Application of Information Theory, Lecture 2
Joint & Conditional Entropy, Mutual Information

Iftach Haitner

Tel Aviv University.

Oct 27, 2015

Part I

Joint and Conditional Entropy

Joint entropy

- Recall that the entropy of rv X , is defined by

$$H(X) = -\sum_{x \in \mathcal{X}} \mathrm{P}_X(x) \log \mathrm{P}_X(x)$$

Joint entropy

- Recall that the entropy of rv X , is defined by

$$H(X) = -\sum_{x \in \mathcal{X}} \mathrm{P}_X(x) \log \mathrm{P}_X(x)$$

- Shorter notation: for $X \sim p$, let $H(X) = -\sum_x p(x) \log p(x)$ (where the summation is over the domain of X).

Joint entropy

- Recall that the entropy of rv X , is defined by

$$H(X) = -\sum_{x \in \mathcal{X}} \mathrm{P}_X(x) \log \mathrm{P}_X(x)$$

- Shorter notation: for $X \sim p$, let $H(X) = -\sum_x p(x) \log p(x)$ (where the summation is over the domain of X).
- The joint entropy of (jointly distributed) rvs X and Y with $(X, Y) \sim p$, is

$$H(X, Y) = -\sum_{x,y} p(x, y) \log p(x, y)$$

Joint entropy

- Recall that the entropy of rv X , is defined by

$$H(X) = -\sum_{x \in \mathcal{X}} \mathrm{P}_X(x) \log \mathrm{P}_X(x)$$

- Shorter notation: for $X \sim p$, let $H(X) = -\sum_x p(x) \log p(x)$ (where the summation is over the domain of X).
- The joint entropy of (jointly distributed) rvs X and Y with $(X, Y) \sim p$, is

$$H(X, Y) = -\sum_{x,y} p(x, y) \log p(x, y)$$

Joint entropy

- Recall that the entropy of rv X , is defined by

$$H(X) = -\sum_{x \in \mathcal{X}} \mathrm{P}_X(x) \log \mathrm{P}_X(x)$$

- Shorter notation: for $X \sim p$, let $H(X) = -\sum_x p(x) \log p(x)$ (where the summation is over the domain of X).
- The joint entropy of (jointly distributed) rvs X and Y with $(X, Y) \sim p$, is

$$H(X, Y) = -\sum_{x,y} p(x, y) \log p(x, y)$$

  This is simply the entropy of the rv $Z = (X, Y)$.

Joint entropy

- Recall that the entropy of rv X , is defined by

$$H(X) = -\sum_{x \in \mathcal{X}} \mathrm{P}_X(x) \log \mathrm{P}_X(x)$$

- Shorter notation: for $X \sim p$, let $H(X) = -\sum_x p(x) \log p(x)$ (where the summation is over the domain of X).
- The joint entropy of (jointly distributed) rvs X and Y with $(X, Y) \sim p$, is

$$H(X, Y) = -\sum_{x,y} p(x, y) \log p(x, y)$$

  This is simply the entropy of the rv $Z = (X, Y)$.
- Example:

Joint entropy

- Recall that the entropy of rv X , is defined by

$$H(X) = -\sum_{x \in \mathcal{X}} \mathrm{P}_X(x) \log \mathrm{P}_X(x)$$

- Shorter notation: for $X \sim p$, let $H(X) = -\sum_x p(x) \log p(x)$ (where the summation is over the domain of X).
- The joint entropy of (jointly distributed) rvs X and Y with $(X, Y) \sim p$, is

$$H(X, Y) = -\sum_{x,y} p(x, y) \log p(x, y)$$

This is simply the entropy of the rv $Z = (X, Y)$.
- Example:

Joint entropy

- Recall that the entropy of rv X , is defined by

$$H(X) = -\sum_{x \in \mathcal{X}} \mathrm{P}_X(x) \log \mathrm{P}_X(x)$$

- Shorter notation: for $X \sim p$, let $H(X) = -\sum_x p(x) \log p(x)$ (where the summation is over the domain of X).
- The joint entropy of (jointly distributed) rvs X and Y with $(X, Y) \sim p$, is

$$H(X, Y) = -\sum_{x,y} p(x, y) \log p(x, y)$$

  This is simply the entropy of the rv $Z = (X, Y)$.

- Example:

Joint entropy

- Recall that the entropy of rv X , is defined by

$$H(X) = -\sum_{x \in \mathcal{X}} \mathrm{P}_X(x) \log \mathrm{P}_X(x)$$

- Shorter notation: for $X \sim p$, let $H(X) = -\sum_x p(x) \log p(x)$ (where the summation is over the domain of X).
- The joint entropy of (jointly distributed) rvs X and Y with $(X, Y) \sim p$, is

$$H(X, Y) = -\sum_{x,y} p(x, y) \log p(x, y)$$

This is simply the entropy of the rv $Z = (X, Y)$.

- Example:

$X \backslash Y$	0	1
0	$\frac{1}{4}$	$\frac{1}{4}$
1	$\frac{1}{2}$	0

Joint entropy

- Recall that the entropy of rv X , is defined by

$$H(X) = -\sum_{x \in \mathcal{X}} \mathrm{P}_X(x) \log \mathrm{P}_X(x)$$

- Shorter notation: for $X \sim p$, let $H(X) = -\sum_x p(x) \log p(x)$ (where the summation is over the domain of X).
- The joint entropy of (jointly distributed) rvs X and Y with $(X, Y) \sim p$, is

$$H(X, Y) = -\sum_{x,y} p(x, y) \log p(x, y)$$

  This is simply the entropy of the rv $Z = (X, Y)$.

- Example:

$X \setminus Y$	0	1
0	$\frac{1}{4}$	$\frac{1}{4}$
1	$\frac{1}{2}$	0

Joint entropy

- Recall that the entropy of rv X , is defined by

$$H(X) = -\sum_{x \in \mathcal{X}} \mathrm{P}_X(x) \log \mathrm{P}_X(x)$$

- Shorter notation: for $X \sim p$, let $H(X) = -\sum_x p(x) \log p(x)$ (where the summation is over the domain of X).
- The joint entropy of (jointly distributed) rvs X and Y with $(X, Y) \sim p$, is

$$H(X, Y) = -\sum_{x,y} p(x, y) \log p(x, y)$$

  This is simply the entropy of the rv $Z = (X, Y)$.

- Example:

$X \backslash Y$	0	1
0	$\frac{1}{4}$	$\frac{1}{4}$
1	$\frac{1}{2}$	0

$$\begin{aligned} H(X, Y) &= -\frac{1}{2} \log \frac{1}{\frac{1}{2}} - \frac{1}{4} \log \frac{1}{\frac{1}{4}} - \frac{1}{4} \log \frac{1}{\frac{1}{4}} \\ &= \frac{1}{2} \cdot 1 + \frac{1}{2} \cdot 2 = 1\frac{1}{2} \end{aligned}$$

Joint entropy, cont.

- The joint entropy of $(X_1, \ldots, X_n) \sim p$, is

$$H(X_1, \ldots, X_n) = -\sum_{x_1, \ldots, x_n} p(x_1, \ldots, x_n) \log p(x_1, \ldots, x_n)$$

Joint entropy, cont.

- The joint entropy of $(X_1, \ldots, X_n) \sim p$, is

$$H(X_1, \ldots, X_n) = -\sum_{x_1, \ldots, x_n} p(x_1, \ldots, x_n) \log p(x_1, \ldots, x_n)$$

Joint entropy, cont.

- The joint entropy of $(X_1, \ldots, X_n) \sim p$, is

$$H(X_1, \ldots, X_n) = -\sum_{x_1, \ldots, x_n} p(x_1, \ldots, x_n) \log p(x_1, \ldots, x_n)$$

Conditional entropy

- Let $(X, Y) \sim p$.

Conditional entropy

- Let $(X, Y) \sim p$.
- For $x \in \mathrm{Supp}(X)$, the random variable $Y|X = x$ is well defined.

Conditional entropy

- Let $(X, Y) \sim p$.
- For $x \in \text{Supp}(X)$, the random variable $Y|X = x$ is well defined.
- The entropy of Y conditioned on X, is defined by

$$H(Y|X) := \mathop{\mathrm{E}}_{x \leftarrow X} H(Y|X = x) = \mathop{\mathrm{E}}_{X} H(Y|X)$$

Conditional entropy

- Let $(X, Y) \sim p$.
- For $x \in \mathrm{Supp}(X)$, the random variable $Y|X = x$ is well defined.
- The entropy of Y conditioned on X, is defined by

$$H(Y|X) := \mathop{\mathrm{E}}_{x \leftarrow X} H(Y|X = x) = \mathop{\mathrm{E}}_{X} H(Y|X)$$

- Measures the uncertainty in Y given X.

Conditional entropy

- Let $(X, Y) \sim p$.
- For $x \in \text{Supp}(X)$, the random variable $Y|X = x$ is well defined.
- The entropy of Y conditioned on X, is defined by

$$H(Y|X) := \mathop{\mathrm{E}}_{x \leftarrow X} H(Y|X = x) = \mathop{\mathrm{E}}_{X} H(Y|X)$$

- Measures the uncertainty in Y given X.
- Let p_X & $p_{Y|X}$ be the marginal & conational distributions induced by p.

Conditional entropy

- Let $(X, Y) \sim p$.
- For $x \in \mathrm{Supp}(X)$, the random variable $Y|X = x$ is well defined.
- The entropy of Y conditioned on X, is defined by

$$H(Y|X) := \mathop{\mathrm{E}}_{x \leftarrow X} H(Y|X = x) = \mathop{\mathrm{E}}_{X} H(Y|X)$$

- Measures the uncertainty in Y given X.
- Let p_X & $p_{Y|X}$ be the marginal & conational distributions induced by p.

Conditional entropy

- Let $(X, Y) \sim p$.
- For $x \in \text{Supp}(X)$, the random variable $Y|X = x$ is well defined.
- The entropy of Y conditioned on X, is defined by

$$H(Y|X) := \mathop{\mathrm{E}}_{x \leftarrow X} H(Y|X = x) = \mathop{\mathrm{E}}_{X} H(Y|X)$$

- Measures the uncertainty in Y given X.
- Let p_X & $p_{Y|X}$ be the marginal & conational distributions induced by p.

$$H(Y|X) = \sum_{x \in \mathcal{X}} p_X(x) \cdot H(Y|X = x)$$

Conditional entropy

- Let $(X, Y) \sim p$.
- For $x \in \text{Supp}(X)$, the random variable $Y|X = x$ is well defined.
- The entropy of Y conditioned on X, is defined by

$$H(Y|X) := \mathop{\mathrm{E}}_{x \leftarrow X} H(Y|X = x) = \mathop{\mathrm{E}}_{X} H(Y|X)$$

- Measures the uncertainty in Y given X.
- Let p_X & $p_{Y|X}$ be the marginal & conational distributions induced by p.

$$\begin{aligned} H(Y|X) &= \sum_{x \in \mathcal{X}} p_X(x) \cdot H(Y|X = x) \\ &= -\sum_{x \in \mathcal{X}} p_X(x) \sum_{y \in \mathcal{Y}} p_{Y|X}(y|x) \log p_{Y|X}(y|x) \end{aligned}$$

Conditional entropy

- Let $(X, Y) \sim p$.
- For $x \in \text{Supp}(X)$, the random variable $Y|X = x$ is well defined.
- The entropy of Y conditioned on X, is defined by

$$H(Y|X) := \mathop{\mathrm{E}}_{x \leftarrow X} H(Y|X = x) = \mathop{\mathrm{E}}_{X} H(Y|X)$$

- Measures the uncertainty in Y given X.
- Let p_X & $p_{Y|X}$ be the marginal & conational distributions induced by p.

$$\begin{aligned} H(Y|X) &= \sum_{x \in \mathcal{X}} p_X(x) \cdot H(Y|X = x) \\ &= -\sum_{x \in \mathcal{X}} p_X(x) \sum_{y \in \mathcal{Y}} p_{Y|X}(y|x) \log p_{Y|X}(y|x) \\ &= -\sum_{x \in \mathcal{X}, y \in \mathcal{Y}} p(x, y) \log p_{Y|X}(y|x) \end{aligned}$$

Conditional entropy

- Let $(X, Y) \sim p$.
- For $x \in \mathrm{Supp}(X)$, the random variable $Y|X = x$ is well defined.
- The entropy of Y conditioned on X, is defined by

$$H(Y|X) := \mathop{\mathrm{E}}_{x \leftarrow X} H(Y|X = x) = \mathop{\mathrm{E}}_{X} H(Y|X)$$

- Measures the uncertainty in Y given X.
- Let p_X & $p_{Y|X}$ be the marginal & conational distributions induced by p.

$$\begin{aligned}
H(Y|X) &= \sum_{x \in \mathcal{X}} p_X(x) \cdot H(Y|X = x) \\
&= -\sum_{x \in \mathcal{X}} p_X(x) \sum_{y \in \mathcal{Y}} p_{Y|X}(y|x) \log p_{Y|X}(y|x) \\
&= -\sum_{x \in \mathcal{X}, y \in \mathcal{Y}} p(x, y) \log p_{Y|X}(y|x) \\
&= -\mathop{\mathrm{E}}_{(X,Y)} \log p_{Y|X}(Y|X)
\end{aligned}$$

Conditional entropy

- Let $(X, Y) \sim p$.
- For $x \in \text{Supp}(X)$, the random variable $Y|X = x$ is well defined.
- The entropy of Y conditioned on X, is defined by

$$H(Y|X) := \mathop{\mathrm{E}}_{x \leftarrow X} H(Y|X = x) = \mathop{\mathrm{E}}_{X} H(Y|X)$$

- Measures the uncertainty in Y given X.
- Let p_X & $p_{Y|X}$ be the marginal & conational distributions induced by p.

$$\begin{aligned}
H(Y|X) &= \sum_{x \in \mathcal{X}} p_X(x) \cdot H(Y|X = x) \\
&= -\sum_{x \in \mathcal{X}} p_X(x) \sum_{y \in \mathcal{Y}} p_{Y|X}(y|x) \log p_{Y|X}(y|x) \\
&= -\sum_{x \in \mathcal{X}, y \in \mathcal{Y}} p(x, y) \log p_{Y|X}(y|x) \\
&= -\mathop{\mathrm{E}}_{(X,Y)} \log p_{Y|X}(Y|X) \\
&= -\mathop{\mathrm{E}}_{Z = p_{Y|X}(Y|X)} \log Z
\end{aligned}$$

Conditional entropy, cont.

- Example

$X \backslash Y$	0	1
0	$\frac{1}{4}$	$\frac{1}{4}$
1	$\frac{1}{2}$	0

What is $H(Y|X)$ and $H(X|Y)$?

Conditional entropy, cont.

- Example

$X \backslash Y$	0	1
0	$\frac{1}{4}$	$\frac{1}{4}$
1	$\frac{1}{2}$	0

What is $H(Y|X)$ and $H(X|Y)$?

$$\begin{aligned} H(Y|X) &= \mathop{\mathrm{E}}_{x \leftarrow X} H(Y|X = x) \\ &= \frac{1}{2} H(Y|X = 0) + \frac{1}{2} H(Y|X = 1) \\ &= \qquad \frac{1}{2} H(\frac{1}{2}, \frac{1}{2}) + \frac{1}{2} H(1, 0) = \frac{1}{2}. \end{aligned}$$

Conditional entropy, cont.

- Example

$X \backslash Y$	0	1
0	$\frac{1}{4}$	$\frac{1}{4}$
1	$\frac{1}{2}$	0

What is $H(Y|X)$ and $H(X|Y)$?

$$\begin{aligned} H(Y|X) &= \mathop{\mathrm{E}}_{x \leftarrow X} H(Y|X = x) \\ &= \frac{1}{2} H(Y|X = 0) + \frac{1}{2} H(Y|X = 1) \\ &= \frac{1}{2} H(\frac{1}{2}, \frac{1}{2}) + \frac{1}{2} H(1, 0) = \frac{1}{2}. \end{aligned}$$

$$\begin{aligned} H(X|Y) &= \mathop{\mathrm{E}}_{y \leftarrow Y} H(X|Y = y) \\ &= \frac{3}{4} H(X|Y = 0) + \frac{1}{4} H(X|Y = 1) \\ &= \frac{3}{4} H(\frac{1}{3}, \frac{2}{3}) + \frac{1}{4} H(1, 0) = 0.6887 \neq H(Y|X). \end{aligned}$$

Conditional entropy, cont..

- $$H(X|Y,Z) = \mathop{\mathrm{E}}_{(y,z)\leftarrow(Y,Z)} H(X|Y=y, Z=z)$$

Conditional entropy, cont..

- $$H(X|Y,Z) = \mathop{\mathrm{E}}_{(y,z)\leftarrow(Y,Z)} H(X|Y=y, Z=z)$$

Conditional entropy, cont..

- $$\begin{aligned} H(X|Y,Z) &= \mathop{\mathrm{E}}_{(y,z)\leftarrow(Y,Z)} H(X|Y=y,Z=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z|Y=y} H(X|Y=y,Z=z) \end{aligned}$$

Conditional entropy, cont..

- $$\begin{aligned} H(X|Y,Z) &= \mathop{\mathrm{E}}_{(y,z)\leftarrow(Y,Z)} H(X|Y=y,Z=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z|Y=y} H(X|Y=y,Z=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z|Y=y} H((X|Y=y)|Z=z) \end{aligned}$$

Conditional entropy, cont..

- $$\begin{aligned} H(X|Y,Z) &= \mathop{\mathrm{E}}_{(y,z)\leftarrow(Y,Z)} H(X|Y=y,Z=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z|Y=y} H(X|Y=y,Z=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z|Y=y} H((X|Y=y)|Z=z) \end{aligned}$$

  Let $(X_y, Z_y) = (X,Z)|Y = y$.

Conditional entropy, cont..

- $$\begin{aligned} H(X|Y,Z) &= \mathop{\mathrm{E}}_{(y,z)\leftarrow(Y,Z)} H(X|Y=y,Z=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z|Y=y} H(X|Y=y,Z=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z|Y=y} H((X|Y=y)|Z=z) \end{aligned}$$

  Let $(X_y, Z_y) = (X,Z)|Y = y$. Then

  $$H(X|Y,Z) = \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z_y} H(X_y|Z=z)$$

Conditional entropy, cont..

- $$\begin{aligned} H(X|Y,Z) &= \mathop{\mathrm{E}}_{(y,z)\leftarrow(Y,Z)} H(X|Y=y,Z=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z|Y=y} H(X|Y=y,Z=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z|Y=y} H((X|Y=y)|Z=z) \end{aligned}$$

  Let $(X_y, Z_y) = (X,Z)|Y=y$. Then

  $$\begin{aligned} H(X|Y,Z) &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z_y} H(X_y|Z=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z_y} H(X_y|Z_y=z) \end{aligned}$$

Conditional entropy, cont..

- $$\begin{aligned} H(X|Y,Z) &= \mathop{\mathrm{E}}_{(y,z)\leftarrow(Y,Z)} H(X|Y=y,Z=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z|Y=y} H(X|Y=y,Z=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z|Y=y} H((X|Y=y)|Z=z) \end{aligned}$$

  Let $(X_y,Z_y) = (X,Z)|Y=y$. Then

  $$\begin{aligned} H(X|Y,Z) &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z_y} H(X_y|Z=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y}\ \mathop{\mathrm{E}}_{z\leftarrow Z_y} H(X_y|Z_y=z) \\ &= \mathop{\mathrm{E}}_{y\leftarrow Y} H(X_y|Z_y) \end{aligned}$$

Relating mutual entropy to conditional entropy

- What is the relation between $H(X)$, $H(Y)$, $H(X,Y)$ and $H(Y|Y)$?

Relating mutual entropy to conditional entropy

- What is the relation between $H(X)$, $H(Y)$, $H(X, Y)$ and $H(Y|Y)$?
- Intuitively, $0 \leq H(Y|X) \leq H(Y)$

Relating mutual entropy to conditional entropy

- What is the relation between $H(X)$, $H(Y)$, $H(X, Y)$ and $H(Y|Y)$?
- Intuitively, $0 \leq H(Y|X) \leq H(Y)$

  Non-negativity is immediate.

Relating mutual entropy to conditional entropy

- What is the relation between $H(X)$, $H(Y)$, $H(X,Y)$ and $H(Y|Y)$?
- Intuitively, $0 \leq H(Y|X) \leq H(Y)$

  Non-negativity is immediate.

Relating mutual entropy to conditional entropy

- What is the relation between $H(X)$, $H(Y)$, $H(X,Y)$ and $H(Y|Y)$?
- Intuitively, $0 \leq H(Y|X) \leq H(Y)$

  Non-negativity is immediate. We prove upperbound later.

Relating mutual entropy to conditional entropy

- What is the relation between $H(X)$, $H(Y)$, $H(X, Y)$ and $H(Y|Y)$?
- Intuitively, $0 \leq H(Y|X) \leq H(Y)$

  Non-negativity is immediate. We prove upperbound later.
- We will also see that $H(Y|X) = H(Y)$ iff X and Y are independent.

Relating mutual entropy to conditional entropy

- What is the relation between $H(X)$, $H(Y)$, $H(X,Y)$ and $H(Y|Y)$?
- Intuitively, $0 \leq H(Y|X) \leq H(Y)$

  Non-negativity is immediate. We prove upperbound later.
- We will also see that $H(Y|X) = H(Y)$ iff X and Y are independent.
- In our example, $H(Y) = H(\frac{3}{4}, \frac{1}{4}) > \frac{1}{2} = H(Y|X)$

Relating mutual entropy to conditional entropy

- What is the relation between $H(X)$, $H(Y)$, $H(X,Y)$ and $H(Y|Y)$?
- Intuitively, $0 \leq H(Y|X) \leq H(Y)$

  Non-negativity is immediate. We prove upperbound later.
- We will also see that $H(Y|X) = H(Y)$ iff X and Y are independent.
- In our example, $H(Y) = H(\frac{3}{4}, \frac{1}{4}) > \frac{1}{2} = H(Y|X)$
- Note that $H(Y|X = x)$ might be larger than $H(Y)$ for some $x \in \mathrm{Supp}(X)$.

Relating mutual entropy to conditional entropy

- What is the relation between $H(X)$, $H(Y)$, $H(X,Y)$ and $H(Y|Y)$?
- Intuitively, $0 \leq H(Y|X) \leq H(Y)$

  Non-negativity is immediate. We prove upperbound later.
- We will also see that $H(Y|X) = H(Y)$ iff X and Y are independent.
- In our example, $H(Y) = H(\frac{3}{4}, \frac{1}{4}) > \frac{1}{2} = H(Y|X)$
- Note that $H(Y|X = x)$ might be larger than $H(Y)$ for some $x \in \mathrm{Supp}(X)$.
- Chain rule (proved next). $H(X,Y) = H(X) + H(Y|X)$

Relating mutual entropy to conditional entropy

- What is the relation between $H(X)$, $H(Y)$, $H(X, Y)$ and $H(Y|Y)$?
- Intuitively, $0 \leq H(Y|X) \leq H(Y)$

  Non-negativity is immediate. We prove upperbound later.
- We will also see that $H(Y|X) = H(Y)$ iff X and Y are independent.
- In our example, $H(Y) = H(\frac{3}{4}, \frac{1}{4}) > \frac{1}{2} = H(Y|X)$
- Note that $H(Y|X = x)$ might be larger than $H(Y)$ for some $x \in \mathrm{Supp}(X)$.
- Chain rule (proved next). $H(X, Y) = H(X) + H(Y|X)$
- Intuitively, uncertainty in (X, Y) is the uncertainty in X plus the uncertainty in Y given X.

Relating mutual entropy to conditional entropy

- What is the relation between $H(X)$, $H(Y)$, $H(X, Y)$ and $H(Y|Y)$?
- Intuitively, $0 \leq H(Y|X) \leq H(Y)$

  Non-negativity is immediate. We prove upperbound later.
- We will also see that $H(Y|X) = H(Y)$ iff X and Y are independent.
- In our example, $H(Y) = H(\frac{3}{4}, \frac{1}{4}) > \frac{1}{2} = H(Y|X)$
- Note that $H(Y|X = x)$ might be larger than $H(Y)$ for some $x \in \text{Supp}(X)$.
- Chain rule (proved next). $H(X, Y) = H(X) + H(Y|X)$
- Intuitively, uncertainty in (X, Y) is the uncertainty in X plus the uncertainty in Y given X.
- $H(Y|X) = H(X, Y) - H(X)$ is as an alternative definition for $H(Y|X)$.

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X, Y) = H(X) + H(Y|X)$.

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X,Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X, Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X, Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X, Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

$X \backslash Y$			
	$P_{1,1}$	$\ldots$	$P_{1,n}$
	$\vdots$	$\vdots$	$\vdots$
	$P_{n,1}$	$\ldots$	$P_{n,n}$

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X,Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

$X \backslash Y$			
	$P_{1,1}$	$\ldots$	$P_{1,n}$
	$\vdots$	$\vdots$	$\vdots$
	$P_{n,1}$	$\ldots$	$P_{n,n}$

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X, Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

$X \backslash Y$			
	$P_{1,1}$	$\ldots$	$P_{1,n}$
	$\vdots$	$\vdots$	$\vdots$
	$P_{n,1}$	$\ldots$	$P_{n,n}$

Let $q_i = \sum_{j=1}^{n} p_{i,j}$ $\qquad (= \Pr[X = i]$

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X,Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

$X \backslash Y$			
	$P_{1,1}$	$\ldots$	$P_{1,n}$
	$\vdots$	$\vdots$	$\vdots$
	$P_{n,1}$	$\ldots$	$P_{n,n}$

Let $q_i = \sum_{j=1}^{n} p_{i,j}$ $\quad (= \Pr[X = i]$

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X,Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

$X \backslash Y$			
	$P_{1,1}$	$\ldots$	$P_{1,n}$
	$\vdots$	$\vdots$	$\vdots$
	$P_{n,1}$	$\ldots$	$P_{n,n}$

Let $q_i = \sum_{j=1}^{n} p_{i,j}$ $\qquad (= \Pr[X = i]$

$$
\begin{aligned}
& H(P_{1,1}, \ldots, P_{n,n}) \\
& = H(q_1, \ldots, q_n) + \sum_i q_i H(\frac{P_{i,1}}{q_i}, \ldots, \frac{P_{i,n}}{q_i}) \\
& = H(X) + H(Y|X).
\end{aligned}
$$

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X,Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

$X \diagdown Y$			
	$P_{1,1}$	$\ldots$	$P_{1,n}$
	$\vdots$	$\vdots$	$\vdots$
	$P_{n,1}$	$\ldots$	$P_{n,n}$

Let $q_i = \sum_{j=1}^{n} p_{i,j}$ $\quad (= \Pr[X = i]$

$$\begin{aligned} & H(P_{1,1}, \ldots, P_{n,n}) \\ & = H(q_1, \ldots, q_n) + \sum_i q_i H(\frac{P_{i,1}}{q_i}, \ldots, \frac{P_{i,n}}{q_i}) \\ & = H(X) + H(Y|X). \end{aligned}$$

- Another proof.

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X,Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

$X \backslash Y$			
	$P_{1,1}$	$\ldots$	$P_{1,n}$
	$\vdots$	$\vdots$	$\vdots$
	$P_{n,1}$	$\ldots$	$P_{n,n}$

Let $q_i = \sum_{j=1}^{n} p_{i,j}$ $\quad (= \Pr[X = i]$

$$
\begin{aligned}
& H(P_{1,1}, \ldots, P_{n,n}) \\
& = H(q_1, \ldots, q_n) + \sum_i q_i H(\frac{P_{i,1}}{q_i}, \ldots, \frac{P_{i,n}}{q_i}) \\
& = H(X) + H(Y|X).
\end{aligned}
$$

- Another proof.

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X,Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

$X \backslash Y$			
	$P_{1,1}$	$\dots$	$P_{1,n}$
	$\vdots$	$\vdots$	$\vdots$
	$P_{n,1}$	$\dots$	$P_{n,n}$

Let $q_i = \sum_{j=1}^{n} p_{i,j}$ $\qquad (= \Pr[X = i]$

$$
\begin{aligned}
& H(P_{1,1}, \dots, P_{n,n}) \\
& = H(q_1, \dots, q_n) + \sum_i q_i H(\frac{P_{i,1}}{q_i}, \dots, \frac{P_{i,n}}{q_i}) \\
& = H(X) + H(Y|X).
\end{aligned}
$$

- Another proof. Let $(X, Y) \sim p$.

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X, Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

$X \backslash Y$			
	$P_{1,1}$	$\ldots$	$P_{1,n}$
	$\vdots$	$\vdots$	$\vdots$
	$P_{n,1}$	$\ldots$	$P_{n,n}$

Let $q_i = \sum_{j=1}^{n} p_{i,j}$ $\qquad (= \Pr[X = i]$

$$H(P_{1,1}, \ldots, P_{n,n})$$
$$= H(q_1, \ldots, q_n) + \sum_i q_i H(\frac{P_{i,1}}{q_i}, \ldots, \frac{P_{i,n}}{q_i})$$
$$= H(X) + H(Y|X).$$

- Another proof. Let $(X, Y) \sim p$.
- $p(x, y) = p_X(x) \cdot p_{Y|X}(x|y)$.

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X, Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

$X \backslash Y$			
	$P_{1,1}$	$\ldots$	$P_{1,n}$
	$\vdots$	$\vdots$	$\vdots$
	$P_{n,1}$	$\ldots$	$P_{n,n}$

Let $q_i = \sum_{j=1}^{n} p_{i,j}$ $\qquad (= \Pr[X = i]$

$$
\begin{aligned}
& H(P_{1,1}, \ldots, P_{n,n}) \\
& = H(q_1, \ldots, q_n) + \sum_i q_i H(\frac{P_{i,1}}{q_i}, \ldots, \frac{P_{i,n}}{q_i}) \\
& = H(X) + H(Y|X).
\end{aligned}
$$

- Another proof. Let $(X, Y) \sim p$.
- $p(x, y) = p_X(x) \cdot p_{Y|X}(x|y)$.

$\implies$ $\log p(x, y) = \log p_X(x) + \log p_{Y|X}(y|x)$

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X,Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

$X \backslash Y$			
	$P_{1,1}$	$\ldots$	$P_{1,n}$
	$\vdots$	$\vdots$	$\vdots$
	$P_{n,1}$	$\ldots$	$P_{n,n}$

Let $q_i = \sum_{j=1}^{n} p_{i,j}$ $\quad (= \Pr[X = i]$

$$\begin{aligned} & H(P_{1,1}, \ldots, P_{n,n}) \\ & = H(q_1, \ldots, q_n) + \sum_i q_i H(\frac{P_{i,1}}{q_i}, \ldots, \frac{P_{i,n}}{q_i}) \\ & = H(X) + H(Y|X). \end{aligned}$$

- Another proof. Let $(X,Y) \sim p$.
- $p(x,y) = p_X(x) \cdot p_{Y|X}(x|y)$.

$\implies$ $\log p(x,y) = \log p_X(x) + \log p_{Y|X}(y|x)$

$\implies$ $\mathrm{E} \log p(X,Y) = \mathrm{E} \log p_X(X) + \mathrm{E} \log p_{Y|X}(Y|X)$

Chain rule (for the entropy function)

Claim 1

For rvs X, Y, it holds that $H(X, Y) = H(X) + H(Y|X)$.

- Proof immediately follow by the grouping axiom:

$X \backslash Y$			
	$P_{1,1}$	$\ldots$	$P_{1,n}$
	$\vdots$	$\vdots$	$\vdots$
	$P_{n,1}$	$\ldots$	$P_{n,n}$

Let $q_i = \sum_{j=1}^{n} p_{i,j}$ $\qquad (= \Pr[X = i]$

$$H(P_{1,1}, \ldots, P_{n,n})$$
$$= H(q_1, \ldots, q_n) + \sum_i q_i H(\frac{P_{i,1}}{q_i}, \ldots, \frac{P_{i,n}}{q_i})$$
$$= H(X) + H(Y|X).$$

- Another proof. Let $(X, Y) \sim p$.
- $p(x, y) = p_X(x) \cdot p_{Y|X}(x|y)$.

$\implies$ $\log p(x, y) = \log p_X(x) + \log p_{Y|X}(y|x)$

$\implies$ $\mathrm{E} \log p(X, Y) = \mathrm{E} \log p_X(X) + \mathrm{E} \log p_{Y|X}(Y|X)$

$\implies$ $H(X, Y) = H(X) + H(Y|X)$.

$H(Y|X) \leq H(Y)$

$H(Y|X) \leq H(Y)$

Jensen inequality: for any concave function f, values $t_1, \ldots, t_k$ and $\lambda_1, \ldots, \lambda_k \in [0, 1]$ with $\sum_i \lambda_i = 1$, it holds that $\sum_i \lambda_i f(t_i) \leq f(\sum_i \lambda_i t_i)$.

$H(Y|X) \leq H(Y)$

Jensen inequality: for any concave function f, values $t_1, \ldots, t_k$ and $\lambda_1, \ldots, \lambda_k \in [0, 1]$ with $\sum_i \lambda_i = 1$, it holds that $\sum_i \lambda_i f(t_i) \leq f(\sum_i \lambda_i t_i)$.
Let $(X, Y) \sim p$.

$H(Y|X) \leq H(Y)$

Jensen inequality: for any concave function f, values $t_1, \ldots, t_k$ and $\lambda_1, \ldots, \lambda_k \in [0,1]$ with $\sum_i \lambda_i = 1$, it holds that $\sum_i \lambda_i f(t_i) \leq f(\sum_i \lambda_i t_i)$.
Let $(X, Y) \sim p$.

$$H(Y|X) = -\sum_{x,y} p(x,y) \log p_{Y|X}(y|x)$$

$H(Y|X) \leq H(Y)$

Jensen inequality: for any concave function f, values $t_1, \ldots, t_k$ and $\lambda_1, \ldots, \lambda_k \in [0,1]$ with $\sum_i \lambda_i = 1$, it holds that $\sum_i \lambda_i f(t_i) \leq f(\sum_i \lambda_i t_i)$.
Let $(X, Y) \sim p$.

$$\begin{aligned} H(Y|X) &= -\sum_{x,y} p(x,y) \log p_{Y|X}(y|x) \\ &= \sum_{x,y} p(x,y) \log \frac{p_X(x)}{p(x,y)} \end{aligned}$$

$H(Y|X) \leq H(Y)$

Jensen inequality: for any concave function f, values $t_1, \ldots, t_k$ and $\lambda_1, \ldots, \lambda_k \in [0,1]$ with $\sum_i \lambda_i = 1$, it holds that $\sum_i \lambda_i f(t_i) \leq f(\sum_i \lambda_i t_i)$.
Let $(X,Y) \sim p$.

$$\begin{aligned}
H(Y|X) &= -\sum_{x,y} p(x,y) \log p_{Y|X}(y|x) \\
&= \sum_{x,y} p(x,y) \log \frac{p_X(x)}{p(x,y)} \\
&= \sum_{x,y} p_Y(y) \cdot \frac{p(x,y)}{p_Y(y)} \log \frac{p_X(x)}{p(x,y)}
\end{aligned}$$

$H(Y|X) \leq H(Y)$

Jensen inequality: for any concave function f, values $t_1, \ldots, t_k$ and $\lambda_1, \ldots, \lambda_k \in [0,1]$ with $\sum_i \lambda_i = 1$, it holds that $\sum_i \lambda_i f(t_i) \leq f(\sum_i \lambda_i t_i)$.
Let $(X, Y) \sim p$.

$$\begin{aligned}
H(Y|X) &= -\sum_{x,y} p(x,y) \log p_{Y|X}(y|x) \\
&= \sum_{x,y} p(x,y) \log \frac{p_X(x)}{p(x,y)} \\
&= \sum_{x,y} p_Y(y) \cdot \frac{p(x,y)}{p_Y(y)} \log \frac{p_X(x)}{p(x,y)} \\
&= \sum_{y} p_Y(y) \sum_{x} \frac{p(x,y)}{p_Y(y)} \log \frac{p_X(x)}{p(x,y)}
\end{aligned}$$

$H(Y|X) \leq H(Y)$

Jensen inequality: for any concave function f, values $t_1, \ldots, t_k$ and $\lambda_1, \ldots, \lambda_k \in [0,1]$ with $\sum_i \lambda_i = 1$, it holds that $\sum_i \lambda_i f(t_i) \leq f(\sum_i \lambda_i t_i)$.
Let $(X,Y) \sim p$.

$$\begin{aligned}
H(Y|X) &= -\sum_{x,y} p(x,y) \log p_{Y|X}(y|x) \\
&= \sum_{x,y} p(x,y) \log \frac{p_X(x)}{p(x,y)} \\
&= \sum_{x,y} p_Y(y) \cdot \frac{p(x,y)}{p_Y(y)} \log \frac{p_X(x)}{p(x,y)} \\
&= \sum_{y} p_Y(y) \sum_{x} \frac{p(x,y)}{p_Y(y)} \log \frac{p_X(x)}{p(x,y)} \\
&\leq \sum_{y} p_Y(y) \log \sum_{x} \frac{p(x,y)}{p_Y(y)} \frac{p_X(x)}{p(x,y)}
\end{aligned}$$

$H(Y|X) \leq H(Y)$

Jensen inequality: for any concave function f, values $t_1, \ldots, t_k$ and $\lambda_1, \ldots, \lambda_k \in [0,1]$ with $\sum_i \lambda_i = 1$, it holds that $\sum_i \lambda_i f(t_i) \leq f(\sum_i \lambda_i t_i)$.
Let $(X, Y) \sim p$.

$$
\begin{aligned}
H(Y|X) &= -\sum_{x,y} p(x,y) \log p_{Y|X}(y|x) \\
&= \sum_{x,y} p(x,y) \log \frac{p_X(x)}{p(x,y)} \\
&= \sum_{x,y} p_Y(y) \cdot \frac{p(x,y)}{p_Y(y)} \log \frac{p_X(x)}{p(x,y)} \\
&= \sum_{y} p_Y(y) \sum_{x} \frac{p(x,y)}{p_Y(y)} \log \frac{p_X(x)}{p(x,y)} \\
&\leq \sum_{y} p_Y(y) \log \sum_{x} \frac{p(x,y)}{p_Y(y)} \frac{p_X(x)}{p(x,y)} \\
&= \sum_{y} p_Y(y) \log \frac{1}{p_Y(y)}
\end{aligned}
$$

$H(Y|X) \leq H(Y)$

Jensen inequality: for any concave function f, values $t_1, \ldots, t_k$ and $\lambda_1, \ldots, \lambda_k \in [0,1]$ with $\sum_i \lambda_i = 1$, it holds that $\sum_i \lambda_i f(t_i) \leq f(\sum_i \lambda_i t_i)$.
Let $(X,Y) \sim p$.

$$
\begin{aligned}
H(Y|X) &= -\sum_{x,y} p(x,y) \log p_{Y|X}(y|x) \\
&= \sum_{x,y} p(x,y) \log \frac{p_X(x)}{p(x,y)} \\
&= \sum_{x,y} p_Y(y) \cdot \frac{p(x,y)}{p_Y(y)} \log \frac{p_X(x)}{p(x,y)} \\
&= \sum_{y} p_Y(y) \sum_{x} \frac{p(x,y)}{p_Y(y)} \log \frac{p_X(x)}{p(x,y)} \\
&\leq \sum_{y} p_Y(y) \log \sum_{x} \frac{p(x,y)}{p_Y(y)} \frac{p_X(x)}{p(x,y)} \\
&= \sum_{y} p_Y(y) \log \frac{1}{p_Y(y)} = H(Y).
\end{aligned}
$$

$H(Y|X) \leq H(Y)$

Jensen inequality: for any concave function f, values $t_1, \ldots, t_k$ and $\lambda_1, \ldots, \lambda_k \in [0,1]$ with $\sum_i \lambda_i = 1$, it holds that $\sum_i \lambda_i f(t_i) \leq f(\sum_i \lambda_i t_i)$.
Let $(X, Y) \sim p$.

$$
\begin{aligned}
H(Y|X) &= -\sum_{x,y} p(x,y) \log p_{Y|X}(y|x) \\
&= \sum_{x,y} p(x,y) \log \frac{p_X(x)}{p(x,y)} \\
&= \sum_{x,y} p_Y(y) \cdot \frac{p(x,y)}{p_Y(y)} \log \frac{p_X(x)}{p(x,y)} \\
&= \sum_{y} p_Y(y) \sum_{x} \frac{p(x,y)}{p_Y(y)} \log \frac{p_X(x)}{p(x,y)} \\
&\leq \sum_{y} p_Y(y) \log \sum_{x} \frac{p(x,y)}{p_Y(y)} \frac{p_X(x)}{p(x,y)} \\
&= \sum_{y} p_Y(y) \log \frac{1}{p_Y(y)} = H(Y).
\end{aligned}
$$

$H(Y|X) \leq H(Y)$ cont.

- Assume X and Y are independent (i.e., $p(x, y) = p_X(x) \cdot p_Y(y)$ for any x, y)

$H(Y|X) \leq H(Y)$ cont.

- Assume X and Y are independent (i.e., $p(x,y) = p_X(x) \cdot p_Y(y)$ for any x, y)

$\implies p_{Y|X} = p_Y$

$H(Y|X) \leq H(Y)$ cont.

- Assume X and Y are independent (i.e., $p(x,y) = p_X(x) \cdot p_Y(y)$ for any x, y)

$$\implies \quad p_{Y|X} = p_Y$$

$$\implies \quad H(Y|X) = H(Y)$$

$H(Y|X) \leq H(Y)$ cont.

- Assume X and Y are independent (i.e., $p(x,y) = p_X(x) \cdot p_Y(y)$ for any x, y)

$\implies \quad p_{Y|X} = p_Y$

$\implies \quad H(Y|X) = H(Y)$

- Is the converse also true: $H(Y|X) = H(Y)$ implies X and Y are independent?

$H(Y|X) \leq H(Y)$ cont.

- Assume X and Y are independent (i.e., $p(x,y) = p_X(x) \cdot p_Y(y)$ for any x, y)

$\implies \quad p_{Y|X} = p_Y$

$\implies \quad H(Y|X) = H(Y)$

- Is the converse also true: $H(Y|X) = H(Y)$ implies X and Y are independent?

  Yes, since log is strictly concave in the range. Equality happens iff all t_i are the same,

$H(Y|X) \leq H(Y)$ cont.

- Assume X and Y are independent (i.e., $p(x,y) = p_X(x) \cdot p_Y(y)$ for any x, y)

$\implies \quad p_{Y|X} = p_Y$

$\implies \quad H(Y|X) = H(Y)$

- Is the converse also true: $H(Y|X) = H(Y)$ implies X and Y are independent?

  Yes, since log is strictly concave in the range. Equality happens iff all t_i are the same,

- which happens iff $p(x,y) = p_X(x)p_Y(y)$ for all x, y

Other inequalities

Other inequalities

- $H(X), H(Y) \leq H(X, Y) \leq H(X) + H(Y)$.

Other inequalities

- $H(X), H(Y) \leq H(X, Y) \leq H(X) + H(Y)$.

Other inequalities

- $H(X), H(Y) \leq H(X,Y) \leq H(X) + H(Y)$.
  Follows from $H(X,Y) = H(X) + H(Y|X)$.

Other inequalities

- $H(X), H(Y) \leq H(X,Y) \leq H(X) + H(Y)$.
  Follows from $H(X,Y) = H(X) + H(Y|X)$.
  - Left inequality since $H(Y|X)$ is non negative.

Other inequalities

- $H(X), H(Y) \leq H(X,Y) \leq H(X) + H(Y)$.
  Follows from $H(X,Y) = H(X) + H(Y|X)$.
  - Left inequality since $H(Y|X)$ is non negative.
  - Right inequality since $H(Y|X) \leq H(Y)$.

Other inequalities

- $H(X), H(Y) \leq H(X, Y) \leq H(X) + H(Y)$.
  Follows from $H(X, Y) = H(X) + H(Y|X)$.
  - Left inequality since $H(Y|X)$ is non negative.
  - Right inequality since $H(Y|X) \leq H(Y)$.
- $\mathrm{H}(X, |Z) = H(X|Z) + H(Y|X, Z)$ (by chain rule)

Other inequalities

- $H(X), H(Y) \leq H(X,Y) \leq H(X) + H(Y)$.
  Follows from $H(X,Y) = H(X) + H(Y|X)$.
  - Left inequality since $H(Y|X)$ is non negative.
  - Right inequality since $H(Y|X) \leq H(Y)$.
- $\mathrm{H}(X,|Z) = H(X|Z) + H(Y|X,Z)$ (by chain rule)
- $H(X|Y,Z) \leq \mathrm{H}(X|Y)$

Other inequalities

- $H(X), H(Y) \leq H(X,Y) \leq H(X) + H(Y)$.
  Follows from $H(X,Y) = H(X) + H(Y|X)$.
  - Left inequality since $H(Y|X)$ is non negative.
  - Right inequality since $H(Y|X) \leq H(Y)$.
- $\mathrm{H}(X,|Z) = H(X|Z) + H(Y|X,Z)$ (by chain rule)
- $H(X|Y,Z) \leq \mathrm{H}(X|Y)$

Proof:
$$H(X|Y,Z) = \mathop{\mathrm{E}}_{Z,Y} H(X \mid Y,Z)$$

Other inequalities

- $H(X), H(Y) \leq H(X,Y) \leq H(X) + H(Y)$.
  Follows from $H(X,Y) = H(X) + H(Y|X)$.
  - Left inequality since $H(Y|X)$ is non negative.
  - Right inequality since $H(Y|X) \leq H(Y)$.
- $\mathrm{H}(X,|Z) = H(X|Z) + H(Y|X,Z)$ (by chain rule)
- $H(X|Y,Z) \leq \mathrm{H}(X|Y)$

Proof:

$$H(X|Y,Z) = \mathop{\mathrm{E}}_{Z,Y} H(X \mid Y, Z)$$

Other inequalities

- $H(X), H(Y) \leq H(X,Y) \leq H(X) + H(Y)$.
  Follows from $H(X,Y) = H(X) + H(Y|X)$.
  - Left inequality since $H(Y|X)$ is non negative.
  - Right inequality since $H(Y|X) \leq H(Y)$.
- $\mathrm{H}(X,|Z) = H(X|Z) + H(Y|X,Z)$ (by chain rule)
- $H(X|Y,Z) \leq \mathrm{H}(X|Y)$

Proof:

$$\begin{aligned} H(X|Y,Z) &= \mathop{\mathrm{E}}_{Z,Y} H(X \mid Y,Z) \\ &= \mathop{\mathrm{E}}_{Y} \mathop{\mathrm{E}}_{Z|Y} H(X \mid Y,Z) \end{aligned}$$

Other inequalities

- $H(X), H(Y) \leq H(X,Y) \leq H(X) + H(Y)$.
  Follows from $H(X,Y) = H(X) + H(Y|X)$.
  - Left inequality since $H(Y|X)$ is non negative.
  - Right inequality since $H(Y|X) \leq H(Y)$.
- $\mathrm{H}(X,|Z) = H(X|Z) + H(Y|X,Z)$ (by chain rule)
- $H(X|Y,Z) \leq \mathrm{H}(X|Y)$

Proof:

$$\begin{aligned} H(X|Y,Z) &= \mathop{\mathrm{E}}_{Z,Y} H(X \mid Y, Z) \\ &= \mathop{\mathrm{E}}_{Y} \mathop{\mathrm{E}}_{Z|Y} H(X \mid Y, Z) \\ &= \mathop{\mathrm{E}}_{Y} \mathop{\mathrm{E}}_{Z|Y} H((X \mid Y) \mid Z) \end{aligned}$$

Other inequalities

- $H(X), H(Y) \leq H(X, Y) \leq H(X) + H(Y)$.
  Follows from $H(X, Y) = H(X) + H(Y|X)$.
  - Left inequality since $H(Y|X)$ is non negative.
  - Right inequality since $H(Y|X) \leq H(Y)$.
- $\mathrm{H}(X, |Z) = H(X|Z) + H(Y|X, Z)$ (by chain rule)
- $H(X|Y, Z) \leq \mathrm{H}(X|Y)$

Proof:

$$\begin{aligned} H(X|Y,Z) &= \mathop{\mathrm{E}}_{Z,Y} H(X \mid Y, Z) \\ &= \mathop{\mathrm{E}}_{Y} \mathop{\mathrm{E}}_{Z|Y} H(X \mid Y, Z) \\ &= \mathop{\mathrm{E}}_{Y} \mathop{\mathrm{E}}_{Z|Y} H((X \mid Y) \mid Z) \\ &\leq \mathop{\mathrm{E}}_{Y} \mathop{\mathrm{E}}_{Z|Y} H(X \mid Y) \end{aligned}$$

Other inequalities

- $H(X), H(Y) \leq H(X,Y) \leq H(X) + H(Y)$.
  Follows from $H(X,Y) = H(X) + H(Y|X)$.
  - Left inequality since $H(Y|X)$ is non negative.
  - Right inequality since $H(Y|X) \leq H(Y)$.
- $\mathrm{H}(X,|Z) = H(X|Z) + H(Y|X,Z)$ (by chain rule)
- $H(X|Y,Z) \leq \mathrm{H}(X|Y)$

  Proof:
$$\begin{aligned} H(X|Y,Z) &= \mathop{\mathrm{E}}_{Z,Y} H(X \mid Y,Z) \\ &= \mathop{\mathrm{E}}_{Y} \mathop{\mathrm{E}}_{Z|Y} H(X \mid Y,Z) \\ &= \mathop{\mathrm{E}}_{Y} \mathop{\mathrm{E}}_{Z|Y} H((X \mid Y) \mid Z) \\ &\leq \mathop{\mathrm{E}}_{Y} \mathop{\mathrm{E}}_{Z|Y} H(X \mid Y) \\ &= \mathop{\mathrm{E}}_{Y} H(X \mid Y) \end{aligned}$$

Other inequalities

- $H(X), H(Y) \leq H(X, Y) \leq H(X) + H(Y)$.
  Follows from $H(X, Y) = H(X) + H(Y|X)$.
  - Left inequality since $H(Y|X)$ is non negative.
  - Right inequality since $H(Y|X) \leq H(Y)$.
- $\mathrm{H}(X, |Z) = H(X|Z) + H(Y|X, Z)$ (by chain rule)
- $H(X|Y, Z) \leq \mathrm{H}(X|Y)$

  Proof:

$$\begin{aligned} H(X|Y,Z) &= \mathop{\mathrm{E}}_{Z,Y} H(X \mid Y, Z) \\ &= \mathop{\mathrm{E}}_{Y} \mathop{\mathrm{E}}_{Z|Y} H(X \mid Y, Z) \\ &= \mathop{\mathrm{E}}_{Y} \mathop{\mathrm{E}}_{Z|Y} H((X \mid Y) \mid Z) \\ &\leq \mathop{\mathrm{E}}_{Y} \mathop{\mathrm{E}}_{Z|Y} H(X \mid Y) \\ &= \mathop{\mathrm{E}}_{Y} H(X \mid Y) \\ &= H(X|Y). \end{aligned}$$

Chain rule (for the entropy function), general case

Claim 2

For rvs $X_1, \ldots, X_k$, it holds that
$H(X_1, \ldots, X_k) = H(X_i) + H(X_2|X_1) + \ldots + H(X_k|X_1, \ldots, X_{k-1})$.

Chain rule (for the entropy function), general case

Claim 2

For rvs $X_1, \ldots, X_k$, it holds that
$H(X_1, \ldots, X_k) = H(X_i) + H(X_2|X_1) + \ldots + H(X_k|X_1, \ldots, X_{k-1})$.

Proof: ?

Chain rule (for the entropy function), general case

Claim 2

For rvs $X_1, \ldots, X_k$, it holds that
$H(X_1, \ldots, X_k) = H(X_i) + H(X_2|X_1) + \ldots + H(X_k|X_1, \ldots, X_{k-1})$.

Proof: ?

- Extremely useful property!

Chain rule (for the entropy function), general case

Claim 2

For rvs $X_1, \ldots, X_k$, it holds that
$H(X_1, \ldots, X_k) = H(X_i) + H(X_2|X_1) + \ldots + H(X_k|X_1, \ldots, X_{k-1})$.

Proof: ?

- Extremely useful property!
- Analogously to the two variables case, it also holds that:

Chain rule (for the entropy function), general case

Claim 2

For rvs $X_1, \ldots, X_k$, it holds that
$H(X_1, \ldots, X_k) = H(X_i) + H(X_2|X_1) + \ldots + H(X_k|X_1, \ldots, X_{k-1})$.

Proof: ?

- Extremely useful property!
- Analogously to the two variables case, it also holds that:
- $H(X_i) \leq H(X_1, \ldots, X_k) \leq \sum_i H(X_i)$

Chain rule (for the entropy function), general case

Claim 2

For rvs $X_1, \ldots, X_k$, it holds that
$H(X_1, \ldots, X_k) = H(X_i) + H(X_2|X_1) + \ldots + H(X_k|X_1, \ldots, X_{k-1})$.

Proof: ?

- Extremely useful property!
- Analogously to the two variables case, it also holds that:
- $H(X_i) \leq H(X_1, \ldots, X_k) \leq \sum_i H(X_i)$
- $H(X_1, \ldots, X_K|Y) \leq \sum_i H(X_i|Y)$

Examples

Examples

- (from last class) Let $X_1, \ldots, X_n$ be Boolean iid with $X_i \sim (\frac{1}{3}, \frac{2}{3})$. Compute $H(X_1, \ldots, X_n)$

Examples

- (from last class) Let $X_1, \ldots, X_n$ be Boolean iid with $X_i \sim (\frac{1}{3}, \frac{2}{3})$. Compute $H(X_1, \ldots, X_n)$
- As above, but under the condition that $\bigoplus_i X_i = 0$?

Examples

- (from last class) Let $X_1, \ldots, X_n$ be Boolean iid with $X_i \sim (\frac{1}{3}, \frac{2}{3})$. Compute $H(X_1, \ldots, X_n)$
- As above, but under the condition that $\bigoplus_i X_i = 0$?

Examples

- (from last class) Let $X_1, \ldots, X_n$ be Boolean iid with $X_i \sim (\frac{1}{3}, \frac{2}{3})$. Compute $H(X_1, \ldots, X_n)$
- As above, but under the condition that $\bigoplus_i X_i = 0$?
  - Via chain rule?

Examples

- (from last class) Let $X_1, \ldots, X_n$ be Boolean iid with $X_i \sim (\frac{1}{3}, \frac{2}{3})$. Compute $H(X_1, \ldots, X_n)$
- As above, but under the condition that $\bigoplus_i X_i = 0$?
  - Via chain rule?
  - Via mapping?

Applications

- Let $X_1, \ldots, X_n$ be Boolean iids with $X_i \sim (p, 1-p)$ and let $X = X_1, \ldots, X_n$. Let f be such that $\Pr[f(X) = z] = \Pr[f(X) = z']$, for every $k \in \mathbb{N}$ and $z, z' \in \{0,1\}^k$. Let $K = |f(X)|$.

  Prove that $\mathrm{E}\, K \leq n \cdot h(p)$.

Applications

- Let $X_1, \ldots, X_n$ be Boolean iids with $X_i \sim (p, 1-p)$ and let $X = X_1, \ldots, X_n$. Let f be such that $\Pr[f(X) = z] = \Pr[f(X) = z']$, for every $k \in \mathbb{N}$ and $z, z' \in \{0,1\}^k$. Let $K = |f(X)|$.

  Prove that $\mathrm{E}\, K \leq n \cdot h(p)$.

- $$n \cdot h(p) = H(X_1, \ldots, X_n)$$

Applications

- Let $X_1, \ldots, X_n$ be Boolean iids with $X_i \sim (p, 1-p)$ and let $X = X_1, \ldots, X_n$. Let f be such that $\Pr[f(X) = z] = \Pr[f(X) = z']$, for every $k \in \mathbb{N}$ and $z, z' \in \{0,1\}^k$. Let $K = |f(X)|$.

  Prove that $\mathrm{E}\, K \leq n \cdot h(p)$.

- $$n \cdot h(p) = H(X_1, \ldots, X_n)$$

Applications

- Let $X_1, \ldots, X_n$ be Boolean iids with $X_i \sim (p, 1-p)$ and let $X = X_1, \ldots, X_n$. Let f be such that $\Pr[f(X) = z] = \Pr[f(X) = z']$, for every $k \in \mathbb{N}$ and $z, z' \in \{0,1\}^k$. Let $K = |f(X)|$.

  Prove that $\mathrm{E}\, K \le n \cdot h(p)$.

- $$\begin{aligned} n \cdot h(p) &= H(X_1, \ldots, X_n) \\ &\ge H(f(X), K) \end{aligned}$$

Applications

- Let $X_1, \ldots, X_n$ be Boolean iids with $X_i \sim (p, 1-p)$ and let $X = X_1, \ldots, X_n$. Let f be such that $\Pr[f(X) = z] = \Pr[f(X) = z']$, for every $k \in \mathbb{N}$ and $z, z' \in \{0,1\}^k$. Let $K = |f(X)|$.

  Prove that $\mathrm{E}\, K \leq n \cdot h(p)$.

- $$\begin{aligned} n \cdot h(p) &= H(X_1, \ldots, X_n) \\ &\geq H(f(X), K) \\ &= H(K) + H(f(X) \mid K) \end{aligned}$$

Applications

- Let $X_1, \ldots, X_n$ be Boolean iids with $X_i \sim (p, 1-p)$ and let $X = X_1, \ldots, X_n$. Let f be such that $\Pr[f(X) = z] = \Pr[f(X) = z']$, for every $k \in \mathbb{N}$ and $z, z' \in \{0,1\}^k$. Let $K = |f(X)|$.

  Prove that $\mathrm{E}\, K \leq n \cdot h(p)$.

- $$\begin{aligned} n \cdot h(p) &= H(X_1, \ldots, X_n) \\ &\geq H(f(X), K) \\ &= H(K) + H(f(X) \mid K) \\ &= H(K) + \mathrm{E}\, K \end{aligned}$$

Applications

- Let $X_1, \ldots, X_n$ be Boolean iids with $X_i \sim (p, 1-p)$ and let $X = X_1, \ldots, X_n$. Let f be such that $\Pr[f(X) = z] = \Pr[f(X) = z']$, for every $k \in \mathbb{N}$ and $z, z' \in \{0,1\}^k$. Let $K = |f(X)|$.

  Prove that $\mathrm{E}\, K \leq n \cdot h(p)$.

- $$\begin{aligned} n \cdot h(p) &= H(X_1, \ldots, X_n) \\ &\geq H(f(X), K) \\ &= H(K) + H(f(X) \mid K) \\ &= H(K) + \mathrm{E}\, K \\ &\geq \mathrm{E}\, K \end{aligned}$$

Applications

- Let $X_1,\ldots,X_n$ be Boolean iids with $X_i \sim (p, 1-p)$ and let $X = X_1,\ldots,X_n$. Let f be such that $\Pr[f(X) = z] = \Pr[f(X) = z']$, for every $k \in \mathbb{N}$ and $z, z' \in \{0,1\}^k$. Let $K = |f(X)|$.

  Prove that $\operatorname{E} K \leq n \cdot h(p)$.

- $$\begin{aligned} n \cdot h(p) &= H(X_1,\ldots,X_n) \\ &\geq H(f(X), K) \\ &= H(K) + H(f(X) \mid K) \\ &= H(K) + \operatorname{E} K \\ &\geq \operatorname{E} K \end{aligned}$$

- Interpretation

Applications

- Let $X_1, \ldots, X_n$ be Boolean iids with $X_i \sim (p, 1-p)$ and let $X = X_1, \ldots, X_n$. Let f be such that $\Pr[f(X) = z] = \Pr[f(X) = z']$, for every $k \in \mathbb{N}$ and $z, z' \in \{0,1\}^k$. Let $K = |f(X)|$.

  Prove that $\mathrm{E}\, K \leq n \cdot h(p)$.

- $$\begin{aligned} n \cdot h(p) &= H(X_1, \ldots, X_n) \\ &\geq H(f(X), K) \\ &= H(K) + H(f(X) \mid K) \\ &= H(K) + \mathrm{E}\, K \\ &\geq \mathrm{E}\, K \end{aligned}$$

- Interpretation
- Upper bounds

Applications cont.

- How many comparisons it takes to sort n elements?

Applications cont.

- How many comparisons it takes to sort n elements?

  Let A be a sorter for n elements algorithm making t comparisons.

  What can we say about t?

Applications cont.

- How many comparisons it takes to sort n elements?

  Let A be a sorter for n elements algorithm making t comparisons.

  What can we say about t?
- Let X be a uniform random permutation of $[n]$ and let $Y_1, \ldots, Y_t$ be the answers A gets when sorting X.

Applications cont.

- How many comparisons it takes to sort n elements?

  Let A be a sorter for n elements algorithm making t comparisons.

  What can we say about t?
- Let X be a uniform random permutation of $[n]$ and let $Y_1, \ldots, Y_t$ be the answers A gets when sorting X.
- X is determined by $Y_1, \ldots, Y_t$.

Applications cont.

- How many comparisons it takes to sort n elements?

  Let A be a sorter for n elements algorithm making t comparisons.

  What can we say about t?

- Let X be a uniform random permutation of $[n]$ and let $Y_1, \ldots, Y_t$ be the answers A gets when sorting X.

- X is determined by $Y_1, \ldots, Y_t$.

  Namely, $X = f(Y_1, \ldots, Y_t)$ for some function f.

Applications cont.

- How many comparisons it takes to sort n elements?

  Let A be a sorter for n elements algorithm making t comparisons.

  What can we say about t?
- Let X be a uniform random permutation of $[n]$ and let $Y_1, \ldots, Y_t$ be the answers A gets when sorting X.
- X is determined by $Y_1, \ldots, Y_t$.

  Namely, $X = f(Y_1, \ldots, Y_t)$ for some function f.
- $H(X) = \log n!$

Applications cont.

- How many comparisons it takes to sort n elements?

  Let A be a sorter for n elements algorithm making t comparisons.

  What can we say about t?
- Let X be a uniform random permutation of $[n]$ and let $Y_1, \ldots, Y_t$ be the answers A gets when sorting X.
- X is determined by $Y_1, \ldots, Y_t$.

  Namely, $X = f(Y_1, \ldots, Y_t)$ for some function f.
- $H(X) = \log n!$
- $$H(X) = H(f(Y_1, \ldots, Y_n))$$

Applications cont.

- How many comparisons it takes to sort n elements?

  Let A be a sorter for n elements algorithm making t comparisons.

  What can we say about t?
- Let X be a uniform random permutation of $[n]$ and let $Y_1, \ldots, Y_t$ be the answers A gets when sorting X.
- X is determined by $Y_1, \ldots, Y_t$.

  Namely, $X = f(Y_1, \ldots, Y_t)$ for some function f.
- $H(X) = \log n!$
- $$H(X) = H(f(Y_1, \ldots, Y_n))$$

Applications cont.

- How many comparisons it takes to sort n elements?

  Let A be a sorter for n elements algorithm making t comparisons.

  What can we say about t?
- Let X be a uniform random permutation of $[n]$ and let $Y_1, \ldots, Y_t$ be the answers A gets when sorting X.
- X is determined by $Y_1, \ldots, Y_t$.

  Namely, $X = f(Y_1, \ldots, Y_t)$ for some function f.
- $H(X) = \log n!$
- $$\begin{aligned} H(X) &= H(f(Y_1, \ldots, Y_n)) \\ &\leq H(Y_1, \ldots, Y_n) \end{aligned}$$

Applications cont.

- How many comparisons it takes to sort n elements?

  Let A be a sorter for n elements algorithm making t comparisons.

  What can we say about t?
- Let X be a uniform random permutation of $[n]$ and let $Y_1, \ldots, Y_t$ be the answers A gets when sorting X.
- X is determined by $Y_1, \ldots, Y_t$.

  Namely, $X = f(Y_1, \ldots, Y_t)$ for some function f.
- $H(X) = \log n!$
- $$\begin{aligned} H(X) &= H(f(Y_1, \ldots, Y_n)) \\ &\leq H(Y_1, \ldots, Y_n) \\ &\leq \sum_i H(Y_i) \end{aligned}$$

Applications cont.

- How many comparisons it takes to sort n elements?

  Let A be a sorter for n elements algorithm making t comparisons.
  What can we say about t?
- Let X be a uniform random permutation of $[n]$ and let $Y_1, \ldots, Y_t$ be the answers A gets when sorting X.
- X is determined by $Y_1, \ldots, Y_t$.

  Namely, $X = f(Y_1, \ldots, Y_t)$ for some function f.
- $H(X) = \log n!$
- $$\begin{aligned} H(X) &= H(f(Y_1, \ldots, Y_n)) \\ &\leq H(Y_1, \ldots, Y_n) \\ &\leq \sum_i H(Y_i) \\ &= t \end{aligned}$$

Applications cont.

- How many comparisons it takes to sort n elements?

  Let A be a sorter for n elements algorithm making t comparisons.

  What can we say about t?
- Let X be a uniform random permutation of $[n]$ and let $Y_1, \ldots, Y_t$ be the answers A gets when sorting X.
- X is determined by $Y_1, \ldots, Y_t$.

  Namely, $X = f(Y_1, \ldots, Y_t)$ for some function f.
- $H(X) = \log n!$
- $$\begin{aligned} H(X) &= H(f(Y_1, \ldots, Y_n)) \\ &\leq H(Y_1, \ldots, Y_n) \\ &\leq \sum_i H(Y_i) \\ &= t \end{aligned}$$

$\implies \quad t \geq \log n! = \Theta(n \log n)$

Concavity of entropy function

Concavity of entropy function

Let $p = (p_1, \ldots, p_n)$ and $q = (q_1, \ldots, q_n)$ be two distributions, and for $\lambda \in [0, 1]$ consider the distribution $\tau_\lambda = \lambda p + (1 - \lambda)q$.
(i.e., $\tau_\lambda = (\lambda p_1 + (1 - \lambda)q_1, \ldots, \lambda p_n + (1 - \lambda)q_n)$.

Concavity of entropy function

Let $p = (p_1, \ldots, p_n)$ and $q = (q_1, \ldots, q_n)$ be two distributions, and for $\lambda \in [0, 1]$ consider the distribution $\tau_\lambda = \lambda p + (1 - \lambda)q$.
(i.e., $\tau_\lambda = (\lambda p_1 + (1 - \lambda)q_1, \ldots, \lambda p_n + (1 - \lambda)q_n)$).

Claim 3

$H(\tau_\lambda) \geq \lambda H(p) + (1 - \lambda)H(q)$

Concavity of entropy function

Let $p = (p_1, \ldots, p_n)$ and $q = (q_1, \ldots, q_n)$ be two distributions, and for $\lambda \in [0, 1]$ consider the distribution $\tau_\lambda = \lambda p + (1 - \lambda)q$.
(i.e., $\tau_\lambda = (\lambda p_1 + (1 - \lambda)q_1, \ldots, \lambda p_n + (1 - \lambda)q_n)$.

Claim 3

$H(\tau_\lambda) \geq \lambda H(p) + (1 - \lambda)H(q)$

Proof:

Concavity of entropy function

Let $p = (p_1, \ldots, p_n)$ and $q = (q_1, \ldots, q_n)$ be two distributions, and for $\lambda \in [0, 1]$ consider the distribution $\tau_\lambda = \lambda p + (1 - \lambda)q$.
(i.e., $\tau_\lambda = (\lambda p_1 + (1 - \lambda)q_1, \ldots, \lambda p_n + (1 - \lambda)q_n)$.

Claim 3

$H(\tau_\lambda) \geq \lambda H(p) + (1 - \lambda)H(q)$

Proof:

- Let Y over $\{0, 1\}$ be 1 wp λ

Concavity of entropy function

Let $p = (p_1, \ldots, p_n)$ and $q = (q_1, \ldots, q_n)$ be two distributions, and for $\lambda \in [0, 1]$ consider the distribution $\tau_\lambda = \lambda p + (1 - \lambda)q$.
(i.e., $\tau_\lambda = (\lambda p_1 + (1 - \lambda)q_1, \ldots, \lambda p_n + (1 - \lambda)q_n)$.

Claim 3

$H(\tau_\lambda) \geq \lambda H(p) + (1 - \lambda)H(q)$

Proof:

- Let Y over $\{0, 1\}$ be 1 wp λ
- Let X be distributed according to p if $Y = 0$ and according to q otherwise.

Concavity of entropy function

Let $p = (p_1, \ldots, p_n)$ and $q = (q_1, \ldots, q_n)$ be two distributions, and for $\lambda \in [0, 1]$ consider the distribution $\tau_\lambda = \lambda p + (1 - \lambda)q$.
(i.e., $\tau_\lambda = (\lambda p_1 + (1 - \lambda)q_1, \ldots, \lambda p_n + (1 - \lambda)q_n)$.

Claim 3

$H(\tau_\lambda) \geq \lambda H(p) + (1 - \lambda)H(q)$

Proof:

- Let Y over $\{0, 1\}$ be 1 wp λ
- Let X be distributed according to p if $Y = 0$ and according to q otherwise.
- $H(\tau_\lambda) = H(X)$

Concavity of entropy function

Let $p = (p_1, \ldots, p_n)$ and $q = (q_1, \ldots, q_n)$ be two distributions, and for $\lambda \in [0, 1]$ consider the distribution $\tau_\lambda = \lambda p + (1 - \lambda)q$.
(i.e., $\tau_\lambda = (\lambda p_1 + (1 - \lambda)q_1, \ldots, \lambda p_n + (1 - \lambda)q_n)$.

Claim 3

$H(\tau_\lambda) \geq \lambda H(p) + (1 - \lambda)H(q)$

Proof:

- Let Y over $\{0, 1\}$ be 1 wp λ
- Let X be distributed according to p if $Y = 0$ and according to q otherwise.
- $H(\tau_\lambda) = H(X)$

Concavity of entropy function

Let $p = (p_1, \ldots, p_n)$ and $q = (q_1, \ldots, q_n)$ be two distributions, and for $\lambda \in [0,1]$ consider the distribution $\tau_\lambda = \lambda p + (1-\lambda)q$.
(i.e., $\tau_\lambda = (\lambda p_1 + (1-\lambda)q_1, \ldots, \lambda p_n + (1-\lambda)q_n)$.

Claim 3

$H(\tau_\lambda) \geq \lambda H(p) + (1-\lambda)H(q)$

Proof:

- Let Y over $\{0,1\}$ be 1 wp λ
- Let X be distributed according to p if $Y = 0$ and according to q otherwise.
- $H(\tau_\lambda) = H(X) \geq H(X \mid Y)$

Concavity of entropy function

Let $p = (p_1, \ldots, p_n)$ and $q = (q_1, \ldots, q_n)$ be two distributions, and for $\lambda \in [0,1]$ consider the distribution $\tau_\lambda = \lambda p + (1-\lambda)q$.
(i.e., $\tau_\lambda = (\lambda p_1 + (1-\lambda)q_1, \ldots, \lambda p_n + (1-\lambda)q_n)$.

Claim 3

$H(\tau_\lambda) \geq \lambda H(p) + (1-\lambda)H(q)$

Proof:

- Let Y over $\{0,1\}$ be 1 wp λ
- Let X be distributed according to p if $Y = 0$ and according to q otherwise.
- $H(\tau_\lambda) = H(X) \geq H(X \mid Y) = \lambda H(p) + (1-\lambda)H(q)$

Concavity of entropy function

Let $p = (p_1, \ldots, p_n)$ and $q = (q_1, \ldots, q_n)$ be two distributions, and for $\lambda \in [0,1]$ consider the distribution $\tau_\lambda = \lambda p + (1-\lambda)q$.
(i.e., $\tau_\lambda = (\lambda p_1 + (1-\lambda)q_1, \ldots, \lambda p_n + (1-\lambda)q_n)$).

Claim 3

$H(\tau_\lambda) \geq \lambda H(p) + (1-\lambda)H(q)$

Proof:

- Let Y over $\{0,1\}$ be 1 wp λ
- Let X be distributed according to p if $Y = 0$ and according to q otherwise.
- $H(\tau_\lambda) = H(X) \geq H(X \mid Y) = \lambda H(p) + (1-\lambda)H(q)$

We are now certain that we drew the graph of the (two-dimensional) entropy function right...

Part II

Mutual Information

Mutual information

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$I(X;Y) := H(Y) - H(Y|X)$$

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$I(X;Y) := H(Y) - H(Y|X)$$

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$\begin{aligned} I(X;Y) &:= H(Y) - H(Y|X) \\ &= H(Y) - (H(X,Y) - H(X)) \end{aligned}$$

Mutual information

- $I(X; Y)$ — the "information" that X gives on Y
- $$\begin{aligned} I(X;Y) &:= H(Y) - H(Y|X) \\ &= H(Y) - (H(X,Y) - H(X)) \\ &= H(X) + H(Y) - H(X,Y) \end{aligned}$$

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$\begin{aligned} I(X;Y) &:= H(Y) - H(Y|X) \\ &= H(Y) - (H(X,Y) - H(X)) \\ &= H(X) + H(Y) - H(X,Y) \\ &= H(Y) - H(X|Y) \end{aligned}$$

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$\begin{aligned} I(X;Y) &:= H(Y) - H(Y|X) \\ &= H(Y) - (H(X,Y) - H(X)) \\ &= H(X) + H(Y) - H(X,Y) \\ &= H(Y) - H(X|Y) \\ &= I(Y;X). \end{aligned}$$

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$\begin{aligned} I(X;Y) &:= H(Y) - H(Y|X) \\ &= H(Y) - (H(X,Y) - H(X)) \\ &= H(X) + H(Y) - H(X,Y) \\ &= H(Y) - H(X|Y) \\ &= I(Y;X). \end{aligned}$$
- The mutual information that X gives about Y equals the mutual information that Y gives about X.

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$\begin{aligned} I(X;Y) &:= H(Y) - H(Y|X) \\ &= H(Y) - (H(X,Y) - H(X)) \\ &= H(X) + H(Y) - H(X,Y) \\ &= H(Y) - H(X|Y) \\ &= I(Y;X). \end{aligned}$$
- The mutual information that X gives about Y equals the mutual information that Y gives about X.
- $I(X;Y) \geq 0$.

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$\begin{aligned} I(X;Y) &:= H(Y) - H(Y|X) \\ &= H(Y) - (H(X,Y) - H(X)) \\ &= H(X) + H(Y) - H(X,Y) \\ &= H(Y) - H(X|Y) \\ &= I(Y;X). \end{aligned}$$
- The mutual information that X gives about Y equals the mutual information that Y gives about X.
- $I(X;Y) \geq 0$.

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$\begin{aligned} I(X;Y) &:= H(Y) - H(Y|X) \\ &= H(Y) - (H(X,Y) - H(X)) \\ &= H(X) + H(Y) - H(X,Y) \\ &= H(Y) - H(X|Y) \\ &= I(Y;X). \end{aligned}$$
- The mutual information that X gives about Y equals the mutual information that Y gives about X.
- $I(X;Y) \geq 0$. When 0?

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$\begin{aligned} I(X;Y) &:= H(Y) - H(Y|X) \\ &= H(Y) - (H(X,Y) - H(X)) \\ &= H(X) + H(Y) - H(X,Y) \\ &= H(Y) - H(X|Y) \\ &= I(Y;X). \end{aligned}$$
- The mutual information that X gives about Y equals the mutual information that Y gives about X.
- $I(X;Y) \geq 0$. When 0?
- $I(X;X) = H(X)$

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$\begin{aligned} I(X;Y) &:= H(Y) - H(Y|X) \\ &= H(Y) - (H(X,Y) - H(X)) \\ &= H(X) + H(Y) - H(X,Y) \\ &= H(Y) - H(X|Y) \\ &= I(Y;X). \end{aligned}$$
- The mutual information that X gives about Y equals the mutual information that Y gives about X.
- $I(X;Y) \geq 0$. When 0?
- $I(X;X) = H(X)$
- $I(X;f(X)) = H(f(X))$ (and smaller then $H(X)$ if f is non-injective)

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$\begin{aligned} I(X;Y) &:= H(Y) - H(Y|X) \\ &= H(Y) - (H(X,Y) - H(X)) \\ &= H(X) + H(Y) - H(X,Y) \\ &= H(Y) - H(X|Y) \\ &= I(Y;X). \end{aligned}$$
- The mutual information that X gives about Y equals the mutual information that Y gives about X.
- $I(X;Y) \geq 0$. When 0?
- $I(X;X) = H(X)$
- $I(X;f(X)) = H(f(X))$ (and smaller then $H(X)$ if f is non-injective)
- $I(X;Y,Z) \geq I(X;Y), I(X;Z)$ (since $H(X \mid Y,Z) \leq H(X \mid Y), H(X \mid Z)$)

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$\begin{aligned} I(X;Y) &:= H(Y) - H(Y|X) \\ &= H(Y) - (H(X,Y) - H(X)) \\ &= H(X) + H(Y) - H(X,Y) \\ &= H(Y) - H(X|Y) \\ &= I(Y;X). \end{aligned}$$
- The mutual information that X gives about Y equals the mutual information that Y gives about X.
- $I(X;Y) \geq 0$. When 0?
- $I(X;X) = H(X)$
- $I(X;f(X)) = H(f(X))$ (and smaller then $H(X)$ if f is non-injective)
- $I(X;Y,Z) \geq I(X;Y), I(X;Z)$ (since $H(X \mid Y,Z) \leq H(X \mid Y), H(X \mid Z)$)
- $I(X;Y|Z) := H(Y|Z) - H(Y|X,Z)$ ≥ 0

Mutual information

- $I(X;Y)$ — the "information" that X gives on Y
- $$\begin{aligned} I(X;Y) &:= H(Y) - H(Y|X) \\ &= H(Y) - (H(X,Y) - H(X)) \\ &= H(X) + H(Y) - H(X,Y) \\ &= H(Y) - H(X|Y) \\ &= I(Y;X). \end{aligned}$$
- The mutual information that X gives about Y equals the mutual information that Y gives about X.
- $I(X;Y) \geq 0$. When 0?
- $I(X;X) = H(X)$
- $I(X;f(X)) = H(f(X))$ (and smaller then $H(X)$ if f is non-injective)
- $I(X;Y,Z) \geq I(X;Y), I(X;Z)$ (since $H(X \mid Y,Z) \leq H(X \mid Y), H(X \mid Z)$)
- $I(X;Y|Z) := H(Y|Z) - H(Y|X,Z)$ ≥ 0
- $I(X;Y|Z) = I(Y;X|Z)$ (since $I(X';Y') = I(Y';X')$)

Numerical example

Numerical example

- Example

$X \backslash Y$	0	1
0	$\frac{1}{4}$	$\frac{1}{4}$
1	$\frac{1}{2}$	0

Numerical example

- Example

$X \backslash Y$	0	1
0	$\frac{1}{4}$	$\frac{1}{4}$
1	$\frac{1}{2}$	0

$$I(X;Y) = H(X) - H(X|Y)$$

Numerical example

- Example

$X \backslash Y$	0	1
0	$\frac{1}{4}$	$\frac{1}{4}$
1	$\frac{1}{2}$	0

$$I(X;Y) = H(X) - H(X|Y)$$

Numerical example

- Example

$X \backslash Y$	0	1
0	$\frac{1}{4}$	$\frac{1}{4}$
1	$\frac{1}{2}$	0

$$\begin{aligned} I(X;Y) &= H(X) - H(X|Y) \\ &= 1 - \frac{3}{4} \cdot h(\frac{1}{3}) \end{aligned}$$

Numerical example

- Example

$X \backslash Y$	0	1
0	$\frac{1}{4}$	$\frac{1}{4}$
1	$\frac{1}{2}$	0

$$\begin{aligned} I(X;Y) &= H(X) - H(X|Y) \\ &= 1 - \frac{3}{4} \cdot h(\frac{1}{3}) \\ &= I(Y;X) \end{aligned}$$

Numerical example

- Example

$X \backslash Y$	0	1
0	$\frac{1}{4}$	$\frac{1}{4}$
1	$\frac{1}{2}$	0

$$\begin{aligned} I(X;Y) &= H(X) - H(X|Y) \\ &= 1 - \frac{3}{4} \cdot h(\frac{1}{3}) \\ &= I(Y;X) \\ &= H(Y) - H(Y|X) \end{aligned}$$

Numerical example

- Example

$X \backslash Y$	0	1
0	$\frac{1}{4}$	$\frac{1}{4}$
1	$\frac{1}{2}$	0

$$
\begin{aligned}
I(X;Y) &= H(X) - H(X|Y) \\
&= 1 - \frac{3}{4} \cdot h(\frac{1}{3}) \\
&= I(Y;X) \\
&= H(Y) - H(Y|X) \\
&= h(\frac{1}{4}) - \frac{1}{2} h(\frac{1}{2})
\end{aligned}
$$

Chain rule for mutual information

Chain rule for mutual information

Claim 4 (Chain rule for mutual information)

For rvs $X_1, \ldots, X_k, Y$, it holds that
$I(X_1, \ldots, X_k; Y) = I(X; Y) + I(X_2; Y|X_1) + \ldots + I(X_k; Y|X_1, \ldots, X_{k-1})$.

Chain rule for mutual information

Claim 4 (Chain rule for mutual information)

For rvs $X_1, \ldots, X_k, Y$, it holds that
$I(X_1, \ldots, X_k; Y) = I(X; Y) + I(X_2; Y|X_1) + \ldots + I(X_k; Y|X_1, \ldots, X_{k-1})$.

Proof: ?

Chain rule for mutual information

Claim 4 (Chain rule for mutual information)

For rvs $X_1, \ldots, X_k, Y$, it holds that
$I(X_1, \ldots, X_k; Y) = I(X; Y) + I(X_2; Y|X_1) + \ldots + I(X_k; Y|X_1, \ldots, X_{k-1})$.

Proof: ? HW

Examples

- Let $X_1, \ldots, X_n$ be iid with $X_i \sim (p, 1-p)$, under the condition that $\bigoplus_i x_i = 0$.
  Compute $I(X_1, \ldots, X_{n-1}; X_n)$.

Examples

- Let $X_1, \ldots, X_n$ be iid with $X_i \sim (p, 1-p)$, under the condition that $\bigoplus_i x_i = 0$.

  Compute $I(X_1, \ldots, X_{n-1}; X_n)$.

  By chain rule

$$
\begin{aligned}
&I(X_1, \ldots, X_{n-1}; X_n) \\
&= H(X_1; X_n) + H(X_2; X_n | X_1) + \ldots + H(X_{n-1}; X_n | X_1, \ldots, X_{n-2})
\end{aligned}
$$

Examples

- Let $X_1, \ldots, X_n$ be iid with $X_i \sim (p, 1-p)$, under the condition that $\bigoplus_i x_i = 0$.

  Compute $I(X_1, \ldots, X_{n-1}; X_n)$.

  By chain rule

$$
\begin{aligned}
& I(X_1, \ldots, X_{n-1}; X_n) \\
& = H(X_1; X_n) + H(X_2; X_n|X_1) + \ldots + H(X_{n-1}; X_n|X_1, \ldots, X_{n-2})
\end{aligned}
$$

Examples

- Let $X_1, \ldots, X_n$ be iid with $X_i \sim (p, 1-p)$, under the condition that $\bigoplus_i x_i = 0$.
  Compute $I(X_1, \ldots, X_{n-1}; X_n)$.

  By chain rule

$$
\begin{aligned}
& I(X_1, \ldots, X_{n-1}; X_n) \\
& = H(X_1; X_n) + H(X_2; X_n | X_1) + \ldots + H(X_{n-1}; X_n | X_1, \ldots, X_{n-2}) \\
& = 0 + 0 + \ldots + 1 = 1.
\end{aligned}
$$

Examples

- Let $X_1, \ldots, X_n$ be iid with $X_i \sim (p, 1-p)$, under the condition that $\bigoplus_i x_i = 0$.
  Compute $I(X_1, \ldots, X_{n-1}; X_n)$.

  By chain rule

  $$\begin{aligned}
  &I(X_1, \ldots, X_{n-1}; X_n) \\
  &= H(X_1; X_n) + H(X_2; X_n|X_1) + \ldots + H(X_{n-1}; X_n|X_1, \ldots, X_{n-2}) \\
  &= 0 + 0 + \ldots + 1 = 1.
  \end{aligned}$$

- Let T and F be the top and front side, respectively, of a 6-sided fair dice. Compute $I(T; F)$.

Examples

- Let $X_1, \ldots, X_n$ be iid with $X_i \sim (p, 1-p)$, under the condition that $\bigoplus_i x_i = 0$.
  Compute $I(X_1, \ldots, X_{n-1}; X_n)$.

  By chain rule

  $$\begin{aligned} & I(X_1, \ldots, X_{n-1}; X_n) \\ & = H(X_1; X_n) + H(X_2; X_n | X_1) + \ldots + H(X_{n-1}; X_n | X_1, \ldots, X_{n-2}) \\ & = 0 + 0 + \ldots + 1 = 1. \end{aligned}$$

- Let T and F be the top and front side, respectively, of a 6-sided fair dice. Compute $I(T; F)$.

  $$I(T; F) = H(T) - H(T|F)$$

Examples

- Let $X_1, \ldots, X_n$ be iid with $X_i \sim (p, 1-p)$, under the condition that $\bigoplus_i x_i = 0$.

  Compute $I(X_1, \ldots, X_{n-1}; X_n)$.

  By chain rule

$$
\begin{aligned}
&I(X_1, \ldots, X_{n-1}; X_n) \\
&= H(X_1; X_n) + H(X_2; X_n|X_1) + \ldots + H(X_{n-1}; X_n|X_1, \ldots, X_{n-2}) \\
&= 0 + 0 + \ldots + 1 = 1.
\end{aligned}
$$

- Let T and F be the top and front side, respectively, of a 6-sided fair dice.

  Compute $I(T; F)$.

$$I(T; F) = H(T) - H(T|F)$$

Examples

- Let $X_1, \ldots, X_n$ be iid with $X_i \sim (p, 1-p)$, under the condition that $\bigoplus_i x_i = 0$.
  Compute $I(X_1, \ldots, X_{n-1}; X_n)$.

  By chain rule

$$\begin{aligned}
&I(X_1, \ldots, X_{n-1}; X_n) \\
&= H(X_1; X_n) + H(X_2; X_n|X_1) + \ldots + H(X_{n-1}; X_n|X_1, \ldots, X_{n-2}) \\
&= 0 + 0 + \ldots + 1 = 1.
\end{aligned}$$

- Let T and F be the top and front side, respectively, of a 6-sided fair dice. Compute $I(T; F)$.

$$\begin{aligned}
I(T; F) &= H(T) - H(T|F) \\
&= \log 6 - \log 4
\end{aligned}$$

Examples

- Let $X_1, \ldots, X_n$ be iid with $X_i \sim (p, 1-p)$, under the condition that $\bigoplus_i x_i = 0$.
  Compute $I(X_1, \ldots, X_{n-1}; X_n)$.

  By chain rule

$$
\begin{aligned}
& I(X_1, \ldots, X_{n-1}; X_n) \\
& = H(X_1; X_n) + H(X_2; X_n | X_1) + \ldots + H(X_{n-1}; X_n | X_1, \ldots, X_{n-2}) \\
& = 0 + 0 + \ldots + 1 = 1.
\end{aligned}
$$

- Let T and F be the top and front side, respectively, of a 6-sided fair dice. Compute $I(T; F)$.

$$
\begin{aligned}
I(T; F) &= H(T) - H(T|F) \\
&= \log 6 - \log 4 \\
&= \log 3 - 1.
\end{aligned}
$$

Part III

Data processing

Data processing Inequality

Definition 5 (Markov Chain)

Rvs $(X, Y, Z) \sim p$ form a Markov chain, denoted $X \to Y \to Z$, if $p(x, y, z) = p_X(x) \cdot p_{Y|X}(y|x) \cdot p_{Z|Y}(z|y)$, for all x, y, z.

Data processing Inequality

Definition 5 (Markov Chain)

Rvs $(X, Y, Z) \sim p$ form a Markov chain, denoted $X \to Y \to Z$, if $p(x, y, z) = p_X(x) \cdot p_{Y|X}(y|x) \cdot p_{Z|Y}(z|y)$, for all x, y, z.

Example: random walk on graph.

Data processing Inequality

Definition 5 (Markov Chain)

Rvs $(X, Y, Z) \sim p$ form a Markov chain, denoted $X \to Y \to Z$, if $p(x, y, z) = p_X(x) \cdot p_{Y|X}(y|x) \cdot p_{Z|Y}(z|y)$, for all x, y, z.

Example: random walk on graph.

Claim 6

If $X \to Y \to Z$, then $I(X; Y) \geq I(X; Z)$.

Data processing Inequality

Definition 5 (Markov Chain)

Rvs $(X, Y, Z) \sim p$ form a Markov chain, denoted $X \to Y \to Z$, if $p(x, y, z) = p_X(x) \cdot p_{Y|X}(y|x) \cdot p_{Z|Y}(z|y)$, for all x, y, z.

Example: random walk on graph.

Claim 6

If $X \to Y \to Z$, then $I(X; Y) \geq I(X; Z)$.

- By Chain rule, $I(X; Y, Z) = I(X; Z) + I(X; Y|Z) = I(X; Y) + I(X; Z|Y)$.

Data processing Inequality

Definition 5 (Markov Chain)

Rvs $(X, Y, Z) \sim p$ form a Markov chain, denoted $X \to Y \to Z$, if $p(x, y, z) = p_X(x) \cdot p_{Y|X}(y|x) \cdot p_{Z|Y}(z|y)$, for all x, y, z.

Example: random walk on graph.

Claim 6

If $X \to Y \to Z$, then $I(X; Y) \geq I(X; Z)$.

- By Chain rule, $I(X; Y, Z) = I(X; Z) + I(X; Y|Z) = I(X; Y) + I(X; Z|Y)$.
- $I(X; Z|Y) = 0$

Data processing Inequality

Definition 5 (Markov Chain)

Rvs $(X, Y, Z) \sim p$ form a Markov chain, denoted $X \to Y \to Z$, if $p(x, y, z) = p_X(x) \cdot p_{Y|X}(y|x) \cdot p_{Z|Y}(z|y)$, for all x, y, z.

Example: random walk on graph.

Claim 6

If $X \to Y \to Z$, then $I(X; Y) \geq I(X; Z)$.

- By Chain rule, $I(X; Y, Z) = I(X; Z) + I(X; Y|Z) = I(X; Y) + I(X; Z|Y)$.
- $I(X; Z|Y) = 0$

Data processing Inequality

Definition 5 (Markov Chain)

Rvs $(X, Y, Z) \sim p$ form a Markov chain, denoted $X \to Y \to Z$, if $p(x, y, z) = p_X(x) \cdot p_{Y|X}(y|x) \cdot p_{Z|Y}(z|y)$, for all x, y, z.

Example: random walk on graph.

Claim 6

If $X \to Y \to Z$, then $I(X; Y) \geq I(X; Z)$.

- By Chain rule, $I(X; Y, Z) = I(X; Z) + I(X; Y|Z) = I(X; Y) + I(X; Z|Y)$.
- $I(X; Z|Y) = 0$
  - $p_{Z|Y=y} = p_{Z|Y=y,X=x}$ for any x, y

Data processing Inequality

Definition 5 (Markov Chain)

Rvs $(X, Y, Z) \sim p$ form a Markov chain, denoted $X \to Y \to Z$, if $p(x, y, z) = p_X(x) \cdot p_{Y|X}(y|x) \cdot p_{Z|Y}(z|y)$, for all x, y, z.

Example: random walk on graph.

Claim 6

If $X \to Y \to Z$, then $I(X; Y) \geq I(X; Z)$.

- By Chain rule, $I(X; Y, Z) = I(X; Z) + I(X; Y|Z) = I(X; Y) + I(X; Z|Y)$.
- $I(X; Z|Y) = 0$
  - $p_{Z|Y=y} = p_{Z|Y=y,X=x}$ for any x, y
  - $$I(X; Z|Y) = H(Z|Y) - H(Z|Y, X)$$

Data processing Inequality

Definition 5 (Markov Chain)

Rvs $(X, Y, Z) \sim p$ form a Markov chain, denoted $X \to Y \to Z$, if $p(x, y, z) = p_X(x) \cdot p_{Y|X}(y|x) \cdot p_{Z|Y}(z|y)$, for all x, y, z.

Example: random walk on graph.

Claim 6

If $X \to Y \to Z$, then $I(X; Y) \geq I(X; Z)$.

- By Chain rule, $I(X; Y, Z) = I(X; Z) + I(X; Y|Z) = I(X; Y) + I(X; Z|Y)$.
- $I(X; Z|Y) = 0$
  - $p_{Z|Y=y} = p_{Z|Y=y,X=x}$ for any x, y
  - $$I(X; Z|Y) = H(Z|Y) - H(Z|Y, X)$$

Data processing Inequality

Definition 5 (Markov Chain)

Rvs $(X, Y, Z) \sim p$ form a Markov chain, denoted $X \to Y \to Z$, if $p(x, y, z) = p_X(x) \cdot p_{Y|X}(y|x) \cdot p_{Z|Y}(z|y)$, for all x, y, z.

Example: random walk on graph.

Claim 6

If $X \to Y \to Z$, then $I(X; Y) \geq I(X; Z)$.

- By Chain rule, $I(X; Y, Z) = I(X; Z) + I(X; Y|Z) = I(X; Y) + I(X; Z|Y)$.
- $I(X; Z|Y) = 0$
  - $p_{Z|Y=y} = p_{Z|Y=y,X=x}$ for any x, y
  - $$\begin{aligned} I(X; Z|Y) &= H(Z|Y) - H(Z|Y, X) \\ &= \mathop{\mathrm{E}}_{Y} H(p_{Z|Y=y}) - \mathop{\mathrm{E}}_{Y,X} H(p_{Z|Y=y,X=x}) \end{aligned}$$

Data processing Inequality

Definition 5 (Markov Chain)

Rvs $(X, Y, Z) \sim p$ form a Markov chain, denoted $X \to Y \to Z$, if $p(x, y, z) = p_X(x) \cdot p_{Y|X}(y|x) \cdot p_{Z|Y}(z|y)$, for all x, y, z.

Example: random walk on graph.

Claim 6

If $X \to Y \to Z$, then $I(X; Y) \geq I(X; Z)$.

- By Chain rule, $I(X; Y, Z) = I(X; Z) + I(X; Y|Z) = I(X; Y) + I(X; Z|Y)$.
- $I(X; Z|Y) = 0$
  - $p_{Z|Y=y} = p_{Z|Y=y,X=x}$ for any x, y
  - $$\begin{aligned} I(X; Z|Y) &= H(Z|Y) - H(Z|Y, X) \\ &= \mathop{\mathrm{E}}_{Y} H(p_{Z|Y=y}) - \mathop{\mathrm{E}}_{Y,X} H(p_{Z|Y=y,X=x}) \\ &= \mathop{\mathrm{E}}_{Y} H(p_{Z|Y=y}) - \mathop{\mathrm{E}}_{Y} H(p_{Z|Y=y}) = 0. \end{aligned}$$

Data processing Inequality

Definition 5 (Markov Chain)

Rvs $(X, Y, Z) \sim p$ form a Markov chain, denoted $X \to Y \to Z$, if $p(x, y, z) = p_X(x) \cdot p_{Y|X}(y|x) \cdot p_{Z|Y}(z|y)$, for all x, y, z.

Example: random walk on graph.

Claim 6

If $X \to Y \to Z$, then $I(X; Y) \geq I(X; Z)$.

- By Chain rule, $I(X; Y, Z) = I(X; Z) + I(X; Y|Z) = I(X; Y) + I(X; Z|Y)$.
- $I(X; Z|Y) = 0$
  - $p_{Z|Y=y} = p_{Z|Y=y,X=x}$ for any x, y
  - $$\begin{aligned} I(X; Z|Y) &= H(Z|Y) - H(Z|Y, X) \\ &= \mathop{\mathrm{E}}_{Y} H(p_{Z|Y=y}) - \mathop{\mathrm{E}}_{Y,X} H(p_{Z|Y=y,X=x}) \\ &= \mathop{\mathrm{E}}_{Y} H(p_{Z|Y=y}) - \mathop{\mathrm{E}}_{Y} H(p_{Z|Y=y}) = 0. \end{aligned}$$
- Since $I(X; Y|Z) \geq 0$, we conclude $I(X; Y) \geq I(X; Z)$.$\square$

Fano's Inequality

- How well can we guess X from Y?

Fano's Inequality

- How well can we guess X from Y?
- Could with no error if $H(X|Y) = 0$.

Fano's Inequality

- How well can we guess X from Y?
- Could with no error if $H(X|Y) = 0$.

Fano's Inequality

- How well can we guess X from Y?
- Could with no error if $H(X|Y) = 0$. What if $H(X|Y)$ is small?

Fano's Inequality

- How well can we guess X from Y?
- Could with no error if $H(X|Y) = 0$. What if $H(X|Y)$ is small?

Fano's Inequality

- How well can we guess X from Y?
- Could with no error if $H(X|Y) = 0$. What if $H(X|Y)$ is small?

Theorem 7 (Fano's inequality)

For any rvs X and Y, and any (even random) g, it holds that

$$h(P_e) + P_e \log |\mathcal{X}| \geq H(X|\hat{X}) \geq H(X|Y)$$

for $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

Fano's Inequality

- How well can we guess X from Y?
- Could with no error if $H(X|Y) = 0$. What if $H(X|Y)$ is small?

Theorem 7 (Fano's inequality)

For any rvs X and Y, and any (even random) g, it holds that

$$h(P_e) + P_e \log|\mathcal{X}| \geq H(X|\hat{X}) \geq H(X|Y)$$

for $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

Fano's Inequality

- How well can we guess X from Y?
- Could with no error if $H(X|Y) = 0$. What if $H(X|Y)$ is small?

Theorem 7 (Fano's inequality)

For any rvs X and Y, and any (even random) g, it holds that

$$h(P_e) + P_e \log |\mathcal{X}| \geq H(X|\hat{X}) \geq H(X|Y)$$

for $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Note that $P_e = 0$ implies that $H(X|Y) = 0$

Fano's Inequality

- How well can we guess X from Y?
- Could with no error if $H(X|Y) = 0$. What if $H(X|Y)$ is small?

Theorem 7 (Fano's inequality)

For any rvs X and Y, and any (even random) g, it holds that

$$h(P_e) + P_e \log|\mathcal{X}| \geq H(X|\hat{X}) \geq H(X|Y)$$

for $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Note that $P_e = 0$ implies that $H(X|Y) = 0$
- The inequality can be weekend to $1 + P_e \log|\mathcal{X}| \geq H(X|Y)$,

Fano's Inequality

- How well can we guess X from Y?
- Could with no error if $H(X|Y) = 0$. What if $H(X|Y)$ is small?

Theorem 7 (Fano's inequality)

For any rvs X and Y, and any (even random) g, it holds that

$$h(P_e) + P_e \log |\mathcal{X}| \geq H(X|\hat{X}) \geq H(X|Y)$$

for $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Note that $P_e = 0$ implies that $H(X|Y) = 0$
- The inequality can be weekend to $1 + P_e \log |\mathcal{X}| \geq H(X|Y)$,
- Alternatively, to $P_e \geq \frac{H(X|Y)-1}{\log|\mathcal{X}|}$

Fano's Inequality

- How well can we guess X from Y?
- Could with no error if $H(X|Y) = 0$. What if $H(X|Y)$ is small?

Theorem 7 (Fano's inequality)

For any rvs X and Y, and any (even random) g, it holds that

$$h(P_e) + P_e \log|\mathcal{X}| \geq H(X|\hat{X}) \geq H(X|Y)$$

for $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Note that $P_e = 0$ implies that $H(X|Y) = 0$
- The inequality can be weekend to $1 + P_e \log|\mathcal{X}| \geq H(X|Y)$,
- Alternatively, to $P_e \geq \frac{H(X|Y)-1}{\log|\mathcal{X}|}$
- Intuition for $\propto \frac{1}{\log|\mathcal{X}|}$

Fano's Inequality

- How well can we guess X from Y?
- Could with no error if $H(X|Y) = 0$. What if $H(X|Y)$ is small?

Theorem 7 (Fano's inequality)

For any rvs X and Y, and any (even random) g, it holds that

$$h(P_e) + P_e \log|\mathcal{X}| \geq H(X|\hat{X}) \geq H(X|Y)$$

for $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Note that $P_e = 0$ implies that $H(X|Y) = 0$
- The inequality can be weekend to $1 + P_e \log|\mathcal{X}| \geq H(X|Y)$,
- Alternatively, to $P_e \geq \frac{H(X|Y)-1}{\log|\mathcal{X}|}$
- Intuition for $\propto \frac{1}{\log|\mathcal{X}|}$
- We call $\hat{X}$ an estimator for X (from Y).

Proving Fano's inequality

Let X and Y be rvs, let $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

Proving Fano's inequality

Let X and Y be rvs, let $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Let $E = \begin{cases} 1, & \hat{X} \neq X \\ 0, & \hat{X} = X. \end{cases}$

Proving Fano's inequality

Let X and Y be rvs, let $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Let $E = \begin{cases} 1, & \hat{X} \neq X \\ 0, & \hat{X} = X. \end{cases}$

$$H(E, X|\hat{X}) = H(X|\hat{X}) + \underbrace{H(E|X, \hat{X})}_{=0}$$

Proving Fano's inequality

Let X and Y be rvs, let $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Let $E = \begin{cases} 1, & \hat{X} \neq X \\ 0, & \hat{X} = X. \end{cases}$

$$H(E, X|\hat{X}) = H(X|\hat{X}) + \underbrace{H(E|X,\hat{X})}_{=0}$$

Proving Fano's inequality

Let X and Y be rvs, let $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Let $E = \begin{cases} 1, & \hat{X} \neq X \\ 0, & \hat{X} = X. \end{cases}$

$$\begin{aligned} H(E, X|\hat{X}) &= H(X|\hat{X}) + \underbrace{H(E|X, \hat{X})}_{=0} \\ &= \underbrace{H(E|\hat{X})}_{\leq H(E) = h(P_e)} + \underbrace{H(X|E, \hat{X})}_{\leq P_e \log|\mathcal{X}|(?)} \end{aligned}$$

Proving Fano's inequality

Let X and Y be rvs, let $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Let $E = \begin{cases} 1, & \hat{X} \neq X \\ 0, & \hat{X} = X. \end{cases}$

$$\begin{aligned} H(E, X|\hat{X}) &= H(X|\hat{X}) + \underbrace{H(E|X, \hat{X})}_{=0} \\ &= \underbrace{H(E|\hat{X})}_{\leq H(E) = h(P_e)} + \underbrace{H(X|E, \hat{X})}_{\leq P_e \log|\mathcal{X}|(?)} \end{aligned}$$

- It follows that $h(P_e) + P_e \log|\mathcal{X}| \geq H(X|\hat{X})$

Proving Fano's inequality

Let X and Y be rvs, let $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Let $E = \begin{cases} 1, & \hat{X} \neq X \\ 0, & \hat{X} = X. \end{cases}$

$$\begin{aligned} H(E, X|\hat{X}) &= H(X|\hat{X}) + \underbrace{H(E|X, \hat{X})}_{=0} \\ &= \underbrace{H(E|\hat{X})}_{\leq H(E) = h(P_e)} + \underbrace{H(X|E, \hat{X})}_{\leq P_e \log|\mathcal{X}|(?)} \end{aligned}$$

- It follows that $h(P_e) + P_e \log|\mathcal{X}| \geq H(X|\hat{X})$
- Since $X \to Y \to \hat{X}$,

Proving Fano's inequality

Let X and Y be rvs, let $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Let $E = \begin{cases} 1, & \hat{X} \neq X \\ 0, & \hat{X} = X. \end{cases}$

$$\begin{aligned} H(E, X|\hat{X}) &= H(X|\hat{X}) + \underbrace{H(E|X, \hat{X})}_{=0} \\ &= \underbrace{H(E|\hat{X})}_{\leq H(E) = h(P_e)} + \underbrace{H(X|E, \hat{X})}_{\leq P_e \log|\mathcal{X}|(?)} \end{aligned}$$

- It follows that $h(P_e) + P_e \log|\mathcal{X}| \geq H(X|\hat{X})$
- Since $X \to Y \to \hat{X}$,

Proving Fano's inequality

Let X and Y be rvs, let $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Let $E = \begin{cases} 1, & \hat{X} \neq X \\ 0, & \hat{X} = X. \end{cases}$

$$\begin{aligned} H(E, X|\hat{X}) &= H(X|\hat{X}) + \underbrace{H(E|X,\hat{X})}_{=0} \\ &= \underbrace{H(E|\hat{X})}_{\leq H(E) = h(P_e)} + \underbrace{H(X|E,\hat{X})}_{\leq P_e \log|\mathcal{X}|(?)} \end{aligned}$$

- It follows that $h(P_e) + P_e \log|\mathcal{X}| \geq H(X|\hat{X})$
- Since $X \rightarrow Y \rightarrow \hat{X}$, it holds that $I(X;Y) \geq I(X;\hat{X})$

Proving Fano's inequality

Let X and Y be rvs, let $\hat{X} = g(Y)$ and $P_e = \Pr\left[\hat{X} \neq X\right]$.

- Let $E = \begin{cases} 1, & \hat{X} \neq X \\ 0, & \hat{X} = X. \end{cases}$

$$\begin{aligned} H(E, X|\hat{X}) &= H(X|\hat{X}) + \underbrace{H(E|X, \hat{X})}_{=0} \\ &= \underbrace{H(E|\hat{X})}_{\leq H(E) = h(P_e)} + \underbrace{H(X|E, \hat{X})}_{\leq P_e \log|\mathcal{X}|(?)} \end{aligned}$$

- It follows that $h(P_e) + P_e \log|\mathcal{X}| \geq H(X|\hat{X})$
- Since $X \to Y \to \hat{X}$, it holds that $I(X;Y) \geq I(X;\hat{X})$
  $\implies H(X|\hat{X}) \geq H(X|Y)$