

HANDS OFF MY PROPERTY!

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ABSTRACT

This paper presents a new and complementary approach to unit-testing the implementation of agent-based simulations, called property-based testing which allows to test specifications of the implementation directly in code which is then tested using *automated* test-data generation. We present two different models as case-studies in which we will show how to apply property-based testing to exploratory and explanatory agent-based models and what its limits are.

We conduct our implementations in the pure functional programming language Haskell, which is the origin of property-based testing. Also we show that simply by switching to such a language one gets rid of a large class of run-time bugs and is able to make stronger guarantees of correctness already at compile time without writing tests for some parts. Further, it makes isolated unit-tests quite easier.

Keywords: Agent-Based Simulation, Property-Based Testing, Validation & Verification, Model Checking, Haskell.

1 INTRODUCTION

When implementing an Agent-Based Simulation (ABS) it is of fundamental importance that the implementation is correct up to some specification and that this specification matches the real world in some way. This process is called verification and validation (V&V), where *validation* is the process of ensuring that a model or specification is sufficiently accurate for the purpose at hand whereas *verification* is the process of ensuring that the model design has been transformed into a computer model with sufficient accuracy (Robinson 2014). In other words, validation determines if we are we building the *right model* and verification if we are we building the *model right* (Balci 1998).

One can argue that ABS should require more rigorous programming standards than other computer simulations (Polhill, Izquierdo, and Gotts 2005). Because researchers in ABS look for an emergent behaviour in the dynamics of the simulation, they are always tempted to look for some surprising behaviour and expect something unexpected from their simulation. Also, due to ABS mostly exploratory nature, there exists some amount of uncertainty about the dynamics the simulation will produce before running it. Thus it is often very difficult to judge whether an unexpected outcome can be attributed to the model or has in fact

its roots in a subtle programming error (Galán, Izquierdo, Izquierdo, Santos, del Olmo, López-Paredes, and Edmonds 2009).

In general this implies that we can only *raise the confidence* in the correctness of the simulation: it is not possible to prove that a model is valid, instead one should think of confidence in its validity. Therefore, the process of V&V is not the proof that a model is correct but trying to prove that the model is incorrect. The more tests/checks one carries out which show that it is not incorrect, the more confidence we can place on the models validity. To tackle such a problem in software, software engineers have developed the concept of test-driven development (TDD) (todo: cite).

TODO: dont spend so much on TDD, mention that property-based testing can be put in the TDD framework but is not the focus of this paper Test-Driven Development (TDD) was conceived in the late 90s by Kent Beck (TODO: cite) as a way to a more agile approach to software-engineering where instead of doing each step (requirements, implementation, testing,...) as separated from each other, all of them are combined in shorter cycles. TDD approaches software construction in a way that one writes first unit-tests for the functionality one wants to test and then iteratively implements this functionality until all tests succeed. This is then repeated until the whole software package is finished. The important difference to traditional models, where the steps are done in separation from each other, is that the customer receives a working software package at the end of each short cycle, allowing to change requirements which in turn allows the software-development team to react quickly to changing requirements.

It is important to understand that the unit-tests act both as documentation / specification of what the code / interface which is tested should do and as an insurance against future changes which might break existing code. If the tests cover all possible code paths - there exist tools to measure the test-coverage and visualising the missing code-paths / tests - of the software, then if the tests also succeed after future changes one has very high confidence that these future changes didn't break existing functionality. If though tests break then either the changes are erroneous or the tests are an incomplete specification and need to be adapted to the new features.

Thus we can say that test-driven development in general and unit-testing together with code-coverage in particular, allow to guarantee the correctness of an implementation to some informal degree, which has been proven to be sufficiently enough through years of practice in the software industry all over the world. Also a fundamental strength of such tests is that programmers gain much more confidence when making changes to code - without such tests all bets are off and there is no reliable way to know whether the changes have broken something or not.

TODO: find out what property-based testing makes different from parameter variation - is there any difference? The work of (Figueredo, Siebers, Owen, Reys, and Aickelin 2014) compares System Dynamics approaches to Cancer Cell simulations to ABS. They investigate the statistical differences between a deterministic ODE (ordinary differential equation), a stochastic Gillespie Algorithm and a stochastic ABS approach. The authors reported that the larger the population size is the closer the 3 approaches matches, especially ABS matches the System Dynamics approaches more closely. Still they found that ABS is capable of capturing "rare" patterns which is due to the memory and stochastic variability on the individual level.

In this paper we discuss a complementary method of testing the implementation of an ABS, called *property-based* testing, which allows to directly express model-specifications in code and test them through *automated* test-data generation. We see it as an addition to TDD where it works in combination with unit-testing to verify and validate a simulation to increase the confidence in its correctness.

Property-based testing has its origins (Claessen and Hughes 2000, Claessen and Hughes 2002, Runciman, Naylor, and Lindblad 2008) in the pure functional programming language Haskell (Hudak, Hughes, Peyton Jones, and Wadler 2007) where it was first conceived and implemented and thus we discuss it from that

perspective. It has been successfully used for testing Haskell code for years and also been proven to be useful in the industry (Hughes 2007), thus we investigate its potential for ABS, which to our best knowledge has not been done yet.

Property-based testing has a close connection to model-checking (McMillan 1993), where properties of a system are proved in a formal way. The important difference is that the checking happens directly on code and not on the abstract, formal model, thus one can say that it combines model-checking and unit-testing, embedding it directly in the software-development and TDD process without an intermediary step. We hypothesise that adding it to the already existing testing methods in the field of ABS is of substantial value as it allows to cover a much wider range of test-cases due to automatic data generation. This can be used in two ways: to verify an implementation against a formal specification or to test hypotheses about an implemented simulation. This puts property-based testing on the same level as agent- and system testing, where not technical implementation details of e.g. agents are checked like in unit-tests but their individual complete behaviour and the system behaviour as a whole.

The work (Onggo and Karatas 2016) explicitly mentions the problem of test coverage which would often require to write a large number of tests manually to cover the parameter ranges sufficiently enough - property-based testing addresses exactly this problem by *automating* the test-data generation. Note that this is closely related to data-generators (Gurcan, Dikenelli, and Bernon 2013) and load generators and random testing (Burnstein 2010) but property-based testing goes one step further by integrating this into a specification language directly into code, emphasising a declarative approach and pushing the generators behind the scenes, making them transparent and focusing on the specification rather than on the data-generation.

To substantiate and test our claims, we present two case-studies. First, the agent-based SIR model (Macal 2010), which is of explanatory nature, where we show how to express formal model-specifications in property-tests. Second, the SugarScape model (Epstein and Axtell 1996), which is of exploratory nature, where we show how to express hypotheses in property-tests.

The aim and contribution of this paper is the investigation of the potential of pure functional property-based testing for ABS using Haskell as programming language. Further we will show that by simply using a pure functional programming language removes a large class of run-time errors and allows much stronger guarantees of correctness already at compile time, increasing the confidence in the correctness of the simulation up to a new level.

The structure of the paper is as follows: First we present related work in Section 2. To make this paper sufficiently self-contained, we introduce pure functional programming in Haskell on a conceptual level in Section 3. Then we give a more in-depth explanation of property-based testing in Section 4. Next we shortly present existing research on *how* to implement ABS in Haskell and conceptually apply property-based testing in Section 5. The heart of the paper are the two case-studies, which we present in Section 6 and 7. Finally we conclude in Section 8 and discuss further research in Section 9.

2 RELATED WORK

Research on TDD of ABS is quite new and thus there exist relative few publications. The work (Collier and Ozik 2013) is the first to discuss how to apply the TDD approach to ABS, using unit-testing to verify the correctness of the implementation up to a certain level. They show how to implement unit-tests within the RePast Framework (North, Collier, Ozik, Tatara, Macal, Bragen, and Sydelko 2013) and make the important point that such a software need to be designed to be sufficiently modular otherwise testing becomes too cumbersome and involves too many parts. The paper (Asta, Özcan, and Peer-Olaf 2014) discusses a similar approach to DES in the AnyLogic software toolkit.

The paper (Onggo and Karatas 2016) proposes Test Driven Simulation Modelling (TDSM) which combines techniques from TDD to simulation modelling. The authors present a case study for maritime search-operations where they employ ABS. They emphasise that simulation modelling is an iterative process, where changes are made to existing parts, making a TDD approach to simulation modelling a good match. They present how to validate their model against analytical solutions from theory using unit-tests by running the whole simulation within a unit-test and then perform a statistical comparison against a formal specification. This approach will become of importance later on in our SIR case study.

The paper (Brambilla, Pincioli, Birattari, and Dorigo 2012) propose property-driven design of robot swarms. They propose a top-down approach by specifying properties a swarm of robots should have from which a prescriptive model is created, which properties are verified using model checking. Then a simulation is implemented following this prescriptive and verified model after then the physical robots are implemented. The authors identify the main difficulty of implementing such a system that the engineer must *"think at the collective-level, but develop at the individual-level"*. It is arguably true that this also applies to implementing agent-based models and simulations where the same collective-individual separation exists from which emergent system behaviour of simulations emerges - this is the very foundation of the ABS methodology.

The paper (Gurcan, Dikenelli, and Bernon 2013) gives an in-depth and detailed overview over verification, validation and testing of agent-based models and simulations and proposes a generic framework for it. The authors present a generic UML class model for their framework which they then implement in the two ABS frameworks RePast and MASON. Both of them are implemented in Java and the authors provide a detailed description how their generic testing framework architecture works and how it utilises JUnit to run automated tests. To demonstrate their framework they provide also a case study of an agent-base simulation of synaptic connectivity where they provide an in-depth explanation of their levels of test together with code.

The review of the literature in the field gives the impression, that most research focuses on high-level validation and does not deal too much with verification on a technical, code-base level.

Although the work on TDD is scarce in ABS, there exists quite some research on applying TDD and unit-testing to multi-agent systems (MAS). Although MAS is a different discipline than ABS, the latter one has derived many technical concepts from the former one thus testing concepts applied to MAS might also be applicable to ABS. The paper (Nguyen, Perini, Bernon, Pavón, and Thangarajah 2011) is a survey of testing in MAS. It distinguishes between unit tests which tests units that make up an agent, agent tests which test the combined functionality of units that make up an agent, integration tests which test the interaction of agents within an environment and observe emergent behaviour, system test which test the MAS as a system running at the target environment and acceptance test in which stakeholders verify that the software meets their goal. Although not all ABS simulations need acceptance and system tests, still this classification gives a good direction and can be directly transferred to ABS.

3 PURE FUNCTIONAL PROGRAMMING

Although property-based is now available in a wide range of programming languages and paradigms, including Java, Python and C++, it has its origins in Haskell. We argue that for that reason property-based testing really shines in pure functional programming, thus we conduct all implementation and research of this paper using Haskell. Therefore we give a brief introduction into the concepts of pure functional programming in Haskell without going into too much technical detail. Further we will show that the use of Haskell automatically increases the confidence in the correctness of an ABS implementation due to its fundamentally different nature. Also it emphasises loose coupled programming to a much stronger extent that does object-oriented programming as in Java, Python and C++, something the authors of (Collier and Ozik 2013) emphasise to be able to properly test agent behaviour. We argue that due to its fundamental different

nature, the functional programming paradigm makes making mistakes much harder, resulting in simulations which are more likely to be correct than implementations with existing object-oriented approaches.

Functional programming makes functions the main concept of programming, promoting them to first-class citizens. Its roots lie in the Lambda Calculus which was first described by Alonzo Church (Church 1936). This is a fundamentally different approach to computation than imperative and object-oriented programming which roots lie in the Turing Machine (Turing 1937). Rather than describing *how* something is computed as in the more operational approach of the Turing Machine, due to the more declarative nature of the Lambda Calculus, code in functional programming describes *what* is computed.

In this paper we are using the functional programming language Haskell. The paper of (Hudak, Hughes, Peyton Jones, and Wadler 2007) gives a comprehensive overview over the history of the language, how it developed and its features and is very interesting to read and get accustomed to the background of the language. The main points why we decided to go for Haskell are:

- Rich Feature-Set - it has all fundamental concepts of the pure functional programming paradigm of which we explain the most important below.
- Real-World applications - the strength of Haskell has been proven through a vast amount of highly diverse real-world applications (Hudak, Hughes, Peyton Jones, and Wadler 2007), is applicable to a number of real-world problems (O'Sullivan, Goerzen, and Stewart 2008) and has a large number of libraries available ¹.
- Modern - Haskell is constantly evolving through its community and adapting to keep up with the fast changing field of computer science. Further, the community is the main source of high-quality libraries.

As a short example we give an implementation of the factorial function in Haskell:

```
factorial :: Integer -> Integer
factorial 0 = 1
factorial n = n * factorial (n-1)
```

When looking at this function we can already see the central concepts of functional programming:

1. Declarative - we describe *what* the factorial function is rather than how to compute it. This is supported by *pattern matching* which allows to give multiple equations for the same function, matching on its input.
2. Immutable data - in functional programming we don't have mutable variables - after a variable is assigned, it cannot change its contents. This also means that there is no destructive assignment operator which can re-assign values to a variable. To change values, we employ recursion.
3. Recursion - the function calls itself with a smaller argument and will eventually reach the case of 0. Recursion is the very meat of functional programming because they are the only way to implement loops in this paradigm due to immutable data.
4. Static Types - the first line indicates the name and the types of the function. In this case the function takes one Integer as input and returns an Integer as output. Types are static in Haskell which means that there can be no type-errors at run-time e.g. when one tries to cast one type into another because this is not supported by this kind of type-system.

¹https://wiki.haskell.org/Applications_and_libraries

5. Explicit input and output - all data which are required and produced by the function have to be explicitly passed in and out of it. There exists no global mutable data whatsoever and data-flow is always explicit.
6. Referential transparency - calling this function with the same argument will *always* lead to the same result, meaning one can replace this function by its value. This means that when implementing this function one can not read from a file or open a connection to a server. This is also known as *purity* and is indicated in Haskell in the types which means that it is also guaranteed by the compiler.

It may seem that one runs into efficiency-problems in Haskell when using algorithms which are implemented in imperative languages through mutable data which allows in-place update of memory. The seminal work of (Okasaki 1999) showed that when approaching this problem from a functional mind-set this does not necessarily be the case. The author presents functional data structures which are asymptotically as efficient as the best imperative implementations and discusses the estimation of the complexity of lazy programs.

For an excellent and widely used introduction to programming in Haskell we refer to (Hutton 2016). Other, more exhaustive books on learning Haskell are (Lipovaca 2011, Allen and Moronuki 2016). For an introduction to programming with the Lambda-Calculus we refer to (Michaelson 2011). For more general discussion of functional programming we refer to (Hughes 1989, MacLennan 1990, Hudak, Hughes, Peyton Jones, and Wadler 2007).

3.1 Side-Effects

One of the fundamental strengths of functional programming and Haskell is their way of dealing with side-effects in functions. A function with side-effects has observable interactions with some state outside of its explicit scope. This means that the behaviour depends on history and that it loses its referential transparency character, which makes understanding and debugging much harder. Examples for side-effects are (amongst others): modifying a global variable, modifying a variable through a reference, await an input from the keyboard, read or write to a file, open a connection to a server, drawing random-numbers,...

Obviously, to write real-world programs which interact with the outside-world we need side-effects. Haskell allows to indicate in the *type* of a function that it does or does *not* have side-effects. Further there are a broad range of different effect types available, to restrict the possible effects a function can have to only the required type. This is then ensured by the compiler which means that a program in which one tries to e.g. read a file in a function which only allows drawing random-numbers will fail to compile. Haskell also provides mechanisms to combine multiple effects e.g. one can define a function which can draw random-numbers and modify some global data. The most common side-effect types are: *IO* allows all kind of I/O related side-effects: reading/writing a file, creating threads, write to the standard output, read from the keyboard, opening network-connections, mutable references,...; *Rand* allows to draw random-numbers from a random-number stream, *Reader* allows to read from a read-only environment, *Writer* allows to write to a write-only environment, *State* allows to read and write a read/write environment.

A function with side-effects has to indicate this in their type e.g. if we want to give our factorial function for debugging purposes the ability to write to the standard output, we add *IO* to its type: `factorial :: Integer -> IO Integer`. A function without any side-effect type is called *pure*. A function with a given effect-type needs to be executed with a given effect-runner which takes all necessary parameters depending on the effect and runs a given effectful function returning its return value and depending on the effect also an effect-related result. For example when running a function with a *State*-effect one needs to specify the initial environment which can be read and written. After running such a function with a *State*-effect the effect-runner returns the changed environment in addition with the return value of the function itself. Note

that we cannot call functions of different effect-types from a function with another effect-type, which would violate the guarantees. Calling a *pure* function though is always allowed because it has by definition no side-effects. An effect-runner itself is a *pure* function. The exception to this is the IO effect type which does not have a runner but originates from the *main* function which is always of type IO.

Although it might seem very restrictive at first, we get a number of benefits from making the type of effects we can use explicit. First we can restrict the side-effects a function can have to a very specific type which is guaranteed at compile time. This means we can have much stronger guarantees about our program and the absence of potential errors already at compile-time which implies that we don't need test them with e.g. unit-tests. Second, because effect-runners are themselves *pure*, we can execute effectful functions in a very controlled way by making the effect-context explicit in the parameters to the effect-runner. This allows a much easier approach to isolated testing because the history of the system is made explicit.

For a technical, in-depth discussion of the concept of side-effects and how they are implemented in Haskell using Monads, we refer to the following papers: (Moggi 1989, Wadler 1992, Wadler 1995, Wadler 1997, Jones 2002).

3.2 Why

In general, Types guide us in program construction by restricting the operations we can perform on the data. This means that by choosing types this reveals already a lot of our program and data and prevents us from making mistakes e.g. interpreting some binary data as text instead of a number. In strongly statically typed languages the types can do this already at compile-time which allows to rule out certain bugs already at compile-time. In general, we can say that for all bugs which can be ruled out at compile-time, we don't need to write property- or unit-tests, because those bugs cannot - per definition - occur at run-time, so it won't make sense to test their absence at run-time. Also, as Dijkstra famously put it: "Testing shows the presence, not the absence of bugs" - thus by induction we can say that compile-time guarantees save us from a potentially infinite amount of testing.

In general it is well established, that pure functional programming as in Haskell, allows to express much stronger guarantees about the correctness of a program *already at compile-time*. This is in fundamental contrast to imperative object-oriented languages like Java or Python where only primitive guarantees about types - mostly relationships between type-hierarchies - can be expressed at compile-time which directly implies that one needs to perform much more testing (user testing or unit-testing) at *run-time* to check whether the model is sufficiently correct. Thus guaranteeing properties already at compile-time frees us from writing unit-tests which cover these cases or test them at run time because they are *guaranteed to be correct under all circumstances, for all inputs*.

In this regards we see pure functional programming as truly superior to the traditional object oriented approaches: they lead to implementations of models which are more likely correct because we can express more guarantees already at compile-time which directly leads to less bugs which directly increases the probability of the software being a correct implementation of the model.

In the next section we give a brief discussion of *how* to apply pure functional programming with Haskell to implement ABS.

4 PROPERTY-BASED TESTING

Property-based testing allows to formulate *functional specifications* in code which then a property-based testing library tries to falsify by *automatically* generating test-data with some user-defined coverage. When

a case is found for which the property fails, the library then reduces it to the most simple one. It is clear to see that this kind of testing is especially suited to ABS, because we can formulate specifications, meaning we describe *what* to test instead of *how* to test. Also the deductive nature of falsification in property-based testing suits very well the constructive and exploratory nature of ABS. Further, the automatic test-generation can make testing of large scenarios in ABS feasible as it does not require the programmer to specify all test-cases by hand, as is required in unit-tests.

Property-based testing was invented by the authors of (Claessen and Hughes 2000, Claessen and Hughes 2002) in which they present the QuickCheck library, which tries to falsify the specifications by *randomly* sampling the space. We argue, that the stochastic sampling nature of this approach is particularly well suited to ABS, because it is itself almost always driven by stochastic events and randomness in the agents behaviour, thus this correlation should make it straight-forward to map ABS to property-testing. The main challenge when using QuickCheck, as will be shown later, is to write *custom* test-data generators for agents and the environment which cover the space sufficiently enough to not miss out on important test-cases. According to the authors of QuickCheck *"The major limitation is that there is no measurement of test coverage."* (Claessen and Hughes 2000). QuickCheck provides help to report the distribution of test-cases but still it could be the case that simple test-cases which would fail are never tested.

As a remedy for the potential sampling difficulties of QuickCheck, there exists also a deterministic property-testing library called SmallCheck (Runciman, Naylor, and Lindblad 2008) which instead of randomly sampling the test-space, enumerates test-cases exhaustively up to some depth. It is based on two observations, derived from model-checking, that (1) *"If a program fails to meet its specification in some cases, it almost always fails in some simple case"* and (2) *"If a program does not fail in any simple case, it hardly ever fails in any case"* (Runciman, Naylor, and Lindblad 2008). This non-stochastic approach to property-based testing might be a complementary addition in some cases where the tests are of non-stochastic nature with a search-space which is too large to implement manually by unit-tests but is relatively easy and small enough to enumerate exhaustively. The main difficulty and weakness of using SmallCheck is to reduce the dimensionality of the test-case depth search to prevent combinatorial explosion, which would lead to exponential number of cases. Thus one can see QuickCheck and SmallCheck as complementary instead of in opposition to each other.

Note that in this paper we primarily focus on the use of QuickCheck due to the match of ABS stochastic nature and the random test generation. We refer to SmallCheck in cases where appropriate. Also note that we regard property-based testing as *complementary* to unit-tests and not in opposition - we see it as an addition in the TDD process of developing an ABS.

5 PURE FUNCTIONAL ABS

There does not exist much research on how to implement an ABS from a pure functional perspective using Haskell. The authors of (Thaler, Altenkirch, and Siebers 2019) probably provide the most in-depth and technical discussion by showing *how* to do it and what the benefits and drawbacks are. Their approach is based on Functional Reactive Programming which allows to express discrete- and continuous-time systems in functional programming. It primarily emphasises a time-driven approach to ABS but also allows to implement an event-driven one (Meyer 2014). Following the conclusions of the paper, we can derive the following benefits, which support directly our initial hypothesis and our claims, giving good reasons *why* to do ABS in a functional way and why it makes it more likely to be correct:

1. Run-Time robustness by compile-time guarantees - by expressing stronger guarantees already at compile-time we can restrict the classes of bugs which occur at run-time by a substantial amount due to Haskell's strong and static type system. This implies the lack of dynamic types and dynamic

casts ² which removes a substantial source of bugs. Note that we can still have run-time bugs in Haskell when our functions are partial.

2. Purity - By being explicit and polymorphic in the types about side-effects and the ability to handle side-effects explicitly in a controlled way allows to rule out non-deterministic side-effects which guarantees reproducibility due to guaranteed same initial conditions and deterministic computation. Also by being explicit about side-effects e.g. Random-Numbers and State makes it easier to verify and test.
3. Explicit Data-Flow and Immutable Data - All data must be explicitly passed to functions thus we can rule out implicit data-dependencies because we are excluding IO. This makes reasoning of data-dependencies and data-flow much easier as compared to traditional object-oriented approaches which utilize pointers or references.
4. Declarative - describing *what* a system is, instead of *how* (imperative) it works. In this way it should be easier to reason about a system and its (expected) behaviour because it is more natural to reason about the behaviour of a system instead of thinking of abstract operational details.
5. Concurrency and parallelism - due to its pure and 'stateless' nature, functional programming is extremely well suited for massively large-scale applications as it allows adding parallelism without any side-effects and provides very powerful and convenient facilities for concurrent programming. The paper of (TODO: cite my own paper on STM) explores the use Haskell for concurrent and parallel ABS in a deeper way.

Although pure functional ABS as in Haskell allows us to leverage on the concepts of functional and its benefits (and drawbacks) we still rely heavily on (property-based) testing to ensure correctness of a simulation because our approach still can have run-time bugs.

5.1 Debugging

TODO: haskell-titan TODO: Testing and Debugging Functional Reactive Programming (Perez and Nilsson 2017)

General there are the following basic verification & validation requirements to ABS (Robinson 2014), which all can be addressed in our *pure* functional approach as described in the paper in Appendix ??:

- Fixing random number streams to allow simulations to be repeated under same conditions - ensured by *pure* functional programming and Random Monads
- Rely only on past - guaranteed with *Arrowized* FRP
- Bugs due to implicitly mutable state - reduced using pure functional programming
- Ruling out external sources of non-determinism / randomness - ensured by *pure* functional programming
- Deterministic time-delta - ensured by *pure* functional programming
- Repeated runs lead to same dynamics - ensured by *pure* functional programming

5.2 Property-Based ABS Testing

TODO: general approach to property-based testing in ABS

²Note that there exist casts between different numerical types but they are all safe and can never lead to errors at run-time.

Although (pure) functional programming allows us to have stronger guarantees about the behaviour and absence of bugs of the simulation already at compile-time, we still need to test all the properties of our simulation which we cannot guarantee at compile-time.

We found property-based testing particularly well suited for ABS. Although it is now available in a wide range of programming languages and paradigms, property-based testing has its origins in Haskell (Claessen and Hughes 2000, Claessen and Hughes 2002) and we argue that for that reason it really shines in pure functional programming. Property-based testing allows to formulate *functional specifications* in code which then the property-testing library (e.g. QuickCheck (Claessen and Hughes 2000)) tries to falsify by automatically generating random test-data covering as much cases as possible. When an input is found for which the property fails, the library then reduces it to the most simple one. It is clear to see that this kind of testing is especially suited to ABS, because we can formulate specifications, meaning we describe *what* to test instead of *how* to test (again the declarative nature of functional programming shines through). Also the deductive nature of falsification in property-based testing suits very well the constructive nature of ABS.

Generally we need to distinguish between two types of testing/verification: 1. testing/verification of models for which we have real-world data or an analytical solution which can act as a ground-truth - examples for such models are the SIR model, stock-market simulations, social simulations of all kind and 2. testing/verification of models which are just exploratory and which are only be inspired by real-world phenomena - examples for such models are Epsteins Sugarscape and Agent_Zero.

5.2.1 Black-Box Verification

In black-box Verification one generally feeds input and compares it to expected output. In the case of ABS we have the following examples of black-box test:

1. Isolated Agent Behaviour - test isolated agent behaviour under given inputs using and property-based testing.
2. Interacting Agent Behaviour - test if interaction between agents are correct .
3. Simulation Dynamics - compare emergent dynamics of the ABS as a whole under given inputs to an analytical solution or real-world dynamics in case there exists some using statistical tests.
4. Hypotheses- test whether hypotheses are valid or invalid using and property-based testing.

Using black-box verification and property-based testing we can apply for the following use cases for testing ABS in FRP:

Finding optimal Δt The selection of the right Δt can be quite difficult in FRP because we have to make assumptions about the system a priori. One could just play it safe with a very conservative, small $\Delta t < 0.1$ but the smaller Δt , the lower the performance as it multiplies the number of steps to calculate. Obviously one wants to select the *optimal* Δt , which in the case of ABS is the largest possible Δt for which we still get the correct simulation dynamics. To find out the *optimal* Δt one can make direct use of the black-box tests: start with a large $\Delta t = 1.0$ and reduce it by half every time the tests fail until no more tests fail - if for $\Delta t = 1.0$ tests already pass, increasing it may be an option. It is important to note that although isolated agent behaviour tests might result in larger Δt , in the end when they are run in the aggregate system, one needs to sample the whole system with the smallest Δt found amongst all tests. Another option would be to apply super-sampling to just the parts which need a very small Δt but this is out of scope of this paper.

Agents as signals Agents *might* behave as signals in FRP which means that their behaviour is completely determined by the passing of time: they only change when time changes thus if they are a signal they should

stay constant if time stays constant. This means that they should not change in case one is sampling the system with $\Delta t = 0$. Of course to prove whether this will *always* be the case is strictly speaking impossible with a black-box verification but we can gain a good level of confidence with them also because we are staying pure. It is only through white-box verification that we can really guarantee and prove this property.

5.2.2 White-Box Verification

White-Box verification is necessary when we need to reason about properties like *forever*, *never*, which cannot be guaranteed from black-box tests. Additional help can be coverage tests with which we can show that all code paths have been covered in our tests.

TODO: List of Common Bugs and Programming Practices to avoid them (Vipindeep and Jalote 2005)

We have discussed in this section *how* to approach an ABS implementation from a pure functional perspective using Haskell where we have also briefly touched on *why* one should do so and what the benefits and drawbacks are. In the next two sections we will expand on the *why* by presenting two case-studies which show the benefits of using Haskell in regards of testing and increasing the confidence in the correctness of the implementation.

6 CASE STUDY I: SIR

As an example we discuss the black-box testing for the SIR model using property-testing. We test if the *isolated* behaviour of an agent in all three states Susceptible, Infected and Recovered, corresponds to model specifications. The crucial thing though is that we are dealing with a stochastic system where the agents act *on averages*, which means we need to average our tests as well. We conducted the tests on the implementation found in the paper of Appendix ??.

6.0.1 Black-Box Verification

The interface of the agent behaviours are defined below. When running the SF with a given Δt one has to feed in the state of all the other agents as input and the agent outputs its state it is after this Δt .

```
data SIRState
  = Susceptible
  | Infected
  | Recovered

type SIRAgent = SF [SIRState] SIRState

susceptibleAgent :: RandomGen g => g -> SIRAgent
infectedAgent   :: RandomGen g => g -> SIRAgent
recoveredAgent  :: RandomGen g => g -> SIRAgent
```

Susceptible Behaviour A susceptible agent *may* become infected, depending on the number of infected agents in relation to non-infected the susceptible agent has contact to. To make this property testable we run a susceptible agent for 1.0 time-unit (note that we are sampling the system with a smaller $\Delta t = 0.1$) and then check if it is infected - that is it returns infected as its current state.

Obviously we need to pay attention to the fact that we are dealing with a stochastic system thus we can only talk about averages and thus it does not suffice to only run a single agent but we are repeating this for e.g. $N = 10.000$ agents (all with different RNGs). We then need a formula for the required fraction of the N agents which should have become infected on average. Per 1.0 time-unit, a susceptible agent makes *on average* contact with β other agents where in the case of a contact with an infected agent the susceptible agent becomes infected with a given probability γ . In this description there is another probability hidden, which is the probability of making contact with an infected agent which is simply the ratio of number of infected agents to number not infected agents. The formula for the target fraction of agents which become infected is then: $\beta * \gamma * \frac{\text{number of infected}}{\text{number of non-infected}}$. To check whether this test has passed we compare the required amount of agents which on average should become infected to the one from our tests (simply count the agents which got infected and divide by N) and if the value lies within some small ϵ then we accept the test as passed.

Obviously the input to the susceptible agents which we can vary is the set of agents with which the susceptible agents make contact with. To save us from constructing all possible edge-cases and combinations and testing them with unit-tests we use property-testing with QuickCheck which creates them randomly for us and reduces them also to all relevant edge-cases. This is an example for how to use property-based testing in ABS where QuickCheck can be of immense help generating random test-data to cover all cases.

Infected Behaviour An infected agent *will always* recover after a finite time, which is *on average* after δ time-units. Note that this property involves stochastics too, so to test this property we run a large number of infected agents e.g. $N = 10.000$ (all with different RNGs) until they recover, record the time of each agents recovery and then average over all recovery times. To check whether this test has passed we compare the average recovery times to δ and if they lie within some small ϵ then we accept the test as passed.

We use property-testing with QuickCheck in this case as well to generate the set of other agents as input for the infected agents. Strictly speaking this would not be necessary as an infected agent never makes contact with other agents and simply ignores them - we could as well just feed in an empty list. We opted for using QuickCheck for the following reasons:

- We wanted to stick to the interface specification of the agent-implementation as close as possible which asks to pass the states of all agents as input.
- We shouldn't make any assumptions about the actual implementation and if it REALLY ignores the other agents, so we strictly stick to the interface which requires us to input the states of all the other agents.
- The set of other agents is ignored when determining whether the test has failed or not which indicates by construction that the behaviour of an infected agent does not depend on other agents.
- We are not just running a single replication over 10.000 agents but 100 of them which should give black-box verification more strength.

Recovered Behaviour A recovered agent will stay in the recovered state *forever*. Obviously we cannot write a black-box test that truly verifies that because it had to run in fact forever. In this case we need to resort to white-box verification (see below).

Because we use multiple replications in combination with QuickCheck obviously results in longer test-runs (about 5 minutes on my machine) In our implementation we utilized the FRP paradigm. It seems that functional programming and FRP allow extremely easy testing of individual agent behaviour because FP and FRP compose extremely well which in turn means that there are no global dependencies as e.g. in OOP where we have to be very careful to clean up the system after each test - this is not an issue at all in our *pure* approach to ABS.

Simulation Dynamics We won't go into the details of comparing the dynamics of an ABS to an analytical solution, that has been done already by (Macal 2010). What is important is to note that population-size matters: different population-size results in slightly different dynamics in SD => need same population size in ABS (probably...?). Note that it is utterly difficult to compare the dynamics of an ABS to the one of a SD approach as ABS dynamics are stochastic which explore a much wider spectrum of dynamics e.g. it could be the case, that the infected agent recovers without having infected any other agent, which would lead to an extreme mismatch to the SD approach but is absolutely a valid dynamic in the case of an ABS. The question is then rather if and how far those two are *really* comparable as it seems that the ABS is a more powerful system which presents many more paths through the dynamics.

Finding optimal Δt Obviously the *optimal* Δt of the SIR model depends heavily on the model parameters: contact rate β and illness duration δ . We fixed them in our tests to be $\beta = 5$ and $\delta = 15$. By using the isolated behaviour tests we found an optimal $\Delta t = 0.125$ for the susceptible behaviour and $\Delta t = 0.25$ for the infected behaviour.

Agents as signals Our SIR agents *are* signals due to the underlying continuous nature of the analytical SIR model and to some extent we can guarantee this through black-box testing. For this we write tests for each individual behaviour as previously but instead of checking whether agents got infected or have recovered we assume that they stay constant: they will output always the same state when sampling the system with $\Delta t = 0$. The tests are conceptual the complementary tests of the previous behaviour tests so in conjunction with them we can assume to some extent that agents are signals. To prove it, we need to look into white-box verification as we cannot make guarantees about properties which should hold *forever* in a computational setting.

6.0.2 White-Box Verification

In the case of the SIR model we have the following invariants:

- A susceptible agent will *never* make the transition to recovered.
- An infected agent will *never* make the transition to susceptible.
- A recovered agent will *forever* stay recovered.

All these invariants can be guaranteed when reasoning about the code. An additional help will be then coverage testing with which we can show that an infected agent never returns susceptible, and a susceptible agent never returned infected given all of their functionality was covered which has to imply that it can never occur!

We will only look at the recovered behaviour as it is the simplest one. We leave the susceptible and infected behaviours for further research / the final thesis because the conceptual idea becomes clear from looking at the recovered agent.

Recovered Behaviour The implementation of the recovered behaviour is as follows:

```
recoveredAgent :: SIRAgent
recoveredAgent = arr (const Recovered)
```

Just by looking at the type we can guarantee the following:

- it is pure, no side-effects of any kind can occur
- no stochasticity possible because no RNG is fed in / we don't run in the random monad

The implementation is as concise as it can get and we can reason that it is indeed a correct implementation of the recovered specification: we lift the constant function which returns the Recovered state into an arrow. Per definition and by looking at the implementation, the constant function ignores its input and returns always the same value. This is exactly the behaviour which we need for the recovered agent. Thus we can reason that the recovered agent will return Recovered *forever* which means our implementation is indeed correct.

7 CASE STUDY II: SUGARSCAPE

TODO

8 CONCLUSIONS

TODO

Although property-based testing has its origin in Haskell, frameworks exist now in other languages as well e.g. Java, Python, C++ and we hope that our research sparked an interest in applying property-based testing to the established object-oriented languages in ABS as well.

9 FURTHER RESEARCH

TODO

Generally speaking, a dependent type is a type whose definition depends on a value e.g. when we have a pair of natural numbers where the second one is greater than the first, we speak of a dependent type. With this power, they allow to push compile-time guarantees to a new level where we can express nearly arbitrary complex guarantees at compile-time because we can *compute types at compile-time*. This means that types are first-class citizen of the language and go as far as being formal proofs of the correctness of an implementation, allowing to narrow the gap between specification and implementation substantially.

We hypothesise that the use of dependent types allows us to push the judgement of the correctness of a simulation to new, unprecedented level, not possible with the established object-oriented approaches so far. This has the direct consequence that the development process is very different and can reduce the amount of testing (both unit-testing and manual testing) substantially. Because one is implementing a simulation which is (as much as possible) correct-by-construction, the correctness (of parts) can be guaranteed statically.

Summarizing, we expect the following benefits from adding dependent types to ABS:

1. Narrowing the gap between the model specification and its implementation reduces the potential for conceptual errors in model-to-code translation.
2. Less number of tests required due to guarantees being expressed already at compile time.
3. Higher confidence in correctness due to formal guarantees in code.

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