Parallel Programming COMP5112

Distributed Memory Programming With MPI (2)

Slides adapted from the lecture notes by Peter Pacheco

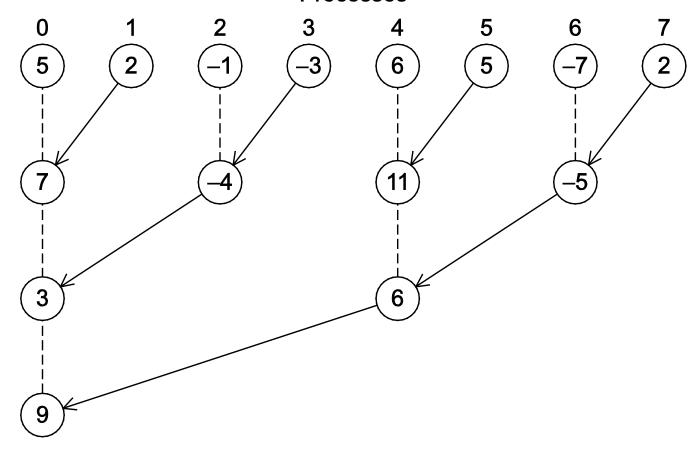
Roadmap

- Writing your first MPI program.
- Using the common MPI functions.
- The Trapezoidal Rule in MPI.
- Collective communication.
- MPI derived datatypes.
- Performance evaluation of MPI programs.
- Parallel sorting.
- Safety in MPI programs.

COLLECTIVE COMMUNICATION

A tree-structured global sum

Processes

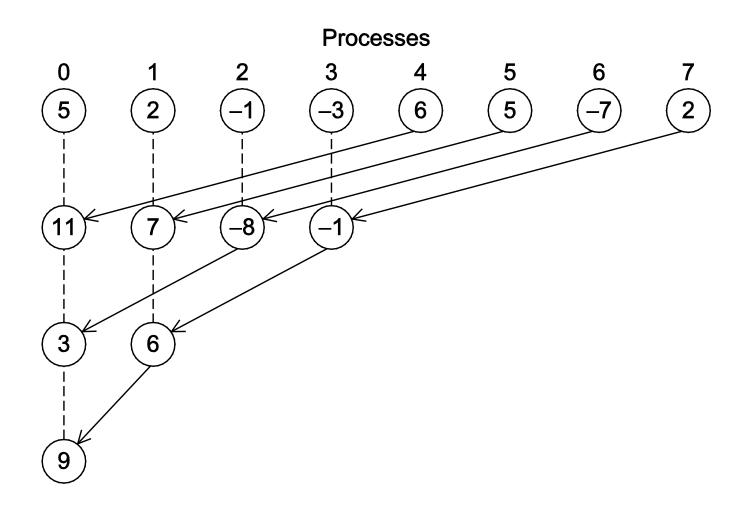


Tree-structured communication

1. In the first phase:

- (a) Process 1 sends to 0, 3 sends to 2, 5 sends to 4, and 7 sends to 6.
- (b) Processes 0, 2, 4, and 6 add in the received values.
- (c) Processes 2 and 6 send their new values to processes 0 and 4, respectively.
- (d) Processes 0 and 4 add the received values into their new values.
- 2. (a) Process 4 sends its newest value to process 0.
 - (b) Process 0 adds the received value to its newest value.

An alternative tree-structured global sum



MPI_Reduce

```
\label{eq:mpi_reduce} \begin{split} \texttt{MPI\_Reduce}(\&\texttt{local\_int}\,,\,\,\&\texttt{total\_int}\,,\,\,1\,,\,\,\texttt{MPI\_DOUBLE}\,,\,\,\texttt{MPI\_SUM}\,,\,\,0\,,\\ \texttt{MPI\_COMM\_WORLD}\,); \end{split}
```

Predefined reduction operators in MPI

Operation Value	Meaning
MPI_MAX	Maximum
MPI_MIN	Minimum
MPI_SUM	Sum
MPI_PROD	Product
MPI_LAND	Logical and
MPI_BAND	Bitwise and
MPI_LOR	Logical or
MPI_BOR	Bitwise or
MPI_LXOR	Logical exclusive or
MPI_BXOR	Bitwise exclusive or
MPI_MAXLOC	Maximum and location of maximum
MPI_MINLOC	Minimum and location of minimum

• <u>All</u> the processes in the communicator must call the same collective function.

 For example, a program that attempts to match a call to MPI_Reduce on one process with a call to MPI_Recv on another process is erroneous, and, in all likelihood, the program will hang or crash.

 The arguments passed by each process to an MPI collective communication must be "compatible."

• For example, if one process passes in 0 as the dest_process and another passes in 1, then the outcome of a call to MPI_Reduce is erroneous, and, once again, the program is likely to hang or crash.

 The output_data_p argument is only used on dest_process.

 However, all of the processes still need to pass in an actual argument corresponding to output_data_p, even if it's just NULL.

 Point-to-point communications are matched on the basis of tags and communicators.

- Collective communications don't use tags.
- They're matched solely on the basis of the communicator and the order in which they're called.

Example (1)

Time	Process 0	Process 1	Process 2
0	a = 1; c = 2	a = 1; c = 2	a = 1; c = 2
1	MPI_Reduce(&a, &b,)	MPI_Reduce(&c, &d,)	MPI_Reduce(&a, &b,)
2	MPI_Reduce(&c, &d,)	MPI_Reduce(&a, &b,)	MPI_Reduce(&c, &d,)

Multiple calls to MPI_Reduce

Example (2)

 Suppose that each process calls MPI_Reduce with operator MPI_SUM, and destination process 0.

 At first glance, it might seem that after the two calls to MPI_Reduce, the value of b will be 3, and the value of d will be 6.

Example (3)

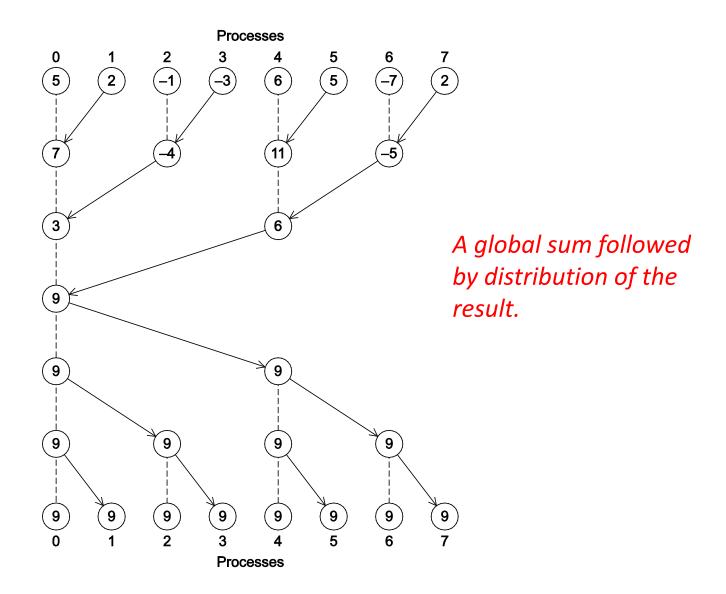
 However, the names of the memory locations are irrelevant to the matching of the calls to MPI_Reduce.

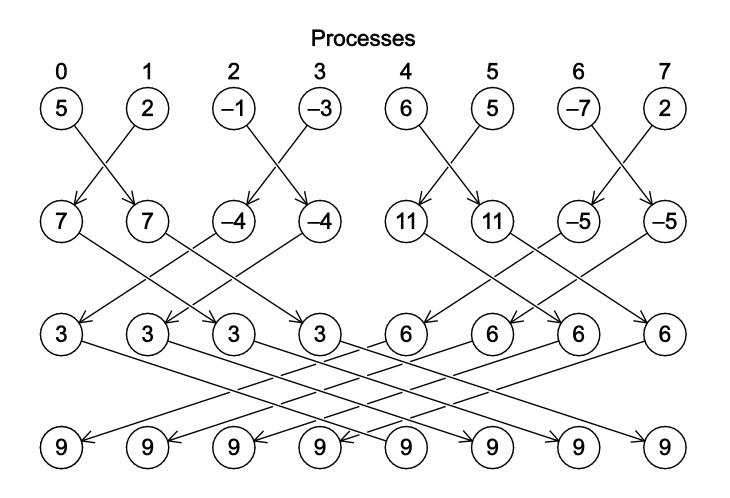
The order of the calls will determine the matching so the value stored in b will be
 1+2+1 = 4, and the value stored in d will be
 2+1+2 = 5.

MPI_Allreduce

 Useful in a situation in which all of the processes need the result of a global sum in order to complete some larger computation.

```
int MPI_Allreduce(
        void*
                    input_data_p /* in */,
        void*
                    output_data_p /* out */,
        int
                                /* in */,
                   count
                             /* in */,
        MPI_Datatype datatype
                                   /* in */,
        qO I 9M
                    operator
                                   /* in */);
        MPI Comm
                    comm
```

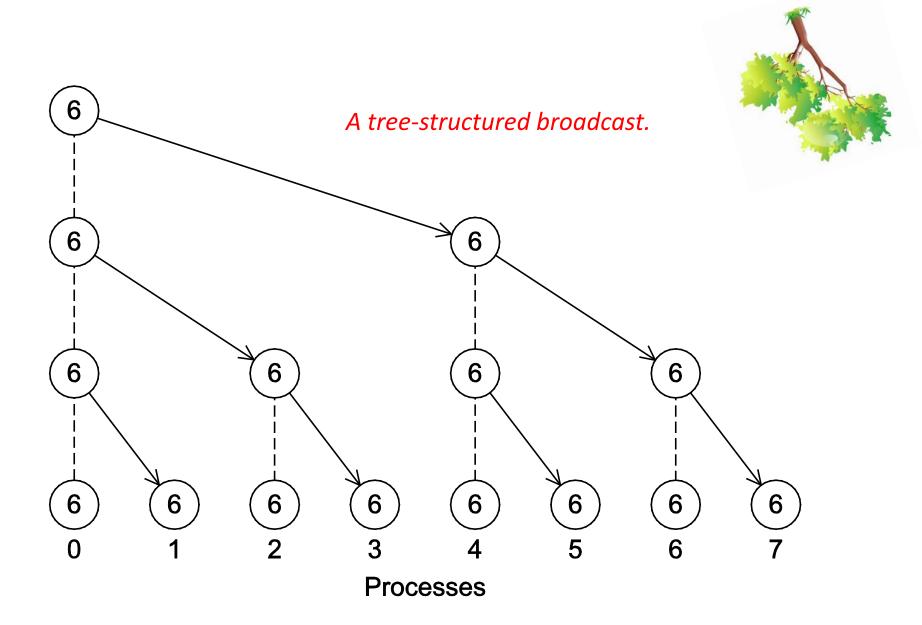




A butterfly-structured global sum.

Broadcast

 Data belonging to a single process is sent to all of the processes in the communicator.



A version of Get_input that uses MPI_Bcast

```
void Get_input(
     int
          my_rank /* in */,
     int comm_sz /* in */,
     double * a_p /* out */,
     double* b_p /* out */,
         n_p /* out */) {
     int*
  if (my_rank == 0) {
     printf("Enter a, b, and n\n");
     scanf("%lf %lf %d", a_p, b_p, n_p);
  MPI_Bcast(a_p, 1, MPI_DOUBLE, 0, MPI_COMM_WORLD);
  MPI_Bcast(b_p, 1, MPI_DOUBLE, 0, MPI_COMM_WORLD);
  MPI_Bcast(n_p, 1, MPI_INT, 0, MPI_COMM_WORLD);
 /* Get_input */
```

Data distributions

$$\mathbf{x} + \mathbf{y} = (x_0, x_1, \dots, x_{n-1}) + (y_0, y_1, \dots, y_{n-1})$$

$$= (x_0 + y_0, x_1 + y_1, \dots, x_{n-1} + y_{n-1})$$

$$= (z_0, z_1, \dots, z_{n-1})$$

$$= \mathbf{z}$$

Compute a vector sum.

Serial implementation of vector addition

```
void Vector_sum(double x[], double y[], double z[], int n) {
  int i;

for (i = 0; i < n; i++)
    z[i] = x[i] + y[i];
} /* Vector_sum */</pre>
```

Different partitions of a 12-component vector among 3 processes

		Components										
		Block-cyclic										
Process	Block			Cyclic			Blocksize = 2					
0	0	1	2	3	0	3	6	9	0	1	6	7
1	4	5	6	7	1	4	7	10	2	3	8	9
2	8	9	10	11	2	5	8	11	4	5	10	11

Partitioning options

- Block partitioning
 - Assign blocks of consecutive components to each process.
- Cyclic partitioning
 - Assign components in a round robin fashion.
- Block-cyclic partitioning
 - Use a cyclic distribution of blocks of components.

Parallel implementation of vector addition

```
void Parallel_vector_sum(
    double local_x[] /* in */,
    double local_y[] /* in */,
    double local_z[] /* out */,
    int local_n /* in */) {
    int local_i;

    for (local_i = 0; local_i < local_n; local_i++)
        local_z[local_i] = local_x[local_i] + local_y[local_i];
} /* Parallel_vector_sum */</pre>
```

Scatter

 MPI_Scatter can be used in a function that reads in an entire vector on process 0 but only sends the needed components to each of the other processes.

```
int MPI Scatter(
    void*
                send buf p /*in */.
    int
                send_count /* in */,
                send_type /* in */,
    MPI_Datatype
    void*
                recv_buf_p /* out */,
    int
                recv count /*in */,
                recv_type /* in */,
    MPI Datatype
    int
                src_proc /* in */,
                comm /* in */):
    MPI Comm
```

Reading and distributing a vector

```
void Read vector(
     double local a[] /* out */,
     int
            local_n /* in */,
     int n /* in */,
     char vec name [] /* in */,
     int my_rank /* in */,
     MPI_Comm comm /* in */) {
  double * a = NULL;
  int i:
  if (my_rank == 0) {
     a = malloc(n*sizeof(double));
     printf("Enter the vector %s\n", vec_name);
     for (i = 0; i < n; i++)
        scanf("%lf", &a[i]);
     MPI Scatter(a, local n, MPI DOUBLE, local a, local n, MPI DOUBLE,
           0. \text{comm}):
     free(a);
  } else {
     MPI Scatter(a, local n, MPI DOUBLE, local a, local n, MPI DOUBLE,
           0. \text{comm}):
  /* Read_vector */
```

Gather

 Collect all of the components of the vector onto process 0, and then process 0 can process all of the components.

Print a distributed vector (1)

```
void Print_vector(
    double local_b[] /* in */,
    int local_n /* in */,
    int
                /* in */,
    char title[] /* in */,
    int
         my_rank /* in */,
                    /* in */) {
    MPI Comm comm
  double*b = NULL;
  int i;
```

Print a distributed vector (2)

```
if (my_rank == 0) 
  b = malloc(n*sizeof(double));
  MPI_Gather(local_b, local_n, MPI_DOUBLE, b, local_n, MPI_DOUBLE,
         0, comm);
  printf("%s\n", title);
   for (i = 0; i < n; i++)
      printf("%f ", b[i]);
  printf("\n");
  free(b);
} else {
  MPI Gather (local b, local n, MPI DOUBLE, b, local n, MPI DOUBLE,
         0, comm);
/* Print_vector */
```

Allgather

- Concatenates the contents of each process' send_buf_p and stores this in each process' recv_buf_p.
- As usual, recv_count is the amount of data being received from each process.

Matrix-vector multiplication

$$A = (a_{ij})$$
 is an $m \times n$ matrix

 \mathbf{x} is a vector with n components

y = Ax is a vector with m components

$$y_i = a_{i0}x_0 + a_{i1}x_1 + a_{i2}x_2 + \cdots + a_{i,n-1}x_{n-1}$$

i-th component of y

Dot product of the ith row of A with x.

Matrix-vector multiplication

				_	
<i>a</i> ₀₀	<i>a</i> ₀₁		$a_{0,n-1}$		
a_{10}	<i>a</i> ₁₁	• • • •	$a_{1,n-1}$	x_0	
:	:		:	x_1	
a_{i0}	a_{i1}	• • • •	$a_{i,n-1}$: =	$y_i =$
:	:		:	x_{n-1}	
$a_{m-1,0}$	$a_{m-1,1}$		$a_{m-1,n-1}$		

	уо
Ī	<i>y</i> ₁
_	:
-	$y_i = a_{i0}x_0 + a_{i1}x_1 + \cdots + a_{i,n-1}x_{n-1}$
	:
	y_{m-1}

Multiply a matrix by a vector

```
/* For each row of A */
for (i = 0; i < m; i++) {
    /* Form dot product of ith row with x */
    v[i] = 0.0;

for (j = 0; j < n; j++)
    y[i] += A[i][j]*x[j];
}</pre>
```

Serial pseudo-code

C style arrays

$$\begin{pmatrix} 0 & 1 & 2 & 3 \\ 4 & 5 & 6 & 7 \\ 8 & 9 & 10 & 11 \end{pmatrix}$$
 stored as

0 1 2 3 4 5 6 7 8 9 10 11

Serial matrix-vector multiplication

```
void Mat_vect_mult(
     double A[] /* in */,
     double x[] /* in */,
     double y[] /* out */,
     int m /*in */,
     int n /* in */) {
  int i, j;
  for (i = 0; i < m; i++)
     y[i] = 0.0;
     for (j = 0; j < n; j++)
        y[i] += A[i*n+j]*x[j];
} /* Mat_vect_mult */
```

An MPI matrix-vector multiplication function (1)

```
void Mat_vect_mult(
     double local_A[] /* in */,
     double local_x[] /* in */,
     double local_y[] /* out */,
     int
             local_m /* in */,
                /* in */,
     int
             n
        local_n /*in */,
     int
     MPI_Comm comm /* in */) {
  double * x;
  int local_i, j;
  int local_ok = 1;
```

An MPI matrix-vector multiplication function (2)