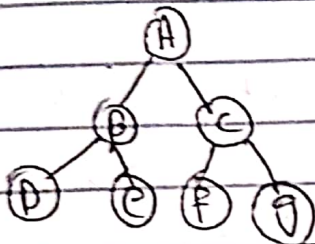


Heap

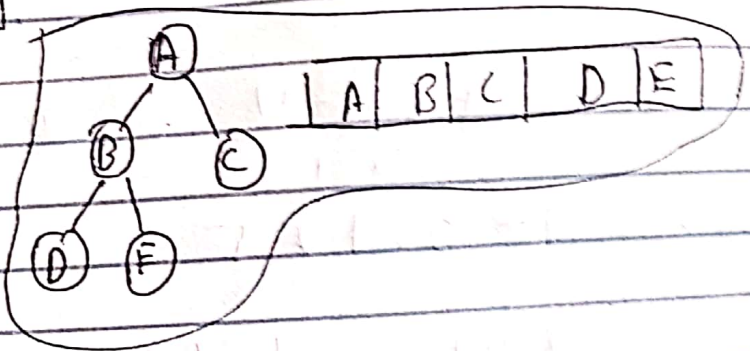
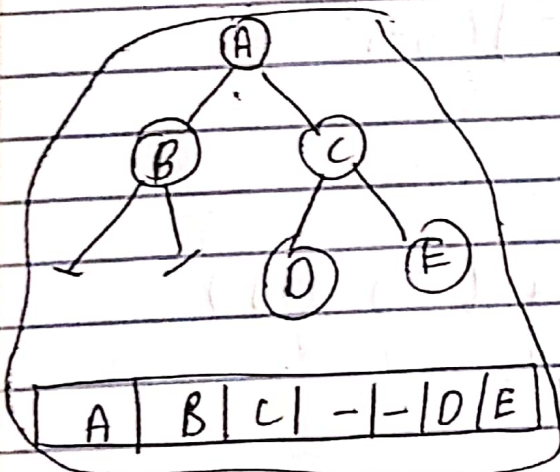
- 1 - Array Representation of BT
- 2 - Complete Binary Tree
- 3 - Heap
- 4 \rightarrow Insert and Delete
- 5 \rightarrow Heap Sort
- 6 \rightarrow Heapify
- 7 \rightarrow Priority Queue

(1)



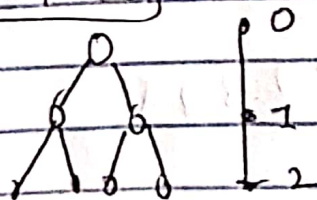
If Node at Index $\rightarrow i$
 Left child $\rightarrow 2 \times i$
 Right child $\rightarrow 2i + 1$
 its parent is $\rightarrow \lfloor \frac{i}{2} \rfloor$

T =	A	B	C	D	E	F	G
	1	2	3	4	5	6	7



Full Binary Tree

Complete



Full Binary Tree are
 $h=2$ Fullfill with nodes
 when we add more
 node then height
 will increase...

Number of Total Node

on Full Binary Tree $= 2^{h+1} - 1$

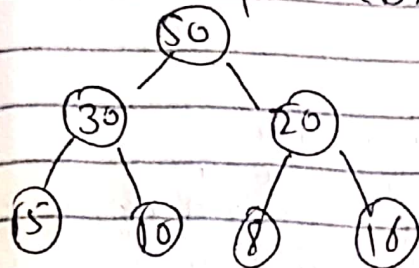
Note = In Array Representation

\rightarrow If any element is missing inbetween
 Array then it's not full
 Binary Tree

the element

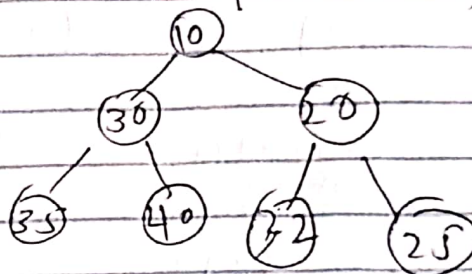
Complete Binary Tree is a full Binary Tree up to height H minus -1

Max Heap (also BT)



$H = [50, 30, 20, 15, 10, 8, 16]$

Min Heap (Also BT)

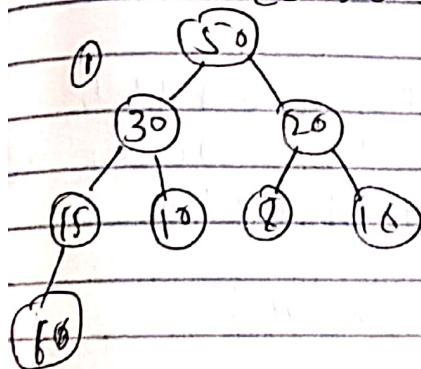


How Insertion

Let us insert (60)

First

$H = [50, 30, 20, 15, 10, 8, 16, 60]$



$60 > 50 \rightarrow$ left side

$60 > 30 \rightarrow$ left side

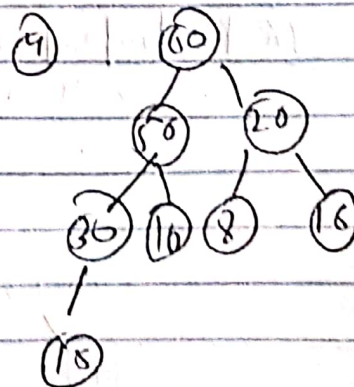
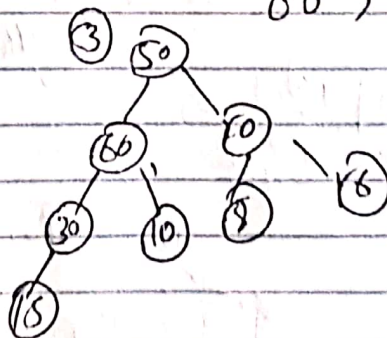
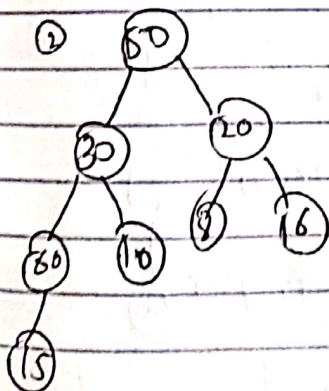
$60 > 15 \rightarrow$ left side

then going on operation with parent A 60

$60 > 15 \rightarrow$ go parent

$60 > 30 \rightarrow$ go parent

$60 > 50 \rightarrow$ go parent



Now its time to Analysis

↳ Depend height

$$\text{height} = O(\log n)$$

$$O(1) \text{ to } O(\log n)$$

Minimum time
if no
swaps occur

For example

Invert (6)

↳ Maximum time

when maximum
swap occurs here

Number of swap

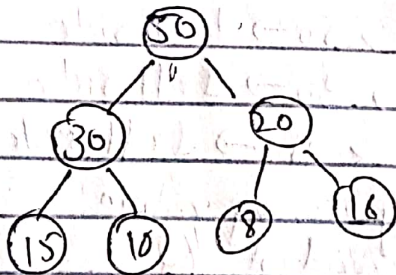
equal to the height

of heap which is

$\log(n)$ which takes

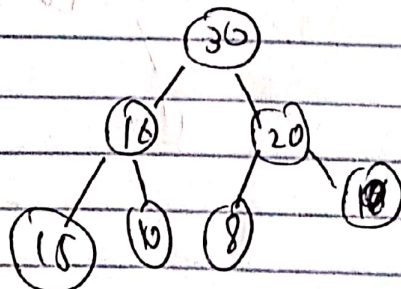
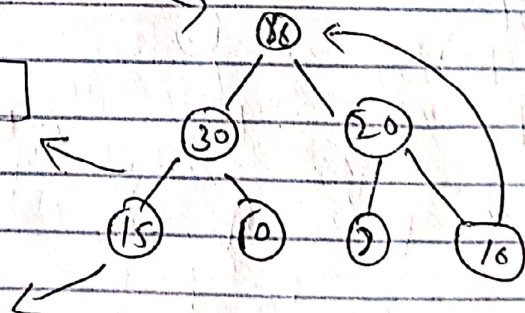
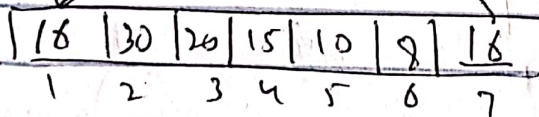
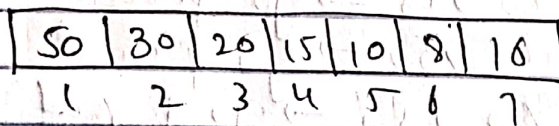
time $\log(n)$

Deletion



Delete 50

↳ Actually Always we delete
the root element because
it's main function



→ deleted root so
still have BST
property

Now its time to Analysis

Swap \rightarrow depend to number node \rightarrow Height

Swap \rightarrow Upward & downward

In Insertion the Swap was moving upward----

Note

When you delete root then get the set then positing but his Actual position save the deleted value when all

the BST delete then you will get Sorted Array but its not BST

"How of Heap sort"

Maximum Time $\log n$

Heap Sort

(10)

(10)
(20)

\rightarrow Property

(20)
(10)

10	20	15	30	40
1	2	3	4	5

10	20		
----	----	--	--

20	10	15
----	----	----

(20)
(10) (15)

20	10	15	30
----	----	----	----

(20)
(10) (15)
(30)

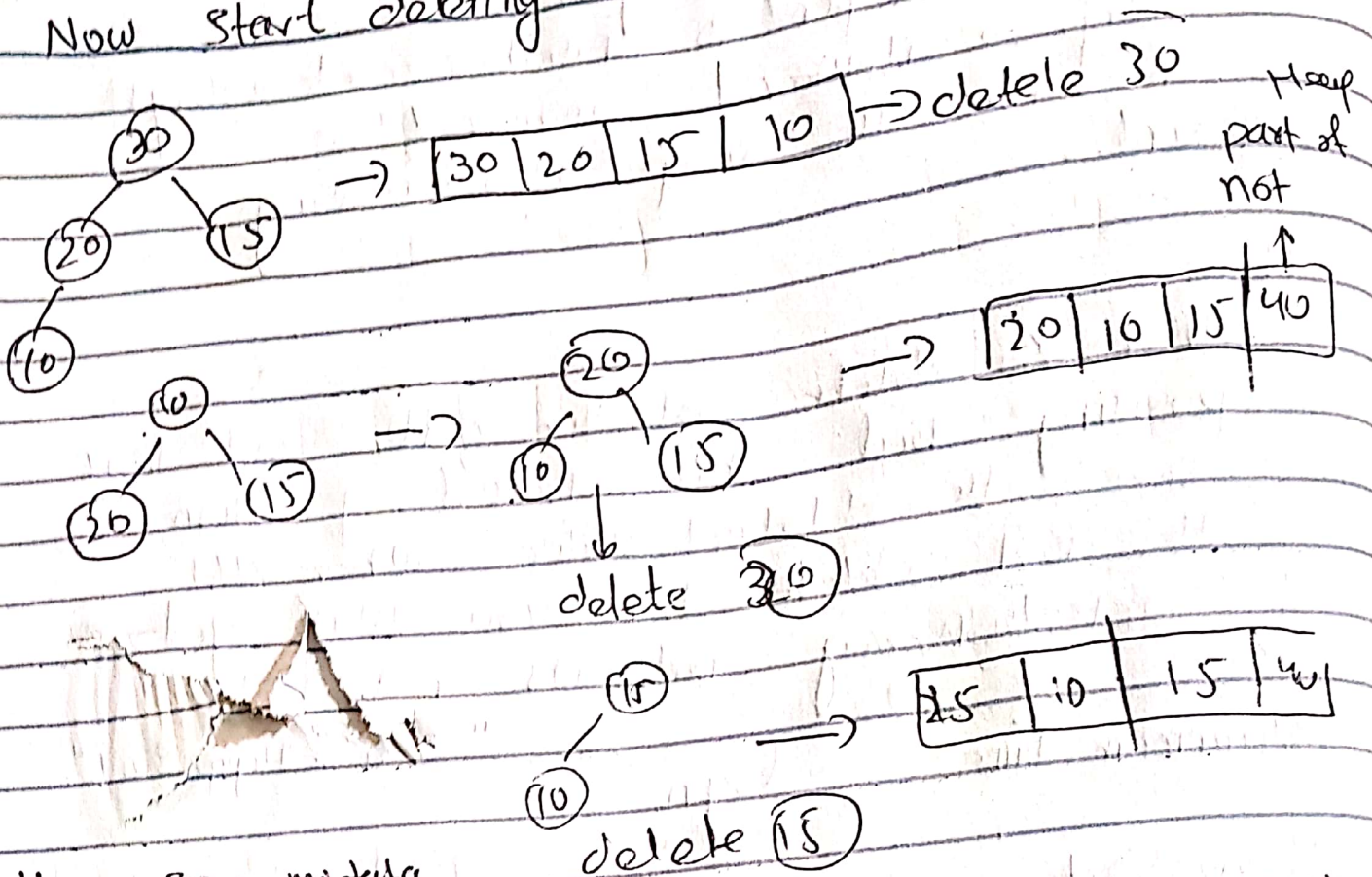
(20)
(30) (15)
(10)

(30)
(20) (15)
(10)

30	20	15	10
----	----	----	----

(30)
(10) (15)
(16) (20)

Now start deleting



Have some mistake
do it when you
again do it

Now this become
we get sorted Array

Delete time = delete one Node time $\times \log n$
 = delete N Node time $\times n(\log n)$
 = Creating the Heap time $\times n \log n$
 deleting the Heap time $\times n \log n$
 Now we get it sorted Array
 time $\times n \log n + n \log n$
 , $\times n \log n$
 time $\times O(n \log n)$

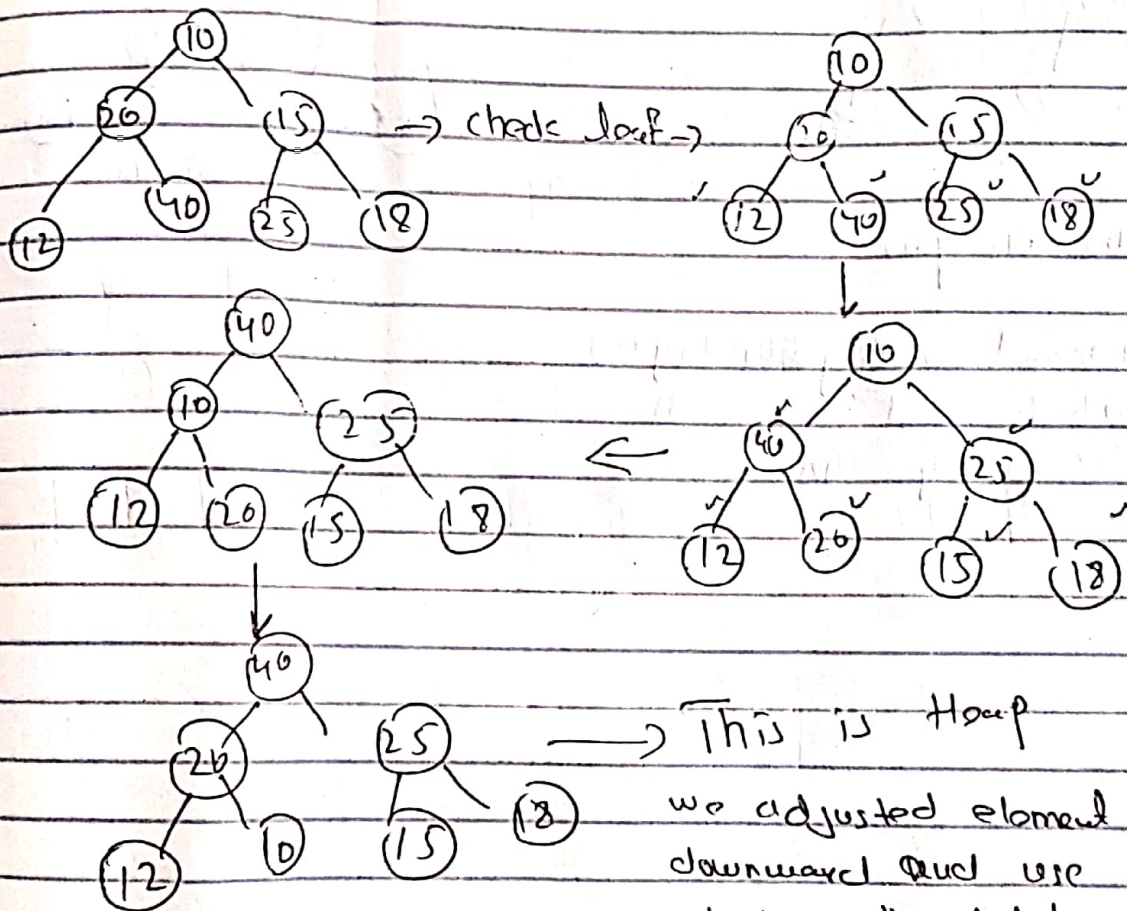
Heapify → Heap property

• Max Heap

For any Given Node:

If P is parent & Node C :
 then $\text{key}(\text{the value})$ of P is greater
 than or equal to key of C

Creating Heap = upward → from leaf
 delete Heap = downward → from root



→ This is Heap
 we adjusted element
 downward and use
 strategy to deleting
 the Node.

Time taken Heapify = $O(n)$

(create and deleting = $n(\log n)$)

$O(n) \neq (n \log n)$

↳ more useful

Priority Queue

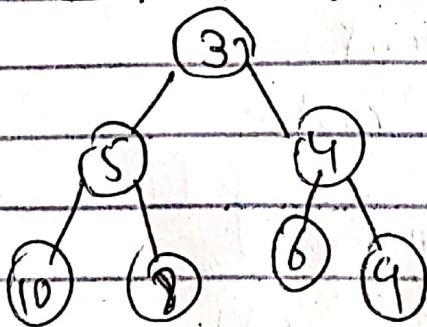
Heap Priority Queue

Small Number High priority

Large Number High priority

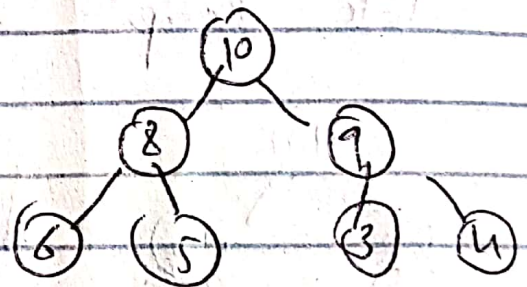
(A)

8 | 6 | 3 | 10 | 5 | 4 | 9



B 8 | 6 | 3 | 10 | 5 | 4 | 9

Max Heap



Time Analysis

delete and creating $O(n \log n)$
which best for the
Priority Queue