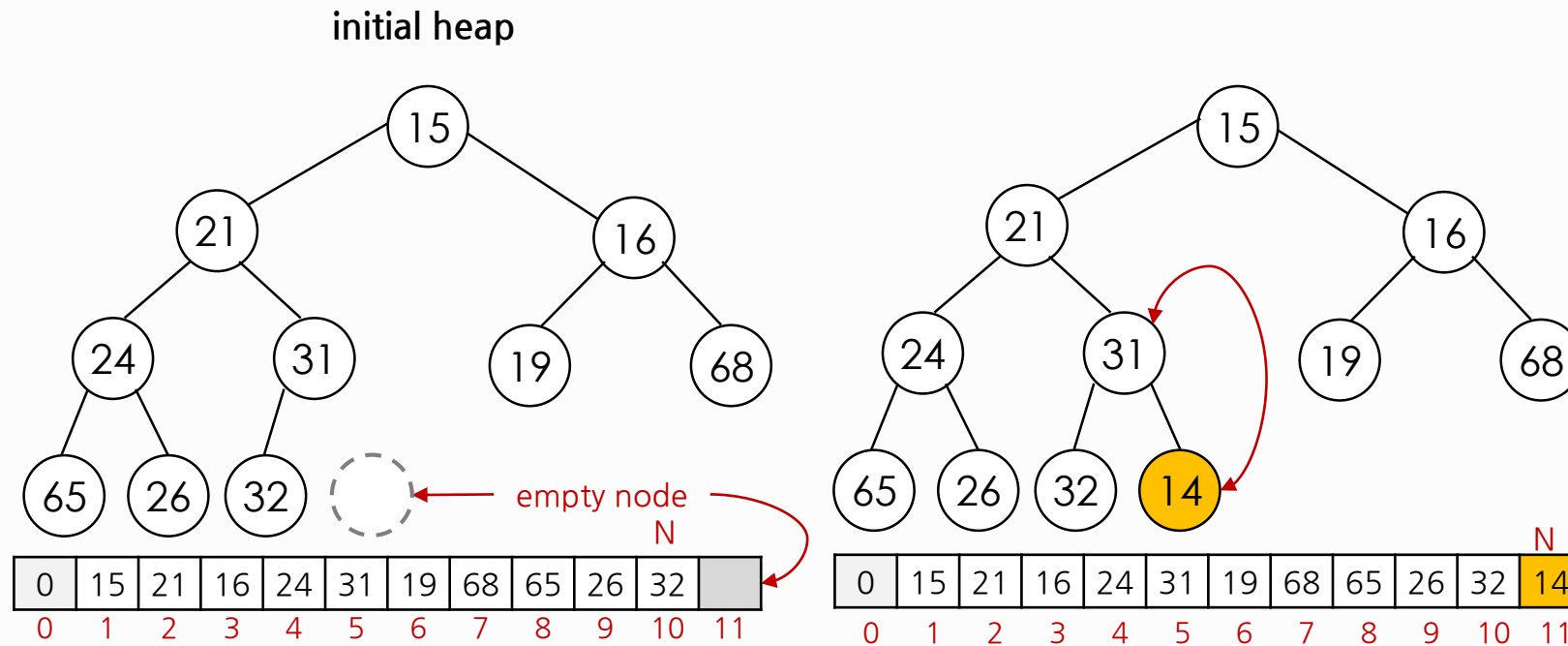
- Insert a new element **while maintaining a heap-structure**
- Move the element up the heap **while not satisfying heap-ordered**



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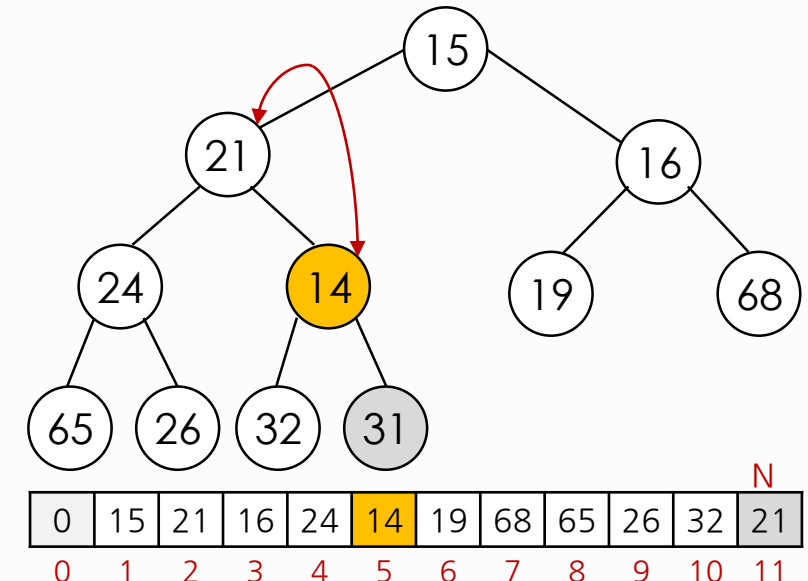
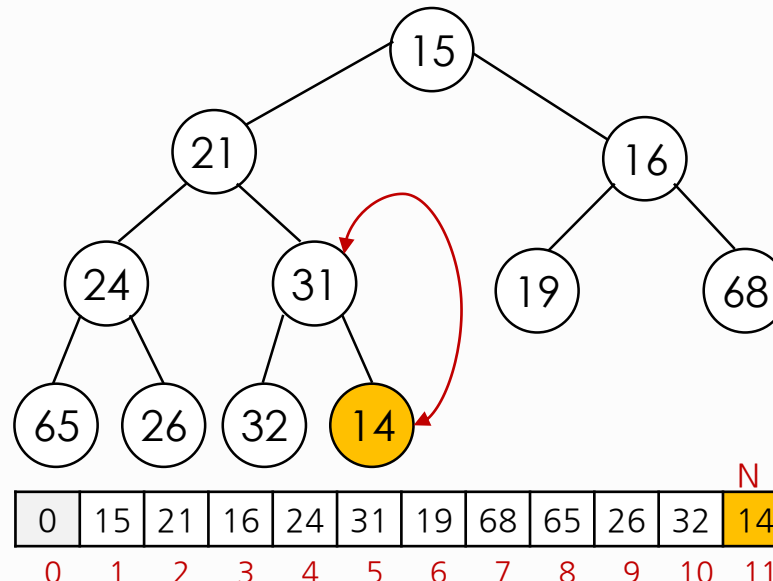
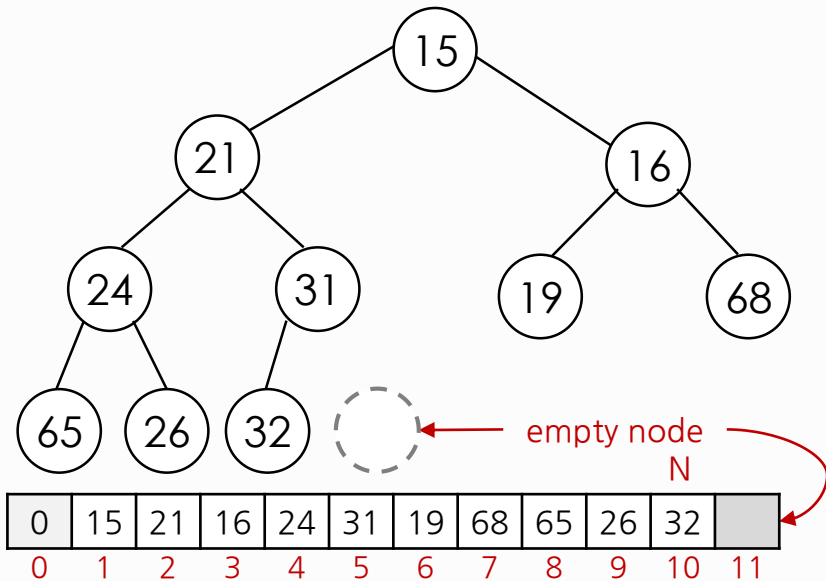


# min-heap: insert(heap, 14)

Algorithm:

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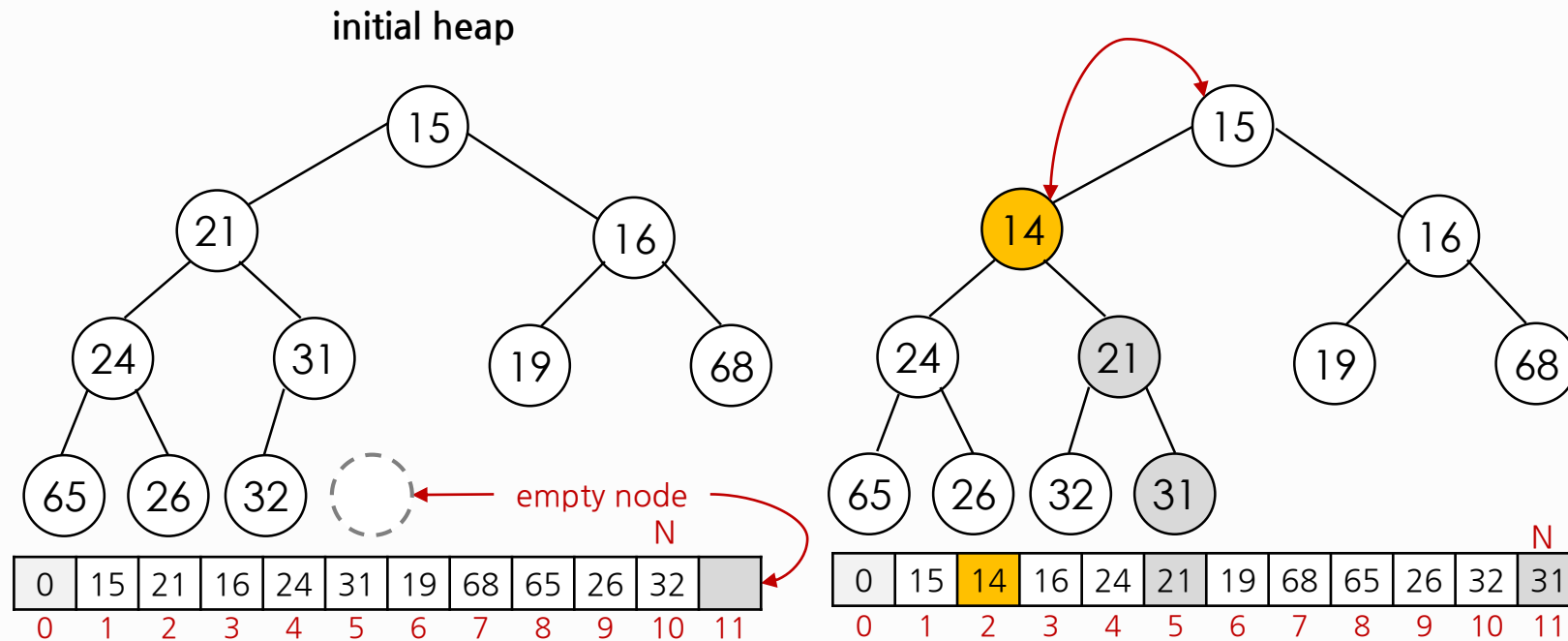
initial heap



## min-heap: insert(heap, 14)

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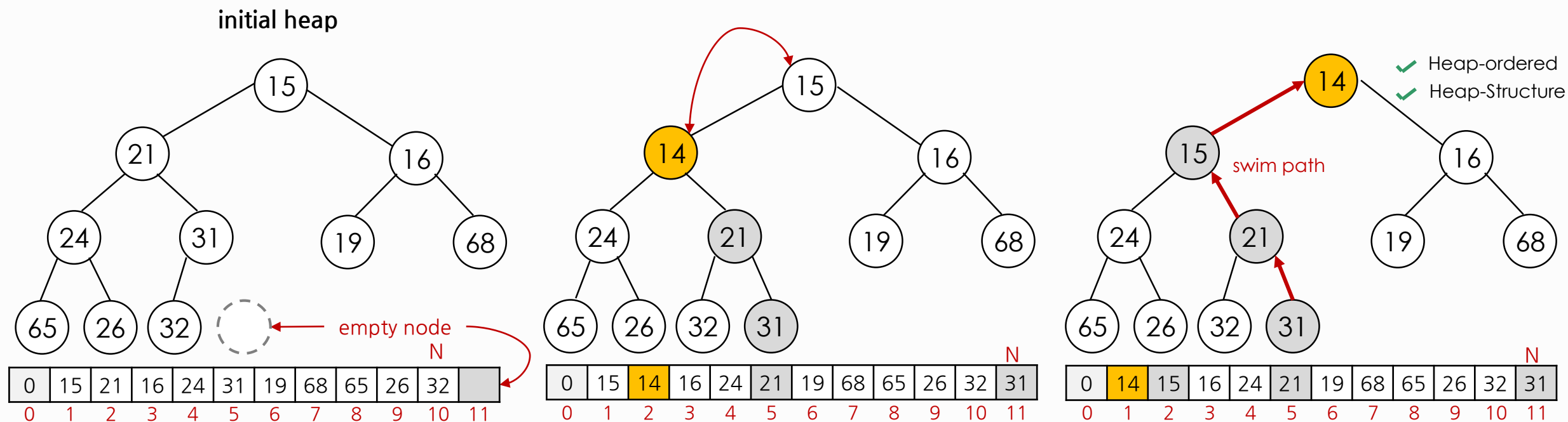




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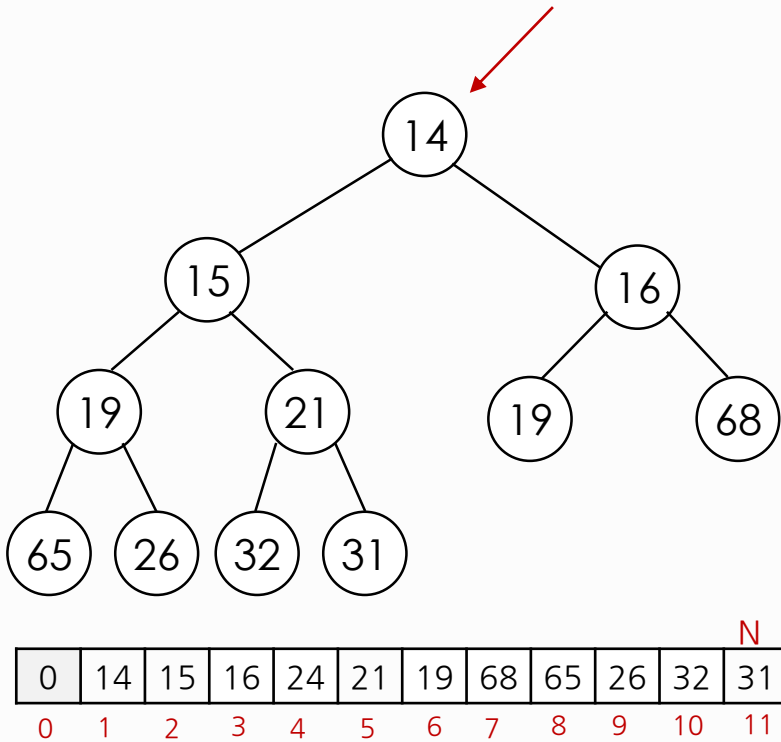
## min-heap: delete() or dequeue()

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- Swap the root and the last element.
- Heap decreases by one in size.
- **Move down (sink) the root** while not satisfying heap-ordered.
  - Minimum element is **always** at the root (by min-heap definition).

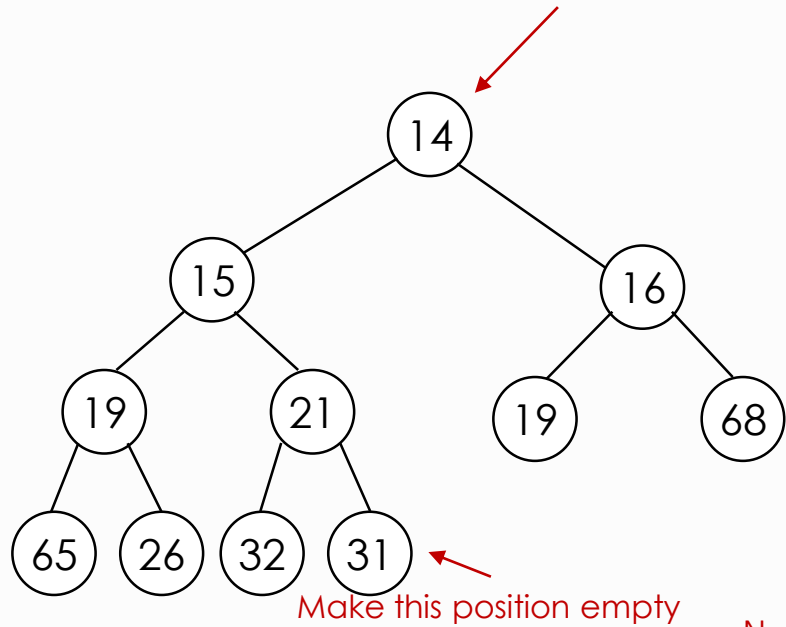
## min-heap: delete() or dequeue()

- Which position of the node will be empty?



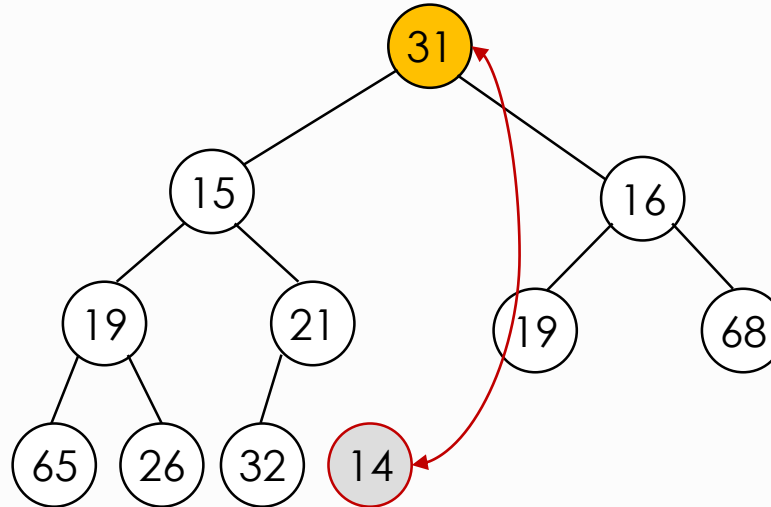
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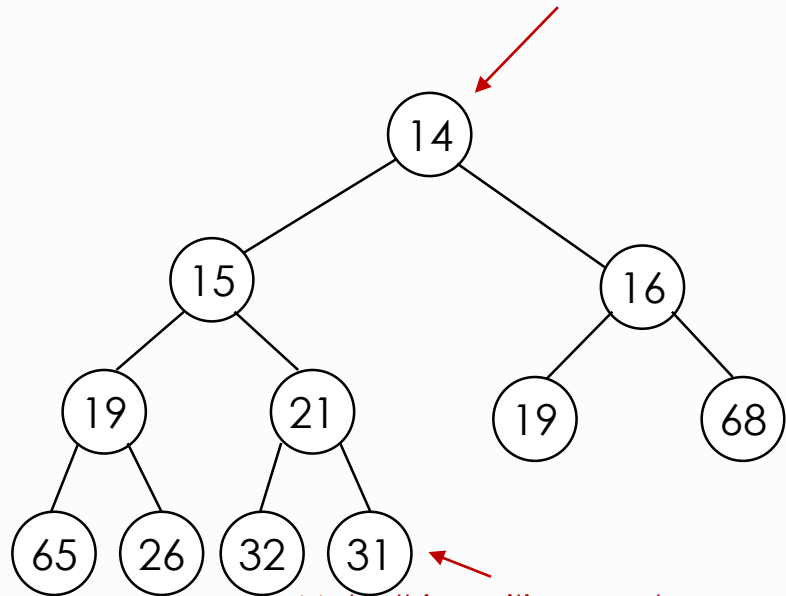
0	14	15	16	24	21	19	68	65	26	32	31
0	1	2	3	4	5	6	7	8	9	10	11

- swap the root & last
- **N--**



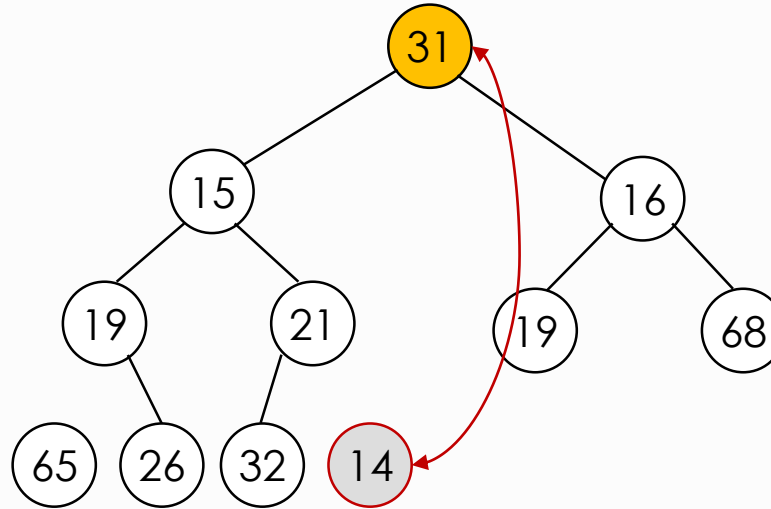
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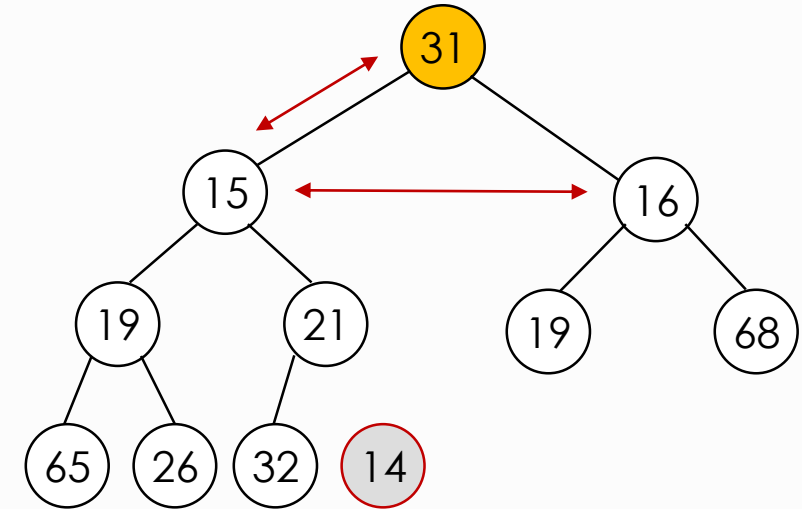
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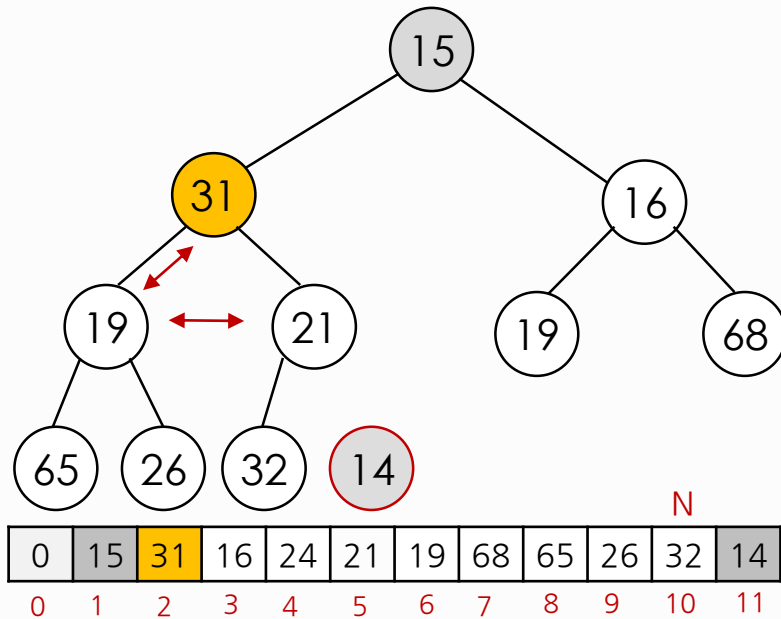
0	31	15	16	24	21	19	68	65	26	32	14
0	1	2	3	4	5	6	7	8	9	10	11

- heap-ordered?
- sink(): select one of two children and compare



# min-heap: delete() or dequeue()

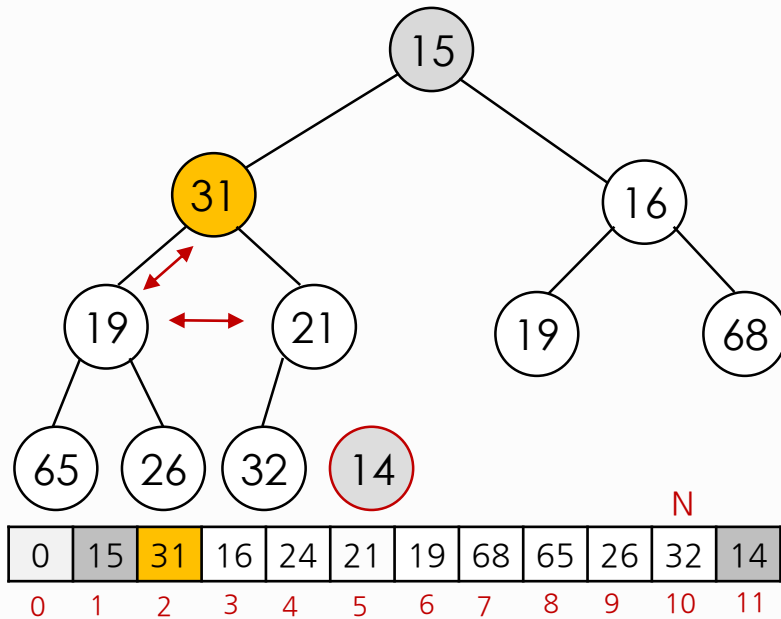
- heap-ordered?
- **sink()**: select one of two children and compare



- Is  $31 > \min(14, 16)$ ?
- Yes - swap 31 with  $\min(14, 16)$

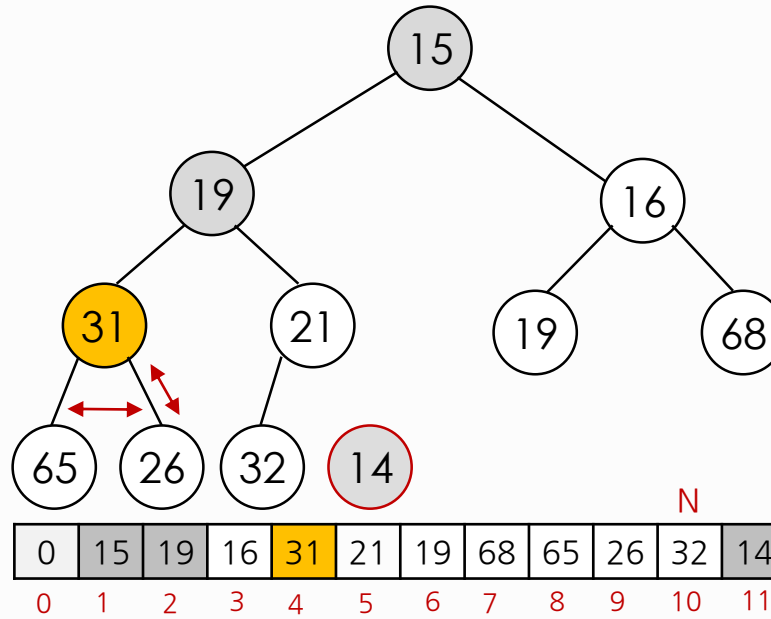
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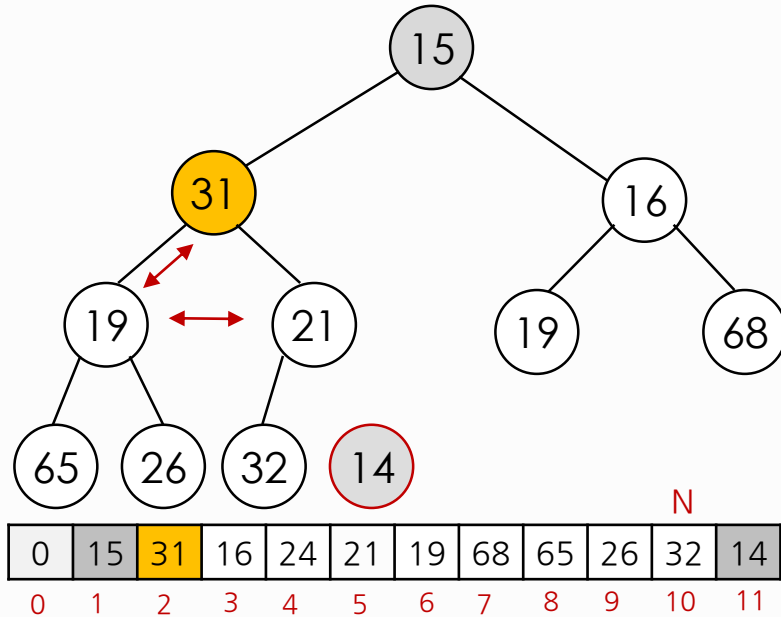
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- Is  $31 > \min(19, 21)$ ?
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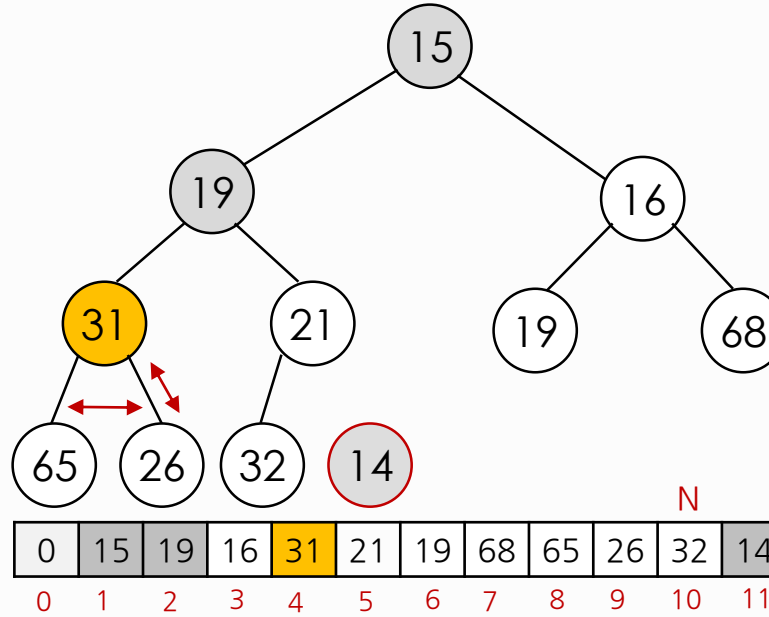
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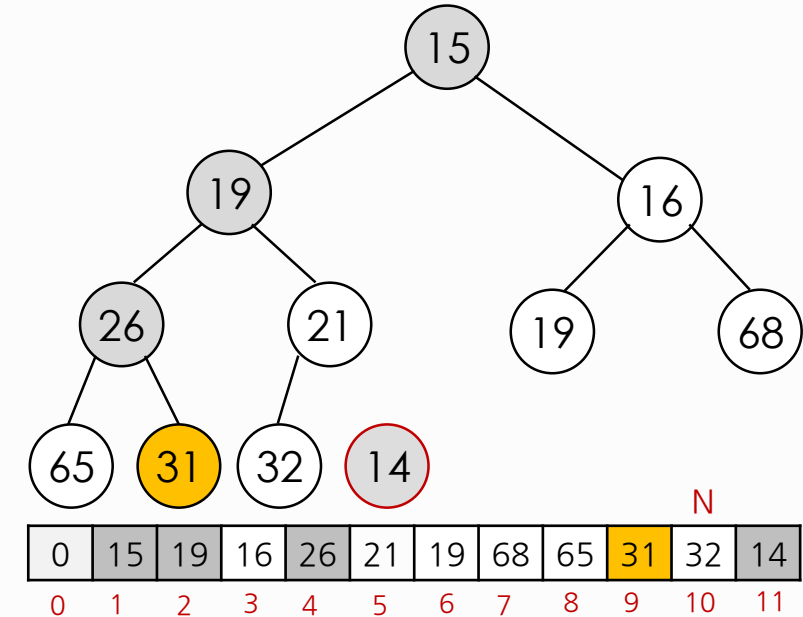
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- heap-ordered?
- **sink()**: select one of two children and compare



- Is  $31 > \min(19, 21)$ ?
- Yes - swap 31 with  $\min(19, 21)$

- heap-ordered?
- **sink()**: select one of two children and compare



- Is  $31 > \min(65, 26)$ ?
- Yes - swap 31 with  $\min(65, 26)$

✓ Heap-ordered  
✓ Heap-Structure



## Binary heap: Time complexity:

- Level of heap is  $\lfloor \log_2 N \rfloor$
- insert:  $O(\log N)$  for each insert
  - In practice, expect less
- delete:  $O(\log N)$  // deleting root node or any node
- increase/decrease key:  $O(\log N)$

Implementation	Insert	Delete	min/max
Unordered array	1	N	N
Ordered array	N	1	1
Binary heap	<b>log N</b>	<b>log N</b>	1 or log N

**Mission Completed**

min/max

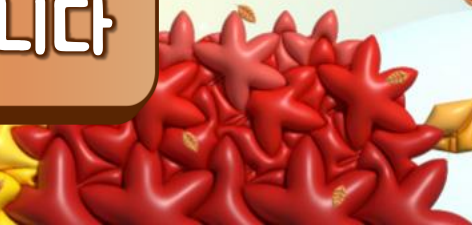
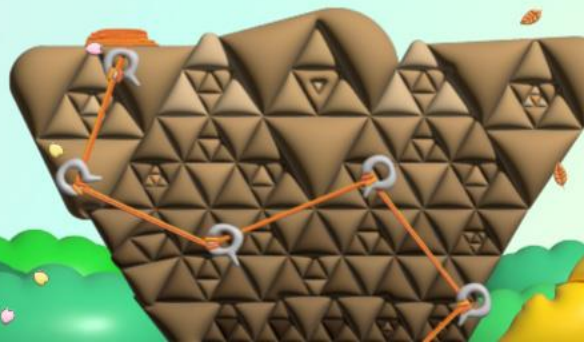
priority queue

# 학습 정리

- 1) 힙(Heap)은 우선순위 큐(Priority queue)에 주로 사용된다
- 2) 힙(Heap)에서의 삽입과 삭제 작업은 최대  $O(\log n)$  시간복잡도로 수행할 수 있다

# 파이썬으로 배우는 데이터 구조

수고했습니다  
곧 다음 시간에  
다시 뵙겠습니다



## Data Structures in Python

- Heap and Priority Queue
- **Heap Coding**
- Min/MaxHeap and Heap sort

### Proof:

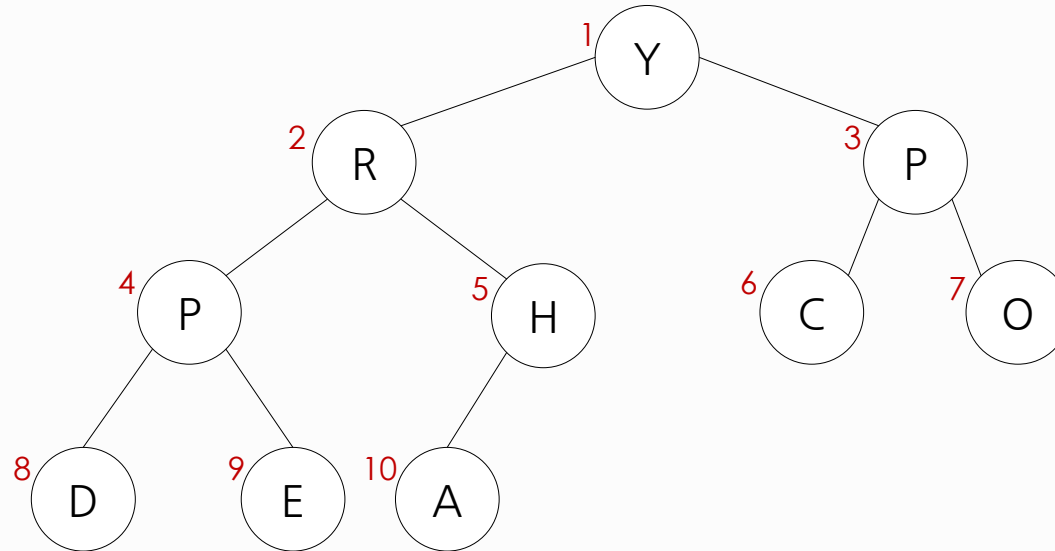
<https://stackoverflow.com/questions/9755721/how-can-building-a-heap-be-on-time-complexity>  
<https://www.insertingwiththeweb.com/data-structures/binary-heap/build-heap-proof/>  
<https://www.quora.com/How-is-the-time-complexity-of-building-a-heap-is-o-n>

### References in Korean:

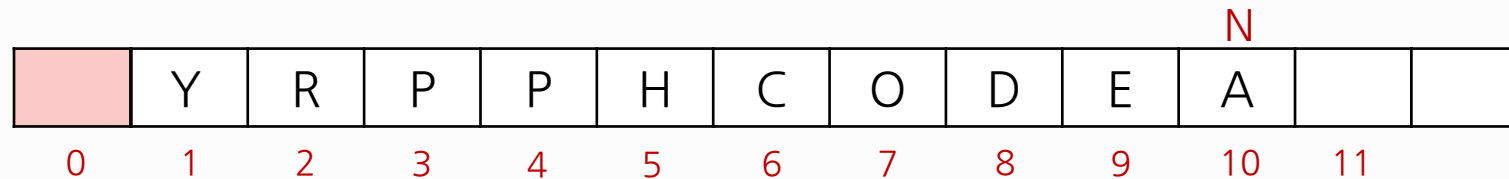
<https://ratsgo.github.io/data%20structure&algorithm/2017/09/27/heapsort/>  
<https://zeddios.tistory.com/56>

## max-heap example

- **Insert:** Add node at end, then swim it up.
- **Remove:** Swap root with last node, then sink down.

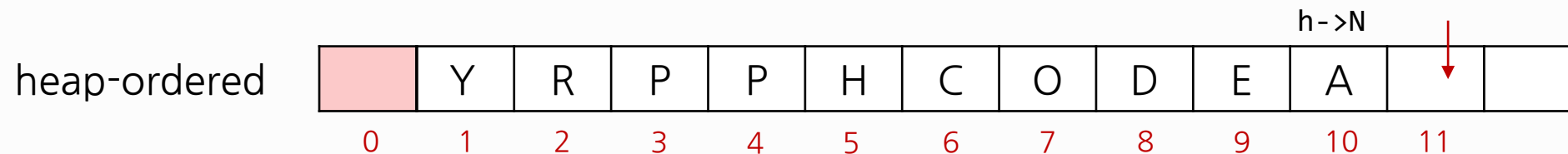
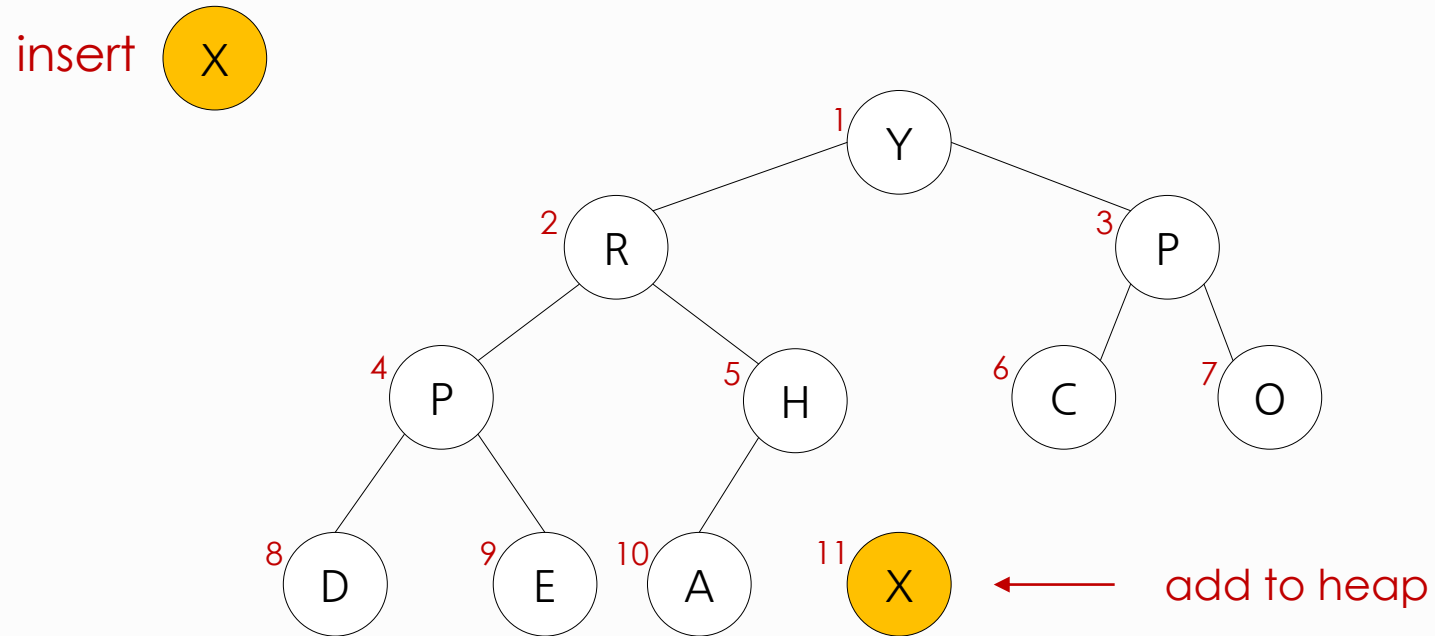


heap-ordered



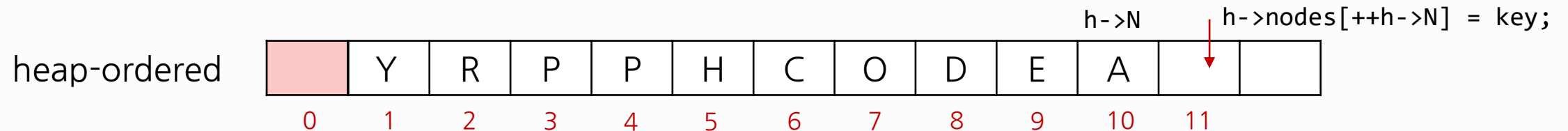
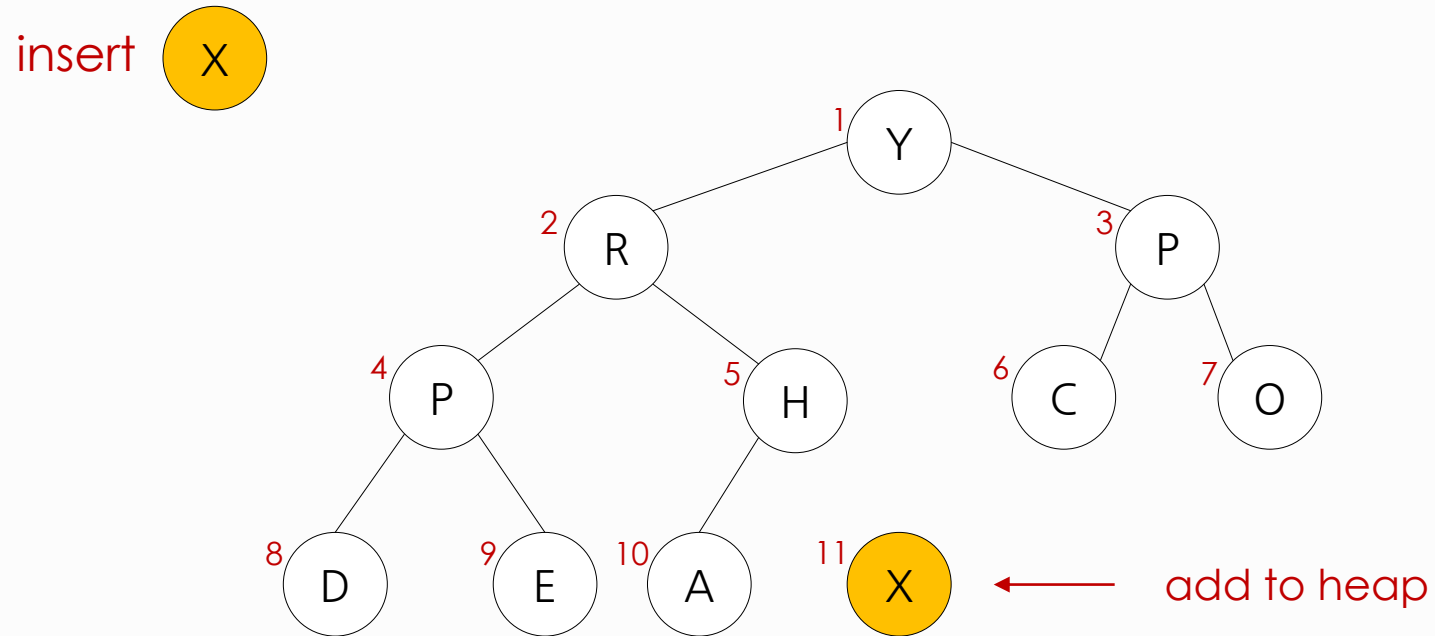
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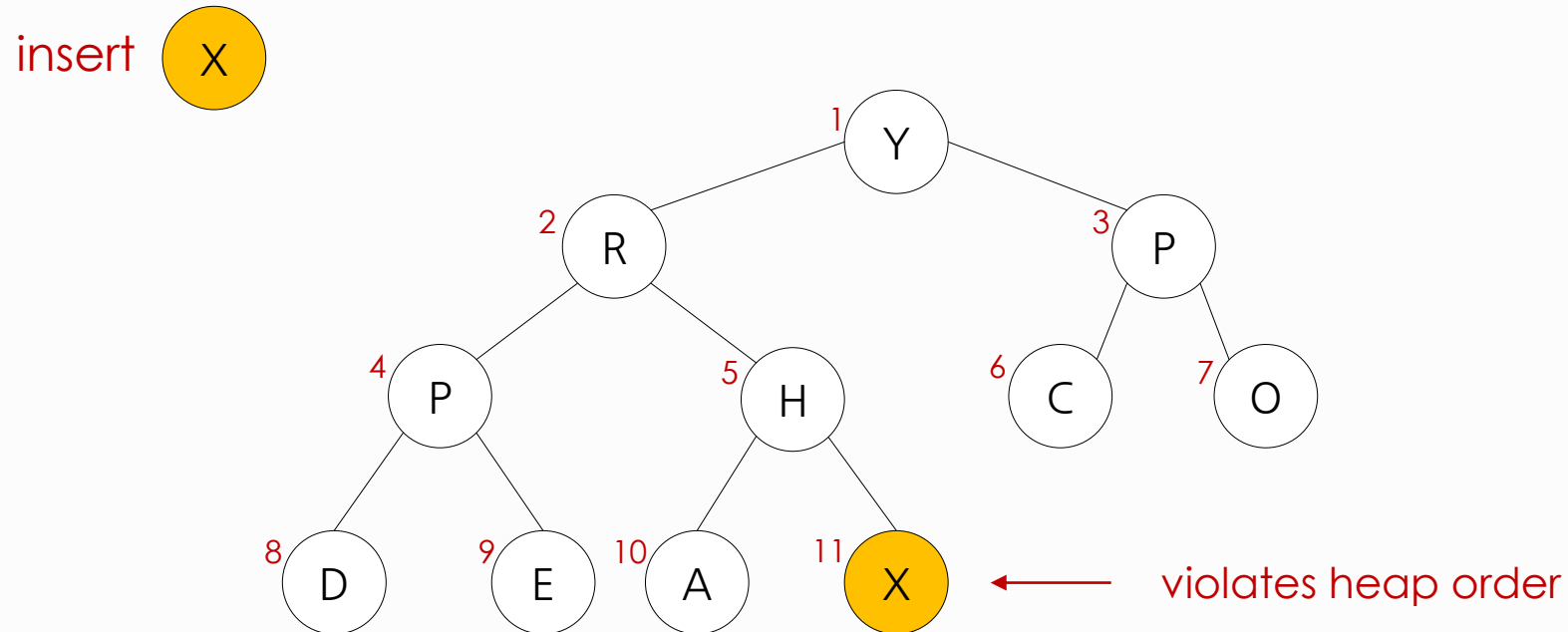
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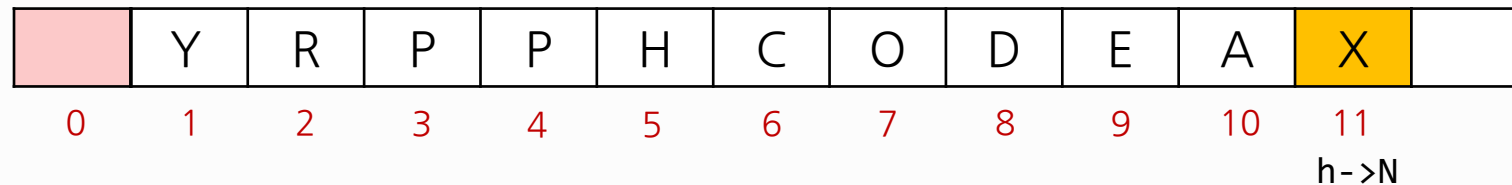


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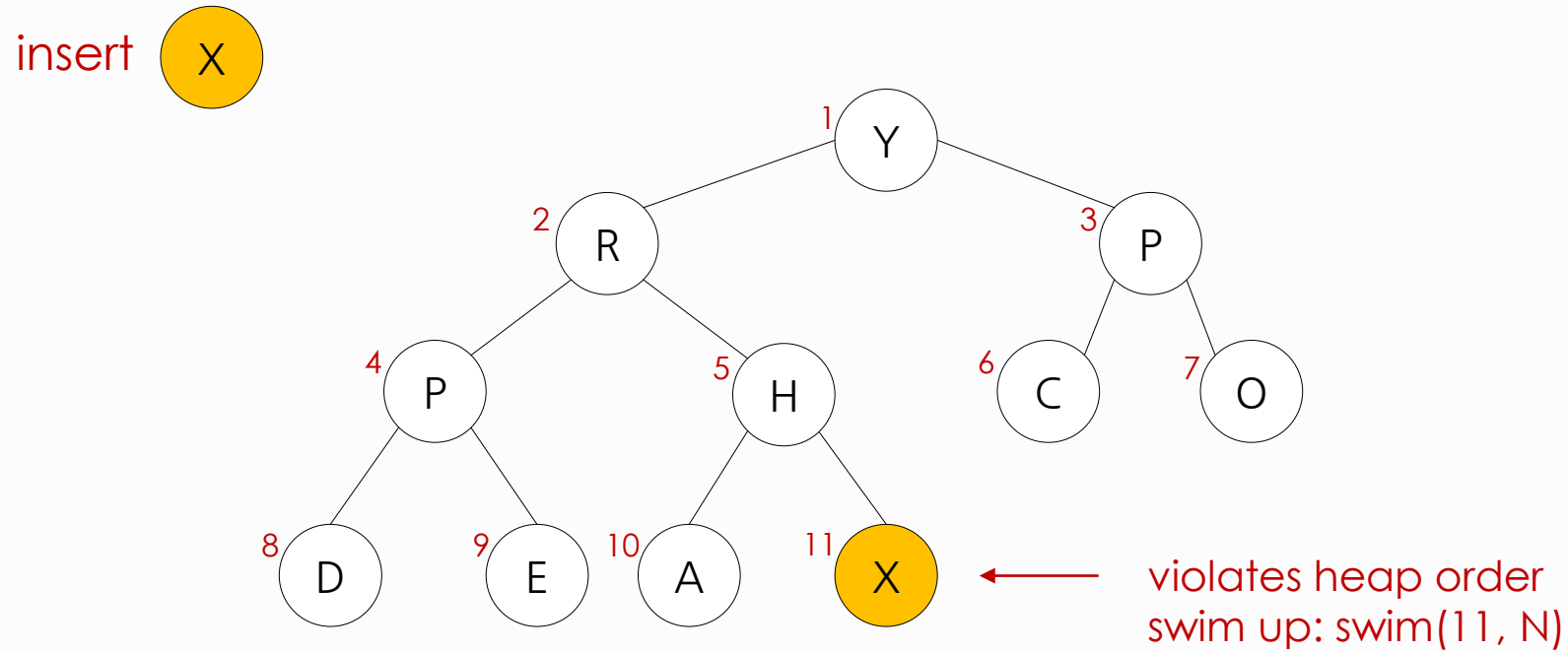
heap-ordered



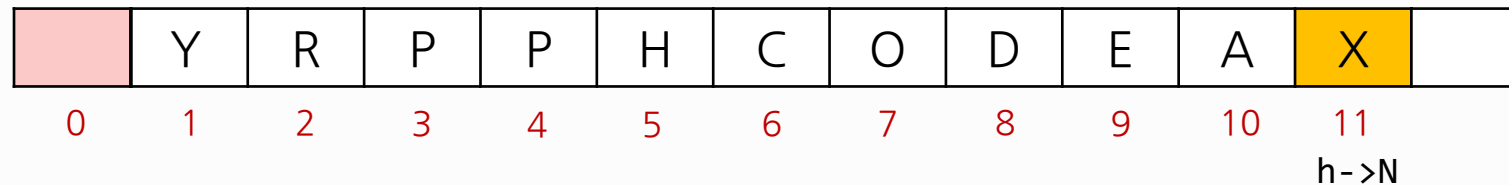


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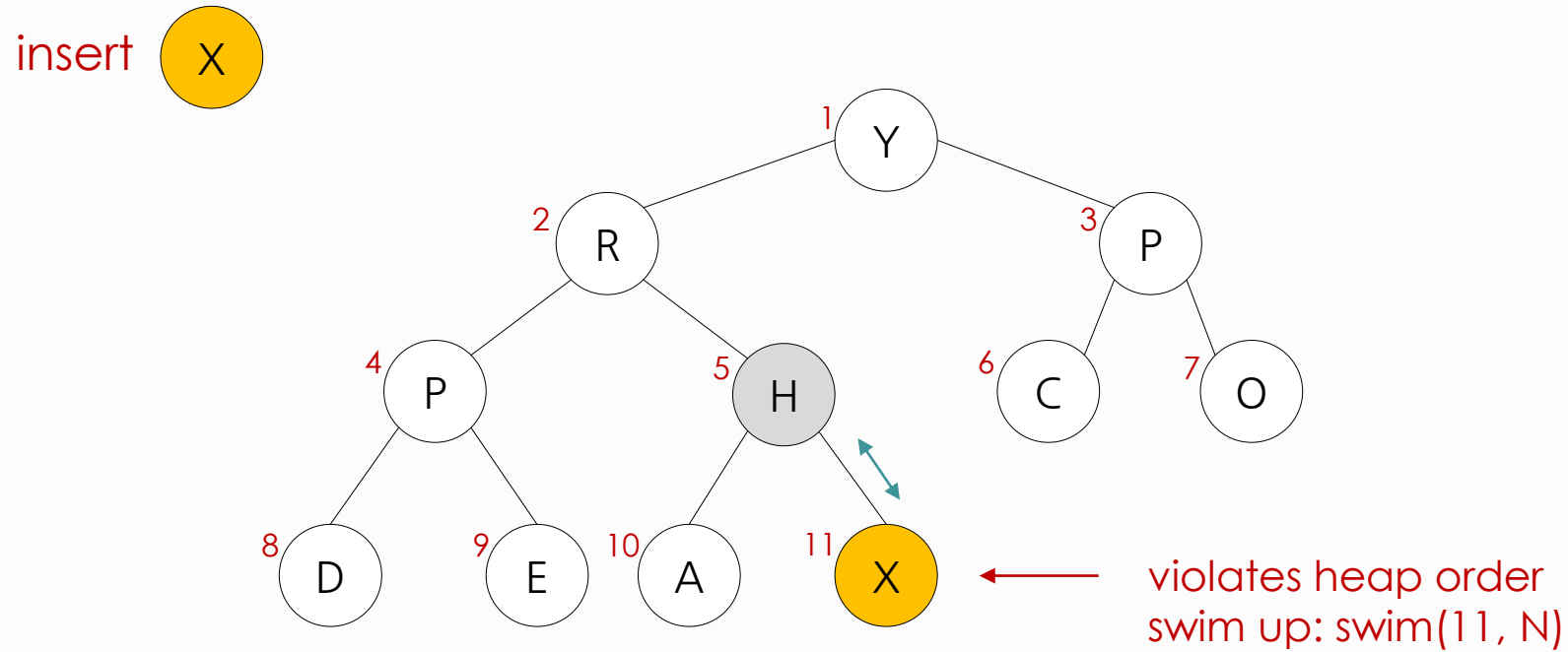


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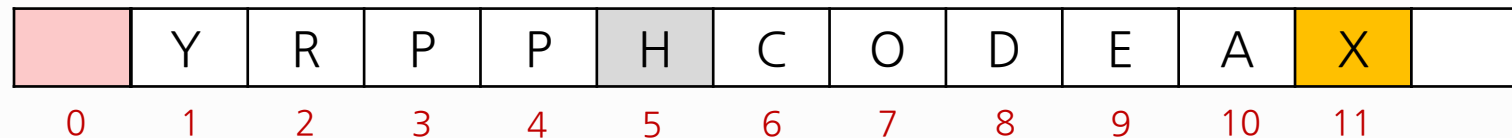


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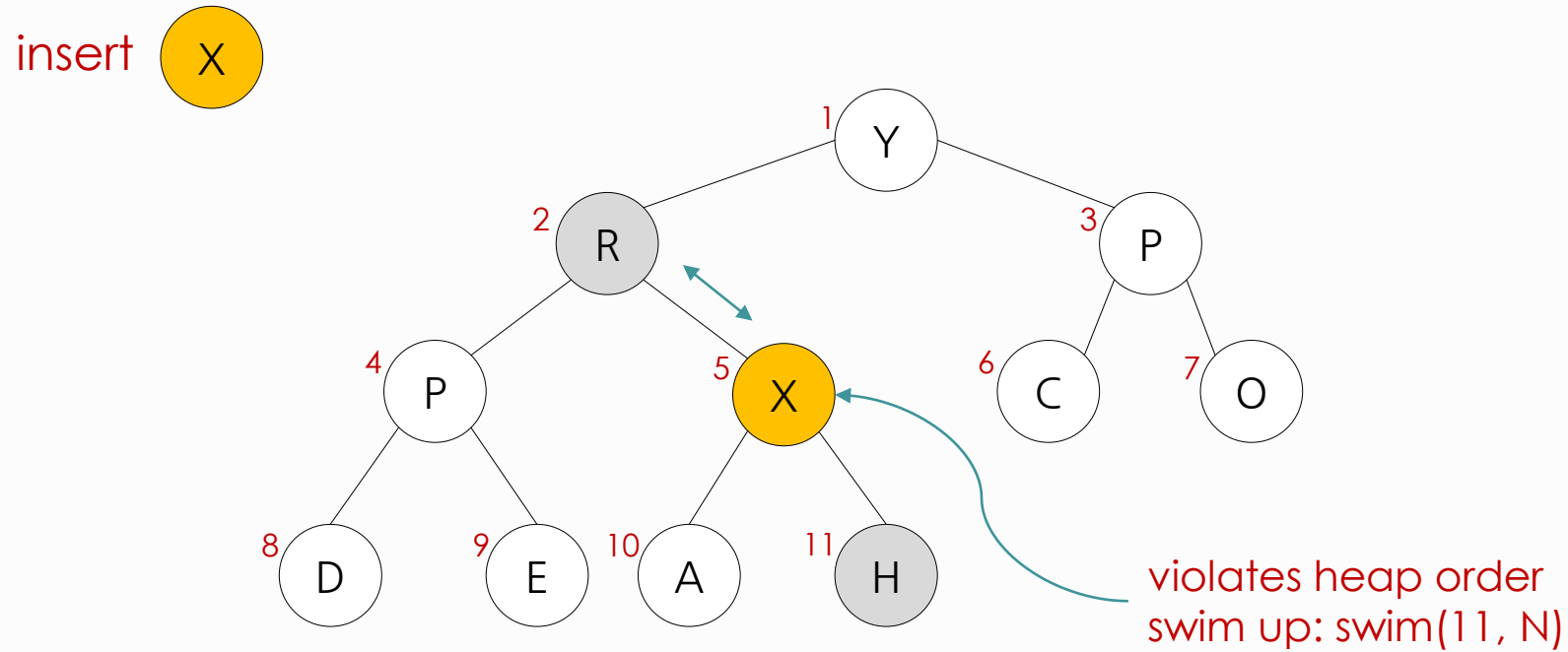


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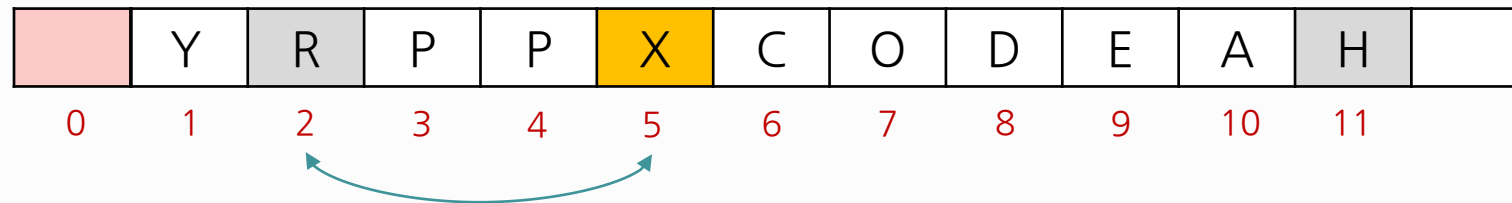


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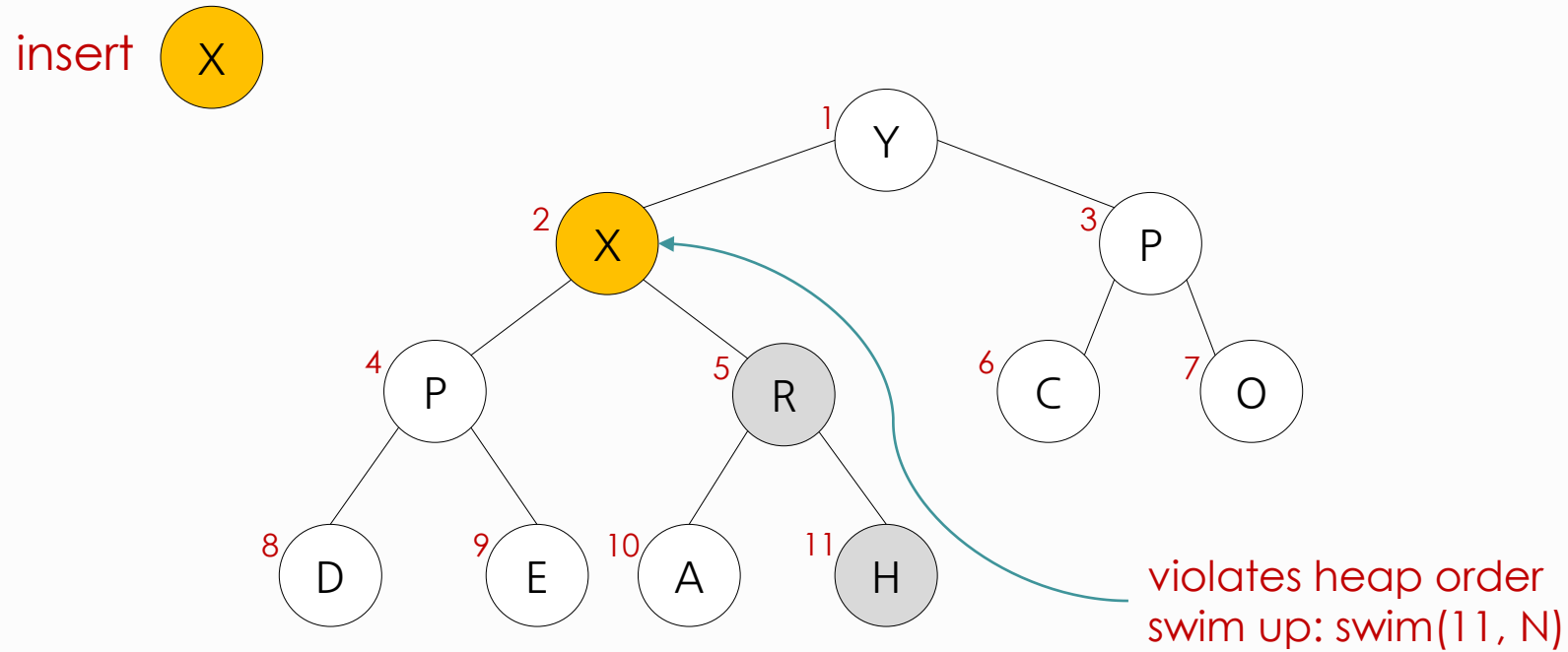


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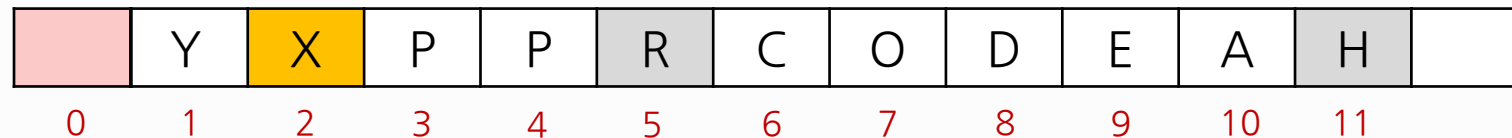


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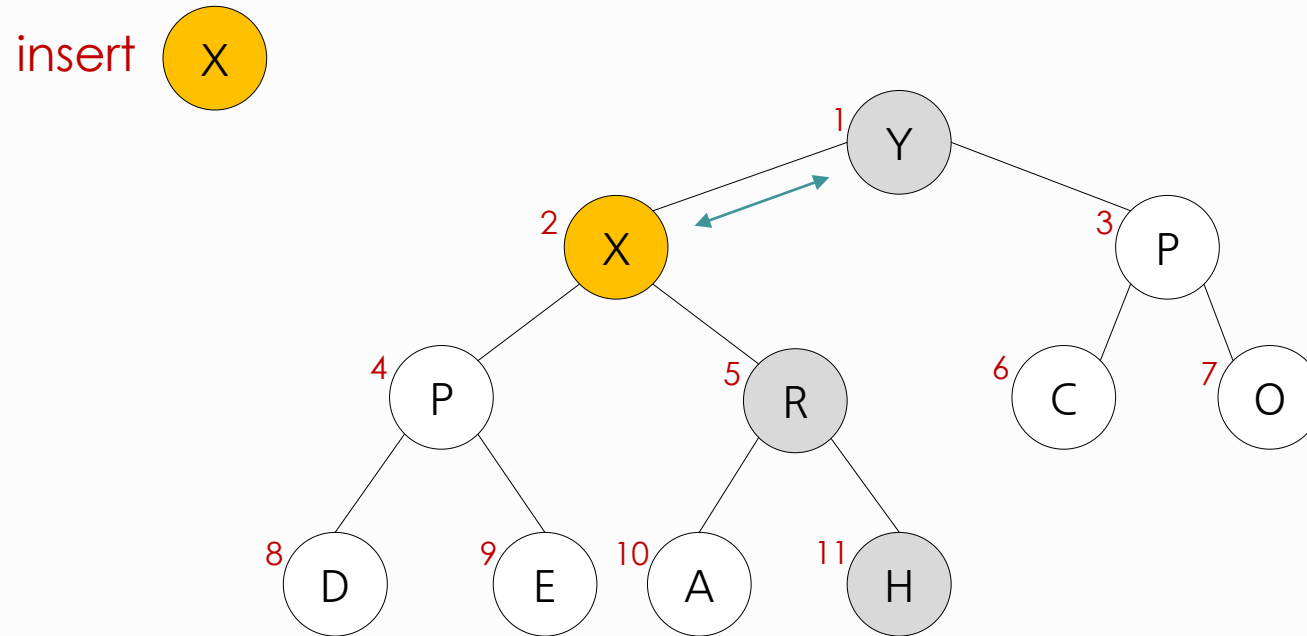


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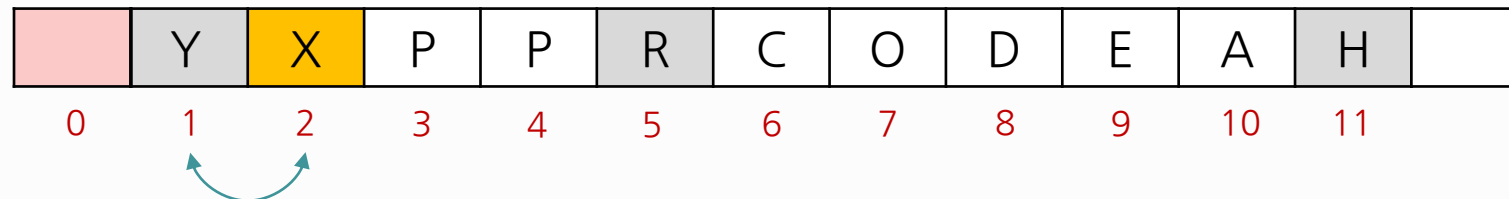


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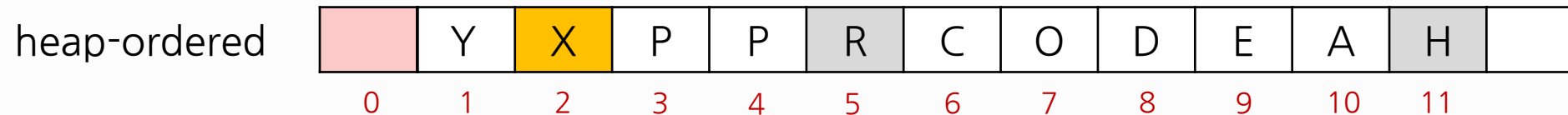
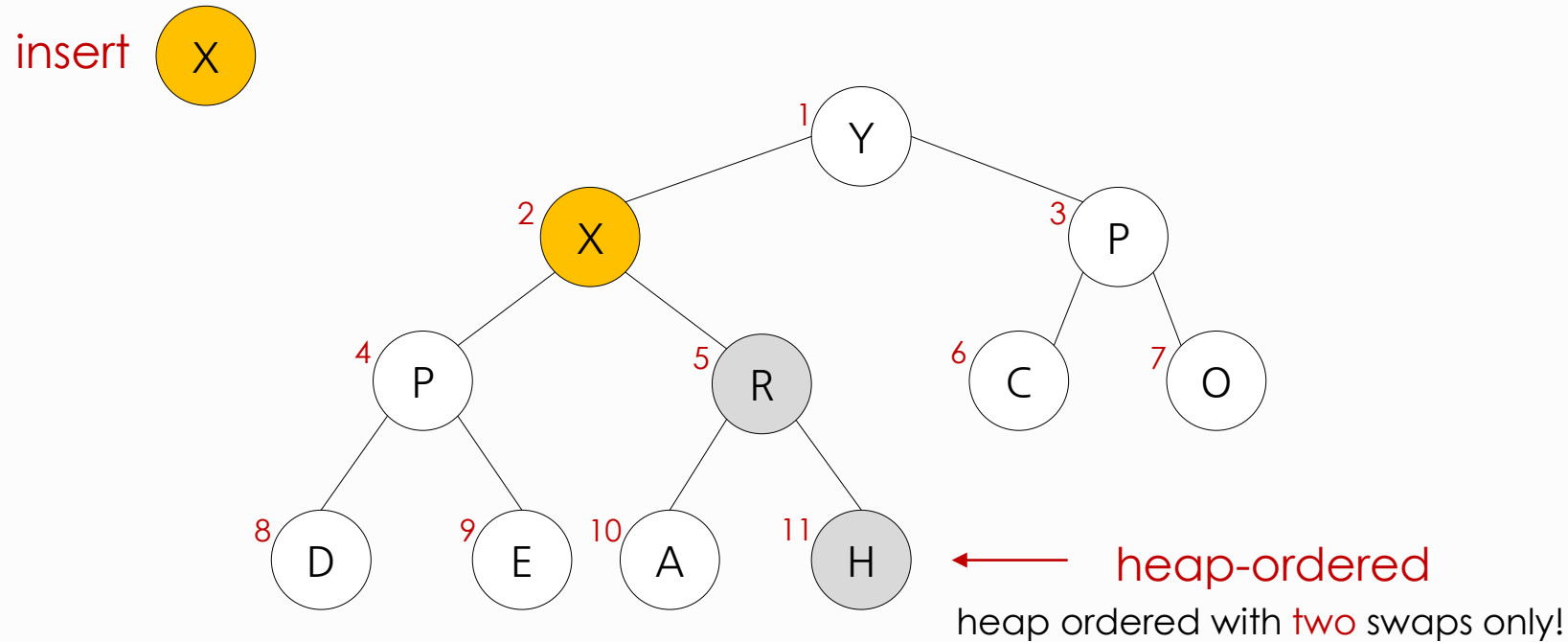


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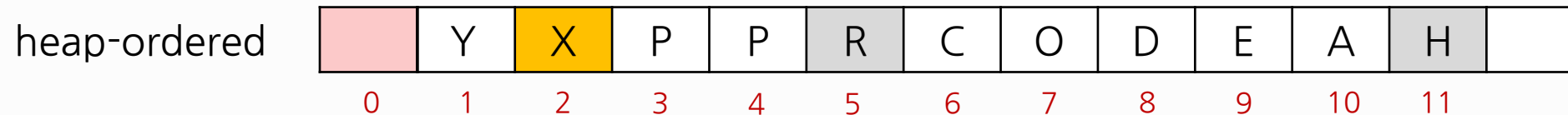
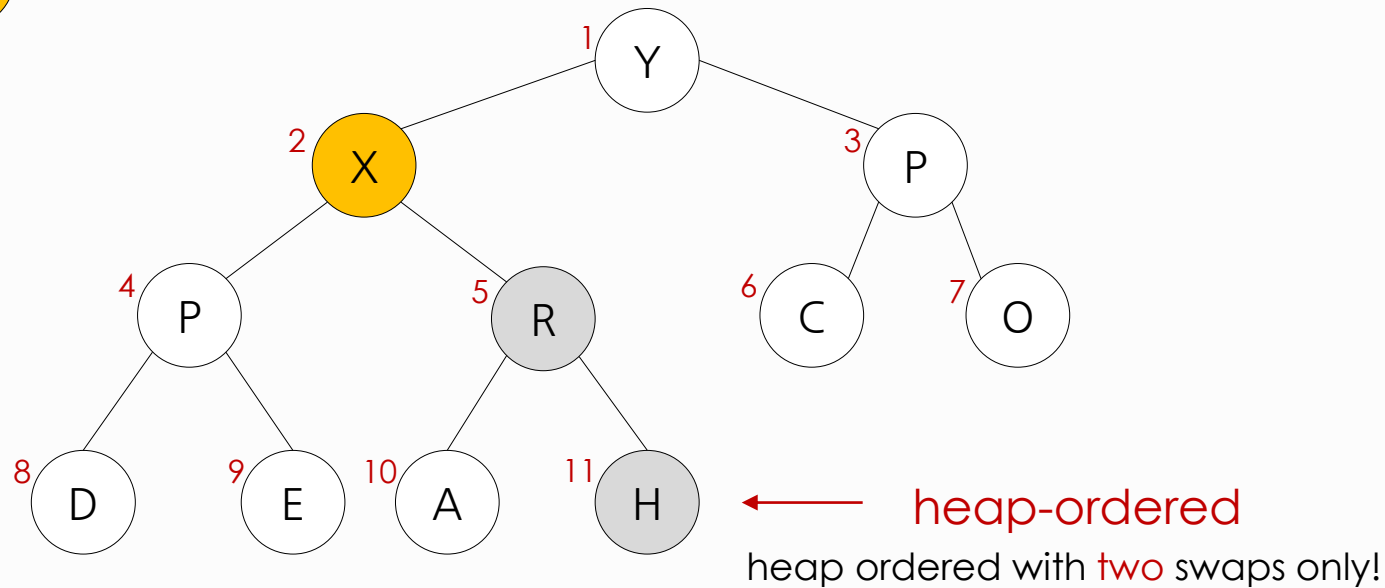
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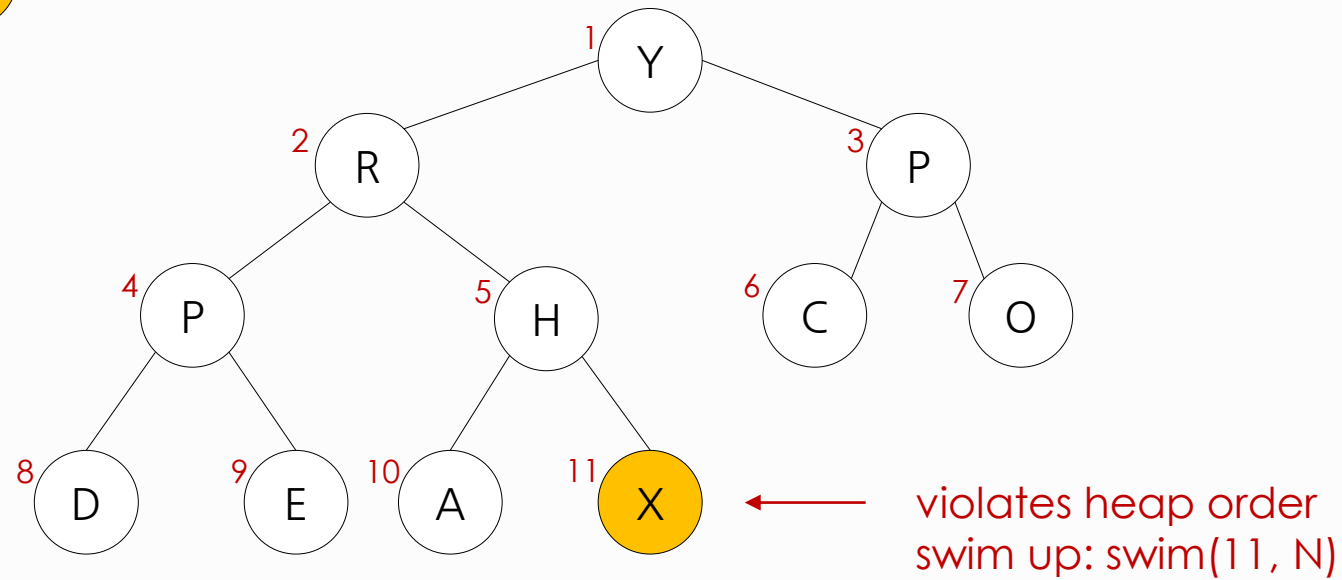
How to code insert 



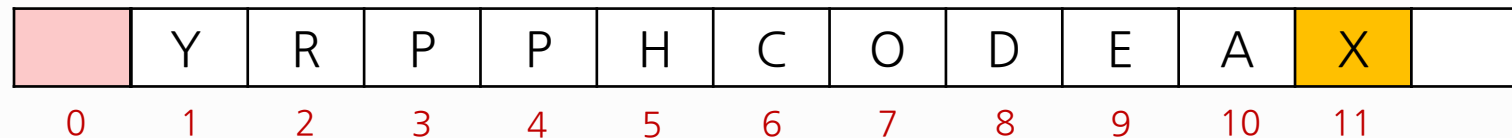
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heap-ordered

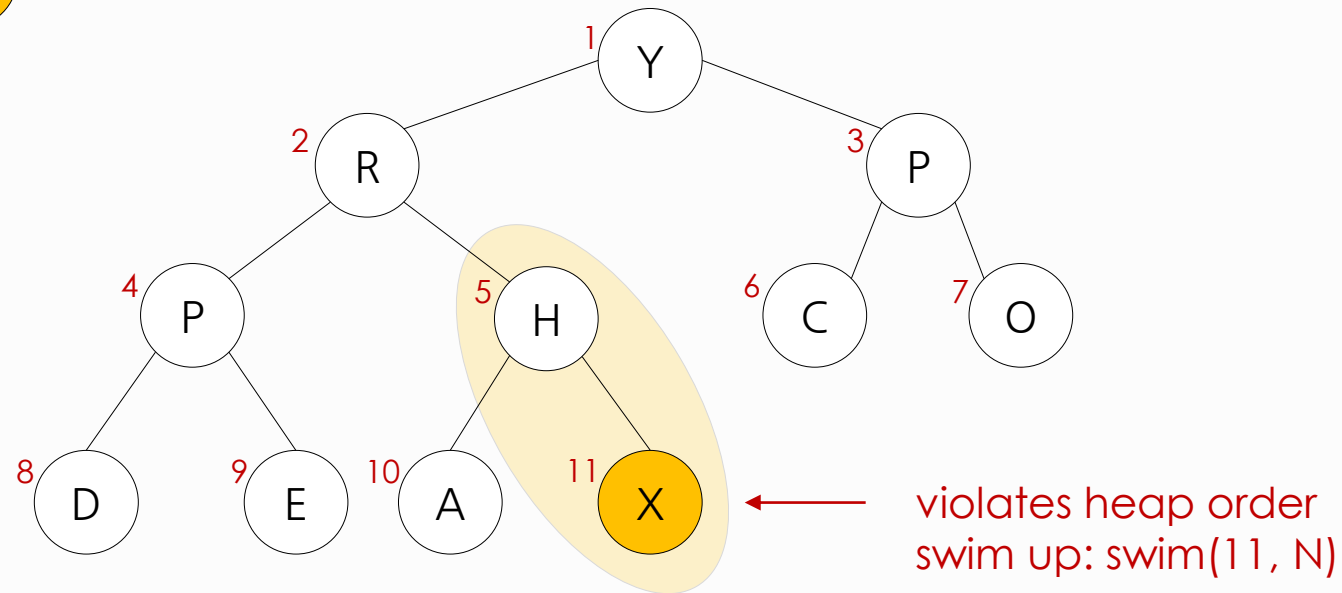




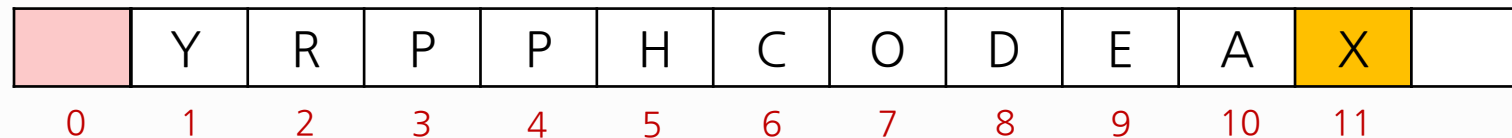
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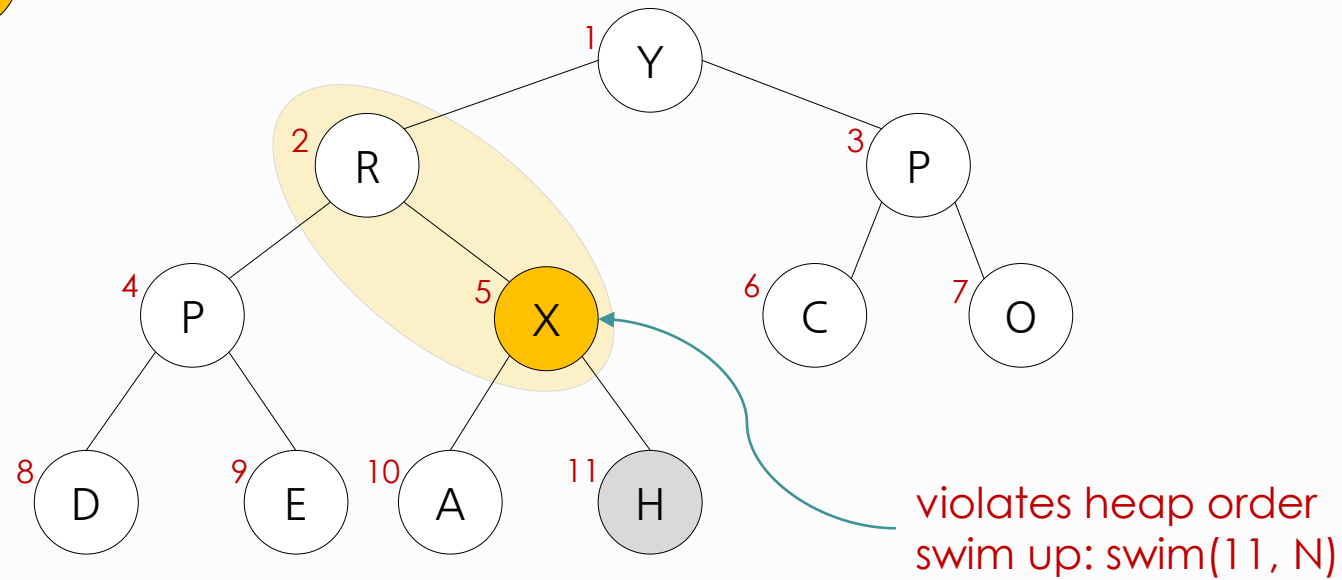
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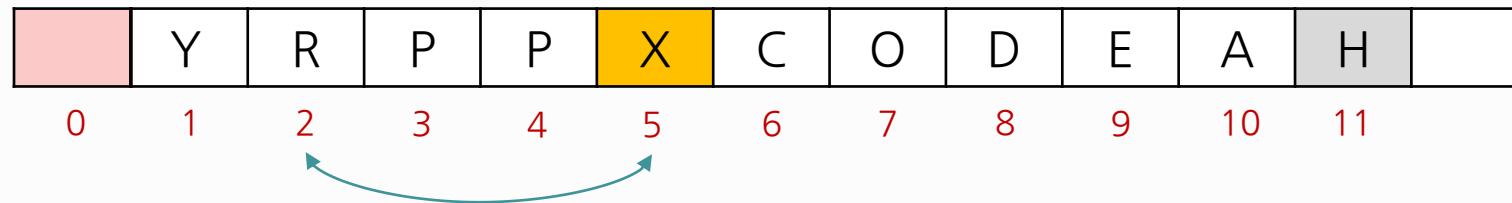
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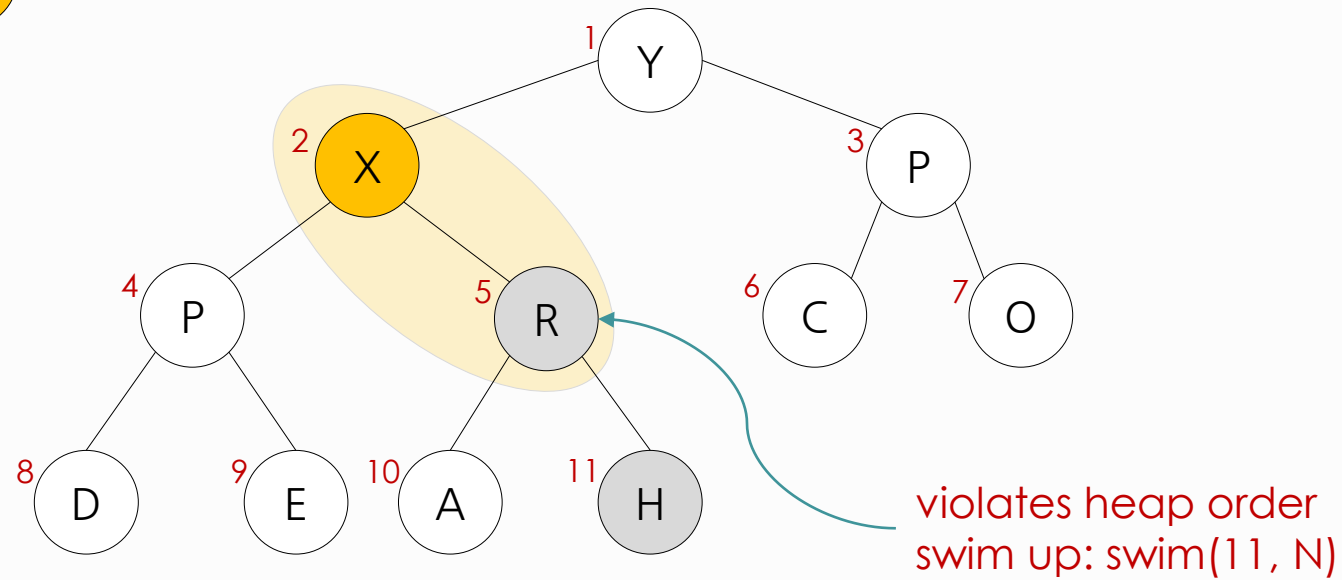
heap-ordered



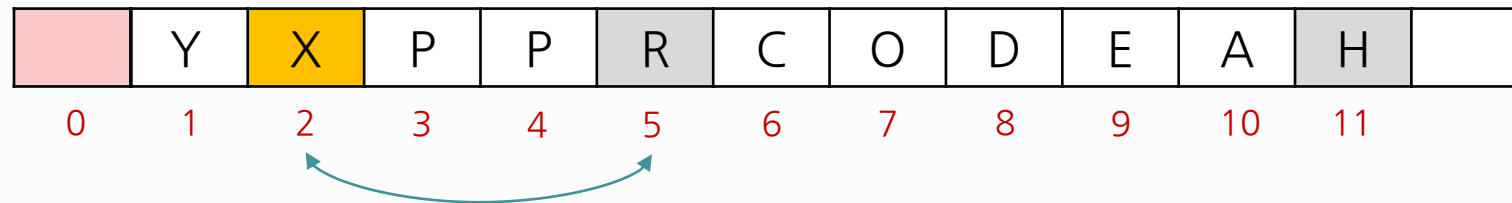
# max-heap example

- **Insert:** Add node at end, then swim it up.
- **Remove:** Swap root with last node, then sink down.

How to code insert 



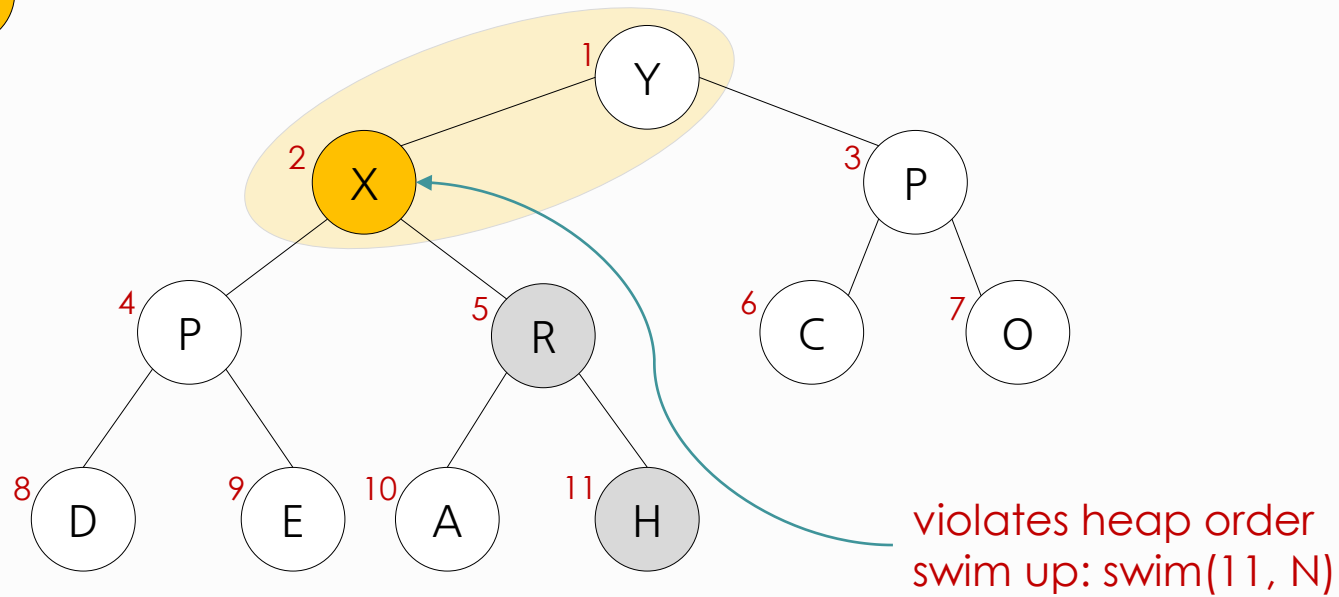
heap-ordered



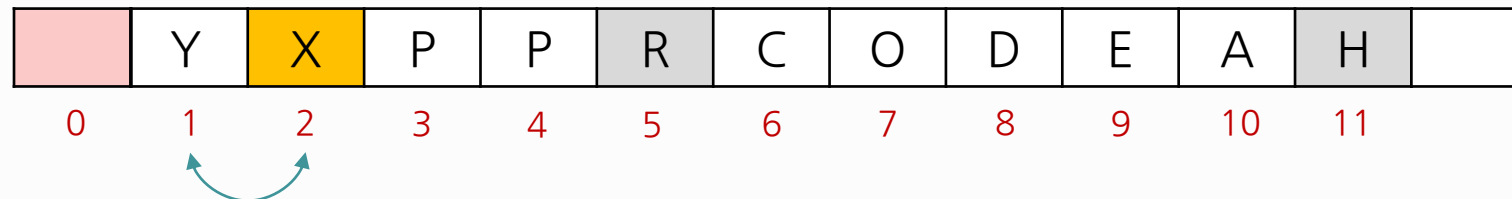
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- **Insert:** Add node at end, then swim it up.
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How to code insert 



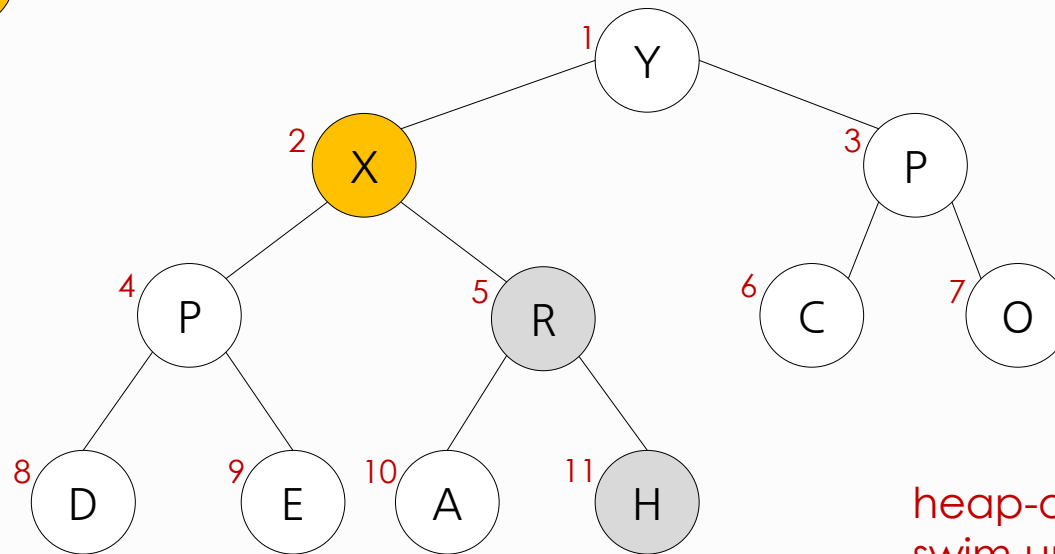
heap-ordered



# max-heap example

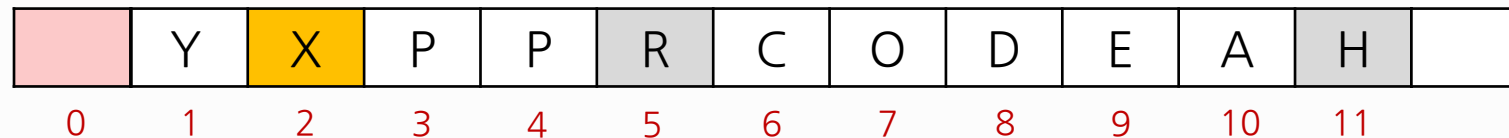
- **Insert:** Add node at end, then swim it up.
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How to code insert 



heap-ordered  
swim up: swim(11, N)

heap-ordered

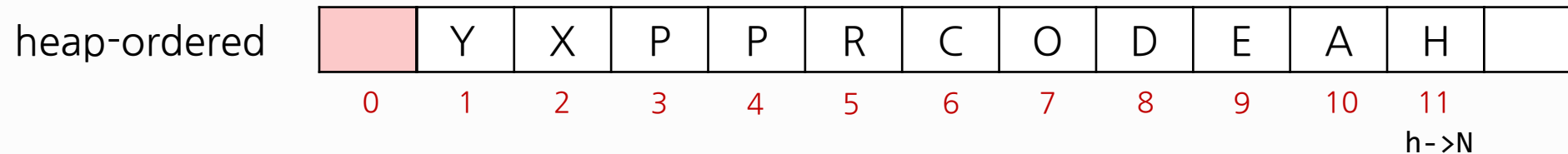


## max-heap example

---

- **Insert:** Add node at end, then swim it up.
- **Remove:** Swap root with last node, then sink down.

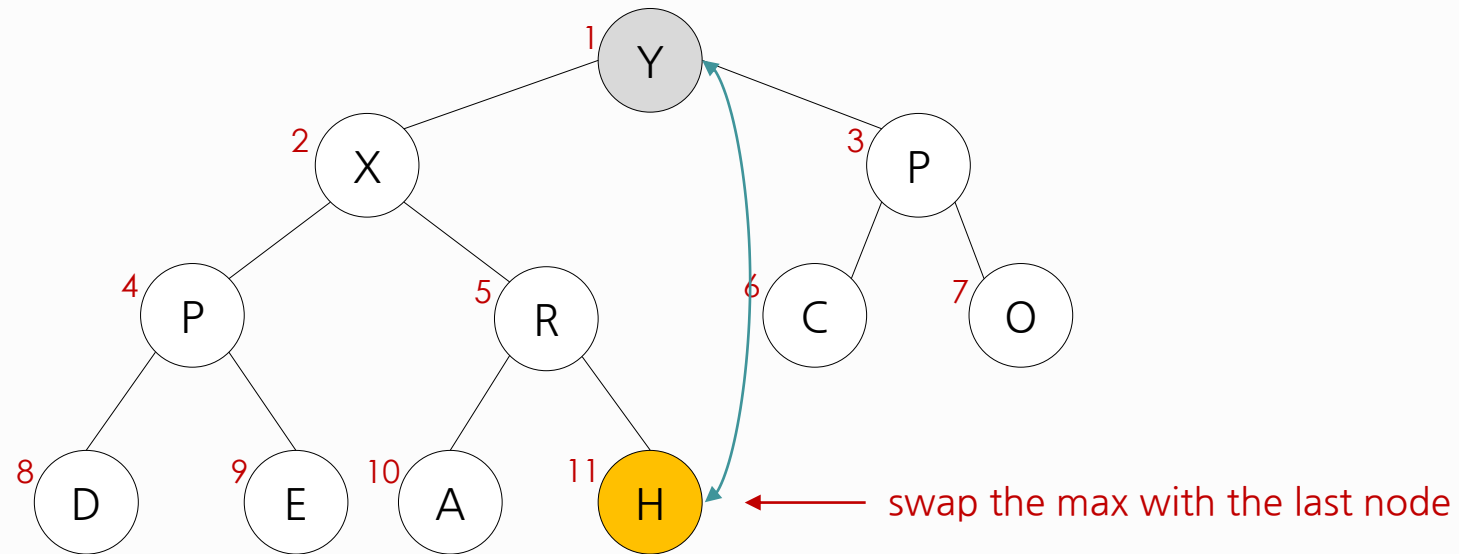
remove the max (root) – pop in priority queue



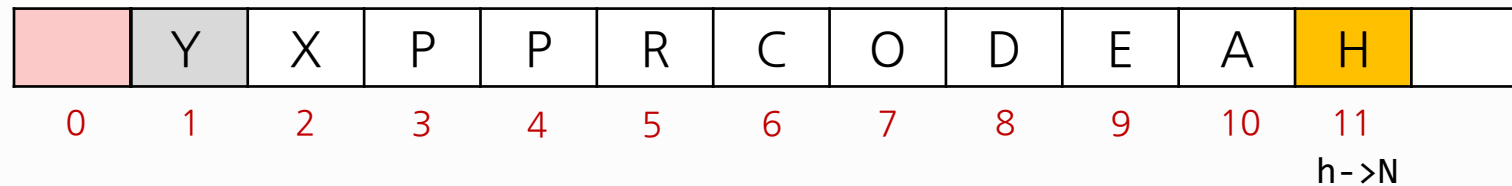
# max-heap example

- **Insert:** Add node at end, then swim it up.
- **Remove:** Swap root with last node, then sink down.

remove the max (root)



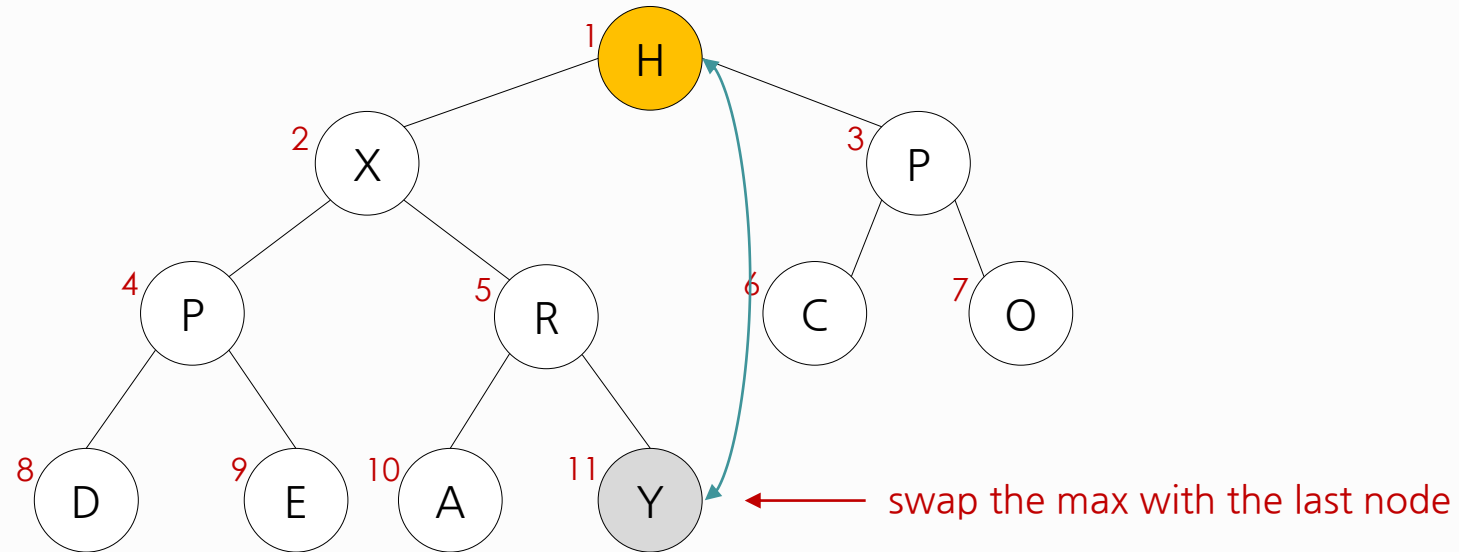
heap-ordered



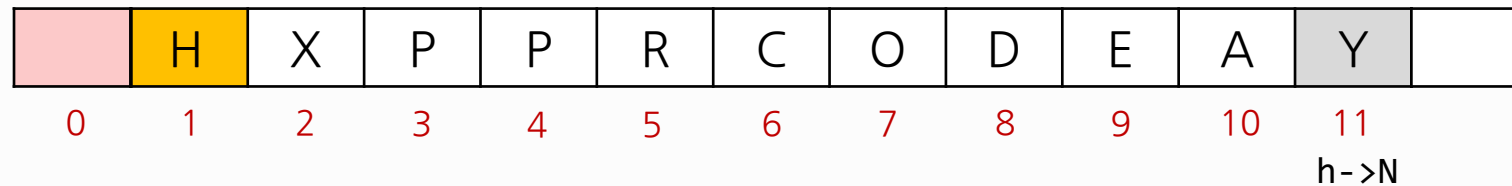
# max-heap example

- **Insert:** Add node at end, then swim it up.
- **Remove:** Swap root with last node, then sink down.

remove the max (root)



heap-ordered

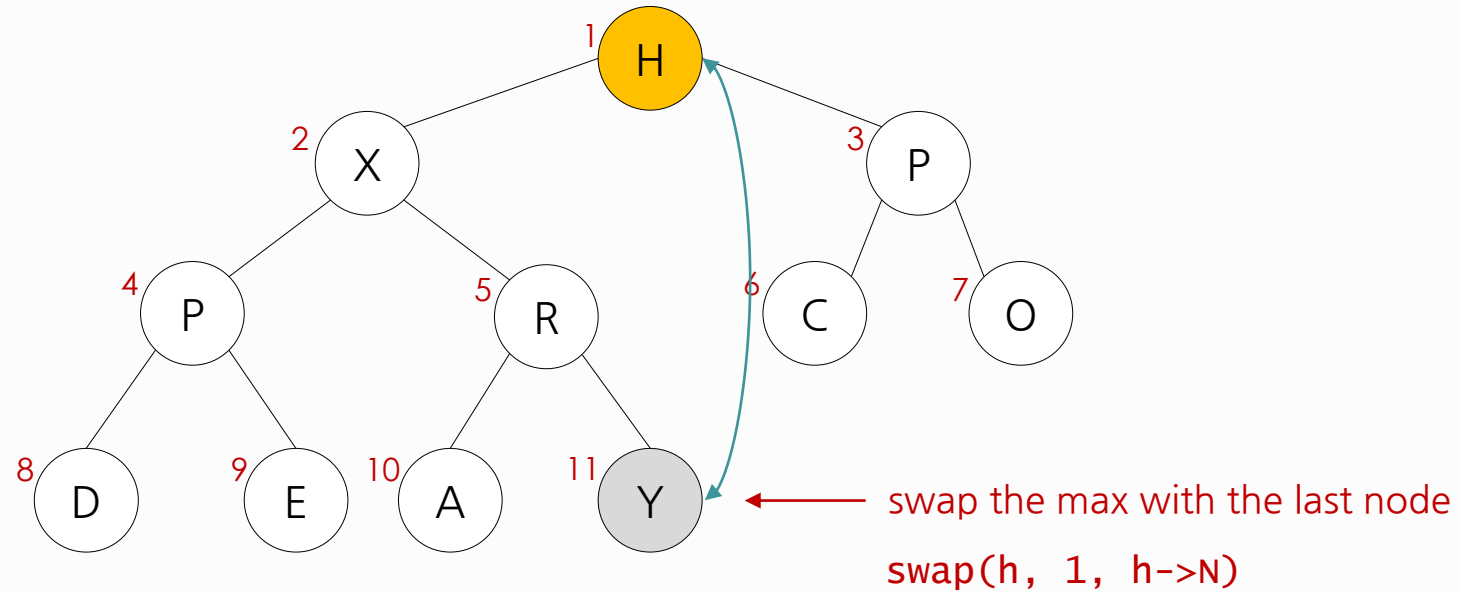




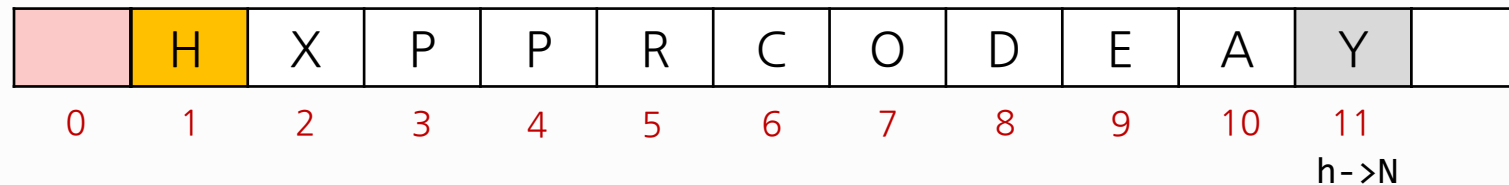
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- **Remove:** Swap root with last node, then sink down.

remove the max (root)



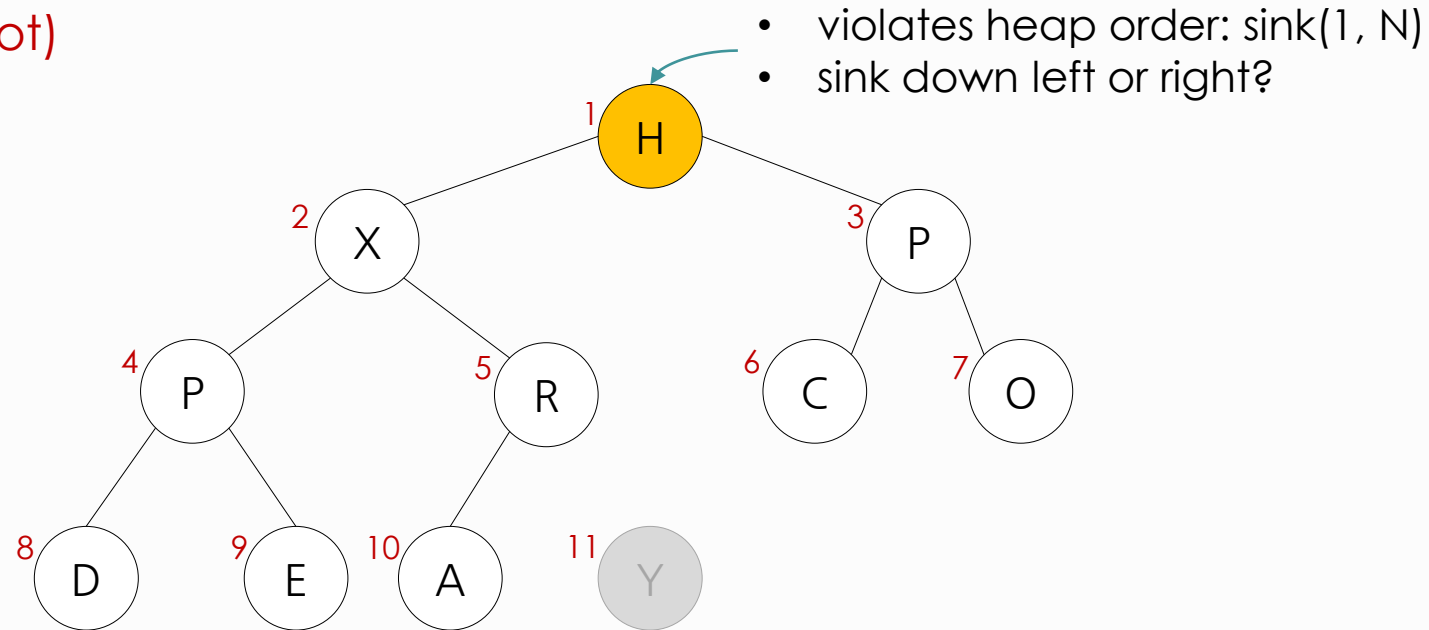
heap-ordered



# max-heap example

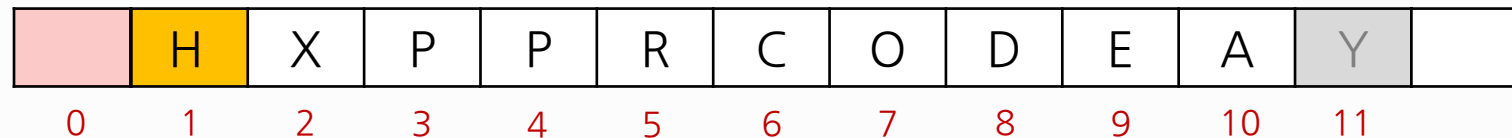
- **Insert:** Add node at end, then swim it up.
- **Remove:** Swap root with last node, then sink down.

remove the max (root)



swap(h, 1, h->N--)

heap-ordered

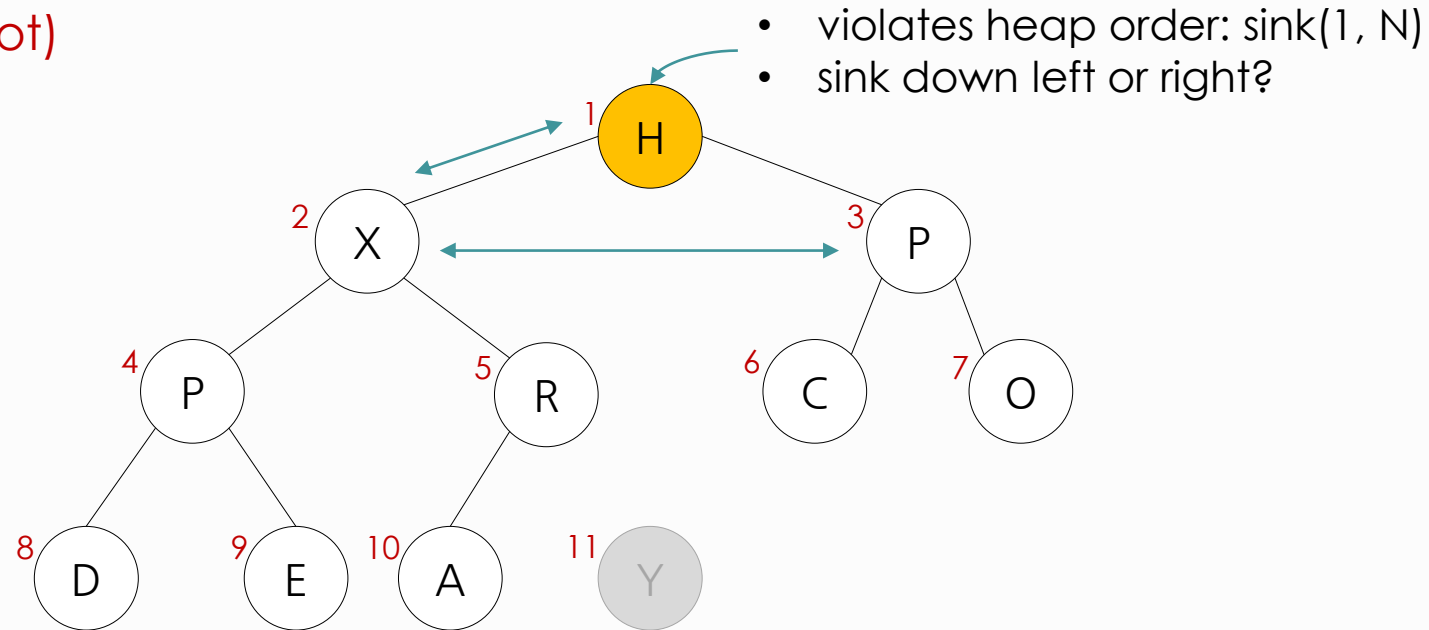


h->N heap size decreased by one

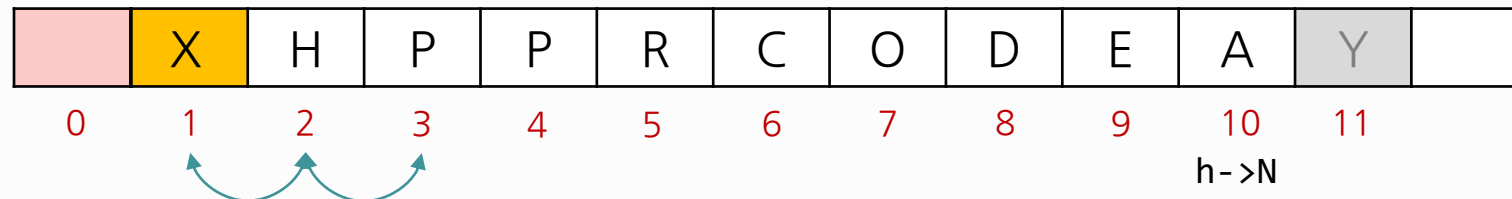
# max-heap example

- **Insert:** Add node at end, then swim it up.
- **Remove:** Swap root with last node, then sink down.

remove the max (root)



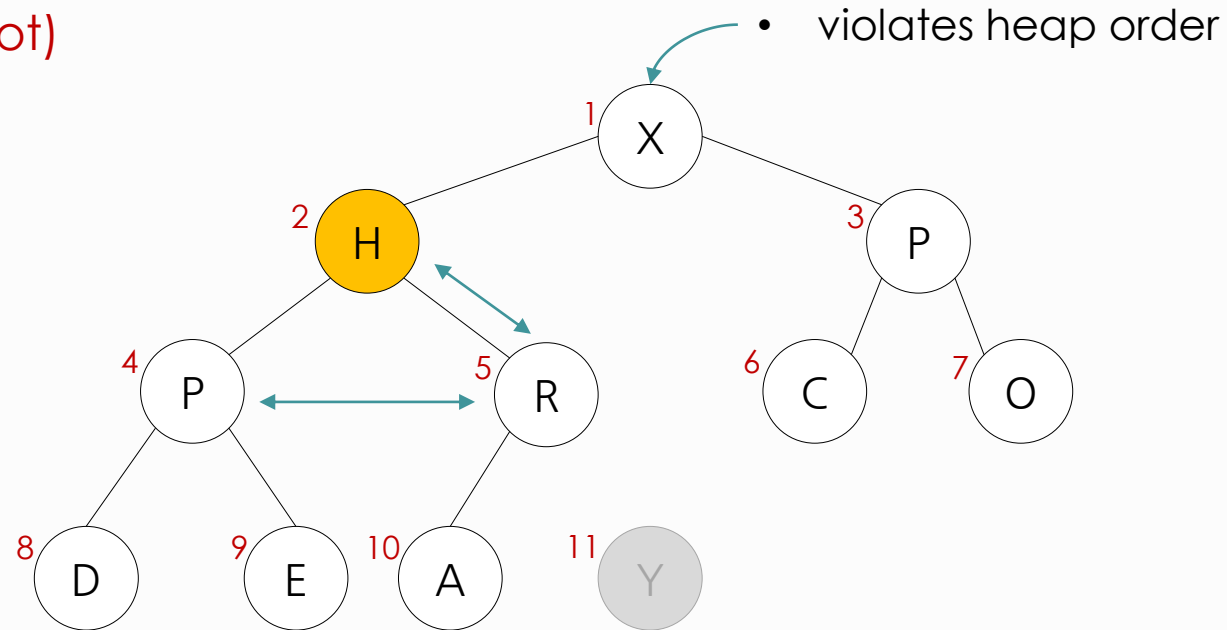
heap-ordered



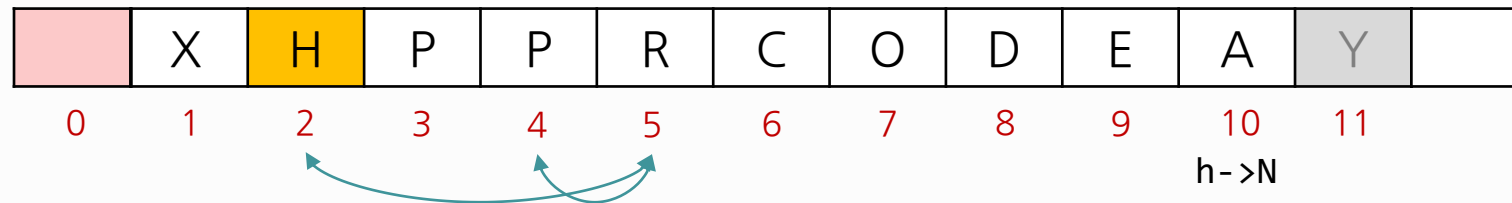
# max-heap example

- **Insert:** Add node at end, then swim it up.
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remove the max (root)



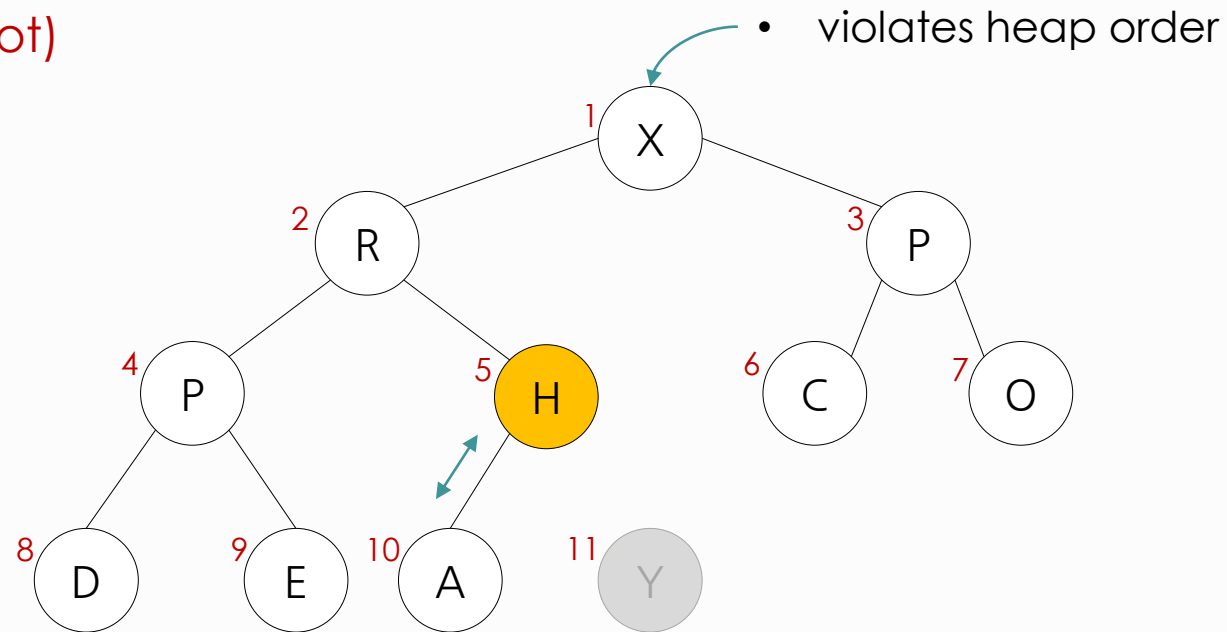
heap-ordered



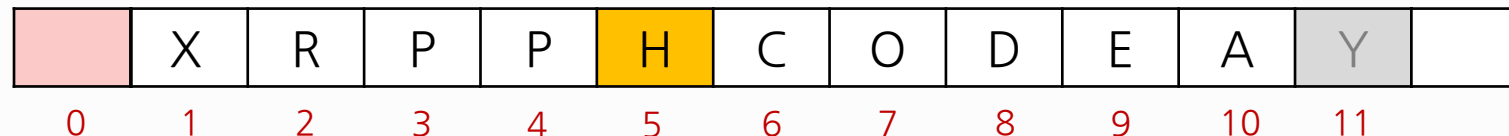
# max-heap example

- **Insert:** Add node at end, then swim it up.
- **Remove:** Swap root with last node, then sink down.

remove the max (root)



heap-ordered

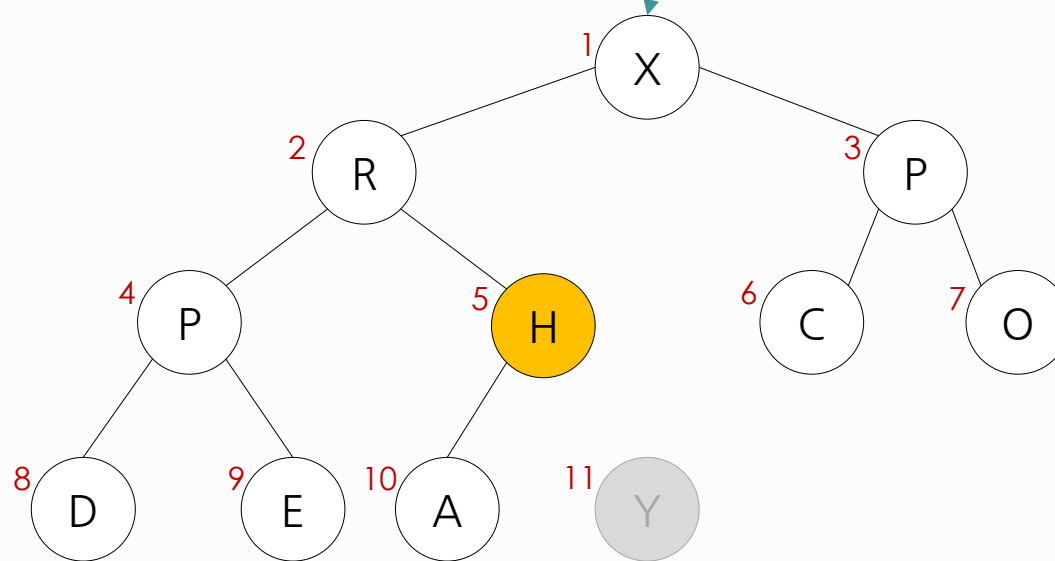


# max-heap example

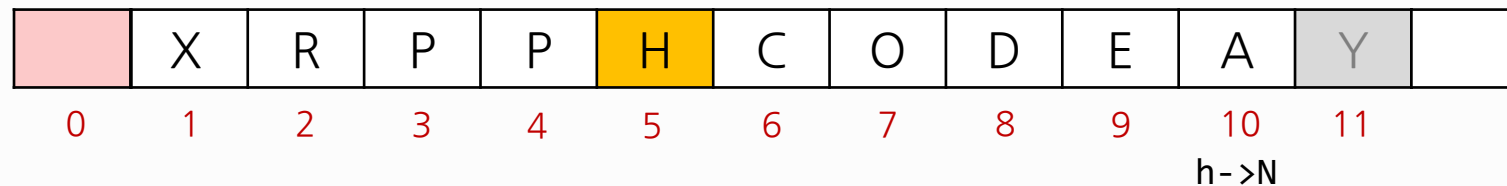
- **Insert:** Add node at end, then swim it up.
- **Remove:** Swap root with last node, then sink down.

remove the max (root)

• heap-ordered



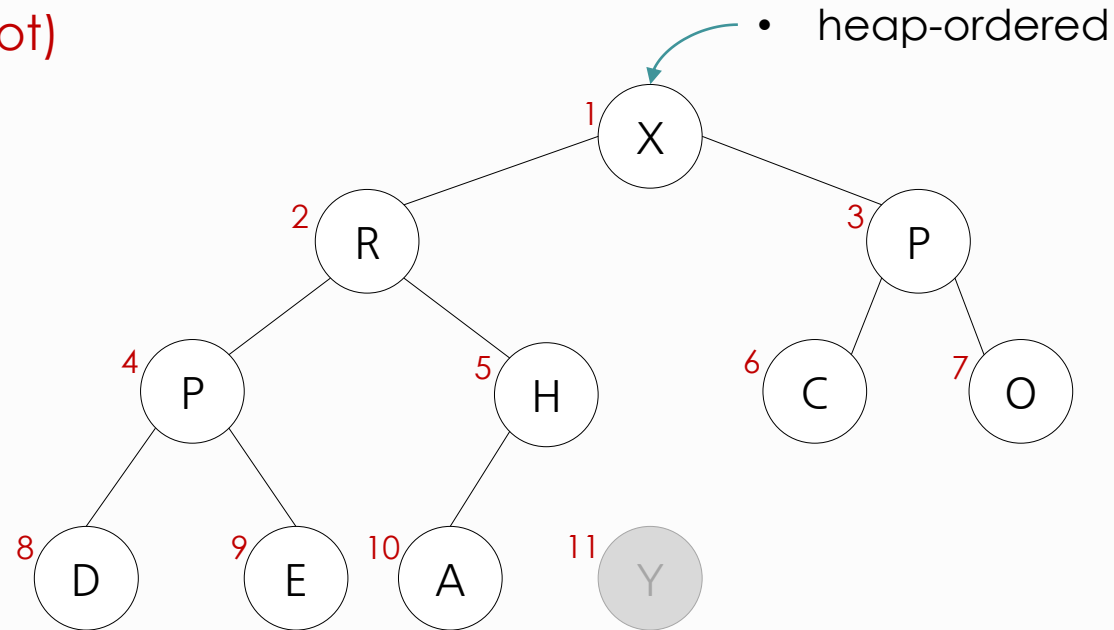
heap-ordered



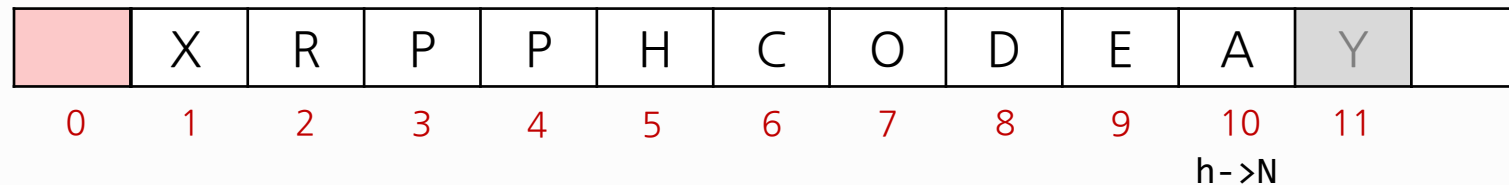
# max-heap example

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remove the max (root)



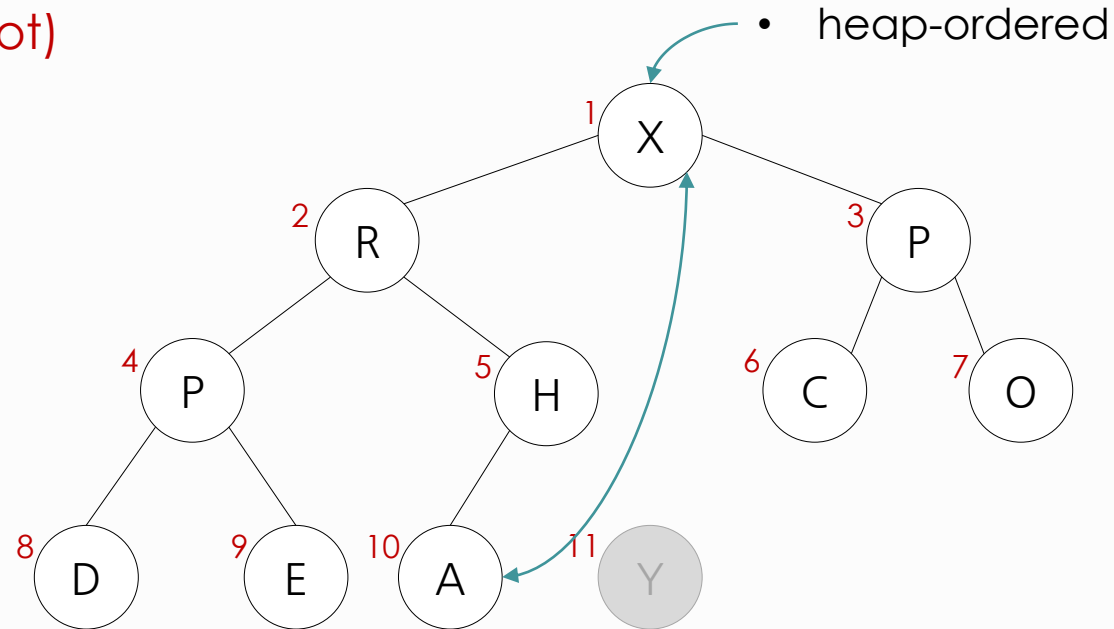
heap-ordered



# max-heap example

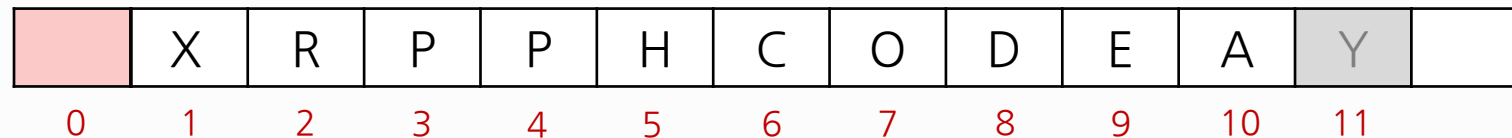
- **Insert:** Add node at end, then swim it up.
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remove the max (root)



swap(h, 1, h->N--)

heap-ordered



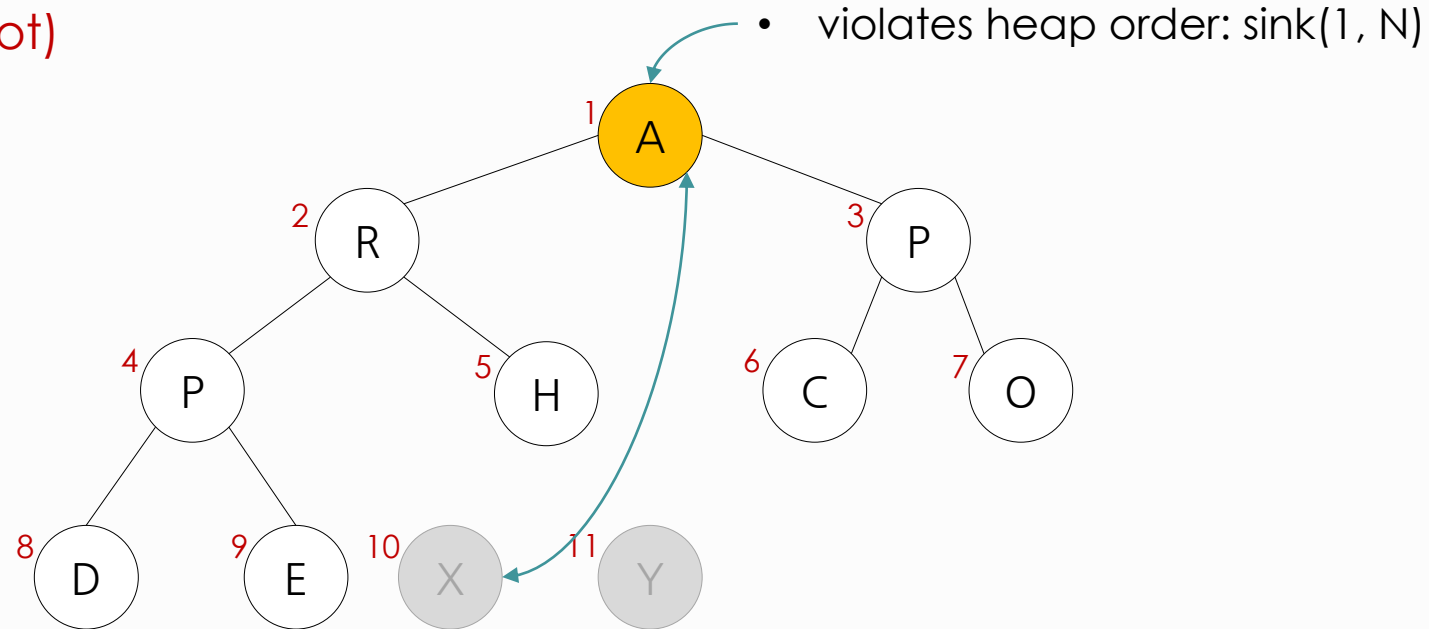
h->N heap size decreased by one



# max-heap example

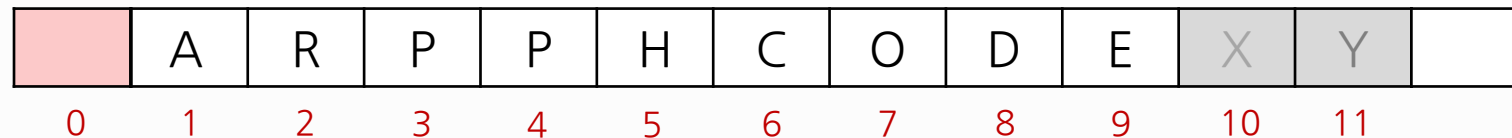
- **Insert:** Add node at end, then swim it up.
- **Remove:** Swap root with last node, then sink down.

remove the max (root)



swap(h, 1, h->N--)

heap-ordered

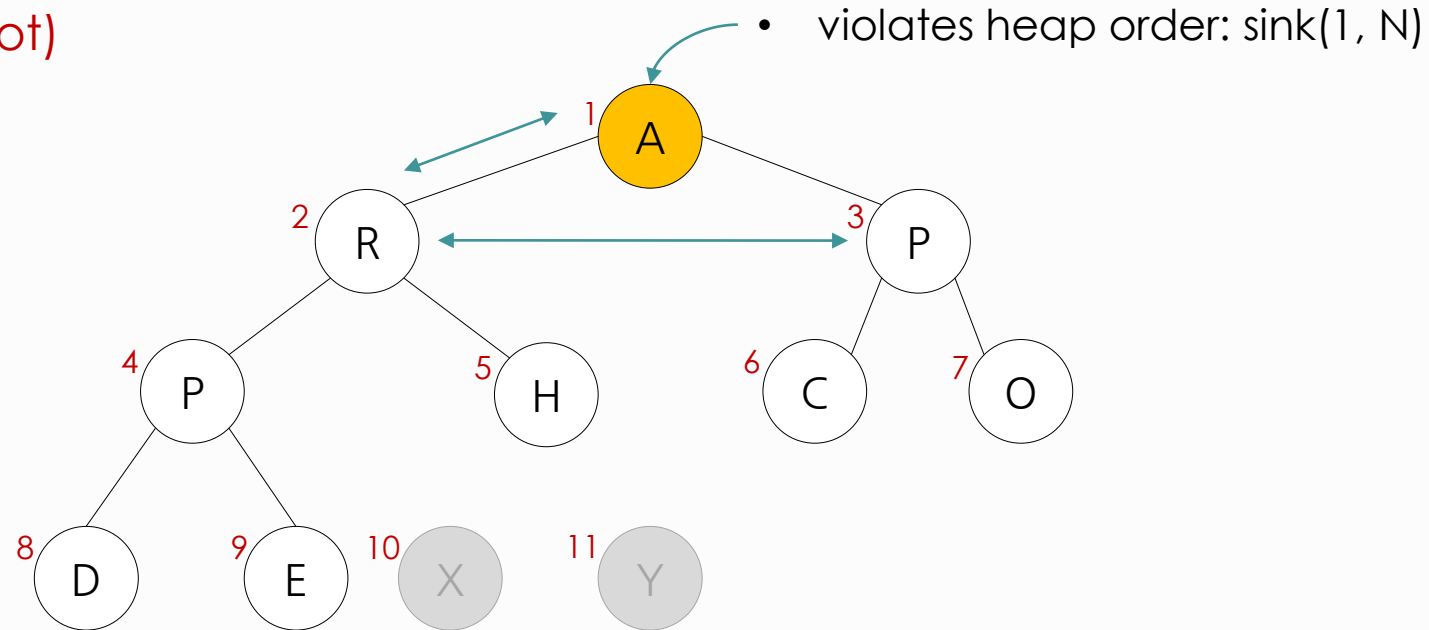


h->N heap size decreased by one

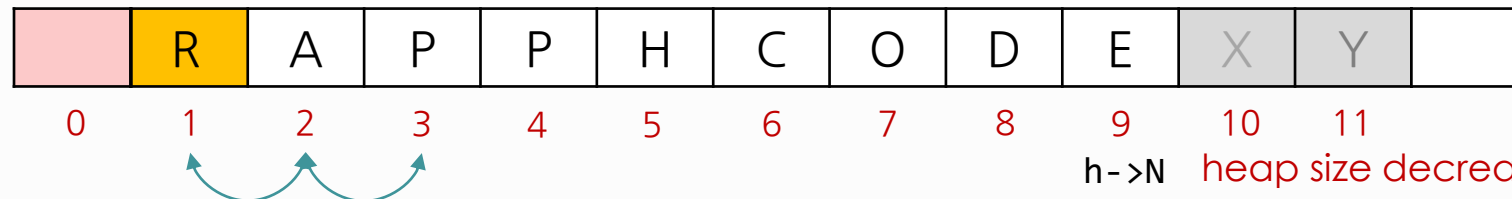
# max-heap example

- **Insert:** Add node at end, then swim it up.
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remove the max (root)



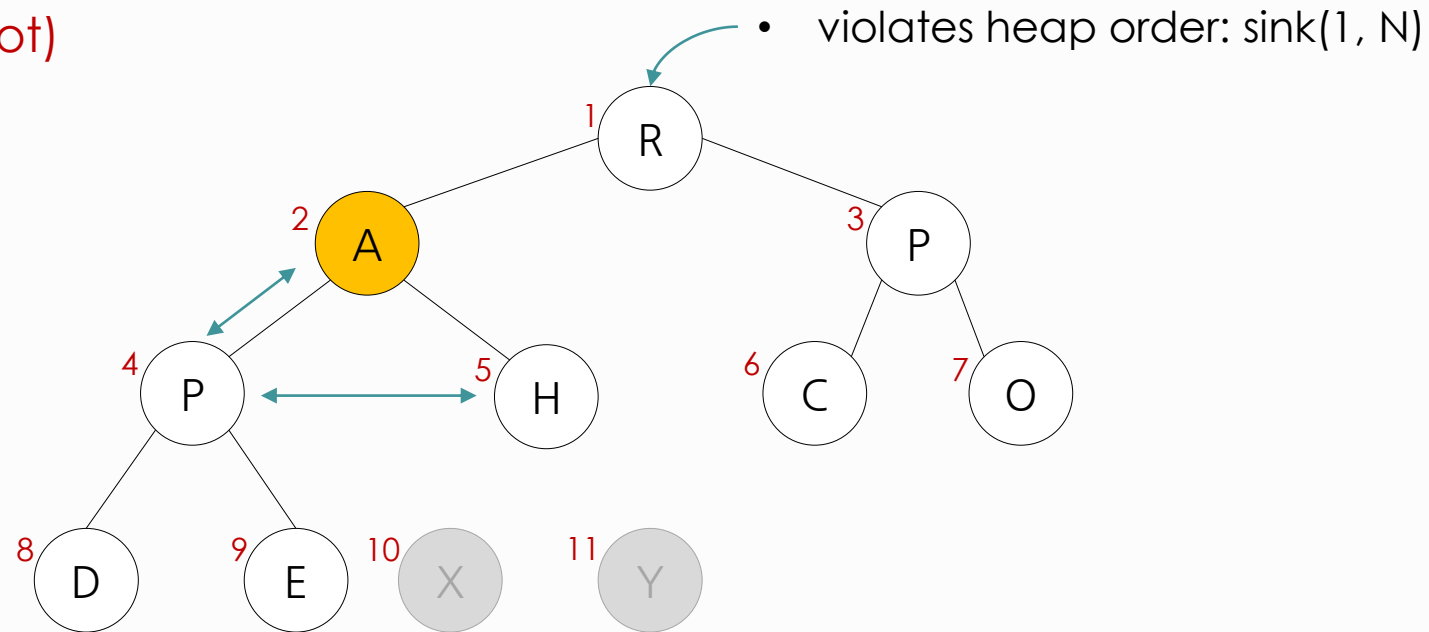
heap-ordered



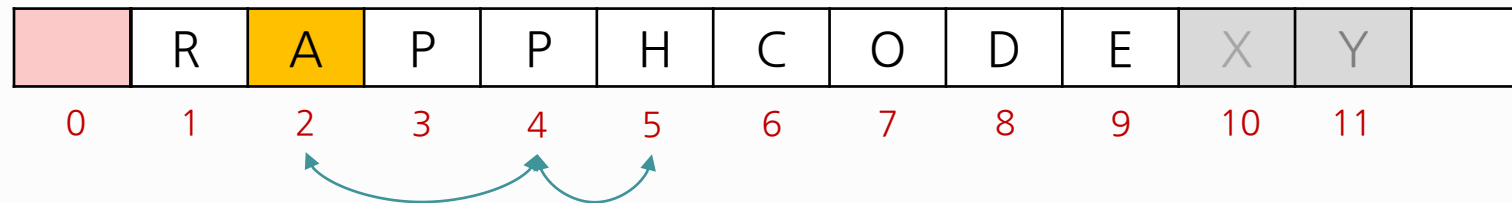
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remove the max (root)



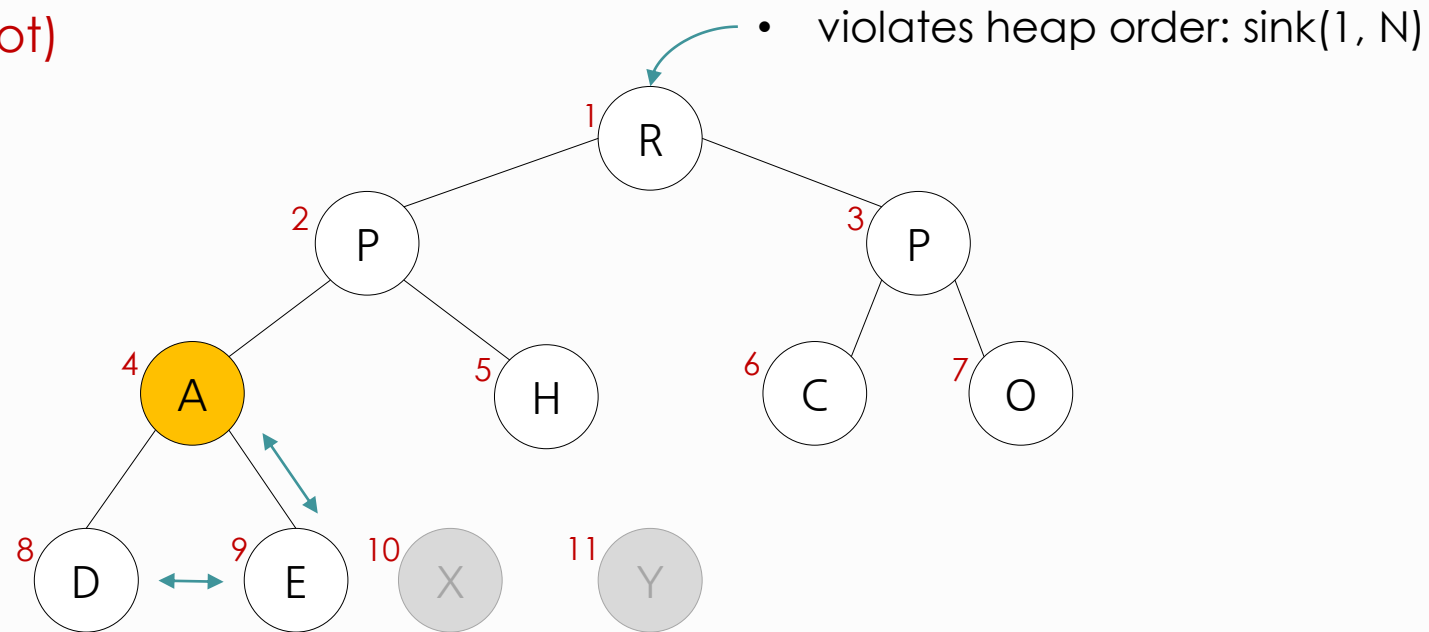
heap-ordered



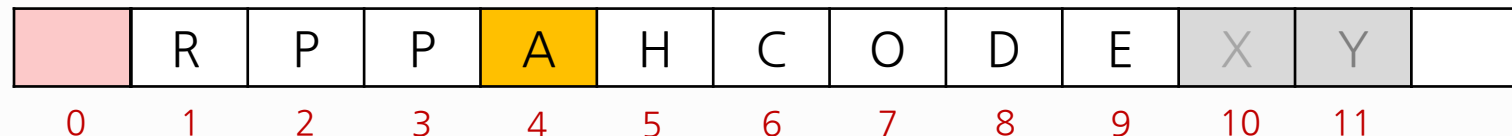
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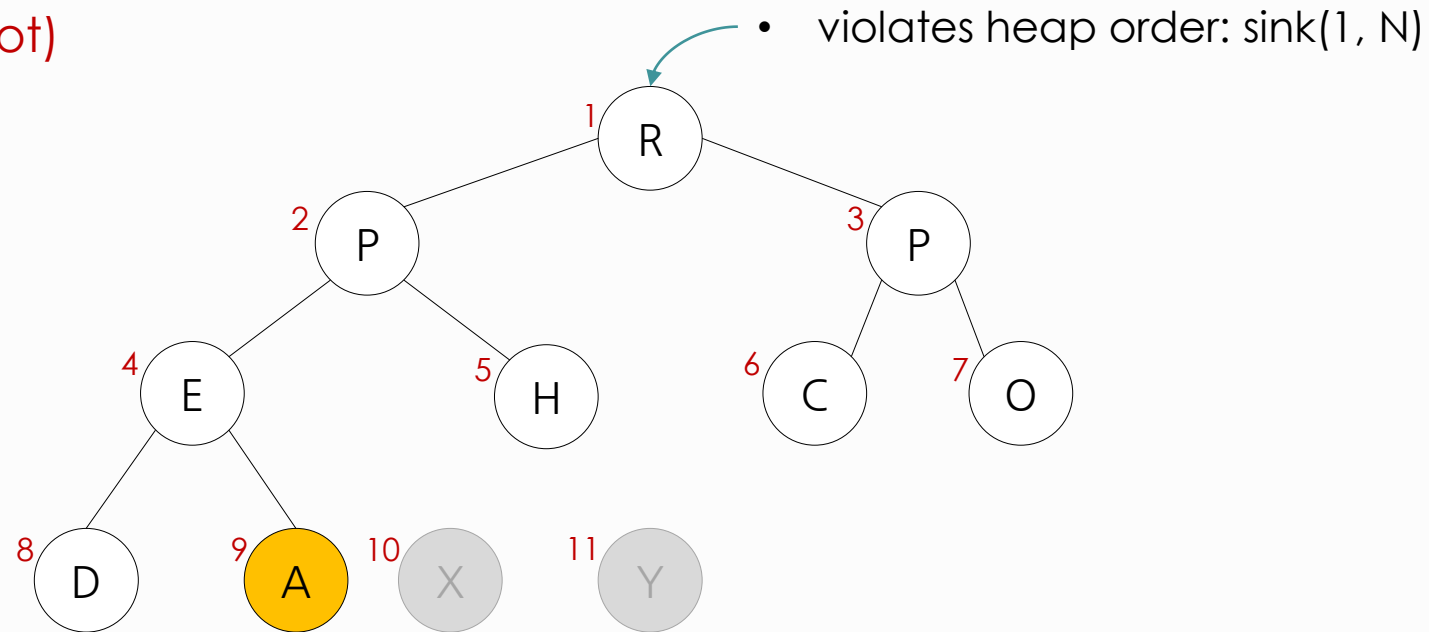
heap-ordered



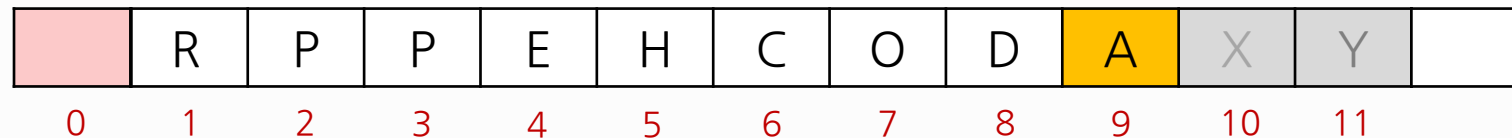
# max-heap example

- **Insert:** Add node at end, then swim it up.
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remove the max (root)



heap-ordered

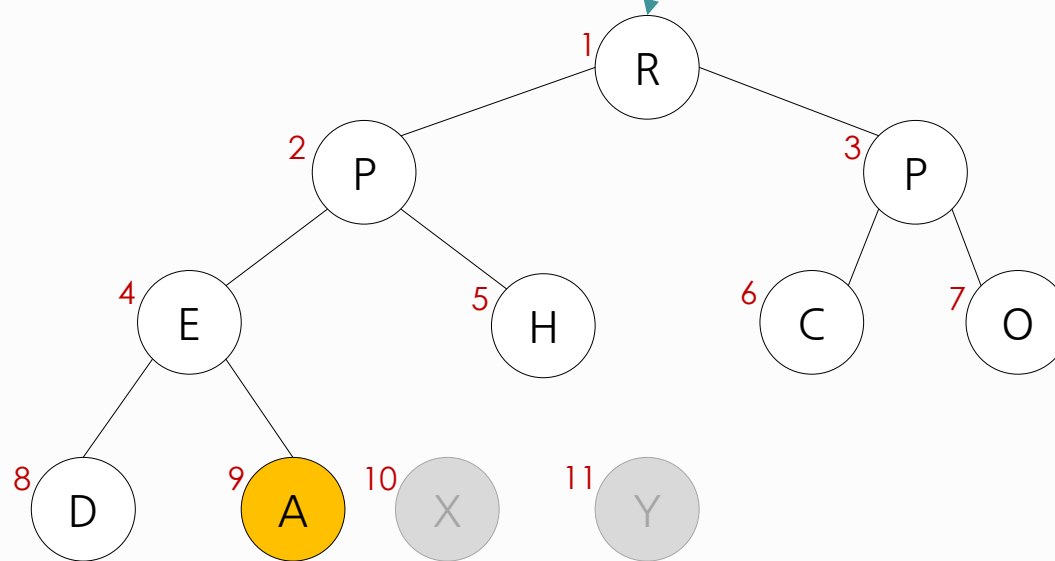


# max-heap example

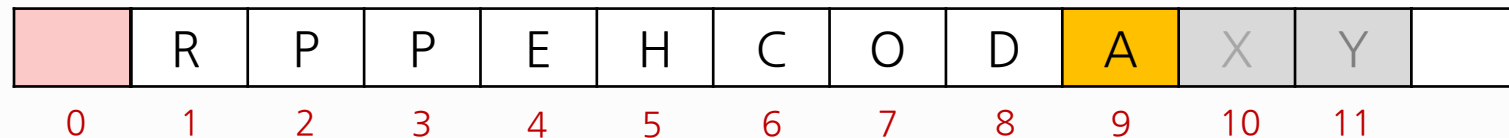
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remove the max (root)

• heap-ordered



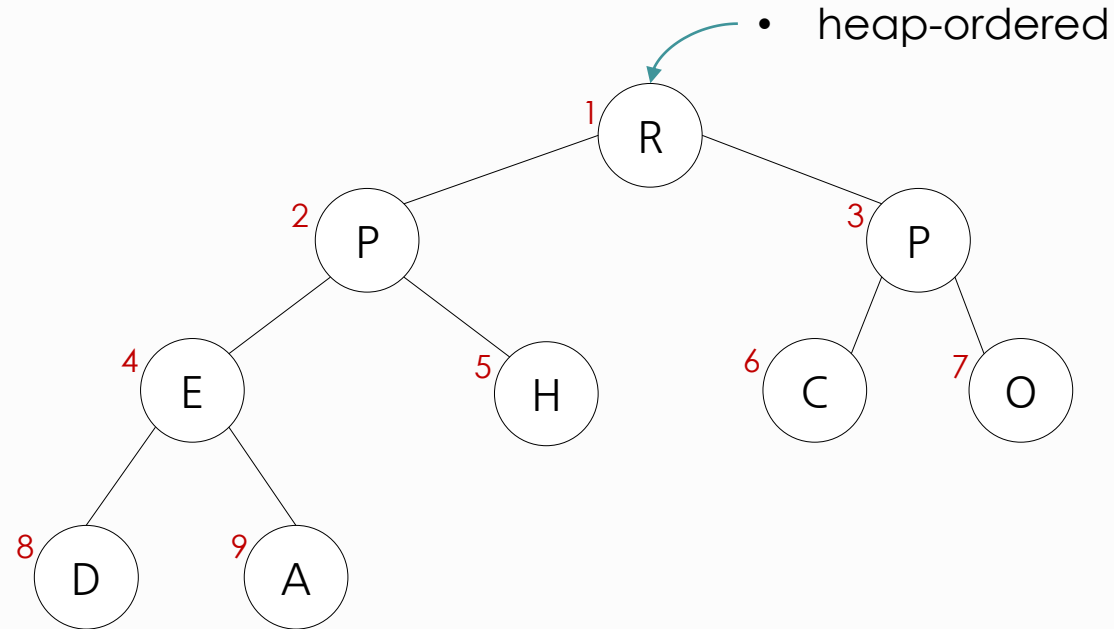
heap-ordered



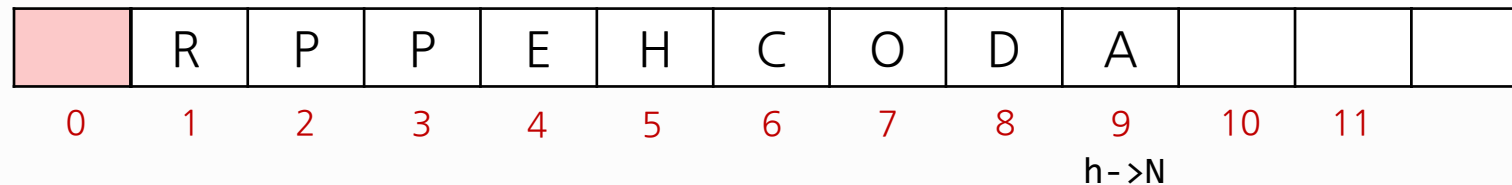
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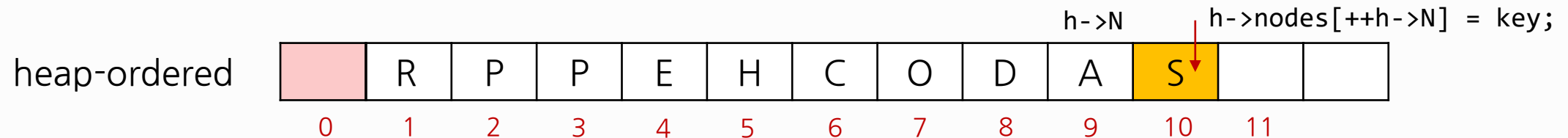
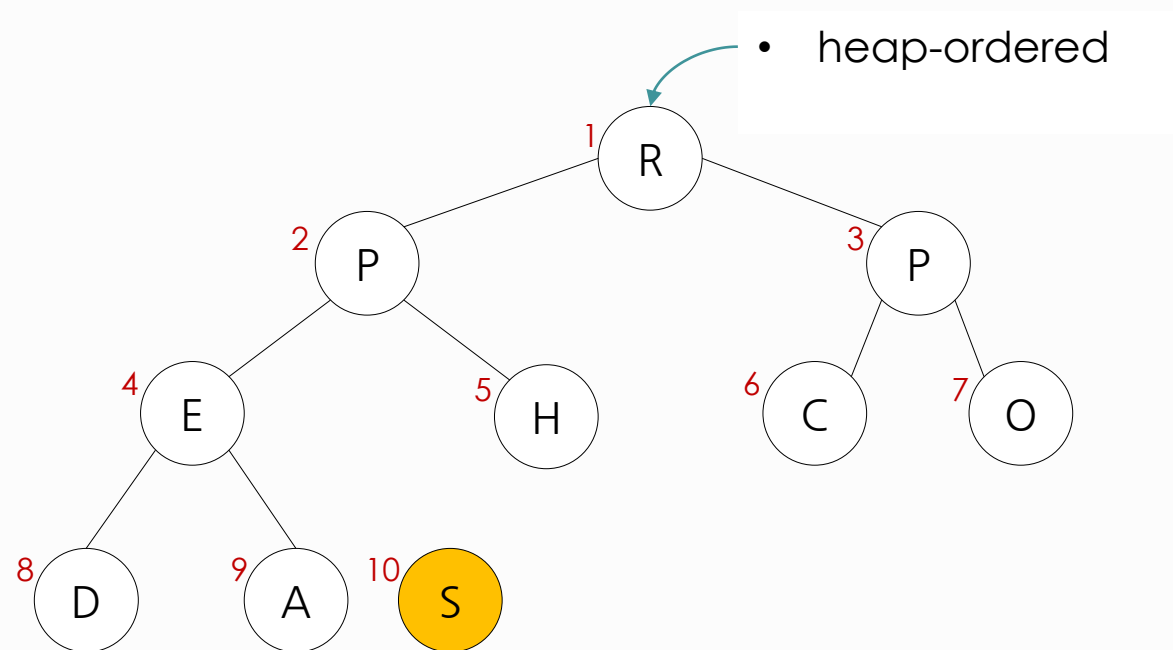
heap-ordered



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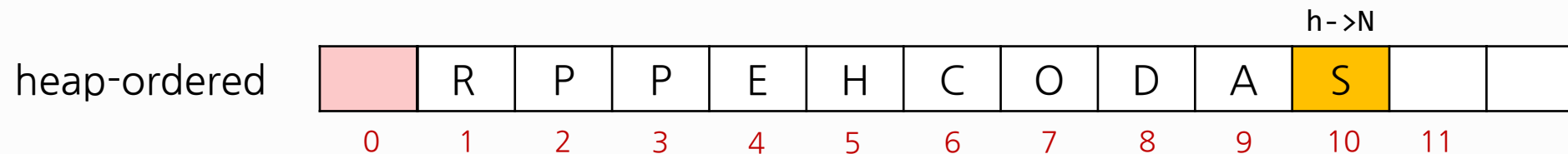
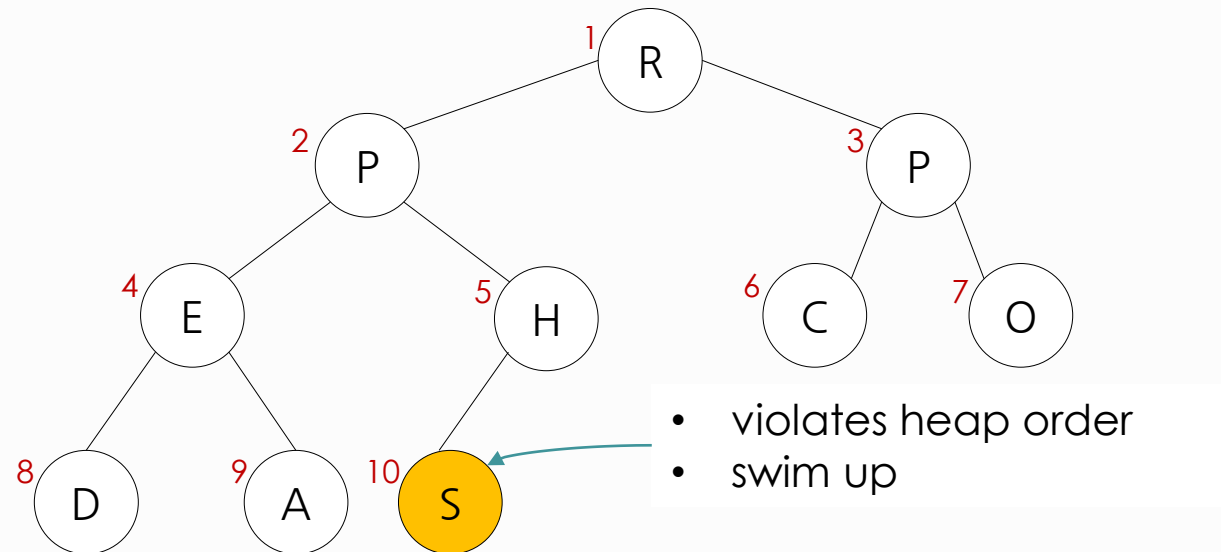




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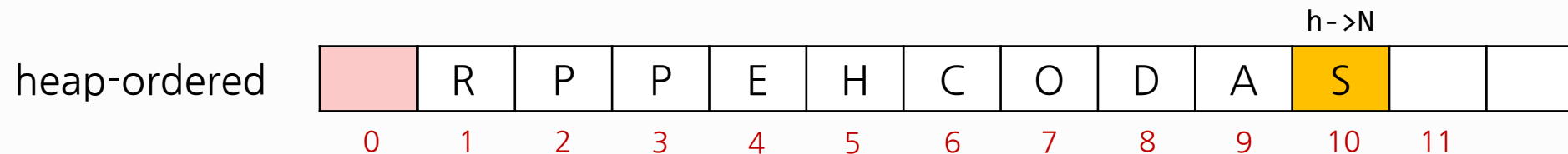
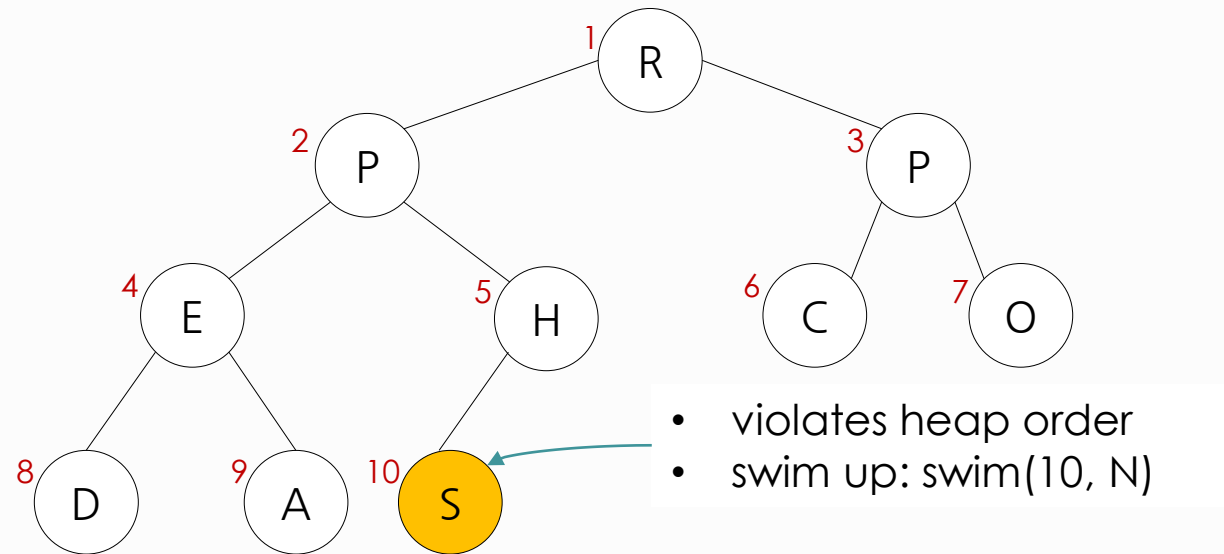
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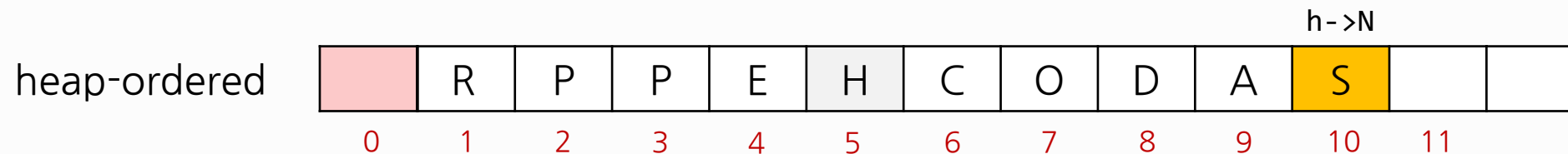
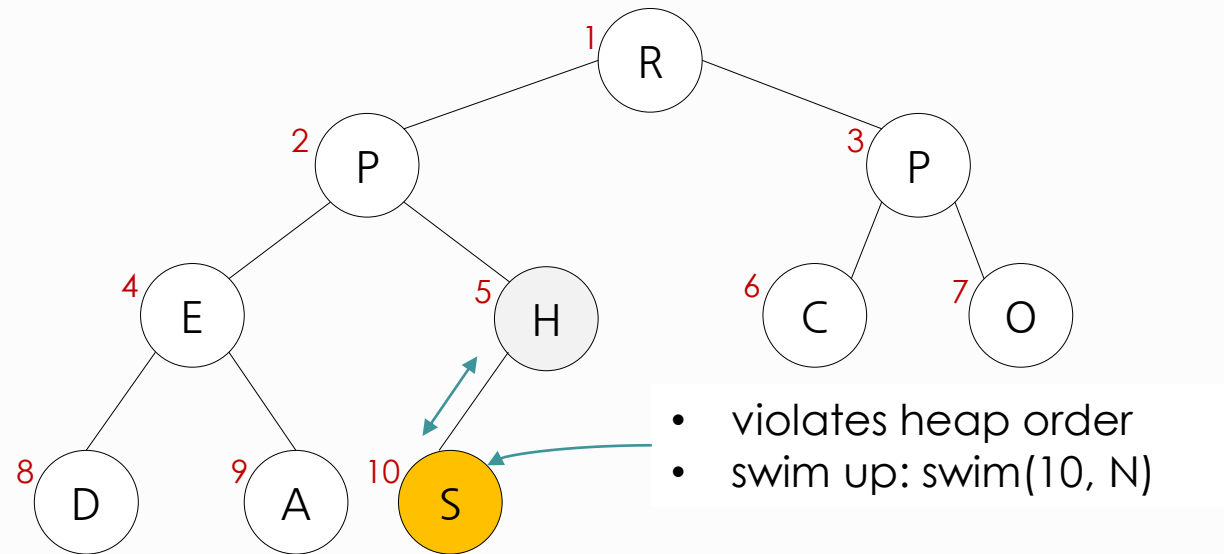
insert 



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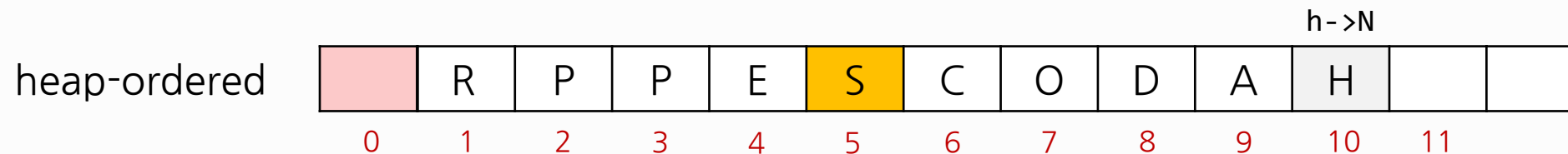
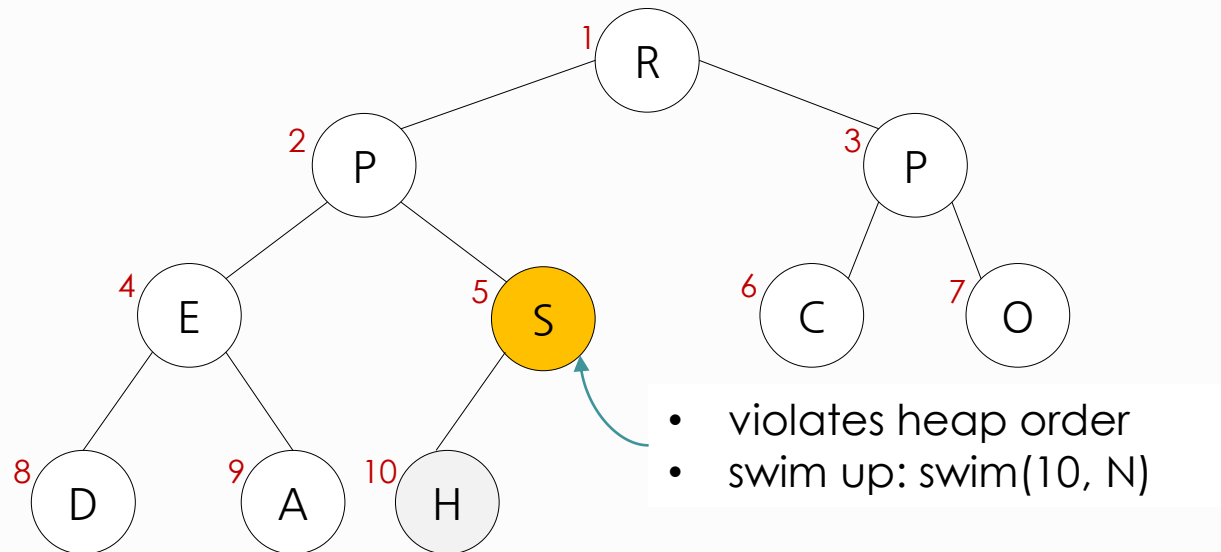
insert 



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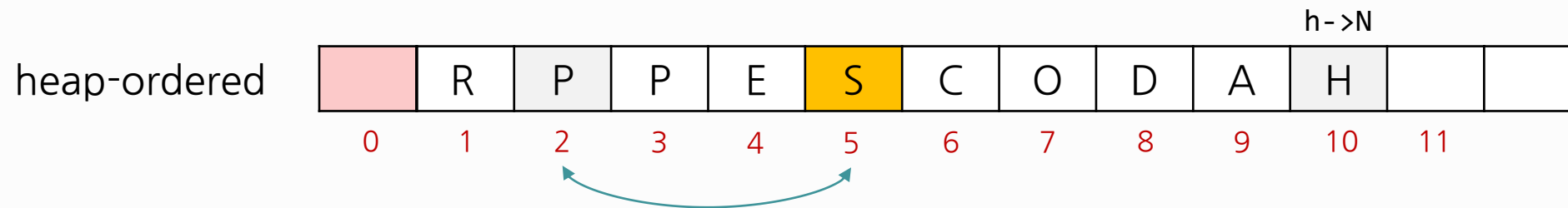
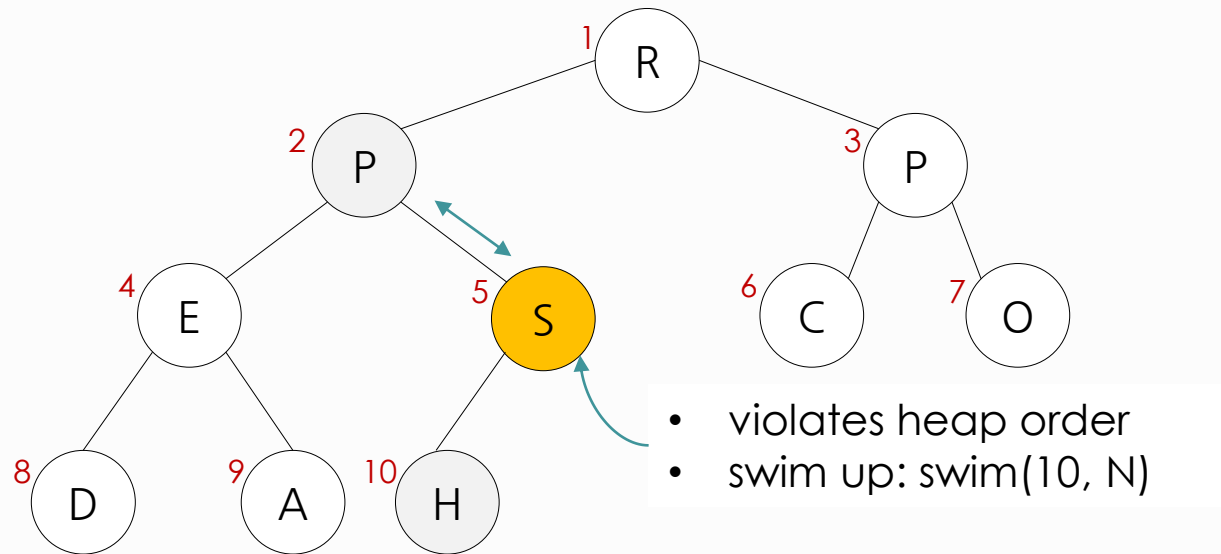
insert 



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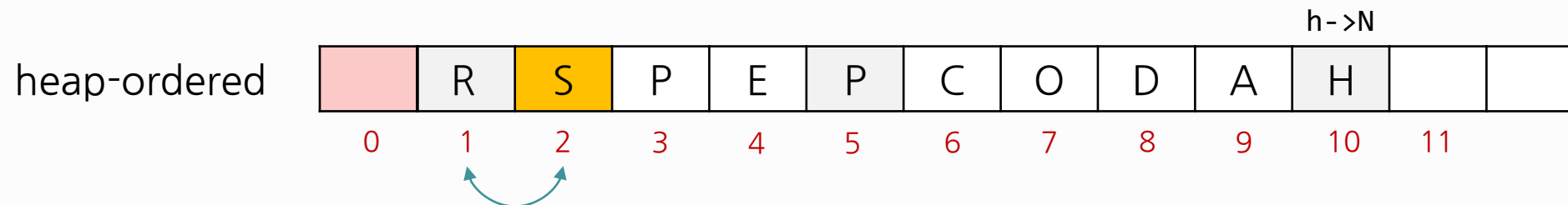
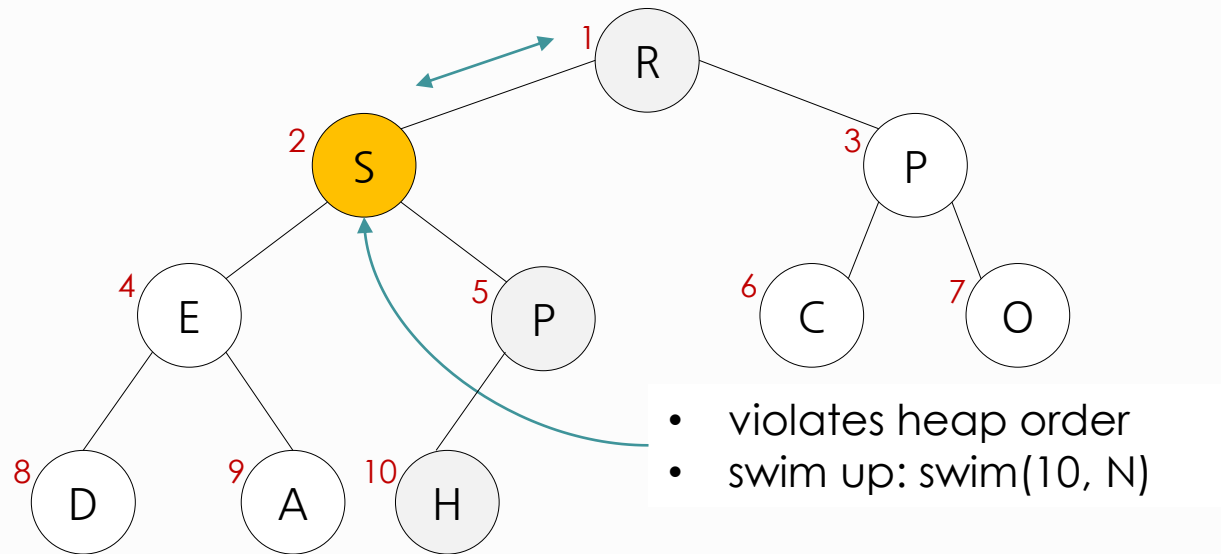
insert 



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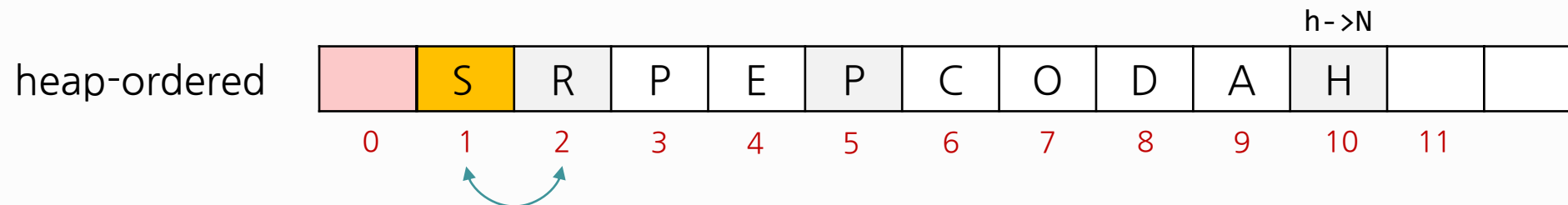
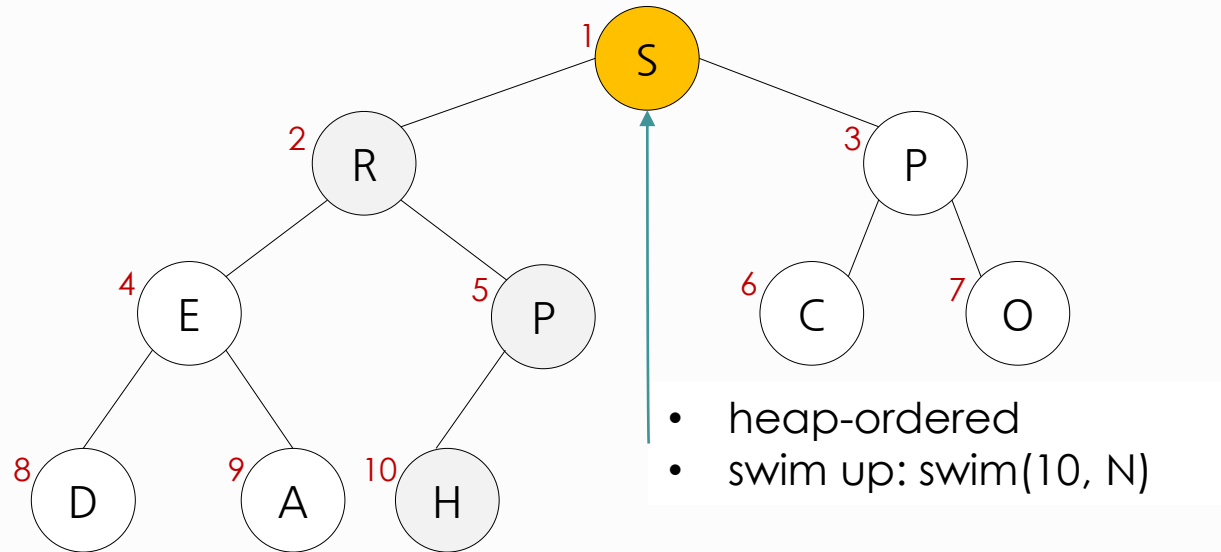
insert 



# max-heap example

- **Insert:** Add node at end, then swim it up.
- **Remove:** Swap root with last node, then sink down.

insert 



## Binary heap operations time complexity with N items:

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- Level of heap is  $\lfloor \log_2 N \rfloor$
- insert:  $O(\log N)$  for each insert
  - In practice, expect less
- delete:  $O(\log N)$  // deleting root node or any node
- increase/decrease key:  $O(\log N)$
- **Heapify():**  $O(N)$
- **Heapsort():**  $O(N \log N)$
- Because  $O(N)$  heapify +  $O(\log N)$  delete =  $O(N \log N)$
  
- **Proof:**
  - <https://stackoverflow.com/questions/9755721/how-can-building-a-heap-be-on-time-complexity>
  - <https://www.insertingwiththeweb.com/data-structures/binary-heap/build-heap-proof/>
  - <https://www.quora.com/How-is-the-time-complexity-of-building-a-heap-is-o-n>
- **References in Korean:**
  - <https://ratsgo.github.io/data%20structure&algorithm/2017/09/27/heapsort/>
  - <https://zeddios.tistory.com/56>