

CS275 – Intro to Databases

Relational Algebra - *Chap.6.3-6.5*

SPECIAL RELATIONAL OPERATORS

- Join operators

There are several kind of join operators.

- (1) Condition Joins

- (2) Equijoins

- (3) Natural Joins

- Outer Join

- Division

Conditional join

- Conditional joins
- Defined as a cross-product followed by a selection

$$R \bowtie_c S = \sigma_c(R \times S)$$

(\bowtie is called the bow-tie) where c is the condition.

Conditional join

Example:

- Given the sample relational instances S1 and R1

| <i>sid</i> | <i>sname</i> | <i>rating</i> | <i>age</i> |
|------------|--------------|---------------|------------|
| 22 | Dustin | 7 | 45.0 |
| 31 | Lubber | 8 | 55.5 |
| 58 | Rusty | 10 | 35.0 |

Figure 4.1 Instance S1 of Sailors

| <i>sid</i> | <i>bid</i> | <i>day</i> |
|------------|------------|------------|
| 22 | 101 | 10/10/96 |
| 58 | 103 | 11/12/96 |

Figure 4.3 Instance R1 of Reserves

The condition join $S1 \bowtie_{S1.sid < R1.sid} R1$ yields

| <i>(sid)</i> | <i>sname</i> | <i>rating</i> | <i>age</i> | <i>(sid)</i> | <i>bid</i> | <i>day</i> |
|--------------|--------------|---------------|------------|--------------|------------|------------|
| 22 | Dustin | 7 | 45.0 | 58 | 103 | 11/12/96 |
| 31 | Lubber | 8 | 55.5 | 58 | 103 | 11/12/96 |

Figure 4.12 $S1 \bowtie_{S1.sid < R1.sid} R1$

Equi Join

- Equijoins

$$R \bowtie_c S = \sigma_c(R \times S)$$

- Example

- S1 Sailors(sid:integer, sname:string, rating:integer, age:real)
- B1 Boats(bid:integer, bname: string, color: string)
- R1 Reserves(sid:integer, bid: integer, day: date)

$$S_1 \bowtie_{S_1.sid=R_1.sid} R_1$$

Equi Join

Equi-join: example

- $R \quad S \quad R \text{ equi-join}_{NR=NS} S$

| NR | Tel |
|------|------|
| Joe | 1234 |
| Jill | 1244 |
| Bill | 2244 |

| NS | Addr |
|------|------|
| Joe | NG7 |
| Jill | NG8 |
| Bob | NG1 |

| NR | Tel | NS | Addr |
|------|------|------|------|
| Joe | 1234 | Joe | NG7 |
| Jill | 1244 | Jill | NG8 |

Natural Join

- Natural joins

$$R \bowtie_c S = \sigma_c(R \times S)$$

- Example

- S1 Sailors(sid:integer, sname:string, rating:integer, age:real)
- B1 Boats(bid:integer, bname: string, color: string)
- R1 Reserves(sid:integer, bid: integer, day: date)

$$S_1 \bowtie R_1$$

or

$$S_1 * R_1$$

Natural Join

Natural join: example

- $R \bowtie_{NR=NS} S$

| NR | Tel |
|------|------|
| Joe | 1234 |
| Jill | 1244 |
| Bill | 2244 |

| NS | Addr |
|------|------|
| Joe | NG7 |
| Jill | NG8 |
| Bob | NG1 |

| NR | Tel | Addr |
|------|------|------|
| Joe | 1234 | NG7 |
| Jill | 1244 | NG8 |

Natural join is definable

- Natural join is a very useful operator, but it can be defined using \times , π and σ .
- To define equi-join of R and S over attributes X and Y we take a Cartesian product of R and S and then select tuples which satisfy $X = Y$.
- To obtain natural join, we project the resulting relation on the set of all attributes apart from Y. Let A be the set of all attributes in R and S:

$$R \bowtie_{X=Y} S = \pi_{A - Y} \sigma_{X=Y} (R \times S)$$

Joins

- More examples

| <i>sid</i> | <i>sname</i> | <i>rating</i> | <i>age</i> |
|------------|--------------|---------------|------------|
| 22 | Dustin | 7 | 45.0 |
| 29 | Brutus | 1 | 33.0 |
| 31 | Lubber | 8 | 55.5 |
| 32 | Andy | 8 | 25.5 |
| 58 | Rusty | 10 | 35.0 |
| 64 | Horatio | 7 | 35.0 |
| 71 | Zorba | 10 | 16.0 |
| 74 | Horatio | 9 | 35.0 |
| 85 | Art | 3 | 25.5 |
| 95 | Bob | 3 | 63.5 |

Figure 4.15 An Instance *S3* of Sailors

| <i>sid</i> | <i>bid</i> | <i>day</i> |
|------------|------------|------------|
| 22 | 101 | 10/10/98 |
| 22 | 102 | 10/10/98 |
| 22 | 103 | 10/8/98 |
| 22 | 104 | 10/7/98 |
| 31 | 102 | 11/10/98 |
| 31 | 103 | 11/6/98 |
| 31 | 104 | 11/12/98 |
| 64 | 101 | 9/5/98 |
| 64 | 102 | 9/8/98 |
| 74 | 103 | 9/8/98 |

Figure 4.16 An Instance *R2* of Reserves

- Find the names of sailors who have reserved boat 103

Joins

- More examples

| <i>sid</i> | <i>sname</i> | <i>rating</i> | <i>age</i> |
|------------|--------------|---------------|------------|
| 22 | Dustin | 7 | 45.0 |
| 29 | Brutus | 1 | 33.0 |
| 31 | Lubber | 8 | 55.5 |
| 32 | Andy | 8 | 25.5 |
| 58 | Rusty | 10 | 35.0 |
| 64 | Horatio | 7 | 35.0 |
| 71 | Zorba | 10 | 16.0 |
| 74 | Horatio | 9 | 35.0 |
| 85 | Art | 3 | 25.5 |
| 95 | Bob | 3 | 63.5 |

Figure 4.15 An Instance *S3* of Sailors

| <i>sid</i> | <i>bid</i> | <i>day</i> |
|------------|------------|------------|
| 22 | 101 | 10/10/98 |
| 22 | 102 | 10/10/98 |
| 22 | 103 | 10/8/98 |
| 22 | 104 | 10/7/98 |
| 31 | 102 | 11/10/98 |
| 31 | 103 | 11/6/98 |
| 31 | 104 | 11/12/98 |
| 64 | 101 | 9/5/98 |
| 64 | 102 | 9/8/98 |
| 74 | 103 | 9/8/98 |

Figure 4.16 An Instance *R2* of Reserves

Find the names of sailors who have reserved boat 103

$$\pi_{sname}((\sigma_{bid=103} Reserves) \bowtie Sailors)$$

Ans; {<Dustin>, <Lubber>, <Horatio>}

Joins

Find the names of sailors who have reserved a red boat.

| <i>sid</i> | <i>sname</i> | <i>rating</i> | <i>age</i> |
|------------|--------------|---------------|------------|
| 22 | Dustin | 7 | 45.0 |
| 29 | Brutus | 1 | 33.0 |
| 31 | Lubber | 8 | 55.5 |
| 32 | Andy | 8 | 25.5 |
| 58 | Rusty | 10 | 35.0 |
| 64 | Horatio | 7 | 35.0 |
| 71 | Zorba | 10 | 16.0 |
| 74 | Horatio | 9 | 35.0 |
| 85 | Art | 3 | 25.5 |
| 95 | Bob | 3 | 63.5 |

Figure 4.15 An Instance *S3* of Sailors

| <i>sid</i> | <i>bid</i> | <i>day</i> |
|------------|------------|------------|
| 22 | 101 | 10/10/98 |
| 22 | 102 | 10/10/98 |
| 22 | 103 | 10/8/98 |
| 22 | 104 | 10/7/98 |
| 31 | 102 | 11/10/98 |
| 31 | 103 | 11/6/98 |
| 31 | 104 | 11/12/98 |
| 64 | 101 | 9/5/98 |
| 64 | 102 | 9/8/98 |
| 74 | 103 | 9/8/98 |

Figure 4.16 An Instance *R2* of Reserves

| <i>bid</i> | <i>bname</i> | <i>color</i> |
|------------|--------------|--------------|
| 101 | Interlake | blue |
| 102 | Interlake | red |
| 103 | Clipper | green |
| 104 | Marine | red |

Figure 4.17 An Instance *B1* of Boats

$$\pi_{sname}((\sigma_{color='red'}\mathbf{Boats}) \bowtie \mathbf{Reserves} \bowtie \mathbf{Sailors})$$

Joins

Find the names of Sailors who have reserved at least one boat

| <i>sid</i> | <i>sname</i> | <i>rating</i> | <i>age</i> |
|------------|--------------|---------------|------------|
| 22 | Dustin | 7 | 45.0 |
| 29 | Brutus | 1 | 33.0 |
| 31 | Lubber | 8 | 55.5 |
| 32 | Andy | 8 | 25.5 |
| 58 | Rusty | 10 | 35.0 |
| 64 | Horatio | 7 | 35.0 |
| 71 | Zorba | 10 | 16.0 |
| 74 | Horatio | 9 | 35.0 |
| 85 | Art | 3 | 25.5 |
| 95 | Bob | 3 | 63.5 |

Figure 4.15 An Instance *S3* of Sailors

| <i>sid</i> | <i>bid</i> | <i>day</i> |
|------------|------------|------------|
| 22 | 101 | 10/10/98 |
| 22 | 102 | 10/10/98 |
| 22 | 103 | 10/8/98 |
| 22 | 104 | 10/7/98 |
| 31 | 102 | 11/10/98 |
| 31 | 103 | 11/6/98 |
| 31 | 104 | 11/12/98 |
| 64 | 101 | 9/5/98 |
| 64 | 102 | 9/8/98 |
| 74 | 103 | 9/8/98 |

Figure 4.16 An Instance *R2* of Reserves

| <i>bid</i> | <i>bname</i> | <i>color</i> |
|------------|--------------|--------------|
| 101 | Interlake | blue |
| 102 | Interlake | red |
| 103 | Clipper | green |
| 104 | Marine | red |

Figure 4.17 An Instance *B1* of Boats

$\pi_{sname}(\text{Sailors} \bowtie \text{Reserves})$

Joins

Find the names of sailors who have reserved a red or a green boat.

| <i>sid</i> | <i>sname</i> | <i>rating</i> | <i>age</i> |
|------------|--------------|---------------|------------|
| 22 | Dustin | 7 | 45.0 |
| 29 | Brutus | 1 | 33.0 |
| 31 | Lubber | 8 | 55.5 |
| 32 | Andy | 8 | 25.5 |
| 58 | Rusty | 10 | 35.0 |
| 64 | Horatio | 7 | 35.0 |
| 71 | Zorba | 10 | 16.0 |
| 74 | Horatio | 9 | 35.0 |
| 85 | Art | 3 | 25.5 |
| 95 | Bob | 3 | 63.5 |

Figure 4.15 An Instance *S3* of Sailors

| <i>sid</i> | <i>bid</i> | <i>day</i> |
|------------|------------|------------|
| 22 | 101 | 10/10/98 |
| 22 | 102 | 10/10/98 |
| 22 | 103 | 10/8/98 |
| 22 | 104 | 10/7/98 |
| 31 | 102 | 11/10/98 |
| 31 | 103 | 11/6/98 |
| 31 | 104 | 11/12/98 |
| 64 | 101 | 9/5/98 |
| 64 | 102 | 9/8/98 |
| 74 | 103 | 9/8/98 |

Figure 4.16 An Instance *R2* of Reserves

| <i>bid</i> | <i>bname</i> | <i>color</i> |
|------------|--------------|--------------|
| 101 | Interlake | blue |
| 102 | Interlake | red |
| 103 | Clipper | green |
| 104 | Marine | red |

Figure 4.17 An Instance *B1* of Boats

$$\rho(\text{Tempboats}, (\sigma_{\text{color}='red'}\text{Boats}) \cup (\sigma_{\text{color}='green'}\text{Boats}))$$

$$\pi_{\text{sname}}(\text{Tempboats} \bowtie \text{Reserves} \bowtie \text{Sailors})$$

Joins

Find the names of Sailors who have reserved a red and a green boat.

It seems tempting to use the expression used in Q5, replacing simply \cup by \cap . However, this won't work, for such an expression is requesting the names of sailors who have requested a boat that is both red and green! The correct expression is as follows:

$\rho(\textit{Tempred}, \pi_{\textit{sid}}((\sigma_{\textit{color}='red'}\textit{Boats}) \bowtie \textit{Reserves}))$

$\rho(\textit{Tempgreen}, \pi_{\textit{sid}}((\sigma_{\textit{color}='green'}\textit{Boats}) \bowtie \textit{Reserves}))$

$\pi_{\textit{sname}} ((\textit{Tempred} \cap \textit{Tempgreen}) \bowtie \textit{Sailors})$

Outer Join

- ✿ Join selects only tuples satisfying the join condition
- ✿ *Outer Join:*
 - Left outer join $(r \bowtie_{\infty} s)$ keeps every tuple in first or left relation
 - Right outer join $(r \bowtie_{\infty} s)$ keeps every tuple in second or right relation
 - Full outer join $(r \bowtie_{\infty} s)$ keeps every tuple
- ✿ Attributes of tuples with no matching tuples are set to NULL

Outer Join

relation r

| A | B | C |
|---|---|---|
| a | b | c |
| d | a | f |
| c | b | d |

relation s

| A | B | D |
|---|---|---|
| b | g | a |
| d | a | f |

$r \Join^{\infty} s$

| A | B | C | D |
|---|---|---|------|
| a | b | c | Null |
| d | a | f | f |
| c | b | d | Null |

$r \Join^{\infty} [s$

| A | B | C | D |
|---|---|------|------|
| a | b | c | Null |
| d | a | f | f |
| c | b | d | Null |
| b | g | Null | a |

$r \Join^{\infty} [s$

| A | B | C | D |
|---|---|------|---|
| d | a | f | f |
| b | g | Null | a |

Aggregate Functions

- Mathematical and Statistical aggregate functions on collections of values

- SUM, MAXIMUM, MINIMUM, AVERAGE
- COUNT number of tuples (cardinality)

\bar{f} **<function list>** (<relation>)

- Function list is a list of pairs
(< function, attribute>)

- E.g., $\bar{f}_{\text{count SSN, AVERAGE QPA}}$ (STUDENT)

DIVISION

The division operator is used for queries which involve the 'all' qualifier such as

“Find the names of sailors who have reserved all boats”.

The formal definition of division is as follows:

$$A/B = \pi_X(A) - \pi_X((\pi_X(A) \times B) - A)$$

Division: example

- R

| Lecturer | Module |
|----------|-----------|
| Brown | Compilers |
| Brown | Databases |
| Green | Prolog |
| Green | Databases |
| Lewis | Prolog |
| Smith | Databases |

S

| Subject |
|---------|
| Prolog |

R | S

| Lecturer |
|----------|
| Green |
| Lewis |

- R

| Lecturer | Module |
|----------|-----------|
| Brown | Compilers |
| Brown | Databases |
| Green | Prolog |
| Green | Databases |
| Lewis | Prolog |
| Smith | Databases |

S

| Subject |
|-----------|
| Databases |
| Prolog |

R | S

| Lecturer |
|----------|
| Green |

- R

S

R | S

| Lecturer | Module |
|----------|-----------|
| Brown | Compilers |
| Brown | Databases |
| Green | Prolog |
| Green | Databases |
| Lewis | Prolog |
| Smith | Databases |

| Subject |
|-----------|
| Compilers |
| Prolog |

| Lecturer |
|----------|
| |

EXAMPLES OF DIVISION

| | | | | | | |
|----------|------------|------------|-----------|------------|-------------|------------|
| A | <i>sno</i> | <i>pno</i> | B1 | <i>pno</i> | A/B1 | <i>sno</i> |
| | s1 | p1 | | p2 | | s1 |
| | s1 | p2 | | | | s2 |
| | s1 | p3 | B2 | <i>pno</i> | | s3 |
| | s1 | p4 | | p2 | | s4 |
| | s2 | p1 | | p4 | A/B2 | <i>sno</i> |
| | s2 | p2 | B3 | <i>pno</i> | | s1 |
| | s3 | p2 | | p1 | | s4 |
| | s4 | p2 | | p2 | A/B3 | <i>sno</i> |
| | s4 | p4 | | p4 | | s1 |

DIVISION

Interpretation of the division operation A/B :

- Divide the attributes of A into 2 sets: A1 and A2.
- Divide the attributes of B into 2 sets: B2 and B3.
- Where the sets A2 and B2 have the same attributes.
- For each set of values in B2:
 - Search in A2 for the sets of rows (having the same A1 values) whose A2 values (taken together) form a set which is the same as the set of B2's.
 - For all the set of rows in A which satisfy the above search, pick out their A1 values and put them in the answer.

DIVISION

Example: Find the names of sailors who have reserved all boats:

(1) $A = \pi_{\text{sid}, \text{bid}}(\text{Reserves})$. $A1 = \pi_{\text{sid}}(\text{Reserves})$ $A2 = \pi_{\text{bid}}(\text{Reserves})$

(2) $B2 = \pi_{\text{bid}}(\text{Boats})$ $B3$ is the rest of B .

Thus, $B2 = \{101, 102, 103, 104\}$

(3) Find the rows of A such that their $A.\text{sid}$ is the same and their combined $A.\text{bid}$ is the set $B2$.

Thus we find $A1 = \{22\}$

(4) Get the set of $A2$ corresponding to $A1$: $A2 = \{\text{Dustin}\}$

Given a relational database schema as follows:

- SUPPLIER (S#, SNAME, STATUS, CITY),
- PART(P#, PNAME, COLOR, WEIGHT, CITY),
- PROJECT(J#, JNAME, CITY),
- SPJ(S#, P#, J#, QTY).

Please write the following query in Relational Algebra :

1. Find all suppliers who locate in Atlanta.

$$\sigma_{\text{CITY} = \text{'Atlanta'}}(\text{SUPPLIER})$$

2. Find all names of those suppliers who locate in Atlanta.

$$\pi_{\text{NAME}} (\sigma_{\text{CITY} = \text{'Atlanta'}} (\text{SUPPLIER}))$$

3. Find all parts that are red in color and weight more than 1 lb.

$$\pi_{\text{P\#}, \text{PNAME}, \text{CITY}} (\sigma_{\text{COLOR} = \text{'Red'} \text{ and } \text{WEIGHT} = \text{'1 lb'}} (\text{PART}))$$

4. Find all suppliers that have supplied P2.

$$\pi_{\text{S\#}, \text{SNAME}, \text{CITY},} (\sigma_{\text{P\#} = \text{'P2'}} (\text{SPJ} \bowtie \text{SUPPLIER}))$$

5. Find supplier names for suppliers who do not supply part P2.

$$\pi_{\text{S\#}, \text{SNAME}, \text{CITY}} (\sigma_{\text{P\#} \neq \text{'P2'}} (\text{SPJ} \bowtie \text{PART}))$$

6. list the S# of suppliers that have supplied all the parts listed in SPJ?

$$\pi_{S\#}(\text{SPJ}) \div \pi_{P\#}(\text{SPJ})$$

7. list the S# of suppliers who have supplied to all the projects with the project name 'CAD/CAM'

$$\pi_{S\#}(\text{SPJ} \div \pi_{J\#}(\sigma_{PNAME='CAD/CAM'}(\text{PROJECT})))$$

8. Find part numbers for parts supplied to any project in London.

$$\pi_{J\#}(\sigma_{CITY='LONDON'}(\text{PROJECT}) \bowtie \text{SPJ})$$

9. Find part numbers for parts supplied to all projects in London.

$$\pi_{P\#}(\text{SPJ} \div \pi_{J\#}(\sigma_{CITY='LONDON'}(\text{PROJECT})))$$

10. Get a list containing (S#, J#) for all suppliers who live in Atlanta but are not supplied to the project CAD/CAM?

$$(\pi_{S\#}(\sigma_{CITY='ATLANTA'}(SUPPLIER)) \times \pi_{J\#}(\sigma_{PNAME='CAD/CAM'}(PROJECT))) - \pi_{S\#,J\#}(SPJ)$$

11. Find names and city of all suppliers that have supplied more than 100 units of P2.

$$\pi_{SNAME, CITY}(\sigma_{P\#='P2' \text{ and } QTY > 100}(SPJ) \bowtie_{S\#} SUPPLIER)$$