




Laboratoire  
Méthodes  
Formelles

# A FLOATING-POINT NUMBERS THEORY FOR EVENT-B

 The LMF Lab Seminar

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# OUTLINE

- The context of the work
- The motivating example
- The proposed approach
- Revisiting the motivating example
- Conclusion and future works

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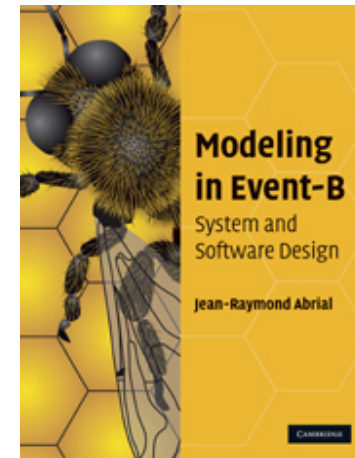
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# THE EVENT-B METHOD

- The **Event-B method** is an evolution of the **classical B method**.
  - modeling a system by a **set of events** instead of **operations**.
- The **Event-B method** is a **formal method** based on **first-order logic** and **set theory**.
- The **Event-B method** is based on :
  - the notions of **pre-conditions** and **post-conditions** (**Hoare**),
  - the **weakest pre-condition** (**Dijkstra**),
  - and the **calculus of substitution** (**Abrial**).



# USING EVENT-B METHOD

- The **Rodin** platform (an **Eclipse-based IDE**) is intended to support the construction and verification of **Event-B models**.
- The use of the **Event-B method** has continued to increase.
  - applied to various applications and domains.
  - railway, automotive, aeronautics, cybersecurity, nuclear-energy, ...
- The **Event-B method** is adapted to analyse **discrete systems**.
  - offers the possibility of modelling **discrete behaviors**.

# THE EVENT-B METHOD

**CONTEXT**  $ctx_1$   
**EXTENDS**  $ctx_2$

**SETS**  $s$   
**CONSTANTS**  $c$   
**AXIOMS**  
     $A(s, c)$   
**THEOREMS**  
     $T(s, c)$   
**END**

**MACHINE**  $mch_1$   
**REFINES**  $mch_2$   
**SEES**  $ctx_i$

**VARIABLES**  $v$   
**INVARIANTS**  
     $I(s, c, v)$   
**THEOREMS**  
     $T(s, c, v)$   
**EVENTS**  
     $[events\_list]$   
**END**

$event \hat{=}$   
    **any**  $x$   
    **where**  
         $G(s, c, v, x)$   
    **then**  
         $BA(s, c, v, x, v')$   
    **end**

$A(s, c) \Rightarrow T(s, c)$   
 $A(s, c) \wedge I(s, c, v) \Rightarrow T(s, c, v)$   
 $A(s, c) \wedge I(s, c, v) \wedge G(s, c, v, x) \Rightarrow \exists v'. BA(s, c, v, x, v')$   
 $A(s, c) \wedge I(s, c, v) \wedge G(s, c, v, x) \wedge BA(s, c, v, x, v') \Rightarrow I(s, c, v')$   
...

# THE THEORY PLUGIN

- **Theory Plug-in** provides capabilities to **extend the Event-B mathematical language** and **the Rodin proving infrastructure**.
- An **Event-B theory** can contain :
  - new datatype definitions,
  - new polymorphic operator definitions,
  - axiomatic definitions,
  - theorems,
  - associated rewrite and inference rules.

- Michael J. Butler and Issam Maamria.  
*Practical theory extension in Event-B*. [Theories of Programming and Formal Methods](#). 2013.
- Thai Son Hoang, Laurent Voisin, Asieh Salehi, Michael J. Butler, Toby Wilkinson, and Nicolas Beauger.  
*Theory plug-in for Rodin 3.x*. [CoRR](#), [abs/1701.08625](#), 2017.

# THE EVENT-B METHOD

**THEORY**  $thy_1$   
**IMPORT**  $thy_2$

**DATATYPES**  
 $DT_1, \dots, DT_n$

**OPERATORS**  
 $OP_{11}, \dots, OP_{1n}$

**AXIOMATIC DEFINITIONS**  
**operators**  
 $OP_{21}, \dots, OP_{2n}$   
**axioms**  
 $A$

**THEOREMS**  
 $T$

**PROOF RULES**  
 $PR$

**END**

**CONTEXT**  $ctx_1$   
**EXTENDS**  $ctx_2$

**SETS**  $s$

**CONSTANTS**  $c$

**AXIOMS**  
 $A(s, c)$

**THEOREMS**  
 $T(s, c)$

**END**

**MACHINE**  $mch_1$   
**REFINES**  $mch_2$   
**SEES**  $ctx_i$

**VARIABLES**  $v$

**INVARIANTS**  
 $I(s, c, v)$

**THEOREMS**  
 $T(s, c, v)$

**EVENTS**  
 $[events\_list]$

**END**



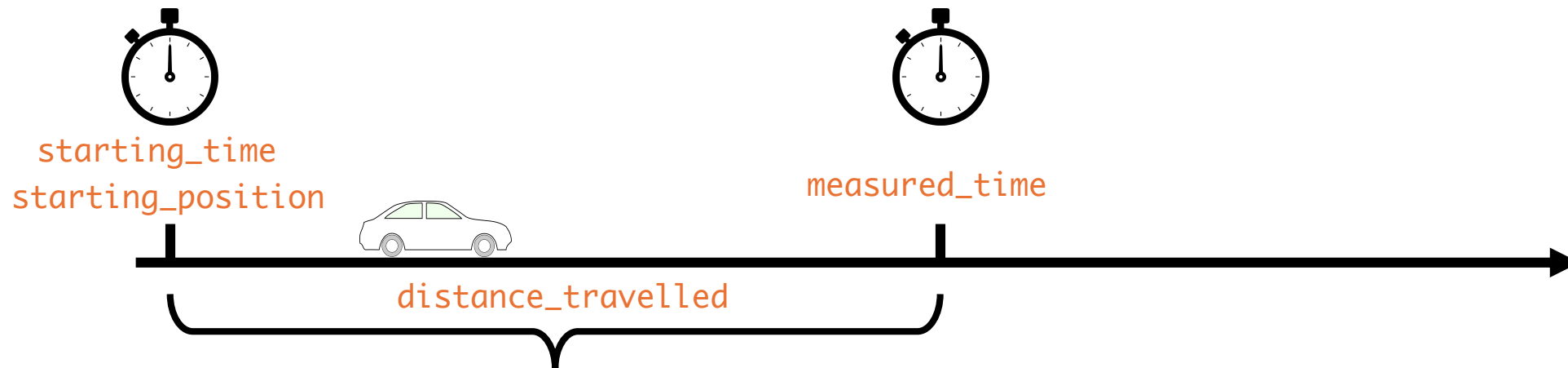
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# A SIMPLE EXAMPLE

- System that continuously calculates **a moving object's speed**



# A SIMPLE EXAMPLE

- Analysing **two functional properties**:
  - **PROP-1** : **the speed of the moving object** is equal to the *distance\_travelled* divided by the *measured\_time* ( $v = d/t$ ).
  - **PROP-2** : when the *distance\_travelled* is **strictly positive**, the *speed* of the moving object must also be **strictly positive**.
    - **the object moves** when its *speed* is different from zero.

**Objectives** → showing some **modelling and verification problems** :

- analysing **physical phenomena**.
  - expressions that come from **the physics laws**.
- using **integer** variables to handle **small values**.

# THE EVENT-B MODEL

- Analysing **two functional properties**:
  - **PROP-1** : **the speed of the moving object** is equal to the *distance\_travelled* divided by the *measured\_time* ( $v = d/t$ ).
  - **PROP-2** : when the *distance\_travelled* is **strictly positive**, the *speed* of the moving object must also be **strictly positive**.
    - **the object moves** when its *speed* is different from zero.

```
MACHINE mch_integer_version
```

```
...
```

```
INVARIANTS
```

```
@inv1: distance_travelled ∈ ℕ
```

```
@inv2: measured_time ∈ ℕ1
```

```
@inv3: speed ∈ ℕ
```

```
@inv4: starting_position ∈ ℕ
```

```
@inv5: starting_time ∈ ℕ
```

```
@inv6: speed = distance_travelled ÷ measured_time // PROP-1
```

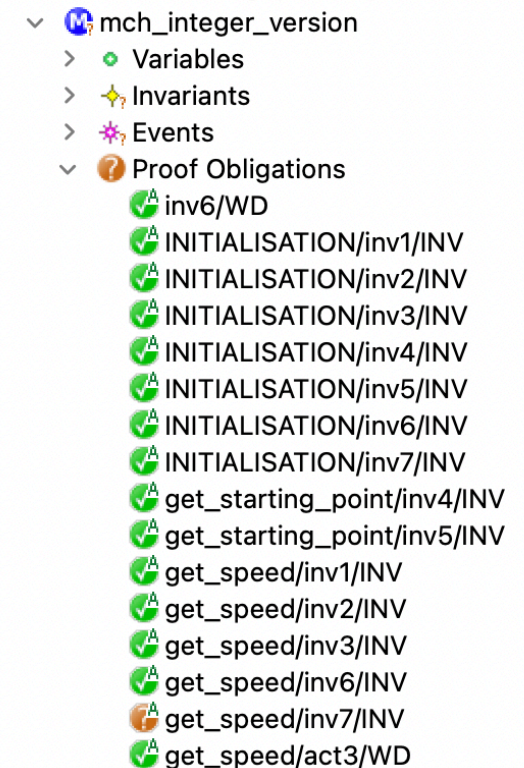
```
@inv7: distance_travelled > 0 ⇒ speed > 0 // PROP-2
```

# THE EVENT-B MODEL

```
MACHINE mch_integer_version
...
EVENTS
...
get_speed  $\hat{=}$ 
  any p t
  where
    @grd1:  $p \in \mathbb{N}_1 \wedge p > \text{starting\_position}$ 
    @grd2:  $t \in \mathbb{N}_1 \wedge t > \text{starting\_time}$ 
  then
    @act1:  $\text{distance\_travelled} := p - \text{starting\_position}$ 
    @act2:  $\text{measured\_time} := t - \text{starting\_time}$ 
    @act3:  $\text{speed} := (p - \text{starting\_position}) \div (t - \text{starting\_time})$ 
  end
END
```

# GENERATED AND PROVEN POS

- All POs are green **except** the one for maintaining the *@inv7* invariant by the *get\_speed* event.
- This invariant formalises the **PROP 2** property.
  - the object moves (*distance\_travelled*  $\neq 0$ ) when *speed*  $\neq 0$ .
- The *get\_speed* event calculates the new value of *distance\_travelled* that can be  $<$  the new value of *measured\_time*.
  - ➡ the new value of *speed*  
(*distance\_travelled*  $\div$  *measured\_time*)  
can be  $= 0$  while *distance\_travelled*  $\neq 0$
  - ➡  $\div$  makes an integer division



# CONCLUSION

**The basic types and operators of the Event-B language  
are not adapted to our needs**

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# FLOATING-POINT NUMBERS

$$x = 3.14159265359 = \underbrace{314159265359}_{\text{significand}} \times \underbrace{10}_{\text{base}}^{\text{exponent} \atop -11}$$

We have chosen that the base always equals ten in our models.

$$x = s(x) \times 10^{e(x)}$$

- The proposed theory **does not model limited precision**.
- The **operators** defined in the theory involve **no precision loss**.

# THE PROPOSED APPROACH

- To allow the **Event-B language** to **embed** this **FP representation**, we need to define two theories:
  1. the first theory formalises **the power operator**.
    - ✗ **^ operator** is **not implemented** in the provers besides **^0** and **^1**.
  2. the second theory formalises **floating-point numbers** by specifying:
    - ➡ the corresponding **data type**,
    - ➡ the supported **arithmetic operators**,
    - ➡ some **axioms** and **theorems** that characterize the proposed modelling.

# THE POWER OPERATOR

**THEORY** thy\_power\_operator

## AXIOMATIC DEFINITIONS

### operators

pow( $x \in \mathbb{Z}, n \in \mathbb{N}$ ) :  $\mathbb{Z}$  **INFIX** //  $x$  pow  $n = x^n$

**wd condition** :  $\neg (x = 0 \wedge n = 0)$  //  $0^0$  is not defined

### axioms

@axm1:  $\forall n \cdot n \in \mathbb{N}_1 \Rightarrow 0 \text{ pow } n = 0$

@axm2:  $\forall x \cdot x \in \mathbb{Z} \wedge x \neq 0 \Rightarrow x \text{ pow } 0 = 1$

@axm3:  $\forall x, n \cdot x \in \mathbb{Z} \wedge x \neq 0 \wedge n \in \mathbb{N}_1 \Rightarrow x \text{ pow } n = x \times (x \text{ pow } (n - 1))$

...

## THEOREMS

@thm1:  $\forall x, n, m \cdot \dots \Rightarrow x \text{ pow } (n + m) = (x \text{ pow } n) \times (x \text{ pow } m)$

@thm2:  $\forall x, n, m \cdot \dots \Rightarrow (x \text{ pow } n) \text{ pow } m = x \text{ pow } (n \times m)$

@thm3:  $\forall x, y, n \cdot \dots \Rightarrow (x \times y) \text{ pow } n = (x \text{ pow } n) \times (y \text{ pow } n)$

...

**END**

# SOME REMARKS

- By using this theory, it **becomes possible to prove**, for example, that  $5 \text{ pow } 3 = 125$
- **The proofs** of all theorems were made by **induction** (following the rules defined by **Cervelle and Gervais - ABZ 2023**).
- We have chosen to define the **pow** operator in a **single theory** to offer the possibility of **reusing it** in other **Event-B developments**.

- Julien Cervelle and Frédéric Gervais.  
*Introducing Inductive Construction in B with the Theory Plugin*. [ABZ, 2023](#).

# THE FLOATING-POINT NUMBERS THEORY

**THEORY** thy\_floating\_point\_numbers

## DATATYPES

$\text{FLOAT\_Type} \hat{=} \text{NEW\_FLOAT}(s \in \mathbb{Z}, e \in \mathbb{Z}) \text{ // } x = s(x) \times 10^{e(x)}$

## OPERATORS

$F0 \hat{=} \text{NEW\_FLOAT}(0, 0) \text{ // } 0 = 0 \times 10^0$

$F1 \hat{=} \text{NEW\_FLOAT}(1, 0) \text{ // } 1 = 1 \times 10^0$

$\text{FLOAT1\_Type} \hat{=} \{ x \cdot x \in \text{FLOAT\_Type} \wedge s(x) \neq 0 \mid x \}$

$\text{FLOAT}(x \in \mathbb{Z}) \hat{=} \text{NEW\_FLOAT}(x, 0) \text{ // } x = x \times 10^0$

$\text{l\_shift}(x \in \text{FLOAT\_Type}, \text{offset} \in \mathbb{N}) \hat{=} \text{NEW\_FLOAT}(s(x) \times (10 \text{ pow offset}), e(x) - \text{offset})$

$\text{eq}(x \in \text{FLOAT\_Type}, y \in \text{FLOAT\_Type}) \text{ **INFIX** } \hat{=} s(\text{l\_shift}(x, e(x) - \min(\{e(x), e(y)\}))) = s(\text{l\_shift}(y, e(y) - \min(\{e(x), e(y)\})))$

$\text{gt}(x \in \text{FLOAT\_Type}, y \in \text{FLOAT\_Type}) \text{ **INFIX** } \hat{=} \dots$

$\text{geq}(x \in \text{FLOAT\_Type}, y \in \text{FLOAT\_Type}) \text{ **INFIX** } \hat{=} \dots$

$\text{lt}(x \in \text{FLOAT\_Type}, y \in \text{FLOAT\_Type}) \text{ **INFIX** } \hat{=} \dots$

$\text{leq}(x \in \text{FLOAT\_Type}, y \in \text{FLOAT\_Type}) \text{ **INFIX** } \hat{=} \dots$

...

**END**

# THE FLOATING-POINT NUMBERS THEORY

**THEORY** thy\_floating\_point\_numbers

...

**OPERATORS**

...

plus( $x \in \text{FLOAT\_Type}$ ,  $y \in \text{FLOAT\_Type}$ ) **INFIX**  $\hat{=}$   
NEW\_FLOAT( $s(l\_shift(x, e(x) - \min(\{e(x), e(y)\})) + s(l\_shift(y, e(y) - \min(\{e(x), e(y)\})))$ ,  
min( $\{e(x), e(y)\}$ ))

minus( $x \in \text{FLOAT\_Type}$ ,  $y \in \text{FLOAT\_Type}$ ) **INFIX**  $\hat{=}$  ...

neg( $x \in \text{FLOAT\_Type}$ )  $\hat{=}$  ...

mult( $x \in \text{FLOAT\_Type}$ ,  $y \in \text{FLOAT\_Type}$ ) **INFIX**  $\hat{=}$  NEW\_FLOAT( $s(x) \times s(y)$  ,  $e(x) + e(y)$ )

f\_pow( $x \in \text{FLOAT\_Type}$ ,  $n \in \mathbb{N}$ ) **INFIX**  $\hat{=}$  NEW\_FLOAT( $s(x)$  pow  $n$ ,  $e(x) \times n$ )

floor( $x \in \text{FLOAT\_Type}$ )  $\hat{=}$  ...

ceiling( $x \in \text{FLOAT\_Type}$ )  $\hat{=}$  ...

integer( $x \in \text{FLOAT\_Type}$ )  $\hat{=}$  ...

frac( $x \in \text{FLOAT\_Type}$ )  $\hat{=}$  ...

...

**END**

# THE CASE OF *inv* AND *div* OPERATORS

- The proposed theory involves **no precision loss** for *plus* and *mult* operators.
- The **division** sometimes **induces a precision loss**.
  - ✗ ex. we cannot precisely represent the result of  $1/3$  or  $2/3$
- For the case of *inv* and *div* operators, we have defined **the well-definedness conditions**.
  - To calculate *inv*( $x$ ), we must find a  $z$ , with  $10^n = z \times s(x)$ .
    - ✓  $inv(2.5) = 1/2.5 = 0.4 = 4 \times 10^{-1}$  ( $z = 4$  because  $100 = 4 \times 25$ )
    - ✗  $inv(3) = 1/3 = 0.3333\dots$  ( $z$  **does not exist**)
  - To calculate  $x \text{ div } y$ , we must find a  $z$ , with  $10^n \times s(x) = z \times s(y)$ .
    - ✓  $2 \text{ div } 5 = 2/5 = 0.4 = 4 \times 10^{-1}$  ( $z = 4$  because  $10 \times 2 = 4 \times 5$ )
    - ✗  $2 \text{ div } 3 = 2/3 = 0.6666\dots$  ( $z$  **does not exist**)

# THE CASE OF *inv* AND *div* OPERATORS

**THEORY** thy\_floating\_point\_numbers

...

## **OPERATORS**

...

$\text{inv\_WD}(a \in \text{FLOAT1\_Type}) \hat{=}$

$\exists n, z \cdot n \in \mathbb{N} \wedge z \in \mathbb{Z} \wedge 10 \text{ pow } n = s(a) \times z$

$\text{div\_WD}(a \in \text{FLOAT\_Type}, b \in \text{FLOAT1\_Type}) \hat{=}$

$\exists n, z \cdot n \in \mathbb{N} \wedge z \in \mathbb{Z} \wedge s(a) \times (10 \text{ pow } n) = s(b) \times z$

...

## **AXIOMATIC DEFINITIONS**

...

### **operators**

$\text{div}(x \in \text{FLOAT\_Type}, y \in \text{FLOAT\_Type}) : \text{FLOAT\_Type}$  **INFIX**

**wd condition** :  $\text{div\_WD}(x)$

### **axioms**

**axm1**:  $\forall x, y, z \cdot (... \Rightarrow ((y \text{ mult } z) = x \Leftrightarrow (x \text{ div } y) = z))$

**axm2**:  $\forall x, y, z \cdot (... \Rightarrow ((y \text{ mult } z) \text{ eq } x \Leftrightarrow (x \text{ div } y) \text{ eq } z))$

**axm3**:  $\forall x, y \cdot (... \Rightarrow x \text{ mult inv}(y) = x \text{ div } y)$

...

**END**



# THE FLOATING-POINT NUMBERS THEORY

**THEORY** thy\_floating\_point\_numbers

...

## **THEOREMS**

@thm1:  $\forall x, y \cdot (\dots \Rightarrow x \text{ eq } y \Leftrightarrow y \text{ eq } x)$   
@thm2:  $\forall x \cdot (\dots \Rightarrow x \text{ geq } x \wedge x \text{ leq } x)$   
@thm3:  $\forall x, y \cdot (\dots x \text{ leq } y \wedge y \text{ leq } x \Rightarrow x \text{ eq } y)$   
@thm4:  $\forall x, y \cdot (\dots \Rightarrow x \text{ leq } y \vee y \text{ leq } x)$   
@thm5:  $\forall x, y, z \cdot (\dots x \text{ leq } y \wedge y \text{ leq } z \Rightarrow x \text{ leq } z)$   
@thm6:  $\forall x, y, z \cdot (\dots x \text{ leq } y \Rightarrow (x \text{ plus } z) \text{ leq } (y \text{ plus } z))$   
@thm7:  $\forall x, y, z \cdot (\dots x \text{ leq } y \Rightarrow (x \text{ mult } z) \text{ leq } (y \text{ mult } z))$   
@thm8:  $\forall x \cdot (\dots \Rightarrow x \text{ plus } F0 \text{ eq } x)$   
@thm9:  $\forall x, y \cdot (\dots \Rightarrow x \text{ plus } y = y \text{ plus } x)$   
@thm10:  $\forall x, y \cdot (\dots \Rightarrow x \text{ plus neg}(y) = y \text{ minus } x)$   
@thm11:  $\forall x \cdot (\dots \Rightarrow x \text{ minus } F0 \text{ eq } x)$   
@thm12:  $\forall x \cdot (\dots \Rightarrow x \text{ minus } x \text{ eq } F0)$   
@thm13:  $\forall x \cdot (\dots \Rightarrow x \text{ mult } F0 \text{ eq } F0)$   
@thm14:  $\forall x \cdot (\dots \Rightarrow x \text{ mult } F1 = x)$   
@thm15:  $\forall x, y \cdot (\dots \Rightarrow x \text{ mult } y = y \text{ mult } x)$   
@thm16:  $\forall x \cdot (\dots \Rightarrow \text{inv}(x) = F1 \text{ div } x)$   
@thm17:  $\forall x \cdot (\dots \Rightarrow x \text{ div } F1 = x)$   
@thm18:  $\forall x \cdot (\dots \Rightarrow x \text{ div } x = F1)$   
@thm19:  $\forall x \cdot (\dots \Rightarrow x \text{ mult inv}(x) = F1)$

...

**END**

## SOME REMARKS

- Due to our choice to formalise **unlimited precision FP** numbers, some **properties** that are **not true** in the FP numbers world **can be deduced**.
  - the associativity of addition and multiplication, for example
- If this theory **is refined** (towards the **IEEE Standard 754**, for example), the developer must **pay attention** to this point.

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# NATURAL VARIABLES

All **NATURAL** variables are typed by **PFLOAT\_Type** set containing **positive floating-point numbers**.

**THEORY** thy\_floating\_point\_numbers

...

$\text{PFLOAT\_Type} = \{ x \cdot x \in \text{FLOAT\_Type} \wedge s(x) \geq 0 \mid x \}$

$\text{PFLOAT1\_Type} = \{ x \cdot x \in \text{FLOAT\_Type} \wedge s(x) > 0 \mid x \}$

...

**END**

# REVISITING OUR EXAMPLE I

**MACHINE** mch\_floating\_point\_version

...

## **INVARIANTS**

@inv1: distance\_travelled  $\in$  PFLOAT\_Type

@inv2: measured\_time  $\in$  PLOAT1\_Type

@inv3: speed  $\in$  PFLOAT\_Type

@inv4: starting\_position  $\in$  PFLOAT\_Type

@inv5: starting\_time  $\in$  PFLOAT\_Type

@inv6:  $\text{div\_WD}(\text{distance\_travelled}, \text{measured\_time})$

@inv7: speed eq distance\_travelled div measured\_time

@inv8: distance\_travelled  $\text{gt } F0 \Rightarrow$  speed  $\text{gt } F0$

...

**END**

# REVISITING OUR EXAMPLE II

**MACHINE** mch\_floating\_point\_version

...

**EVENTS**

...

get\_speed  $\hat{=}$

**any** p t

**where**

@grd1:  $p \in \text{PFLOAT\_Type} \wedge p \text{ gt starting\_position}$

@grd2:  $t \in \text{PFLOAT\_Type} \wedge t \text{ gt starting\_time}$

@grd3:  $\text{div\_WD}(p \text{ minus starting\_position}, t \text{ minus starting\_time})$

**then**

@act1:  $\text{distance\_travelled} := p \text{ minus starting\_position}$






@act2:  $\text{measured\_time} := t \text{ minus starting\_time}$

@act3:  $\text{speed} := (p \text{ minus starting\_position}) \text{ div } (t \text{ minus starting\_time})$

**end**

**END**

# GENERATED AND PROVEN POS

- ✓  mch\_floating\_point\_speed
  - >  Variables
  - >  Invariants
  - >  Events
  - ✓  Proof Obligations
    - ✓ inv6/WD
    - ✓ inv7/WD
    - ✓ INITIALISATION/inv1/INV
    - ✓ INITIALISATION/inv2/INV
    - ✓ INITIALISATION/inv3/INV
    - ✓ INITIALISATION/inv4/INV
    - ✓ INITIALISATION/inv5/INV
    - ✓ INITIALISATION/inv6/INV
    - ✓ INITIALISATION/inv7/INV
    - ✓  INITIALISATION/inv8/INV
    - ✓  get\_starting\_point/inv4/INV
    - ✓  get\_starting\_point/inv5/INV
    - ✓ get\_speed/grd5/WD
    - ✓ get\_speed/inv1/INV
    - ✓ get\_speed/inv2/INV
    - ✓ get\_speed/inv3/INV
    - ✓  get\_speed/inv6/INV
    - ✓ get\_speed/inv7/INV
    - ✓ get\_speed/inv8/INV
    - ✓ get\_speed/act3/WD

- All generated POs have been proven.
- The *get\_speed/inv8/INV* PO becomes ✓.
  - ➡ thanks to handling small values (`]0..1[`),
  - ➡ and to the new *div* operator specification.

The floating-point numbers theory is more suitable than the basic integers of Event-B.

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# CONCLUSION

- Extending the **Event-B type-checking system** by an approach using the **theory plugin**.
- Development of a **floating point number theory** formalizing floating point numbers.
  - an extension of the **Event-B power operator**.
  - an **abstract representation** of the **floating-point numbers**.
  - a set of theorems and associated **rewrite** and **inference rules**.

# FUTURE WORKS

- Refining the proposed theory to any **more concrete implementation** (the **IEEE standard 754**, for example).
- Developing a **more general theory** formalizing the standard units of **measurement** defined by the **International System of Units (SI)**.
  - ▢ extends the **floating point number theory**.
  - ▢ helpful in **modelling cyber-physical/hybrid** systems.

# THANK YOU

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