Section 1

Commitment Schemes

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An efficient two-stage protocol (S, R) .

Commit The sender S has private input $b \in \{0,1\}^*$ and the common input is 1^n. The commitment stage result in a joint output c, the *commitment*, and a private output d to S, the *decommitment*.

Reveal S sends the pair (d, b) to R, and R either accepts or rejects.

Completeness: R always accepts in an honest execution.

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Hiding:. In commit stage: $\forall$ R^*, $m \in \mathbb{N}$ and $b \neq b' \in \{0,1\}^m$, $\{\mathrm{View}_{\mathrm{R}^*}(\mathrm{S}(b), \mathrm{R}^*)(1^n)\}_{n \in \mathbb{N}} \approx_c \{\mathrm{View}_{\mathrm{R}^*}(\mathrm{S}(b'), \mathrm{R}^*)(1^n)\}_{n \in \mathbb{N}}$.

Commitment Schemes cont.

Binding: “Any" S^* succeeds in the following game with negligible probability in n:

> *On security parameter* 1^n*,* S^* *interacts with* R *in the commit stage resulting in a commitment* c*, and then output two pairs* (d, b) *and* (d', b') *with* $b \neq b'$ *such that* $\mathsf{R}(c, d, b) = \mathsf{R}(c, d', b') = \texttt{Accept}$

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- (non-uniform) OWFs imply statistically binding, and statistically hiding commitments

Perfectly Binding Commitment from OWP

Let $f\colon \{0,1\}^n \mapsto \{0,1\}^n$ be a permutation and let P be a (non-uniform) hardcore predicate for f

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Protocol 2 ((S,R))

Commit:

S's input: $b \in \{0,1\}$

S chooses a random $x \in \{0,1\}^n$, and sends $c = (f(x), P(x) \oplus b)$ to R

Reveal:

S sends (x, b) to R, and R accepts iff (x, b) is consistent with c (i.e., $P(x) \oplus b = c$)

Perfectly Binding Commitment from OWP cont.

Claim 3

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$$\Delta_n^{\mathsf{A}} = |\Pr[\mathsf{A}(f(U_n), P(U_n) \oplus 0) = 1] - \Pr[\mathsf{A}(f(U_n), P(U_n) \oplus 1) = 1]|$$

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Thus, Δ_n^A is negligible for any PPT

Statistically Binding Commitment from OWF.

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Protocol 4 ((S, R))

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Common input: 1^n

S's input: $b \in \{0,1\}$

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1. R chooses a random $r \leftarrow \{0,1\}^{3n}$ to S
2. S chooses a random $x \in \{0,1\}^n$, and send $g(x)$ to S in case $b = 0$ and $c = g(x) \oplus r$ otherwise.

Reveal: S sends (b, x) to R, and R accepts iff (b, x) is consistent with r and c

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