Application of Information Theory, Lecture 10

Hardcore Predicates

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Part I

Motivation and Definition

Hardcore predicates

- Let $f\colon \{0,1\}^n \mapsto \{0,1\}^n$ be a “hard to invert" function, how unpredictable is x given $f(x)$

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- Parts of x might be (totally) predictable
- It turns out that there is an hardcore part in x.

Hardcore predicates, cont.

Definition 1 (hardcore predicates)

A predicate $b\colon \{0,1\}^n \mapsto \{0,1\}$ is (s,ε)-hardcore predicate of $f\colon \{0,1\}^n \mapsto \{0,1\}^n$, if $\Pr_{x\leftarrow\{0,1\}^n}[\mathsf{P}(f(x)) = b(x)] \leq \frac{1}{2} + \varepsilon$, for any s-size P.

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- Does the existence of hardcore predicate for f implies that f is hard to invert?

Part II

The Information Theoretic Settings

Some definitions

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- In both examples $\mathrm{H}_\infty(Z) = k$

2-universal families

Definition 2 (2-universal families)

A function family $\mathcal{G} = \{g \colon \mathcal{D} \mapsto \mathcal{R}\}$ is 2-universal, if $\forall\ x \neq x' \in \mathcal{D}$ it holds that $\Pr_{g \leftarrow \mathcal{G}}\left[g(x) = g(x')\right] = \frac{1}{|\mathcal{R}|}$.

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Example: $\mathcal{D} = \{0,1\}^n$, $\mathcal{R} = \{0,1\}^m$ and $\mathcal{G} = \{A \in \{0,1\}^{m \times n}\}$ with $A(x) = A \times x \bmod 2$.

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Lemma 3 (leftover hash lemma)

Let X be a rv over $\{0,1\}^n$ with $\mathrm{H}_2(X) \geq k$ let $\mathcal{G} = \{g \colon \{0,1\}^n \mapsto \{0,1\}^m\}$ be 2-universal and let $G \leftarrow \mathcal{G}$. Then $\mathrm{SD}((G, G(X)), (G, \sim \{0,1\}^m)) \leq \frac{1}{2} \cdot 2^{(m-k)/2}$.

Hardcore predicate for regular functions

Lemma 4

Let $f\colon \{0,1\}^n \mapsto \{0,1\}^n$ *be* 2^k*-regular function, let* $\mathcal{G} = \{g\colon \{0,1\}^n \mapsto \{0,1\}\}$ *be* 2*-universal and let* $v\colon \{0,1\}^n \times \mathcal{G}_n \mapsto \{0,1\}^n \times \mathcal{G}_n$ *be defined by* $v(x,h) = (f(x), g)$*. Then* $b(x,g) = g(x)$ *is* $(\infty, 2^{-(k-1)/2})$ *hardcore-predicated of* g*.*

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- b is an hardcore predicate of v (not of f)

Proving Lemma 4

Claim 5

$\mathrm{SD}\left((f(X), G, G(X)), (f(X), G, U)\right) \leq 2^{-(k-1)/2}$,
for $G \leftarrow \mathcal{G}$, $X \leftarrow \{0,1\}^n$ and $U \leftarrow \{0,1\}$.

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Lemma 6 (predicting to distinguishing)

Let Y, Z be rvs over $\{0,1\}^* \times \{0,1\}$ and let P be an algorithm with $\Pr\left[\mathsf{P}(Y) = Z\right] \geq \frac{1}{2} + \varepsilon$. Then $\exists$ algorithm D, with essentially the same complexity as P, with $\Pr\left[\mathsf{D}(Y, Z) = 1\right] - \Pr\left[\mathsf{D}(Y, U) = 1\right] \geq \varepsilon$.

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Corollary 7

If $\mathrm{SD}((Y, Z), (Y, U)) < \varepsilon$, then $\Pr\left[\mathsf{P}(Y) = Z\right] < \frac{1}{2} + \varepsilon$ for any predictor P.

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$$\mathrm{SD}((f(X), G, G(X)), (f(X), G, U))$$
$$= \sum_{y \in \mathrm{Im}(f)} \Pr[f(X) = y] \cdot \mathrm{SD}\big((y, G, G(X)|_{f(X)=y}), (y, G, U)\big) \quad \text{(board)}$$

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\end{aligned}$$

Since $\mathrm{H}_\infty(X_y) = k$ for every $y \in Im(f)$, the leftover hash lemma yields that

$$\begin{aligned}
\mathrm{SD}((G, G(X_y)), (G, U)) &\le \frac{1}{2} \cdot 2^{(1 - \mathrm{H}_\infty(X_y)))} \\
&= 2^{(-k-1)/2}. \square
\end{aligned}$$

Part III

The Computational Settings

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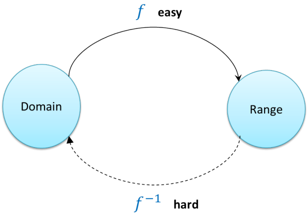

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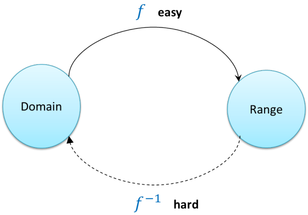

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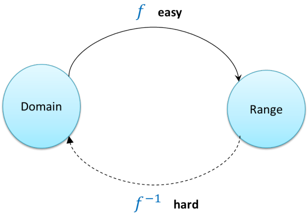

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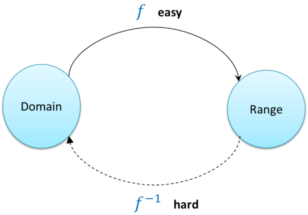

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- Sufficient for many cryptographic primitives

One-way functions, cont.

Definition 8 (one-way functions (OWFs))

A poly-time $f\colon \{0,1\}^n \mapsto \{0,1\}^n$ is (s,ε)-one-way, if $\Pr_{x \leftarrow \{0,1\}^n}\left[\mathsf{Inv}(f(x)) \in f^{-1}(f(x))\right] = \varepsilon$ for any s-size Inv.

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- Such hardcore predicates have many cryptographic applications

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- Inv can "flip" coins
- f is one-way $\implies$ predicting x from $f(x)$ is hard.
- But does any one-way function has an hardcore predicate?
- Such hardcore predicates have many cryptographic applications
- f is injective and not one-way $\implies$ f has no hardcore predicate.

Direct product predicate

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For $f\colon \{0,1\}^n \mapsto \{0,1\}^n$, *define* $g\colon \{0,1\}^n \times [n] \mapsto \{0,1\}^n \times [n]$ *by* $g(x,i) = (f(x),i)$. *Assuming* f *is* $(s,\frac{1}{2})$*-one way, then* $\Pr_{x \leftarrow \{0,1\}^n, i \leftarrow [n]}[\mathsf{P}(f(x),i) = x_i] \leq 1 - 1/2n$ *for any* $\frac{s}{n}$*-size* P.

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3. Construction is “inefficient"

The Goldreich-Levin predicate

For $x, r \in \{0,1\}^n$, let $\langle x, r\rangle_2 := (\sum_{i=1}^{n} x_i \cdot r_i) \bmod 2 = \bigoplus_{i=1}^{n} x_i \cdot r_i$.

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Theorem 10 (Goldreich-Levin)

For $f\colon \{0,1\}^n \mapsto \{0,1\}^n$, define $g\colon \{0,1\}^n \times \{0,1\}^n \mapsto \{0,1\}^n \times \{0,1\}^n$ by $g(x,r) = (f(x), r)$. Assume f is (s,ε)-one-way, then $b(x,r) := \langle x, r\rangle_2$ is an $(\sqrt[3]{n\varepsilon}, \frac{\varepsilon}{n^2} \cdot s)$-hardcore predicate of g.

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- We prove $\exists\ \frac{n^2}{\delta^2}$-size Inv with $\Pr[\mathsf{Inv}(f(X))=X] \in \Omega(\delta^3/n)$.

Focusing on a good set

Claim 11

There exists set $\mathcal{S} \subseteq \{0,1\}^n$ with

1. $\frac{|\mathcal{S}|}{2^n} \geq \frac{\delta}{2}$, and
2. $\Pr\left[\mathsf{P}(f(x), R) = b(x, R)\right] \geq \frac{1}{2} + \frac{\delta}{2}, \qquad \forall x \in \mathcal{S}$.

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We conclude the theorem's proof showing that there exists a $\frac{n^2}{\delta^2}$-size Inv with

$$\Pr\left[\mathsf{Inv}(f(x)) = x\right] \in \Omega(\delta^2/n)$$

for every $x \in \mathcal{S}$.

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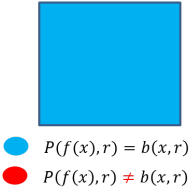
for every $x \in \mathcal{S}$. In the following we fix $x \in \mathcal{S}$.

The perfect case

$$\Pr\left[\mathsf{P}(f(x), R) = b(x, R)\right] = 1$$



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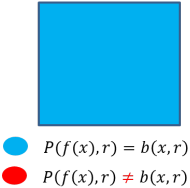

In particular, $\mathsf{P}(f(x), e^i) = b(x, e^i)$ for every $i \in [n]$, for $e^i = (\underbrace{0, \ldots, 0}_{i-1}, 1, \underbrace{0, \ldots, 0}_{n-i})$.

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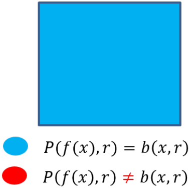

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Hence, $x_i = \langle x, e^i \rangle_2$

The perfect case

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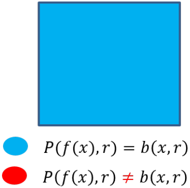

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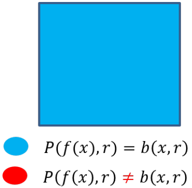

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Algorithm 12 (Inverter Inv on input $y \in \operatorname{Im}(f)$)

Return $(\mathsf{P}(y, e^1), \ldots, \mathsf{P}(y, e^n))$.

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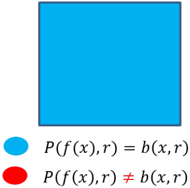

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$\mathsf{Inv}(f(x)) = x$.

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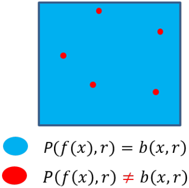
$$\Pr\left[\mathrm{P}(f(x), R) = b(x, R)\right] \geq 1 - \frac{1}{4n}$$



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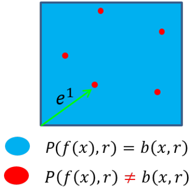

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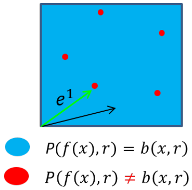

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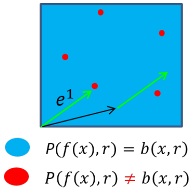

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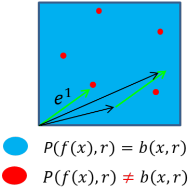

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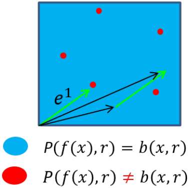

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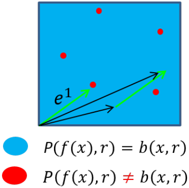

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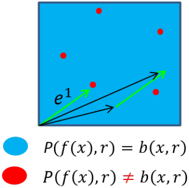

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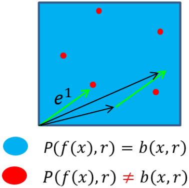

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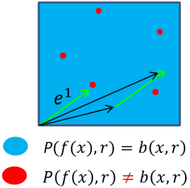

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Algorithm 14 (Inverter Inv on input y)

Return $(\mathrm{P}(y, R) \oplus \mathrm{P}(y, R \oplus e^1)), \ldots, \mathrm{P}(y, R) \oplus \mathrm{P}(y, R \oplus e^n))$.

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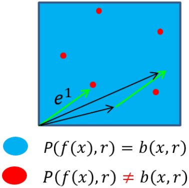

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$$
\begin{aligned}
b(x,y) \oplus b(x,w) &= \left(\bigoplus_{i=1}^{n} x_i \cdot y_i\right) \oplus \left(\bigoplus_{i=1}^{n} x_i \cdot w_i\right) \\
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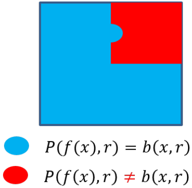
Intermediate Case

$\Pr\left[\mathrm{P}(f(x), R) = b(x, R)\right] \geq \frac{3}{4} + \delta$



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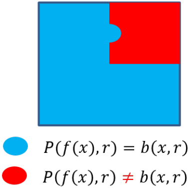

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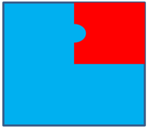

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Algorithm 15 ($\mathsf{Inv}(y)$)

For every $i \in [n]$:

1. Sample $r^1, \ldots, r^v \in \{0, 1\}^n$ uniformly at random
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Output $(m_1, \ldots, m_n)$

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The following claim holds for “large enough" v.

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Let $X^1, \ldots, X^v$ be iids over $[0,1]$ with expectation μ. Then,
$\Pr\left[\left|\frac{\sum_{j=i}^{v} X^j}{v} - \mu\right| \geq \alpha\right] \leq 2 \cdot \exp(-2\alpha^2 v)$ *for every $\alpha > 0$.*

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- Hence, we complete the proof taking $X^j = W^j$, $\alpha = 2\delta$ and $v = \left\lceil \log(n) \cdot \frac{1}{2\alpha^2} \right\rceil + 1$.

The actual (hard) case

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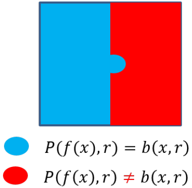

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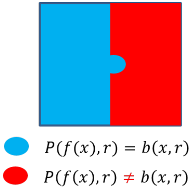

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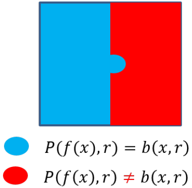

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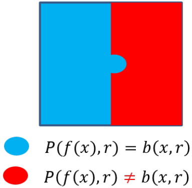

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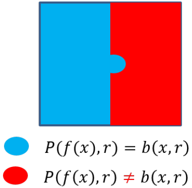

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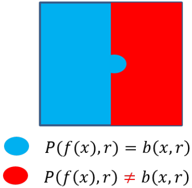

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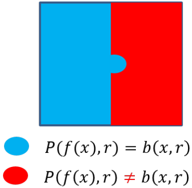

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- Solution: choose the samples in a correlated manner

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- Fix $i \in [n]$, and let $W^{\mathcal{L}}$ be 1 iff $\mathsf{P}(f(x), r^{\mathcal{L}} \oplus e^i) \oplus b(x, r^{\mathcal{L}}) = x_i$.
- We need to lowerbound $\Pr\left[\sum_{\mathcal{L} \subseteq [\ell]} W^{\mathcal{L}} > \frac{v}{2}\right]$

Algorithm Inv

- For $\ell \in \mathbb{N}$ ($\approx \log \frac{n}{\delta}$, to be determined later), let $v = 2^\ell - 1$.
- In the following $\mathcal{L} \subseteq [\ell]$ stands for a non empty subset

Algorithm 18 (Inverter Inv on $y = f(x) \in \{0,1\}^n$)

1. Sample uniformly (and independently) $t^1, \ldots, t^\ell \in \{0,1\}^n$
2. Guess the value of $\{b(x, t^i)\}_{i \in [\ell]}$
3. For all $\mathcal{L} \subseteq [\ell]$: set $r^{\mathcal{L}} = \bigoplus_{i \in \mathcal{L}} t^i$ and compute $b(x, r^{\mathcal{L}}) = \bigoplus_{i \in \mathcal{L}} b(x, t^i)$.
4. For all $i \in [n]$, let $m_i = \mathrm{maj}_{\mathcal{L} \subseteq [\ell]}\{\mathsf{P}(f(x), r^{\mathcal{L}} \oplus e^i) \oplus b(x, r^{\mathcal{L}})\}$
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- We need to lowerbound $\Pr\left[\sum_{\mathcal{L} \subseteq [\ell]} W^{\mathcal{L}} > \frac{v}{2}\right]$
- Problem: the $W^{\mathcal{L}}$'s are dependent!

Analyzing Inv's success probability

1. Let $T^1, \ldots, T^\ell$ be iid and uniform over $\{0,1\}^n$.
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1. $\forall \mathcal{L} \subseteq [\ell]$, $R^{\mathcal{L}}$ is uniformly distributed over $\{0,1\}^n$.
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$$= \sum_{(t^2,\ldots,t^\ell) \in \{0,1\}^{(\ell-1)n}} \Pr[(T^2,\ldots,T^\ell) = (t^2,\ldots,t^\ell)] \cdot \Pr[R^{\mathcal{L}} = w \wedge R^{\mathcal{L}'} = w' \mid (T^2,\ldots,T^\ell) = (t^2,\ldots,t^\ell)]$$

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Pairwise independence variables

Definition 20 (pairwise independent random variables)

A sequence of rv's $X^1, \ldots, X^v$ is pairwise independent, if $\forall i \neq j \in [v]$ and $\forall a, b$, it holds that $\Pr[X^i = a \wedge X^j = b] = \Pr[X^i = a] \cdot \Pr[X^j = b]$.

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Lemma 21 (Chebyshev's inequality)

Let $X^1, \ldots, X^v$ be pairwise-independent random variables with expectation μ and variance σ^2. Then, for every $\alpha > 0$: $\Pr\left[\left|\frac{\sum_{j=1}^{v} X^j}{v} - \mu\right| \geq \alpha\right] \leq \frac{\sigma^2}{\alpha^2 v}$.

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- Recalling that we guaranteed to work well on $\frac{\delta}{2}$ of the x's. We conclude that $\Pr[\mathrm{Inv}(f(x)) = x] \in \Theta(\delta^3/n)$.

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- Hardcore functions:

  Similar ideas allows to output $\log n$ “pseudorandom bits"

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  Let X be a rv with over $\{0,1\}^n$ with $H_\infty(X) \geq t$, and assume $SD((R, \langle R, X\rangle_2), (R, U)) > \alpha = 2^{-c \cdot t}$ for some universal $c > 0$.

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$\Longrightarrow$ (by GL) $\exists$ Inv that guesses X from nothing, with prob $\alpha^{O(1)} > 2^{-t}$

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  - The code we used here is known as the Hadamard code
- LPN - learning parity with noise:

  Given polynomially many samples of the form $(R_i, \langle x, R_i\rangle_2 + \theta)$, for $R_i \leftarrow \{0,1\}^n$ and boolean $\theta_i \sim (\frac{1}{2} - \delta, \frac{1}{2} - \delta)$, find x.

Reflections cont.

- List decoding:
  - Encoder $f\colon \{0,1\}^n \mapsto \{0,1\}^m$ and decoder g, such that for any $x \in \{0,1\}^n$ and c of hamming distance at most $(\frac{1}{2} - \delta)$ from $f(x)$: g examines $\text{poly}(1/\delta)$ symbols of c and outputs a $\text{poly}(1/\delta)$-size list that whp contains x
  - The code we used here is known as the Hadamard code
- LPN - learning parity with noise:

  Given polynomially many samples of the form $(R_i, \langle x, R_i\rangle_2 + \theta)$, for $R_i \leftarrow \{0,1\}^n$ and boolean $\theta_i \sim (\frac{1}{2} - \delta, \frac{1}{2} - \delta)$, find x.
- The difference comparing to Goldreich-Levin — no control over the R's.