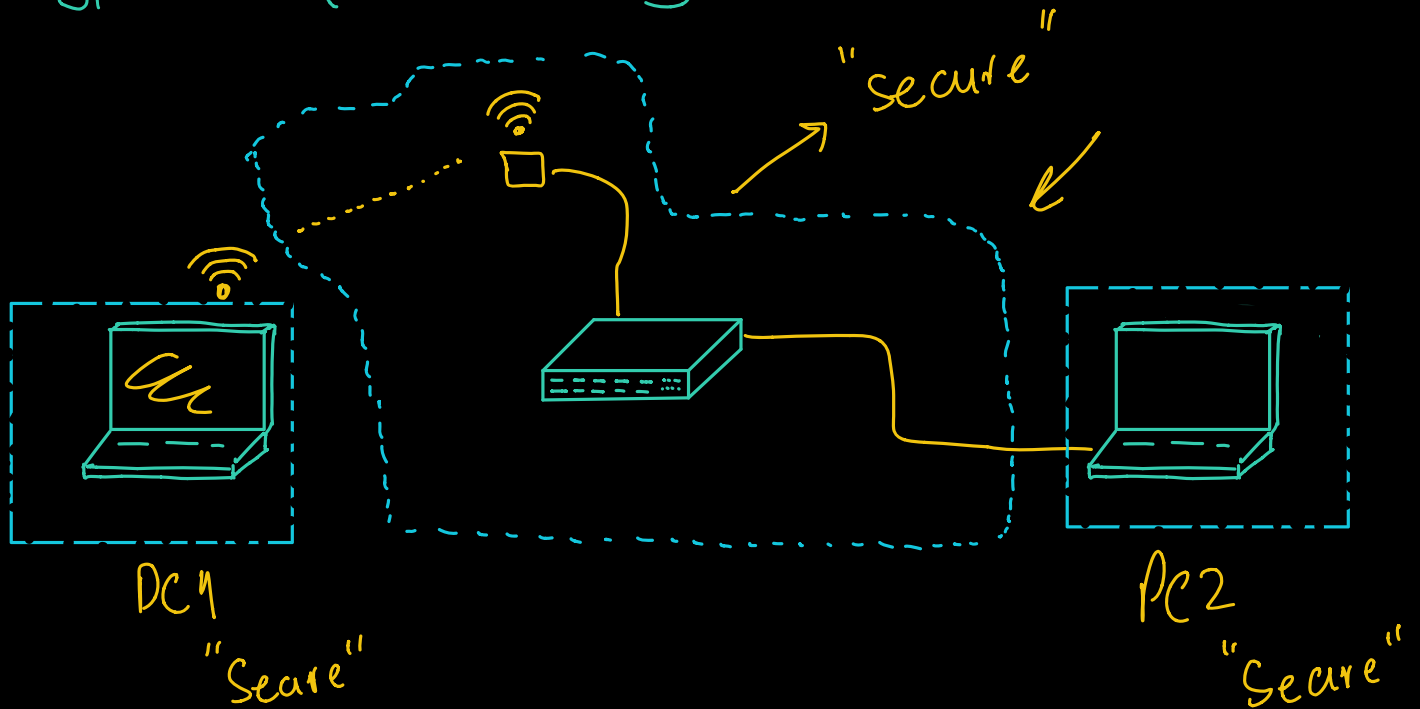


Operating Systems Design

19. Protection

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Types of security:



— Network security ✓

— Systems security

— "Authorization" ✓

— Malware
"stuff" ||

Protection & Security

- Security Policy
 - Prevention of unauthorized access to a system
 - Malicious or accidental access →
 - “access” may be:
 - user login, a process accessing things it shouldn't, physical access
 - The access operations may be reading, destruction, or alteration
- Protection Mechanism
 - The mechanism that provides and enforces controlled access of resources to processes
 - A protection mechanism enforces security policies

Principle of Least Privilege

- At each abstraction layer, every element (user, process, function) should be able to access only the resources necessary to perform its task
- Even if an element is compromised, the scope of damage is limited



Security Goals

" كون "

- **Authentication**
 - Ensure that users, machines, programs, and resources are properly identified
- **Confidentiality** ✓
 - Prevent unauthorized access to data
- **Integrity** ✓
 - Verify that data has not been compromised: deleted, modified, added
- **Availability** ✓
 - Ensure that the system is accessible

The Operating System

Subject

Object

- The OS provides processes with access to resources

Action
Right

Resource <i>objects</i>	OS component <i>(Subjects)</i>
Processor(s) ✓	Process scheduler
Memory ✓	Memory Management + MMU
Peripheral devices ✓	Device drivers & buffer cache
Logical persistent data	File systems
Communication networks	Sockets

S
O
R

- Resource access attempts go through the OS
- OS decides whether access should be granted
 - Rules that guide the decision = policy

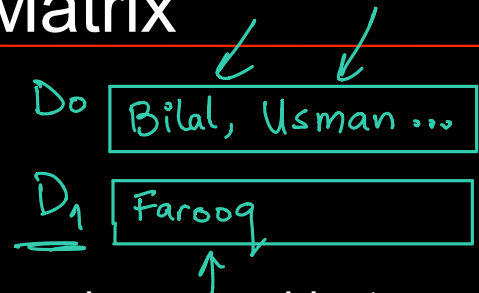
App ^{sysd}
↓
OS

Domains of protection

- subjects Processes interact with objects
 - Objects:
 - hardware (CPU, memory, I/O devices)
 - software: files, semaphores, messages, signals
- A process should be allowed to access only objects that it is authorized to access
 - A process operates in a **protection domain** D_0
 - Protection domain defines the objects the process may access and how it may access them

Modeling Protection: Access Matrix

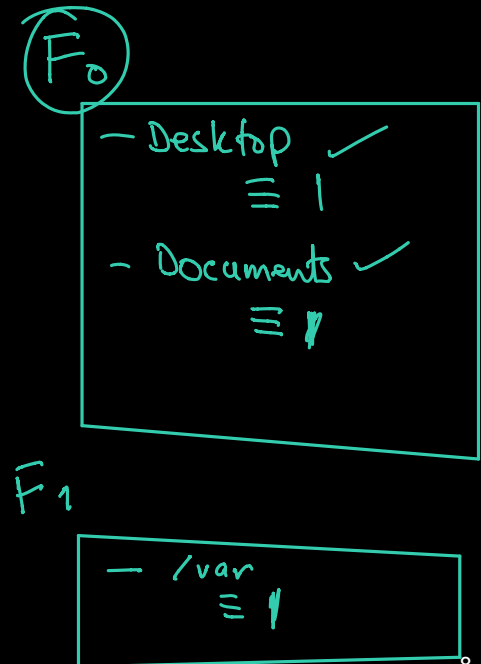
- Rows: domains
- Columns: objects
- Each entry represents an access right of a domain on an object



objects

	F_0	F_1	Printer
D_0	read	read-write	print
D_1	read-write-execute	read	
D_2	read-execute		
D_3		read	print
D_4			print

domains of protection



Access Matrix: Domain Transfers

- Switching from one domain to another is a configurable policy

$D_0 \rightarrow \text{HoD}$

$D_1 \rightarrow \text{Instructor}$

objects

domains of protection

	F_0	F_1	Printer	D_0	D_1	D_2	D_3	D_4
D_0	read	read-write	print	—	switch	switch		
D_1	read-write-execute	read			—			
D_2	read-execute				switch	—		
D_3		read	print					
D_4			print					

Access Matrix: Additional operations

- **Owner**: allow new rights to be added or removed
 - ✗ An object may be identified as being *owned* by the domain
 - ✗ Owner can add and remove any right in any column of the object

objects

domains of protection

	F_0	F_1
D_0	read owner	

Implementing an access matrix

- A single table is usually impractical
 - Big size: # domains (users) × # objects (files)
 - Objects may come and go frequently
- Access Control List
 - Associate a column of the table with each object

Attend Evaluation
|
Sparse

Implementing an access matrix

- Access Control List (ACL) *Default Deny*
 - Associate a column of the table with each object

domains of protection

	<i>objects</i>	<i>F</i> ₀	<i>F</i> ₁	Printer	<i>D</i> ₀	<i>D</i> ₁	<i>D</i> ₂	<i>D</i> ₃	<i>D</i> ₄
<i>D</i> ₀	read owner	read-write	print						
<i>D</i> ₁	read-write-execute	read*				–			
<i>D</i> ₂	read-execute					switch	–		
<i>D</i> ₃		read	print						
<i>D</i> ₄			print						

ACL for file *F*₀

F_0



(Printer)

$= \{ \begin{array}{l} \text{'D}_0\text{' } : [\text{owner, read}], \\ \text{'D}_1\text{' } : [\text{read, write, exec}], \\ \vdots \\ \end{array} \}$

$F_0 = \{ \}$

100
100
100,000

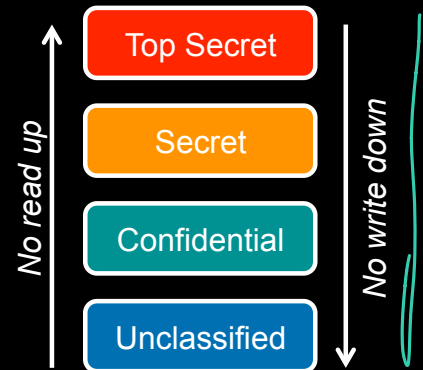
Access Control Models: MAC vs. DAC

- **DAC: Discretionary Access Control** 
 - A subject (domain) can pass information onto any other subject
 - In some cases, access rights may be transferred
 - *Most systems use this* (er... not really... only OSs)
- **MAC: Mandatory Access Control** 
 - Policy is centrally controlled
 - Users cannot override the policy

Multi-level Access Control

(skip)

- Typical MAC implementations use a **Multi-Level Secure (MLS)** access model
- **Bell-LaPadula** model
 - Identifies the ability to access and communicate data
 - Objects are classified into a hierarchy of sensitivity levels
 - Unclassified, Confidential, Secret, Top Secret
 - Users are assigned a clearance
 - “**No read up; no write down**”
 - Cannot read from a higher clearance level
 - Cannot write to a lower clearance level
- Works well for government information
- Does not translate well to civilian life

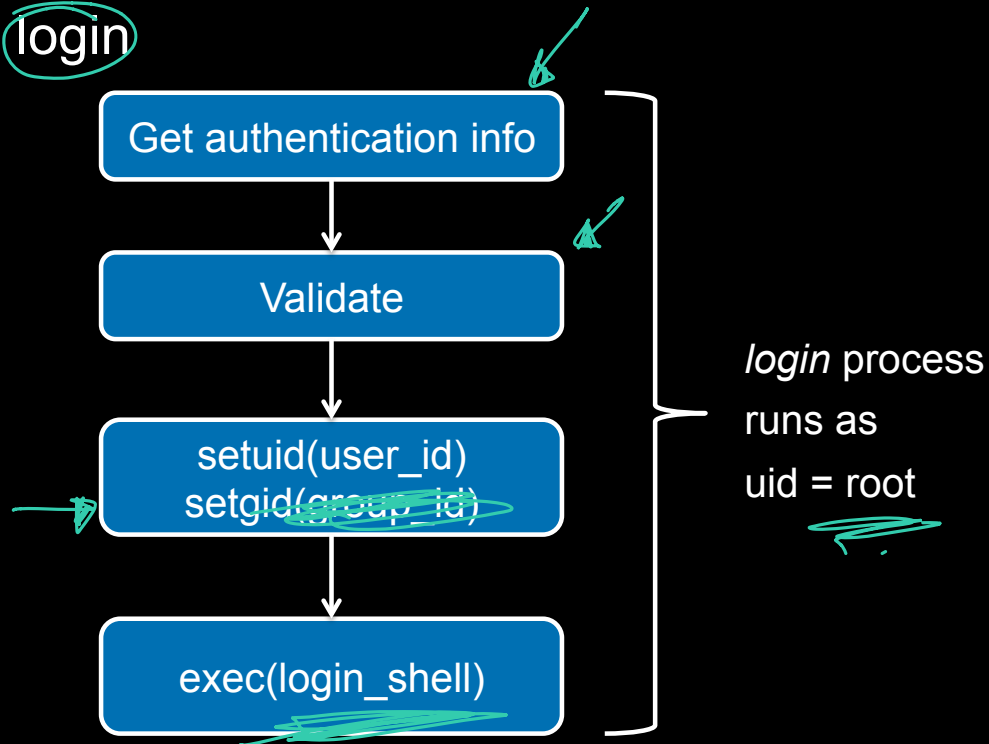


*Confidential cannot read Secret
Confidential cannot write Unclassified*

Authentication

Authentication

- Establish & verify identity
 - Then decide whether to allow access to resources
- Example: login



Password Authentication Protocol (PAP)

- Reusable passwords
- (Server keeps a database of *username:password* mappings) *Shalat*
- Prompt client/user for a login name & password
- To authenticate, use the login name as a key to look up the corresponding password in a database (file) to authenticate

match

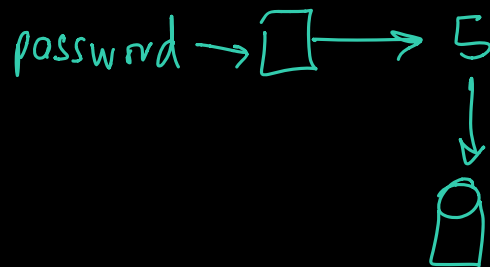
```
if (supplied_password == retrieved_password)
    then user is authenticated
```

$$\text{hash}(\text{hash}(\overrightarrow{sp}) + L) \stackrel{\text{hash}}{=} (\text{h}(\overrightarrow{rp}) + L)$$

PAP: Reusable passwords

password breach

One problem: what if the password file isn't sufficiently protected and an intruder gets hold of it, he gets all the passwords!



Enhancement:

Store a hash of the password in a file

- given a file, you don't get the passwords
- have to resort to a dictionary or brute-force attack
- Unix approach
 - Password encrypted with 3DES hashes; then MD5 hashes; now SHA512 hashes
 - Salt used to guard against **dictionary attacks**

"MySecurePassword123#"

Adversary
malicious!

Passwords file

nam: MySecure...
...
...
...

1

Solved!

HASH

SHA-256

nam: AE9CD729...

MySecu

AE9CD

3

Solved!

Dictionary
attack

(A): A9c29...
(Apple): C29D7...
...
MySecu...: AE9C...

Solved!

hash (AE9CD729)

(1296929) → Salt

nam: 192AC9DE...

?!?
hash (C29D7 + 1296)

2

4

Hashes

$$9 \bmod 9 \equiv 0$$

$$5 \bmod 9 \equiv 5$$

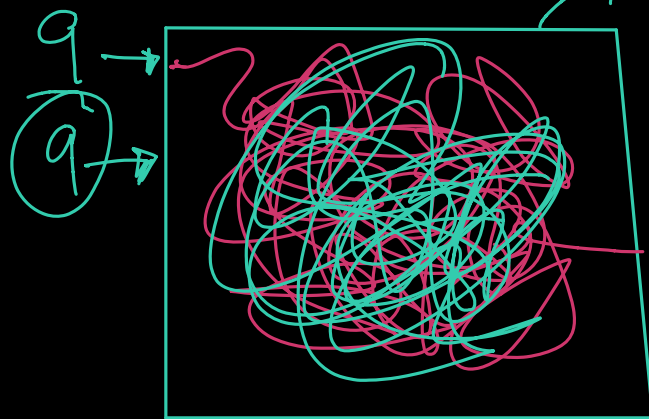
$$14 \bmod 9 \equiv 5$$

A Hash function

— It's bad because it has "collisions"...

— And it's reversible

Given 5, you can predict what the original number was



Secure hash function

SHA-1
SHA-256
SHA-512

A9C29DE47

20 bytes

— Same output for same input
↪ Irreversible!

*

NEVER

EVER

EVER

STORE

PASSWORDS

IN

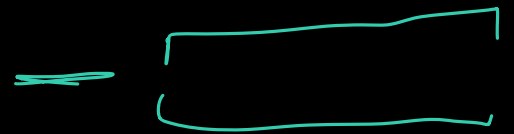
PLAINTEXT!

— EVER!



INSERT _____

VALUE (Username,
password (pass));



Username	password
nam	ABCDG

Authentication

Three factors:

– something you have

- can be stolen

key, card, phone!

– something you know

- can be guessed, shared, stolen

passwords

– something you are

- costly, can be copied (sometimes)

biometrics

Authentication

factors may be combined

– ATM machine: 2-factor authentication

- **ATM card** *something you have*
- **PIN** *something you know*

Versus Authorization

Authorization defines access control

Usman — policy

Once we know a user's identity:

- Allow/disallow request
- Operating system enforces system access based on user's credentials
 - Network services usually run in another context
 - Network server may not know of the user
 - Application takes responsibility
- May contact an authorization server
 - Trusted third party that will grant credentials
 - Kerberos ticket granting service , LDAP, ActiveDirectory
 - RADIUS (centralized authentication/authorization)

Three (Four?) A's of Security

- Authentication ✓
 - Validate an identity or a message

- Authorization (Access Control) ✓
 - Enforce policy

- Accounting ✓

logging

- Auditing

source

Accounting

If security has been compromised

... *what happened?*

... *who did it?*

... *how did they do it?*

Log transactions

- Logins
- Commands
- Database operations
- *Who looks at audits?*

Log to remote systems

- Minimize chances for intruders to delete logs

Auditing

Go through software source code and search for security holes

- Need access to source
 - Some operating systems > 50 million lines!
- Experienced staff + time
- E.g., OpenBSD

Complex systems will have more bugs

- And will be harder to audit

The End