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		</

```

33     return ret;
34 }

K-Dimension Tree

1 struct Point { double x,y; int id; }P[maxn];
2 struct Rectangle
3 {
4     double lx,rx,ly,ry;
5     inline Rectangle() = default;
6     inline Rectangle(double _lx,double _rx,double _ly,double _ry):lx(_lx),rx(_rx),ly(_ly),ry(_ry) {}
7     inline void set(const Point &p) { lx = rx = p.x; ly = ry = p.y; }
8     inline void merge(const Point &p)
9     {
10         lx = min(lx,p.x); rx = max(rx,p.x);
11         ly = min(ly,p.y); ry = max(ry,p.y);
12     }
13     inline void merge(const Rectangle &r)
14     {
15         lx = min(lx,r.lx); rx = max(rx,r.rx);
16         ly = min(ly,r.ly); ry = max(ry,r.ry);
17     }
18     // 最小距离, 到 4 个角和 4 条边距离
19     inline double dist(const Point &p)
20     {
21         if (p.x <= lx&&p.y <= ly) return (p-Point(lx,ly)).norm();
22         else if (p.x <= rx&&p.y <= ly) return p.y-ly;
23         else if (p.x >= rx&&p.y <= ly) return (p-Point(rx,ly)).norm();
24         else if (p.x >= rx&&p.y <= ry) return p.x-rx;
25         else if (p.x >= rx&&p.y >= ry) return (p-Point(rx,ry)).norm();
26         else if (p.x >= lx&&p.y >= ry) return p.y-ry;
27         else if (p.x <= lx&&p.y >= ry) return (p-Point(lx,ry)).norm();
28         else if (p.x <= lx&&p.y >= ly) return p.x-lx;
29         return 0;
30     }
31     // 最大距离, 到 4 个角的距离
32     inline double dist(const Point &p)
33     {
34         double ret = 0;
35         ret += max((rx-p.x)*(rx-p.x),(lx-p.x)*(lx-p.x));
36         ret += max((ry-p.y)*(ry-p.y),(ly-p.y)*(ly-p.y));
37         return ret;
38     }
39 };

struct Node {
40     int child[2]; Point p; Rectangle r;
41     inline Node() = default;
42     inline Node(const Point &p,const Rectangle &r):p(p),r(r) { r.set(p); memset(child,0,8); }
43     inline void set(const Point &p) { p = _p; r.set(p); memset(child,0,8); }
44 }tree[maxn];

45 inline bool cmpx(const Point &a,const Point &b) {
46     if (a.x != b.x) return a.x < b.x;
47     else return a.y < b.y;
48 }
49 inline bool cmpy(const Point &a,const Point &b) {
50     if (a.y != b.y) return a.y < b.y;
51     else return a.x < b.x;
52 }
53 inline bool cmp(pair <double,int> a,pair <double,int> b) {
54     int sgn = dcmp(a.first-b.first);
55     if (sgn) return sgn < 0;
56     else return a.second < b.second;
57 }
58 // 查询 k 大/小
59 inline void query(int now,const Point &p,int k,pair <double,int> ret[],bool dim = false) {
60     if (dcmp(tree[now].r.dist(p)-ret[k].first) > 0) return;
61     pair <double,int> val = make_pair((p-tree[now].p).norm(),tree[now].p.id);
62     for (int i = 1;i <= k;++i) if (cmp(val,ret[i])) {
63         for (int j = k+1;j > i;--j) ret[j] = ret[j-1];
64         ret[i] = val; break;
65     }
66     if ((dim&&cmpx(p,tree[now].p))||(!dim&&cmpy(p,tree[now].p))) {
67         if (tree[now].child[0]) query(tree[now].child[0],p,k,ret,dim^1);
68         if (tree[now].child[1]) query(tree[now].child[1],p,k,ret,dim^1);
69     } else {
70         if (tree[now].child[1]) query(tree[now].child[1],p,k,ret,dim^1);
71         if (tree[now].child[0]) query(tree[now].child[0],p,k,ret,dim^1);
72     }
73 }
74 inline int build(int l,int r,bool dim) {
75     // ...
76 }
77 
```

```

78 int now = ++size,mid = (l+r)>>1;
79 nth_element(vec.begin()+l-1,vec.begin()+mid-1,vec.begin()+r,dim?cmpx:cmpy);
80 tree[now].set(vec[mid-1]);
81 if (l < mid) {
82     tree[now].child[0] = build(l,mid-1,dim^1);
83     tree[now].r.merge(tree[tree[now].child[0]].r);
84 }
85 if (r > mid) {
86     tree[now].child[1] = build(mid+1,r,dim^1);
87     tree[now].r.merge(tree[tree[now].child[1]].r);
88 }
89 return now;
90 }

```

## KM

```

1 // Truly O(n^3), 最大权匹配
2 // 邻接矩阵, 不能连的边设为 -INF, 求最小权匹配时边权取负, 但不能连的还是 -INF, 使用时先对 1 -> n 调用 hungary(), 再
3 struct KM
4 {
5     int w[maxn][maxn],lx[maxn],ly[maxn],match[maxn],way[maxn],slack[maxn];
6     bool used[maxn];
7
8     inline void init()
9     {
10         for (int i = 1;i <= N;++i)
11             match[i] = lx[i] = ly[i] = way[i] = 0;
12     }
13
14     inline void hungary(int x)
15     {
16         match[0] = x; int j0 = 0;
17         for (int j = 0;j <= N;++j)
18             slack[j] = inf,used[j] = false;
19         do
20         {
21             used[j0] = true;
22             int i0 = match[j0],delta = inf,j1 = 0;
23             for (int j = 1;j <= N;++j)
24                 if (!used[j])
25                 {
26                     int cur = -w[i0][j]-lx[i0]-ly[j];
27                     if (cur < slack[j])
28                         slack[j] = cur,way[j] = j0;
29                     if (slack[j] < delta)
30                         delta = slack[j],j1 = j;
31                 }
32             for (int j = 0;j <= N;++j)
33             {
34                 if (used[j]) lx[match[j]] += delta,ly[j] -= delta;
35                 else slack[j] -= delta;
36             }
37             j0 = j1;
38         } while (match[j0]);
39         do
40         {
41             int j1 = way[j0];
42             match[j0] = match[j1];
43             j0 = j1;
44         } while (j0);
45     }
46
47     inline void work() { for (int i = 1;i <= N;++i) hungary(i); }
48
49     inline int get_ans()
50     {
51         int sum = 0;
52         for (int i = 1;i <= N;++i)
53         {
54             // if (w[match[i]][i] == -inf); //无解
55             if (match[i] > 0) sum += w[match[i]][i];
56         }
57         return sum;
58     }
59 }km;
60
61 
```

## Link Cut Tree

```

1 inline bool isroot(int a) { return ch[fa[a]][0] != a&&ch[fa[a]][1] != a; }
2 
```

```

3 inline void update(int x) { val[x] = (val[ch[x][0]]+val[ch[x][1]]).merge(x); }
4 inline void pushdown(int x)
5 {
6     if (rev[x])
7     {
8         int &lc = ch[x][0], &rc = ch[x][1];
9         swap(lc, rc);
10        if (lc) rev[lc] ^= 1;
11        if (rc) rev[rc] ^= 1;
12        rev[x] = false;
13    }
14 }
15
16 inline void rotate(int x)
17 {
18     int y = fa[x], z = fa[y], l = ch[y][1] == x, r = l^1;
19     if (!isroot(y)) ch[z][ch[z][1] == y] = x; fa[x] = z;
20     if (ch[x][r]) fa[ch[x][r]] = y; ch[y][l] = ch[x][r];
21     fa[y] = x; ch[x][r] = y; update(y); update(x);
22 }
23
24 inline void splay(int x)
25 {
26     int top = 0, i;
27     for (i = x; !isroot(i); i = fa[i]) stk[++top] = i; stk[++top] = i;
28     while (top) pushdown(stk[top--]);
29     while (!isroot(x))
30     {
31         int y = fa[x], z = fa[y];
32         if (!isroot(y))
33         {
34             if ((ch[y][0] == x) ^ (ch[z][0] == y)) rotate(x);
35             else rotate(y);
36         }
37         rotate(x);
38     }
39 }
40
41 inline int access(int x)
42 {
43     int t = 0;
44     for (t = 0; x; t = x, x = fa[x])
45         splay(x), ch[x][1] = t, update(x);
46     return t;
47 }
48
49 inline int evert(int x) { int t; rev[t = access(x)] ^= 1; return t; }
50
51 inline int find(int x)
52 {
53     x = access(x);
54     while (pushdown(x), ch[x][0]) x = ch[x][0];
55     return x;
56 }
57
58 inline void cut(int x, int y)
59 {
60     evert(x); access(y); splay(y);
61     if (ch[y][0] != x || ch[x][1] != 0) return;
62     ch[y][0] = fa[x] = 0; update(x); update(y);
63 }
64
65 inline void link(int x, int y) { fa[evert(x)] = y; }

```

## Numerical Integration

```

1 //self-adapt simpson
2 inline long double simpson(long double l, long double r, long double mid, long double C1, long double
   ↪ Cr, long double Cm)
3 {
4     long double tC1 = calc((l+mid)/2), tCr = calc((mid+r)/2);
5     long double ans = (r-l)*(C1+Cr+4*Cm)/6, lans = (mid-l)*(C1+Cm+4*tC1)/6, rans = (r-mid)*(Cr+Cm+4*tCr)/6;
6     if (r-l <= 1e-3 && fabs(lans+rans-ans) < eps) return ans;
7     // if (dep > lim && fabs(lans+rans-ans) < eps) return ans;
8     else return simpson(l, mid, (l+mid)/2, C1, Cm, tC1) + simpson(mid, r, (mid+r)/2, Cm, Cr, tCr);
9 }

```

## Splay

```

1 const int maxn = 500010, inf = 1<<29;
2 int N, M, root, cnt, arr[maxn], tag[maxn], key[maxn], fa[maxn], ch[maxn][2], lb[maxn], rb[maxn];
3 int wb[maxn], sum[maxn], size[maxn], stk[maxn]; bool rev[maxn]; char cmd[20]; queue<int> team;
4
5 inline int gi()
6 {
7     char ch; int ret = 0, f = 1;
8     do ch = getchar(); while (!(ch >= '0' && ch <= '9') && ch != '-');

```

```

9     if (ch == '-') f = -1, ch = getchar();
10    do ret = ret*10 + ch - '0', ch = getchar(); while (ch >= '0' && ch <= '9');
11    return ret*f;
12 }
13
14 inline int newnode(int x = 0)
15 {
16     int ret;
17     if (!team.empty())
18         ret = team.front(), team.pop();
19     else ret = ++cnt;
20     key[ret] = sum[ret] = lb[ret] = rb[ret] = wb[ret] = x;
21     rev[ret] = false; tag[ret] = inf; size[ret] = 1;
22     return ret;
23 }
24
25 inline void pushdown(int now) { }
26
27 inline void update(int now)
28 {
29     // pushdown(now);
30     int lc = ch[now][0], rc = ch[now][1];
31     size[now] = size[lc] + size[rc] + 1;
32     sum[now] = sum[lc] + sum[rc] + key[now];
33     if (lc && rc) { }
34     else if (lc) { }
35     else if (rc) { }
36     else { }
37 }
38
39 inline int build(int l, int r)
40 {
41     int mid = (l+r) >> 1, ret = newnode(arr[mid]);
42     if (l < mid) ch[ret][0] = build(l, mid-1), fa[ch[ret][0]] = ret;
43     if (r > mid) ch[ret][1] = build(mid+1, r), fa[ch[ret][1]] = ret;
44     update(ret); return ret;
45 }
46
47 inline void init()
48 {
49     root = newnode(); ch[root][1] = newnode(); fa[2] = 1;
50     for (int i = 1; i <= N; ++i) arr[i] = gi();
51     ch[2][0] = build(1, N); fa[ch[2][0]] = 2;
52     update(2); update(1);
53 }
54
55 inline int find(int rk)
56 {
57     for (int now = root;;)
58     {
59         pushdown(now);
60         if (rk == size[ch[now][0]] + 1) return now;
61         else if (rk > size[ch[now][0]] + 1)
62             rk -= size[ch[now][0]] + 1, now = ch[now][1];
63         else now = ch[now][0];
64     }
65     return 0;
66 }
67
68 inline void rotate(int x)
69 {
70     int y = fa[x], z = fa[y], l = ch[y][0] != x, r = l^1;
71     if (z) ch[z][ch[z][0] != y] = x;
72     fa[x] = z; fa[y] = x; fa[ch[x][r]] = y;
73     ch[y][l] = ch[x][r]; ch[x][r] = y;
74     update(y); update(x);
75 }
76
77 inline void splay(int x, int aim)
78 {
79     int top = 0;
80     for (int i = x; i != fa[i]; i = fa[i]) stk[++top] = i;
81     while (top) pushdown(stk[top--]);
82     while (fa[x] != aim)
83     {
84         int y = fa[x], z = fa[y];
85         if (z != aim)
86         {
87             if ((ch[y][0] == x) ^ (ch[z][0] == y)) rotate(x);
88             else rotate(y);
89         }
90         rotate(x);

```

```

90     }
91     if (!aim) root = x;
92 }
93
94 inline void Delete(int &now)
95 {
96     if (!now) return;
97     Delete(ch[now][0]);
98     Delete(ch[now][1]);
99     team.push(now); now = 0;
100 }
101
102 inline void print()
103 {
104     for (int i = 1; i <= cnt; ++i)
105         printf("%d:%d %d\n", i, ch[i][0], ch[i][1]);
106     for (int i = 1; i <= cnt; ++i)
107         printf("%d:%d\n", i, fa[i]);
108 }
109 }
110
111 inline void laydown(int now)
112 {
113     if (!now) return;
114     pushdown(now);
115     laydown(ch[now][0]);
116     printf("%d ", key[now]);
117     laydown(ch[now][1]);
118     update(now);
119 }

```

## 几何基础

```

1 //计算几何常用公式
2 struct Point { double x,y; inline Point unit() const {
3     double len = norm(); if (!dcmp(len)) return Point(1,0); else return *this/len; }
4     inline Point reflect(const Point &p) const {
5         Point v = *this-p; double len = v.norm();
6         v = v/len; return p+v*(1/len); }
7     inline Point vertical() const { return Point(y,-x); }
8     inline double angle() const {
9         double ret = atan2(y,x); if (ret < 0) ret += 2*pi; return ret; }
10 };
11 struct Line
12 {
13     Point p,v; double slop; inline Line() = default;
14     inline Line(const Point &_p,const Point &_v):p(_p),v(_v) {}
15     inline void update() { slop = v.alpha(); }
16     friend inline bool operator <(const Line &l1,const Line &l2)
17     { return l1.slop < l2.slop; }
18     inline double dis(const Point &a) { fabs((a-p)/v)/(v.len()); } //点到直线距离
19 };
20
21 inline bool OnLine(const Line &l,const Point &p) { return !dcmp(l.v/(p-l.p)); } //点在直线上
22
23 inline Point CrossPoint(const Line &a,const Line &b) //直线交点
24 { Point u = a.p - b.p; double t = (b.v/u)/(a.v/b.v); return a.p+a.v*t; }
25
26 inline bool parallel(const Line &a,const Line &b) { return !dcmp(a.v/b.v); } //直线平行

```

## 凸包

```

1 struct Point
2 {
3     inline Point() = default;
4     inline Point(double _x,double _y):x(_x),y(_y) {}
5     inline Point unit() const
6     {
7         double len = norm();
8         if (!dcmp(len)) return Point(1,0);
9         else return *this/len;
10    }
11    inline double norm() const { return sqrt(x*x+y*y); }
12    inline Point reflect(const Point &p) const
13    {
14        Point v = *this-p; double len = v.norm();
15        v = v/len; return p+v*(1/len);
16    }
17    inline void read() { scanf("%lf %lf",&x,&y); }
18    inline Point vertical() const { return Point(y,-x); }
19    inline double angle() const
20    {

```

```

21     double ret = atan2(y,x);
22     if (ret < 0) ret += 2*pi;
23     return ret;
24 }
25 friend inline bool operator ==(const Point &a,const Point &b) { return
    ↪ !dcmp(a.x-b.x)&&!dcmp(a.y-b.y); }
26 friend inline Point operator -(const Point &a,const Point &b) { return Point(a.x-b.x,a.y-b.y); }
27 friend inline Point operator +(const Point &a,const Point &b) { return Point(a.x+b.x,a.y+b.y); }
28 friend inline Point operator /(const Point &a,double b) { return Point(a.x/b,a.y/b); }
29 friend inline Point operator *(const Point &a,double b) { return Point(a.x*b,a.y*b); }
30 friend inline Point operator *(double b,const Point &a) { return Point(a.x*b,a.y*b); }
31 friend inline double operator /(const Point &a,const Point &b) { return a.x*b.y-a.y*b.x; }
32 friend inline bool operator <(const Point &a,const Point &b)
33 {
34     if (a.x != b.x) return a.x < b.x;
35     else return a.y < b.y;
36 }
37 }P[maxn],convex[maxn];
38
39 inline void ConvexHull()
40 {
41     sort(P+1,P+N+1); //x 第一关键字, y 第二关键字从小到大排序
42     for (int i = 1; i <= N; ++i)
43     {
44         while (m > 1&&(convex[m]-convex[m-1])/(P[i]-convex[m-1]) <= 0) --m;
45         convex[++m] = P[i];
46     }
47     int k = m;
48     for (int i = N-1; i; --i)
49     {
50         while (m > k&&(convex[m]-convex[m-1])/(P[i]-convex[m-1]) <= 0) --m;
51         convex[++m] = P[i];
52     }
53     if (N > 1) m--;
54 }

```

## 无向图最小割

```

1 int node[N], dist[N];
2 bool visit[N];
3 int solve(int n) {
4     int answer = INT_MAX;
5     for (int i = 0; i < n; ++i) node[i] = i;
6     while (n > 1) {
7         int max = 1;
8         for (int i = 0; i < n; ++i) {
9             dist[node[i]] = graph[node[0]][node[i]];
10            if (dist[node[i]] > dist[node[max]]) max = i;
11        }
12        int prev = 0;
13        memset(visit, 0, sizeof(visit));
14        visit[node[0]] = true;
15        for (int i = 1; i < n; ++i) {
16            if (i == n-1) {
17                answer = std::min(answer, dist[node[max]]);
18                for (int k = 0; k < n; ++k) {
19                    graph[node[k]][node[prev]] = (graph[node[prev]][node[k]] += graph[node[k]][node[max]]);
20                }
21                node[max] = node[--n];
22            }
23            visit[node[max]] = true;
24            prev = max; max = -1;
25            for (int j = 1; j < n; ++j) {
26                if (!visit[node[j]]) {
27                    dist[node[j]] += graph[node[prev]][node[j]];
28                    if (max == -1 || dist[node[max]] < dist[node[j]]) {
29                        max = j;
30                    }
31                }
32            }
33        }
34    }
35    return answer;
36 }

```

## 匈牙利

```

1 //匈牙利算法
2 //Version1
3 inline bool find(int x)
4 {

```

```

5   if (cor[x]) return false;
6   for (int i = side[x]; i = next[i]) if (!used[toit[i]])
7   {
8       used[toit[i]] = true;
9       if (!cho[toit[i]] || find(cho[toit[i]]))
10      {
11          cho[toit[i]] = x; map[x] = toit[i];
12          return true;
13      }
14  }
15  return false;
16 }
17
18 inline void hungry()
19 {
20     for (int i = 1; i <= p; ++i)
21         memset(used, false, sizeof(used)), find(i);
22     for (int i = 1; i <= m; ++i)
23     {
24         memset(used, false, sizeof(used)), cho[map[i]] = 0;
25         find(i), cor[i] = true;
26     }
27 }
28 //Version2
29 inline int find(int x)
30 {
31     for (int i = 1; i <= n; ++i)
32         if (f[x][i] && !used[i])
33         {
34             used[i] = true;
35             if (!cho[i] || find(cho[i])) { cho[i] = x; return true; }
36         }
37     return false;
38 }
39
40 inline int hungry()
41 {
42     int ret = 0;
43     for (int i = 1; i <= n; ++i)
44     {
45         memset(used, false, sizeof(used));
46         if (find(i)) ret++;
47     }
48     return ret;
49 }

```

## 半平面交

```

1 //半平面交，直线左侧半平面，注意最后是 tail-head <= 0 还是 tail-head <= 1
2 inline int dcmp(double a)
3 {
4     if (-eps <= a && a <= eps) return 0;
5     else if (a > 0) return 1; else return -1;
6 }
7
8 struct Point
9 {
10     double x, y;
11     inline Point() = default;
12     inline Point(double _x, double _y): x(_x), y(_y) {}
13     inline void read() { x = gi(), y = gi(); }
14     inline Point vertical() const { return Point(-y, x); }
15     inline Point unit() const
16     {
17         double len = norm();
18         if (!dcmp(len)) return Point(1, 0);
19         else return *this / len;
20     }
21     inline double norm() const { return sqrt(x*x+y*y); }
22     inline double angle() const { return atan2(y, x); }
23     friend inline Point operator+(const Point &a, const Point &b) { return Point(a.x+b.x, a.y+b.y); }
24     friend inline Point operator-(const Point &a, const Point &b) { return Point(a.x-b.x, a.y-b.y); }
25     friend inline Point operator*(const Point &a, double b) { return Point(a.x*b, a.y*b); }
26     friend inline Point operator*(double b, const Point &a) { return Point(a.x*b, a.y*b); }
27     friend inline double operator/(const Point &a, const Point &b) { return a.x*b.y-a.y*b.x; }
28 } P[maxn], pp[maxn], pol[maxn];
29
30 struct Line
31 {
32     Point p, v;
33     inline Line(const Point _p = Point(), const Point _v = Point()): p(_p), v(_v) {}
34     inline double slop() const { return v.angle(); }

```

```

35     friend inline bool operator<(const Line &a, const Line &b) { return a.slop() < b.slop(); }
36 } line[maxn], qq[maxn];
37
38 inline bool onleft(const Line &L, const Point &p)
39 {
40     return dcmp(L.v/(p-L.p)) > 0;
41 }
42 inline bool parallel(const Line &a, const Line &b) { return !dcmp(a.v/b.v); }
43 inline Point crosspoint(const Line &a, const Line &b)
44 {
45     Point u = a.p-b.p;
46     double t = (b.v/u)/(a.v/b.v);
47     return a.p+(a.v*t);
48 }
49
50 inline int half_plane_intersection()
51 {
52     sort(lines+1, lines+tot+1); //直线按斜率排序
53     int head, tail;
54     qq[head = tail = 1] = lines[1];
55     for (int i = 2; i <= tot; ++i)
56     {
57         while (head < tail && !onleft(lines[i], pp[tail-1])) --tail;
58         while (head < tail && !onleft(lines[i], pp[head])) ++head;
59         qq[++tail] = lines[i];
60         if (parallel(qq[tail], qq[tail-1]))
61         {
62             tail--;
63             if (onleft(qq[tail], lines[i].p)) qq[tail] = lines[i];
64         }
65         if (head < tail) pp[tail-1] = crosspoint(qq[tail], qq[tail-1]);
66     }
67     while (head < tail && !onleft(qq[head], pp[tail-1])) --tail;
68     if (tail-head <= 0) return 0;
69     pp[tail] = crosspoint(qq[tail], qq[head]);
70     for (int i = head; i <= tail; ++i) pol[++m] = pp[i]; //半平面交点
71     pol[0] = pol[m];
72     return m;
73 }

```

## 后缀数组

```

1 inline void build(char *buf, int *Sa, int *Rank, int *Height, int n, int now, int m)
2 {
3     int i, j, k, *x = t1, *y = t2;
4     memset(c, 0, 4*m);
5     for (i = 0; i < n; ++i) c[x[i] = buf[i] - 'A']++;
6     for (i = 1; i < m; ++i) c[i] += c[i-1];
7     for (i = n-1; i >= 0; --i) Sa[--c[x[i]]] = i;
8     for (k = 1; k < n; k <= 1)
9     {
10         int p = 0;
11         for (i = n-k; i < n; ++i) y[p++] = i;
12         for (i = 0; i < n; ++i) if (Sa[i] >= k) y[p++] = Sa[i] - k;
13         memset(c, 0, 4*m);
14         for (i = 0; i < n; ++i) c[x[y[i]]]++;
15         for (i = 1; i < m; ++i) c[i] += c[i-1];
16         for (i = n-1; i >= 0; --i) Sa[--c[x[y[i]]]] = y[i];
17         swap(x, y); p = 1; x[Sa[0]] = 0;
18         for (i = 1; i < n; ++i)
19             x[Sa[i]] = y[Sa[i-1]] == y[Sa[i]] && y[Sa[i-1]+k] == y[Sa[i]+k] ? p-1 : p++;
20         if (p >= n) break; m = p;
21     }
22     for (i = 0; i < n; ++i) Rank[Sa[i]] = i;
23     for (i = k = 0; i < n; ++i)
24     {
25         if (k) --k; if (!Rank[i]) continue;
26         j = Sa[Rank[i]-1];
27         while (i+k < n && j+k < n && buf[i+k] == buf[j+k]) ++k;
28         Height[Rank[i]] = k;
29     }
30 }

```

## 圆与多边形面积交

```

1 const int maxn = 510;
2 const double eps = 1e-9;
3
4 inline int dcmp(double a)
5 {
6     if (a > eps) return 1;

```



```

7     else if (a < -eps) return -1;
8     else return 0;
9 }
10
11 struct Point
12 {
13     double x,y;
14     Point() = default;
15     Point(double _x,double _y):x(_x),y(_y) {}
16     inline double norm() const { return sqrt(x*x+y*y); }
17     inline Point unit() const { double len = norm(); return Point(x/len,y/len); }
18     friend Point operator +(const Point &a,const Point &b) { return Point(a.x+b.x,a.y+b.y); }
19     friend Point operator -(const Point &a,const Point &b) { return Point(a.x-b.x,a.y-b.y); }
20     friend Point operator *(const Point &a,double b) { return Point(a.x*b,a.y*b); }
21     friend Point operator *(double b,const Point &a) { return Point(a.x*b,a.y*b); }
22     friend Point operator /(const Point &a,double b) { return Point(a.x/b,a.y/b); }
23     friend double operator /(const Point &a,const Point &b) { return a.x*b.y-b.x*a.y; }
24     friend double operator *(const Point &a,const Point &b) { return a.x*b.x+a.y*b.y; }
25     inline void read() { scanf("%lf %lf",&x,&y); }
26 }P[maxn],A,B;
27 int N; double K;
28
29 inline double getSectorArea(const Point &a,const Point &b,double r)
30 {
31     double c = (2*r*r-((a-b)*(a-b)))/(2*r*r);
32     double alpha = acos(c);
33     return r*r*alpha/2.0;
34 }
35
36 inline pair <double,double> getSolution(double a,double b,double c)
37 {
38     double delta = b*b-4*a*c;
39     if (dcmp(delta) < 0) return make_pair(0,0);
40     else return make_pair((-b-sqrt(delta))/(2*a),(-b+sqrt(delta))/(2*a));
41 }
42
43 inline pair <Point,Point> getIntersection(const Point &a,const Point &b,double r)
44 {
45     Point d = b-a;
46     double A = d*d,B = 2*(d*a),C = (a*a)-r*r;
47     pair <double,double> s = getSolution(A,B,C);
48     return make_pair(a+(d*s.first),a+(d*s.second));
49 }
50
51 inline double getPointDist(const Point &a,const Point &b)
52 {
53     Point d = b-a;
54     int sA = dcmp(a*d),sB = dcmp(b*d);
55     if (sA*sB <= 0) return (a/b)/((a-b).norm());
56     else return min(a.norm(),b.norm());
57 }
58
59 double getArea(const Point &a,const Point &b,double r)
60 {
61     double dA = a*a, dB = b*b, dC = getPointDist(a,b), ans = 0;
62     if (dcmp(dA-r*r) <= 0 && dcmp(dB-r*r) <= 0) return (a/b)/2;
63     Point tA = a.unit()*r, tB = b.unit()*r;
64     if (dcmp(dC-r) > 0) return getSectorArea(tA,tB,r);
65     pair <Point,Point> ret = getIntersection(a,b,r);
66     if (dcmp(dA-r*r) > 0 && dcmp(dB-r*r) > 0)
67     {
68         ans += getSectorArea(tA,ret.first,r);
69         ans += (ret.first/ret.second)/2;
70         ans += getSectorArea(ret.second,tB,r);
71     }
72     if (dcmp(dA-r*r) > 0) return (ret.first/b)/2+getSectorArea(tA,ret.first,r);
73     else return (a/ret.second)/2.0+getSectorArea(ret.second,tB,r);
74 }
75
76 double getArea(int n,Point *p,const Point &c,double r)
77 {
78     double ret = 0;
79     for (int i = 0;i < n;++i)
80     {
81         int sgn = dcmp((p[i]-c)/(p[(i+1)%n]-c));
82         if (sgn > 0) ret += getArea(p[i]-c,p[(i+1)%n]-c,r);
83         else ret -= getArea(p[(i+1)%n]-c,p[i]-c,r);
84     }
85     return fabs(ret);
86 }
87 }

```

## 圆的 K 次交

```

1 //modified
2 const double eps = 1e-7,pi = acos(-1.0);
3 int N,M; double area[maxn]; // area[k] -> area of intersections >= k.
4
5 inline int dcmp(double a)
6 {
7     if (-eps <= a && a <= eps) return 0;
8     else if (a > 0) return 1; else return -1;
9 }
10
11 struct Point
12 {
13     double x,y;
14     inline Point() = default;
15     inline Point(double _x,double _y):x(_x),y(_y) {}
16     inline void read() { x = gi(),y = gi(); }
17     inline double norm() const { return sqrt(x*x+y*y); }
18     inline double angle() const { return atan2(y,x); }
19     inline Point unit() const { double len = norm(); return Point(x/len,y/len); }
20     friend inline Point operator-(const Point &a,const Point &b) { return Point(a.x-b.x,a.y-b.y); }
21     friend inline Point operator+(const Point &a,const Point &b) { return Point(a.x+b.x,a.y+b.y); }
22     friend inline Point operator*(const Point &a,double b) { return Point(a.x*b,a.y*b); }
23     friend inline Point operator*(double b,const Point &a) { return Point(a.x*b,a.y*b); }
24     friend inline Point operator/(const Point &a,double b) { return Point(a.x/b,a.y/b); }
25     friend inline double operator/(const Point &a,const Point &b) { return a.x*b.y-a.y*b.x; }
26 };
27 struct Circle
28 {
29     Point C; double r; int sgn;
30     inline Circle() = default;
31     inline Circle(const Point &_C,double _r,int _sgn):C(_C),r(_r),sgn(_sgn) {}
32     // sgn 代表该圆的权值，默认 1
33     friend inline bool operator==(const Circle &a,const Circle &b)
34     {
35         if (dcmp(a.r-b.r)) return false;
36         if (dcmp(a.C.x-b.C.x)) return false;
37         if (dcmp(a.C.y-b.C.y)) return false;
38         if (a.sgn != b.sgn) return false;
39         return true;
40     }
41     friend inline bool operator!=(const Circle &a,const Circle &b) { return !(a == b); }
42 }cir[maxn];
43
44 inline Point rotate(const Point &p,double cost,double sint)
45 {
46     double x = p.x,y = p.y;
47     return Point(x*cost-y*sint,x*sint+y*cost);
48 }
49 inline pair <Point,Point> crosspoint(const Point &ap,double ar,const Point &bp,double br)
50 {
51     double d = (ap-bp).norm(),cost = (ar*ar+d*br*br)/(2*ar*d),sint = sqrt(1-cost*cost);
52     Point v = ((bp-ap).unit()*ar;
53     return make_pair(ap+rotate(v,br,br),ap+rotate(v,-sint,br));
54 }
55 inline pair <Point,Point> crosspoint(const Circle &a,const Circle &b) { return
56     crosspoint(a.C,a.r,b.C,b.r); }
57
58 inline bool overlap(const Circle &a,const Circle &b) { return dcmp(a.r-b.r-(a.C-b.C).norm()) >= 0; }
59 // b 是不是在 a 里面
60 inline bool intersect(const Circle &a,const Circle &b)
61 {
62     if (overlap(a,b)) return false;
63     if (overlap(b,a)) return false;
64     return dcmp((a.C-b.C).norm()-a.r-b.r) < 0;
65 }
66
67 struct Event
68 {
69     Point p; double a; int d;
70     inline Event() = default;
71     inline Event(const Point &p,double _a,double _d):p(_p),a(_a),d(_d) {}
72     friend inline bool operator <(const Event &a,const Event &b) { return a.a < b.a; }
73 };
74
75 inline void solve()
76 {
77     for (int i = 1;i <= M;++i) area[i] = 0;
78     for (int i = 1;i <= M;++i)

```

```

76 {
77     int cnt = cir[i].sgn; if (cnt < 0) cnt = 0; vector<Event> event;
78     for (int j = 1; j < i; ++j) if (cir[i] == cir[j]) cnt += cir[j].sgn;
79     for (int j = 1; j <= M; ++j)
80         if (j != i && cir[i] != cir[j] && !overlap(cir[j], cir[i])) cnt += cir[j].sgn;
81     for (int j = 1; j <= M; ++j)
82         if (j != i && !intersect(cir[i], cir[j]))
83         {
84             pair<Point, Point> res = crosspoint(cir[i], cir[j]); swap(res.first, res.second);
85             double alpha1 = (res.first - cir[i].C).angle(), alpha2 = (res.second - cir[i].C).angle();
86             event.push_back(Event(res.second, alpha2, cir[j].sgn));
87             event.push_back(Event(res.first, alpha1, -cir[j].sgn));
88             cnt += (alpha2 > alpha1) * cir[j].sgn;
89         }
90     if (!event.size()) area[cnt] += pi * cir[i].r * cir[i].r * cir[i].sgn;
91     else
92     {
93         sort(event.begin(), event.end());
94         event.push_back(event.front());
95         for (int j = 0; j+1 < (int)event.size(); ++j)
96         {
97             cnt += event[j].d;
98             area[cnt] += event[j].p / event[j+1].p / 2 * cir[i].sgn;
99             double alpha = event[j+1].a - event[j].a;
100             if (alpha < 0) alpha += 2 * pi;
101             if (!dcmp(alpha)) continue;
102             area[cnt] += alpha * cir[i].r * cir[i].r / 2 * cir[i].sgn;
103             area[cnt] += -sin(alpha) * cir[i].r * cir[i].r / 2 * cir[i].sgn;
104         }
105     }
106 }
107 }
108
109 // origin
110 struct Event {
111     Point p;
112     double ang;
113     int delta;
114     Event(Point p = Point(0, 0), double ang = 0, double delta = 0) : p(p), ang(ang), delta(delta) {}
115 };
116 bool operator < (const Event &a, const Event &b) {
117     return a.ang < b.ang;
118 }
119 void addEvent(const Circle &a, const Circle &b, vector<Event> &evt, int &cnt) {
120     double d2 = (a.o - b.o).len2(),
121     dRatio = ((a.r - b.r) * (a.r + b.r) / d2 + 1) / 2,
122     pRatio = sqrt(-(d2 - sqrt(a.r - b.r)) * (d2 - sqrt(a.r + b.r)) / (d2 * d2 * 4));
123     Point d = b.o - a.o, p = d.rotate(PI / 2),
124     q0 = a.o + d * dRatio + p * pRatio,
125     q1 = a.o + d * dRatio - p * pRatio;
126     double ang0 = (q0 - a.o).ang(),
127     ang1 = (q1 - a.o).ang();
128     evt.push_back(Event(q1, ang1, 1));
129     evt.push_back(Event(q0, ang0, -1));
130     cnt += ang1 > ang0;
131 }
132 bool issame(const Circle &a, const Circle &b) { return sign((a.o - b.o).len()) == 0 && sign(a.r - b.r)
    == 0; }
133 bool overlap(const Circle &a, const Circle &b) { return sign(a.r - b.r - (a.o - b.o).len()) >= 0; }
134 bool intersect(const Circle &a, const Circle &b) { return sign((a.o - b.o).len() - a.r - b.r) < 0; }
135 Circle c[N];
136 double area[N]; // area[k] -> area of intersections >= k.
137 Point centroid[N]; // k 次圆的质心
138 bool keep[N];
139 void add(int cnt, DB a, Point c) {
140     area[cnt] += a;
141     centroid[cnt] = centroid[cnt] + c * a;
142 }
143 void solve(int C) {
144     for (int i = 1; i <= C; ++i) {
145         area[i] = 0;
146         centroid[i] = Point(0, 0);
147     }
148     for (int i = 0; i < C; ++i) {
149         int cnt = 1;
150         vector<Event> evt;
151         for (int j = 0; j < i; ++j) if (issame(c[i], c[j])) ++cnt;
152         for (int j = 0; j < C; ++j) {
153             if (j != i && !issame(c[i], c[j]) && !overlap(c[j], c[i])) {
154                 ++cnt;
155             }

```

```

156     }
157     for (int j = 0; j < C; ++j) {
158         if (j != i && !overlap(c[j], c[i]) && !overlap(c[i], c[j]) && intersect(c[i], c[j])) {
159             addEvent(c[i], c[j], evt, cnt);
160         }
161     }
162     if (evt.size() == 0) {
163         add(cnt, PI * c[i].r * c[i].r, c[i].o);
164     } else {
165         sort(evt.begin(), evt.end());
166         evt.push_back(evt.front());
167         for (int j = 0; j + 1 < (int)evt.size(); ++j) {
168             cnt += evt[j].delta;
169             add(cnt, det(evt[j].p, evt[j+1].p) / 2, (evt[j].p + evt[j+1].p) / 3);
170             double ang = evt[j+1].ang - evt[j].ang;
171             if (ang < 0) {
172                 ang += PI * 2;
173             }
174             if (sign(ang) == 0) continue;
175             double ang0 = evt[j].a, ang1 = evt[j+1].a;
176             add(cnt, ang * c[i].r * c[i].r / 2, c[i].o +
177                 Point(sin(ang1) - sin(ang0), -cos(ang1) + cos(ang0)) * (2 / (3 * ang) * c[i].r));
178             add(cnt, -sin(ang) * c[i].r * c[i].r / 2, (c[i].o + evt[j].p + evt[j+1].p) / 3);
179         }
180     }
181 }
182 for (int i = 1; i <= C; ++i)
183     if (sign(area[i])) {
184         centroid[i] = centroid[i] / area[i];
185     }
186 }

```

### 最小覆盖圆

```

1 circle minimum_circle(vector<point> p) {
2     circle ret;
3     random_shuffle(p.begin(), p.end());
4     for (int i = 0; i < (int)p.size(); ++i)
5         if (!in_circle(p[i], ret)) {
6             ret = circle(p[i], 0);
7             for (int j = 0; j < i; ++j)
8                 if (!in_circle(p[j], ret)) {
9                     ret = make_circle(p[j], p[i]);
10                    for (int k = 0; k < j; ++k)
11                        if (!in_circle(p[k], ret)) ret = make_circle(p[i], p[j], p[k]);
12                }
13        }
14     return ret;
15 }

```

### 三维凸包

```

1 struct Triangle { // Construction function removed.
2     TPoint a, b, c;
3     double getArea() {
4         TPoint ret = det(b - a, c - a);
5         return dist(ret) / 2.0;
6     }
7 };
8 namespace Convex_Hull {
9     struct Face { // Construction function removed.
10         int a, b, c;
11         bool isOnConvex;
12     };
13     int nFace, left, right, whe[MAXN][MAXN];
14     Face queue[MAXF], tmp[MAXF];
15     bool isVisible(const std::vector<TPoint> &p, const Face &f, const TPoint &a) {
16         return dcmp(detdot(p[f.a], p[f.b], p[f.c], a)) > 0;
17     }
18     bool init(std::vector<TPoint> &p) {
19         bool check = false;
20         for (int i = 1; i < (int)p.size(); ++i) {
21             if (dcmp(sqrdist(p[0], p[i])) > 0) {
22                 std::swap(p[1], p[i]);
23                 check = true;
24                 break;
25             }
26         }
27         if (!check) return false;
28         check = false;
29         for (int i = 2; i < (int)p.size(); ++i) {
30             if (dcmp(sqrdist(det(p[i] - p[0], p[1] - p[0]))) > 0) {

```

```

31     std::swap(p[2], p[i]);
32     check = true;
33     break;
34 }
35 }
36 if (!check) return false;
37 check = false;
38 for (int i = 3; i < (int)p.size(); i++) {
39     if (dcmp(detdot(p[0], p[1], p[2], p[i]))) {
40         std::swap(p[3], p[i]);
41         check = true;
42         break;
43     }
44 }
45 if (!check) return false;
46 for (int i = 0; i < (int)p.size(); i++)
47     for (int j = 0; j < (int)p.size(); j++) {
48         whe[i][j] = -1;
49     }
50 return true;
51 }
52 void pushface(const int &a, const int &b, const int &c) {
53     nFace++;
54     tmp[nFace] = Face(a, b, c);
55     tmp[nFace].isOnConvex = true;
56     whe[a][b] = nFace;
57     whe[b][c] = nFace;
58     whe[c][a] = nFace;
59 }
60 bool deal(const std::vector<TPoint> &p, const std::pair<int, int> &now, const TPoint &base) {
61     int id = whe[now.second][now.first];
62     if (!tmp[id].isOnConvex) return true;
63     if (isVisible(p, tmp[id], base)) {
64         queue[++right] = tmp[id];
65         tmp[id].isOnConvex = false;
66         return true;
67     }
68     return false;
69 }
70 std::vector<Triangle> getConvex(std::vector<TPoint> &p) {
71     static std::vector<Triangle> ret;
72     ret.clear();
73     if (!init(p)) return ret;
74     if (!isVisible(p, Face(0,1,2),p[3])) pushface(0,1,2); else pushface(0,2,1);
75     if (!isVisible(p, Face(0,1,3),p[2])) pushface(0,1,3); else pushface(0,3,1);
76     if (!isVisible(p, Face(0,2,3),p[1])) pushface(0,2,3); else pushface(0,3,2);
77     if (!isVisible(p, Face(1,2,3),p[0])) pushface(1,2,3); else pushface(1,3,2);
78     for (int a = 4; a < (int)p.size(); a++) {
79         TPoint base = p[a];
80         for (int i = 1; i <= nFace; i++) {
81             if (tmp[i].isOnConvex && isVisible(p, tmp[i], base)) {
82                 left = 0, right = 0;
83                 queue[++right] = tmp[i];
84                 tmp[i].isOnConvex = false;
85                 while (left < right) {
86                     Face now = queue[++left];
87                     if (!deal(p, std::make_pair(now.a, now.b), base)) pushface(now.a, now.b, a);
88                     if (!deal(p, std::make_pair(now.b, now.c), base)) pushface(now.b, now.c, a);
89                     if (!deal(p, std::make_pair(now.c, now.a), base)) pushface(now.c, now.a, a);
90                 }
91                 break;
92             }
93         }
94     }
95     for (int i = 1; i <= nFace; i++) {
96         Face now = tmp[i];
97         if (now.isOnConvex) ret.push_back(Triangle(p[now.a], p[now.b], p[now.c]));
98     }
99     return ret;
100 }
101 }
102 // Usage
103 std::vector<TPoint> p;
104 std::vector<Triangle> answer;
105 answer = Convex_Hull::getConvex(p);

```

### 三维绕轴旋转

注意事项：以右手拇指为向量方向，逆时针绕轴（剩下四根手指方向）旋转  $\theta$  角的右乘矩阵。

```

1 Matrix getTrans(const double &a, const double &b, const double &c) {
2     Matrix ret;

```

```

3     ret.a[0][0] = 1; ret.a[0][1] = 0; ret.a[0][2] = 0; ret.a[0][3] = 0;
4     ret.a[1][0] = 0; ret.a[1][1] = 1; ret.a[1][2] = 0; ret.a[1][3] = 0;
5     ret.a[2][0] = 0; ret.a[2][1] = 0; ret.a[2][2] = 1; ret.a[2][3] = 0;
6     ret.a[3][0] = a; ret.a[3][1] = b; ret.a[3][2] = c; ret.a[3][3] = 1;
7     return ret;
8 }
9 Matrix getRotate(const double &a, const double &b, const double &c, const double &theta) {
10     Matrix ret;
11     ret.a[0][0] = a * a * (1 - cos(theta)) + cos(theta);
12     ret.a[0][1] = a * b * (1 - cos(theta)) + c * sin(theta);
13     ret.a[0][2] = a * c * (1 - cos(theta)) - b * sin(theta);
14     ret.a[0][3] = 0;
15     ret.a[1][0] = b * a * (1 - cos(theta)) - c * sin(theta);
16     ret.a[1][1] = b * b * (1 - cos(theta)) + cos(theta);
17     ret.a[1][2] = b * c * (1 - cos(theta)) + a * sin(theta);
18     ret.a[1][3] = 0;
19     ret.a[2][0] = c * a * (1 - cos(theta)) + b * sin(theta);
20     ret.a[2][1] = c * b * (1 - cos(theta)) - a * sin(theta);
21     ret.a[2][2] = c * c * (1 - cos(theta)) + cos(theta);
22     ret.a[2][3] = 0;
23     ret.a[3][0] = 0; ret.a[3][1] = 0; ret.a[3][2] = 0; ret.a[3][3] = 1;
24     return ret;
25 }
26 Matrix getRotate(const double &ax, const double &ay, const double &az, const double &bx, const double
    ↳ &by, const double &bz, const double &theta) {
27     double l = dist(Point(0, 0, 0), Point(bx, by, bz));
28     Matrix ret = getTrans(-ax, -ay, -az);
29     ret = ret * getRotate(bx / l, by / l, bz / l, theta);
30     ret = ret * getTrans(ax, ay, az);
31     return ret;
32 }

```

### 平面图

```

1 // 包括平面图转对偶图
2 inline int dcmp(double a)
3 {
4     if (fabs(a) <= eps) return 0;
5     else if (a > 0) return 1;
6     else return -1;
7 }
8 struct Point
9 {
10     double x, y;
11     inline Point(double _x = 0, double _y = 0):x(_x), y(_y) {}
12     inline void read() { x = gi(), y = gi(); }
13     friend inline Point operator-(const Point &a, const Point &b) { return Point(a.x-b.x, a.y-b.y); }
14     friend inline double operator/(const Point &a, const Point &b) { return a.x*b.y-a.y*b.x; }
15     inline double angle() { return atan2(y, x); }
16 } pp[maxn];
17 struct Segment
18 {
19     int from, to, h, id, sur; // from 号点到 to 号点, h 为边权, suf 为这条有向边出来的平面编号。
20     inline Segment(int _from = 0, int _to = 0, int _h = 0, int _id = 0, int _sur =
        ↳ 0):from(_from), to(_to), h(_h), id(_id), sur(_sur) {}
21     friend inline bool operator<(const Segment &a, const Segment &b) { return
        ↳ (pp[a.to]-pp[a.from]).angle() < (pp[b.to]-pp[b.from]).angle(); }
22 } edge[maxm*2];
23 vector<int> G[maxn];
24
25 inline void nadd(int u, int v, int h) { ++ncnt; G[u].push_back(ncnt); edge[ncnt] = Segment(u, v, h); }
26 inline void nins(int u, int v, int h) { nadd(u, v, h); nadd(v, u, h); }
27
28 inline bool cmp(int a, int b) { return edge[a] < edge[b]; }
29
30 inline void find_surface()
31 {
32     for (int i = 1; i <= N; ++i) sort(G[i].begin(), G[i].end(), cmp);
33     for (int i = 1; i <= N; ++i)
34     {
35         int nn = G[i].size();
36         for (int j = 0; j < nn; ++j)
37             edge[G[i][j]].id = j;
38     }
39     for (int i = 2; i <= ncnt; ++i)
40         if (!edge[i].sur)
41         {
42             ++tot; int j = i, p, nn; vector<Point> vec;
43             while (!edge[j].sur)
44             {
45                 edge[j].sur = tot; vec.push_back(pp[edge[j].from]);

```



```

46     p = edge[j].to; nn = G[p].size();
47     j ^= 1; j = G[p][(edge[j].id+1)%nn];
48 }
49 double res = 0; nn = vec.size();
50 for (j = 0; j < nn; ++j)
51     res += (vec[j]-vec[0])/(vec[(j+1)%nn]-vec[0]);
52 res /= 2; space[tot] = res;
53 // 第 tot 个平面的有向面积, 外面的大平面面积为正, 其余为负, 大平面可能有多个 (平面图不连通)
54 }
55 // 开始建边, 以 mst 为例
56 // for (int i = 2; i <= cnt; i += 2)
57 // {
58 //     if (space[edge[i].sur]<0&&space[edge[i^1].sur]<0)
59 //         arr[++all] = (ARR) { edge[i].sur, edge[i^1].sur, edge[i].h };
60 //     else arr[++all] = (ARR) { edge[i].sur, edge[i^1].sur, inf };
61 // }
62 }
63 // 点定位
64 struct Scan
65 {
66     double x,y; int bel,sign;
67     inline Scan(double _x = 0, double _y = 0, int _bel = 0, int _sign = 0):x(_x),y(_y),bel(_bel),sign(_sign)
68     {}
69     friend inline bool operator < (const Scan &a, const Scan &b)
70     {
71         if (a.x != b.x) return a.x < b.x;
72         else return a.sign > b.sign;
73     }
74 }bac[maxn*4];
75 struct Splay
76 {
77     int num, root, ch[maxn][2], fa[maxn], key[maxn]; queue <int> team;
78     inline int newnode()
79     {
80         int ret;
81         if (team.empty()) ret = ++num;
82         else ret = team.front(), team.pop();
83         fa[ret] = ch[ret][0] = ch[ret][1] = 0;
84         return ret;
85     }
86 }
87 inline void init() { num = 0; root = newnode(); key[root] = cnt; }
88 inline void rotate(int x)
89 {
90     int y = fa[x], z = fa[y], l = ch[y][1] == x ? r : l^1;
91     if (z != 0) ch[z][ch[z][1] == y] = x;
92     fa[x] = z; fa[y] = x; fa[ch[x][r]] = y;
93     ch[y][l] = ch[x][r]; ch[x][r] = y;
94 }
95 inline void splay(int x)
96 {
97     while (fa[x] != 0)
98     {
99         int y = fa[x], z = fa[y];
100         if (fa[y] != 0)
101         {
102             if ((ch[y][0] == x)^(ch[z][0] == y)) rotate(x);
103             else rotate(y);
104         }
105         rotate(x);
106     }
107     root = x;
108 }
109 inline int lower_bound(const Point &p)
110 {
111     int now = root, ret = 0;
112     while (now)
113     {
114         int k = key[now];
115         if ((p-pp[edge[k].from])/(pp[edge[k].to]-pp[edge[k].from]) >= 0)
116             ret = k, now = ch[now][0];
117         else now = ch[now][1];
118     }
119     return ret;
120 }

```

```

125 inline int find(int w)
126 {
127     int now = root;
128     double x = pp[edge[w].to].x, y = pp[edge[w].to].y;
129     double ang = (pp[edge[w].to] - pp[edge[w].from]).angle();
130     while (now)
131     {
132         int k = key[now];
133         if (k == w) return now;
134         NODE p = pp[edge[k].to] - pp[edge[k].from], q = pp[edge[k].from];
135         double xx = x - q.x, yy = q.y+xx/p.x*p.y;
136         if (equal(yy, y))
137         {
138             double t = p.angle();
139             now = ch[now][ang < t];
140         }
141         else now = ch[now][y > yy];
142     }
143 }
144 inline void erase(int w)
145 {
146     int p = find(w);
147     while (ch[p][0] || ch[p][1])
148     {
149         if (ch[p][0])
150         {
151             rotate(ch[p][0]);
152             if (p == root) root = fa[p];
153         }
154         else
155         {
156             rotate(ch[p][1]);
157             if (p == root) root = fa[p];
158         }
159     }
160     team.push(p);
161     ch[fa[p]][ch[fa[p]][1] == p] = 0;
162     fa[p] = 0;
163 }
164 inline void insert(int w)
165 {
166     int now = root, pre;
167     double x = pp[edge[w].from].x, y = pp[edge[w].from].y;
168     double ang = (pp[edge[w].to] - pp[edge[w].from]).angle();
169     double xx, yy;
170     while (true)
171     {
172         int k = key[now];
173         NODE p = pp[edge[k].to] - pp[edge[k].from], q = pp[edge[k].from];
174         xx = x - q.x, yy = q.y+xx/p.x*p.y;
175         if (equal(yy, y))
176         {
177             double t = p.angle();
178             pre = now, now = ch[now][ang > t];
179             if (!now)
180             {
181                 now = newnode();
182                 fa[now] = pre; ch[pre][ang > t] = now; key[now] = w;
183                 break;
184             }
185         }
186         else
187         {
188             pre = now, now = ch[now][y > yy];
189             if (!now)
190             {
191                 now = newnode();
192                 fa[now] = pre; ch[pre][y>yy] = now; key[now] = w;
193                 break;
194             }
195         }
196     }
197     splay(now);
198 }
199 }
200 }S;
201 inline void locate()
202 {
203     // ...
204 }
205

```

```

206 int nn = 0;
207 for (int i = 2; i <= cnt; i += 2)
208 {
209     if (!dcmp(pp[edge[i].from].x - pp[edge[i].to].x)) continue;
210     bac[+nn] = Scan(pp[edge[i].from].x, pp[edge[i].from].y, i, 2);
211     bac[+nn] = Scan(pp[edge[i].to].x, pp[edge[i].to].y, i, 3);
212 }
213 scanf("%d", &T); double x, y;
214 // 查询 (x, y) 所在平面
215 for (int i = 1; i <= T; ++i)
216 {
217     scanf("%lf %lf", &x, &y);
218     bac[+nn] = Scan(x, y, i, 0);
219     scanf("%lf %lf", &x, &y);
220     bac[+nn] = Scan(x, y, i, 1);
221 }
222 sort(bac + 1, bac + nn + 1);
223 pp[+n] = Point(-oo, -oo); pp[+n] = (oo, oo);
224 edge[+cnt] = Edge(n - 1, n);
225 S.init(); int p;
226 for (int i = 1; i <= nn; ++i)
227 {
228     if (bac[i].sign == 2 || bac[i].sign == 3)
229     {
230         if (bac[i].sign == 2) S.insert(bac[i].bel);
231         else S.erase(bac[i].bel);
232     }
233     else
234     {
235         p = S.lower_bound(Point(bac[i].x, bac[i].y));
236         query[bac[i].bel][bac[i].sign] = edge[p].sur;
237     }
238 }
239 }

```

## 弦图染色最大势

```

1 #include <algorithm>
2 #include <queue>
3 #include <cstdio>
4 #include <cstdlib>
5 #include <set>
6 using namespace std;
7
8 #define maxn 10010
9 #define maxc 510
10 #define maxm 1000010
11 int tot, n, m, cnt, color[maxn][maxc], label[maxn], all;
12 int side[maxn], next[maxm * 2], toit[maxm * 2], per[maxn];
13 bool in[maxn];
14 struct node
15 {
16     int key, ord;
17     friend bool operator < (node a, node b) {return a.key > b.key; }
18 };
19 multiset <node> S;
20
21 inline void add(int a, int b)
22 {
23     next[+cnt] = side[a]; side[a] = cnt; toit[cnt] = b;
24 }
25
26 inline void ins(int a, int b) {add(a, b); add(b, a);}
27
28 inline void mcs()
29 {
30     int i, u;
31     for (i = 1; i <= n; ++i) S.insert((node){0, i});
32     while (all < n)
33     {
34         u = (*S.begin()).ord; S.erase(S.begin()); if (in[u]) continue;
35         in[u] = true; per[+all] = u;
36         for (i = side[u]; i; i = next[i])
37             if (!in[toit[i]])
38             {
39                 label[toit[i]]++;
40                 S.insert((node){label[toit[i]], toit[i]});
41             }
42     }
43 }
44
45 inline void paint()

```

```

46 {
47     int p, i, j, t;
48     for (p = 1; p <= n; ++p)
49     {
50         i = per[p];
51         for (j = 1; j <= tot; ++j)
52             if (!color[i][j]) {t = j; break; }
53         if (j == tot + 1) t = ++tot;
54         for (j = side[i]; j; j = next[j])
55             color[toit[j]][t] = true;
56     }
57 }
58
59 int main()
60 {
61     freopen("1006.in", "r", stdin);
62     freopen("1006.out", "w", stdout);
63     scanf("%d %d", &n, &m);
64     for (int i = 1; i <= m; ++i)
65     { int a, b; scanf("%d %d", &a, &b); ins(a, b); }
66     mcs();
67     paint();
68     printf("%d", tot);
69     fclose(stdin); fclose(stdout);
70     return 0;
71 }

```

## 强连通分量

```

1 int dfn[maxn], low[maxn], timestamp;
2 stack <int> stk; vector <int> scc[maxn];
3 void tarjan(int now)
4 {
5     dfn[now] = low[now] = ++timestamp;
6     stk.push(now);
7     for (int i = side[now]; i; i = nxt[i])
8     {
9         if (!dfn[toit[i]])
10             tarjan(toit[i]), low[now] = min(low[now], low[toit[i]]);
11         else if (!bel[toit[i]]) low[now] = min(low[now], dfn[toit[i]]);
12     }
13     if (dfn[now] == low[now])
14     {
15         ++tot;
16         while (stk.top() != now)
17         {
18             scc[tot].push_back(stk.top());
19             bel[stk.top()] = tot; stk.pop();
20         }
21         scc[tot].push_back(stk.top());
22         bel[stk.top()] = tot; stk.pop();
23     }
24 }

```

## 支配树

```

1 //建出来的树点的编号 i 在原图中是 redfn[i]
2 int N, M, Ts, cnt, side[maxn], nxt[maxn], toit[maxn], dfn[maxn], redfn[maxn], idom[maxn], best[maxn], semi[maxn];
3 int ans[maxn], anc[maxn], fa[maxn], child[maxn], size[maxn]; vector <int>
4     ↪ prod[maxn], bucket[maxn], son[maxn];
5
6 inline void init()
7 {
8     cnt = 1; memset(side, 0, sizeof side); memset(ans, 0, sizeof ans);
9     for (int i = 0; i <= N; ++i) prod[i].clear(), bucket[i].clear(), son[i].clear();
10 }
11
12 inline void add(int a, int b) {nxt[+cnt] = side[a]; side[a] = cnt; toit[cnt] = b; }
13
14 inline int gi()
15 {
16     char ch; int ret = 0, f = 1;
17     do ch = getchar(); while (!(ch >= '0' && ch <= '9') && ch != '-');
18     if (ch == '-') f = -1, ch = getchar();
19     do ret = ret * 10 + ch - '0', ch = getchar(); while (ch >= '0' && ch <= '9');
20     return ret * f;
21 }
22
23 inline void dfs(int now)
24 {
25     dfn[now] = ++Ts; redfn[Ts] = now;

```

```

25 anc[Ts] = idom[Ts] = child[Ts] = size[Ts] = 0;
26 semi[Ts] = best[Ts] = Ts;
27 for (int i = side[now]; i != nxt[i])
28 {
29     if (!dfn[toit[i]])
30         dfs(toit[i]), fa[dfn[toit[i]]] = dfn[now];
31     prod[dfn[toit[i]]].push_back(dfn[now]);
32 }
33 }
34
35 inline void compress(int now)
36 {
37     if (anc[anc[now]] != 0)
38     {
39         compress(anc[now]);
40         if (semi[best[now]] > semi[best[anc[now]]])
41             best[now] = best[anc[now]];
42         anc[now] = anc[anc[now]];
43     }
44 }
45
46 inline int eval(int now)
47 {
48     if (!anc[now]) return now;
49     else
50     {
51         compress(now);
52         return semi[best[anc[now]]] >= semi[best[now]] ? best[now] : best[anc[now]];
53     }
54 }
55
56 inline void link(int v, int w)
57 {
58     int s = w;
59     while (semi[best[w]] < semi[best[child[w]]])
60     {
61         if (size[s] + size[child[child[s]]] >= 2 * size[child[s]])
62             anc[child[s]] = s, child[s] = child[child[s]];
63         else size[child[s]] = size[s], s = anc[s] = child[s];
64     }
65     best[s] = best[w]; size[v] += size[w];
66     if (size[v] < 2 * size[w]) swap(s, child[v]);
67     while (s) anc[s] = v, s = child[s];
68 }
69
70 inline void lengauer_tarjan()
71 {
72     memset(dfn, 0, sizeof dfn); memset(fa, -1, sizeof fa); Ts = 0;
73     dfs(N); fa[1] = 0;
74     for (int w = Ts; w > 1; --w)
75     {
76         for (auto x: prod[w])
77         {
78             int u = eval(x);
79             if (semi[w] > semi[u]) semi[w] = semi[u];
80         }
81         bucket[semi[w]].push_back(w);
82         link(fa[w], w); if (!fa[w]) continue;
83         for (auto x: bucket[fa[w]])
84         {
85             int u = eval(x);
86             if (semi[u] < fa[w]) idom[x] = u;
87             else idom[x] = fa[w];
88         }
89         bucket[fa[w]].clear();
90     }
91     for (int w = 2; w <= Ts; ++w)
92         if (idom[w] != semi[w])
93             idom[w] = idom[idom[w]];
94     idom[1] = 0;
95     for (int i = Ts; i > 1; --i)
96     {
97         if (fa[i] == -1) continue;
98         son[idom[i]].push_back(i);
99     }
100 }

```

## 点双联通分量

```

1 const int maxn = 400010;
2 int N, M, Q, cnt = 1, side[maxn], toit[maxn], nxt[maxn], f[maxn][25], father[maxn], low[maxn];
3 int tot, dep[maxn], dfn[maxn], nside[maxn], ntoit[maxn], nnxt[maxn]; bool cut[maxn];

```

```

4 stack<int> S; vector<int> bel[maxn], bcc[maxn]; bool vis[maxn];
5
6 inline int find(int a) { if (father[a] != a) father[a] = find(father[a]); return father[a]; }
7 inline void add(int a, int b) { nxt[++cnt] = side[a]; side[a] = cnt; toit[cnt] = b; }
8 inline void ins(int a, int b) { add(a, b); add(b, a); }
9 inline void nadd(int a, int b) { nnxt[++cnt] = nside[a]; nside[a] = cnt; ntoit[cnt] = b; }
10 inline void nins(int a, int b) { nadd(a, b); nadd(b, a); }
11
12 inline void tj(int now, int fa)
13 {
14     dfn[now] = low[now] = ++cnt; int child = 0;
15     for (int i = side[now]; i != nxt[i])
16     {
17         if (toit[i] == fa) continue;
18         if (!dfn[toit[i]])
19         {
20             S.push(i >> 1); tj(toit[i], now); ++child;
21             low[now] = min(low[now], low[toit[i]]);
22             if (low[toit[i]] >= dfn[now])
23             {
24                 cut[now] = true; ++tot;
25                 while (true)
26                 {
27                     int t = S.top(); S.pop();
28                     bel[toit[t << 1]].push_back(tot); bel[toit[t << 1 | 1]].push_back(tot);
29                     bcc[tot].push_back(toit[t << 1]); bcc[tot].push_back(toit[t << 1 | 1]);
30                     if (t == (i >> 1)) break;
31                 }
32             }
33         }
34         else low[now] = min(low[now], dfn[toit[i]]);
35     }
36     if (!fa && child == 1) cut[now] = false;
37 }
38
39 inline void build()
40 {
41     vector<int> cuts; cnt = 1;
42     for (int i = 1; i <= tot; ++i)
43     {
44         sort(bcc[i].begin(), bcc[i].end());
45         bcc[i].erase(unique(bcc[i].begin(), bcc[i].end()), bcc[i].end());
46     }
47     for (int i = 1; i <= N; ++i) if (cut[i]) cuts.push_back(i);
48     for (auto x: cuts)
49     {
50         sort(bel[x].begin(), bel[x].end());
51         bel[x].erase(unique(bel[x].begin(), bel[x].end()), bel[x].end());
52         ++tot; for (auto y: bel[x]) nins(tot, y);
53         bel[x].clear(); bel[x].push_back(tot); bcc[tot].push_back(x);
54     }
55 }

```

## 线性规划

```

1 #include<iostream>
2 #include<cstdio>
3 #include<cstdlib>
4 using namespace std;
5
6 #define maxn (30)
7 #define eps (1e-8)
8
9 int N, M, op, tot, q[maxn], idx[maxn], idy[maxn]; double a[maxn][maxn], A[maxn];
10
11 inline void pivot(int x, int y)
12 {
13     swap(idy[x], idx[y]);
14     double tmp = a[x][y]; a[x][y] = 1/a[x][y];
15     for (int i = 0; i <= N; ++i) if (y != i) a[x][i] /= tmp;
16     tot = 0; for (int i = 0; i <= N; ++i) if (i != y && (a[x][i] > eps || a[x][i] < -eps)) q[++tot] = i;
17     for (int i = 0; i <= M; ++i)
18     {
19         if ((x == i) || (a[i][y] < eps && a[i][y] > -eps)) continue;
20         for (int j = 1; j <= tot; ++j) a[i][q[j]] -= a[x][q[j]] * a[i][y];
21         a[i][y] = -a[i][y] / tmp;
22     }
23 }
24
25 int main()
26 {

```

```

27 freopen("179.in", "r", stdin);
28 freopen("179.out", "w", stdout);
29 scanf("%d %d %d", &N, &M, &op); srand(233);
30 for (int i = 1; i <= N; ++i) scanf("%lf", a[0] + i);
31 for (int i = 1; i <= M; ++i)
32 {
33     for (int j = 1; j <= N; ++j) scanf("%lf", a[i] + j);
34     scanf("%lf", a[i]);
35 }
36 for (int i = 1; i <= N; ++i) idx[i] = i;
37 for (int i = 1; i <= M; ++i) idy[i] = i + N;
38 while (true)
39 {
40     int x = 0, y = 0;
41     for (int i = 1; i <= M; ++i) if (a[i][0] < -eps && ((x || (rand() < 1))) x = i; if (!x) break;
42     for (int i = 1; i <= N; ++i) if (a[x][i] < -eps && ((y || (rand() < 1))) y = i; if (!y) return
↪ puts("Infeasible"), 0;
43     pivot(x, y);
44 }
45 while (true)
46 {
47     int x = 0, y = 0; double mn = 1e15;
48     for (int i = 1; i <= N; ++i) if (a[0][i] > eps) { y = i; break; } if (!y) break;
49     for (int i = 1; i <= M; ++i) if (a[i][y] > eps && a[i][0] / a[i][y] < mn) mn = a[i][0] / a[i][y], x = i;
↪ if (!x) return puts("Unbounded"), 0;
50     pivot(x, y);
51 }
52 printf("%.8lf\n", -a[0][0]); if (!op) return 0;
53 for (int i = 1; i <= M; ++i) if (idy[i] <= N) A[idy[i]] = a[i][0];
54 for (int i = 1; i <= N; ++i) printf("%.8lf ", A[i]);
55 fclose(stdin); fclose(stdout);
56 return 0;
57 }

```

### 费用流

```

1 int side[maxv], nxt[maxe], tolt[maxe], cost[maxe], pre[maxv];
2 int cap[maxv], arr[maxv], dis[maxv]; bool in[maxv];
3 int source, sink;
4
5 inline void add(int a, int b, int c, int d) { nxt[++cnt] = side[a]; side[a] = cnt; tolt[cnt] = b; cap[cnt]
↪ = c; cost[cnt] = d; }
6 inline void ins(int a, int b, int c, int d) { add(a, b, c, d); add(b, a, 0, -d); }
7
8 inline bool spfa(int &Flow, int &Cost)
9 {
10     queue<int> team; team.push(source);
11     memset(dis, 0x7f, 4 * (sink + 5));
12     dis[source] = 0; in[source] = true;
13     arr[source] = inf; arr[sink] = 0;
14     while (!team.empty())
15     {
16         int now = team.front(); team.pop();
17         for (int i = side[now]; i; i = nxt[i])
18         {
19             if (!cap[i]) continue;
20             if (dis[tolt[i]] > dis[now] + cost[i])
21             {
22                 arr[tolt[i]] = min(cap[i], arr[now]); pre[tolt[i]] = i;
23                 dis[tolt[i]] = dis[now] + cost[i];
24                 if (!in[tolt[i]]) in[tolt[i]] = true, team.push(tolt[i]);
25             }
26         }
27         in[now] = false;
28     }
29     if (!arr[sink]) return false;
30     Flow += arr[sink];
31     for (int now = sink; now != source; now = tolt[i^1])
32     {
33         i = pre[now]; Cost += cost[pre[now]] * arr[sink];
34         cap[i] -= arr[sink]; cap[i^1] += arr[sink];
35     }
36     return true;
37 }

```

### 直线下整点个数

注意事项：返回结果为： $\sum_{0 \leq i < n} \lfloor \frac{a+b \cdot i}{m} \rfloor$  即直线下整点个数。

```

1 long long solve(const long long &n, const long long &a,
2               const long long &b, const long long &m) {
3     if (b == 0) return n * (a / m);
4     if (a >= m) return n * (a / m) + solve(n, a % m, b, m);
5     if (b >= m) return (n - 1) * n / 2 * (b / m) + solve(n, a, b % m, m);
6     return solve((a + b * n) / m, (a + b * n) % m, b);
7 }

```

### 闪电素数判定

```

1 const int BASE[12] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37};
2 bool check(const long long &prime, const long long &base) {
3     long long number = prime - 1;
4     for (; -number & 1; number >= 1);
5     long long result = power_mod(base, number, prime);
6     for (; number != prime - 1 && result != 1 && result != prime - 1; number <= 1) {
7         result = multiply_mod(result, result, prime);
8     }
9     return result == prime - 1 || (number & 1) == 1;
10 }
11 bool miller_rabin(const long long &number) {
12     if (number < 2) return false;
13     if (number < 4) return true;
14     if (-number & 1) return false;
15     for (int i = 0; i < 12 && BASE[i] < number; ++i) {
16         if (!check(number, BASE[i])) {
17             return false;
18         }
19     }
20     return true;
21 }

```

### 闪电质因数分解

```

1 long long pollard_rho(const long long &number, const long long &seed) {
2     long long x = rand() % (number - 1) + 1, y = x;
3     for (int head = 1, tail = 2; ; ) {
4         x = multiply_mod(x, x, number);
5         x = add_mod(x, seed, number);
6         if (x == y) return number;
7         long long answer = std::__gcd(abs(x - y), number);
8         if (answer > 1 && answer < number) return answer;
9         if (++head == tail) {
10             y = x;
11             tail <= 1;
12         }
13     }
14 }
15 void factorize(const long long &number, std::vector<long long> &divisor) {
16     if (number > 1) {
17         if (miller_rabin(number)) divisor.push_back(number);
18         else {
19             long long factor = number;
20             for (; factor >= number; factor = pollard_rho(number, rand() % (number - 1) + 1));
21             factorize(number / factor, divisor);
22             factorize(factor, divisor);
23         }
24     }
25 }

```

### 自适应辛普森

```

1 double area(const double &left, const double &right) {
2     double mid = (left + right) / 2;
3     return (right - left) * (calc(left) + 4 * calc(mid) + calc(right)) / 6;
4 }
5 double simpson(const double &left, const double &right,
6               const double &eps, const double &area_sum) {
7     double mid = (left + right) / 2;
8     double area_left = area(left, mid), area_right = area(mid, right);
9     double area_total = area_left + area_right;
10    if (std::abs(area_total - area_sum) < 15 * eps) {
11        return area_total + (area_total - area_sum) / 15;
12    }
13    return simpson(left, mid, eps / 2, area_left) + simpson(mid, right, eps / 2, area_right);
14 }
15 double simpson(const double &left, const double &right, const double &eps) {
16     return simpson(left, right, eps, area(left, right));
17 }

```

## 孙司宇 FFT

```

1  #include<iostream>
2  #include<cstdio>
3  #include<cmath>
4  using namespace std;
5  const double eps=1e-8;
6  const double PI=acos(-1.0);
7  struct Complex
8  {
9      double real,image;
10     Complex(double _real,double _image)
11     {
12         real=_real;
13         image=_image;
14     }
15     Complex(){real=0;image=0;}
16 };
17
18 Complex operator + (const Complex &c1, const Complex &c2)
19 {
20     return Complex(c1.real + c2.real, c1.image + c2.image);
21 }
22
23 Complex operator - (const Complex &c1, const Complex &c2)
24 {
25     return Complex(c1.real - c2.real, c1.image - c2.image);
26 }
27
28 Complex operator * (const Complex &c1, const Complex &c2)
29 {
30     return Complex(c1.real*c2.real - c1.image*c2.image, c1.real*c2.image + c1.image*c2.real);
31 }
32
33 int rev(int id,int len)
34 {
35     int ret=0;
36     for(int i=0;(1<<i)<len;i++)
37     {
38         ret<<=1;
39         if(id&(1<<i))
40             ret|=1;
41     }
42     return ret;
43 }
44
45 Complex* IterativeFFT(Complex* a,int len,int DFT)
46 {
47     Complex* A=new Complex[len];
48     for(int i=0;i<len;i++)
49         A[rev(i,len)]=a[i];
50     for(int s=1;(1<<s)<=len;s++)
51     {
52         int m=(1<<s);
53         Complex wm=Complex(cos(DFT*2*PI/m),sin(DFT*2*PI/m));
54         for(int k=0;k<len;k+=m)
55         {
56             Complex w=Complex(1,0);
57             for(int j=0;j<(m>>1);j++)
58             {
59                 Complex t=w*A[k+j+(m>>1)];
60                 Complex u=A[k+j];
61                 A[k+j]=u+t;
62                 A[k+j+(m>>1)]=u-t;
63                 w=w*wm;
64             }
65         }
66         if(DFT==-1)
67             for(int i=0;i<len;i++)
68             {
69                 A[i].real/=len;
70                 A[i].image/=len;
71             }
72         return A;
73     }
74     char s[101010],t[101010];
75     Complex a[202020],b[202020],c[202020];
76     int pr[202020];
77     int main()

```

```

78 {
79     int len;
80     scanf("%d",&len);
81     scanf("%s",s);
82     scanf("%s",t);
83     for(int i=0;i<len;i++)
84         a[i]=Complex(s[len-i-1]-'0',0);
85     for(int i=0;i<len;i++)
86         b[i]=Complex(t[len-i-1]-'0',0);
87     int tmp=1;
88     while(tmp<=len)
89     {
90         tmp*=2;
91         len=tmp*2;
92         Complex* aa=IterativeFFT(a,len,1);
93         Complex* bb=IterativeFFT(b,len,1);
94         for(int i=0;i<len;i++)
95             c[i]=aa[i]*bb[i];
96         Complex* ans=IterativeFFT(c,len,-1);
97         for(int i=0;i<len;i++)
98             pr[i]=round(ans[i].real);
99         for(int i=0;i<len;i++)
100         {
101             pr[i+1]+=pr[i]/10;
102             pr[i]%=10;
103         }
104         bool flag=0;
105         for(int i=len-1;i>=0;i--)
106         {
107             if(pr[i]>0)
108                 flag=1;
109             if(flag)
110                 printf("%d",pr[i]);
111         }
112         printf("\n");
113         return 0;
114     }
115 }

```

## NTT

```

1  #include <iostream>
2  #include <cstdio>
3  #include <cstring>
4  #include <algorithm>
5  #include <cmath>
6  using namespace std;
7  const int N=(1<<18)+5, INF=1e9;
8  const double PI=acos(-1);
9  long long P=1004535809;
10 long long Pow(long long a, long long b,long long P)
11 {
12     long long ans=1;
13     for(; b; b>>=1, a=a*a%P)
14         if(b&1) ans=ans*a%P;
15     return ans;
16 }
17
18 struct NumberTheoreticTransform {
19     int n, rev[N];
20     long long g;
21     void ini(int lim) {
22         g=3;
23         n=1; int k=0;
24         while(n<lim) n<<=1, k++;
25         for(int i=0; i<n; i++) rev[i] = (rev[i>>1]>>1) | ((i&1)<<(k-1));
26     }
27     void dft(long long *a, int flag) {
28         for(int i=0; i<n; i++) if(i<rev[i]) swap(a[i], a[rev[i]]);
29         for(int l=2; l<=n; l<<=1) {
30             int m=l>>1;
31             long long wn = Pow(g, flag==1 ? (P-1)/l : P-1-(P-1)/l, P);
32             for(long long *p=a; p!=a+n; p+=l) {
33                 long long w=1;
34                 for(int k=0; k<m; k++) {
35                     long long t = w * p[k+m]%P;
36                     p[k+m]=(p[k]-t+P)%P;
37                     p[k]=(p[k]+t)%P;
38                     w=w*wn%P;
39                 }
40             }
41         }
42         if(flag==1) {
43             long long inv=Pow(n, P-2, P);
44             for(int i=0; i<n; i++) a[i]=a[i]*inv%P;
45         }
46     }
47 }

```



```

44     }
45 }
46 void mul(long long *a, long long *b, int m) {
47     ini(m);
48     dft(a, 1); dft(b, 1);
49     for(int i=0; i<n; i++) a[i]=a[i]*b[i];
50     dft(a, -1);
51 }
52 }f;
53
54 int n1, n2, m, c[N];
55 long long a[N], b[N];
56 char s1[N], s2[N];
57 int main()
58 {
59     int n;
60     scanf("%d",&n);
61     scanf("%s%s",s1,s2);
62     n1=strlen(s1); n2=strlen(s2);
63     for(int i=0;i<n1;i++)
64         a[i]=s1[n1-i-1]-'0';
65     for(int i=0;i<n2;i++)
66         b[i]=s2[n2-i-1]-'0';
67     m=n1+n2-1;
68     f.mul(a,b,m);
69     for(int i=0;i<m;i++) c[i]=a[i];
70     for(int i=0;i<m;i++) c[i+1]+=c[i]/10, c[i]%=10;
71     if(c[m])
72         m++;
73     for(int i=m-1; i>=0; i--)
74         printf("%d",c[i]);
75     return 0;
76 }

```

## SAM

```

1  #include<iostream>
2  #include<cstring>
3  using namespace std;
4  const int MaxPoint=1010101;
5  struct Suffix_AutoMachine{
6      int son[MaxPoint][27],pre[MaxPoint],step[MaxPoint],right[MaxPoint],last,root,num;
7      int NewNode(int stp)
8      {
9          num++;
10         memset(son[num],0,sizeof(son[num]));
11         pre[num]=0;
12         step[num]=stp;
13         return num;
14     }
15     Suffix_AutoMachine()
16     {
17         num=0;
18         root=last=NewNode(0);
19     }
20     void push_back(int ch)
21     {
22         int np=NewNode(step[last]+1);
23         right[np]=1;
24         step[np]=step[last]+1;
25         int p=last;
26         while(p&&!son[p][ch])
27         {
28             son[p][ch]=np;
29             p=pre[p];
30         }
31         if(!p)
32             pre[np]=root;
33         else
34         {
35             int q=son[p][ch];
36             if(step[q]==step[p]+1)
37                 pre[np]=q;
38             else
39             {
40                 int nq=NewNode(step[p]+1);
41                 memcpy(son[nq],son[q],sizeof(son[q]));
42                 step[nq]=step[p]+1;
43                 pre[nq]=pre[p];
44                 pre[q]=pre[np]=nq;
45                 while(p&&son[p][ch]==q)
46                     {

```

```

47             son[p][ch]=nq;
48             p=pre[p];
49         }
50     }
51     last=np;
52 }
53 };
54 /*
55
56 int arr[1010101];
57 bool Step_Cmp(int x,int y)
58 {
59     return S.step[x]<S.step[y];
60 }
61 void Get_Right()
62 {
63     for(int i=1;i<=S.num;i++)
64         arr[i]=i;
65     sort(arr+1,arr+S.num+1,Step_Cmp);
66     for(int i=S.num;i>=2;i--)
67         S.right[S.pre[arr[i]]]+=S.right[arr[i]];
68 }
69 */
70 int main()
71 {
72 }
73
74 return 0;
75 }

```

## manacher

```

1  #include<iostream>
2  #include<cstring>
3  using namespace std;
4  char Mana[202020];
5  int cher[202020];
6  int Manacher(char *S)
7  {
8      int len=strlen(S),id=0,mx=0,ret=0;
9      Mana[0]='$';
10     Mana[1]='#';
11     for(int i=0;i<len;i++)
12     {
13         Mana[2*i+2]=S[i];
14         Mana[2*i+3]='#';
15     }
16     Mana[2*len+2]=0;
17     for(int i=1;i<=2*len+1;i++)
18     {
19         if(i<mx)
20             cher[i]=min(cher[2*id-i],mx-i);
21         else
22             cher[i]=0;
23         while(Mana[i+cher[i]+1]==Mana[i-cher[i]-1])
24             cher[i]++;
25         if(cher[i]>mx)
26         {
27             mx=cher[i]+i;
28             id=i;
29         }
30         ret=max(ret,cher[i]);
31     }
32     return ret;
33 }
34 char S[101010];
35 int main()
36 {
37     ios::sync_with_stdio(false);
38     cin.tie(0);
39     cout.tie(0);
40     cin>>S;
41     cout<<Manacher(S)<<endl;
42     return 0;
43 }

```

## 中国剩余定理

```

1  // 51nod 1079
2  #include<iostream>
3  using namespace std;
4  int gcd(int x,int y)

```

```

5  {
6      if(x==0)
7          return y;
8      if(y==0)
9          return x;
10     return gcd(y,x/y);
11 }
12 long long exgcd(long long a,long long b,long long &x,long long &y)
13 {
14     if(b==0)
15     {
16         x=1;
17         y=0;
18         return a;
19     }
20     long long ans=exgcd(b,a%b,x,y);
21     long long temp=x;
22     x=y;
23     y=temp-a/b*y;
24     return ans;
25 }
26 void fix(long long &x,long long &y)
27 {
28     x%=y;
29     if(x<0)
30         x+=y;
31 }
32 bool solve(int n, std::pair<long long, long long> input[],std::pair<long long, long long> &output)
33 {
34     output = std::make_pair(1, 1);
35     for(int i = 0; i < n; ++i)
36     {
37         long long number, useless;
38         exgcd(output.second, input[i].second, number, useless);
39         long long divisor = gcd(output.second, input[i].second);
40         if((input[i].first - output.first) % divisor)
41         {
42             return false;
43         }
44         number *= (input[i].first - output.first) / divisor;
45         fix(number, input[i].second);
46         output.first += output.second * number;
47         output.second *= input[i].second / divisor;
48         fix(output.first, output.second);
49     }
50     return true;
51 }
52 pair<long long,long long> input[101010],output;
53 int main()
54 {
55     int n;
56     cin>>n;
57     for(int i=0;i<n;i++)
58         cin>>input[i].second>>input[i].first;
59     solve(n,input,output);
60     cout<<output.first<<endl;
61     return 0;
62 }

```

## 回文自动机

```

1  //Tsinsen A1280 最长双回文串
2  #include<iostream>
3  #include<cstring>
4  using namespace std;
5
6  const int maxn = 100005;// n(空间复杂度  $O(n*ALP)$ ), 实际开  $n$  即可
7  const int ALP = 26;
8
9  struct PAM{ // 每个节点代表一个回文串
10     int next[maxn][ALP]; // next 指针, 参照 Trie 树
11     int fail[maxn]; // fail 失配后缀链接
12     int cnt[maxn]; // 此回文串出现个数
13     int num[maxn];
14     int len[maxn]; // 回文串长度
15     int s[maxn]; // 存放添加的字符
16     int last; //指向上一个字符所在的节点, 方便下一次 add
17     int n; // 已添加字符个数
18     int p; // 节点个数
19
20     int newnode(int w)

```

```

21     { // 初始化节点, w= 长度
22         for(int i=0;i<ALP;i++)
23             next[p][i] = 0;
24         cnt[p] = 0;
25         num[p] = 0;
26         len[p] = w;
27         return p++;
28     }
29     void init()
30     {
31         p = 0;
32         newnode(0);
33         newnode(-1);
34         last = 0;
35         n = 0;
36         s[n] = -1; // 开头放一个字符集中没有的字符, 减少特判
37         fail[0] = 1;
38     }
39     int get_fail(int x)
40     { // 和 KMP 一样, 失配后找一个尽量最长的
41         while(s[n-len[x]-1] != s[n]) x = fail[x];
42         return x;
43     }
44     int add(int c)
45     {
46         c -= 'a';
47         s[++n] = c;
48         int cur = get_fail(last);
49         if(!next[cur][c])
50         {
51             int now = newnode(len[cur]+2);
52             fail[now] = next[get_fail(fail[cur])][c];
53             next[cur][c] = now;
54             num[now] = num[fail[now]] + 1;
55         }
56         last = next[cur][c];
57         cnt[last]++;
58         return len[last];
59     }
60     void count()
61     {
62         // 最后统计一遍每个节点出现个数
63         // 父亲累加儿子的 cnt, 类似 SAM 中 parent 树
64         // 满足 parent 拓扑关系
65         for(int i=p-1;i>0;i--)
66             cnt[fail[i]] += cnt[i];
67     }
68 }pam;
69 char S[101010];
70 int l[101010],r[101010];
71 int main()
72 {
73     cin>>S;
74     int len=strlen(S);
75     pam.init();
76     for(int i=0;i<len;i++)
77         l[i]=pam.add(S[i]);
78     pam.init();
79     for(int i=len-1;i>=0;i--)
80         r[i]=pam.add(S[i]);
81     pam.init();
82     int ans=0;
83     for(int i=0;i<len-1;i++)
84         ans=max(ans,l[i]+r[i+1]);
85     cout<<ans<<endl;
86     return 0;
87 }

```

## 多项式开方

```

1  //  $\sqrt{a}$ 
2  //  $N \log^2 N$ 
3  #include <cstdio>
4  #include <algorithm>
5  #define FOR(i,j,k) for(i=j;i<=k;++i)
6  #define rep(i,j,k) for(i=j;i<k;++i)
7  #define gmod(i) ((i)%mod+mod)%mod
8  const int N = 262144, mod = 998244353, inv2 = 499122177;
9  using namespace std;
10 typedef long long ll;
11 ll qpow(ll x, int y) {

```

```

12 ll z = 1;
13 for (; y; x = x * x % mod, y /= 2)
14     if (y & 1) z = z * x % mod;
15 return z;
16 }
17 namespace NTT {
18     int n, rev[N], inv_n, m = -1;
19     void init(int c) {
20         int k = -1, i;
21         if (m == c) return; else m = c;
22         for (n = 1; n <= m; n <= 1) ++k;
23         inv_n = qpow(n, mod - 2);
24         rep(i, 0, n) rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << k);
25     }
26     void ntt(int *a, int f) {
27         int h, i, j;
28         rep(i, 0, n) if (i < rev[i]) swap(a[i], a[rev[i]]);
29         for (h = 2; h <= n; h *= 2) {
30             int wn = qpow(3, (mod - 1) / h);
31             for (i = 0; i < n; i += h) {
32                 int w = 1;
33                 rep(j, 0, h/2) {
34                     int u = a[i + j], t = 1ll * a[i + j + h / 2] * w % mod;
35                     a[i + j + h / 2] = (u - t + mod) % mod;
36                     a[i + j] = (u + t) % mod;
37                     w = 1ll * w * wn % mod;
38                 }
39             }
40         }
41         if (f) {
42             rep(i, 1, n/2) swap(a[i], a[n - i]);
43             rep(i, 0, n) a[i] = 1ll * a[i] * inv_n % mod;
44         }
45     }
46     void inv(int *a, int *b, int n) {
47         static int t[N];
48         int i;
49         if (n == 1) { b[0] = qpow(a[0], mod - 2); return; }
50         inv(a, b, n / 2);
51         rep(i, 0, n) t[i] = a[i]; rep(i, n, 2*n) t[i] = 0;
52         NTT::init(n);
53         NTT::ntt(t, 0); NTT::ntt(b, 0);
54         rep(i, 0, NTT::n) t[i] = (1ll) b[i] * gmod(2ll - (1ll) t[i] * b[i] % mod) % mod;
55         NTT::ntt(t, 1);
56         rep(i, 0, n) b[i] = t[i]; rep(i, n, 2*n) b[i] = 0;
57     }
58     void sqrt(int *a, int *b, int n) {
59         static int t[N], b1[N];
60         if (n == 1) { b[0] = 1; return; }
61         int i;
62         sqrt(a, b, n / 2);
63         rep(i, 0, n) b1[i] = 0;
64         inv(b, b1, n);
65         rep(i, 0, n) t[i] = a[i]; rep(i, n, 2*n) t[i] = 0;
66         NTT::init(n);
67         NTT::ntt(t, 0); NTT::ntt(b, 0); NTT::ntt(b1, 0);
68         rep(i, 0, NTT::n) t[i] = inv2 * ((b[i] + (1ll) b1[i] * t[i] % mod) % mod) % mod;
69         NTT::ntt(t, 1);
70         rep(i, 0, n) b[i] = t[i]; rep(i, n, 2*n) b[i] = 0;
71     }
72 }
73 int main() {
74     static int c[N], sc[N], ic[N];
75     int i, x, n, m, l;
76     scanf("%d%d", &n, &m);
77     FOR(i, 1, n) scanf("%d", &x), ++c[x];
78     c[0] = gmod(1 - c[0]);
79     FOR(i, 1, m) c[i] = gmod(-4 * c[i]);
80     for (l = 1; l <= m; l <= 1);
81     sqrt(c, sc, l);
82     (++sc[0]) %= mod;
83     inv(sc, ic, l);
84     FOR(i, 0, m) ic[i] = 2ll * ic[i] % mod;
85     FOR(i, 1, m) printf("%d\n", ic[i]);
86     return 0;
87 }

```

## 多项式求逆

```

1 //³  bzoj3456
2 #include<iostream>
3 #include<cstdio>

```

```

4 #include<algorithm>
5 #include<cstring>
6 #include<cmath>
7 #define N 5000003
8 #define LL long long
9 #define p 1004535809
10 using namespace std;
11 int n, m;
12 int a[N], b[N], c[N], jc[N], inv_j[N], wn[N];
13 LL quickpow(LL num, LL x)
14 {
15     LL base=num%p; LL ans=1;
16     while (x) {
17         if (x&1) ans=ans*base%p;
18         x>>=1;
19         base=base*base%p;
20     }
21     return ans;
22 }
23 void init()
24 {
25     jc[0]=1; inv_j[0]=quickpow(jc[0], p-2);
26     for (int i=1; i<=n; i++)
27         jc[i]=(LL)jc[i-1]*i%p, inv_j[i]=quickpow(jc[i], p-2);
28     for (int i=1; i<=n*8; i<=1)
29         wn[i]=quickpow(3, (p-1)/(i<<1));
30 }
31 void NTT(int n, int *a, int opt)
32 {
33     for (int i=0, j=0; i<n; i++) {
34         if (i>j) swap(a[i], a[j]);
35         for (int l=n>>1; (j^=l)<l; l>>=1);
36     }
37     for (int i=1; i<n; i<=1) {
38         LL wn1=wn[i];
39         for (int p1=i<<1, j=0; j<n; j+=p1) {
40             LL w=1;
41             for (int k=0; k<i; k++, w=(LL)w*wn1%p) {
42                 int x=a[j+k]; int y=(LL)a[j+k+i]*w%p;
43                 a[j+k]=(x+y)%p; a[j+k+i]=(x-y+p)%p;
44             }
45         }
46     }
47     if (opt==-1) reverse(a+1, a+n);
48 }
49 void inverse(int n, int *a, int *b, int *c)
50 {
51     if (n==1) b[0]=quickpow(a[0], p-2);
52     else {
53         inverse((n+1)>>1, a, b, c);
54         int k=0;
55         for (k=1; k<=(n<<1); k<=1);
56         for (int i=0; i<n; i++) c[i]=a[i];
57         for (int i=n; i<k; i++) c[i]=0;
58         NTT(k, c, 1);
59         NTT(k, b, 1);
60         for (int i=0; i<k; i++) {
61             b[i]=(LL)(2-(LL)c[i]*b[i]%p)*b[i]%p;
62             if (b[i]<0) b[i]+=p;
63         }
64         NTT(k, b, -1);
65         int inv=quickpow(k, p-2);
66         for (int i=0; i<k; i++) b[i]=(LL)b[i]*inv%p;
67         for (int i=n; i<k; i++) b[i]=0;
68     }
69 }
70 int main()
71 {
72     scanf("%d", &n); init();
73     int n1=0;
74     for (n1=1; n1<=n*2; n1<=1);
75     a[0]=1;
76     for (int i=1; i<=n; i++) a[i]=(LL)quickpow(2, (LL)i*(i-1)/2)*inv_j[i]%p;
77     inverse(n1, a, b, c);
78     memset(c, 0, sizeof(c));
79     for (int i=1; i<=n; i++) c[i]=(LL)quickpow(2, (LL)i*(i-1)/2)*inv_j[i-1]%p;
80     NTT(n1, b, 1); NTT(n1, c, 1);
81     for (int i=0; i<=n1; i++) b[i]=(LL)b[i]*c[i]%p;
82     NTT(n1, b, -1);
83     LL inv=quickpow(n1, p-2);
84     for (int i=0; i<=n1; i++) b[i]=(LL)b[i]*inv%p;

```

```

85     printf("%d\n", (LL)b[n]*jc[n-1]%p);
86     return 0;
87 }

```

## 广义 SAM

```

1  #include<iostream>
2  #include<cstring>
3  using namespace std;
4  const int MaxPoint=1010101;
5  struct Suffix_AutoMachine{
6      int son[MaxPoint][27], pre[MaxPoint], step[MaxPoint], right[MaxPoint], root, num;
7      int NewNode(int stp)
8      {
9          num++;
10         memset(son[num], 0, sizeof(son[num]));
11         pre[num]=0;
12         step[num]=stp;
13         return num;
14     }
15     Suffix_AutoMachine()
16     {
17         num=0;
18         root=NewNode(0);
19     }
20     int push_back(int ch, int p)
21     {
22         int np=NewNode(step[p]+1);
23         right[np]=1;
24         step[np]=step[p]+1;
25         while(p&&!son[p][ch])
26         {
27             son[p][ch]=np;
28             p=pre[p];
29         }
30         if(!p)
31             pre[np]=root;
32         else
33         {
34             int q=son[p][ch];
35             if(step[q]==step[p]+1)
36                 pre[np]=q;
37             else
38             {
39                 int nq=NewNode(step[p]+1);
40                 memcpy(son[nq], son[q], sizeof(son[q]));
41                 step[nq]=step[p]+1;
42                 pre[nq]=pre[q];
43                 pre[q]=pre[np]=nq;
44                 while(p&&son[p][ch]==q)
45                 {
46                     son[p][ch]=nq;
47                     p=pre[p];
48                 }
49             }
50         }
51         return np;
52     }
53 };
54 int main()
55 {
56     return 0;
57 }
58 }

```

## 循环串最小表示

```

1  int getmin( char s[] )
2  {
3      int i, j, k, m, t;
4      m = strlen( s );
5      i = 0; j = 1; k = 0;
6      while( i < m && j < m && k < m )
7      {
8          t = s[( i + k ) % m] - s[( j + k ) % m];
9          if( !t )
10             ++k;
11         else
12         {
13             if( t > 0 )
14                 i += k + 1;
15             else

```

```

16             j += k + 1;
17             if( i == j )
18                 j++;
19             k = 0;
20         }
21     }
22     return min(i, j);
23 }
24 }

```

## 最大团搜索

```

1  #include<iostream>
2  using namespace std;
3  int ans;
4  int num[1010];
5  int path[1010];
6  int a[1010][1010], n;
7  bool dfs(int *adj, int total, int cnt)
8  {
9      int i, j, k;
10     int t[1010];
11     if(total==0)
12     {
13         if(ans<cnt)
14         {
15             ans=cnt;
16             return 1;
17         }
18         return 0;
19     }
20     for(i=0; i<total; i++)
21     {
22         if(cnt+(total-i)<=ans)
23             return 0;
24         if(cnt+num[adj[i]]<=ans)
25             return 0;
26         for(k=0, j=i+1; j<total; j++)
27             if(a[adj[i]][adj[j]])
28                 t[k++]=adj[j];
29         if(dfs(t, k, cnt+1))
30             return 1;
31     }
32     return 0;
33 }
34 int MaxClique()
35 {
36     int i, j, k;
37     int adj[1010];
38     if(n<=0)
39         return 0;
40     ans=1;
41     for(i=n-1; i>=0; i--)
42     {
43         for(k=0, j=i+1; j<n; j++)
44             if(a[i][j])
45                 adj[k++]=j;
46         dfs(adj, k, 1);
47         num[i]=ans;
48     }
49     return ans;
50 }
51 int main()
52 {
53     ios::sync_with_stdio(0);
54     cin.tie(0);
55     cout.tie(0);
56     while(cin>>n)
57     {
58         if(n==0)
59             break;
60         for(int i=0; i<n; i++)
61             for(int j=0; j<n; j++)
62                 cin>>a[i][j];
63         cout<<MaxClique()<<endl;
64     }
65     return 0;
66 }

```

## 求原根

```

1  //51Nod - 1135
2  #include <iostream>

```

```

3  #include <string.h>
4  #include <algorithm>
5  #include <stdio.h>
6  #include <math.h>
7  #include <bitset>
8
9  using namespace std;
10 typedef long long LL;
11
12 const int N = 1000010;
13
14 bitset<N> prime;
15 int p[N], pri[N];
16 int k, cnt;
17
18 void isprime()
19 {
20     prime.set();
21     for(int i=2; i<N; i++)
22     {
23         if(prime[i])
24         {
25             p[k++] = i;
26             for(int j=i+i; j<N; j+=i)
27                 prime[j] = false;
28         }
29     }
30 }
31
32 void Divide(int n)
33 {
34     cnt = 0;
35     int t = (int)sqrt(1.0*n);
36     for(int i=0; p[i]<=t; i++)
37     {
38         if(n%p[i]==0)
39         {
40             pri[cnt++] = p[i];
41             while(n%p[i]==0) n /= p[i];
42         }
43     }
44     if(n > 1)
45         pri[cnt++] = n;
46 }
47
48 LL quick_mod(LL a, LL b, LL m)
49 {
50     LL ans = 1;
51     a %= m;
52     while(b)
53     {
54         if(b&1)
55         {
56             ans = ans * a % m;
57             b--;
58         }
59         b >>= 1;
60         a = a * a % m;
61     }
62     return ans;
63 }
64
65 int main()
66 {
67     int P;
68     isprime();
69     while(cin >> P)
70     {
71         Divide(P-1);
72         for(int g=2; g<P; g++)
73         {
74             bool flag = true;
75             for(int i=0; i<cnt; i++)
76             {
77                 int t = (P - 1) / pri[i];
78                 if(quick_mod(g, t, P) == 1)
79                 {
80                     flag = false;
81                     break;
82                 }
83             }
84             if(flag)

```

```

85         {
86             int root = g;
87             cout << root << endl;
88             break;
89         }
90     }
91     return 0;
92 }

```

## 线性递推多项式

```

1  //³µ
2  void linear_recurrence(long long n, int m, int a[], int c[], int p)
3  {
4      long long v[M] = {1 % p}, u[M << 1], msk = !n;
5      for(long long i(n); i > 1; i >= 1)
6      {
7          msk <<= 1;
8      }
9      for(long long x(0); msk; msk >>= 1, x <<= 1)
10     {
11         fill_n(u, m << 1, 0);
12         int b(!(n & msk));
13         x |= b;
14         if(x < m)
15         {
16             u[x] = 1 % p;
17         }
18         else
19         {
20             for(int i(0); i < m; i++)
21             {
22                 for(int j(0), t(i + b); j < m; j++, t++)
23                 {
24                     u[t] = (u[t] + v[i] * v[j]) % p;
25                 }
26             }
27             for(int i((m << 1) - 1); i >= m; i--)
28             {
29                 for(int j(0), t(i - m); j < m; j++, t++)
30                 {
31                     u[t] = (u[t] + c[j] * u[i]) % p;
32                 }
33             }
34         }
35         copy(u, u + m, v);
36     }
37     for(int i(m); i < 2 * m; i++)
38     {
39         a[i] = 0;
40         for(int j(0); j < m; j++)
41         {
42             a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
43         }
44     }
45     for(int j(0); j < m; j++)
46     {
47         b[j] = 0;
48         for(int i(0); i < m; i++)
49         {
50             b[j] = (b[j] + v[i] * a[i + j]) % p;
51         }
52     }
53     for(int j(0); j < m; j++)
54     {
55         a[j] = b[j];
56     }
57 }

```

## 经纬度球面距离

```

1  double sphereDis(double lon1, double lat1, double lon2, double lat2, double R) {
2      return R * acos(cos(lat1) * cos(lat2) * cos(lon1 - lon2) + sin(lat1) * sin(lat2));
3  }

```

## 日期公式

```

1  int zeller(int y, int m, int d) { // y 年 m 月 d 日是星期几
2      if (m <= 2) y--, m += 12; int c = y / 100; y %= 100;
3      int w = ((c >> 2) - (c << 1) + y + (y >> 2) + (13 * (m + 1) / 5) + d - 1) % 7;

```



```

4     if (w < 0) w += 7; return w;
5 }
6 int getId(int y, int m, int d) { // y 年 m 月 d 日的日期编号
7     if (m < 3) {y--; m += 12;}
8     return 365 * y + y / 4 - y / 100 + y / 400 + (153 * m + 2) / 5 + d;
9 }

```

## Manacher

注意事项：1-based 算法，请注意下标。

```

1 int manacher(char *text, int length, int *palindrome) {
2     static char buffer[MAXN];
3     for (int i = 1; i <= length; i++) {
4         buffer[2 * i - 1] = text[i];
5         if (i != 0) buffer[2 * i] = '#';
6     }
7     palindrome[1] = 1;
8     for (int i = 2, j = 0; i <= 2 * length - 1; ++i) {
9         if (j + palindrome[j] <= i) palindrome[i] = 0;
10        else palindrome[i] = std::min(palindrome[(j << 1) - i], j + palindrome[j] - i);
11        while (i - palindrome[i] >= 1 && i + palindrome[i] <= 2 * length - 1 && buffer[i - palindrome[i]]
↵ == buffer[i + palindrome[i]]) {
12            palindrome[i]++;
13        }
14        if (i + palindrome[i] > j + palindrome[j]) j = i;
15    }
16    int answer = 0;
17    for (int i = 1; i < 2 * length; i++) {
18        if (i & 1) answer = std::max(answer, 2 * (palindrome[i] - 1 >> 1) + 1);
19        else answer = std::max(answer, 2 * (palindrome[i] >> 1));
20    }
21    return answer;
22 }

```

## 丁尧尧

### kth.shortest.path

```

1 #include <stdio>
2 #include <cstring>
3 #include <queue>
4 using namespace std;
5
6 const int N = 1010;
7 const int M = 100010;
8 const int oo = 0x3f3f3f3f;
9
10 struct Elist {
11     int _head[N], _dest[M], _dist[M], _last[M], etot;
12     inline void add( int u, int v, int d ) {
13         etot++;
14         _dest[etot] = v;
15         _dist[etot] = d;
16         _last[etot] = _head[u];
17         _head[u] = etot;
18     }
19     inline int head( int u ) { return _head[u]; }
20     inline int dest( int t ) { return _dest[t]; }
21     inline int dist( int t ) { return _dist[t]; }
22     inline int last( int t ) { return _last[t]; }
23 };
24 struct Stat {
25     int u, d;
26     Stat(){}
27     Stat( int u, int d ):u(u),d(d){}
28 };
29
30 int n, m, K;
31 Elist e, re;
32 int src, dst;
33 int rdis[N];
34 bool done[N];
35
36 bool operator<( const Stat &r, const Stat &s ) {
37     return r.d + rdis[r.u] > s.d + rdis[s.u];
38 }
39 void dijkstra() {
40     memset( done, false, sizeof(done) );
41     memset( rdis, 0x3f, sizeof(rdis) );
42     priority_queue<pair<int,int>> > Q;
43     Q.push( make_pair(0,dst) );
44     rdis[dst] = 0;

```

```

45 while( !Q.empty() ) {
46     int u = Q.top().second;
47     Q.pop();
48     if( done[u] ) continue;
49     done[u] = true;
50     for( int t = re.head(u); t; t = re.last(t) ) {
51         int v = re.dest(t), d = re.dist(t);
52         if( done[v] ) continue;
53         if( rdis[v] > rdis[u] + d ) {
54             rdis[v] = rdis[u] + d;
55             Q.push( make_pair( -rdis[v], v ) );
56         }
57     }
58 }
59
60 int astar() {
61     int pcnt = 0;
62     priority_queue<Stat> Q;
63
64     if( rdis[src] == oo ) return -1;
65     if( src == dst ) K++;
66     Q.push( Stat( src, 0 ) );
67     while( !Q.empty() ) {
68         Stat s = Q.top();
69         Q.pop();
70         if( s.u == dst ) {
71             pcnt++;
72             if( pcnt == K )
73                 return s.d;
74         }
75         for( int t = e.head(s.u); t; t = e.last(t) ) {
76             int v = e.dest(t), d = e.dist(t);
77             if( rdis[v] == oo ) continue;
78             Q.push( Stat( v, s.d + d ) );
79         }
80     }
81     return -1;
82 }
83
84 int main() {
85     scanf( "%d%d", &n, &m );
86     for( int i = 1; i <= m; i++ ) {
87         int u, v, d;
88         scanf( "%d%d%d", &u, &v, &d );
89         e.adde( u, v, d );
90         re.adde( v, u, d );
91     }
92     scanf( "%d%d%d", &src, &dst, &K );
93     dijkstra();
94     printf( "%d\n", astar() );
95 }

```

## 弦图

```

1 /*
2 一些定义：
3 弦图是一种特殊图：它的所有极小环都只有 3 个顶点。
4 单纯点：该顶点与其邻接点在原图中的导出子图是一个完全图。
5 图 G 的完美消去序列：一个顶点序列  $a_1 a_2 a_3 \dots a_n$ ，使得对于每个元素  $a_i$ ， $a_i$  在  $a_i, a_{i+1}, a_{i+2} \dots a_n$  的导出子图中是一个单纯点。
6 弦图有一个性质：任何一个弦图都至少存在一个单纯点（该点及其邻接点组成一个完全图）
7 弦图另一个性质：一个图是弦图当且仅当其存在完美消去序列。（归纳证明）
8 最大势算法（msc）：若原图是弦图，则该算法计算出的序列是完美消去序列。
9 算法大致思想：从后往前计算序列，每次选择点  $v$  作为序列中的元素， $v$  是还未选的点中与已经选了的点连边最多的点。
10 然后检查该序列是否是完美消去序列。
11 */
12 #include <stdio>
13 #include <cstring>
14 #define N 1010
15 #define M N*N*2
16 int n, m;
17 bool c[N][N];
18 int qu[N], inq[N], dgr[N];
19 int stk[N], top;
20 void msc() {
21     dgr[0] = -1;
22     for( int i=n; i>=1; i-- ) {
23         int s = 0;
24         for( int u=1; u<=n; u++ )
25             if( !inq[u] && dgr[u]>dgr[s] ) s=u;
26         qu[i] = s;
27         inq[s] = true;

```

```

28     for( int u=1; u<=n; u++ )
29         if( !inq[u] && c[s][u] ) dgr[u]++;
30     }
31 }
32 bool check() {
33     for( int i=n; i>=1; i-- ) {
34         int s=qu[i];
35         top = 0;
36         for( int j=i+1; j<=n; j++ )
37             if( c[s][qu[j]] ) stk[++top] = qu[j];
38         if( top==0 ) continue;
39         for( int j=2; j<=top; j++ )
40             if( !c[stk[i]][stk[j]] ) return false;
41     }
42     return true;
43 }
44 int main() {
45     scanf( "%d%d", &n, &m );
46     for( int i=1,u,v; i<=m; i++ ) {
47         scanf( "%d%d", &u, &v );
48         c[u][v] = c[v][u] = 1;
49     }
50     msc();
51     printf( "%s\n", check() ? "Perfect" : "Imperfect" );
52 }
53
54 /*
55 给定一个弦图，问最少染色数。
56 对于弦图的一个完美消去序列，从后往前染色，每次染可以染的最小编号的颜色，
57 由完美消去序列的定义，序列任一后缀的点的导出子图中，由该后缀第一个元素
58 及其邻接点导出的子图一定是完全图，所以，序列中 某一元素染的颜色编号是该
59 完全图的大小。所以最小染色数小于等于最大团的点数，而显然前者又大于等于后者，
60 故弦图的最小染色数等于最大团的大小。
61 */
62 #include <cstdio>
63 #include <vector>
64 #define maxn 10010
65 using namespace std;
66 int n, m;
67 vector<int> g[maxn];
68 bool done[maxn];
69 int label[maxn], pos[maxn];
70 int msc() {
71     int rt = 0;
72     for( int i=n; i>=1; i-- ) {
73         int mu = 0;
74         for( int u=1; u<=n; u++ ) {
75             if( !done[u] ) {
76                 if( !mu || label[u]>label[mu] )
77                     mu = u;
78             }
79         }
80         done[mu] = true;
81         pos[mu] = i;
82         int cnt = 0;
83         for( int t=0; t<g[mu].size(); t++ ) {
84             int v = g[mu][t];
85             if( done[v] ) {
86                 cnt++;
87             } else {
88                 label[v]++;
89             }
90         }
91         rt = max( rt, cnt+1 );
92     }
93     return rt;
94 }
95 int main() {
96     scanf( "%d%d", &n, &m );
97     for( int i=1,u,v; i<=m; i++ ) {
98         scanf( "%d%d", &u, &v );
99         g[u].push_back(v);
100        g[v].push_back(u);
101    }
102    printf( "%d\n", msc() );
103 }

```

## 集合幂级数

```

1 #include <cstdio>
2 const int N = 10;

```

```

3 int n, U;
4 int a[1<<N], b[1<<N], c[1<<N];
5 void trans( int a[], int flag ) {
6     for( int b=0; b<n; b++ ) {
7         int u = U ^ (1<<b);
8         for( int s=u,t=1<<(n-1); t; s=(s-1)&u,t-- ) {
9             int l=a[s], r=a[s|(1<<b)];
10            /* NOT AND
11             if( flag==1 ) { a[s] = l+r; a[s|(1<<b)] = r;
12             } else { a[s] = r; a[s|(1<<b)] = l-r; } */
13            /* NOT XOR
14             if( flag==1 ) { a[s] = l+r; a[s|(1<<b)] = l-r;
15             } else { a[s] = (l-r)/2; a[s|(1<<b)] = (l+r)/2; } */
16            /* NOT OR
17             if( flag==1 ) { a[s] = l; a[s|(1<<b)] = l+r;
18             } else { a[s] = r-l; a[s|(1<<b)] = l; } */
19            /* OR
20             if( flag==1 ) { a[s] = l; a[s|(1<<b)] = l+r;
21             } else { a[s] = l; a[s|(1<<b)] = r-l; } */
22            /* AND
23             if( flag==1 ) { a[s] = l+r; a[s|(1<<b)] = r;
24             } else { a[s] = l-r; a[s|(1<<b)] = r; } */
25            /* XOR
26             if( flag==1 ) { a[s] = l+r; a[s|(1<<b)] = l-r;
27             } else { a[s] = (l+r)/2; a[s|(1<<b)] = (l-r)/2; } */
28        }
29    }
30 }
31 int main() {
32     scanf( "%d", &n );
33     U = (1<<n)-1;
34     for( int i=0; i<=U; i++ ) scanf( "%d", a+i );
35     for( int i=0; i<=U; i++ ) scanf( "%d", b+i );
36     trans(a,1); trans(b,1);
37     for( int s=0; s<=U; s++ ) c[s] = a[s]*b[s];
38     trans(c,-1);
39     for( int s=0; s<=U; s++ ) printf( "%d ", c[s] );
40     printf( "\n" );
41 }

```

## 其他

### Java Hints

```

1 import java.util.*;
2 import java.math.*;
3 import java.io.*;
4 public class Main{
5     static class Task{
6         void solve(int testId, InputReader cin, PrintWriter cout) {
7             // Write down the code you want
8         }
9     };
10    public static void main(String args[]) {
11        InputStream inputStream = System.in;
12        OutputStream outputStream = System.out;
13        InputReader in = new InputReader(inputStream);
14        PrintWriter out = new PrintWriter(outputStream);
15        TaskA solver = new TaskA();
16        solver.solve(1, in, out);
17        out.close();
18    }
19    static class InputReader {
20        public BufferedReader reader;
21        public StringTokenizer tokenizer;
22        public InputReader(InputStream stream) {
23            reader = new BufferedReader(new InputStreamReader(stream), 32768);
24            tokenizer = null;
25        }
26        public String next() {
27            while (tokenizer == null || !tokenizer.hasMoreTokens()) {
28                try {
29                    tokenizer = new StringTokenizer(reader.readLine());
30                } catch (IOException e) {
31                    throw new RuntimeException(e);
32                }
33            }
34            return tokenizer.nextToken();
35        }
36        public int nextInt() {
37            return Integer.parseInt(next());
38        }
39    }

```

```

39     }
40 };
41 // Arrays
42 int a[];
43 .fill(a[,int fromIndex,int toIndex],val);|.sort(a[, int fromIndex, int toIndex])
44 // String
45 String s;
46 .charAt(int i);|compareTo(String)|compareToIgnoreCase ()|contains(String)|
47 length ()|substring(int l, int len)
48 // BigInteger
49 .abs ()|.add ()|bitLength ()|subtract ()|divide ()|remainder ()|divideAndRemainder ()|
50 modPow(b, c)|pow(int) | multiply () | compareTo () |
51 gcd () | intValue () | longValue () | isProbablePrime(int c) (1 - 1/2^c) |
52 nextProbablePrime () | shiftLeft(int) | valueOf ()
53 // BigDecimal
54 .ROUND_CEILING | ROUND_DOWN_FLOOR | ROUND_HALF_DOWN | ROUND_HALF_EVEN | ROUND_HALF_UP | ROUND_UP
55 .divide(BigDecimal b, int scale , int round_mode) | doubleValue () | movePointLeft(int) | pow(int) |
56 setScale(int scale , int round_mode) | stripTrailingZeros ()
57 // StringBuilder
58 StringBuilder sb = new StringBuilder ();
59 sb.append(elem) | out.println(sb)
60 // TODO Java STL 的使用方法以及上面这些方法的检验

```

## 常用结论

### 上下界网络流

$B(u, v)$  表示边  $(u, v)$  流量的下界,  $C(u, v)$  表示边  $(u, v)$  流量的上界,  $F(u, v)$  表示边  $(u, v)$  的流量。设  $G(u, v) = F(u, v) - B(u, v)$ , 显然有:  $0 \leq G(u, v) \leq C(u, v) - B(u, v)$

#### 无源汇的上下界可行流

建立超级源点  $S^*$  和超级汇点  $T^*$ , 对于原图每条边  $(u, v)$  在新网络中连如下三条边:  $S^* \rightarrow v$ , 容量为  $B(u, v)$ ;  $u \rightarrow T^*$ , 容量为  $B(u, v)$ ;  $u \rightarrow v$ , 容量为  $C(u, v) - B(u, v)$ 。最后求新网络的最大流, 判断从超级源点  $S^*$  出发的边是否都满流即可, 边  $(u, v)$  的最终解中的实际流量为  $G(u, v) + B(u, v)$ 。

#### 有源汇的上下界可行流

从汇点  $T$  到源点  $S$  连一条上界为  $\infty$ , 下界为 0 的边。按照无源汇的上下界可行流一样做即可, 流量即为  $T \rightarrow S$  边上的流量。

#### 有源汇的上下界最大流

1. 在有源汇的上下界可行流中, 从汇点  $T$  到源点  $S$  的边改为连一条上界为  $\infty$ , 下届为  $x$  的边。 $x$  满足二分性质, 找到最大的  $x$  使得新网络存在无源汇的上下界可行流即为原图的最大流。
2. 从汇点  $T$  到源点  $S$  连一条上界为  $\infty$ , 下界为 0 的边, 变成无源汇的网络。按照无源汇的上下界可行流的方法, 建立超级源点  $S^*$  和超级汇点  $T^*$ , 求一遍  $S^* \rightarrow T^*$  的最大流, 再将汇点  $T$  到源点  $S$  的这条边拆掉, 求一次  $S \rightarrow T$  的最大流即可。

#### 有源汇的上下界最小流

1. 在有源汇的上下界可行流中, 从汇点  $T$  到源点  $S$  的边改为连一条上界为  $x$ , 下界为 0 的边。 $x$  满足二分性质, 找到最小的  $x$  使得新网络存在无源汇的上下界可行流即为原图的最小流。
2. 按照无源汇的上下界可行流的方法, 建立超级源点  $S^*$  与超级汇点  $T^*$ , 求一遍  $S^* \rightarrow T^*$  的最大流, 但是注意这一次不加上汇点  $T$  到源点  $S$  的这条边, 即不使之改为无源汇的网络去求解。求完后, 再加上那条汇点  $T$  到源点  $S$  上界  $\infty$  的边。因为这条边下界为 0, 所以  $S^*$ ,  $T^*$  无影响, 再直接求一次  $S^* \rightarrow T^*$  的最大流。若超级源点  $S^*$  出发的边全部满流, 则  $T \rightarrow S$  边上的流量即为原图的最小流, 否则无解。

### 上下界费用流

来源: BZOJ 3876 设汇  $t$ , 源  $s$ , 超级源  $S$ , 超级汇  $T$ , 本质是每条边的下界为 1, 上界为  $MAX$ , 跑一遍有源汇的上下界最小费用最小流。(因为上界无穷大, 所以只要满足所有下界的最小费用最小流)

1. 对每个点  $x$ : 从  $x$  到  $t$  连一条费用为 0, 流量为  $MAX$  的边, 表示可以任意停止当前的剧情 (接下来的剧情从更优的路径去走, 画个样例就知道了)
2. 对于每一条边权为  $z$  的边  $x \rightarrow y$ :
  - 从  $S$  到  $y$  连一条流量为 1, 费用为  $z$  的边, 代表这条边至少要被走一次。
  - 从  $x$  到  $y$  连一条流量为  $MAX$ , 费用为  $z$  的边, 代表这条边除了至少走的一次之外还可以随便走。
  - 从  $x$  到  $T$  连一条流量为 1, 费用为 0 的边。(注意是每一条  $x \rightarrow y$  的边都连, 或者你可以记下  $x$  的出边数  $K_x$ , 连一次流量为  $K_x$ , 费用为 0 的边)。

建完图后从  $S$  到  $T$  跑一遍费用流, 即可。(当前跑出来的就是满足上下界的最小费用最小流了)

### 弦图相关

1. 团数  $\leq$  色数, 弦图团数 = 色数
2. 设  $next(v)$  表示  $N(v)$  中最前的点. 令  $w^*$  表示所有满足  $A \in B$  的  $w$  中最后的一个点, 判断  $v \cup N(v)$  是否为极大团, 只需判断是否存在一个  $w$ , 满足  $Next(w) = v$  且  $|N(v)| + 1 \leq |N(w)|$  即可.
3. 最小染色: 完美消除序列从后往前依次给每个点染色, 给每个点染上可以染的最小的颜色
4. 最大独立集: 完美消除序列从前往后能选就选
5. 弦图最大独立集数 = 最小团覆盖数, 最小团覆盖: 设最大独立集为  $\{p_1, p_2, \dots, p_t\}$ , 则  $\{p_1 \cup N(p_1), \dots, p_t \cup N(p_t)\}$  为最小团覆盖

### Bernoulli 数

1. 初始化:  $B_0(n) = 1$
2. 递推公式:  $B_m(n) = n^m - \sum_{k=0}^{m-1} \binom{m}{k} \frac{B_k(n)}{m-k+1}$
3. 应用:  $\sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} n^{m+1-k}$

### 常见错误

1. 数组或者变量类型开错, 例如将 double 开成 int;
2. 函数忘记返回返回值;
3. 初始化数组没有初始化完全;
4. 对空间限制判断不足导致 MLE

### 博弈游戏

#### 巴什博弈

1. 只有一堆  $n$  个物品, 两个人轮流从这堆物品中取物, 规定每次至少取一个, 最多取  $m$  个。最后取光者得胜。
2. 显然, 如果  $n = m + 1$ , 那么由于一次最多只能取  $m$  个, 所以, 无论先取者拿走多少个, 后取者都能够一次拿走剩余的物品, 后者取胜。因此我们发现了如何取胜的法则: 如果  $n = m + 1r + s$ , ( $r$  为任意自然数,  $s \leq m$ ), 那么先取者要拿走  $s$  个物品, 如果后取者拿走  $k$  ( $k \leq m$ ) 个, 那么先取者再拿走  $m + 1 - k$  个, 结果剩下  $(m + 1)(r - 1)$  个, 以后保持这样的取法, 那么先取者肯定获胜。总之, 要保持给对手留下  $(m + 1)$  的倍数, 就能最后获胜。

#### 威佐夫博弈

1. 有两堆各若干个物品, 两个人轮流从某一堆或同时从两堆中取同样多的物品, 规定每次至少取一个, 多者不限, 最后取光者得胜。
2. 判断一个局势  $(a, b)$  为奇异局势 (必败态) 的方法:  $a_k = [k(1 + \sqrt{5})/2]$   $b_k = a_k + k$

#### 阶梯博弈

1. 博弈在一列阶梯上进行, 每个阶梯上放着自然数个点, 两个人进行阶梯博弈, 每一步则是将一个阶梯上的若干个 (至少一个) 移到前面去, 最后没有点可以移动的人输。
2. 解决方法: 把所有奇数阶梯看成  $N$  堆石子, 做 NIM。(把石子从奇数堆移动到偶数堆可以理解为拿走石子, 就相当于几个奇数堆的石子在做 Nim)

### 图上删边游戏

#### 链的删边游戏

1. 游戏规则: 对于一条链, 其中一个端点是根, 两人轮流删边, 脱离根的部分也算被删去, 最后没边可删的人输。
2. 做法:  $sg[i] = n - dist(i) - 1$  (其中  $n$  表示总点数,  $dist(i)$  表示离根的距离)

#### 树的删边游戏

1. 游戏规则: 对于一棵有根树, 两人轮流删边, 脱离根的部分也算被删去, 没边可删的人输。
2. 做法: 叶子结点的  $sg = 0$ , 其他节点的  $sg$  等于儿子结点的  $sg + 1$  的异或和。

#### 局部连通图的删边游戏

1. 游戏规则: 在一个局部连通图上, 两人轮流删边, 脱离根的部分也算被删去, 没边可删的人输。局部连通图的构图规则是, 在一棵基础树上加边得到, 所有形成的环保证不共用边, 且只与基础树有一个公共点。
2. 做法: 去掉所有的偶环, 将所有的奇环变为长度为 1 的链, 然后做树的删边游戏。

## 常用数学公式

### 求和公式

- $\sum_{k=1}^n (2k-1)^2 = \frac{n(4n^2-1)}{3}$
- $\sum_{k=1}^n k^3 = [\frac{n(n+1)}{2}]^2$
- $\sum_{k=1}^n (2k-1)^3 = n^2(2n^2-1)$
- $\sum_{k=1}^n k^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$
- $\sum_{k=1}^n k^5 = \frac{n^2(n+1)^2(2n^2+2n-1)}{12}$
- $\sum_{k=1}^n k(k+1) = \frac{n(n+1)(n+2)}{3}$
- $\sum_{k=1}^n k(k+1)(k+2) = \frac{n(n+1)(n+2)(n+3)}{4}$
- $\sum_{k=1}^n k(k+1)(k+2)(k+3) = \frac{n(n+1)(n+2)(n+3)(n+4)}{5}$
- $\frac{1}{(1-x)^{n+1}} = \sum_{i=0}^n \binom{i+n}{i} x^i$
- $\frac{1}{\sqrt{1-4x}} = \sum_{i=0}^n \binom{2i}{i} x^i$

### 斐波那契数列

- $fib_0 = 0, fib_1 = 1, fib_n = fib_{n-1} + fib_{n-2}$
- $fib_{n+2} \cdot fib_n - fib_{n+1}^2 = (-1)^{n+1}$
- $fib_{-n} = (-1)^{n-1} fib_n$
- $fib_{n+k} = fib_k \cdot fib_{n+1} + fib_{k-1} \cdot fib_n$
- $gcd(fib_m, fib_n) = fib_{gcd(m,n)}$
- $fib_m | fib_n^2 \Leftrightarrow n | fib_m$

### 错排公式

$$D_n = (n-1)(D_{n-2} - D_{n-1}) == n! \cdot (1 - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + \frac{(-1)^n}{n!})$$

### 莫比乌斯函数

$$\mu(n) = \begin{cases} 1 & \text{若 } n=1 \\ (-1)^k & \text{若 } n \text{ 无平方数因子, 且 } n=p_1 p_2 \dots p_k \\ 0 & \text{若 } n \text{ 有大于1的平方数因数} \end{cases}$$

$$\sum_{d|n} \mu(d) = \begin{cases} 1 & \text{若 } n=1 \\ 0 & \text{其他情况} \end{cases}$$

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu(d) g(\frac{n}{d})$$

$$g(x) = \sum_{n=1}^{[x]} f(\frac{x}{n}) \Leftrightarrow f(x) = \sum_{n=1}^{[x]} \mu(n) g(\frac{x}{n})$$

### Burnside 引理

设  $G$  是一个有限群, 作用在集合  $X$  上. 对每个  $g$  属于  $G$ , 令  $X^g$  表示  $X$  中在  $g$  作用下的不动元素, 轨道数 (记作  $|X/G|$ ) 为  $|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$ .

### 五边形数定理

$$\text{设 } p(n) \text{ 是 } n \text{ 的拆分数, 有 } p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k-1} p\left(n - \frac{k(3k-1)}{2}\right)$$

## 树的计数

- 有根树计数:  $n+1$  个结点的有根树的个数为  $a_{n+1} = \frac{\sum_{j=1}^n j \cdot a_j \cdot S_{n,j}}{n}$ , 其中,  $S_{n,j} = \sum_{i=1}^{n/j} a_{n+1-ij} = S_{n-j,j} + a_{n+1-j}$
- 无根树计数: 当  $n$  为奇数时,  $n$  个结点的无根树的个数为  $a_n - \sum_{i=1}^{n/2} a_i a_{n-i}$ , 当  $n$  为偶数时,  $n$  个结点的无根树的个数为  $a_n - \sum_{i=1}^{n/2} a_i a_{n-i} + \frac{1}{2} a_{\frac{n}{2}} (a_{\frac{n}{2}} + 1)$
- $n$  个结点的完全图的生成树个数为:  $n^{n-2}$
- 矩阵 - 树定理: 图  $G$  由  $n$  个结点构成, 设  $\mathbf{A}[G]$  为图  $G$  的邻接矩阵,  $\mathbf{D}[G]$  为图  $G$  的度数矩阵, 则图  $G$  的不同生成树的个数为  $\mathbf{C}[G] = \mathbf{D}[G] - \mathbf{A}[G]$  的任意一个  $n-1$  阶主子式的行列式值。

## 欧拉公式

平面图的顶点个数、边数和面的个数有如下关系:  $V - E + F = C + 1$

其中,  $V$  是顶点的数目,  $E$  是边的数目,  $F$  是面的数目,  $C$  是组成图形的连通部分的数目。当图是单连通图的时候, 公式简化为:  $V - E + F = 2$

## 皮克定理

给定顶点坐标均是整数 (或正方形格点) 的简单多边形, 其面积  $A$  和内部格点数目  $i$ 、边上格点数目  $b$  的关系:  $A = i + \frac{b}{2} - 1$

## 牛顿恒等式

设

$$\prod_{i=1}^n (x - x_i) = a_n + a_{n-1}x + \dots + a_1x^{n-1} + a_0x^n$$

$$p_k = \sum_{i=1}^n x_i^k$$

则

$$a_0 p_k + a_1 p_{k-1} + \dots + a_{k-1} p_1 + k a_k = 0$$

特别地, 对于

$$|\mathbf{A} - \lambda \mathbf{E}| = (-1)^n (a_n + a_{n-1} \lambda + \dots + a_1 \lambda^{n-1} + a_0 \lambda^n)$$

有

$$p_k = \text{Tr}(\mathbf{A}^k)$$

## 平面几何公式

### 三角形

- 面积:  $S = \frac{a \cdot H_a}{2} = \frac{ab \cdot \sin C}{2} = \sqrt{p(p-a)(p-b)(p-c)} \left( \frac{a+b+c}{2} \right)$
- 中线:  $M_a = \frac{\sqrt{2(b^2+c^2)-a^2}}{2} = \frac{\sqrt{b^2+c^2+2bc \cdot \cos A}}{2}$
- 角平分线:  $T_a = \frac{\sqrt{bc \cdot [(b+c)^2 - a^2]}}{b+c} = \frac{2bc \cos \frac{A}{2}}{b+c}$
- 高线:  $H_a = b \sin C = c \sin B = \sqrt{b^2 - (\frac{a^2+b^2-c^2}{2a})^2}$
- 内切圆半径

$$r = \frac{S}{p} = \frac{\arcsin \frac{B}{2} \cdot \sin \frac{C}{2}}{\sin \frac{B+C}{2}} = 4R \cdot \sin \frac{A}{2} \sin \frac{B}{2} \sin \frac{C}{2}$$

$$= \sqrt{\frac{(p-a)(p-b)(p-c)}{p}} = p \cdot \tan \frac{A}{2} \tan \frac{B}{2} \tan \frac{C}{2}$$

- 外接圆半径:  $R = \frac{abc}{4S} = \frac{a}{2 \sin A} = \frac{b}{2 \sin B} = \frac{c}{2 \sin C}$

**四边形**

$D_1, D_2$  为对角线,  $M$  为对角线中点连线,  $A$  为对角线夹角,  $p$  为半周长

1.  $a^2 + b^2 + c^2 + d^2 = D_1^2 + D_2^2 + 4M^2$
2.  $S = \frac{1}{2}D_1D_2\sin A$
3. 对于圆内接四边形:  $ac + bd = D_1D_2$
4. 对于圆内接四边形:  $S = \sqrt{(p-a)(p-b)(p-c)(p-d)}$

**正  $n$  边形**

$R$  为外接圆半径,  $r$  为内切圆半径

1. 中心角:  $A = \frac{2\pi}{n}$
2. 内角:  $C = \frac{n-2}{n}\pi$
3. 边长:  $a = 2\sqrt{R^2 - r^2} = 2R \cdot \sin \frac{A}{2} = 2r \cdot \tan \frac{A}{2}$
4. 面积:  $S = \frac{nar}{2} = nr^2 \cdot \tan \frac{A}{2} = \frac{nR^2}{2} \cdot \sin A = \frac{na^2}{4 \cdot \tan \frac{A}{2}}$

**圆**

1. 弧长:  $l = rA$
2. 弦长:  $a = 2\sqrt{2hr - h^2} = 2r \cdot \sin \frac{A}{2}$
3. 弓形高:  $h = r - \sqrt{r^2 - \frac{a^2}{4}} = r(1 - \cos \frac{A}{2}) = \frac{1}{2} \cdot \arctan \frac{A}{4}$
4. 扇形面积:  $S_1 = \frac{rl}{2} = \frac{r^2 A}{2}$
5. 弓形面积:  $S_2 = \frac{rl - a(r-h)}{2} = \frac{r^2}{2}(A - \sin A)$

**棱柱**

1. 体积 ( $A$  为底面积,  $h$  为高):  $V = Ah$
2. 侧面积 ( $l$  为棱长,  $p$  为直截面周长):  $S = lp$
3. 全面积:  $T = S + 2A$

**棱锥**

1. 体积 ( $A$  为底面积,  $h$  为高):  $V = \frac{1}{3}Ah$
2. 正棱锥侧面积 ( $l$  为棱长,  $p$  为直截面周长):  $S = lp$
3. 正棱锥全面积:  $T = S + 2A$

**棱台**

1. 体积 ( $A_1, A_2$  为上下底面积,  $h$  为高):  $V = (A_1 + A_2 + \sqrt{A_1A_2}) \cdot \frac{h}{3}$
2. 正棱台侧面积 ( $p_1, p_2$  为上下底面周长,  $l$  为斜高):  $S = \frac{p_1 + p_2}{2}l$
3. 正棱台全面积:  $T = S + A_1 + A_2$

**圆柱**

1. 侧面积:  $S = 2\pi rh$
2. 全面积:  $T = 2\pi r(h + r)$
3. 体积:  $V = \pi r^2 h$

**圆锥**

1. 母线:  $l = \sqrt{h^2 + r^2}$
2. 侧面积:  $S = \pi rl$
3. 全面积:  $T = \pi r(l + r)$
4. 体积:  $V = \frac{\pi}{3}r^2 h$

**圆台**

1. 母线:  $l = \sqrt{h^2 + (r_1 - r_2)^2}$
2. 侧面积:  $S = \pi(r_1 + r_2)l$
3. 全面积:  $T = \pi r_1(l + r_1) + \pi r_2(l + r_2)$
4. 体积:  $V = \frac{\pi}{3}(r_1^2 + r_2^2 + r_1r_2)h$

**球**

1. 全面积:  $T = 4\pi r^2$
2. 体积:  $V = \frac{4}{3}\pi r^3$

**球台**

1. 侧面积:  $S = 2\pi rh$
2. 全面积:  $T = \pi(2rh + r_1^2 + r_2^2)$
3. 体积:  $V = \frac{\pi h[3(r_1^2 + r_2^2) + h^2]}{6}$

**球扇形**

1. 全面积 ( $h$  为球冠高,  $r_0$  为球冠底面半径):  $T = \pi r(2h + r_0)$
2. 体积:  $V = \frac{2}{3}\pi r^2 h$

**立体几何公式****球面三角公式**

设  $a, b, c$  是边长,  $A, B, C$  是所对的二面角, 有余弦定理

$$\cos a = \cos b \cdot \cos c + \sin b \cdot \sin c \cdot \cos A$$

正弦定理

$$\frac{\sin A}{\sin a} = \frac{\sin B}{\sin b} = \frac{\sin C}{\sin c}$$

三角形面积是  $A + B + C - \pi$

**四面体体积公式**

$U, V, W, u, v, w$  是四面体的 6 条棱,  $U, V, W$  构成三角形,  $(U, u), (V, v), (W, w)$  互为对棱, 则

$$V = \frac{\sqrt{(s-2a)(s-2b)(s-2c)(s-2d)}}{192uvw}$$

其中

$$\begin{cases} a &= \sqrt{xYZ}, \\ b &= \sqrt{yZX}, \\ c &= \sqrt{zXY}, \\ d &= \sqrt{xyz}, \\ s &= a + b + c + d, \\ X &= (w - U + v)(U + v + w), \\ x &= (U - v + w)(v - w + U), \\ Y &= (u - V + w)(V + w + u), \\ y &= (V - w + u)(w - u + V), \\ Z &= (v - W + u)(W + u + v), \\ z &= (W - u + v)(u - v + W) \end{cases}$$



附录  
NTT 素数及原根列表

Id	Primes	PRT	Id	Primes	PRT	Id	Primes	PRT
1	7340033	3	38	311427073	7	75	786432001	7
2	13631489	15	39	330301441	22	76	799014913	13
3	23068673	3	40	347078657	3	77	800063489	3
4	26214401	3	41	359661569	3	78	802160641	11
5	28311553	5	42	361758721	29	79	818937857	5
6	69206017	5	43	377487361	7	80	824180737	5
7	70254593	3	44	383778817	5	81	833617921	13
8	81788929	7	45	387973121	6	82	850395137	3
9	101711873	3	46	399507457	5	83	862978049	3
10	104857601	3	47	409993217	3	84	880803841	26
11	111149057	3	48	415236097	5	85	883949569	7
12	113246209	7	49	447741953	3	86	897581057	3
13	120586241	6	50	459276289	11	87	899678209	7
14	132120577	5	51	463470593	3	88	907018241	3
15	136314881	3	52	468713473	5	89	913309697	3
16	138412033	5	53	469762049	3	90	918552577	5
17	141557761	26	54	493879297	10	91	919601153	3
18	147849217	5	55	531628033	5	92	924844033	5
19	155189249	6	56	576716801	6	93	925892609	3
20	158334977	3	57	581959681	11	94	935329793	3
21	163577857	23	58	595591169	3	95	938475521	3
22	167772161	3	59	597688321	11	96	940572673	7
23	169869313	5	60	605028353	3	97	943718401	7
24	185597953	5	61	635437057	11	98	950009857	7
25	186646529	3	62	639631361	6	99	957349889	6
26	199229441	3	63	645922817	3	100	962592769	7
27	204472321	19	64	648019969	17	101	972029953	10
28	211812353	3	65	655360001	3	102	975175681	17
29	221249537	3	66	666894337	5	103	976224257	3
30	230686721	6	67	683671553	3	104	985661441	3
31	246415361	3	68	710934529	17	105	998244353	3
32	249561089	3	69	715128833	3	106	1004535809	3
33	257949697	5	70	718274561	3	107	1007681537	3
34	270532609	22	71	740294657	3	108	1012924417	5
35	274726913	3	72	745537537	5	109	1045430273	3
36	290455553	3	73	754974721	11	110	1051721729	6
37	305135617	5	74	770703361	11	111	1053818881	7