



# SwapVM + Aqua Security Audit Report

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- None

## **Table of contents**

See table

- SwapVM + Aqua Security Audit Report
- Table of contents
- About Vadim Fadeev
- Disclaimer
- Risk Classification
- Audit Details
  - Scope
  - Methodology
  - Threat Model and Trust Assumptions
- System Overview

- Aqua (Shared Liquidity Layer)
- SwapVM (Programmed Swaps)
- Executive Summary
  - Issues found
  - Key Takeaways
- Findings
  - Medium
    - \* [M-1] Division-by-zero if `feeBps == BPS` in `ExactOut` fee calculation
    - \* [M-2] Division-by-zero if `initialLiquidity == 0` in `XYCCconcentrate`
    - \* [M-3] `Decay::decayXD` underflow when offset exceeds balance
    - \* [M-4] `VM::runLoop` lacks bounds checking for malformed programs
    - \* [M-5] Invalid opcode causes array out-of-bounds panic
  - Low
    - \* [L-1] `Fee` does not enforce `feeBps <= BPS` at runtime for flat fees
    - \* [L-2] `Decay::decayPeriod == 0` should be validated
    - \* [L-3] `Controls::onlyTakerTokenSupplyShareGte` has potential overflow
    - \* [L-4] Unused imports
    - \* [L-5] State changes without event emission in `Invalidators`
- Appendix
  - Hardening Checklist
  - References

## About Vadim Fadeev

## Disclaimer

The Vadim Fadeev team makes all effort to find as many vulnerabilities in the code in the given time period, but holds no responsibilities for the findings provided in this document. A security audit by the team is not an endorsement of the underlying business or product. The audit was time-boxed and the review of the code was solely on the security aspects of the solidity implementation of the contracts.

## Risk Classification

Likelihood	High Impact	Medium Impact	Low Impact
High	H	H/M	M
Medium	H/M	M	M/L
Low	M	M/L	L

## Audit Details

The findings described in this document correspond the following commit hash:

```
1 89cedfa608aa9fdc60d8cd936b76a540e4b9346f
```

## Scope

This report covers the SwapVM + Aqua integration codebase.

## SwapVM

**In scope** - `src/*.sol` - `src/routers/AquaSwapVMRouter.sol` - `src/opcodes/AquaOpcodes.sol` - `src/instructions/Balances.sol` - `src/instructions/Controls.sol` - `src/instructions/Decay.sol` - `src/instructions/Extraction.sol` - `src/instructions/Fee.sol` - `src/instructions/PeggedSwap.sol` - `src/instructions/XYCCConcentrate.sol` - `src/instructions/XYCSwap.sol` - `src/libs/MakerTraits.sol` - `src/libs/PeggedSwapMath.sol` - `src/libs/TakerTraits.sol` - `src/libs/VM.sol`

## Methodology

- Manual review of instruction contracts and shared libraries
- Reasoning about invariants and state transitions (especially: balances bookkeeping, fee layering, swap math)
- Threat modeling around “maker-programmed” execution and “taker-executed” swaps
- Cross-reference with Solodit Security Checklist
- Identification of:
  - panic reverts / implicit underflow/overflow reliance
  - initialization sentinels and state ambiguity

- external call surfaces (fee providers, extraction targets)
- denial-of-service vectors (gas griefing / malformed programs)

## Threat Model and Trust Assumptions

SwapVM is explicitly a **programmable execution environment**: - Makers define a program (bytecode-like instruction stream). - Takers execute it (often via routers) and must assume the maker program may be adversarial.

Therefore: - Many issues are “**order-level misconfiguration or malicious program risks**” rather than protocol-wide vulnerabilities. - Wherever possible, the protocol should fail with **clear errors**, enforce **basic numeric bounds**, and prevent **state ambiguity** that can be exploited across calls.

External call surfaces: - `Fee._dynamicProtocolFeeAmountInXD` calls an untrusted fee provider via `staticcall`. - `Extraction._extraction` calls an arbitrary target contract (view or non-view depending on context). These are inherently risky and should be treated as part of the trusted computing base *for that specific program/order*.

## System Overview

### Aqua (Shared Liquidity Layer)

Aqua is a shared liquidity layer that manages balances/allowances between makers and “apps/strategies”, allowing apps to pull from maker wallets while tracking balances.

### SwapVM (Programmed Swaps)

SwapVM provides: - A **VM** (`Context`, `VM`, `ContextLib::runLoop`) executing an instruction stream from calldata. - A set of **instructions** (e.g., `Balances`, `Fee`, `XYCSwap`, `PeggedSwap`, `Decay`, `Controls`, `Extraction`). - Routers/opcode registries (e.g., `AquaOpcodes`, `AquaSwapVMRouter`) assembling instruction tables.

Stateful instructions: - `Balances` maintains per-order token balances. - `Decay` maintains per-order decaying offsets. - `XYCConcentrate` tracks per-order liquidity scaling.

## Executive Summary

### Issues found

Severity	Number of issues found
High	0
Medium	5
Low	5
Gas Optimizations	0
<b>Total</b>	<b>10</b>

## Key Takeaways

- Two **Medium severity** division-by-zero vulnerabilities were identified in `Fee.sol` and `XYCConcentrate.sol` that can permanently brick orders.
- Several core instructions rely on **implicit panics** (division-by-zero, underflow) rather than explicit validation with custom errors, which harms:
  - debuggability
  - integrator UX
  - security monitoring / incident response
- External call surfaces (`dynamic fee provider, extraction`) are legitimate features and should be clearly documented as **trust boundaries**.

## Findings

### Medium

#### [M-1] Division-by-zero if `feeBps == BPS` in ExactOut fee calculation

**File:** `src/instructions/Fee.sol:196`

**Description:**

In the ExactOut path, `_feeAmountIn()` computes:

```
1 feeAmountIn = ctx.swap.amountIn * feeBps / (BPS - feeBps);
```

If `feeBps == BPS` (100% fee), the denominator becomes zero, causing a panic revert.

**Root Cause:**

- `FeeArgsBuilder.buildFlatFee()` validates `feeBps <= BPS` (allows equality) - `parseFlatFee()` has no runtime validation - Programs can encode `feeBps = BPS` directly

**Impact:**

- Orders with 100% fee permanently bricked in ExactOut mode - Panic revert instead of clear custom error

**Verification:**

```
1 // Fee.sol:196
2 feeAmountIn = ctx.swap.amountIn * feeBps / (BPS - feeBps);
3 // If feeBps = 1e9 (BPS), then (BPS - feeBps) = 0 → division by zero
```

**Recommended Mitigation:**

```
1 function _feeAmountIn(Context memory ctx, uint256 feeBps) internal
2     returns (uint256 feeAmountIn) {
3     require(ctx.swap.amountIn == 0 || ctx.swap.amountOut == 0, ...);
4     + require(feeBps < BPS, FeeBpsCannotBe100PercentForExactOut(feeBps));
5     // ...
6 }
```

**[M-2] Division-by-zero if `initialLiquidity == 0` in XYCConcentrate**

**File:** `src/instructions/XYCConcentrate.sol:113`

**Description:**

The `concentratedBalance()` function divides by `initialLiquidity`:

```
1 function concentratedBalance(..., uint256 initialLiquidity) public view
2     returns (uint256) {
3     uint256 currentLiquidity = liquidity[orderHash];
4     return currentLiquidity == 0
5         ? balance + delta
6         : balance + delta * currentLiquidity / initialLiquidity; // ←
7             Division!
8 }
```

**Attack Scenario:** 1. Maker creates order with `initialLiquidity = 0` in program args 2. First trade succeeds (`currentLiquidity == 0` takes first branch) 3. `_updateScales()` sets `liquidity[orderHash]` to non-zero value 4. Second trade: `currentLiquidity != 0`, divides by `initialLiquidity = 0` → **PANIC**

**Impact:**

- Order permanently bricked after first trade - Maker loses access to remaining liquidity - No recovery mechanism

**Verification:**

```

1 // XYCConcentrate.sol:111-114
2 function concentratedBalance(bytes32 orderHash, uint256 balance,
3     uint256 delta, uint256 initialLiquidity) public view returns (
4     uint256) {
5     uint256 currentLiquidity = liquidity[orderHash];
6     return currentLiquidity == 0 ? balance + delta : balance + delta *
7         currentLiquidity / initialLiquidity;
8 }
9 // Line 157: liquidity[ctx.query.orderHash] = Math.sqrt(newInv); //
10 Sets non-zero after first trade

```

**Recommended Mitigation:**

```

1 function parse2D(bytes calldata args, address tokenIn, address tokenOut
2     ) internal pure returns (...) {
3     // ... existing parsing ...
4     liquidity = uint256(bytes32(args.slice(64, 96, ...)));
5     + require(liquidity > 0, ConcentrateInitialLiquidityMustBeNonZero());
6 }

```

**[M-3] Decay::\_decayXD underflow when offset exceeds balance**

**File:** [src/instructions/Decay.sol:87](#)

**Description:**

```

1 ctx.swap.balanceOut -= _offsets[ctx.query.orderHash][ctx.query.tokenOut
2 ][false].getOffset(period);

```

If the decayed offset exceeds `balanceOut`, subtraction underflows → panic.

**How this happens:** 1. Initial `balanceOut` = 100 2. Large swap adds offset of 80 to `tokenOut` 3. Another swap reduces `balanceOut` to 50 via Balances instruction 4. Offset decays to 60, but 50 – 60 underflows

**Impact:** - Orders become temporarily or permanently unusable - Depends on interaction between Decay and balance-modifying instructions

**Recommended Mitigation:**

```

1 uint256 offsetOut = _offsets[ctx.query.orderHash][ctx.query.tokenOut][
    false].getOffset(period);
2 ctx.swap.balanceOut = ctx.swap.balanceOut > offsetOut ? ctx.swap.
    balanceOut - offsetOut : 0;

```

## [M-4] VM::runLoop lacks bounds checking for malformed programs

**File:** `src/libs/VM.sol:97-107`

### Description:

```

1 for (uint256 pc = ctx.vm.nextPC; pc < programBytes.length; ) {
2     unchecked {
3         uint256 opcode = uint8(programBytes[pc++]);           // Read byte 1
4         uint256 argsLength = uint8(programBytes[pc++]);      // Read byte 2
            - 00B if only 1 byte left!
5         uint256 nextPC = pc + argsLength;
6         bytes calldata args = programBytes[pc:nextPC];      // 00B if
            argsLength exceeds remaining
7         // ...
8     }
9 }

```

**Problem scenarios:** 1. **Single trailing byte:** Loop enters with 1 byte remaining, reading `argsLength` panics  
2. **argsLength exceeds remaining bytes:** Slice operation panics

**Impact:** - Malformed programs cause opaque panic reverts - No way for takers to validate program well-formedness before execution

### Recommended Mitigation:

```

1 for (uint256 pc = ctx.vm.nextPC; pc < programBytes.length; ) {
2     + require(pc + 1 < programBytes.length, RunLoopMalformedInstruction(
        pc));
3
4     uint256 opcode = uint8(programBytes[pc]);
5     uint256 argsLength = uint8(programBytes[pc + 1]);
6     uint256 nextPC = pc + 2 + argsLength;
7
8     + require(nextPC <= programBytes.length, RunLoopArgsExceedProgram(pc,
        argsLength));
9     + require(opcode < ctx.vm.opcodes.length, RunLoopInvalidOpcode(opcode));
10
11     bytes calldata args = programBytes[pc + 2:nextPC];

```

```
12      // ...
13 }
```

---

### [M-5] Invalid opcode causes array out-of-bounds panic

**File:** `src/libs/VM.sol:105`

**Description:**

```
1 ctx.vm.opcodes[opcode](ctx, args);
```

If `opcode >= ctx.vm.opcodes.length`, this causes an array index out-of-bounds panic.

**Impact:** - Malformed programs with invalid opcodes revert with opaque panics - No clear error message for debugging

**Recommended Mitigation:**

```
1 require(opcode < ctx.vm.opcodes.length, RunLoopInvalidOpcode(opcode,
                  ctx.vm.opcodes.length));
2 ctx.vm.opcodes[opcode](ctx, args);
```

---

### Low

### [L-1] Fee does not enforce feeBps <= BPS at runtime for flat fees

**File:** `src/instructions/Fee.sol:42-44`

**Description:**

`parseFlatFee()` has no bounds check; only `buildFlatFee()` validates. Programs can bypass the builder with arbitrary `feeBps` values.

**Impact:** - `feeBps > BPS` could cause underflow in fee subtraction - Opaque panics instead of clear errors

**Recommended Mitigation:**

```
1 function _feeAmountIn(Context memory ctx, uint256 feeBps) internal
    returns (uint256 feeAmountIn) {
2 +   require(feeBps <= BPS, FeeBpsOutOfRange(feeBps));
```

```

3      // ...
4 }
```

### [L-2] Decay::decayPeriod == 0 should be validated

**File:** [src/instructions/Decay.sol:45](#)

**Description:**

While analysis shows division-by-zero is practically unreachable due to timestamp logic, a `decayPeriod = 0` value is clearly invalid and should be rejected at parse time for defense-in-depth.

**Recommended Mitigation:**

```

1 function parse(bytes calldata args) internal pure returns (uint16
    period) {
2     period = uint16(bytes2(args.slice(0, 2, DecayMissingPeriodArg.
        selector)));
3 +    require(period > 0, DecayPeriodMustBeNonZero());
4 }
```

### [L-3] Controls::\_onlyTakerTokenSupplyShareGte has potential overflow

**File:** [src/instructions/Controls.sol](#)

**Description:**

The share calculation multiplies balance by 1e18:

```

1 require(totalSupply > 0 && balance * 1e18 >= minShareE18 * totalSupply,
    ...);
```

If `balance > type(uint256).max / 1e18 (~1.15e59)`, multiplication overflows.

**Impact:** - Tokens with extremely high decimals (27+) or supply could cause unexpected reverts - Panic error instead of clear message

**Recommended Mitigation:**

```

1 uint256 share = Math.mulDiv(balance, 1e18, totalSupply);
2 require(share >= minShareE18, ...);
```

## [L-4] Unused imports

**Files:** - `src/instructions/Fee.sol:7`- `src/instructions/FeeExperimental.sol:12`- `src/libs/TakerTraits.sol:11`

### Description:

Redundant import statements that are not used in the contracts:

```
1 // Fee.sol
2 import { Math } from "@openzeppelin/contracts/utils/math/Math.sol";  
3  
4 // FeeExperimental.sol
5 import { IProtocolFeeProvider } from "./interfaces/IProtocolFeeProvider.sol";  
6  
7 // TakerTraits.sol
8 import { ITakerCallbacks } from "../interfaces/ITakerCallbacks.sol";
```

**Impact:** - Slightly increases bytecode size - Reduces code clarity - May confuse auditors about expected functionality

**Recommended Mitigation:** Remove the unused import statements.

---

## [L-5] State changes without event emission in Invalidators

**File:** `src/instructions/Invalidators.sol:52-60`

### Description:

Three functions modify state but emit no events:

```
1 function invalidateBit(uint256 bitIndex) external {
2     // State change, no event
3 }
4
5 function invalidateTokenIn(bytes32 orderHash, address tokenIn) external
6     {
7     // State change, no event
8 }
9 function invalidateTokenOut(bytes32 orderHash, address tokenOut)
10    external {
11    // State change, no event
```

```
11 }
```

**Impact:** - Off-chain indexers cannot track invalidation changes - Reduced transparency for order lifecycle monitoring - Harder to debug invalidation-related issues

**Recommended Mitigation:** Add events for each state change:

```
1 event BitInvalidated(address indexed maker, uint256 bitIndex);
2 event TokenInInvalidated(address indexed maker, bytes32 indexed
    orderHash, address tokenIn);
3 event TokenOutInvalidated(address indexed maker, bytes32 indexed
    orderHash, address tokenOut);
```

---

## Appendix

### Hardening Checklist

1. **Add explicit initialization tracking** for `Balances::dynamicBalancesXD` and `XYCConcentrate`
2. **Add runtime bounds checks** for:
  - `feeBps < BPS` for ExactOut paths
  - `feeBps <= BPS` for all fee instructions
  - `decayPeriod > 0`
  - `initialLiquidity > 0`
3. **Add program encoding validation** in `VM::runLoop`:
  - Instruction header bounds
  - Args bounds
  - Opcode bounds
4. **Cap returndata handling** for dynamic fee providers
5. **Use `Math.mulDiv`** for overflow-safe calculations in:
  - `Controls::_onlyTakerTokenSupplyShareGte`
  - `XYCConcentrate::concentratedBalance`
6. **Add fuzz tests** for:

- Fee extremes (0, near BPS, BPS)
- Malformed programs (odd lengths, invalid opcodes)
- Zero-amount swaps with `allowZeroAmountIn`
- Multi-swap sequences with Decay
- Concentrate liquidity scaling over multiple trades

## References

- Solodit Security Checklist
- OpenZeppelin Math Library