

1 Classical linear codes 15.07.22

1.1 Repetition code

Repetition code. Encoding

$$0 \rightarrow 000, 1 \rightarrow 111 \quad (1)$$

Decoding: majority vote, e.g. $010 \rightarrow 0, 011 \rightarrow 1$. Succeeds when 0 or 1 bit is flipped, fails when 2 or 3 bits are flipped. Error probability is reduced from p to $3p^2$ (for $p \ll 1$).

1.2 Linear codes

Repetition code is an example of a linear code. A linear code (n, k, d) is defined by a generator matrix G of size $k \times n$ and a parity check matrix H of size $(n - k) \times n$. Code subspace is bG , for all k -bitstrings b . Alternatively, code subspace consists of all n -bitstrings b' satisfying $H(b')^T = 0$. For the repetition code

$$G = \begin{pmatrix} 1 & 1 & 1 \end{pmatrix}, \quad H = \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \end{pmatrix}, \quad H^T G = 0 \quad (2)$$

1.3 Code distance

Code distance is the minimum Hamming distance between any two code words. For a linear code, code distance is the minimum (non-zero) weight of n -bitstrings in the code space. For the repetition code $d = 3$. A code that can correct t errors has distance $d = 2t + 1$.

1.4 Dual code

For a linear code with generator matrix G and parity check matrix H the dual code is another linear code with G and H swapped $G^\perp = H, H^\perp = G$. For the dual repetition code

$$G = \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \end{pmatrix}, \quad H = \begin{pmatrix} 1 & 1 & 1 \end{pmatrix}, \quad H^T G = 0 \quad (3)$$

Dual repetition code encodes 2 bits into 3 bits

$$00 \rightarrow 000 \quad (4)$$

$$01 \rightarrow 011 \quad (5)$$

$$10 \rightarrow 110 \quad (6)$$

$$11 \rightarrow 101 \quad (7)$$

Distance of the code is 2. It can detect one error, but correct none.

1.5 Hamming bound

”Good enough” codes exist. Encoding k -bitstrings and being able to correct t errors (distance at least $d \geq 2t + 1$) requires the embedding space to have dimension at least n

$$2^n \geq 2^k \sum_{i=0}^t C_n^i \quad (8)$$

1.6 Gilbert–Varshamov bound

Codes need not be ”too bad”. Given n physical bits a code with distance d exists encoding k logical bits with

$$2^k \geq \frac{2^n}{\sum_{i=0}^{d-1} C_n^i} \quad (9)$$