Data Structures Chapter 5 Tree

- 1. Introduction
- 2. Binary Tree
- 3. Binary Search Tree
- 4. Balancing Tree
 - AVL Tree Introduction
 - AVL Operations
 - AVL Coding

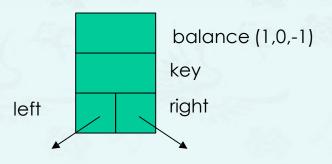


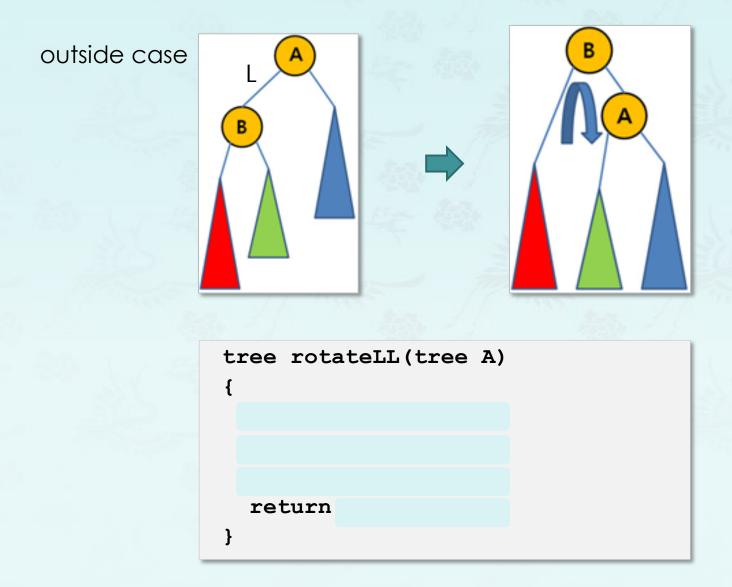
모든 성경은 하나님의 감동으로 된 것으로 교훈과 책망과 바르게 함과 의로 교육하기에 유익하니이는 하나님의 사람으로 온전하게 하며 모든 선한 일을 행할 능력을 갖추게 하려 함이라 (딤후3:16-17)

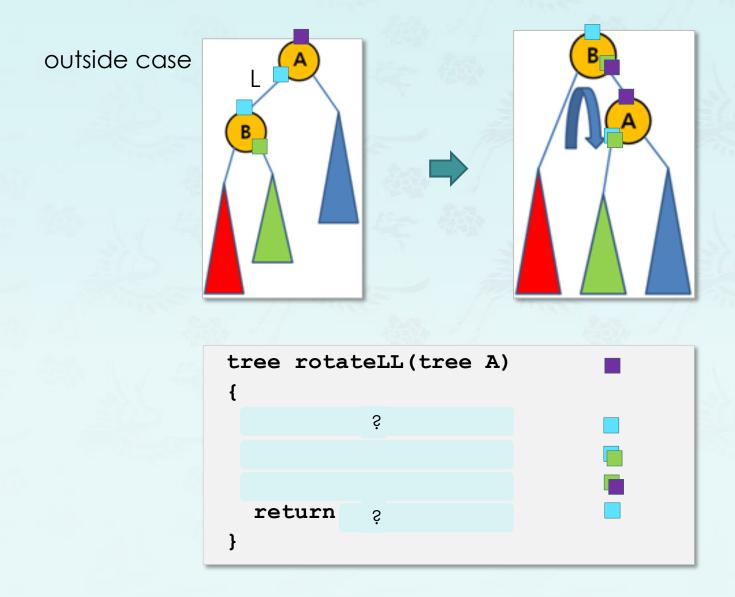
우리는 그가 만드신 바라 그리스도 예수 안에서 선한 일을 위하여 지으심을 받은 자니 이일은 하나님이 전에 예비하사 우리로 그 가운데서 행하게 하려 하심이니라 (엡2:10)

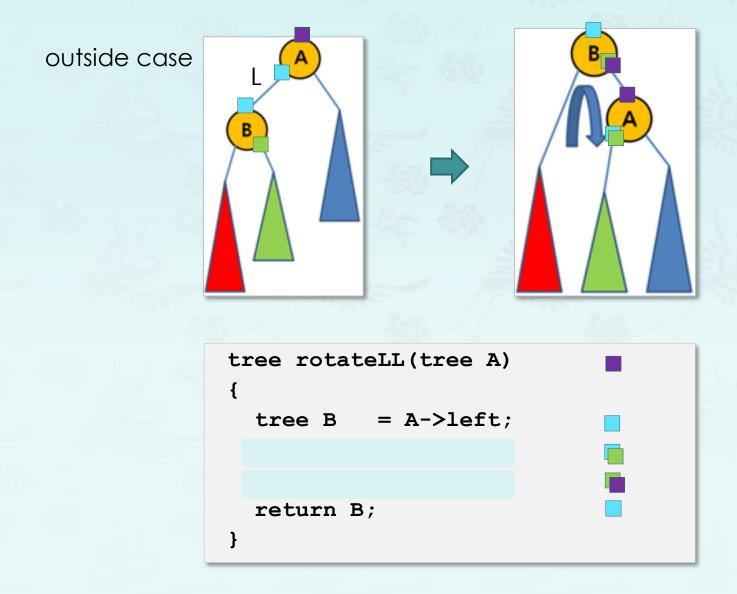
Coding

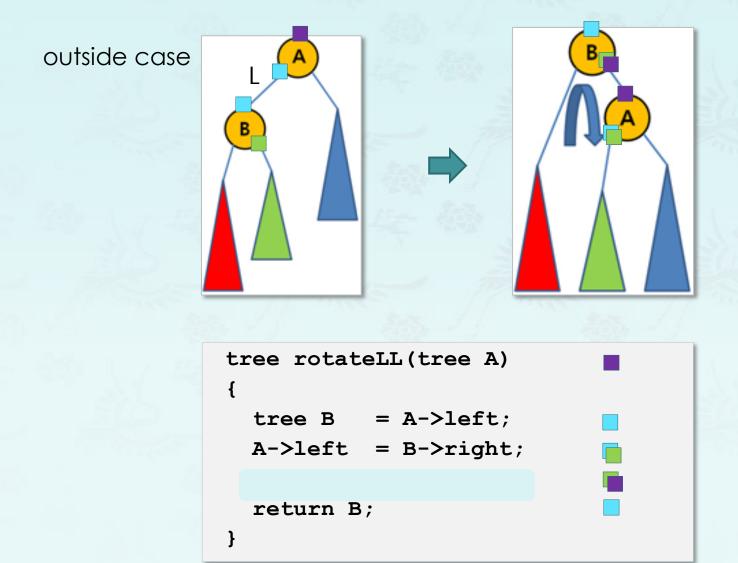
- You can either keep the height or just the difference in height, i.e. the balance factor; this has to be modified on the path of insertion even if you don't perform rotations.
 - Once you have performed a rotation (single or double) you won't need to go back up the tree for the computation.
- You may compute the balance factor on the fly after the insert is done during the recursion.

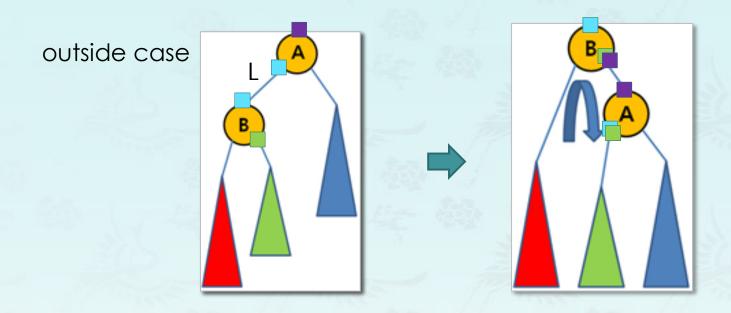




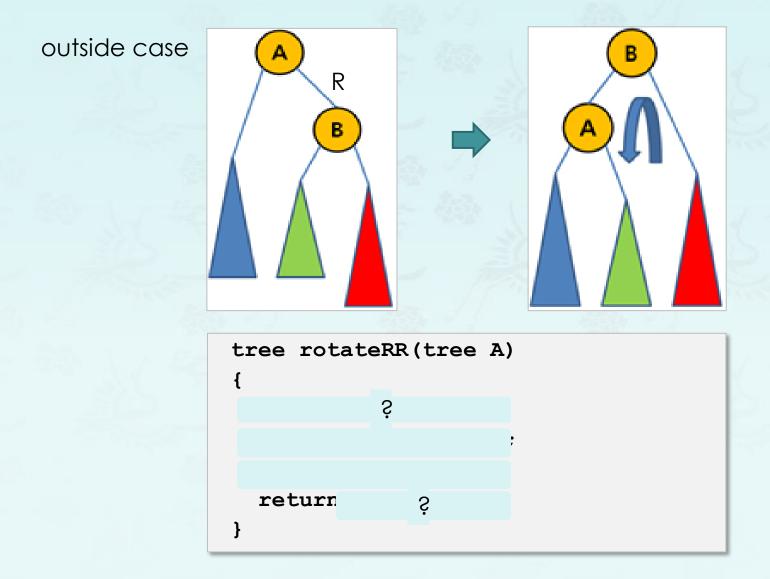


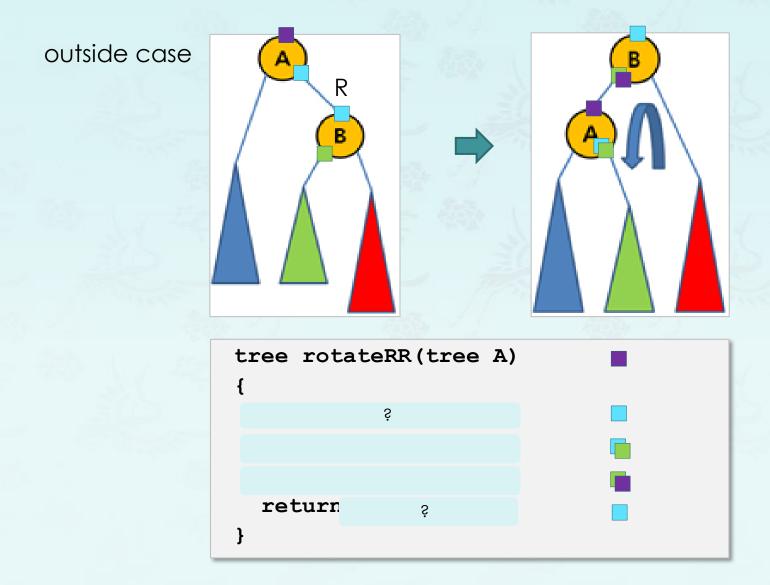


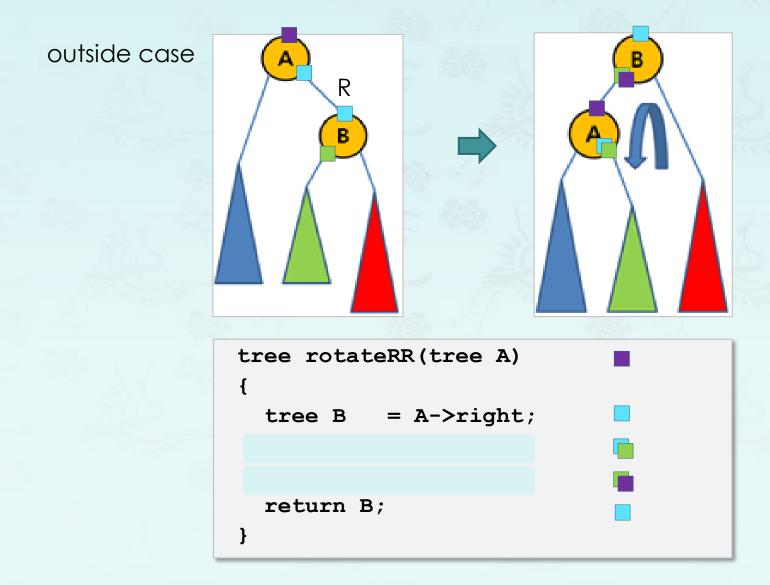


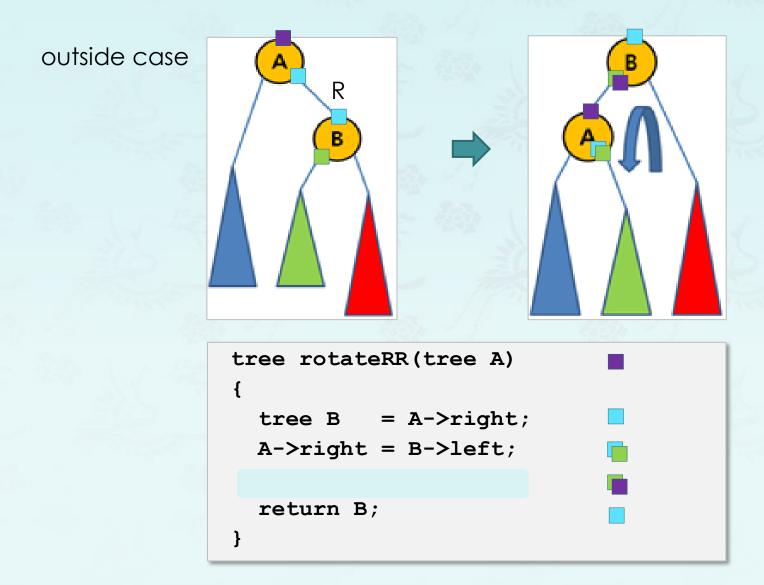


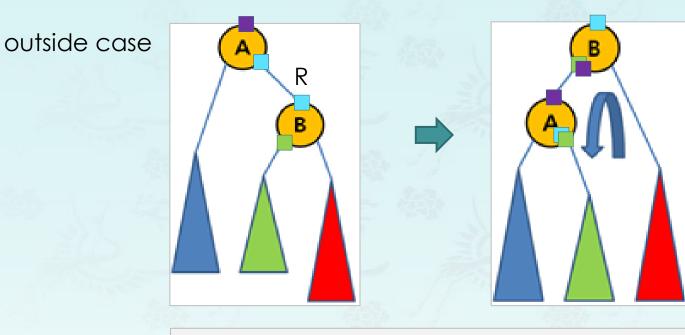
```
tree rotateLL(tree A)
{
  tree B = A->left;
  A->left = B->right;
  B->right = A;
  return B;
}
```





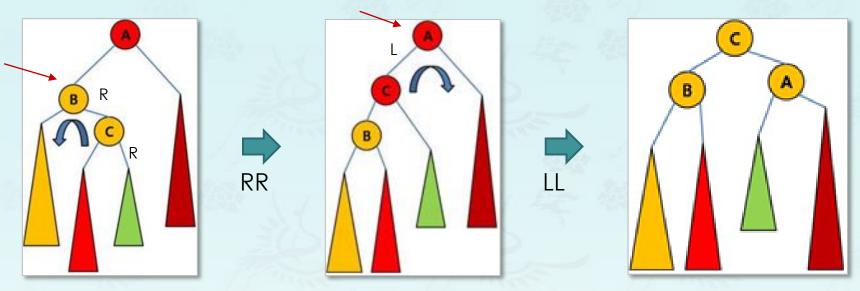




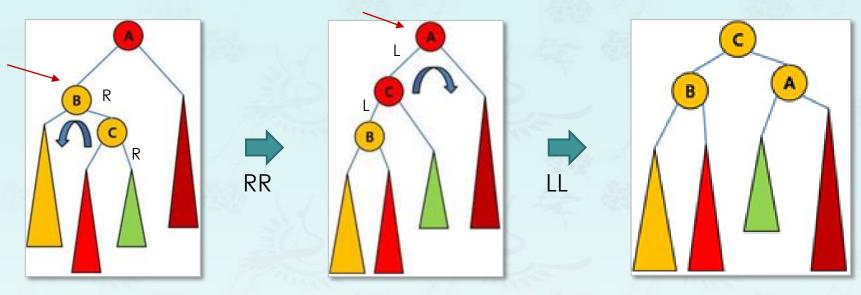


```
tree rotateRR(tree A)
{
  tree B = A->right;
  A->right = B->left;
  B->left = A;
  return B;
}
```

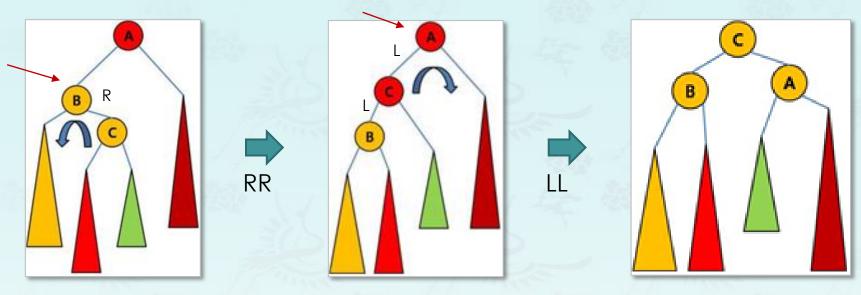
inside case RR tree rotateLR(tree A) // RR and LL



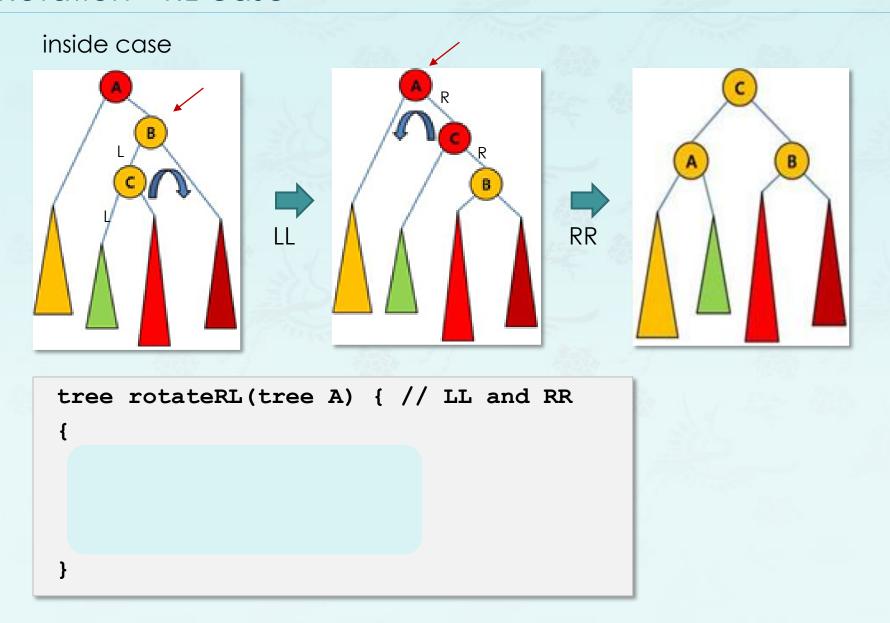
```
tree rotateLR(tree A) // RR and LL
{
  tree B = A->left;
} What will return eventually?
```

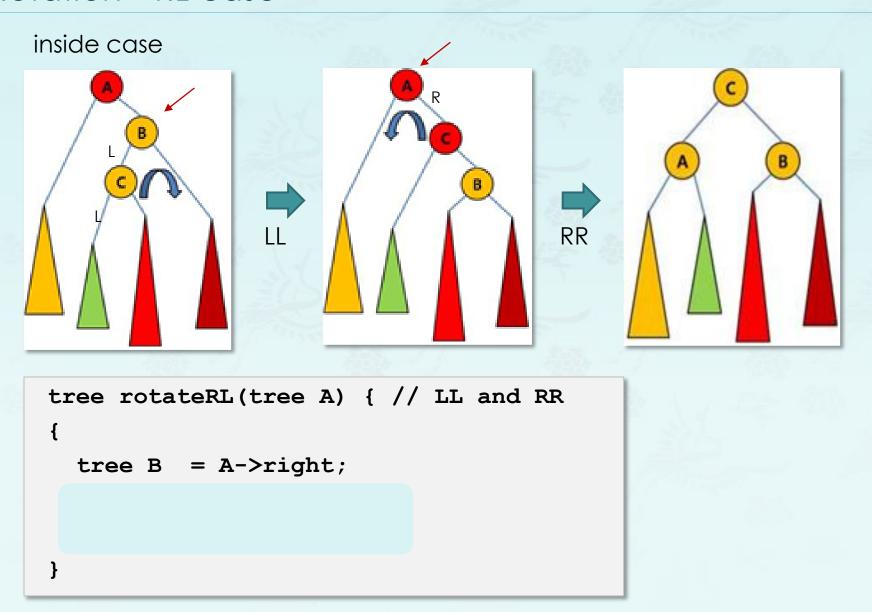


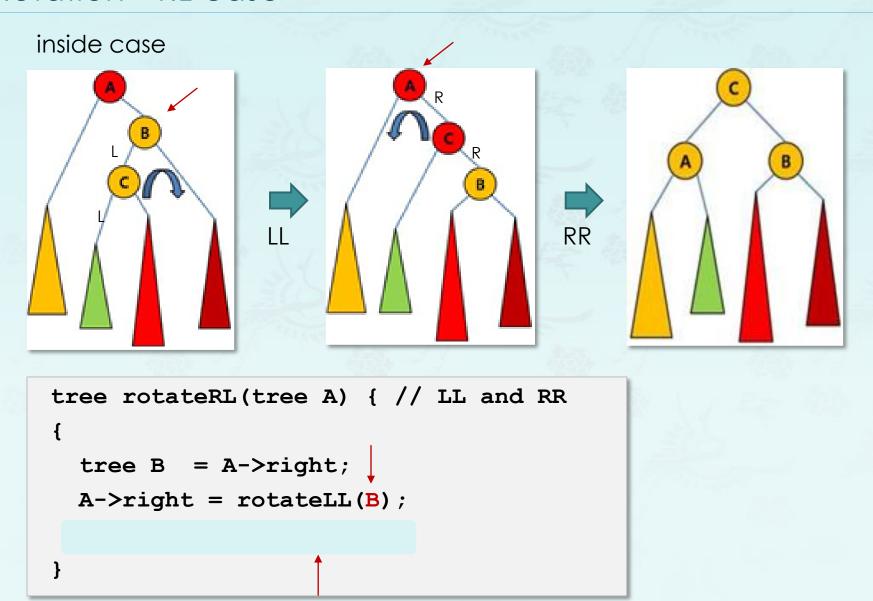
```
tree rotateLR(tree A) // RR and LL
{
   tree B = A->left;
   A->left = rotateRR(B);
}
What will return eventually?
```

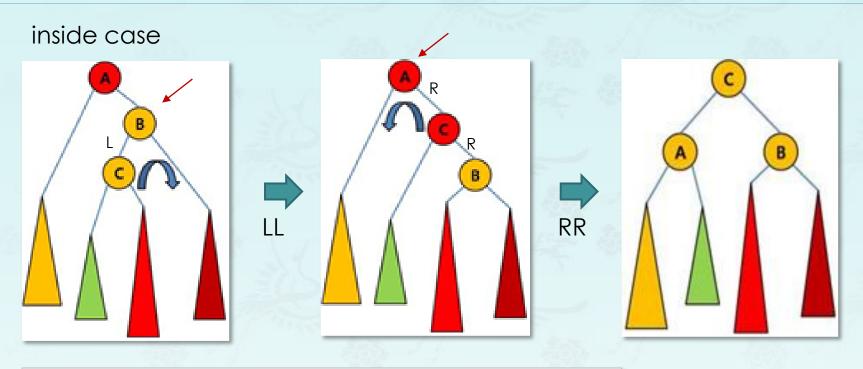


```
tree rotateLR(tree A) // RR and LL
{
  tree B = A->left;
  A->left = rotateRR(B);
  return rotateLL(A);
} What will return eventually?
```





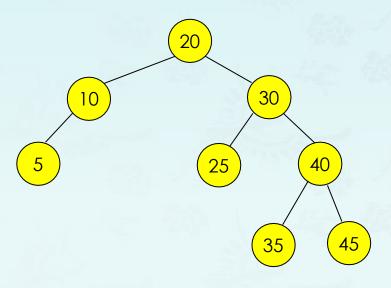




```
tree rotateRL(tree A) { // LL and RR
{
  tree B = A->right;
  A->right = rotateLL(B);
  return rotateRR(A);
}
```

Double Rotation -??case

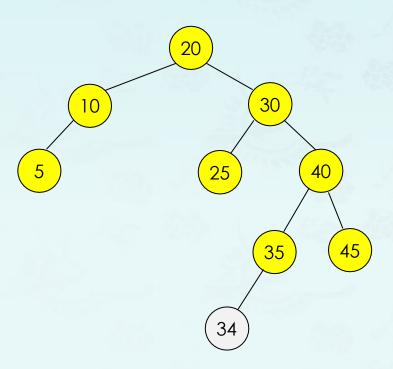
- Insertion of 34
- Imbalance at ?
- Balance factor ??
- Rotation ___??__ case



AVL balanced tree

Double Rotation -??case

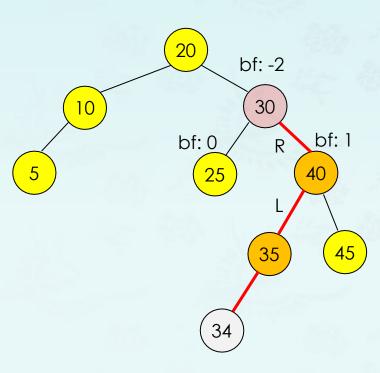
- Insertion of 34
- Imbalance at ?
- Balance factor ??
- Rotation ___??__ case



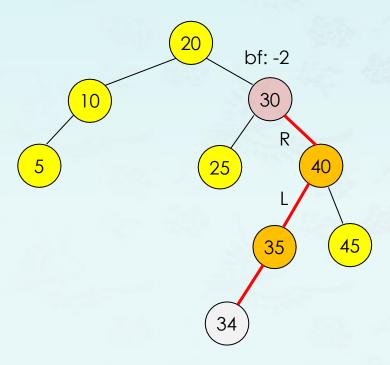
After insertion, AVL imbalanced tree

Double Rotation -??case

- Insertion of 34
- Imbalance at 30
- Balance factor -2
- Rotation ___??__ case

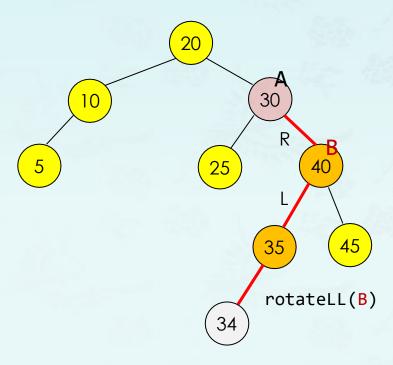


- Insertion of 34
- Imbalance at 30
- Balance factor -2
- Rotation ___RL__ case



```
tree rotateRL(tree A) {
  tree B = A->right;
  A->right = rotateLL(B);
  return rotateRR(A);
}
```

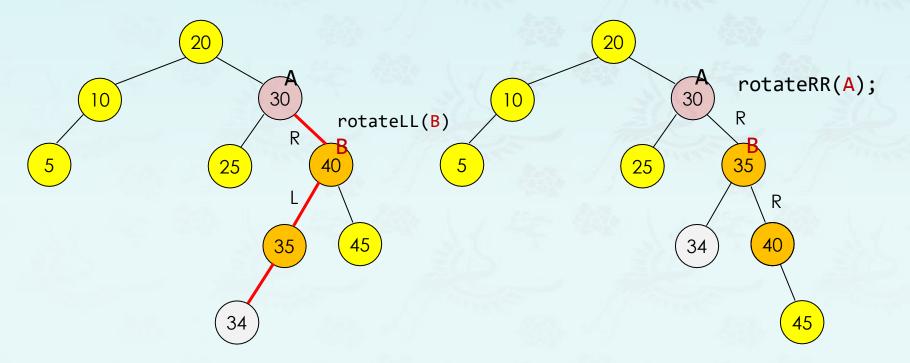
- Insertion of 34
- Imbalance at 30
- Balance factor -2
- Rotation ___RL__ case



```
tree rotateRL(tree A) {
  tree B = A->right;
  A->right = rotateLL(B);
  return rotateRR(A);
}
```

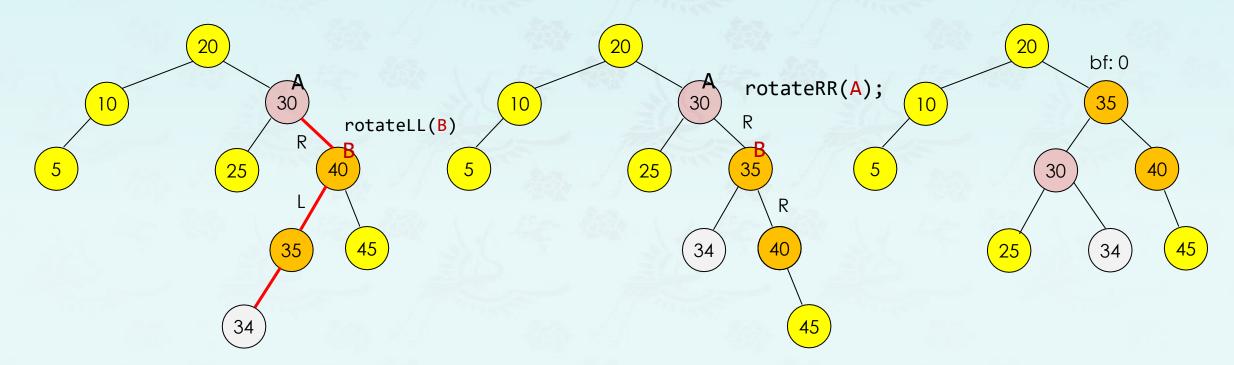
- Insertion of 34
- Imbalance at 30
- Balance factor -2
- Rotation ___RL__ case

```
tree rotateRL(tree A) {
  tree B = A->right;
  A->right = rotateLL(B);
  return rotateRR(A);
}
```



- Insertion of 34
- Imbalance at 30
- Balance factor -2
- Rotation ___RL__ case

```
tree rotateRL(tree A) {
  tree B = A->right;
  A->right = rotateLL(B);
  return rotateRR(A);
}
```



After insertion, AVL imbalanced tree

After insertion, AVL balanced tree

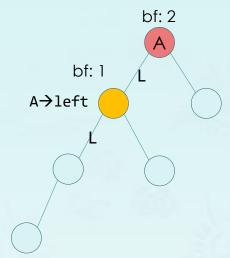
Balance Factor and Height

```
int height(tree node) {
   if (empty(node)) return -1;
   int left = height(node->left);
   int right = height(node->right);
   return max(left, right) + 1;
}
```

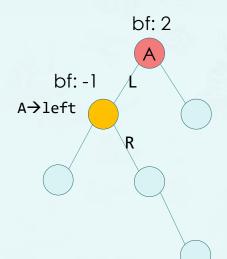
```
int balanceFactor(tree node) {
  if (node == NULL) return 0;
  int left = height(node->left);
  int right = height(node->right);
  return left - right;
}
```

Rebalance

outside case



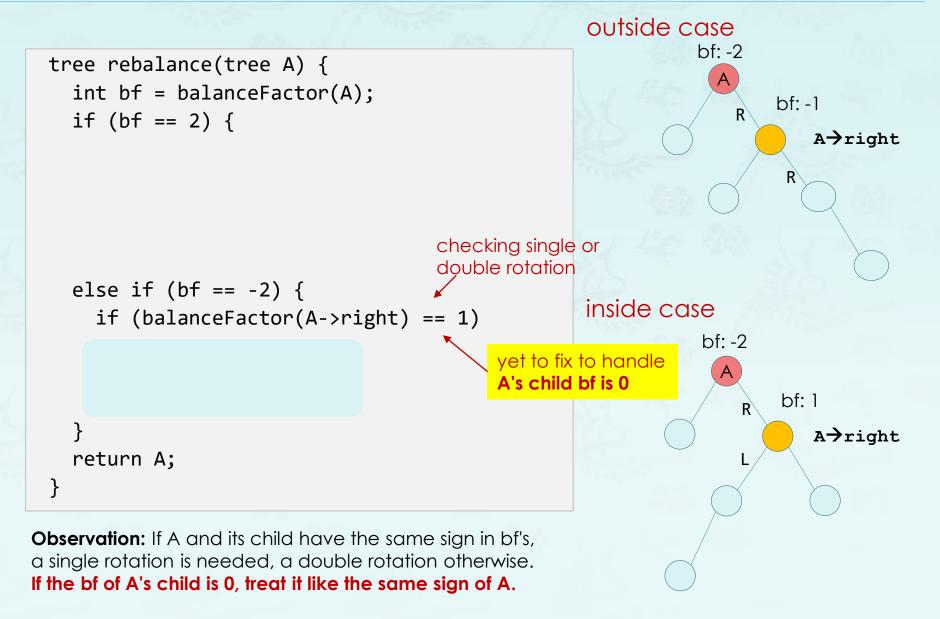
inside case



```
tree rebalance(tree A) {
                                   checking single or
  int bf = balanceFactor(A);
                                   double rotation
  if (bf == 2) {
                                                 yet to fix to handle
    if (balanceFactor(A->left) == 1) ◀
                                                 A's child bf is 0
```

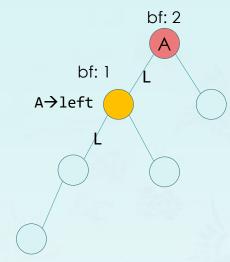
Observation: If A and its child have the same sign in bf's, a single rotation is needed, a double rotation otherwise. If the bf of A's child is 0, treat it like the same sign of A.

Rebalance

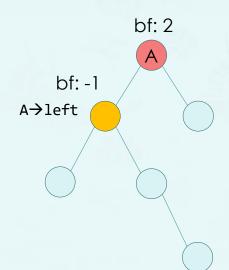


Rebalance

outside case



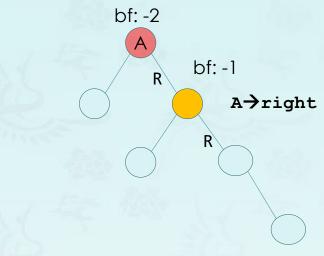
inside case

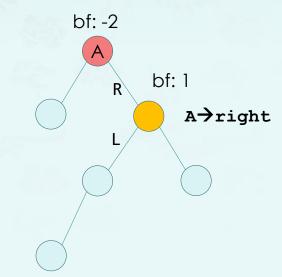


```
tree rebalance(tree A) {
  int bf = balanceFactor(A);
  if (bf == 2) {
    if (balanceFactor(A->left) == 1)
                                yet to fix to handle
                                A's child bf is 0
  else if (bf == -2) {
    if (balanceFactor(A->right) == -1)
  return A; // no rebalanced needed
```

Observation: If A and its child have the same sign in bf's, a single rotation is needed, a double rotation otherwise. If the bf of A's child is 0, treat it like the same sign of A.

outside case





growAVL() & trimAVL()

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