Data Structures Chapter 7: Graph

- 1. Introduction
 - Terminology, Representation, ADT
- 2. Basic Operations
 - DFS, CC, BFS, Processing
- 3. Digraph and Applications
- 4. Minimum Spanning Tree(MST)



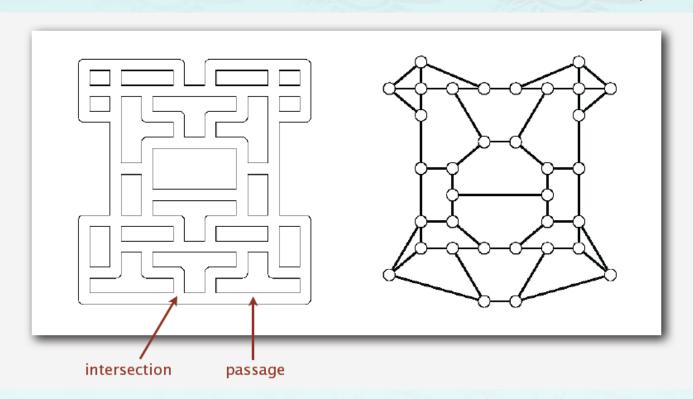
죄의 삯은 사망이요 하나님의 은사는 그리스도 예수 우리 주 안에 있는 영생이니라 (로마서 6:23)

모든 사람이 죄를 범하였으매 하나님의 영광에 이르지 못하더니 그리스도 예수 안에 있는 속량으로 말미암아 하나님의 은혜로 값없이 의롭다 하심을 얻은 자 되었느니라 (로마서 3:23-24)

Algorithm:

- Vertex = intersection
- Edge = passage

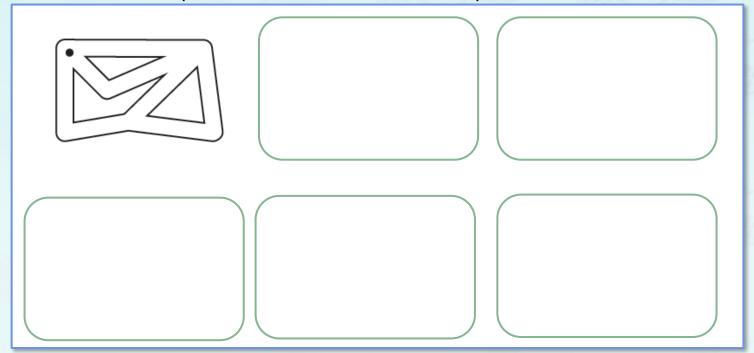
pacman



Maze Goal: Explore every intersection in the maze.

Maze graph:

- Unroll a ball of string behind you.
- Mark each visited intersection and each visited passage.
- Retrace steps when no unvisited options



Maze Goal: Explore every intersection in the maze.

Good Visualization: https://www.cs.usfca.edu/~galles/visualization/DFS.html

Maze graph:

- Unroll a ball of string behind you.
- Mark each visited intersection and each visited passage.
- Retrace steps when no unvisited options



Theseus, a hero of Greek mythology, is best known for slaying a monster called the Minotaur. When Theseus entered the Labyrinth where the Minotaur lived, he took a ball of <u>yarn</u> to unwind and mark his route. Once he found the Minotaur and killed it, Theseus used the string to find his way out of the maze.

Read more: http://www.mythencyclopedia.com/Sp-Tl/Theseus.html#ixzz30wFO3ofe

Maze Goal: Explore every intersection in the maze.

Maze graph:

- Unroll a ball of string behind you.
- Mark each visited intersection and each visited passage.
- Retrace steps when no unvisited options



Shannon and his famous <u>electromechanical</u> mouse Theseus (named after <u>Theseus</u> from Greek mythology) which he tried to have solve the maze in one of the first experiments in <u>artificial intelligence</u>.

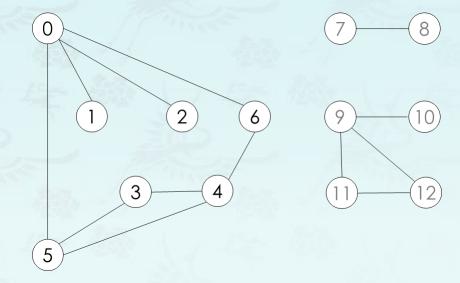
The Las Vegas connection: Shannon and his wife Betty also used to go on weekends to <u>Las Vegas</u> with <u>MIT</u> mathematician <u>Ed Thorp</u>, and made very successful forays in <u>blackjack</u> using <u>game theory</u>.

Maze Goal: Explore every intersection in the maze.

- Design pattern: Decouple graph data type
- Idea: Mimic maze exploration

DFS (to visit a vertex v)

- Mark v as visited.
- Recursively visit all unmarked vertices w adjacent to v.



Typical applications:

- Find all vertices connected to a given source vertex.
- Find a path between two vertices.

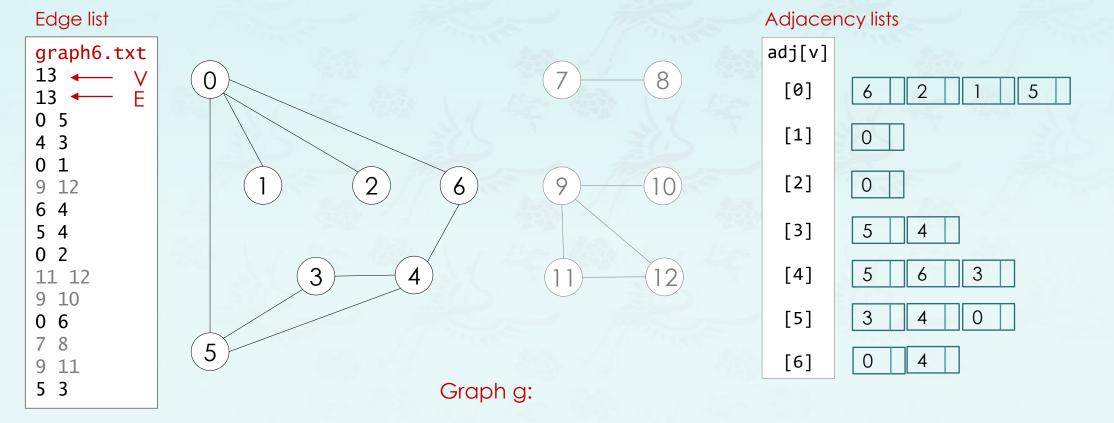
Challenge:

How to implement?

Goal: Systematically search through a graph from graph processing

- Create a graph object
- Pass the graph to a graph processing routine
- Query the graph-processing routine
 - path from v to w
 - distance from v to w
 - connected
 - bipartite
 - cyclic

- For each edge(v, w) in the list
- Insert front each vertex both (adj[v], w) and (adj[w], v) addEdgeFromTo(g, v, w); // from v to w.



Challenge: build adjacency lists?

Graph - ADT

```
// a structure to represent an adjacency list of vertices
struct Gnode {
  int    item;
    Gnode* next;
  Gnode (int i, Gnode *p = nullptr) {
    item = i;    next = p;
  }
    ~Gnode() {}
};
using gnode = Gnode *;
```

Graph ADT – graph.h

```
struct Graph {
 int V; // N vertices
 int E; // N edges
 gnode adj; // array of linked lists of vertices
 Graph(int v = 0) \{ // constructs a graph with v vertices \}
   V = V;
   E = 0;
   adj = new (nothrow) Gnode[v];
   assert(adj != nullptr);
   for (int i = 0; i < v; i++) { // initialize adj list as empty;
       adj[i].item = i;
                                 to begin with
                    unused:
                      but may store the degree of vertex i.
 ~Graph() {}
                      graph g = new Graph(v);
using graph = Graph *;
                      for (int i = 0; i < E; i++)
                        addEdge(g, from[i], to[i]);
```

Graph ADT – graph.cpp

```
// add an edge from v to w to an undirected graph
// A new vertex is added at the beginning of adj list of v.
void addEdgeFromTo(graph g, int v, int w) {
 gnode node = new Gnode(w);
 g->adj[v].next = node;
 g->E++;
   With a bug
// add an edge to an undirected graph
void addEdge(graph g, int v, int w) {
  addEdgeFromTo(g, v, w); // edge from v to w
 addEdgeFromTo(g, w, v);  // since undirected
```

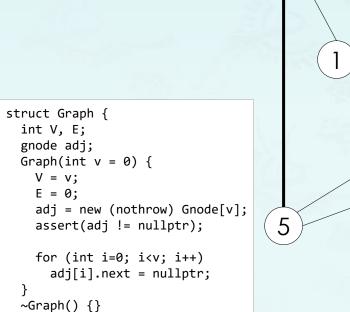
Graph ADT – graph.cpp

```
// add an edge from v to w to an undirected graph
// A new vertex is added at the beginning of adj list of v.
void addEdgeFromTo(graph g, int v, int w) {
                                                     instantiate a node w and
 gnode node = new Gnode(w);
                                                     make a link with vertex v.
 g->adj[v].next = node;
                                                     what is wrong?
 g->E++;
   With a bug
// add an edge to an undirected graph
void addEdge(graph g, int v, int w) {
  addEdgeFromTo(g, v, w); // edge from v to w
 addEdgeFromTo(g, w, v);  // since undirected
```

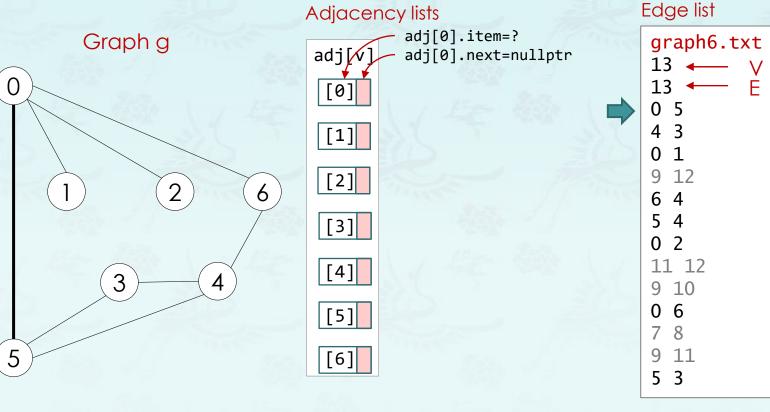
Graph ADT – graph.cpp

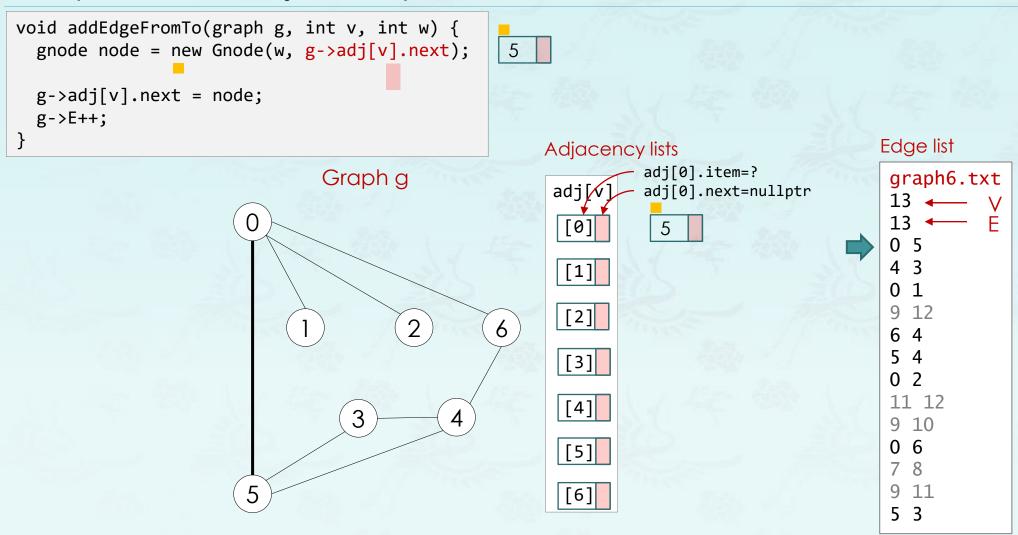
```
// add an edge from v to w to an undirected graph
// A new vertex is added at the beginning of adj list of v.
void addEdgeFromTo(graph g, int v, int w) {
                                                     instantiate a node w and
 gnode node = new Gnode(w, g->adj[v].next);
                                                     insert it at the front of adj[v]
 g->adj[v].next = node;
 g->E++;
// add an edge to an undirected graph
void addEdge(graph g, int v, int w) {
 addEdgeFromTo(g, v, w); // edge from v to w
 addEdgeFromTo(g, w, v);  // since undirected
```

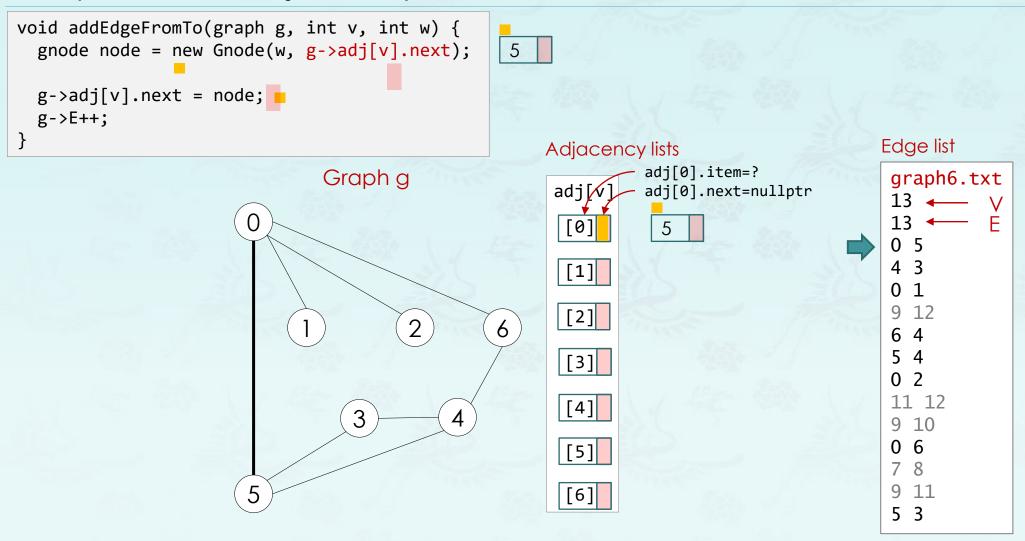
```
void addEdgeFromTo(graph g, int v, int w) {
  gnode node = new Gnode(w, g->adj[v].next);
  g->adj[v].next = node;
  g->E++;
}
```



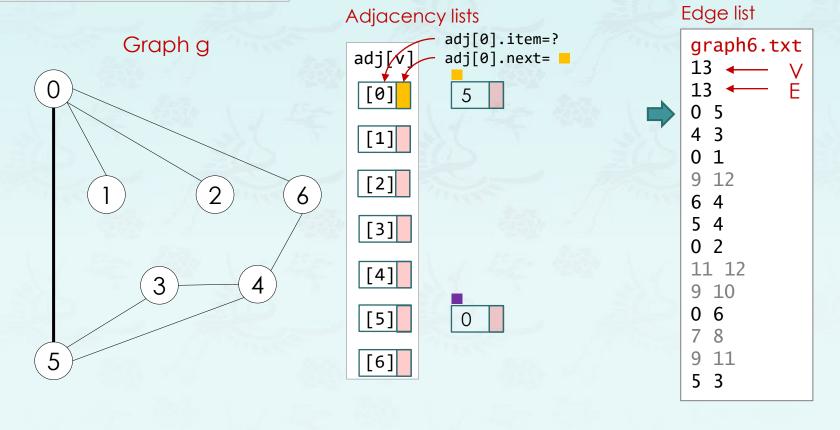
using graph = Graph *;



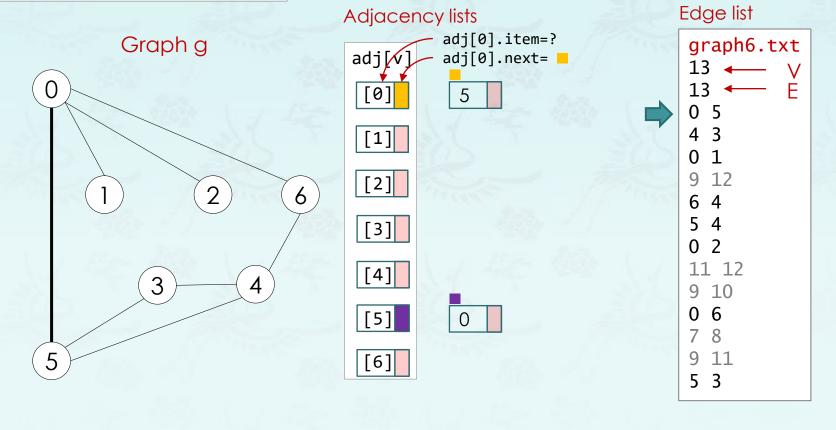




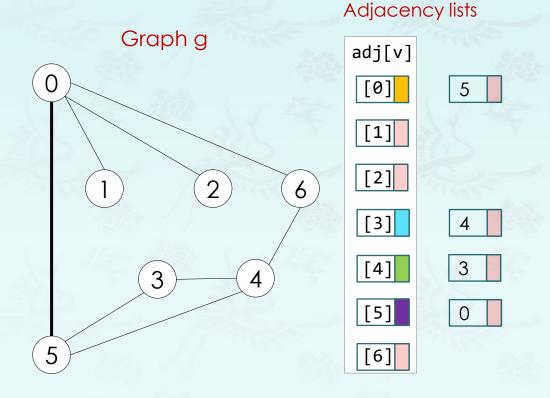
```
void addEdgeFromTo(graph g, int v, int w) {
  gnode node = new Gnode(w, g->adj[v].next);
  g->adj[v].next = node;
  g->E++;
}
```

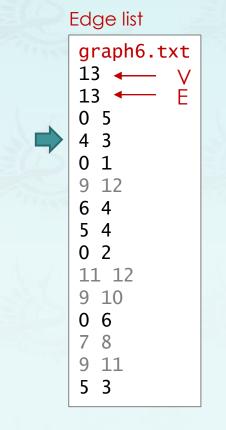


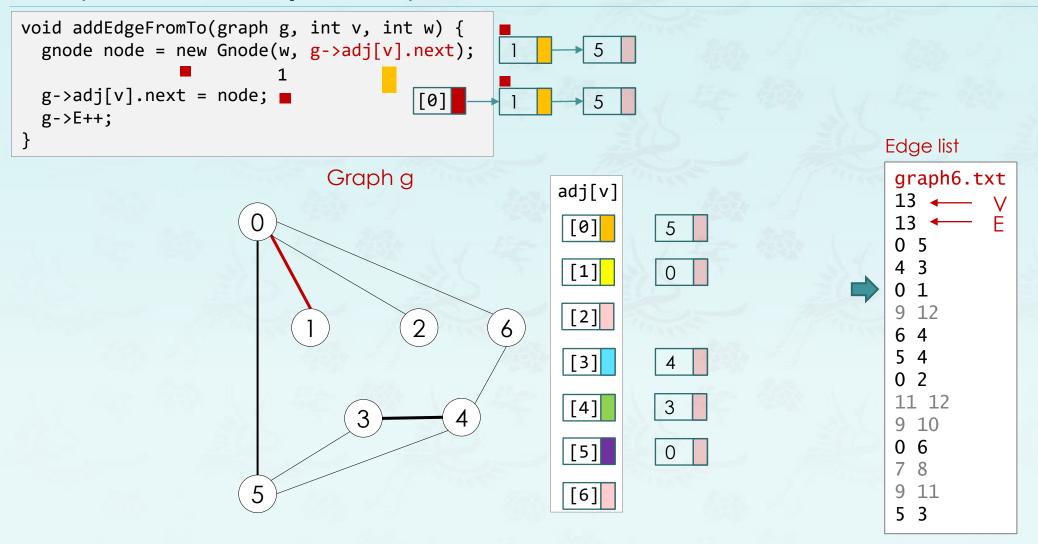
```
void addEdgeFromTo(graph g, int v, int w) {
  gnode node = new Gnode(w, g->adj[v].next);
  g->adj[v].next = node;
  g->E++;
}
```



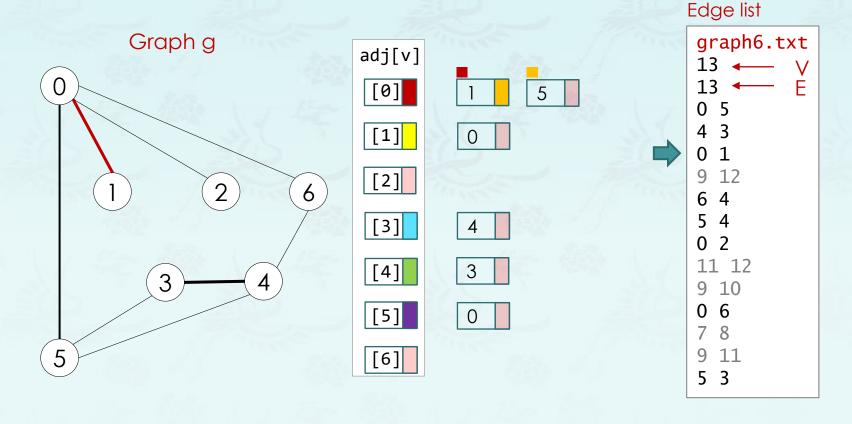
```
void addEdgeFromTo(graph g, int v, int w) {
  gnode node = new Gnode(w, g->adj[v].next);
  g->adj[v].next = node;
  g->E++;
}
```



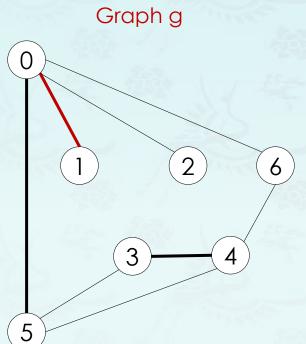




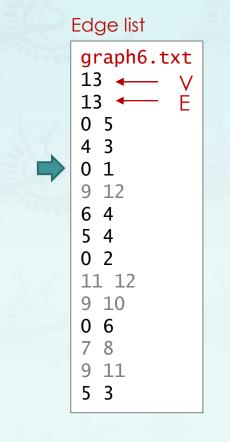
```
void addEdgeFromTo(graph g, int v, int w) {
  gnode node = new Gnode(w, g->adj[v].next);
  g->adj[v].next = node;
  g->E++;
}
```



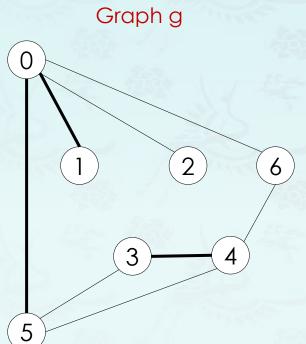
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void addEdgeFromTo(graph g, int v, int w) {
  gnode node = new Gnode(w, g->adj[v].next);
  g->adj[v].next = node;
  g->E++;
}
```



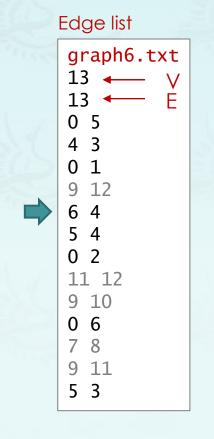
Adjacency lists adj[v] [0] [1] 0 [2] [3] 4 [4] 3 [5] 0 [6]



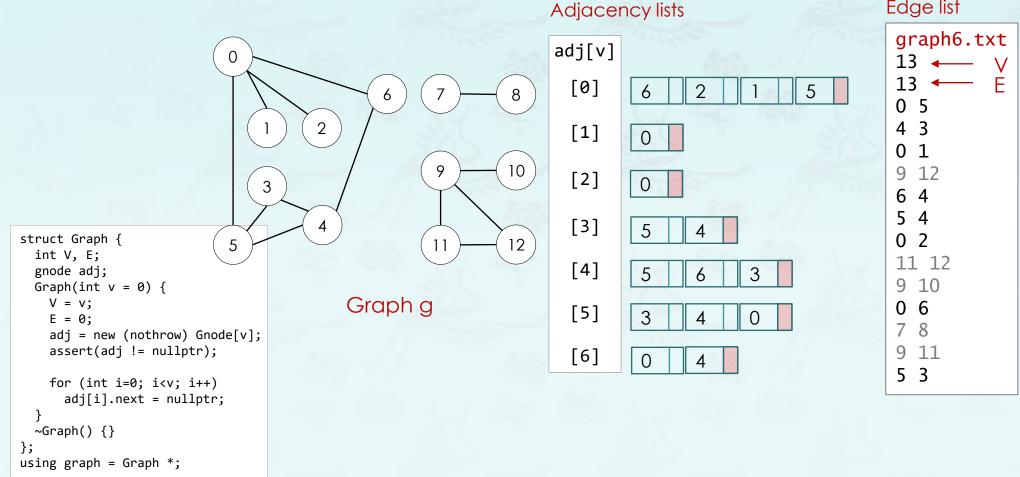
```
void addEdgeFromTo(graph g, int v, int w) {
  gnode node = new Gnode(w, g->adj[v].next);
  g->adj[v].next = node;
  g->E++;
}
```



Adjacency lists adj[v] [0] [1] 0 [2] [3] 4 [4] 3 [5] 0 [6]



Warming-up: degree()



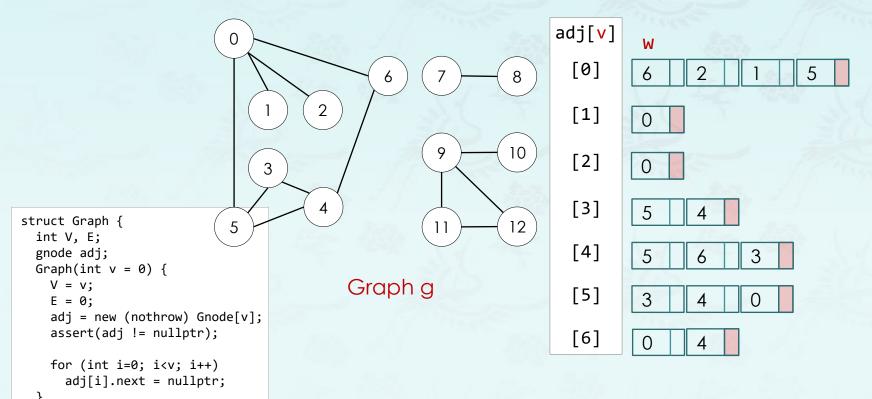
Edge list

Warming-up: degree()

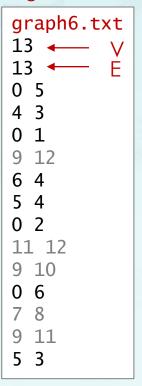
~Graph() {}

using graph = Graph *;

```
int degree(graph g, int v) {
  int deg = 0;
  for (gnode w = g->adj[v].next; w != nullptr; w = w->next)
    deg++;
  return deg;
}
```



Edge list

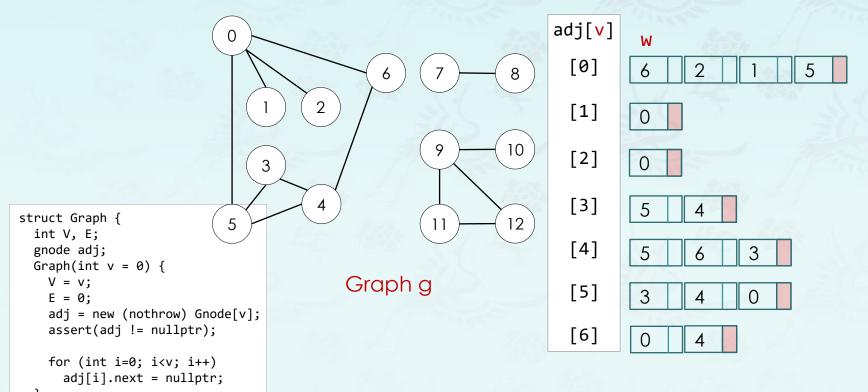


Warming-up: degree()

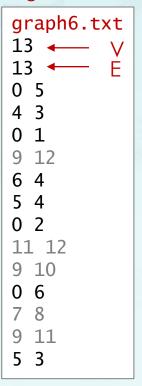
~Graph() {}

using graph = Graph *;

```
int degree(graph g) {
  int deg = 0;
  for (int v = 0; v < V(g); v++)
    deg = max(degree(g, v), deg);
  return deg;
}</pre>
```

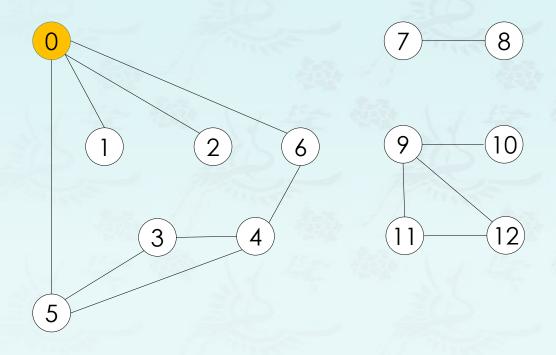


Edge list



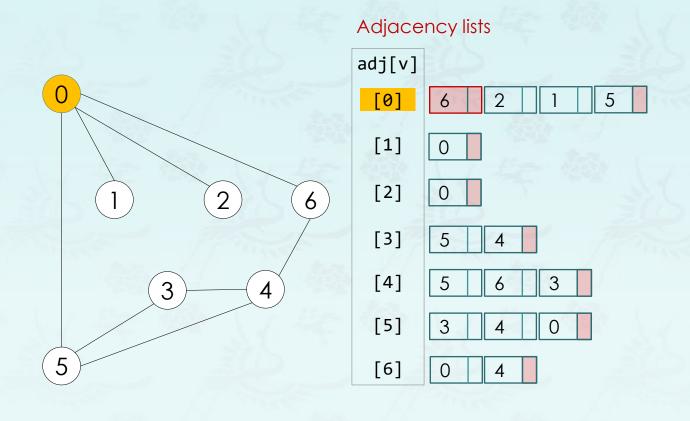
To visit a vertex v:

- Mark vertex v as visited.
- Recursively visit all unmarked vertices adjacent to v.



V	marked[]	parent[v]
0	Т	-
1	F	-
2	F	-
3	F	-
4	F	-
5	F	-
6	F	-
7	F	-
8	F	-
9	F	-
16) F	-
11	L F	-
12	2 F	-

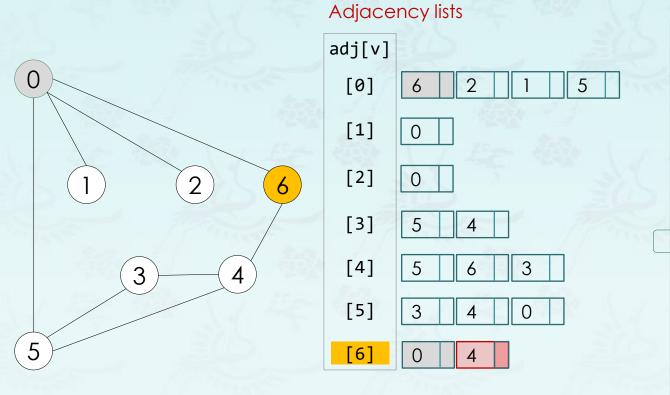
Which one first? visit 0:

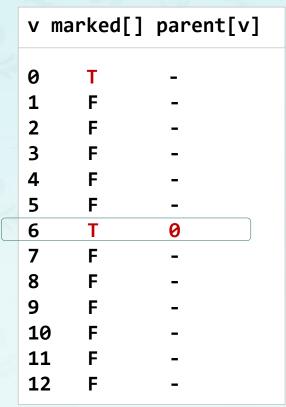


	V	marked[]	parent[v]	
	0	т	-	
	1	F	-	
	2	F	-	
	3	F	-	
	4	F	-	
	5	F	-	
	6	F	_	
	7	F	_	
	8	F	-	
	9	F	-	
	16) F	-	
	11	L F	-	
	12	2 F	-	
L				

visit 0: check 6, check 2, check 1, check 5

DFS 0





visit 6: check 0, check 4

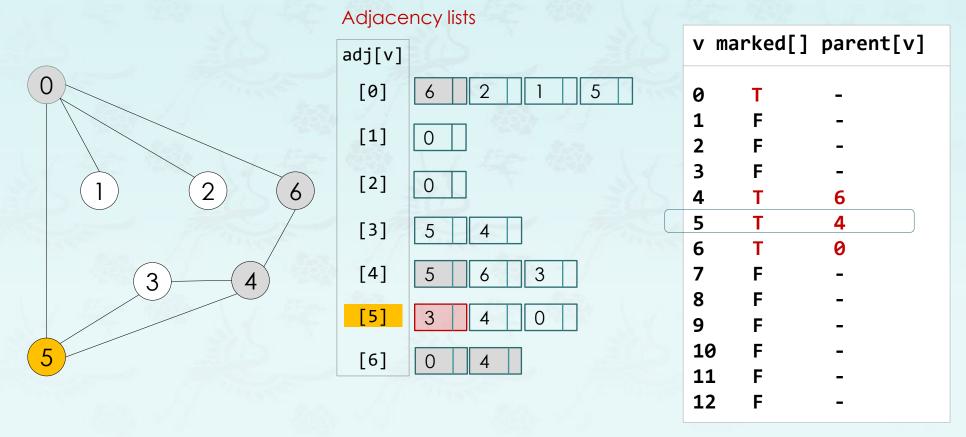
DFS 0 6



6

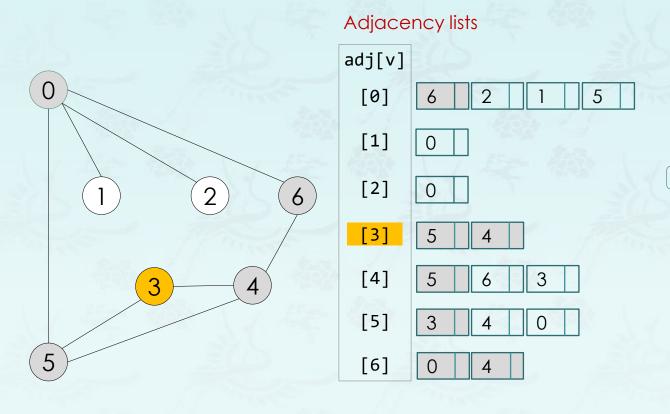
visit 4: check 5, check 6, check 3

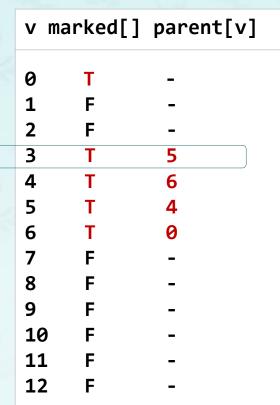
DFS 0 6 4



visit 5: check 3, check 4, check 0

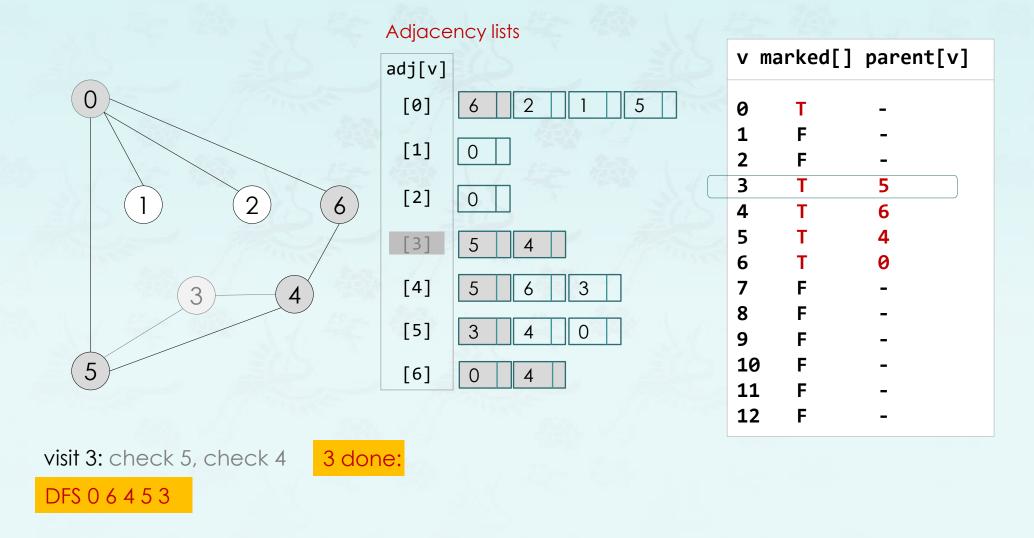
DFS 0 6 4 5

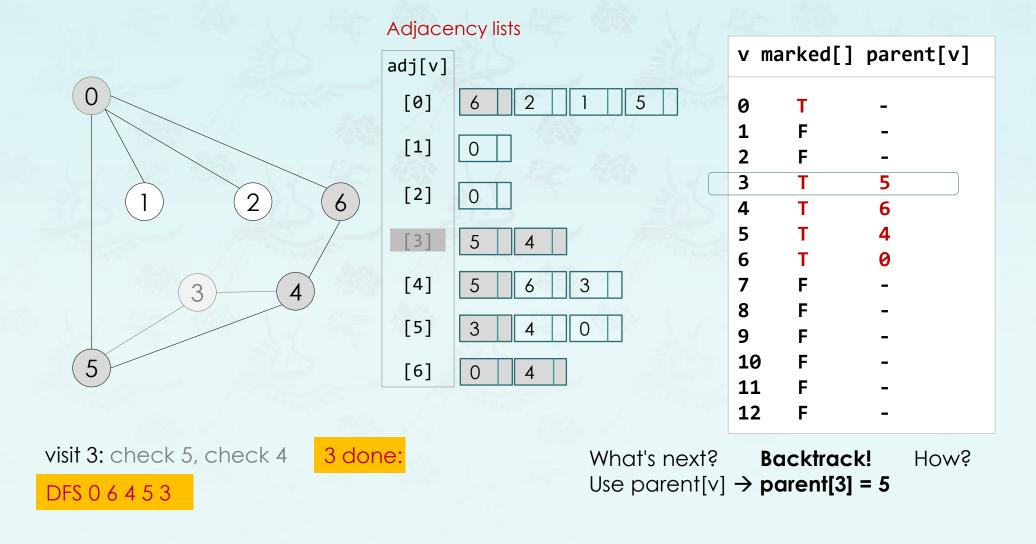




visit 3: check 5, check 4

DFS 0 6 4 5 3

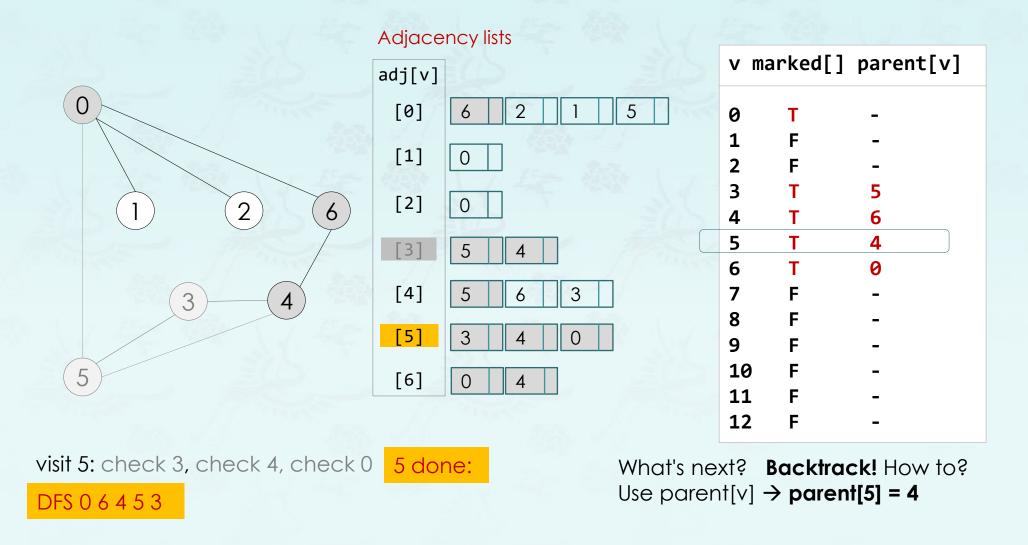


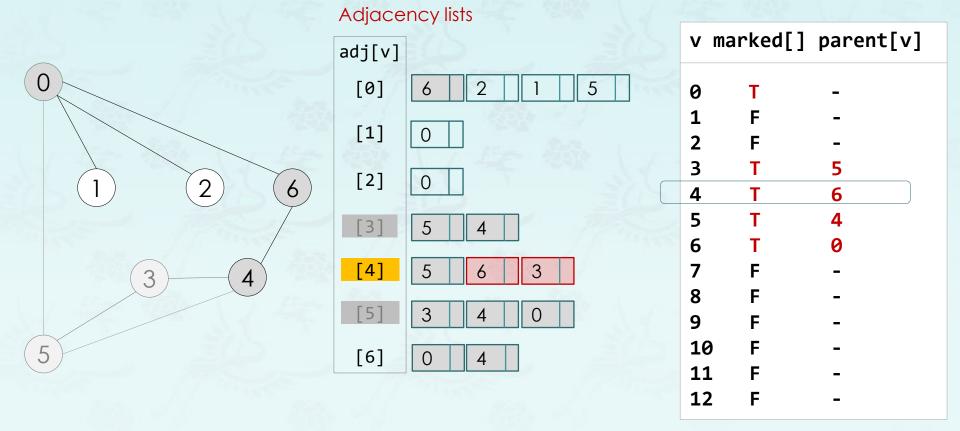




visit 5: check 3, check 4, check 0

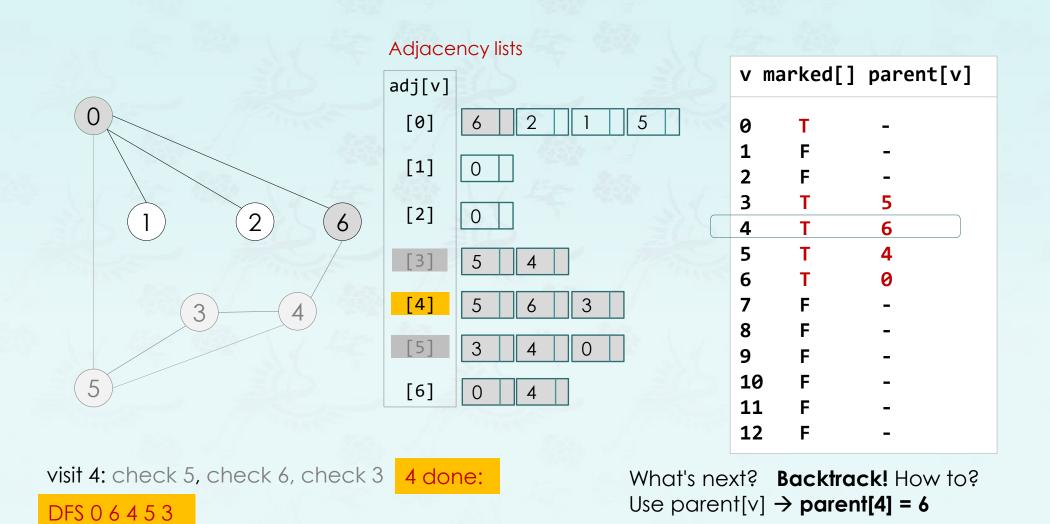
DFS 0 6 4 5 3

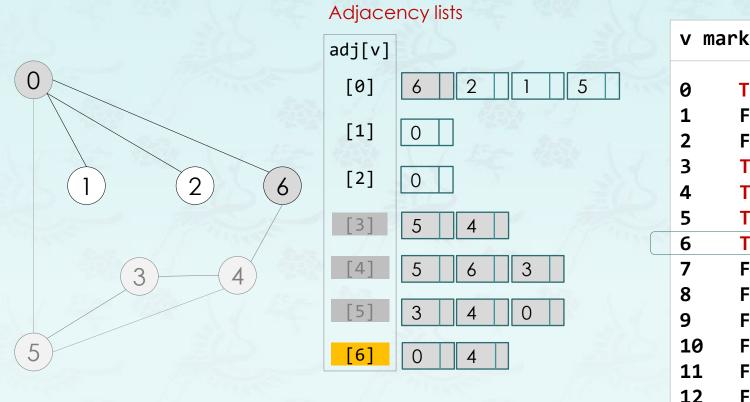




visit 4: check 5, check 6, check 3

DFS 0 6 4 5 3

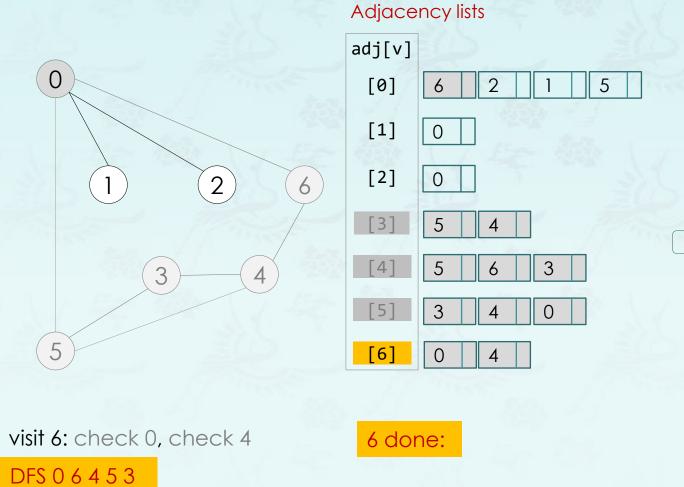






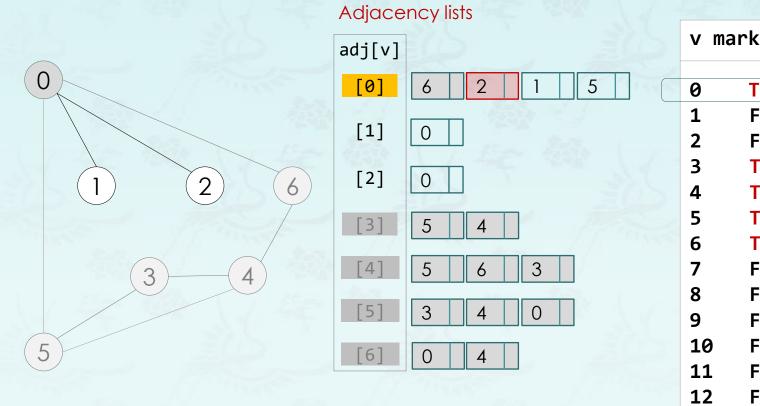
visit 6: check 0, check 4

DFS 0 6 4 5 3





What's next? Backtrack! parent[6] = 0

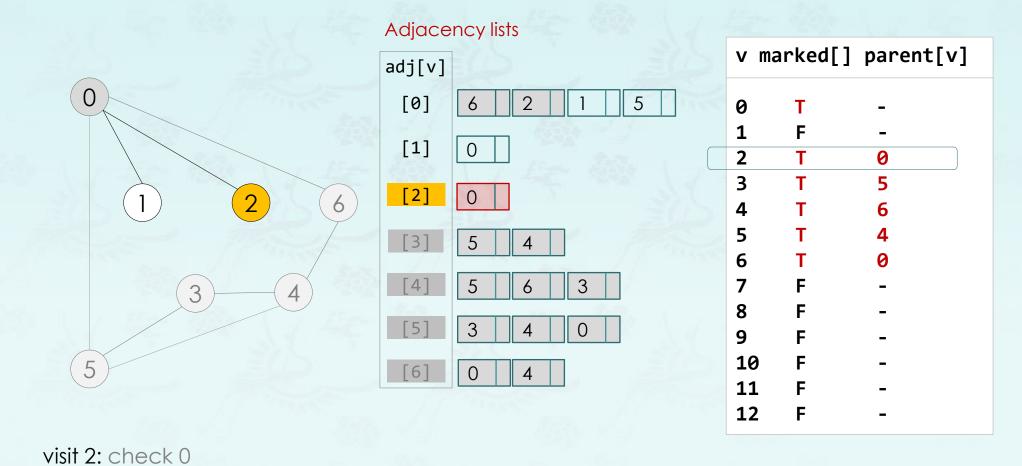


V	marked[]	<pre>parent[v]</pre>
0	Т	-
1	F	-
2	F	-
3	T	5
4	T	6
5	T	4
6	T	0
7	F	-
8	F	-
9	F	-
16		-
11		-
12	2 F	-

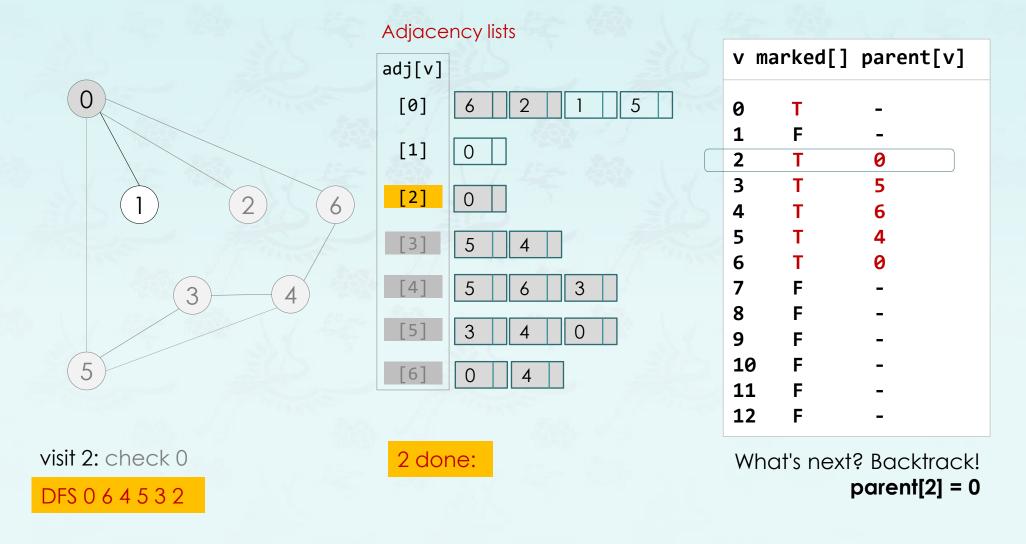
visit 0: check 6, check 2, check 1, check 5

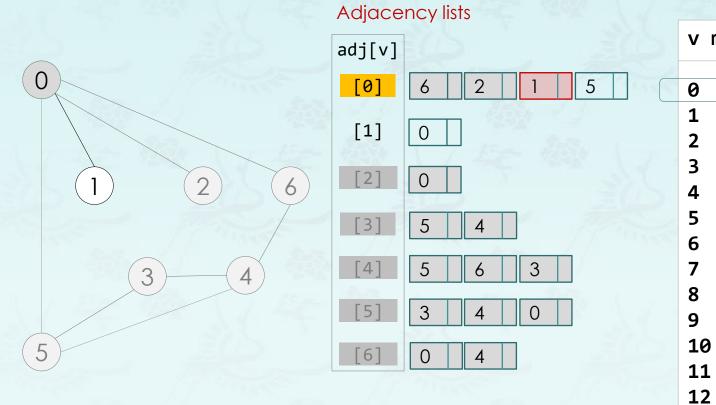
DFS 0 6 4 5 3

DFS 0 6 4 5 3 2



43

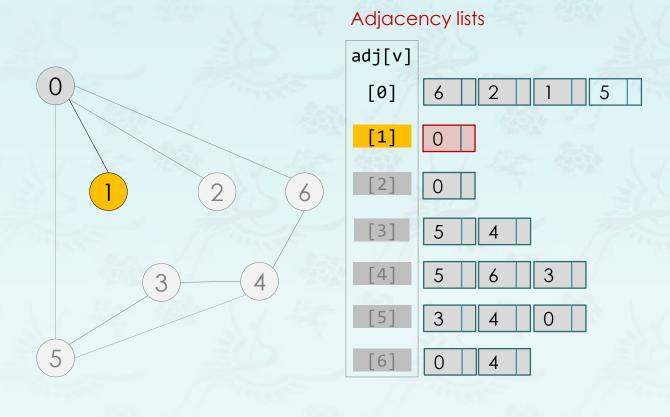


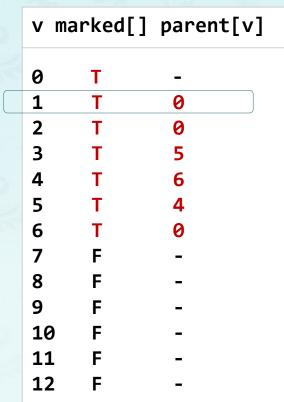


V	marked[]	<pre>parent[v]</pre>	
0	T	-	
1	F	-	
2	T	0	
3	T	5	
4	Т	6	
5	Т	4	
6	Т	0	
7	F	-	
8	F	-	
9	F	-	
16) F	-	
11	L F	-	
12	2 F	-	

visit 0: check 6, check 2, check 1, check 5

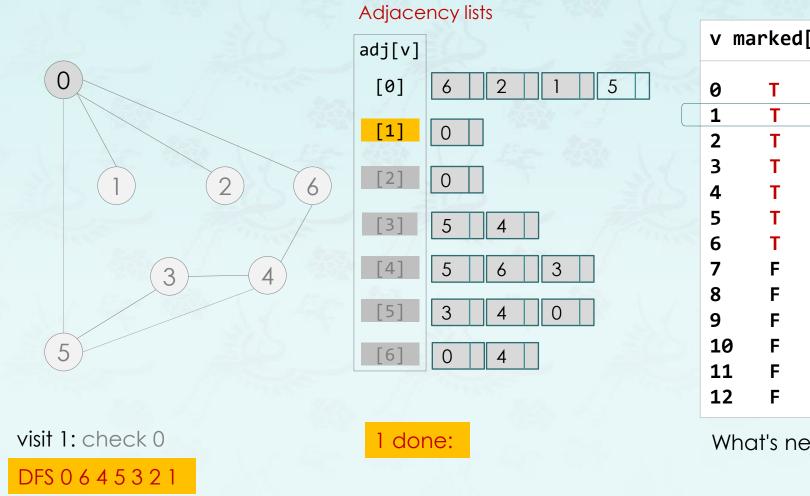
DFS 0 6 4 5 3 2





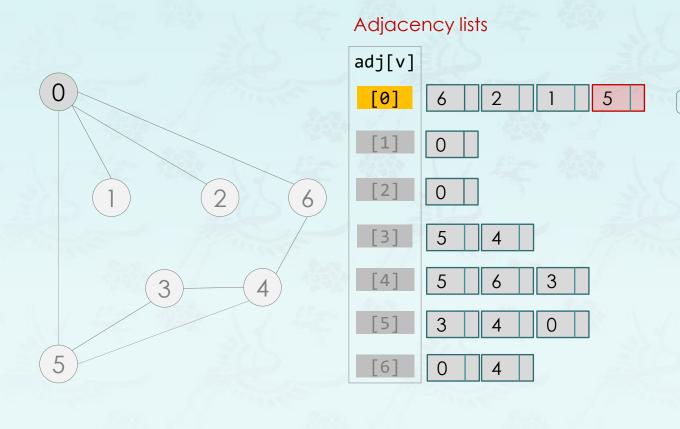
visit 1: check 0

DFS 0 6 4 5 3 2 1



V	marked[]	<pre>parent[v]</pre>
0	Т	-
1	Т	0
2	T	0
3	T	5
4	T	6
5	T	4
6	T	0
7	F	-
8	F	-
9	F	-
16	9 F	-
1:	l F	-
12	2 F	-

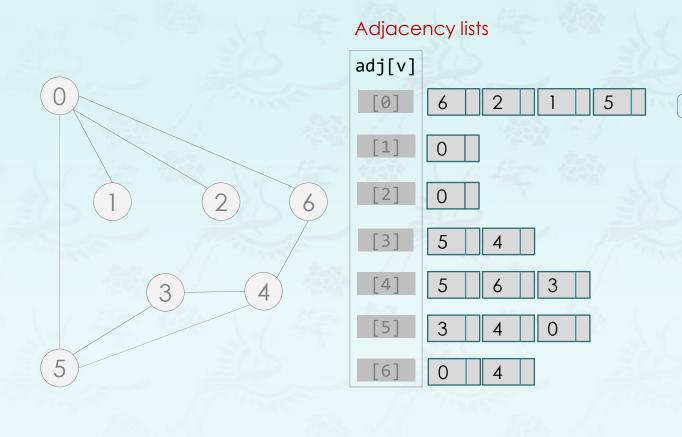
What's next? Backtrack! parent[1] = 0



V	marked[]	parent[v]	
0	T	_	
1	Т	0	
2	Т	0	
3	Т	5	
4	Т	6	
5	Т	4	
6	T	0	
7	F	_	
8	F	-	
9	F	-	
16) F	-	
11	. F	-	
12	2 F	-	

visit 1: check 6, check 2, check 1, check 5

DFS 0 6 4 5 3 2 1



V	marked[]	<pre>parent[v]</pre>	
0	Т	-	
1	Т	0	
2	T	0	
3	T	5	
4	T	6	
5	T	4	
6	T	0	
7	F	-	
8	F	-	
9	F	-	
16) F	-	
11	L F	-	
12	2 F	-	

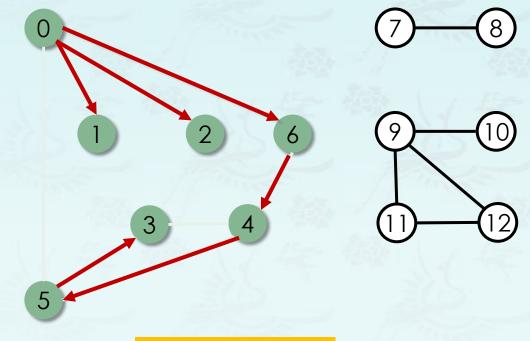
visit 1: check 6, check 2, check 1, check 5

0 done:

DFS 0 6 4 5 3 2 1

To visit a vertex v:

- Mark vertex v as visited.
- Recursively visit all unmarked vertices adjacent to v.



DFS Output:	DFS: 0 6 4 5 3 2 1
--------------------	--------------------

- found vertices reachable from 0
- build a data structure parent[v]

V	marked[]	<pre>parent[v]</pre>	
0	T	-	
1	Т	0	
2	Т	0	
3	T	5	
4	T	6	
5	T	4	
6	Т	0	
7	F	-	
8	F	-	
9	F	-	
16) F	-	
11	L F	_	
12	2 F	_	

Goal: Find all vertices connected to s (and a corresponding path).

Idea: Mimic maze exploration

Algorithm:

- Use recursion (ball of string).
- Mark each visited vertex (and keep track of edge taken to visit it).
- Return (retrace steps) when no unvisited options.

Data Structures:

- Boolean[] marked to mark visited vetices.
- int[] parent to keep tree of paths.
 (parent[w] == v) means that edge v-w taken to visit w for first time

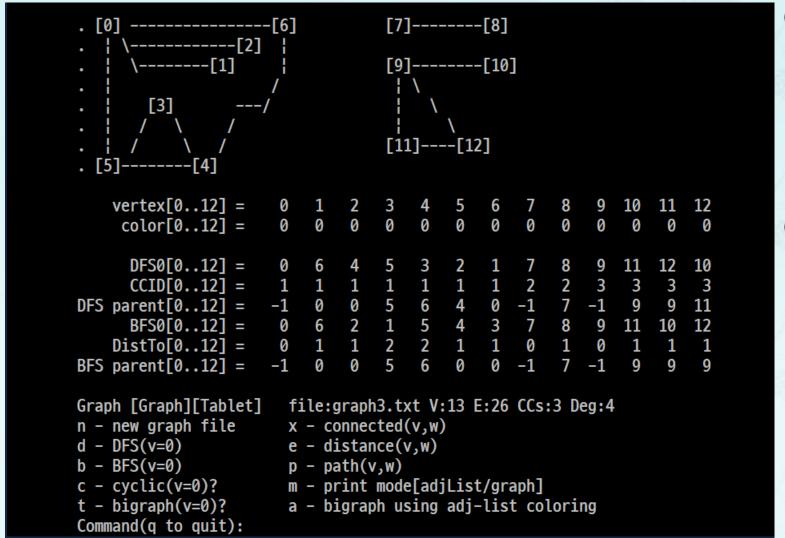
DFS: Depth-First Search Coding

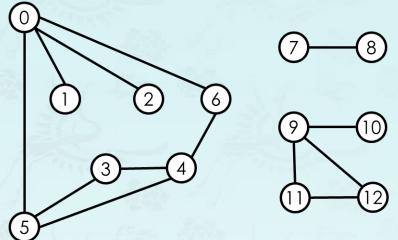
Adjacency lists adj[] 0 1 3 4 6

DFS: Depth-First Search Coding

```
// runs DFS for all components and produces DFS0[], CCID[] & parentDFS[].
void DFS CCs(graph g) {
  if (empty(g)) return;
  for (int i = 0; i < V(g); i++) {
    g->marked[i] = false;
    g->parentDFS[i] = -1;
    g \rightarrow CCID[i] = 0;
                              It works only for the 1st CC at 0.
                              Make it work for all CC's.
  queue<int> que;
                              Need for-loop to do it for all v's.
  DFS(g, 0, que);
  setDFS0(g, 0, que);
                              // set results into DFS0[] and CCID[]
 g \rightarrow DFSv = \{\};
                                     // runs DFS for at vertex v recursively.
                                     // Only que, g->marked[v] and g->parentDFS[] are updated here.
                                     void DFS(graph g, int v, queue<int>& que) {
                                       g->marked[v] = true;  // visited
                                       que.push(v);
                                                                          // save the path
                                       for (gnode w = g->adj[v].next; w; w = w->next) {
                                         cout << "your code here (recursion) \n";</pre>
```

DFS: Depth-First Search Coding





DFS: Depth-First Search Properties

- Proposition: After DFS, can find vertices connected to s in constant time and can find a path to s (if one exists) in time proportional to its length.
- Proof: parent[] is parent-link representation of a tree rooted at s.

٧	marked[]	<pre>parent[v]</pre>
0	Т	-
1	T	0
2	Т	0
3	Т	5
4	Т	6
5	Т	4
6	Т	0
7	F	-
8	F	-
9	F	-
10) F	-
1:	l F	-
12	2 F	-

DFS: Depth-First Search Properties

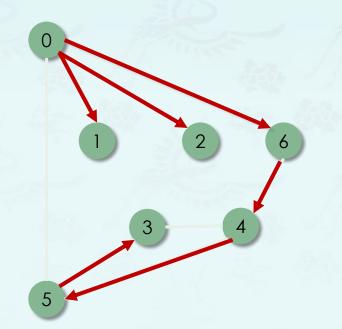
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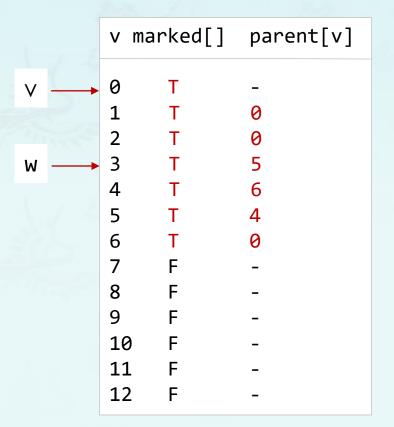
```
// returns a path from v to w using the DFS result or parentDFS[].
// It has to use a stack to retrace the path back to the source.
// Once the client(caller) gets a stack returned,
void DFSpath(graph g, int v, int w, stack<int>& path) {
  if (empty(g)) return;
 for (int i = 0; i < V(g); i++) {
   g->marked[i] = false;
   g->parentDFS[i] = -1;
 queue<int> q;
 DFS(g, v, q); // DFS at v, starting vertex
 g->DFSv = q; // DFS result at v
 path = {};
 cout << "your code here\n"; // push v to w path to the stack path</pre>
```

V	marked[]	parent[v]
0	Т	-
1	T	0
2	T	0
3	Т	5
4	Т	6
5	Т	4
6	Т	0
7	F	-
8	F	-
9	F	-
10	9 F	_
1:	1 F	-
1	2 F	-

DFS: Depth-First Search Properties

- Proposition: After DFS, can find vertices connected to sin constant time and can find a path to s (if one exists) in time proportional to its length.
- Proof: parent[] is parent-link representation of a tree rooted at s.
- What is the path from vertex 0 to vertex 3?
- In this case, what is in the stack when parent[] returns?
 void DFSpath(graph g, int v, int w, stack<int>& path)

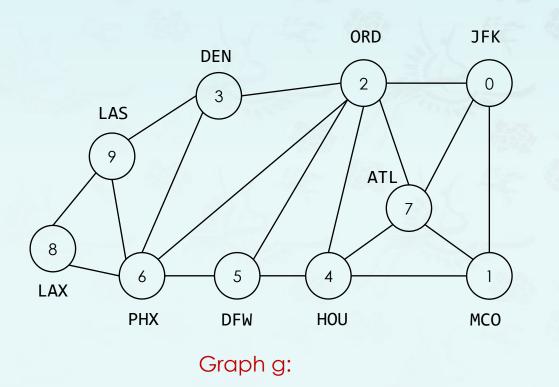




DFS/BFS - Exercise

To visit a vertex v:

- Mark vertex v as visited.
- Recursively visit all unmarked vertices adjacent to v.



routes.txt edge list

16)	
18	3	
0	1	
2	3	
2	4	
5	6	
0	7	
0	2	
2	5	
2	6	

Adjacency list

DFS/BFS - Exercise

- Get familiar with using graphx.exe and using ~.txt graph files provided in pset.
- Create an routes.txt such that it generates the results as shown in the screen capture. Run DFS/BFS and get familiar with the results, its meaning and menu items.

```
Adjacency-list:
V[0]: 2 7 1
V[1]: 4 7 0
V[2]: 7 6 5 0 4 3
V[3]: 9 6 2
V[4]: 5 1 7 2
V[5]: 4 2 6
V[6]: 8 9 3 2 5
V[7]: 1 2 4 0
V[8]: 9 6
V[9]: 8 6 3
   [8]--[6]----[5]----[4]----[1]
    vertex[0..9] =
     color[0..9] =
     DFS0[0..9] =
      CCID[0..9] =
DFS parent[0..9] =
      BFS0[0..9] =
    DistTo[0..9] =
BFS parent[0..9] =
Graph [AdjList][Graph][Tablet]
                                file:routes.txt V:10 E:36 CCs
n - new graph file
                        x - connected(v, w)
d - DFS(v=0)
                        e - distance(v,w)
b - BFS(v=0)
                        p - path(v,w)
m - print mode[adjList/graph]
Command(q to quit):
```

Data Structures Chapter 7: Graph

- 1. Introduction
 - Terminology, Representation, ADT
- 2. Basic Operations
 - DFS, CC, BFS, Processing
- 3. Digraph and Applications
- 4. Minimum Spanning Tree(MST)