

1. (10 %) **Free Blocks Management Using a Linked List** Consider a file system managing free blocks by using linked lists. The table below shows the final two blocks storing free blocks. Fill the empty tables below to show the changes which occur in the tables after the following scenarios. Highlight the changes using a color pencil.

- (a) Five new blocks are allocated
- (b) The block 22 is freed
- (c) Another 5 blocks are allocated
- (d) Another block is allocated
- (e) Another three blocks are allocated
- (f) Four blocks (23456, 8345345, 56, and 634534) are freed

Bloc k #	17	18							
Next Bloc k	18	0							
	4589	24353							
	43546	98745							
	718	76345							
	345	9877							
	23456	7345							
	8345345	34535							
	634534	154698							
	3478	967							
	56	8657							
Bloc k #	17	18	Bloc k #	17	18	Block #	17	18	
Next Bloc k	18	0	Next Bloc k	18	0	Next Block	18	0	
	4589	24353		4589	24353			24353	
	43546	98745		43546	98745			98745	
	718	76345		718	76345			76345	
	345	9877		345	9877			9877	
		7345		22	7345			7345	
		34535			34535			34535	
		154698			154698			154698	
		967			967			967	
		8657			8657			8657	
Bloc k #	18		Bloc k #	18		Block #			
Next Bloc k	0		Next Bloc k	0		Next Block			

	24353			24353				24353
	98745			98745				98745
	76345			76345				76345
	9877			9877				9877
	7345			7345				7345
	34535			34535				34535
	154698							23456
	967							8345345
	17							56

2. **Free Blocks Management — Comparision** Given the two memory footprint scenarios for Free Blocks Management as presented in class. State the condition under which the linked list approach uses less space than the bitmap approach.

Wenn die disc schon zb fast voll ist, und das free block management einfach nur mehr aus ein paar blocks besteht und dadurch weniger platz braucht wie die ganze bitmap.