Information Theory, Fall 2015	Iftach Haitner
Problem set	t 4
December 27 2015	Due: December 22

- Please submit the handout in class, or email the grader (omer.rotem1 at gmail.com ).
- Write clearly and shortly using sub-claims if needed. The emphasize in most questions is on the proofs (no much point is writing a "solution" w/o proving its correctness)
- For Latex users, a solution example can be found in the course web site.
- It is allowed to work in (small) groups, but please write the id list of your partners in the solution file, and each student should write his solution by *himself* (joint effort is only allowed in the "thinking phase")

- 1. Show that  $SD(p,q) = \max_{S \subseteq [m]} \left( \sum_{i \in S} p_i \sum_{i \in S} q_i \right)$  for any two distributions p,q over [m].
- 2. Use the above to prove that  $SD(p,q) = \max_{D} (\Pr_{X \sim p} [D(X) = 1] \Pr_{X \sim q} [D(X) = 1])$ , where the max is over all deterministic distinguishers. Try and extend the above to randomized distinguishers?
- 3. Relative entropy is not symmetric: give two distributions p,q such that  $D(p||q) \neq D(q||p)$ , and  $D(p||q), D(q||p) < \infty$ .
- 4. Relative entropy does not obey the triangle inequality: give three distributions  $p_1, p_2, p_3$  such that  $D(p_1||p_2) + D(p_2||p_3) < D(p_1||p_3)$
- 5. Relative entropy is non-negative: given two distributions p, q, show that  $D(p||q) \ge 0$ , with equality only if p = q.
- 6. Does Theorem 4 in Lecture 7 hold for any prefix code C with  $E_{i \leftarrow q}[|C(i)|] \leq H(q) + 1$ ? (and not only for a code C with  $C(x) = \left\lceil \log \frac{1}{q(x)} \right\rceil$ , as stated in the theorem)
- 7. Prove the data processing inequality of relative entropy (Claim 7 in Lecture 7) for randomized functions.