

Preface to bloom filter

Q: Compare two sets

How can we create a function used to judge list A ?= list B ? List A and B are made up of unique elements.

1. Using sort comparison:

Copy A and B to arrays/vectors

Sort them

compare one by one:

Time complexity: $O(n \log n)$, Space complexity: $O(n)$

2. Hashing

Create a hash set S for A

For each element in B check whether they are in S or not

Time complexity: $O(n)$, Space complexity: $O(n)$

3. methods with probability to be wrong:

Compare sum $\sum_{i,i \in A} i \stackrel{?}{=} \sum_{i,i \in B}$,

Compute product $\prod_{i,i \in A} i \stackrel{?}{=} \prod_{i,i \in B}$

We call the reduced sum / product as "**sketch**" of set A and B.

A wish to use the power of probability

Can we define a randomized `isEqual(A, B)` which satisfies:

$$\begin{cases} P\{isEqual(A, B)\} = 1 & \text{if } A = B \\ P\{isEqual(A, B)\} \rightarrow 0 & \text{if } A \neq B \end{cases}$$

A naive solution

Suppose we draw a hash function $h : u \rightarrow \{0, \dots, 2^d - 1\}$, $h \in \mathcal{H}$, and \mathcal{H} is a ROM. Define the following sketch for A and B:

$$\begin{aligned} S_A &= \sum h(a) \\ S_B &= \sum h(b) \\ isEqual(A, B) &= S_A \stackrel{?}{=} S_B \end{aligned}$$

If length of sets: $|A| \neq |B|$, then they are easily differentiable.

Otherwise:

1. If A is a permutation of B -> `isEqual(A, B) = true`
2. If A is not a permutation of B:

suppose $a^* \in A \setminus B$ (set difference), then:

$$\begin{aligned}
& P\{\sum h(a) = \sum h(b)\} \\
&= P\{\underbrace{h(a^*)}_L = \underbrace{\sum_{b \in B} h(b) - \sum_{a \in A \setminus a^*} h(a)}_R\} \\
&= 1/2^d
\end{aligned}$$

because:

1. R is a random variable, independent of $h(a^*)$
2. L and R are independent, supported by the random oracle model
3. L is uniformly distributed

Summary

This question introduces us to bloom filters.