Information Theory, Fall 2014	Iftach Haitner
Problem se	et 3
December 9, 2014	Due: December 23

- Please submit the handout in class, or email the grader.
- Write clearly and shortly using sub-claims if needed. The emphasize in most questions is on the proofs (no much point is writing a "solution" w/o proving its correctness)
- For Latex users, a solution example can be found in the course web site.
- It is allowed to work in (small) groups, but please write the id list of your partners in the solution file, and each student should write his solution by *himself* (joint effort is only allowed in the "thinking phase")

- 1. Let $Q: \{0,1\} \mapsto \{0,1\} \cup \{\bot\}$ be the random function with $\Pr\left[Q(x) = \bot\right] = p$ and $\Pr\left[Q(x) = x\right] = 1 p$ for any $x \in \{0,1\}$. Find the capacity of the channel described by Q?
 - That is, find the right value of C_p for which the natural adjustment of Shannon's theorem (Theorem 1 in lecture 5) for the noise model described by Q (i.e., Q is applied independently to each transmitted bit) can be proven.
- 2. Let G be the graph with set of nodes $\{0,1,2\}^n$, where two nodes $(x,y) \in \{0,1,2\}^n$ are connected by an edge iff they differ in exactly one coordinate. State and prove an isoperimetric inequality for the graph G. (Similar to the isoperimetric inequality for the hyper-cube we did in class).
- 3. Prove or give a counter example: For every rv's X_1, X_2, X_3, X_4 :

$$H(X_1, X_2, X_3) + H(X_2, X_3, X_4) + H(X_3, X_4, X_1) + H(X_4, X_1, X_2)$$

$$\leq \frac{3}{2} [H(X_1, X_2) + H(X_2, X_3) + H(X_3, X_4) + H(X_4, X_1)]$$

- 4. Show that $SD(p,q) = \max_{S \subseteq [m]} \left(\sum_{i \in S} p_i \sum_{i \in S} q_i \right)$ for any two distributions p,q over [m].
- 5. Relative entropy is not symmetric: given two distributions p, q such that $D(p||q) \neq D(q||p)$, and $D(p||q), D(q||p) < \infty$.
- 6. Relative entropy does not obey the triangle inequality: give three distributions p_1, p_2, p_3 such that $D(p_1||p_2) + D(p_2||p_3) < D(p_1||p_3)$
- 7. Relative entropy is non-negative: given two distributions p, q, show that $D(p||q) \ge 0$, with equality only if p = q.
- 8. Does Theorem 1 in Lecture 7 hold for any prefix code C with $E_{i \leftarrow q}[|C(i)|] \le H(q) + 1$? (and not only for a code C with $C(x) = \left\lceil \log \frac{1}{q(x)} \right\rceil$, as stated in the theorem)