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Message Passing Interface (MPI)

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Abstract

The Message Passing Interface Standard (MPI) is a message passing library standard based on the consensus of the MPI Forum, which has over 40 participating organizations, including vendors, researchers, software library developers, and users. The goal of the Message Passing Interface is to establish a portable, efficient, and flexible standard for message passing that will be widely used for writing message passing programs. As such, MPI is the first standardized, vendor independent, message passing library. The advantages of developing message passing software using MPI closely match the design goals of portability, efficiency, and flexibility. MPI is not an IEEE or ISO standard, but has in fact, become the "industry standard" for writing message passing programs on HPC platforms.

The goal of this tutorial is to teach those unfamiliar with MPI how to develop and run parallel programs according to the MPI standard. The primary topics that are presented focus on those which are the most useful for new MPI programmers. The tutorial begins with an introduction, background, and basic information for getting started with MPI. This is followed by a detailed look at the MPI routines that are most useful for new MPI programmers, including MPI Environment Management, Point-to-Point Communications, and Collective Communications routines. Numerous examples in both C and Fortran are provided, as well as a lab exercise.

The tutorial materials also include more advanced topics such as Derived Data Types, Group and Communicator Management Routines, and Virtual Topologies. However, these are not actually presented during the lecture, but are meant to serve as "further reading" for those who are interested.

Level/Prerequisites: This tutorial is ideal for those who are new to parallel programming with MPI. A basic understanding of parallel programming in C or Fortran is required. For those who are unfamiliar with Parallel Programming in general, the material covered in EC3500: Introduction To Parallel Computing would be helpful.

What is MPI?

An Interface Specification:

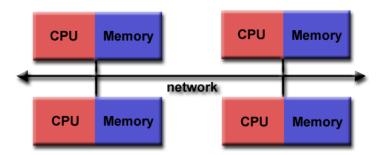
- M P I = Message Passing Interface
- MPI is a specification for the developers and users of message passing libraries. By itself, it is NOT a library - but rather the specification of what such a library should be.
- MPI primarily addresses the message-passing parallel programming model: data is moved from the address space of one process to that of another process through cooperative operations on each process.



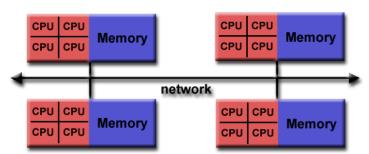
- Simply stated, the goal of the Message Passing Interface is to provide a widely used standard for writing message passing programs. The interface attempts to be:
 - Practical
 - o Portable
 - Efficient
 - Flexible
- The MPI standard has gone through a number of revisions, with the most recent version being MPI-3.x
- Interface specifications have been defined for C and Fortran90 language bindings:
 - C++ bindings from MPI-1 are removed in MPI-3
 - MPI-3 also provides support for Fortran 2003 and 2008 features
- Actual MPI library implementations differ in which version and features of the MPI standard they support. Developers/users will need to be aware of this.

Programming Model:

Originally, MPI was designed for distributed memory architectures, which were becoming increasingly popular at that time (1980s - early 1990s).



- As architecture trends changed, shared memory SMPs were combined over networks creating hybrid distributed memory / shared memory systems.
- MPI implementors adapted their libraries to handle both types of underlying memory architectures seamlessly. They also adapted/developed ways of handling different interconnects and protocols.



- Today, MPI runs on virtually any hardware platform:
 - Distributed Memory
 - Shared Memory
 - Hybrid
- The programming model <u>clearly remains a distributed memory model</u> however, regardless of the underlying physical architecture
 of the machine.
- All parallelism is explicit: the programmer is responsible for correctly identifying parallelism and implementing parallel algorithms
 using MPI constructs.

Reasons for Using MPI:

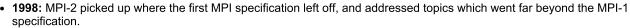
- Standardization MPI is the only message passing library that can be considered a standard. It is supported on virtually all HPC platforms. Practically, it has replaced all previous message passing libraries.
- Portability There is little or no need to modify your source code when you port your application to a different platform that supports (and is compliant with) the MPI standard.
- Performance Opportunities Vendor implementations should be able to exploit native hardware features to optimize
 performance. Any implementation is free to develop optimized algorithms.
- Functionality There are over 430 routines defined in MPI-3, which includes the majority of those in MPI-2 and MPI-1.

Note: Most MPI programs can be written using a dozen or less routines

• Availability - A variety of implementations are available, both vendor and public domain.

History and Evolution: (for those interested)

- MPI has resulted from the efforts of numerous individuals and groups that began in 1992. Some history:
- 1980s early 1990s: Distributed memory, parallel computing develops, as do a number of incompatible software tools for writing such programs usually with tradeoffs between portability, performance, functionality and price. Recognition of the need for a standard arose.
- Apr 1992: Workshop on Standards for Message Passing in a
 Distributed Memory Environment, sponsored by the Center
 for Research on Parallel Computing, Williamsburg, Virginia.
 The basic features essential to a standard message passing
 interface were discussed, and a working group established to
 continue the standardization process. Preliminary draft
 proposal developed subsequently.
- Nov 1992: Working group meets in Minneapolis. MPI draft proposal (MPI1) from ORNL presented. Group adopts procedures and organization to form the MPI Forum. It eventually comprised of about 175 individuals from 40 organizations including parallel computer vendors, software writers, academia and application scientists.
- Nov 1993: Supercomputing 93 conference draft MPI standard presented.
- May 1994: Final version of MPI-1.0 released
 - o MPI-1.1 (Jun 1995)
 - MPI-1.2 (Jul 1997)
 - o MPI-1.3 (May 2008).



- o MPI-2.1 (Sep 2008)
- MPI-2.2 (Sep 2009)
- Sep 2012: The MPI-3.0 standard was approved.
 - o MPI-3.1 (Jun 2015)

Documentation:

Documentation for all versions of the MPI standard is available at: http://www.mpi-forum.org/docs/.

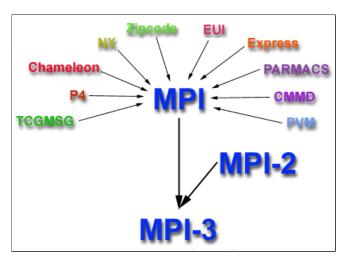
LLNL MPI Implementations and Compilers

- Although the MPI programming interface has been standardized, actual library implementations will differ in which version and features of the standard they support.
- The way MPI programs are compiled and run on different platforms may also vary.
- Currently, LC supports these MPI implementations:
 - o MVAPICH Linux clusters
 - o Open MPI Linux clusters
 - Intel MPI Linux clusters
 - IBM BG/Q MPI BG/Q clusters
 - o IBM Spectrum MPI Coral Early Access (Power8 + Pascal) clusters
- A summary of each is provided below, along with links to additional detailed information.

MVAPICH

General Info:

- MVAPICH MPI is developed and supported by the Network-Based Computing Lab at Ohio State University.
- · Available on all of LC's Linux clusters.
- MVAPICH 1.2 (TOSS 2 clusters only)
 - o Default version of MPI
 - MPI-1 implementation that also includes support for MPI-I/O
 - Based on MPICH-1.2.7 MPI library from Argonne National Laboratory
 - o Not thread-safe. All MPI calls should be made by the master thread in a multi-threaded MPI program.





- See /usr/local/docs/mpi.mvapich.basics for LC usage details.
- MVAPICH2 (TOSS 2 and TOSS 3 clusters)
 - Multiple versions available
 - MPI-2 and MPI-3 implementations based on MPICH MPI library from Argonne National Laboratory. Versions 1.9 and later implement MPI-3 according to the developer's documentation.
 - TOSS 3: Default MPI implementation
 - TOSS 2: Not the default requires the "use" command to load the selected dotkit package. For example:

```
use -1 mvapich (list available packages)
use mvapich2-intel-2.1 (use the package of interest)
```

o Thread-safe

Compiling:

• See the MPI Build Scripts table below.

Running:

• MPI executables are launched using the SLURM srun command with the appropriate options. For example, to launch an 8-process MPI job split across two different nodes in the pdebug pool:

```
srun -N2 -n8 -ppdebug a.out
```

• The srun command is discussed in detail in the Running Jobs section of the Linux Clusters Overview tutorial.

Documentation:

- MVAPICH home page: mvapich.cse.ohio-state.edu/
- MVAPICH2 User Guides: http://mvapich.cse.ohio-state.edu/userguide/
- MVAPICH 1.2 User Guide: available HERE
- MPICH home page: http://www.mpich.org/
- /usr/local/docs on LC's TOSS 2 clusters:
 - o mpi.basics
 - o mpi.mvapich.basics
 - mpi.mvapich2.basics

Open MPI

General Information:

dotkit (TOSS 2 clusters):

- Open MPI is a thread-safe, open source MPI implementation developed and supported by a consortium of academic, research, and industry partners.
- Available on all LC Linux clusters. However, you'll need to load the desired <u>dotkit</u> or <u>module</u> first. For example:



```
use -1 openmpi (list available packages)
use openmpi-gnu-1.8.4 (use the package of interest)

module (TOSS 3 clusters):

module avail (list available modules)
module load openmpi/2.0.0 (use the module of interest)
```

This ensures that LC's MPI wrapper scripts point to the desired version of Open MPI.

Compiling:

See the MPI Build Scripts table below.

Running:

- Be sure to load the same Open MPI dotkit/module that you used to build your executable. If you are running a batch job, you will need to load the dotkit/module in your batch script.
- Launching an Open MPI job can be done using the following commands. For example, to run a 48 process MPI job:

```
mpirun -np 48 a.out
mpiexec -np 48 a.out
```

srun -n 48 a.out

Documentation:

- Open MPI home page: http://www.open-mpi.org/
- /usr/local/docs/mpi.openmpi.basics on LC's TOSS 2 clusters:

Intel MPI:

- Available on LC's Linux clusters.
- Based on MPICH3. Supports MPI-3 functionality.
- Thread-safe
- Compiling and running Intel MPI programs: see the LC documentation at: https://lc.llnl.gov/confluence/pages/viewpage.action?
 pageId=137725526

IBM BG/Q Clusters:

- The IBM BG/Q MPI library is the only supported implementation on these clusters.
- Default version is based on MPICH2, which includes MPI-2 functionality minus dynamic processes.
- A version supporting MPI-3 functionality is available.
- Thread-safe
- Compiling and running IBM BG/Q MPI programs: see the BG/Q Tutorial: computing.llnl.gov/tutorials/bgg/

CORAL Early Access Clusters:

- The IBM Spectrum MPI library is the only supported implementation on these clusters.
- · Based on Open MPI. Includes MPI-3 functionality.
- · Thread-safe
- NVIDIA GPU support
- Compiling and running IBM Spectrum MPI programs: see the CORAL EA MPI documentation at: https://lc.llnl.gov/confluence/display/CORALEA/MPI

MPI Build Scripts

- LC developed MPI compiler wrapper scripts are used to compile MPI programs
- Automatically perform some error checks, include the appropriate MPI #include files, link to the necessary MPI libraries, and pass
 options to the underlying compiler.



Note: you may need to load a dotkit/module for the desired MPI implementation, as discussed above. Failing to do this will result in getting the default implementation.

- For additional information:
 - See the man page (if it exists)
 - Issue the script name with the -help option
 - View the script yourself directly

MPI Build Scripts						
Implementation	Implementation Language Script Name Underlying Compiler					
MVAPCH 1.2 (TOSS 2 only)		mpicc	gcc - GNU			
		mpigcc	gcc - GNU			
		mpiicc	icc - Intel			

incode in the index (in i)				
		mpipgcc	pgcc - PGI	
	C++	mpiCC	g++ - GNU	
		mpig++	g++ - GNU	
		mpiicpc	icpc - Intel	
		mpipgCC	pgCC - PGI	
		mpif77	g77 - GNU	
		mpigfortran	gfortran - GNU	
	Fortran	mpiifort	ifort - Intel	
		mpipgf77	pgf77 - PGI	
		mpipgf90	pgf90 - PGI	
	С	mpicc	C compiler of dotkit package loaded	
MVAPCH2	C++	mpicxx	C++ compiler of dotkit package loaded	
WVAPCHZ	Fortran	mpif77	Fortran77 compiler of dotkit package loaded	
		mpif90	Fortran90 compiler of dotkit package loaded	
	С	mpicc	C compiler of dotkit package loaded	
Open MPI	C++	mpiCC mpic++ mpicxx	C++ compiler of dotkit package loaded	
	Fortran	mpif77	Fortran77 compiler of dotkit package loaded	
		mpif90	Fortran90 compiler of dotkit package loaded	

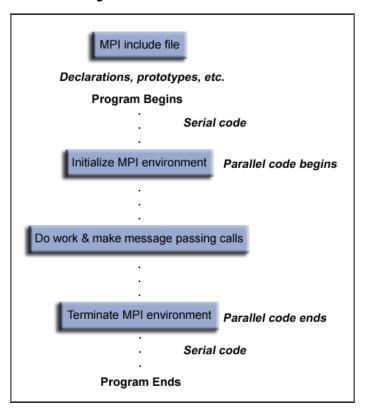
Level of Thread Support

- MPI libraries vary in their level of thread support:
 - MPI_THREAD_SINGLE Level 0: Only one thread will execute.
 - MPI_THREAD_FUNNELED Level 1: The process may be multi-threaded, but only the main thread will make MPI calls all MPI calls are funneled to the main thread.
 - MPI_THREAD_SERIALIZED Level 2: The process may be multi-threaded, and multiple threads may make MPI calls, but only one at a time. That is, calls are not made concurrently from two distinct threads as all MPI calls are serialized.
 - MPI THREAD MULTIPLE Level 3: Multiple threads may call MPI with no restrictions.
- Consult the MPI Init thread() man page for details.
- A simple C language example for determining thread level support is shown below.

```
#include "mpi.h"
#include <stdio.h>
int main( int argc, char *argv[] )
{
    int provided, claimed;
/*** Select one of the following
    MPI_Init_thread( 0, 0, MPI_THREAD_SINGLE, &provided );
    MPI_Init_thread( 0, 0, MPI_THREAD_FUNNELED, &provided );
MPI_Init_thread( 0, 0, MPI_THREAD_SERIALIZED, &provided );
    MPI_Init_thread( 0, 0, MPI_THREAD_MULTIPLE, &provided );
    MPI_Init_thread(0, 0, MPI_THREAD_MULTIPLE, &provided);
    MPI_Query_thread( &claimed );
        printf( "Query thread level= %d Init_thread level= %d\n", claimed, provided );
    MPI Finalize();
}
Sample output:
Query thread level= 3 Init thread level= 3
```

Getting Started

General MPI Program Structure:



Header File:

· Required for all programs that make MPI library calls.

C include file	Fortran include file	
#include "mpi.h"	include 'mpif.h'	

• With MPI-3 Fortran, the USE mpi f08 module is preferred over using the include file shown above.

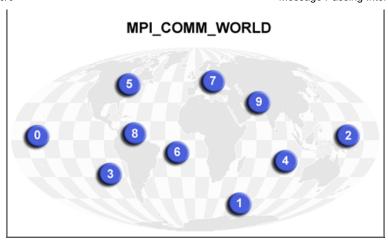
Format of MPI Calls:

- C names are case sensitive; Fortran names are not.
- Programs must not declare variables or functions with names beginning with the prefix MPI_ or PMPI_ (profiling interface).

C Binding				
Format:	rc = MPI_Xxxxx(parameter,)			
Example:	rc = MPI_Bsend(&buf,count,type,dest,tag,comm)			
Error code:	Returned as "rc". MPI_SUCCESS if successful			
	Fortran Binding			
Format:	CALL MPI_XXXXX (parameter,, ierr) call mpi_xxxxx (parameter,, ierr)			
Example:	CALL MPI_BSEND(buf,count,type,dest,tag,comm,ierr)			
Error code:	Returned as "ierr" parameter. MPI_SUCCESS if successful			

Communicators and Groups:

- MPI uses objects called communicators and groups to define which collection of processes may communicate with each other.
- Most MPI routines require you to specify a communicator as an argument.
- Communicators and groups will be covered in more detail later. For now, simply use MPI_COMM_WORLD whenever a
 communicator is required it is the predefined communicator that includes all of your MPI processes.



Rank:

- Within a communicator, every process has its own unique, integer identifier assigned by the system when the process initializes. A rank is sometimes also called a "task ID". Ranks are contiguous and begin at zero.
- Used by the programmer to specify the source and destination of messages. Often used conditionally by the application to control program execution (if rank=0 do this / if rank=1 do that).

Error Handling:

- Most MPI routines include a return/error code parameter, as described in the "Format of MPI Calls" section above.
- However, according to the MPI standard, the default behavior of an MPI call is to abort if there is an error. This means you will
 probably not be able to capture a return/error code other than MPI SUCCESS (zero).
- The standard does provide a means to override this default error handler. A discussion on how to do this is available <u>HERE</u>. You can also consult the error handling section of the relevant MPI Standard documentation located at http://www.mpi-forum.org/docs/.
- The types of errors displayed to the user are implementation dependent.

Environment Management Routines

This group of routines is used for interrogating and setting the MPI execution environment, and covers an assortment of purposes, such as initializing and terminating the MPI environment, querying a rank's identity, querying the MPI library's version, etc. Most of the commonly used ones are described below.

MPI_Init

Initializes the MPI execution environment. This function must be called in every MPI program, must be called before any other MPI functions and must be called only once in an MPI program. For C programs, MPI_Init may be used to pass the command line arguments to all processes, although this is not required by the standard and is implementation dependent.

```
MPI_Init (&argc,&argv)
MPI_INIT (ierr)
```

MPI Comm size

Returns the total number of MPI processes in the specified communicator, such as MPI_COMM_WORLD. If the communicator is MPI_COMM_WORLD, then it represents the number of MPI tasks available to your application.

```
MPI_Comm_size (comm,&size)
MPI_COMM_SIZE (comm,size,ierr)
```

MPI_Comm_rank

Returns the rank of the calling MPI process within the specified communicator. Initially, each process will be assigned a unique integer rank between 0 and number of tasks - 1 within the communicator MPI_COMM_WORLD. This rank is often referred to as a task ID. If a process becomes associated with other communicators, it will have a unique rank within each of these as well.

```
MPI_COMm_rank (comm,&rank)
MPI_COMM_RANK (comm,rank,ierr)
```

MPI Abort

Terminates all MPI processes associated with the communicator. In most MPI implementations it terminates ALL processes regardless of the communicator specified.

```
MPI_Abort (comm,errorcode)
MPI_ABORT (comm,errorcode,ierr)
```

MPI Get processor name

Returns the processor name. Also returns the length of the name. The buffer for "name" must be at least MPI_MAX_PROCESSOR_NAME characters in size. What is returned into "name" is implementation dependent - may not be the same as the output of the "hostname" or "host" shell commands.

```
MPI_Get_processor_name (&name,&resultlength)
MPI_GET_PROCESSOR_NAME (name,resultlength,ierr)
```

MPI_Get_version

Returns the version and subversion of the MPI standard that's implemented by the library.

```
MPI_Get_version (&version,&subversion)
MPI_GET_VERSION (version,subversion,ierr)
```

MPI Initialized

Indicates whether MPI_Init has been called - returns flag as either logical true (1) or false(0). MPI requires that MPI_Init be called once and only once by each process. This may pose a problem for modules that want to use MPI and are prepared to call MPI_Init if necessary. MPI_Initialized solves this problem.

```
MPI_Initialized (&flag)
MPI_INITIALIZED (flag,ierr)
```

MPI Wtime

Returns an elapsed wall clock time in seconds (double precision) on the calling processor.

```
MPI_Wtime ()
MPI_WTIME ()
```

MPI Wtick

Returns the resolution in seconds (double precision) of MPI_Wtime.

```
MPI_Wtick ()
MPI_WTICK ()
```

MPI_Finalize

Terminates the MPI execution environment. This function should be the last MPI routine called in every MPI program - no other MPI routines may be called after it.

```
MPI_Finalize ()
MPI_FINALIZE (ierr)
```

Examples: Environment Management Routines



C Language - Environment Management Routines

```
1 // required MPI include file
```

- 2 #include "mpi.h"
- 3 #include <stdio.h>

```
5
     int main(int argc, char *argv[]) {
     int numtasks, rank, len, rc;
 6
 7
     char hostname[MPI_MAX_PROCESSOR_NAME];
 8
 9
     // initialize MPI
10
     MPI_Init(&argc,&argv);
11
12
     // get number of tasks
13
     MPI Comm size(MPI COMM WORLD, &numtasks);
14
15
     // get my rank
16
     MPI_Comm_rank(MPI_COMM_WORLD,&rank);
17
     // this one is obvious
18
     MPI Get processor name (hostname, &len);
19
     printf ("Number of tasks= %d My rank= %d Running on %s\n", numtasks,rank,hostname);
20
21
22
23
          // do some work with message passing
24
25
26
     // done with MPI
27
     MPI_Finalize();
28
```

```
Fortran - Environment Management Routines
 1
     program simple
 2
 3
     ! required MPI include file
 4
     include 'mpif.h'
 5
 6
     integer numtasks, rank, len, ierr
 7
     character(MPI_MAX_PROCESSOR_NAME) hostname
 8
 9
     ! initialize MPI
     call MPI INIT(ierr)
10
11
12
     ! get number of tasks
     call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)
13
14
15
     ! get my rank
16
     call MPI COMM RANK (MPI COMM WORLD, rank, ierr)
17
18
     ! this one is obvious
     call MPI_GET_PROCESSOR_NAME(hostname, len, ierr)
19
20
               'Number of tasks=',numtasks,' My rank=',rank,' Running on=',hostname
21
22
23
           ! do some work with message passing
24
25
26
     ! done with MPI
27
     call MPI_FINALIZE(ierr)
28
29
     end
```

MPI Exercise 1

Getting Started

Overview:

- Login to an LC cluster using your workshop username and OTP token
- Copy the exercise files to your home directory

- Familiarize yourself with LC's MPI compilers
- Write a simple "Hello World" MPI program using several MPI Environment Management routines
- · Successfully compile your program
- Successfully run your program several different ways



GO TO THE EXERCISE HERE

Approx. 20 minutes

Point to Point Communication Routines

General Concepts

First, a Simple Example:

- The value of PI can be calculated in various ways. Consider the Monte Carlo method of approximating PI:
 - Inscribe a circle with radius **r** in a square with side length of **2r**
 - The area of the circle is Πr^2 and the area of the square is $4r^2$
 - The ratio of the area of the circle to the area of the square is: $\Pi r^2 / 4r^2 = \Pi / 4$
 - If you randomly generate N points inside the square, approximately
 - $\mathbf{N} * \mathbf{\Pi} / \mathbf{4}$ of those points (**M**) should fall inside the circle.
 - Π is then approximated as:
 - $N * \Pi / 4 = M$
 - $\Pi/4 = M/N$
 - $\Pi = 4 * M / N$
 - Note that increasing the number of points generated improves the approximation.
- · Serial pseudo code for this procedure:

```
npoints = 10000
circle_count = 0

do j = 1,npoints
  generate 2 random numbers between 0 and 1
  xcoordinate = random1
  ycoordinate = random2
  if (xcoordinate, ycoordinate) inside circle
  then circle_count = circle_count + 1
end do

PI = 4.0*circle count/npoints
```

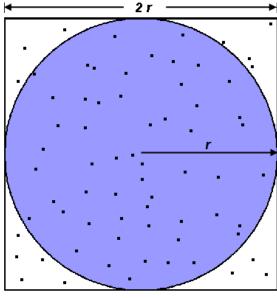
- Leads to an "embarassingly parallel" solution:
 - Break the loop iterations into chunks that can be executed by different tasks simultaneously.
 - Each task executes its portion of the loop a number of times.
 - Each task can do its work without requiring any information from the other tasks (there are no data dependencies).
 - Master task recieves results from other tasks using send/receive point-to-point operations.
- Pseudo code solution: red highlights changes for parallelism.

```
npoints = 10000
circle_count = 0

p = number of tasks
num = npoints/p

find out if I am MASTER or WORKER

do j = 1,num
```



$$A_S = (2r)^2 = 4r^2$$

$$A_C = \pi r^2$$

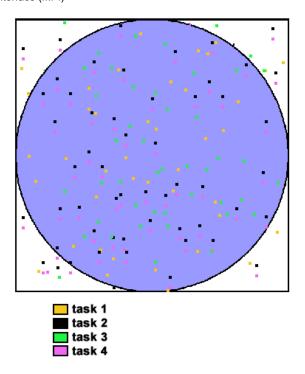
$$\pi = 4 \times \frac{A_C}{A_S}$$

```
generate 2 random numbers between 0 and 1
xcoordinate = random1
ycoordinate = random2
if (xcoordinate, ycoordinate) inside circle
then circle_count = circle_count + 1
end do

if I am MASTER
receive from WORKERS their circle_counts
compute PI (use MASTER and WORKER calculations)
else if I am WORKER
send to MASTER circle_count
endif

Example MPI Program in C: mpi_pi_reduce.c
Example MPI Program in Fortran: mpi_pi_reduce.f
```

 Key Concept: Divide work between available tasks which communicate data via point-to-point message passing calls.

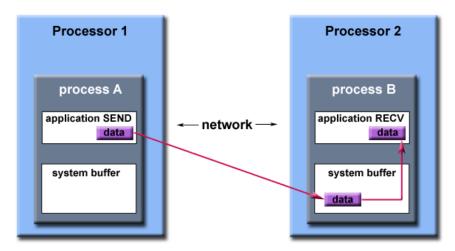


Types of Point-to-Point Operations:

- MPI point-to-point operations typically involve message passing between two, and only two, different MPI tasks. One task is
 performing a send operation and the other task is performing a matching receive operation.
- There are different types of send and receive routines used for different purposes. For example:
 - Synchronous send
 - Blocking send / blocking receive
 - Non-blocking send / non-blocking receive
 - o Buffered send
 - o Combined send/receive
 - o "Ready" send
- Any type of send routine can be paired with any type of receive routine.
- MPI also provides several routines associated with send receive operations, such as those used to wait for a message's arrival or probe to find out if a message has arrived.

Buffering:

- In a perfect world, every send operation would be perfectly synchronized with its matching receive. This is rarely the case. Somehow or other, the MPI implementation must be able to deal with storing data when the two tasks are out of sync.
- · Consider the following two cases:
 - o A send operation occurs 5 seconds before the receive is ready where is the message while the receive is pending?
 - Multiple sends arrive at the same receiving task which can only accept one send at a time what happens to the messages
 that are "backing up"?
- The MPI implementation (not the MPI standard) decides what happens to data in these types of cases. Typically, a **system buffer** area is reserved to hold data in transit. For example:



Path of a message buffered at the receiving process

- · System buffer space is:
 - Opaque to the programmer and managed entirely by the MPI library
 - A finite resource that can be easy to exhaust
 - Often mysterious and not well documented
 - o Able to exist on the sending side, the receiving side, or both
 - o Something that may improve program performance because it allows send receive operations to be asynchronous.
- User managed address space (i.e. your program variables) is called the application buffer. MPI also provides for a user managed send buffer.

Blocking vs. Non-blocking:

• Most of the MPI point-to-point routines can be used in either blocking or non-blocking mode.

Blocking:

- A blocking send routine will only "return" after it is safe to modify the application buffer (your send data) for reuse. Safe
 means that modifications will not affect the data intended for the receive task. Safe does not imply that the data was actually
 received it may very well be sitting in a system buffer.
- A blocking send can be synchronous which means there is handshaking occurring with the receive task to confirm a safe send.
- o A blocking send can be asynchronous if a system buffer is used to hold the data for eventual delivery to the receive.
- o A blocking receive only "returns" after the data has arrived and is ready for use by the program.

· Non-blocking:

- Non-blocking send and receive routines behave similarly they will return almost immediately. They do not wait for any
 communication events to complete, such as message copying from user memory to system buffer space or the actual
 arrival of message.
- Non-blocking operations simply "request" the MPI library to perform the operation when it is able. The user can not predict
 when that will happen.
- It is unsafe to modify the application buffer (your variable space) until you know for a fact the requested non-blocking
 operation was actually performed by the library. There are "wait" routines used to do this.
- Non-blocking communications are primarily used to overlap computation with communication and exploit possible performance gains.

Blocking Send	Non-blocking Send
myvar = 0;	myvar = 0;
<pre>for (i=1; i<ntasks; (&myvar="" *="" +="" 2="" do="" i++)="" mpi_send="" myvar="myvar" pre="" some="" task="i;" task);="" work="" {="" }<=""></ntasks;></pre>	<pre>for (i=1; i<ntasks; (&myvar="" ();="" *="" +="" 2;="" do="" i++)="" mpi_isend="" mpi_wait="" myvar="myvar" pre="" some="" task="i;" task);="" work="" {="" }<=""></ntasks;></pre>
Safe. Why?	Unsafe. Why?

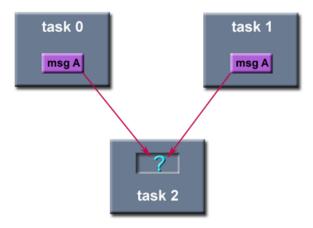
Order and Fairness:

• Order:

- MPI guarantees that messages will not overtake each other.
- If a sender sends two messages (Message 1 and Message 2) in succession to the same destination, and both match the same receive, the receive operation will receive Message 1 before Message 2.
- If a receiver posts two receives (Receive 1 and Receive 2), in succession, and both are looking for the same message, Receive 1 will receive the message before Receive 2.
- o Order rules do not apply if there are multiple threads participating in the communication operations.

• Fairness:

- MPI does not guarantee fairness it's up to the programmer to prevent "operation starvation".
- Example: task 0 sends a message to task 2. However, task 1 sends a competing message that matches task 2's receive.
 Only one of the sends will complete.



Point to Point Communication Routines

MPI Message Passing Routine Arguments

MPI point-to-point communication routines generally have an argument list that takes one of the following formats:

Blocking sends	MPI_Send(buffer,count,type,dest,tag,comm)	
Non-blocking sends	MPI_Isend(buffer,count,type,dest,tag,comm,request)	
Blocking receive	MPI_Recv(buffer,count,type,source,tag,comm,status)	
Non-blocking receive	MPI_Irecv(buffer,count,type,source,tag,comm,request)	

Buffer

Program (application) address space that references the data that is to be sent or received. In most cases, this is simply the variable name that is be sent/received. For C programs, this argument is passed by reference and usually must be prepended with an ampersand: &var1

Data Count

Indicates the number of data elements of a particular type to be sent.

Data Type

For reasons of portability, MPI predefines its elementary data types. The table below lists those required by the standard.

C Data Types		Fortran D	Fortran Data Types	
MPI_CHAR	char	MPI_CHARACTER	character(1)	
MPI_WCHAR	wchar_t - wide character			
MPI_SHORT	signed short int			
MPI_INT	signed int	MPI_INTEGER MPI_INTEGER1 MPI_INTEGER2 MPI_INTEGER4	integer integer*1 integer*2 integer*4	
MPI_LONG	signed long int			
MPI_LONG_LONG_INT MPI_LONG_LONG	signed long long int			
MPI_SIGNED_CHAR	signed char			
MPI_UNSIGNED_CHAR	unsigned char			
MPI_UNSIGNED_SHORT	unsigned short int			
MPI_UNSIGNED	unsigned int			
MPI_UNSIGNED_LONG	unsigned long int			
MPI_UNSIGNED_LONG_LONG	unsigned long long int			
MPI_FLOAT	float	MPI_REAL MPI_REAL2 MPI_REAL4 MPI_REAL8	real real*2 real*4 real*8	

MPI_DOUBLE	double	MPI_DOUBLE_PRECISION	double precision
MPI_LONG_DOUBLE	long double		
MPI_C_COMPLEX MPI_C_FLOAT_COMPLEX	float _Complex	MPI_COMPLEX	complex
MPI_C_DOUBLE_COMPLEX	double _Complex	MPI_DOUBLE_COMPLEX	double complex
MPI_C_LONG_DOUBLE_COMPLEX	long double _Complex		
MPI_C_BOOL	_Bool	MPI_LOGICAL	logical
MPI_INT8_T MPI_INT16_T MPI_INT32_T MPI_INT64_T	int8_t int16_t int32_t int64_t		
MPI_UINT8_T MPI_UINT16_T MPI_UINT32_T MPI_UINT64_T	uint8_t uint16_t uint32_t uint64_t		
MPI_BYTE	8 binary digits	MPI_BYTE	8 binary digits
MPI_PACKED	data packed or unpacked with MPI_Pack()/ MPI_Unpack	MPI_PACKED	data packed or unpacked with MPI_Pack()/ MPI_Unpack

Notes:

- Programmers may also create their own data types (see <u>Derived Data Types</u>).
- MPI_BYTE and MPI_PACKED do not correspond to standard C or Fortran types.
- Types shown in **GRAY FONT** are recommended if possible.
- Some implementations may include additional elementary data types (MPI_LOGICAL2, MPI_COMPLEX32, etc.). Check the MPI header file.

Destination

An argument to send routines that indicates the process where a message should be delivered. Specified as the rank of the receiving process.

Source

An argument to receive routines that indicates the originating process of the message. Specified as the rank of the sending process. This may be set to the wild card MPI_ANY_SOURCE to receive a message from any task.

Tag

Arbitrary non-negative integer assigned by the programmer to uniquely identify a message. Send and receive operations should match message tags. For a receive operation, the wild card MPI_ANY_TAG can be used to receive any message regardless of its tag. The MPI standard guarantees that integers 0-32767 can be used as tags, but most implementations allow a much larger range than this.

Communicator

Indicates the communication context, or set of processes for which the source or destination fields are valid. Unless the programmer is explicitly creating new communicators, the predefined communicator MPI_COMM_WORLD is usually used.

Status

For a receive operation, indicates the source of the message and the tag of the message. In C, this argument is a pointer to a predefined structure MPI_Status (ex. stat.MPI_SOURCE stat.MPI_TAG). In Fortran, it is an integer array of size MPI_STATUS_SIZE (ex. stat(MPI_SOURCE) stat(MPI_TAG)). Additionally, the actual number of bytes received is obtainable from Status via the MPI_Get_count routine. The constants MPI_STATUS_IGNORE and MPI_STATUSES_IGNORE can be substituted if a message's source, tag or size will be be queried later.

Request

Used by non-blocking send and receive operations. Since non-blocking operations may return before the requested system buffer space is obtained, the system issues a unique "request number". The programmer uses this system assigned "handle" later (in a WAIT type routine) to determine completion of the non-blocking operation. In C, this argument is a pointer to a predefined structure MPI Request. In Fortran, it is an integer.

Point to Point Communication Routines

Blocking Message Passing Routines

The more commonly used MPI blocking message passing routines are described below.

MPI Send

Basic blocking send operation. Routine returns only after the application buffer in the sending task is free for reuse. Note that this routine may be implemented differently on different systems. The MPI standard permits the use of a system buffer but does not require it. Some implementations may actually use a synchronous send (discussed below) to implement the basic blocking send.

```
MPI_Send (&buf,count,datatype,dest,tag,comm)
MPI_SEND (buf,count,datatype,dest,tag,comm,ierr)
```

MPI Recv

Receive a message and block until the requested data is available in the application buffer in the receiving task.

```
MPI_Recv (&buf,count,datatype,source,tag,comm,&status)
MPI_RECV (buf,count,datatype,source,tag,comm,status,ierr)
```

MPI Ssend

Synchronous blocking send: Send a message and block until the application buffer in the sending task is free for reuse and the destination process has started to receive the message.

```
MPI_Ssend (&buf,count,datatype,dest,tag,comm)
MPI_SSEND (buf,count,datatype,dest,tag,comm,ierr)
```

MPI Sendrecy

Send a message and post a receive before blocking. Will block until the sending application buffer is free for reuse and until the receiving application buffer contains the received message.

```
MPI_Sendrecv (&sendbuf, sendcount, sendtype, dest, sendtag,
..... &recvbuf, recvcount, recvtype, source, recvtag,
..... comm, &status)

MPI_SENDRECV (sendbuf, sendcount, sendtype, dest, sendtag,
..... recvbuf, recvcount, recvtype, source, recvtag,
..... comm, status, ierr)
```

MPI_Wait
MPI_Waitany
MPI_Waitall
MPI_Waitsome

MPI_Wait blocks until a specified non-blocking send or receive operation has completed. For multiple non-blocking operations, the programmer can specify any, all or some completions.

```
MPI_Wait (&request, &status)
MPI_Waitany (count, &array_of_requests, &index, &status)
MPI_Waitall (count, &array_of_requests, &array_of_statuses)
MPI_Waitsome (incount, &array_of_requests, &outcount,
..... &array_of_offsets, &array_of_statuses)
MPI_WAIT (request, status, ierr)
MPI_WAITANY (count, array_of_requests, index, status, ierr)
MPI_WAITALL (count, array_of_requests, array_of_statuses,
..... ierr)
MPI_WAITSOME (incount, array_of_requests, outcount,
..... array_of_offsets, array_of_statuses, ierr)
```

MPI_Probe

Performs a blocking test for a message. The "wildcards" MPI_ANY_SOURCE and MPI_ANY_TAG may be used to test for a message from any source or with any tag. For the C routine, the actual source and tag will be returned in the status structure as status.MPI_SOURCE and status.MPI_TAG. For the Fortran routine, they will be returned in the integer array status (MPI_SOURCE) and status (MPI_TAG). Example Use

```
MPI_Probe (source,tag,comm,&status)
MPI_PROBE (source,tag,comm,status,ierr)
```

MPI_Get_count

Returns the source, tag and number of elements of datatype received. Can be used with both blocking and non-blocking receive operations. For the C routine, the actual source and tag will be returned in the status structure as status.MPI_SOURCE and status.MPI_TAG. For the Fortran routine, they will be returned in the integer array status (MPI_SOURCE) and status (MPI_TAG).

```
MPI_Get_count (&status,datatype,&count)
MPI_GET_COUNT (status,datatype,count,ierr)
```

Examples: Blocking Message Passing Routines

Task 0 pings task 1 and awaits return ping

```
C Language - Blocking Message Passing Example
      #include "mpi.h"
 1
 2
      #include <stdio.h>
 3
      main(int argc, char *argv[])
 5
      int numtasks, rank, dest, source, rc, count, tag=1;
 6
      char inmsg, outmsg='x';
 7
                        // required variable for receive routines
      MPI Status Stat;
 8
 9
      MPI_Init(&argc, &argv);
10
      MPI_Comm_size(MPI_COMM_WORLD, &numtasks);
11
      MPI Comm rank(MPI COMM WORLD, &rank);
12
13
      // task 0 sends to task 1 and waits to receive a return message
      if (rank == 0) {
14
        dest = 1;
15
16
        source = 1;
        MPI_Send(&outmsg, 1, MPI_CHAR, dest, tag, MPI COMM WORLD);
17
        MPI Recv(&inmsg, 1, MPI CHAR, source, tag, MPI COMM WORLD, &Stat);
18
19
20
21
      // task 1 waits for task 0 message then returns a message
22
      else if (rank == 1) {
23
        dest = 0;
24
        source = 0;
        MPI_Recv(&inmsg, 1, MPI_CHAR, source, tag, MPI_COMM_WORLD, &Stat);
25
26
        MPI Send(&outmsg, 1, MPI CHAR, dest, tag, MPI COMM WORLD);
27
28
      // query recieve Stat variable and print message details
29
30
      MPI Get count(&Stat, MPI CHAR, &count);
31
      printf("Task %d: Received %d char(s) from task %d with tag %d n",
32
             rank, count, Stat.MPI_SOURCE, Stat.MPI_TAG);
33
34
      MPI Finalize();
35
```

```
Fortran - Blocking Message Passing Example
 1
     program ping
 2
     include 'mpif.h'
 3
 4
     integer numtasks, rank, dest, source, count, tag, ierr
 5
     integer stat (MPI_STATUS_SIZE) ! required variable for receive routines
 6
     character inmsg, outmsg
 7
     outmsg = 'x'
 8
     tag = 1
 9
10
     call MPI_INIT(ierr)
     call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
11
12
     call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)
13
```

```
! task 0 sends to task 1 and waits to receive a return message
15
     if (rank .eq. 0) then
        dest = 1
16
17
        source = 1
        call MPI SEND (outmsq, 1, MPI CHARACTER, dest, taq, MPI COMM WORLD, ierr)
18
        call MPI RECV (inmsg, 1, MPI CHARACTER, source, tag, MPI COMM WORLD, stat, ierr)
19
20
21
     ! task 1 waits for task 0 message then returns a message
     else if (rank .eq. 1) then
22
23
        dest = 0
24
        source = 0
25
        call MPI RECV(inmsg, 1, MPI CHARACTER, source, tag, MPI COMM WORLD, stat, err)
        call MPI SEND (outmsg, 1, MPI CHARACTER, dest, tag, MPI COMM WORLD, err)
26
27
     endif
28
29
     ! query recieve Stat variable and print message details
     call MPI GET COUNT(stat, MPI CHARACTER, count, ierr)
30
31
              'Task ',rank,': Received', count, 'char(s) from task', &
32
              stat(MPI SOURCE), 'with tag',stat(MPI TAG)
33
34
     call MPI FINALIZE (ierr)
35
36
     end
```

Point to Point Communication Routines

Non-blocking Message Passing Routines

The more commonly used MPI non-blocking message passing routines are described below.

MPI Isend

Identifies an area in memory to serve as a send buffer. Processing continues immediately without waiting for the message to be copied out from the application buffer. A communication request handle is returned for handling the pending message status. The program should not modify the application buffer until subsequent calls to MPI_Wait or MPI_Test indicate that the non-blocking send has completed.

```
MPI_Isend (&buf,count,datatype,dest,tag,comm,&request)
MPI_ISEND (buf,count,datatype,dest,tag,comm,request,ierr)
```

MPI Irecv

Identifies an area in memory to serve as a receive buffer. Processing continues immediately without actually waiting for the message to be received and copied into the the application buffer. A communication request handle is returned for handling the pending message status. The program must use calls to MPI_Wait or MPI_Test to determine when the non-blocking receive operation completes and the requested message is available in the application buffer.

```
MPI_Irecv (&buf,count,datatype,source,tag,comm,&request)
MPI_IRECV (buf,count,datatype,source,tag,comm,request,ierr)
```

MPI Issend

Non-blocking synchronous send. Similar to MPI_Isend(), except MPI_Wait() or MPI_Test() indicates when the destination process has received the message.

```
MPI_Issend (&buf,count,datatype,dest,tag,comm,&request)
MPI_ISSEND (buf,count,datatype,dest,tag,comm,request,ierr)
```

```
MPI_Test
MPI_Testany
MPI_Testall
MPI_Testsome
```

MPI_Test checks the status of a specified non-blocking send or receive operation. The "flag" parameter is returned logical true (1) if the operation has completed, and logical false (0) if not. For multiple non-blocking operations, the programmer can specify any, all or some completions.

```
MPI_Test (&request,&flag,&status)
MPI_Testany (count,&array_of_requests,&index,&flag,&status)
```

```
MPI_Testall (count,&array_of_requests,&flag,&array_of_statuses)
MPI_Testsome (incount,&array_of_requests,&outcount,
..... &array_of_offsets, &array_of_statuses)
MPI_TEST (request,flag,status,ierr)
MPI_TESTANY (count,array_of_requests,index,flag,status,ierr)
MPI_TESTALL (count,array_of_requests,flag,array_of_statuses,ierr)
MPI_TESTSOME (incount,array_of_requests,outcount,
..... array_of_offsets, array_of_statuses,ierr)
```

MPI_Iprobe

Performs a non-blocking test for a message. The "wildcards" MPI_ANY_SOURCE and MPI_ANY_TAG may be used to test for a message from any source or with any tag. The integer "flag" parameter is returned logical true (1) if a message has arrived, and logical false (0) if not. For the C routine, the actual source and tag will be returned in the status structure as status.MPI_TAG. For the Fortran routine, they will be returned in the integer array status (MPI_SOURCE) and status (MPI_TAG).

```
MPI_Iprobe (source,tag,comm,&flag,&status)
MPI_IPROBE (source,tag,comm,flag,status,ierr)
```

Examples: Non-blocking Message Passing Routines

Nearest neighbor exchange in a ring topology



```
₽
     C Language - Non-blocking Message Passing Example
      #include "mpi.h"
 2
      #include <stdio.h>
 3
 4
      main(int argc, char *argv[]) {
 5
      int numtasks, rank, next, prev, buf[2], tag1=1, tag2=2;
 6
      MPI Request reqs[4];
                            // required variable for non-blocking calls
 7
      MPI Status stats[4];
                              // required variable for Waitall routine
 8
 9
      MPI_Init(&argc, &argv);
10
      MPI Comm size(MPI COMM WORLD, &numtasks);
11
      MPI_Comm_rank(MPI_COMM_WORLD, &rank);
12
      // determine left and right neighbors
13
14
      prev = rank-1;
15
      next = rank+1;
16
      if (rank == 0) prev = numtasks - 1;
      if (rank == (numtasks - 1)) next = 0;
17
18
19
      // post non-blocking receives and sends for neighbors
20
      MPI_Irecv(&buf[0], 1, MPI_INT, prev, tag1, MPI_COMM_WORLD, &reqs[0]);
21
      MPI Irecv(&buf[1], 1, MPI INT, next, tag2, MPI COMM WORLD, &reqs[1]);
22
23
      MPI_Isend(&rank, 1, MPI_INT, prev, tag2, MPI_COMM_WORLD, &reqs[2]);
24
      MPI_Isend(&rank, 1, MPI_INT, next, tag1, MPI_COMM_WORLD, &reqs[3]);
25
         // do some work while sends/receives progress in background
26
27
28
      // wait for all non-blocking operations to complete
29
      MPI Waitall(4, reqs, stats);
30
31
         // continue - do more work
32
33
      MPI_Finalize();
34
      }
```

```
Fortran - Non-blocking Message Passing Example
     program ringtopo
 1
 2
     include 'mpif.h'
 3
     integer numtasks, rank, next, prev, buf(2), tag1, tag2, ierr
 5
     integer reqs(4) ! required variable for non-blocking calls
 6
     integer stats(MPI_STATUS_SIZE,4)
                                          ! required variable for WAITALL routine
 7
     tag1 = 1
 8
     tag2 = 2
 9
10
     call MPI_INIT(ierr)
11
     call MPI COMM RANK (MPI COMM WORLD, rank, ierr)
     call MPI_COMM_SIZE (MPI_COMM_WORLD, numtasks, ierr)
12
13
     ! determine left and right neighbors
14
15
     prev = rank - 1
     next = rank + 1
16
17
     if (rank .eq. 0) then
        prev = numtasks - 1
18
19
     endif
20
     if (rank .eq. numtasks - 1) then
21
        next = 0
22
     endif
23
24
     ! post non-blocking receives and sends for neighbors
     call MPI_IRECV(buf(1), 1, MPI_INTEGER, prev, tag1, MPI_COMM_WORLD, reqs(1), ierr)
25
     call MPI IRECV(buf(2), 1, MPI INTEGER, next, tag2, MPI COMM WORLD, reqs(2), ierr)
26
27
     call MPI_ISEND(rank, 1, MPI_INTEGER, prev, tag2, MPI_COMM_WORLD, reqs(3), ierr)
call MPI_ISEND(rank, 1, MPI_INTEGER, next, tag1, MPI_COMM_WORLD, reqs(4), ierr)
28
29
30
31
         ! do some work while sends/receives progress in background
32
33
     ! wait for all non-blocking operations to complete
34
     call MPI_WAITALL(4, reqs, stats, ierr);
35
36
         ! continue - do more work
37
     call MPI FINALIZE(ierr)
38
39
40
     end
```

MPI Exercise 2

Point-to-Point Message Passing

Overview:

- . Login to the LC workshop cluster, if you are not already logged in
- Using your "Hello World" MPI program from Exercise 1, add MPI blocking point-to-point routines to send and receive messages
- · Successfully compile your program
- Successfully run your program several different ways
- Try the same thing with nonblocking send/receive routines



GO TO THE EXERCISE HERE

Approx. 20 minutes

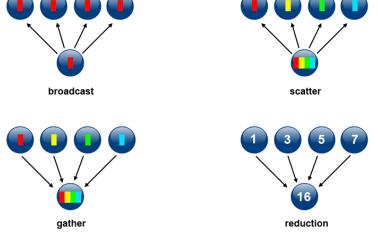
Collective Communication Routines

Types of Collective Operations:

- Synchronization processes wait until all members of the group have reached the synchronization point.
- Data Movement broadcast, scatter/gather, all to all.
- Collective Computation (reductions) one member of the group collects data from the other members and performs an operation (min, max, add, multiply, etc.) on that data.

Scope:

- Collective communication routines must involve all processes within the scope of a communicator.
 - All processes are by default, members in the communicator MPI_COMM_WORLD.
 - Additional communicators can be defined by the programmer. See the <u>Group and Communicator Management Routines</u> section for details.



- Unexpected behavior, including program failure, can occur if even one task in the communicator doesn't participate.
- It is the programmer's responsibility to ensure that all processes within a communicator participate in any collective operations.

Programming Considerations and Restrictions:

- · Collective communication routines do not take message tag arguments.
- Collective operations within subsets of processes are accomplished by first partitioning the subsets into new groups and then
 attaching the new groups to new communicators (discussed in the <u>Group and Communicator Management Routines</u> section).
- Can only be used with MPI predefined datatypes not with MPI Derived Data Types.
- MPI-2 extended most collective operations to allow data movement between intercommunicators (not covered here).
- With MPI-3, collective operations can be blocking or non-blocking. Only blocking operations are covered in this tutorial.

Collective Communication Routines

MPI_Barrier

Synchronization operation. Creates a barrier synchronization in a group. Each task, when reaching the MPI_Barrier call, blocks until all tasks in the group reach the same MPI_Barrier call. Then all tasks are free to proceed.

```
MPI_Barrier (comm)
MPI_BARRIER (comm,ierr)
```

MPI Bcast

Data movement operation. Broadcasts (sends) a message from the process with rank "root" to all other processes in the group.

Diagram Here

```
MPI_Bcast (&buffer,count,datatype,root,comm)
MPI_BCAST (buffer,count,datatype,root,comm,ierr)
```

MPI Scatter

Data movement operation. Distributes distinct messages from a single source task to each task in the group.

Diagram Here

```
MPI_Scatter (&sendbuf,sendcnt,sendtype,&recvbuf,
..... recvcnt,recvtype,root,comm)
MPI_SCATTER (sendbuf,sendcnt,sendtype,recvbuf,
..... recvcnt,recvtype,root,comm,ierr)
```

MPI_Gather

Data movement operation. Gathers distinct messages from each task in the group to a single destination task. This routine is the reverse operation of MPI_Scatter.

Diagram Here

```
MPI_Gather (&sendbuf,sendcnt,sendtype,&recvbuf,
..... recvcount,recvtype,root,comm)
MPI_GATHER (sendbuf,sendcnt,sendtype,recvbuf,
..... recvcount,recvtype,root,comm,ierr)
```

MPI Allgather

Data movement operation. Concatenation of data to all tasks in a group. Each task in the group, in effect, performs a one-to-all broadcasting operation within the group.

Diagram Here

```
MPI_Allgather (&sendbuf,sendcount,sendtype,&recvbuf,
..... recvcount,recvtype,comm)
MPI_ALLGATHER (sendbuf,sendcount,sendtype,recvbuf,
..... recvcount,recvtype,comm,info)
```

MPI Reduce

Collective computation operation. Applies a reduction operation on all tasks in the group and places the result in one task.

Diagram Here

```
MPI_Reduce (&sendbuf,&recvbuf,count,datatype,op,root,comm)
MPI_REDUCE (sendbuf,recvbuf,count,datatype,op,root,comm,ierr)
```

The predefined MPI reduction operations appear below. Users can also define their own reduction functions by using the MPI Op create routine.

MPI Reduction Operation		C Data Types	Fortran Data Type
MPI_MAX	maximum	integer, float	integer, real, complex
MPI_MIN	minimum	integer, float	integer, real, complex
MPI_SUM	sum	integer, float	integer, real, complex
MPI_PROD	product	integer, float	integer, real, complex
MPI_LAND	logical AND	integer	logical
MPI_BAND	bit-wise AND	integer, MPI_BYTE	integer, MPI_BYTE
MPI_LOR	logical OR	integer	logical
MPI_BOR	bit-wise OR	integer, MPI_BYTE	integer, MPI_BYTE
MPI_LXOR	logical XOR	integer	logical
MPI_BXOR	bit-wise XOR	integer, MPI_BYTE	integer, MPI_BYTE
MPI_MAXLOC	max value and location	float, double and long double	real, complex,double precision
MPI_MINLOC	min value and location	float, double and long double	real, complex, double precision

• Note from the MPI_Reduce man page: The operation is always assumed to be associative. All predefined operations are also assumed to be commutative. Users may define operations that are assumed to be associative, but not commutative. The "canonical" evaluation order of a reduction is determined by the ranks of the processes in the group. However, the implementation can take advantage of associativity, or associativity and commutativity in order to change the order of evaluation. This may change the result of the reduction for operations that are not strictly associative and commutative, such as floating point addition. [Advice to implementors] It is strongly recommended that MPI_REDUCE be implemented so that the same result be obtained whenever the function is applied on the same arguments, appearing in the same order. Note that this may prevent optimizations that take advantage of the physical location of processors. [End of advice to implementors]

MPI_Allreduce

Collective computation operation + data movement. Applies a reduction operation and places the result in all tasks in the group. This is equivalent to an MPI_Reduce followed by an MPI_Bcast.

Diagram Here

```
MPI_Allreduce (&sendbuf,&recvbuf,count,datatype,op,comm)
MPI_ALLREDUCE (sendbuf,recvbuf,count,datatype,op,comm,ierr)
```

MPI Reduce scatter

Collective computation operation + data movement. First does an element-wise reduction on a vector across all tasks in the group. Next, the result vector is split into disjoint segments and distributed across the tasks. This is equivalent to an MPI_Reduce followed by an MPI_Scatter operation.

Diagram Here

MPI_Alltoall

Data movement operation. Each task in a group performs a scatter operation, sending a distinct message to all the tasks in the group in order by index.

Diagram Here

```
MPI_Alltoall (&sendbuf,sendcount,sendtype,&recvbuf,
..... recvcnt,recvtype,comm)
MPI_ALLTOALL (sendbuf,sendcount,sendtype,recvbuf,
..... recvcnt,recvtype,comm,ierr)
```

MPI_Scan

Performs a scan operation with respect to a reduction operation across a task group.

Diagram Here

```
MPI_Scan (&sendbuf,&recvbuf,count,datatype,op,comm)
MPI_SCAN (sendbuf,recvbuf,count,datatype,op,comm,ierr)
```

Examples: Collective Communications

Perform a scatter operation on the rows of an array

```
C Language - Collective Communications Example
     #include "mpi.h"
 1
 2
     #include <stdio.h>
 3
     #define SIZE 4
 5
     main(int argc, char *argv[])
 6
     int numtasks, rank, sendcount, recvcount, source;
 7
     float sendbuf[SIZE][SIZE] = {
 8
       {1.0, 2.0, 3.0, 4.0},
 9
       {5.0, 6.0, 7.0, 8.0},
10
       \{9.0, 10.0, 11.0, 12.0\},\
11
       {13.0, 14.0, 15.0, 16.0} };
12
     float recvbuf[SIZE];
13
14
     MPI_Init(&argc,&argv);
     MPI Comm rank (MPI_COMM_WORLD, &rank);
15
     MPI Comm size(MPI COMM WORLD, &numtasks);
16
17
18
     if (numtasks == SIZE) {
19
       // define source task and elements to send/receive, then perform collective scatter
20
       source = 1;
21
       sendcount = SIZE;
22
       recvcount = SIZE;
23
       MPI Scatter (sendbuf, sendcount, MPI FLOAT, recvbuf, recvcount,
24
                    MPI FLOAT, source, MPI COMM WORLD);
25
26
       printf("rank= %d Results: %f %f %f %f\n",rank,recvbuf[0],
27
              recvbuf[1],recvbuf[2],recvbuf[3]);
28
       }
29
     else
30
       printf("Must specify %d processors. Terminating.\n",SIZE);
31
```

```
32 MPI_Finalize();
33 }
```

```
Fortran - Collective Communications Example
 1
     program scatter
 2
     include 'mpif.h'
 3
 4
     integer SIZE
 5
     parameter(SIZE=4)
     integer numtasks, rank, sendcount, recvcount, source, ierr
 7
     real*4 sendbuf(SIZE,SIZE), recvbuf(SIZE)
     ! Fortran stores this array in column major order, so the
 9
      ! scatter will actually scatter columns, not rows.
10
     data sendbuf /1.0, 2.0, 3.0, 4.0, & 5.0, 6.0, 7.0, 8.0, & 9.0, 10.0, 11.0, 12.0, &
11
12
13
14
                     13.0, 14.0, 15.0, 16.0 /
15
16
     call MPI_INIT(ierr)
     call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)
17
18
19
20
     if (numtasks .eq. SIZE) then
21
         ! define source task and elements to send/receive, then perform collective scatter
         source = 1
22
23
         sendcount = SIZE
24
         recvcount = SIZE
         call MPI_SCATTER(sendbuf, sendcount, MPI_REAL, recvbuf, recvcount, MPI_REAL, &
25
26
                            source, MPI_COMM_WORLD, ierr)
27
28
         print *, 'rank= ',rank,' Results: ',recvbuf
29
30
        print *, 'Must specify',SIZE,' processors. Terminating.'
31
32
     endif
33
34
     call MPI FINALIZE (ierr)
35
36
     end
```

```
rank= 0 Results: 1.000000 2.000000 3.000000 4.000000
rank= 1 Results: 5.000000 6.000000 7.000000 8.000000
rank= 2 Results: 9.000000 10.000000 11.000000 12.000000
rank= 3 Results: 13.000000 14.000000 15.000000 16.000000
```

Derived Data Types

• As previously mentioned, MPI predefines its primitive data types:

	C Data Types	Fortran Data Types
MPI_CHAR	MPI_C_COMPLEX	MPI_CHARACTER
MPI_WCHAR	MPI_C_FLOAT_COMPLEX	MPI_INTEGER
MPI SHORT	MPI C DOUBLE COMPLEX	MPI INTEGER1
MPI INT	MPI C LONG DOUBLE COMPLEX	MPI INTEGER2
MPI LONG	MPI C BOOL	MPI INTEGER4
MPI LONG LONG INT	MPI LOGICAL	MPI REAL
MPI LONG LONG	MPI C LONG DOUBLE COMPLEX	MPI REAL2
MPI SIGNED CHAR	MPI INT8 T	MPI REAL4
MPI UNSIGNED CHAR	MPI INT16 T	MPI REAL8
MPI_UNSIGNED_SHORT	MPI_INT32_T	MPI_DOUBLE_PRECISION

```
MPI UNSIGNED LONG
                           MPI INT64 T
                                                                        MPI COMPLEX
MPI UNSIGNED
                           MPI UINT8 T
                                                                        MPI DOUBLE COMPLEX
MPI FLOAT
                           MPI UINT16 T
                                                                        MPI LOGICAL
MPI_DOUBLE
                           MPI UINT32 T
                                                                       MPI BYTE
MPI LONG DOUBLE
                           MPI_UINT64_T
                                                                        MPI PACKED
                           MPI BYTE
                           MPI PACKED
```

- MPI also provides facilities for you to define your own data structures based upon sequences of the MPI primitive data types.
 Such user defined structures are called derived data types.
- Primitive data types are contiguous. Derived data types allow you to specify non-contiguous data in a convenient manner and to treat it as though it was contiguous.
- MPI provides several methods for constructing derived data types:
 - Contiguous
 - Vector
 - Indexed
 - o Struct

Derived Data Type Routines

MPI_Type_contiguous

The simplest constructor. Produces a new data type by making count copies of an existing data type.

```
MPI_Type_contiguous (count,oldtype,&newtype)
MPI_TYPE_CONTIGUOUS (count,oldtype,newtype,ierr)
```

MPI_Type_vector MPI Type hvector

Similar to contiguous, but allows for regular gaps (stride) in the displacements. MPI_Type_hvector is identical to MPI_Type_vector except that stride is specified in bytes.

```
MPI_Type_vector (count,blocklength,stride,oldtype,&newtype)
MPI_TYPE_VECTOR (count,blocklength,stride,oldtype,newtype,ierr)
```

MPI_Type_indexed MPI_Type_hindexed

An array of displacements of the input data type is provided as the map for the new data type. MPI_Type_hindexed is identical to MPI_Type_indexed except that offsets are specified in bytes.

```
MPI_Type_indexed (count,blocklens[],offsets[],old_type,&newtype)
MPI_TYPE_INDEXED (count,blocklens(),offsets(),old_type,newtype,ierr)
```

MPI_Type_struct

The new data type is formed according to completely defined map of the component data types. **NOTE:** This function is deprecated in MPI-2.0 and replaced by MPI_Type_create_struct in MPI-3.0

```
MPI_Type_struct (count,blocklens[],offsets[],old_types,&newtype)
MPI_TYPE_STRUCT (count,blocklens(),offsets(),old_types,newtype,ierr)
```

MPI_Type_extent

Returns the size in bytes of the specified data type. Useful for the MPI subroutines that require specification of offsets in bytes. **NOTE:** This function is deprecated in MPI-2.0 and replaced by MPI_Type_get_extent in MPI-3.0

```
MPI_Type_extent (datatype,&extent)
MPI_TYPE_EXTENT (datatype,extent,ierr)
```

MPI Type commit

Commits new datatype to the system. Required for all user constructed (derived) datatypes.

```
MPI_Type_commit (&datatype)
MPI_TYPE_COMMIT (datatype,ierr)
```

MPI Type free

Deallocates the specified datatype object. Use of this routine is especially important to prevent memory exhaustion if many datatype objects are created, as in a loop.

```
MPI_Type_free (&datatype)
MPI_TYPE_FREE (datatype,ierr)
```

Examples: Contiguous Derived Data Type

Create a data type representing a row of an array and distribute a different row to all processes.

Diagram Here

```
H
     C Language - Contiguous Derived Data Type Example
      #include "mpi.h"
 1
      #include <stdio.h>
 2
 3
      #define SIZE 4
 4
 5
      main(int argc, char *argv[])
 6
      int numtasks, rank, source=0, dest, tag=1, i;
 7
      float a[SIZE][SIZE] =
 8
        {1.0, 2.0, 3.0, 4.0,
 9
         5.0, 6.0, 7.0, 8.0,
10
         9.0, 10.0, 11.0, 12.0,
11
         13.0, 14.0, 15.0, 16.0};
12
      float b[SIZE];
13
14
      MPI Status stat;
15
                               // required variable
      MPI_Datatype rowtype;
16
17
      MPI Init(&argc,&argv);
18
      MPI Comm rank (MPI COMM WORLD, &rank);
19
      MPI_Comm_size(MPI_COMM_WORLD, &numtasks);
20
21
      // create contiguous derived data type
22
      MPI Type contiguous (SIZE, MPI FLOAT, &rowtype);
      MPI_Type_commit(&rowtype);
23
24
25
      if (numtasks == SIZE) {
26
          // task 0 sends one element of rowtype to all tasks
27
          if (rank == 0) {
28
             for (i=0; i<numtasks; i++)</pre>
29
               MPI Send(&a[i][0], 1, rowtype, i, tag, MPI COMM WORLD);
30
31
32
          // all tasks receive rowtype data from task 0
33
          MPI Recv(b, SIZE, MPI FLOAT, source, tag, MPI COMM WORLD, &stat);
         printf("rank= %d b= \frac{1}{8}3.1f %3.1f %3.1f \frac{1}{8}3.1f\n",
34
35
                 rank,b[0],b[1],b[2],b[3]);
36
         }
37
      else
         printf("Must specify %d processors. Terminating.\n",SIZE);
38
39
40
      // free datatype when done using it
41
      MPI_Type_free(&rowtype);
42
      MPI Finalize();
43
      }
```

```
Fortran - Contiguous Derived Data Type Example
```

1 program contiguous

```
4
      integer SIZE
      parameter (SIZE=4)
 6
      integer numtasks, rank, source, dest, tag, i, ierr
 7
      real*4 a(0:SIZE-1,0:SIZE-1), b(0:SIZE-1)
      integer stat(MPI STATUS SIZE)
 8
 9
      integer columntype
                          ! required variable
10
      tag = 1
11
12
       ! Fortran stores this array in column major order
      data a /1.0, 2.0, 3.0, 4.0, &
13
14
               5.0, 6.0, 7.0, 8.0, &
15
               9.0, 10.0, 11.0, 12.0, &
                13.0, 14.0, 15.0, 16.0 /
16
17
18
      call MPI INIT(ierr)
      call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
19
20
      call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)
21
      ! create contiguous derived data type
22
23
      call MPI_TYPE_CONTIGUOUS(SIZE, MPI_REAL, columntype, ierr)
      call MPI_TYPE_COMMIT(columntype, ierr)
24
25
26
      \quad \hbox{if (numtasks .eq. SIZE) then} \\
27
          ! task 0 sends one element of columntype to all tasks
28
         if (rank .eq. 0) then
29
            do i=0, numtasks-1
            call MPI SEND(a(0,i), 1, columntype, i, tag, MPI COMM WORLD, ierr)
30
31
            end do
         endif
32
33
34
         ! all tasks receive columntype data from task 0
35
         source = 0
36
         call MPI RECV(b, SIZE, MPI REAL, source, tag, MPI COMM WORLD, stat, ierr)
37
         print *, 'rank= ',rank,' b= ',b
38
      else
         print *, 'Must specify',SIZE,' processors. Terminating.'
39
40
      endif
41
42
      ! free datatype when done using it
      call MPI TYPE FREE(columntype, ierr)
43
      call MPI FINALIZE (ierr)
44
45
46
      end
```

```
rank= 0 b= 1.0 2.0 3.0 4.0

rank= 1 b= 5.0 6.0 7.0 8.0

rank= 2 b= 9.0 10.0 11.0 12.0

rank= 3 b= 13.0 14.0 15.0 16.0
```

Examples: Vector Derived Data Type

Create a data type representing a column of an array and distribute different columns to all processes.

Diagram Here

```
1
     C Language - Vector Derived Data Type Example
      #include "mpi.h"
1
      #include <stdio.h>
2
3
      #define SIZE 4
4
5
      main(int argc, char *argv[])
      int numtasks, rank, source=0, dest, tag=1, i;
6
7
      float a[SIZE][SIZE] =
        {1.0, 2.0, 3.0, 4.0,
```

```
5.0, 6.0, 7.0, 8.0,
9.0, 10.0, 11.0, 12.0,
10
        13.0, 14.0, 15.0, 16.0};
11
12
      float b[SIZE];
13
14
      MPI Status stat;
      MPI_Datatype columntype; // required variable
15
16
17
18
      MPI Init(&argc,&argv);
      MPI_Comm_rank(MPI_COMM_WORLD, &rank);
19
20
      MPI Comm size(MPI COMM WORLD, &numtasks);
21
22
      // create vector derived data type
      MPI Type vector (SIZE, 1, SIZE, MPI FLOAT, &columntype);
23
24
      MPI Type commit(&columntype);
25
26
      if (numtasks == SIZE) {
27
         // task 0 sends one element of columntype to all tasks
         if (rank == 0) {
28
29
            for (i=0; i<numtasks; i++)</pre>
30
               MPI_Send(&a[0][i], 1, columntype, i, tag, MPI_COMM_WORLD);
31
            }
32
         // all tasks receive columntype data from task 0
33
34
         MPI_Recv(b, SIZE, MPI_FLOAT, source, tag, MPI_COMM_WORLD, &stat);
35
         36
                rank,b[0],b[1],b[2],b[3]);
37
38
      else
39
         printf("Must specify %d processors. Terminating.\n",SIZE);
40
41
      // free datatype when done using it
42
      MPI Type free(&columntype);
43
      MPI Finalize();
44
      }
```

```
\blacksquare
     Fortran - Vector Derived Data Type Example
 1
      program vector
 2
      include 'mpif.h'
 3
 4
      integer SIZE
 5
      parameter (SIZE=4)
      integer numtasks, rank, source, dest, tag, i, ierr
 6
 7
      real*4 a(0:SIZE-1,0:SIZE-1), b(0:SIZE-1)
 8
      integer stat(MPI_STATUS_SIZE)
 9
      integer rowtype
                        ! required variable
10
      tag = 1
11
12
      ! Fortran stores this array in column major order
13
      data a /1.0, 2.0, 3.0, 4.0, &
14
                5.0, 6.0, 7.0, 8.0,
15
                9.0, 10.0, 11.0, 12.0, &
16
                13.0, 14.0, 15.0, 16.0 /
17
18
      call MPI INIT(ierr)
      call MPI COMM RANK (MPI COMM WORLD, rank, ierr)
19
20
      call MPI COMM SIZE (MPI COMM WORLD, numtasks, ierr)
21
22
      ! create vector derived data type
      call MPI TYPE VECTOR(SIZE, 1, SIZE, MPI REAL, rowtype, ierr)
23
24
      call MPI TYPE COMMIT (rowtype, ierr)
25
26
      if (numtasks .eq. SIZE) then
27
          ! task 0 sends one element of rowtype to all tasks
         if (rank .eq. 0) then
28
29
             do i=0, numtasks-1
30
             call MPI_SEND(a(i,0), 1, rowtype, i, tag, MPI_COMM_WORLD, ierr)
31
             end do
         endif
32
```

```
! all tasks receive rowtype data from task 0
35
         source = 0
         call MPI_RECV(b, SIZE, MPI_REAL, source, tag, MPI COMM WORLD, stat, ierr)
         print *, 'rank= ',rank,' b= ',b
37
38
      else
39
         print *, 'Must specify', SIZE,' processors. Terminating.'
40
      endif
41
42
      ! free datatype when done using it
43
      call MPI TYPE FREE(rowtype, ierr)
      call MPI_FINALIZE(ierr)
44
45
46
      end
```

```
rank= 0 b= 1.0 5.0 9.0 13.0

rank= 1 b= 2.0 6.0 10.0 14.0

rank= 2 b= 3.0 7.0 11.0 15.0

rank= 3 b= 4.0 8.0 12.0 16.0
```

Examples: Indexed Derived Data Type

Create a datatype by extracting variable portions of an array and distribute to all tasks.

Diagram Here

```
Ħ
      C Language - Indexed Derived Data Type Example
       #include "mpi.h"
 1
 2
       #include <stdio.h>
 3
       #define NELEMENTS 6
       main(int argc, char *argv[]) {
int numtasks, rank, source=0, dest, tag=1, i;
 5
 6
 7
       int blocklengths[2], displacements[2];
 8
       float a[16] =
         \{1.0, 2.0, 3.0, 4.0, 5.0, 6.0, 7.0, 8.0,
 9
10
          9.0, 10.0, 11.0, 12.0, 13.0, 14.0, 15.0, 16.0};
11
       float b[NELEMENTS];
12
13
       MPI Status stat;
      MPI_Datatype indextype; // required variable
14
15
16
       MPI_Init(&argc,&argv);
17
       MPI Comm rank (MPI COMM WORLD, &rank);
       MPI_Comm_size(MPI_COMM_WORLD, &numtasks);
18
19
20
       blocklengths[0] = 4;
21
      blocklengths[1] = 2;
22
       displacements[0] = 5;
23
       displacements[1] = 12;
24
25
       // create indexed derived data type
26
       MPI Type indexed(2, blocklengths, displacements, MPI FLOAT, &indextype);
27
       MPI_Type_commit(&indextype);
28
       if (rank == 0) {
29
30
         for (i=0; i<numtasks; i++)</pre>
31
          // task 0 sends one element of indextype to all tasks
32
            MPI_Send(a, 1, indextype, i, tag, MPI_COMM_WORLD);
33
34
35
       // all tasks receive indextype data from task 0
      MPI_Recv(b, NELEMENTS, MPI_FLOAT, source, tag, MPI_COMM_WORLD, &stat);
printf("rank= %d b= %3.1f %3.1f %3.1f %3.1f %3.1f %3.1f\n",
36
37
38
               rank,b[0],b[1],b[2],b[3],b[4],b[5]);
39
40
       // free datatype when done using it
```

```
41 MPI_Type_free(&indextype);
42 MPI_Finalize();
43 }
```

```
Fortran - Indexed Derived Data Type Example
1
      program indexed
 2
      include 'mpif.h'
 3
      integer NELEMENTS
 4
      parameter (NELEMENTS=6)
      integer numtasks, rank, source, dest, tag, i,
 6
                                                       ierr
      integer blocklengths(0:1), displacements(0:1)
 8
      real*4 a(0:15), b(0:NELEMENTS-1)
      integer stat(MPI STATUS SIZE)
10
      integer indextype
                         ! required variable
11
      tag = 1
12
      data a /1.0, 2.0, 3.0, 4.0, 5.0, 6.0, 7.0, 8.0, &
13
14
                9.0, 10.0, 11.0, 12.0, 13.0, 14.0, 15.0, 16.0 /
15
16
      call MPI INIT(ierr)
      call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
17
      call MPI COMM SIZE (MPI COMM WORLD, numtasks, ierr)
18
19
20
      blocklengths(0) = 4
21
      blocklengths(1) = 2
22
      displacements(0) = 5
23
      displacements(1) = 12
24
25
      ! create indexed derived data type
26
      call MPI_TYPE_INDEXED(2, blocklengths, displacements, MPI_REAL, &
27
                             indextype, ierr)
28
      call MPI_TYPE_COMMIT(indextype, ierr)
29
      if (rank .eq. 0) then
30
31
         ! task 0 sends one element of indextype to all tasks
32
         do i=0, numtasks-1
33
         call MPI SEND(a, 1, indextype, i, tag, MPI COMM WORLD, ierr)
34
         end do
35
      endif
36
      ! all tasks receive indextype data from task 0
37
38
      call MPI_RECV(b, NELEMENTS, MPI_REAL, source, tag, MPI_COMM_WORLD, &
39
40
                     stat, ierr)
      print *, 'rank= ',rank,' b= ',b
41
42
43
      ! free datatype when done using it
      call MPI TYPE FREE (indextype, ierr)
44
45
      call MPI_FINALIZE(ierr)
46
47
      end
```

```
rank= 0 b= 6.0 7.0 8.0 9.0 13.0 14.0 rank= 1 b= 6.0 7.0 8.0 9.0 13.0 14.0 rank= 2 b= 6.0 7.0 8.0 9.0 13.0 14.0 rank= 3 b= 6.0 7.0 8.0 9.0 13.0 14.0
```

Examples: Struct Derived Data Type

Create a data type that represents a particle and distribute an array of such particles to all processes.

```
₽
     C Language - Struct Derived Data Type Example
      #include "mpi.h"
 1
      #include <stdio.h>
 2
 3
      #define NELEM 25
 4
 5
      main(int argc, char *argv[])
      int numtasks, rank, source=0, dest, tag=1, i;
 6
 7
 8
      typedef struct {
 9
        float x, y, z;
10
        float velocity;
11
        int n, type;
12
                    Particle:
13
      Particle
                   p[NELEM], particles[NELEM];
      MPI Datatype particletype, oldtypes[2]; // required variables
14
15
      int
                   blockcounts[2];
16
17
      // MPI Aint type used to be consistent with syntax of
18
      // MPI Type extent routine
19
                  offsets[2], extent;
      MPI_Aint
20
21
      MPI_Status stat;
22
23
      MPI_Init(&argc,&argv);
      MPI Comm rank (MPI_COMM_WORLD, &rank);
24
25
      MPI Comm size(MPI COMM WORLD, &numtasks);
26
27
      // setup description of the 4 MPI FLOAT fields x, y, z, velocity
28
      offsets[0] = 0;
      oldtypes[0] = MPI FLOAT;
29
30
      blockcounts[0] = \overline{4};
31
32
      // setup description of the 2 MPI_INT fields n, type
33
      // need to first figure offset by getting size of MPI_FLOAT
      MPI_Type_extent(MPI_FLOAT, &extent);
34
35
      offsets[1] = 4 * extent;
      oldtypes[1] = MPI INT;
36
37
      blockcounts[1] = \overline{2};
38
39
      // define structured type and commit it
      MPI Type struct(2, blockcounts, offsets, oldtypes, &particletype);
40
41
      MPI Type commit(&particletype);
42
43
      // task 0 initializes the particle array and then sends it to each task
      if (rank == 0) {
44
        for (i=0; i<NELEM; i++) {
45
46
           particles[i].x = i * 1.0;
           particles[i].y = i * -1.0;
47
48
           particles[i].z = i * 1.0;
49
           particles[i].velocity = 0.25;
50
           particles[i].n = i;
51
           particles[i].type = i % 2;
52
53
        for (i=0; i<numtasks; i++)</pre>
54
           MPI_Send(particles, NELEM, particletype, i, tag, MPI_COMM_WORLD);
55
56
57
      // all tasks receive particletype data
58
      MPI_Recv(p, NELEM, particletype, source, tag, MPI_COMM_WORLD, &stat);
59
60
      printf("rank= %d %3.2f %3.2f %3.2f %d %d\n", rank,p[3].x,
61
           p[3].y,p[3].z,p[3].velocity,p[3].n,p[3].type);
62
63
      // free datatype when done using it
64
      MPI Type free(&particletype);
65
      MPI_Finalize();
66
      }
```

```
Fortan - Struct Derived Data Type Example
      program struct
 1
 2
      include 'mpif.h'
 3
      integer NELEM
 5
      parameter (NELEM=25)
      integer numtasks, rank, source, dest, tag, i, ierr
 6
      integer stat(MPI STATUS SIZE)
 7
 9
      type Particle
10
      sequence
11
      real*4 x, y, z, velocity
12
      integer n, type
13
      end type Particle
14
15
      type (Particle) p(NELEM), particles(NELEM)
16
      integer particletype, oldtypes(0:1) ! required variables
17
      integer blockcounts(0:1), offsets(0:1), extent
18
      tag = 1
19
      call MPI_INIT(ierr)
20
      call MPI COMM RANK (MPI COMM WORLD, rank, ierr)
21
22
      call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)
23
24
      ! setup description of the 4 MPI_REAL fields x, y, z, velocity
25
      offsets(0) = 0
26
      oldtypes(0) = MPI REAL
27
      blockcounts(0) = \overline{4}
28
      ! setup description of the 2 MPI INTEGER fields n, type
29
      ! need to first figure offset by getting size of MPI REAL
31
      call MPI_TYPE_EXTENT(MPI_REAL, extent, ierr)
32
      offsets(1) = 4 * extent
      oldtypes(1) = MPI INTEGER
33
34
      blockcounts(1) = \overline{2}
35
36
      ! define structured type and commit it
37
      call MPI_TYPE_STRUCT(2, blockcounts, offsets, oldtypes, &
38
                            particletype, ierr)
39
      call MPI TYPE COMMIT (particletype, ierr)
40
41
       ! task 0 initializes the particle array and then sends it to each task
42
      if (rank .eq. 0) then
         do i=0, NELEM-1
43
         particles(i) = Particle ( 1.0*i, -1.0*i, 1.0*i, 0.25, i, mod(i,2) )
44
45
         end do
46
47
         do i=0, numtasks-1
48
         call MPI_SEND(particles, NELEM, particletype, i, tag, &
49
                        MPI COMM WORLD, ierr)
50
         end do
      endif
51
52
53
      ! all tasks receive particletype data
54
      source = 0
      call MPI_RECV(p, NELEM, particletype, source, tag, &
55
                     MPI_COMM_WORLD, stat, ierr)
56
57
58
      print *, 'rank= ',rank,' p(3)= ',p(3)
59
      ! free datatype when done using it
60
61
      call MPI TYPE FREE (particletype, ierr)
62
      call MPI FINALIZE (ierr)
63
      end
```

```
rank= 0 3.00 -3.00 3.00 0.25 3 1
rank= 2 3.00 -3.00 3.00 0.25 3 1
rank= 1 3.00 -3.00 3.00 0.25 3 1
rank= 3 3.00 -3.00 3.00 0.25 3 1
```

Group and Communicator Management Routines

Groups vs. Communicators:

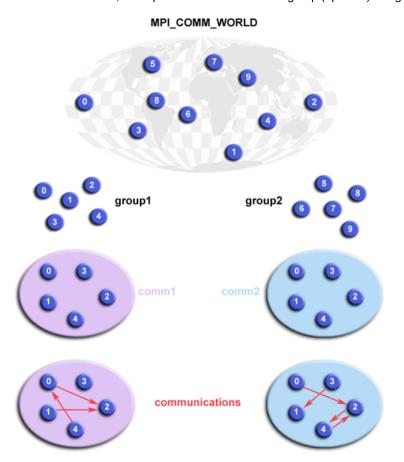
- A group is an ordered set of processes. Each process in a group is associated with a unique integer rank. Rank values start at zero and go to N-1, where N is the number of processes in the group. In MPI, a group is represented within system memory as an object. It is accessible to the programmer only by a "handle". A group is always associated with a communicator object.
- A communicator encompasses a group of processes that may communicate with each other. All MPI messages must specify a
 communicator. In the simplest sense, the communicator is an extra "tag" that must be included with MPI calls. Like groups,
 communicators are represented within system memory as objects and are accessible to the programmer only by "handles". For
 example, the handle for the communicator that comprises all tasks is MPI COMM WORLD.
- From the programmer's perspective, a group and a communicator are one. The group routines are primarily used to specify which processes should be used to construct a communicator.

Primary Purposes of Group and Communicator Objects:

- 1. Allow you to organize tasks, based upon function, into task groups.
- 2. Enable Collective Communications operations across a subset of related tasks.
- 3. Provide basis for implementing user defined virtual topologies
- 4. Provide for safe communications

Programming Considerations and Restrictions:

- Groups/communicators are dynamic they can be created and destroyed during program execution.
- · Processes may be in more than one group/communicator. They will have a unique rank within each group/communicator.
- MPI provides over 40 routines related to groups, communicators, and virtual topologies.
- Typical usage:
 - 1. Extract handle of global group from MPI COMM WORLD using MPI Comm group
 - 2. Form new group as a subset of global group using MPI Group incl
 - 3. Create new communicator for new group using MPI_Comm_create
 - 4. Determine new rank in new communicator using MPI_Comm_rank
 - 5. Conduct communications using any MPI message passing routine
 - 6. When finished, free up new communicator and group (optional) using MPI_Comm_free and MPI_Group_free



Group and Communicator Management Routines

Create two different process groups for separate collective communications exchange. Requires creating new communicators

```
Ħ
     C Language - Group and Communicator Example
      #include "mpi.h"
 1
 2
      #include <stdio.h>
      #define NPROCS 8
 3
 5
      main(int argc, char *argv[])
 6
                 rank, new_rank, sendbuf, recvbuf, numtasks,
 7
                 ranks1[4]={0,1,2,3}, ranks2[4]={4,5,6,7};
 8
      MPI_Group orig_group, new_group; // required variables
 9
      MPI Comm new comm;
                              // required variable
10
11
      MPI_Init(&argc,&argv);
12
      MPI Comm rank (MPI COMM WORLD, &rank);
13
      MPI Comm size(MPI COMM WORLD, &numtasks);
14
15
      if (numtasks != NPROCS) {
        printf("Must specify MP PROCS= %d. Terminating.\n",NPROCS);
16
17
        MPI Finalize();
18
        exit(0);
19
        }
20
21
      sendbuf = rank;
22
23
      // extract the original group handle
24
      MPI Comm group (MPI COMM WORLD, &orig group);
25
26
      // divide tasks into two distinct groups based upon rank
27
      if (rank < NPROCS/2) {</pre>
28
        MPI_Group_incl(orig_group, NPROCS/2, ranks1, &new_group);
29
30
      else {
        MPI Group incl(orig group, NPROCS/2, ranks2, &new group);
31
32
33
      // create new new communicator and then perform collective communications
34
35
      MPI_Comm_create(MPI_COMM_WORLD, new_group, &new_comm);
36
      MPI_Allreduce(&sendbuf, &recvbuf, 1, MPI_INT, MPI_SUM, new_comm);
37
38
      // get rank in new group
39
      MPI Group rank (new group, &new rank);
40
      printf("rank= %d newrank= %d recvbuf= %d\n",rank,new_rank,recvbuf);
41
42
      MPI_Finalize();
43
      }
```

```
Fortran - Group and Communicator Example
 1
      program group
 2
      include 'mpif.h'
 3
      integer NPROCS
 4
      parameter(NPROCS=8)
 6
      integer rank, new_rank, sendbuf, recvbuf, numtasks
 7
      integer ranks1(4), ranks2(4), ierr
 8
      integer orig_group, new_group, new_comm ! r
data ranks1 /0, 1, 2, 3/, ranks2 /4, 5, 6, 7/
                                                    ! required variables
10
11
      call MPI_INIT(ierr)
12
      call MPI COMM RANK (MPI COMM WORLD, rank, ierr)
      call MPI COMM SIZE (MPI COMM WORLD, numtasks, ierr)
13
14
15
      if (numtasks .ne. NPROCS) then
```

```
print *, 'Must specify NPROCS= ',NPROCS,' Terminating.'
       call MPI_FINALIZE(ierr)
17
18
       stop
     endif
19
20
21
     sendbuf = rank
22
23
     ! extract the original group handle
24
     call MPI_COMM_GROUP(MPI_COMM_WORLD, orig_group, ierr)
25
26
     ! divide tasks into two distinct groups based upon rank
27
     if (rank .1t. NPROCS/2) then
28
        call MPI_GROUP_INCL(orig_group, NPROCS/2, ranks1, new_group, ierr)
29
     else
30
        call MPI GROUP INCL (orig group, NPROCS/2, ranks2, new group, ierr)
31
     endif
32
33
     ! create new new communicator and then perform collective communications
34
     call MPI_COMM_CREATE(MPI_COMM_WORLD, new_group, new_comm, ierr)
35
     call MPI_ALLREDUCE(sendbuf, recvbuf, 1, MPI_INTEGER, MPI_SUM, new_comm, ierr)
36
37
     ! get rank in new group
38
     call MPI_GROUP_RANK(new_group, new_rank, ierr)
     print *, 'rank= ',rank,' newrank= ',new_rank,' recvbuf= ', recvbuf
39
40
41
     call MPI FINALIZE (ierr)
42
     end
```

```
rank= 7 newrank= 3 recvbuf= 22 rank= 0 newrank= 0 recvbuf= 6 rank= 1 newrank= 1 recvbuf= 6 rank= 2 newrank= 2 recvbuf= 6 rank= 6 newrank= 2 recvbuf= 22 rank= 3 newrank= 3 recvbuf= 6 rank= 4 newrank= 0 recvbuf= 22 rank= 5 newrank= 1 recvbuf= 22
```

Virtual Topologies

What Are They?

- In terms of MPI, a virtual topology describes a mapping/ordering of MPI processes into a geometric "shape".
- The two main types of topologies supported by MPI are Cartesian (grid) and Graph.
- MPI topologies are virtual there may be no relation between the physical structure of the parallel machine and the process topology.
- Virtual topologies are built upon MPI communicators and groups.
- Must be "programmed" by the application developer.

Why Use Them?

- Convenience
 - Virtual topologies may be useful for applications with specific communication patterns patterns that match an MPI topology structure.
 - For example, a Cartesian topology might prove convenient for an application that requires 4-way nearest neighbor communications for grid based data.
- · Communication Efficiency
 - Some hardware architectures may impose penalties for communications between successively distant "nodes".
 - A particular implementation may optimize process mapping based upon the physical characteristics of a given parallel machine.
 - The mapping of processes into an MPI virtual topology is dependent upon the MPI implementation, and may be totally ignored.

Example:

A simplified mapping of processes into a Cartesian virtual topology appears below:

0	1 (0,1)	2	3
(0,0)		(0,2)	(0,3)
4	5	6	7
(1,0)	(1,1)	(1,2)	(1,3)
8	9	10	11
(2,0)	(2,1)	(2,2)	(2,3)
12	13	14	15
(3,0)	(3,1)	(3,2)	(3,3)

Virtual Topology Routines

Create a 4 x 4 Cartesian topology from 16 processors and have each process exchange its rank with four neighbors.

```
C Language - Cartesian Virtual Topology Example
     #include "mpi.h"
     #include <stdio.h>
 2
     #define SIZE 16
 3
 4
     #define UP
     #define DOWN
 5
                    1
     #define LEFT 2
 6
 7
     #define RIGHT 3
 8
 9
     main(int argc, char *argv[])
10
     int numtasks, rank, source, dest, outbuf, i, tag=1,
        inbuf[4]={MPI_PROC_NULL,MPI_PROC_NULL,MPI_PROC_NULL,MPI_PROC_NULL,},
11
12
        nbrs[4], dims[2] = {4,4},
13
        periods[2]={0,0}, reorder=0, coords[2];
14
15
     MPI_Request reqs[8];
16
     MPI Status stats[8];
17
     MPI Comm cartcomm;
                            // required variable
18
19
     MPI_Init(&argc,&argv);
     MPI_Comm_size(MPI_COMM_WORLD, &numtasks);
20
21
22
     if (numtasks == SIZE) {
23
        // create cartesian virtual topology, get rank, coordinates, neighbor ranks
        MPI Cart create (MPI COMM WORLD, 2, dims, periods, reorder, &cartcomm);
24
25
        MPI Comm rank(cartcomm, &rank);
26
        MPI_Cart_coords(cartcomm, rank, 2, coords);
        MPI_Cart_shift(cartcomm, 0, 1, &nbrs[UP], &nbrs[DOWN]);
MPI_Cart_shift(cartcomm, 1, 1, &nbrs[LEFT], &nbrs[RIGHT]);
27
28
29
         printf("rank= %d coords= %d %d neighbors(u,d,l,r)= %d %d %d %d n", \\
30
31
                rank, coords[0], coords[1], nbrs[UP], nbrs[DOWN], nbrs[LEFT],
32
                nbrs[RIGHT]);
33
34
        outbuf = rank;
35
36
        // exchange data (rank) with 4 neighbors
        for (i=0; i<4; i++) {
37
38
           dest = nbrs[i];
39
            source = nbrs[i];
40
           MPI_Isend(&outbuf, 1, MPI_INT, dest, tag,
41
                      MPI COMM WORLD, &reqs[i]);
42
           MPI Irecv(&inbuf[i], 1, MPI INT, source, tag,
43
                      MPI_COMM_WORLD, &reqs[i+4]);
44
45
        MPI Waitall(8, reqs, stats);
46
47
48
        printf("rank= %d
                                             inbuf(u,d,l,r)= %d %d %d %d\n",
49
                rank,inbuf[UP],inbuf[DOWN],inbuf[LEFT],inbuf[RIGHT]); }
50
     else
        printf("Must specify %d processors. Terminating.\n",SIZE);
```

```
52
53 MPI_Finalize();
54 }
```

```
Fortran - Cartesian Virtual Topology Example
 1
     program cartesian
 2
     include 'mpif.h'
 3
 4
     integer SIZE, UP, DOWN, LEFT, RIGHT
 5
     parameter(SIZE=16)
 6
     parameter (UP=1)
 7
     parameter (DOWN=2)
 8
     parameter (LEFT=3)
 9
     parameter (RIGHT=4)
10
     integer numtasks, rank, source, dest, outbuf, i, tag, ierr, &
11
              inbuf(4), nbrs(4), dims(2), coords(2), periods(2), reorder
     integer stats(MPI_STATUS_SIZE, 8), reqs(8)
integer cartcomm ! required variable
12
13
14
     data inbuf /MPI PROC NULL, MPI PROC NULL, MPI PROC NULL, MPI PROC NULL/, &
15
           dims /4,4/, tag /1/, periods /0,0/, reorder /0/
16
      call MPI INIT(ierr)
17
     call MPI COMM SIZE (MPI COMM WORLD, numtasks, ierr)
18
19
20
     if (numtasks .eq. SIZE) then
         ! create cartesian virtual topology, get rank, coordinates, neighbor ranks call MPI_CART_CREATE(MPI_COMM_WORLD, 2, dims, periods, reorder, &
21
22
23
                                cartcomm, ierr)
24
         call MPI_COMM_RANK(cartcomm, rank, ierr)
25
         call MPI_CART_COORDS(cartcomm, rank, 2, coords, ierr)
         call MPI_CART_SHIFT(cartcomm, 0, 1, nbrs(UP), nbrs(DOWN), ierr)
26
27
         call MPI CART SHIFT (cartcomm, 1, 1, nbrs(LEFT), nbrs(RIGHT), ierr)
28
29
         write(*,20) rank,coords(1),coords(2),nbrs(UP),nbrs(DOWN), &
30
                      nbrs(LEFT), nbrs(RIGHT)
31
32
         ! exchange data (rank) with 4 neighbors
33
         outbuf = rank
34
         do i=1,4
35
            dest = nbrs(i)
36
            source = nbrs(i)
37
            call MPI_ISEND(outbuf, 1, MPI_INTEGER, dest, tag, &
38
                            MPI COMM WORLD, reqs(i), ierr)
            call MPI_IRECV(inbuf(i), 1, MPI_INTEGER, source, tag, &
39
40
                            MPI COMM WORLD, reqs(i+4), ierr)
41
         enddo
42
43
         call MPI WAITALL(8, reqs, stats, ierr)
44
45
         write(*,30) rank,inbuf
46
47
       print *, 'Must specify',SIZE,' processors. Terminating.'
48
49
      endif
50
51
      call MPI FINALIZE (ierr)
52
53
      20 format('rank= ',I3,' coords= ',I2,I2, &
                 ' neighbors(u,d,1,r)= ',I3,I3,I3,I3 )
54
55
      30 format('rank= ',I3,'
                  inbuf(u,d,1,r)= ',I3,I3,I3,I3)
56
57
58
      end
```

Sample program output: (partial)

```
rank= 0 coords= 0 0 neighbors(u,d,1,r)= -1 4 -1 1
rank= 0 inbuf(u,d,1,r)= -1 4 -1 1
```

```
rank=
        8 coords=
                    2 0 neighbors (u,d,l,r) =
                                                 4 12 -1
rank=
        R
                             inbuf(u,d,l,r) =
rank=
        1 coords=
                     0 1 neighbors (u,d,l,r) =
                                                           2
                                                -1 5
                                                          2
rank=
        1
                             inbuf(u,d,l,r) =
                                                       0
rank=
       13 coords=
                    3 1 neighbors (u,d,l,r) =
                                                 9 -1 12 14
rank=
       13
                             inbuf(u,d,l,r) =
                                                 9 -1 12 14
. . .
. . .
                    0 3 neighbors (u,d,l,r) =
        3 coords=
                                                    7
                                                       2 -1
rank=
                                                -1
                             inbuf(u,d,l,r) =
                                                -1
                                                    7
                                                        2
                                                          -1
rank=
        3
       11 coords=
                    2 3 neighbors (u,d,l,r) =
                                                 7 15 10 -1
rank=
                             inbuf(u,d,l,r) =
rank=
                                                 7 15 10 -1
rank=
                                                 6 14
       10 coords=
                    2 2 neighbors (u,d,l,r) =
                                                        9 11
rank=
       10
                             inbuf(u,d,l,r) =
                                                 6 14
                                                        9
                                                         11
                                                 5 13
rank=
        9 coords=
                    2 1 neighbors (u,d,l,r) =
                                                        8 10
rank=
                             inbuf(u,d,l,r) =
                                                 5 13
                                                        8 10
```

A Brief Word on MPI-2 and MPI-3

► MPI-2:

- Intentionally, the MPI-1 specification did not address several "difficult" issues. For reasons of expediency, these issues were
 deferred to a second specification, called MPI-2 in 1998.
- MPI-2 was a major revision to MPI-1 adding new functionality and corrections.
- · Key areas of new functionality in MPI-2:
 - **Dynamic Processes** extensions that remove the static process model of MPI. Provides routines to create new processes after job startup.
 - **One-Sided Communications** provides routines for one directional communications. Include shared memory operations (put/get) and remote accumulate operations.
 - · Extended Collective Operations allows for the application of collective operations to inter-communicators
 - External Interfaces defines routines that allow developers to layer on top of MPI, such as for debuggers and profilers.
 - o Additional Language Bindings describes C++ bindings and discusses Fortran-90 issues.
 - Parallel I/O describes MPI support for parallel I/O.

► MPI-3:

- The MPI-3 standard was adopted in 2012, and contains significant extensions to MPI-1 and MPI-2 functionality including:
 - Nonblocking Collective Operations permits tasks in a collective to perform operations without blocking, possibly offering
 performance improvements.
 - New One-sided Communication Operations to better handle different memory models.
 - Neighborhood Collectives extends the distributed graph and Cartesian process topologies with additional communication power.
 - o Fortran 2008 Bindings expanded from Fortran 90 bindings
 - **MPIT Tool Interface** allows the MPI implementation to expose certain internal variables, counters, and other states to the user (most likely performance tools).
 - o Matched Probe fixes an old bug in MPI-2 where one could not probe for messages in a multi-threaded environment.

► More Information on MPI-2 and MPI-3:

• MPI Standard documents: http://www.mpi-forum.org/docs/

MPI Exercise 3

Your Choice

Overview:

Login to the LC workshop cluster, if you are not already logged in

- Following the Exercise 3 instructions will take you through all sorts of MPI programs pick any/all that are of interest.
- The intention is review the codes and see what's happening not just compile and run.
- Several codes provide serial examples for a comparison with the parallel MPI versions.
- Check out the "bug" programs.



GO TO THE EXERCISE HERE

This completes the tutorial.



Please complete the online evaluation form - unless you are doing the exercise, in which case please complete it at the end of the exercises.

Where would you like to go now?

- Exercise 3
- Agenda
- Back to the top

References and More Information

- · Author: Blaise Barney, Livermore Computing.
- MPI Standard documents: http://www.mpi-forum.org/docs/
- "Using MPI", Gropp, Lusk and Skjellum. MIT Press, 1994.
- MPI Tutorials: www.mcs.anl.gov/research/projects/mpi/tutorial
- Livermore Computing specific information:
 - Linux Clusters Overview tutorial computing.llnl.gov/tutorials/linux_clusters
 - Using the Sequoia/Vulcan BG/Q Systems tutorial computing.llnl.gov/tutorials/bgq
- "A User's Guide to MPI", Peter S. Pacheco. Department of Mathematics, University of San Francisco.

Appendix A: MPI-1 Routine Index

- These man pages were derived from the MVAPICH 0.9 implementation of MPI and may differ from the man pages of other implementations.
- Not all MPI routines are shown
- * = deprecated in MPI-2.0, replaced in MPI-3.0
- The complete MPI-3 standard (2012) defines over 430 routines.

Environment Management Routines				
MPI_Abort MPI_Errhandler_create* MPI_Errhandler_free MPI_Errhandler_get*				
MPI_Errhandler_set* MPI_Error_class MPI_Error_string MPI_Finalize				

/6	Messag	ge Passing Interface (MPI)	
MPI_Get_processor_name	MPI_Get_version	MPI_Init	MPI_Initialized
MPI_Wtick	MPI_Wtime		
	Point-to-Point Comm	unication Routines	,
MPI_Bsend	MPI_Bsend_init	MPI_Buffer_attach	MPI_Buffer_detach
MPI_Cancel	MPI_Get_count	MPI_Get_elements	MPI_lbsend
MPI_Iprobe	MPI_Irecv	MPI_Irsend	MPI_Isend
MPI_Issend	MPI_Probe	MPI_Recv	MPI_Recv_init
MPI_Request_free	MPI_Rsend	MPI_Rsend_init	MPI_Send
MPI_Send_init	MPI_Sendrecv	MPI_Sendrecv_replace	MPI_Ssend
MPI_Ssend_init	MPI_Start	MPI_Startall	MPI_Test
MPI_Test_cancelled	MPI_Testall	MPI_Testany	MPI_Testsome
MPI_Wait	MPI_Waitall	MPI_Waitany	MPI_Waitsome
	Collective Commur	nication Routines	
MPI_Allgather	MPI_Allgatherv	MPI_Allreduce	MPI_Alltoall
MPI_Alltoallv	MPI_Barrier	MPI_Bcast	MPI_Gather
MPI_Gatherv	MPI_Op_create	MPI_Op_free	MPI_Reduce
MPI_Reduce_scatter	MPI_Scan	MPI_Scatter	MPI_Scatterv
	Process Grou	p Routines	
MPI_Group_compare	MPI_Group_difference	MPI_Group_excl	MPI_Group_free
MPI_Group_incl	MPI_Group_intersection	MPI Group range excl	MPI Group range incl
MPI_Group_rank	MPI_Group_size	MPI_Group_translate_ranks	MPI_Group_union
	Communicato	rs Routines	
MPI_Comm_compare	MPI_Comm_create	MPI_Comm_dup	MPI_Comm_free
MPI_Comm_group	MPI_Comm_rank	MPI_Comm_remote_group	MPI_Comm_remote_size
MPI_Comm_size	MPI_Comm_split	MPI_Comm_test_inter	MPI_Intercomm_create
MPI_Intercomm_merge			
	Derived Type	es Routines	,
MPI_Type_commit	MPI_Type_contiguous	MPI_Type_extent*	MPI_Type_free
MPI_Type_hindexed*	MPI_Type_hvector*	MPI_Type_indexed	MPI_Type_lb
MPI_Type_size	MPI_Type_struct*	MPI_Type_ub*	MPI_Type_vector
	Virtual Topolo	gy Routines	
MPI_Cart_coords	MPI_Cart_create	MPI_Cart_get	MPI_Cart_map
MPI_Cart_rank	MPI_Cart_shift	MPI_Cart_sub	MPI_Cartdim_get
MPI_Dims_create	MPI_Graph_create	MPI_Graph_get	MPI_Graph_map
MPI_Graph_neighbors	MPI_Graph_neighbors_count	MPI_Graphdims_get	MPI_Topo_test
-	Miscellaneou	s Routines	
MPI_Address*	MPI_Attr_delete*	MPI_Attr_get*	MPI_Attr_put*
MPI_Keyval_create*	MPI_Keyval_free*	MPI_Pack	MPI_Pack_size
MPI_Pcontrol	MPI_Unpack		

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