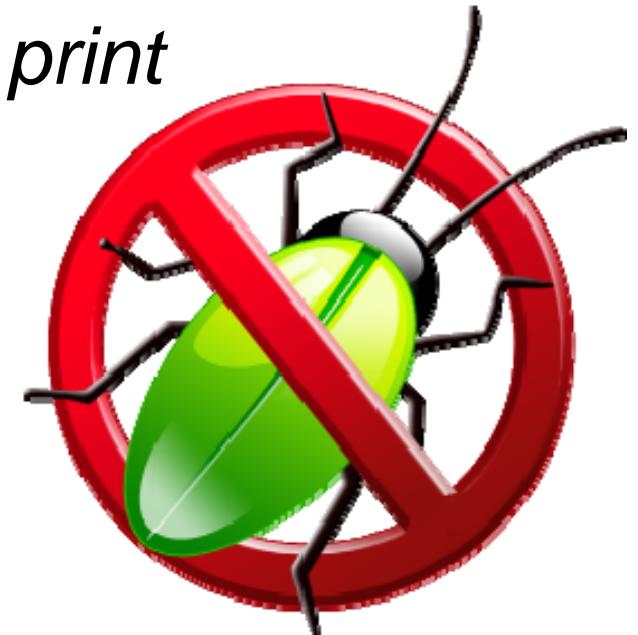


# Popular Ways to Debug

- Using “printf”
  - Often mocked, viewed as unscientific
  - In fact, just an easy way to apply dynamic/log analysis and perform experiments
  - If you ask the right questions, *print* can be a great debugging tool
    - Supports scientific debugging
    - I have no lessons here, other than to print intelligently, not blindly grope around



# Popular Ways to Debug

- Using a debugger
  - Usually thought of as “more scientific”
  - It can be!
  - A debugger is good when you want to:
    - Inspect closely what happens to some values
    - Slow down and carefully watch things during one part of a run
    - Get information across a lot of state at once
  - Hard to make guidelines, but in general printf is for “across time” and debuggers are for “across state”



# Basic Assertions

```
if (divisor == 0) {  
    printf("Division by zero!");  
    abort();  
}
```

---

# Specific Assertions

***assert (divisor != 0);***

# Implementation

```
void assert (int x)
{
    if (!x)
    {
        printf("Assertion failed!\n");
        abort();
    }
}
```

# Execution

```
$ my-program
Assertion failed!
Abort (core dumped)
$
```

# Better Diagnostics

```
$ my-program
divide.c:37:
    assertion 'divisor != 0' failed
Abort (core dumped)
$ _
```

# Assertions as Macros

```
#ifndef NDEBUG
#define assert(ex) \
((ex) ? 1 : (cerr << __FILE__ << ":" << __LINE__ \
      << ": assertion \" #ex \" failed\n", \
      abort(), 0))
#else
#define assert(x) ((void) 0)
#endif
```

# Ensuring Sanity (A Time Class)

```
class Time {  
public:  
    int hour(); // 0..23  
    int minutes(); // 0..59  
    int seconds(); // 0..60 (incl. leap seconds)  
  
    void set_hour(int h);  
    ...  
}
```

***Any time from 00:00:00 to 23:59:60 is valid***

# Ensuring Sanity

```
void Time::set_hour(int h)
{
    // precondition
    assert (0 <= hour() && hour() <= 23) &&
           (0 <= minutes() && minutes() <= 59) &&
           (0 <= seconds() && seconds() <= 60);

    ...
    // postcondition
    assert (0 <= hour() && hour() <= 23) &&
           (0 <= minutes() && minutes() <= 59) &&
           (0 <= seconds() && seconds() <= 60);
}
```

# Ensuring Sanity

```
bool Time::sane()
{
    return (0 <= hour() && hour() <= 23) &&
           (0 <= minutes() && minutes() <= 59) &&
           (0 <= seconds() && seconds() <= 60);
}

void Time::set_hour(int h)
{
    assert (sane()); // precondition
    ...
    assert (sane()); // postcondition
}
```

# Ensuring Sanity

```
bool Time::sane()  
{  
    return (0 <= hour() && hour() <= 23) &&  
        (0 <= minutes() && minutes() <= 59) &&  
        (0 <= seconds() && seconds() <= 60);  
}
```

- **sane()** is the *invariant* of a Time object:
  - holds *before* every public method
  - holds *after* every public method

# Ensuring Sanity

```
bool Time::sane()  
{  
    return (0 <= hour() && hour() <= 23) &&  
        (0 <= minutes() && minutes() <= 59) &&  
        (0 <= seconds() && seconds() <= 60);  
}
```

```
void Time::set_hour(int h)  
{  
    assert (sane());  
    ...  
    assert (sane());  
}
```

*same for  
set\_minute(),  
set\_seconds(), etc.*

# Locating Infections

- Precondition failure = infection *before* method
- Postcondition failure = infection *within* method
- All assertions pass = no infection

```
void Time::set_hour(int h)
{
    assert(sane()); // precondition
    ...
    assert(sane()); // postcondition
}
```

# Relative Debugging

- Rather than checking a spec, we can also compare against a *reference run*:

- The environment has changed—e.g. ports or new interpreters
- The code has changed
- The program has been reimplemented

# Relative Assertions

- We compare two program runs
- A *relative assertion* compares variable values across the two runs:  
`.assert \ p1::perimeter@polygon.java:65 == \ p0::perimeter@polygon.java:65`
- Specifies when and what to compare

# Concepts

- ★ **Assertions catch infections before they propagate too far**
- ★ **Assertions check preconditions, postconditions and invariants**
- ★ **Assertions can serve as specifications**
- ★ **A program can serve as reference to be compared against**

# System Invariants

- Some properties of a program must hold over the entire run:
  - must not access data of other processes
  - must handle mathematical exceptions
  - must not exceed its privileges
- Typically checked by hardware and OS

# Memory Invariants

- Even within a single process, some invariants must hold over the entire run
  - code integrity
  - data integrity
- This is a major issue in C and C++

# Heap Misuse

***s = malloc(30)***

***free(s)***

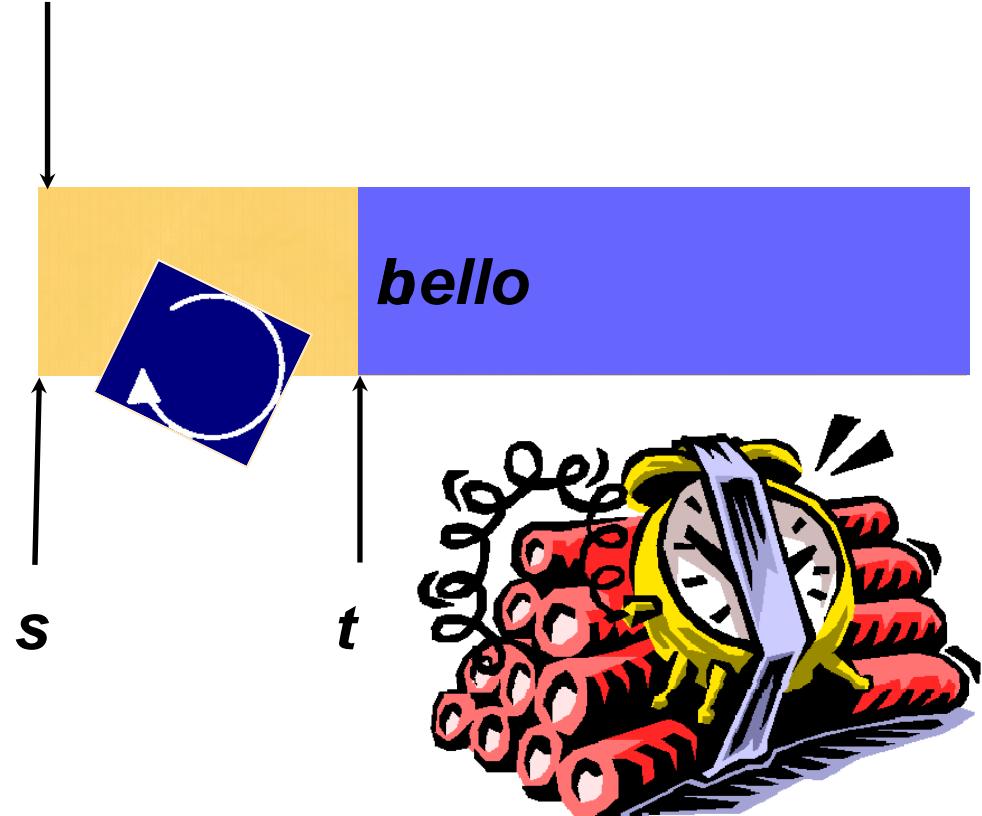
***t = malloc(20)***

***strcpy(t, "hello")***

***s[10] = 'b'***

***free(s)***

***free\_list***



# Heap Assertions

- The *GNU C runtime library* provides a simple check against common errors:

```
$ MALLOC_CHECK_=2 myprogram myargs  
free() called on area that was already free'd()  
Aborted (core dumped)  
$ _
```

# Heap Assertions

`s = malloc(30)`

`free(s)`

`free(s)`



*free() called on area that was already free'd()*

*Aborted (core dumped)*

# Array Assertions

- The *Electric Fence* library checks for array overflows:

```
$ gcc -g -o sample-with-efence sample.c -lefence  
$ ./sample-with-efence 11 14  
Electric Fence 2.1  
Segmentation fault (core dumped)  
$ _
```

# Array Assertions

`s = malloc(30)`

`s[30] = 'x'`



*Segmentation fault (core dumped)*

# Memory Assertions

- The *Valgrind* tool checks *all* memory accesses:

```
$ valgrind sample 11 14
```

*Invalid read of size 4*

*at 0x804851F: shell\_sort (sample.c:18)*

*by 0x8048646: main (sample.c:35)*

*by 0x40220A50: \_\_libc\_start\_main (in /lib/libc.so)*

*by 0x80483D0: (within /home/zeller/sample)*

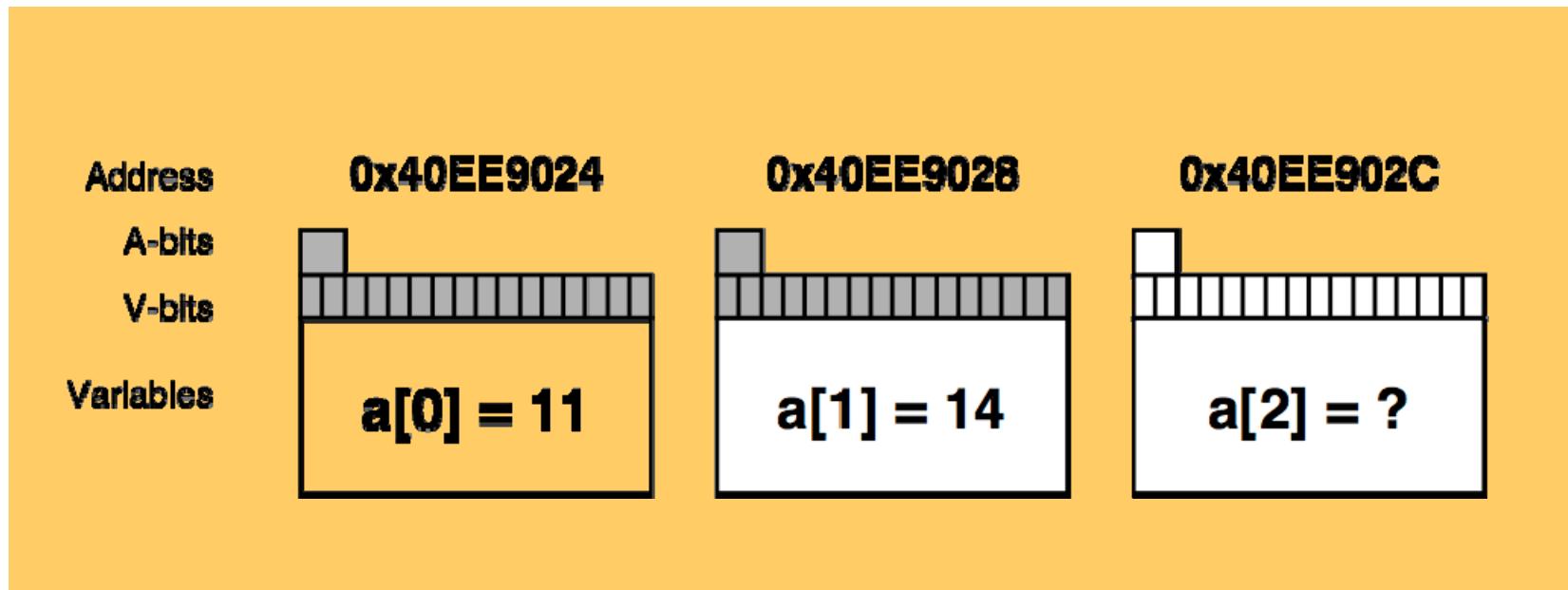
- Valgrind works as an *interpreter* for x86 code

# Valgrind Checks

- ▶ **Use of uninitialized memory**
- ▶ **Accessing free'd memory**
- ▶ **Accessing memory beyond malloc'd block**
- ▶ **Accessing inappropriate stack areas**
- ▶ **Memory leaks: allocated area is not free'd**
- ▶ **Passing uninitialized memory to system calls**

# Shadow Memory

- V-bit set = corresponding bit is initialized
- A-bit set = corresponding byte is accessible



# V-Bits

- When a bit is first written, its V-bit is set
- Simple read accesses to uninitialized memory do not result in warnings:

```
struct S { int x; char c; };
struct S s1, s2;
s1.x = 42;
s1.c = 'z';
s2 = s1;
```

5 bytes *initialized*  
8 bytes *copied*  
*(no warning)*

# V-Bits Warnings

- **Reading uninitialized data causes a warning if**
  - a value is used to generate an *address*
  - a *control flow decision* is to be made
  - a value is passed to a *system call*

# A-Bits

- When the program starts, all global data is marked “accessible” (= A-bits are set)
- `malloc()` sets A-bits for the area returned; `free()` clears them
- Local variables are “accessible” on entry and “non-accessible” on exit
- Accessing “non-accessible” data ⇒ error

# Overhead

Tool	GNU C Library	Electric Fence	Valgrind
Space	2 bytes/ malloc	1 page/malloc	100 %
Time	negligible	negligible	2500 %

# Preventing Misuse

- CYCLONE is a C dialect which prevents common pitfalls of C
- Most important feature: *special pointers*

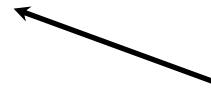
# Non-NULL Pointers

```
int getc(FILE @fp);
```



***fp may not be NULL***

```
extern FILE *fp;  
char c = getc(fp);
```

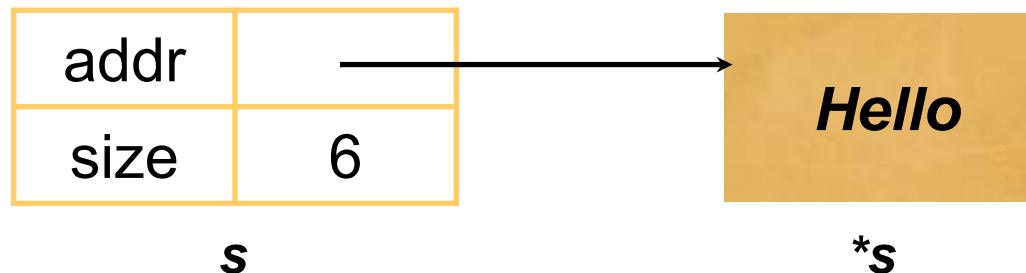


***warning: NULL check inserted***

# Fat Pointers

- A fat pointer holds address *and* size
- All accesses via a fat pointer are automatically bounds-checked

*int strlen(const char? s)*



# CYCLONE Restrictions

- ▶ **NULL checks are inserted**
- ▶ **Pointer arithmetic is restricted**
- ▶ **Pointers must be initialized before use**
- ▶ **Dangling pointers are prevented through region analysis and limitations on free()**
- ▶ **Only “safe” casts and unions are allowed**

---

# Production Code

- Should products ship with active assertions?

# Things to Check

**Critical results.** If lives, health, or money depend on a result, it had better be checked.

**External conditions.** Any conditions which are not within our control must be checked for integrity.

# Points to Consider

- The more active assertions, the greater the chance to catch infections.
- The sooner a program fails, the easier it is to track the defect.
- Defects that escape into the field are the hardest to track.

# More to Consider

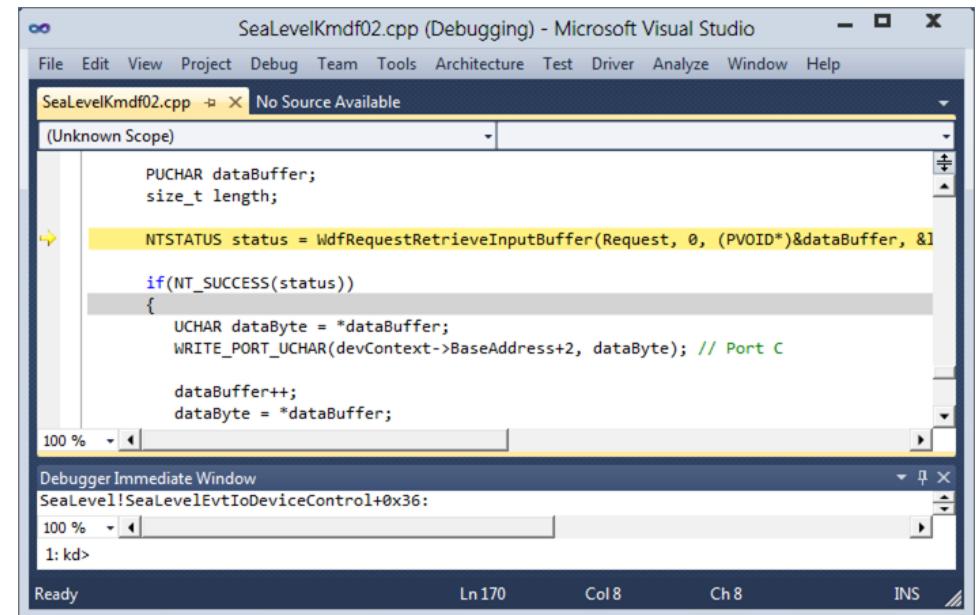
- By default, failing assertions are not user-friendly.
  - *Handle assertions in a user-friendly way*
- Assertions impact performance.
  - *First measure; then turn off only the most time-consuming assertions*

# Concepts

- ★ **To check memory integrity, use specialized tools to detect errors at run time**
- ★ **Apply such tools before any other method**
- ★ **To fully prevent memory errors, use another language (or dialect, e.g. Cyclone)**
- ★ **Turning assertions off seldom justifies the risk of erroneous computation**

# Debuggers

- Let you “take over” a program and run it under more control than usual
- Command line (gdb) and visual debuggers are both popular and widely used



# GDB

- There are many other debuggers out there
- Won't spend a lot of time on GDB specifics
- Major features (common to many debuggers or becoming more common)
  - Single step through a program line at a time
    - GDB: **step** and **next**
      - **next** skips over functions called in a line
    - Inspect memory locations/values
    - Set breakpoints – places to stop execution
      - Can be conditional (break at line XX if y > z)
    - Set watchpoints – events to watch for in execution
    - Change values in memory
    - GDB can also “run a program backwards” a little bit now!

# GDB

- **Sig**  
de

```
/cygdrive/c/Documents and Settings/Alex/Desktop/cs362class/dominion
Alex@groce /cygdrive/c/Documents and Settings/Alex/Desktop/cs362class/dominion
$ ./badTestDrawCard.exe
Testing drawCard.
RANDOM TESTS.
Segmentation fault (core dumped)

• Alex@groce /cygdrive/c/Documents and Settings/Alex/Desktop/cs362class/dominion
$ gdb ./badTestDrawCard.exe
GNU gdb (GDB) 7.3.50.2011026-cvs (cygwin-special)
Copyright (C) 2011 Free Software Foundation, Inc.
License GPLv3+: GNU GPL version 3 or later <http://gnu.org/licenses/gpl.html>
This is free software: you are free to change and redistribute it.
There is NO WARRANTY, to the extent permitted by law. Type "show copying"
and "show warranty" for details.
This GDB was configured as "i686-cygwin".
For bug reporting instructions, please see:
<http://www.gnu.org/software/gdb/bugs/>...
Reading symbols from /cygdrive/c/Documents and Settings/Alex/Desktop/cs362class/
dominion/badTestDrawCard.exe...done.
(gdb) run
Starting program: /cygdrive/c/Documents and Settings/Alex/Desktop/cs362class/dominion/badTestDrawCard.exe
[New Thread 1980.0x46c]
[New Thread 1980.0xf2c]
Testing drawCard.
RANDOM TESTS.

Program received signal SIGSEGV, Segmentation fault.
0x00403348 in drawCard (player=145, state=0x2844f0) at dominion.c:534
534          state->deck[player][i] = state->discard[player][i];
(gdb) bt
#0  0x00403348 in drawCard (player=145, state=0x2844f0) at dominion.c:534
#1  0x004011a4 in checkDrawCard (p=145, post=0x2844f0) at badTestDrawCard.c:14
#2  0x0040141d in main () at badTestDrawCard.c:38
(gdb) print i
$1 = 0
(gdb) print player
$2 = 145
(gdb) |
```

ault

ther

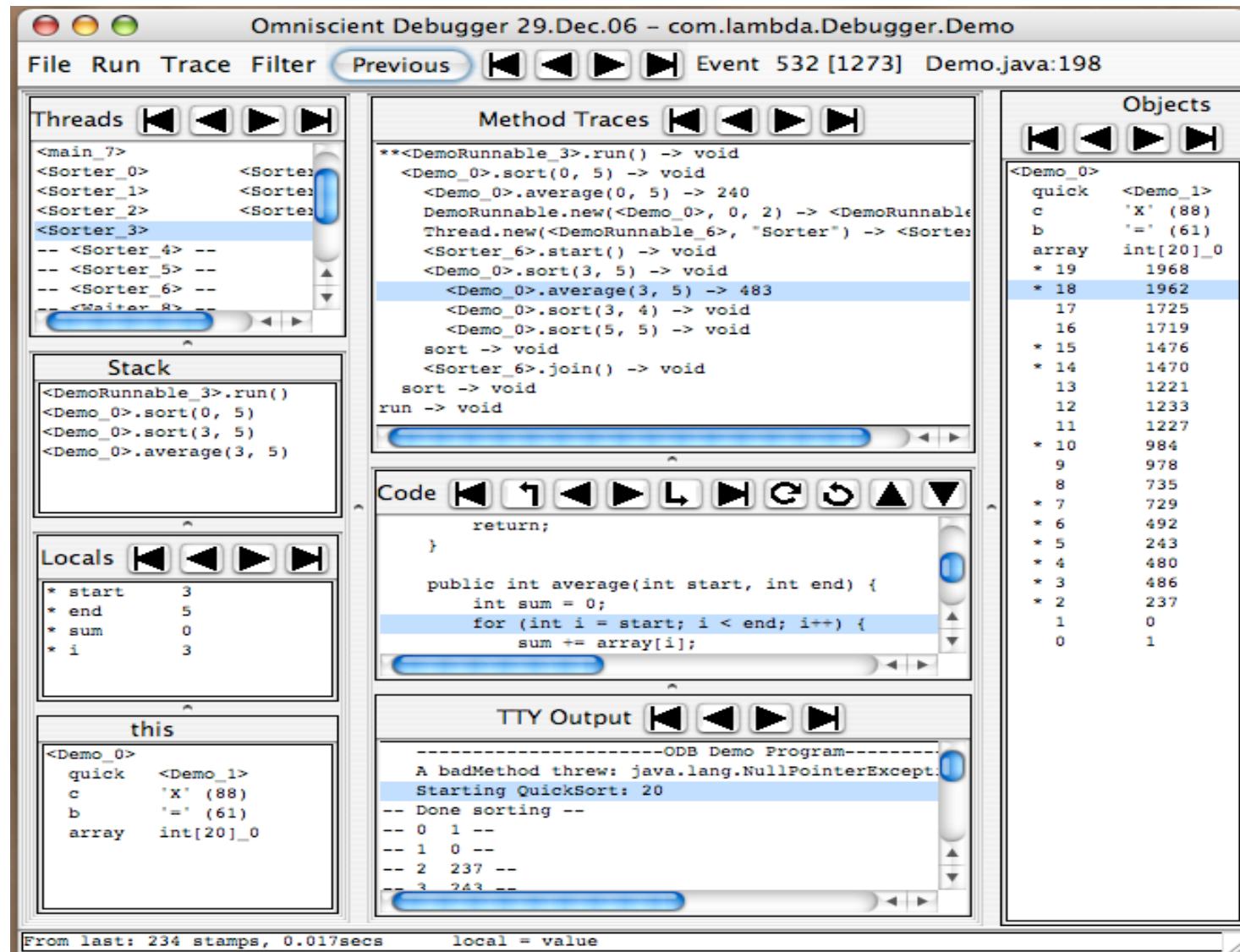
# GDB

```
$ gdb testDrawCard
GNU gdb (GDB) 7.3.50.20111020-cvs (cygwin Special)
Copyright (C) 2011 Free Software Foundation, Inc.
License GPLv3+: GNU GPL version 3 or later <http://gnu.org/licenses/gpl.html>
This is free software: you are free to change and redistribute it.
There is NO WARRANTY, to the extent permitted by law. Type "show copying"
and "show warranty" for details.
This GDB was configured as "i686-cygwin".
For bug reporting instructions, please see:
<http://www.gnu.org/software/gdb/bugs/>...
Reading symbols from /cygdrive/c/Documents and Settings/Alex/Desktop/cs362class/dominion/testDrawCard... do
no
(gdb) break drawCard
Breakpoint 1 at 0x403790: file dominion.c, line 528
(gdb) run
Starting program: /cygdrive/c/Documents and Settings/Alex/Desktop/cs362class/dominion/testDrawCard
[New Thread 1272.0xd0c]
[New Thread 1272.0x558]
Testing drawCard.
RANDOM TESTS.
TEST #0

Breakpoint 1, drawCard (player=0, state=0x284500) at dominion.c:528
warning: Source file is more recent than executable.
528      if (<state->deckCount[player] <= 0){//Deck is empty
(gdb) bt
#0  drawCard (player=0, state=0x284500) at dominion.c:528
#1  0x004011c9 in checkDrawCard (p=0, post=0x284500) at testDrawCard.c:19
#2  0x004018da in main () at testDrawCard.c:69
(gdb) print state
$1 = (struct gameState *) 0x284500
(gdb) print state->whoseTurn
$2 = 1817024117
(gdb) print state->discardCount[player]
$3 = 78
(gdb) print state->deckCount[player]
$4 = 190
(gdb) watch state->deckCount[player]
Hardware watchpoint 2: state->deckCount[player]
(gdb) continue
Continuing.
Hardware watchpoint 2: state->deckCount[player]

old value = 190
New value = 189
drawCard (player=0, state=0x284500) at dominion.c:577
577      state->handCount[player]++;
(gdb) print state->deckCount[player]
$5 = 189
(gdb) list
573
574      deckCounter = state->deckCount[player];//Create holder for the deck count
575      state->hand[player][count] = state->deck[player][deckCounter - 1];//Add card to the hand
576      state->deckCount[player]--;
577      state->handCount[player]++;
578  }
579
580      return 0;
581  }
(gdb)
```

# Omniscient Debugger



# How does it work?

- ODB records a *trace* of the entire execution history
- Slows down programs by a factor of 10
- Records about 100 MB/s
- Now available in commercial tools:
  - RETROVUE
  - CODEGUIDE

# Exploring the Past

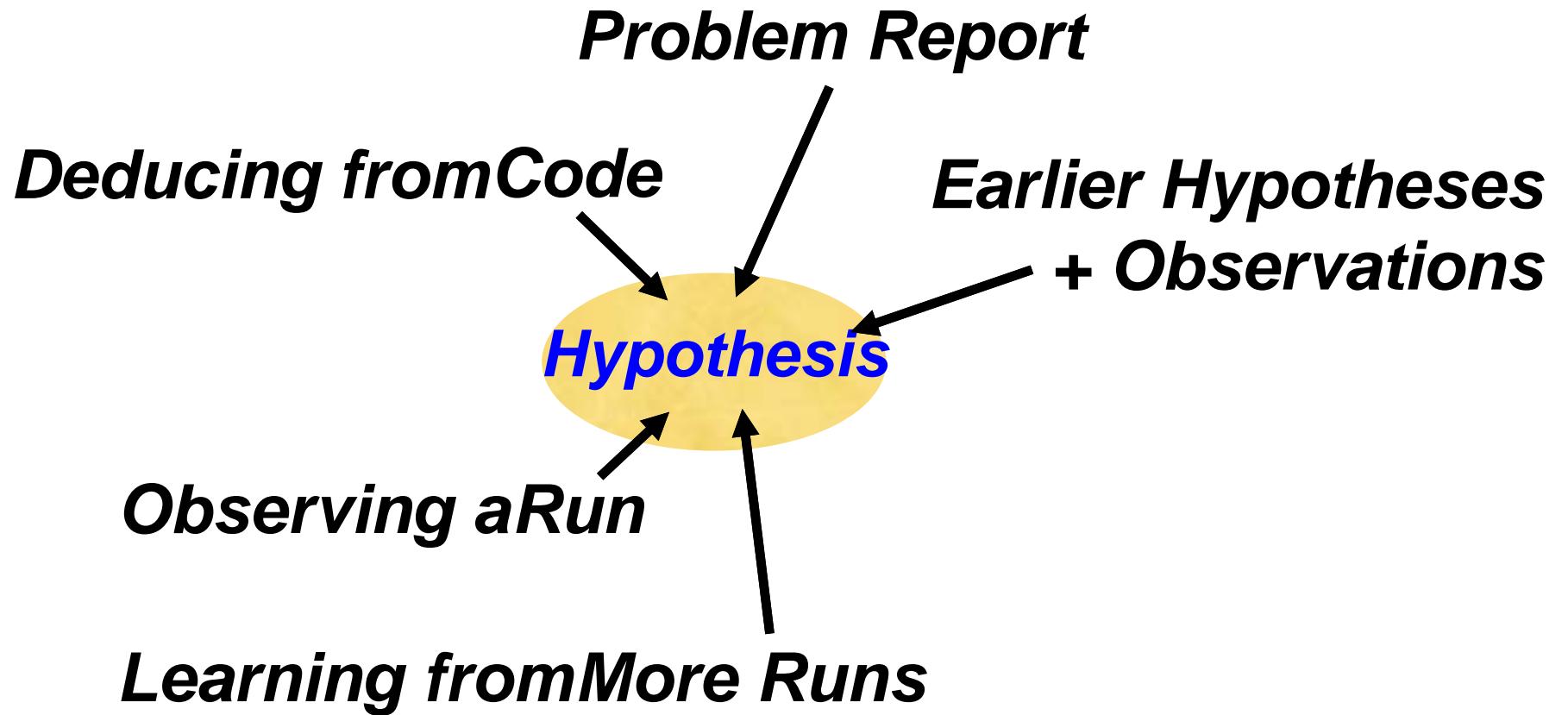
*A typical debugging session looks like this:*

- 1. Set a breakpoint**
- 2. Start program, reaching breakpoint**
- 3. Step, Step, Step, ...**
- 4. Oops! I've gone too far!**

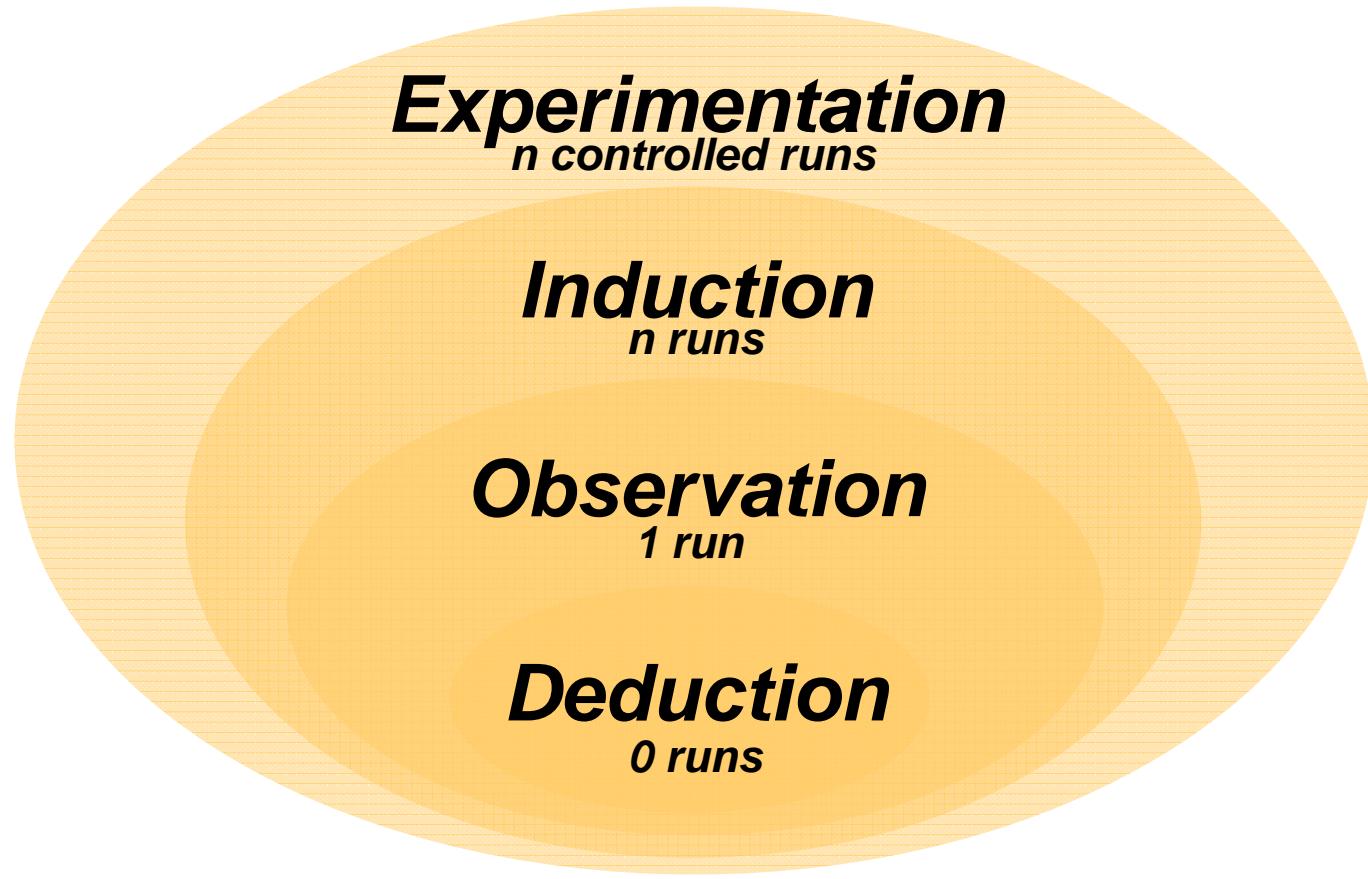
---

# Deducing Errors

# Obtaining a Hypothesis



# Reasoning about Runs



# Reasoning about Runs

***Deduction***  
*0 runs*

# Effects of Statements

**Write.** A statement can change the program state  
(i.e. write to a variable)

**Control.** A statement may determine which  
statement is executed next  
(other than unconditional transfer)

# Affected Statements

**Read.** A statement can read the program state (i.e. from a variable)

**Execution.** To have any effect, a statement must be executed.

# Effects in fibo.c

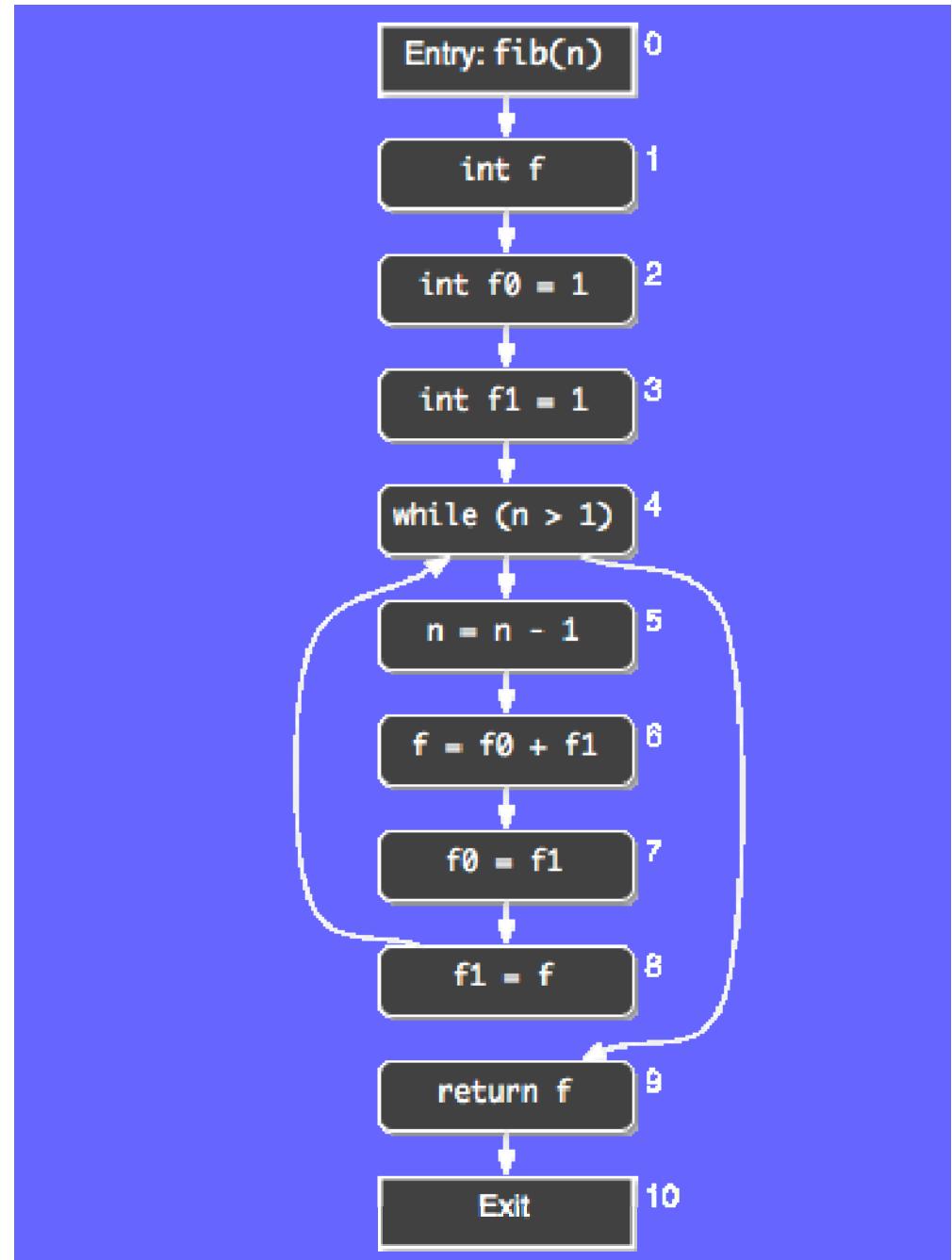
	Statement	Reads	Writes	Controls
0	fib(n)		n	1-10
1	int f		f	
2	f0 = 1		f0	
3	f1 = 1		f1	
4	while (n > 1)	n		5-8
5	n = n - 1	n	n	
6	f = f0 + f1	f0, f1	f	
7	f0 = f1	f1	f0	
8	f1 = f	f	f1	
9	return f	f	<ret>	

# Control Flow

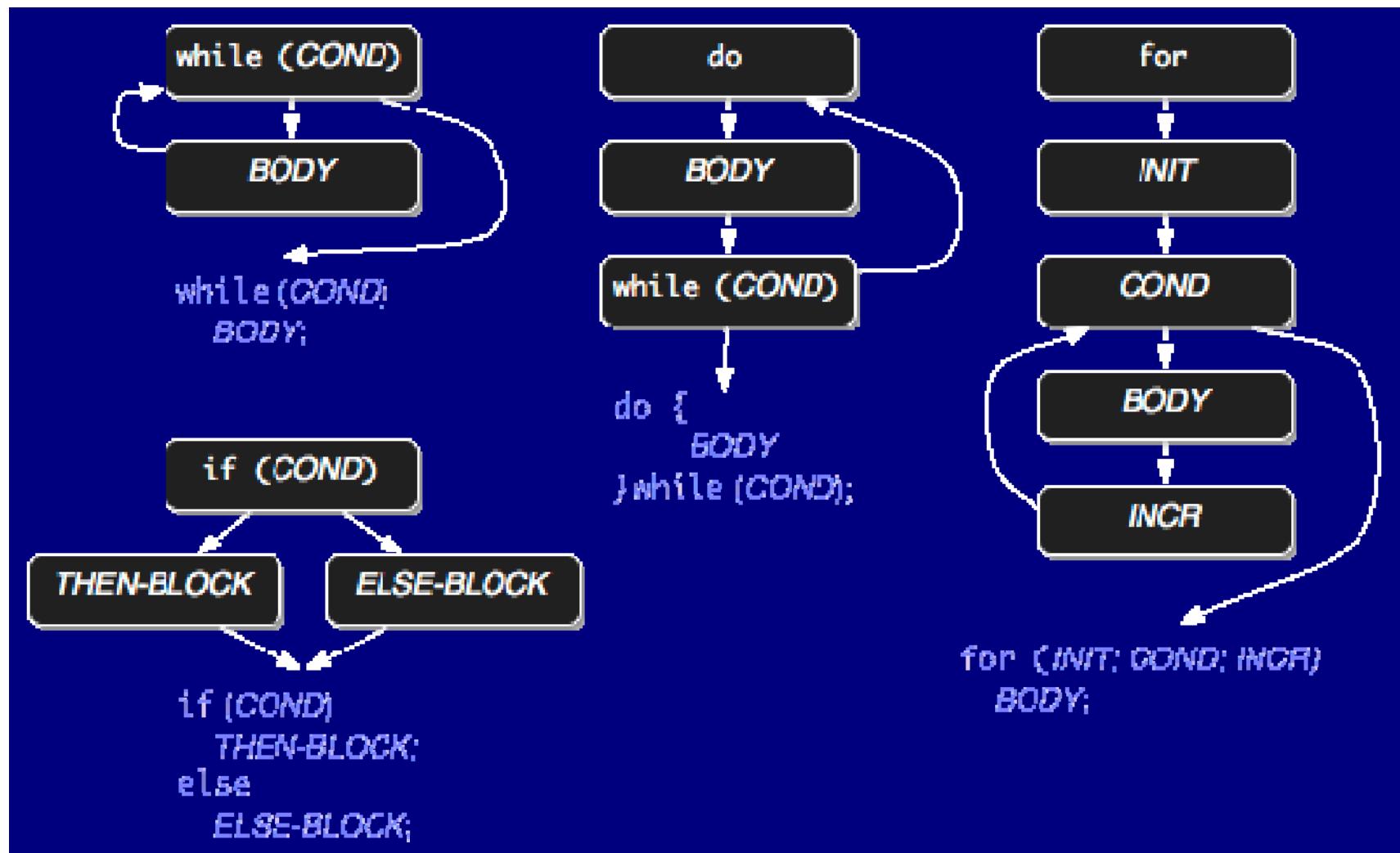
```
int fib(int n)
{
    int f, f0 = 1, f1 = 1;

    while (n > 1) {
        n = n - 1;
        f = f0 + f1;
        f0 = f1;
        f1 = f;
    }

    return f;
}
```



# Control Flow Patterns



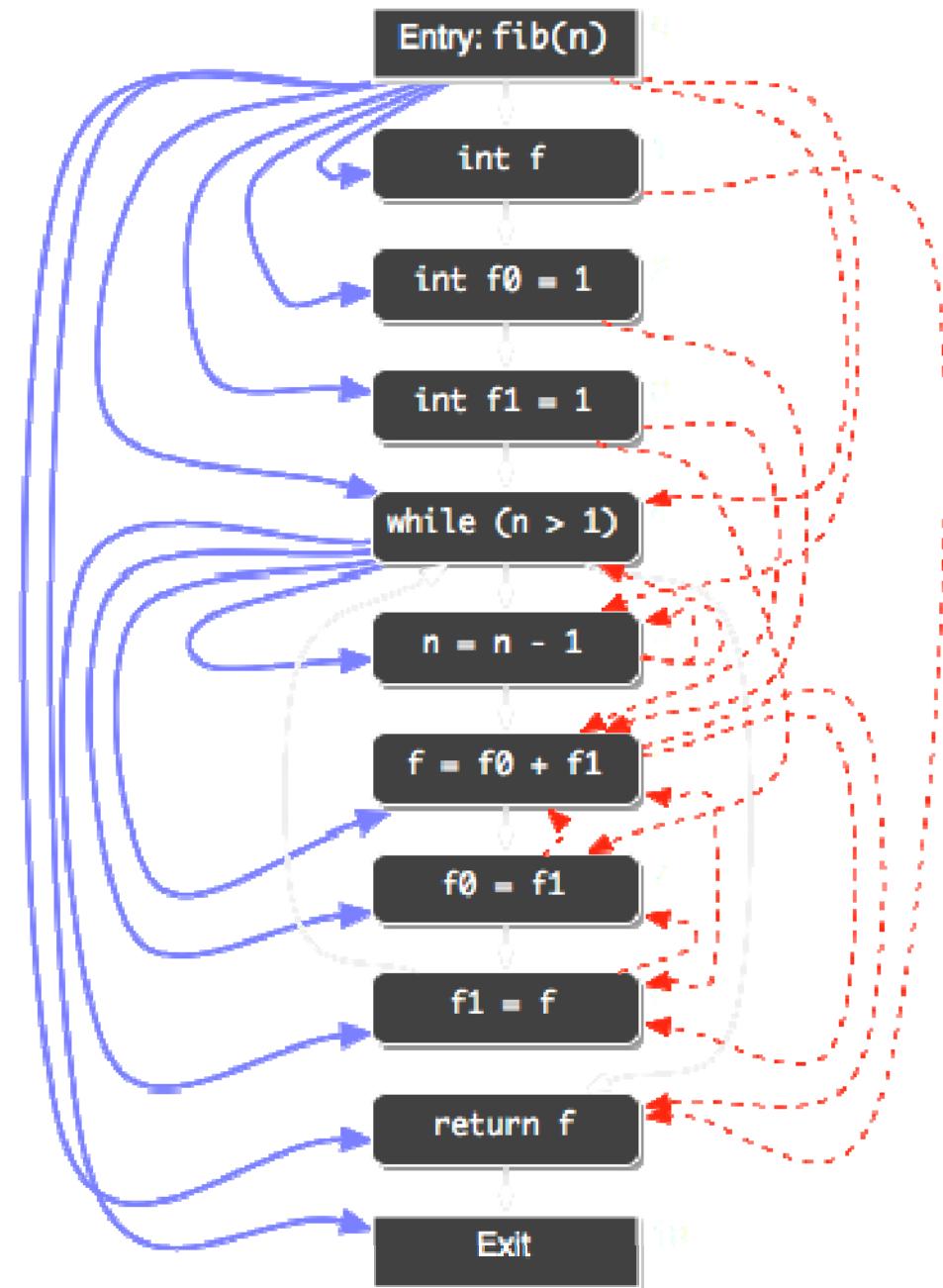
# Dependences



**Data dependency:**  
A depends on B  
B is data dependent on A

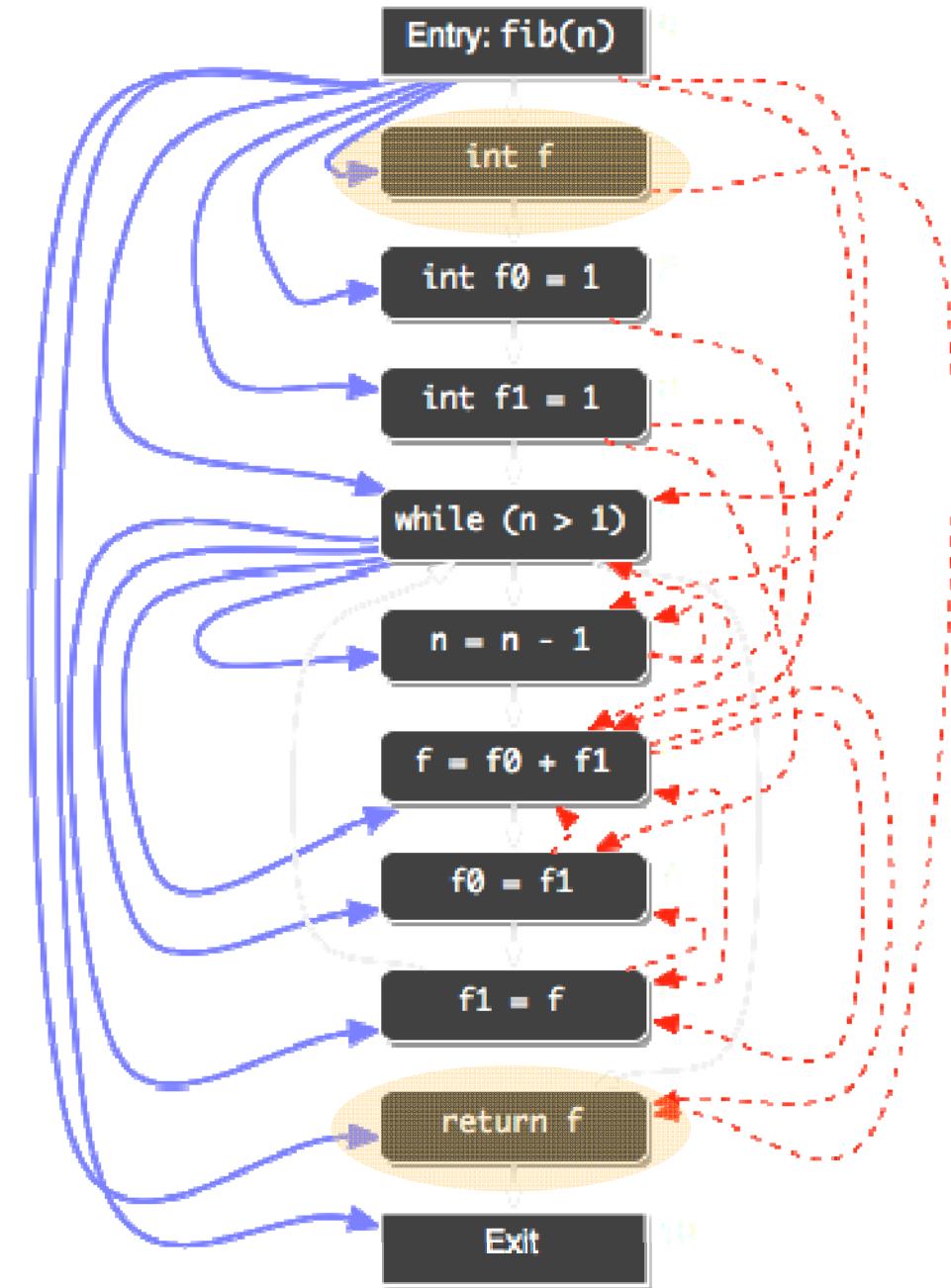


**Control dependency:**  
A controls B  
B is control dependent on A

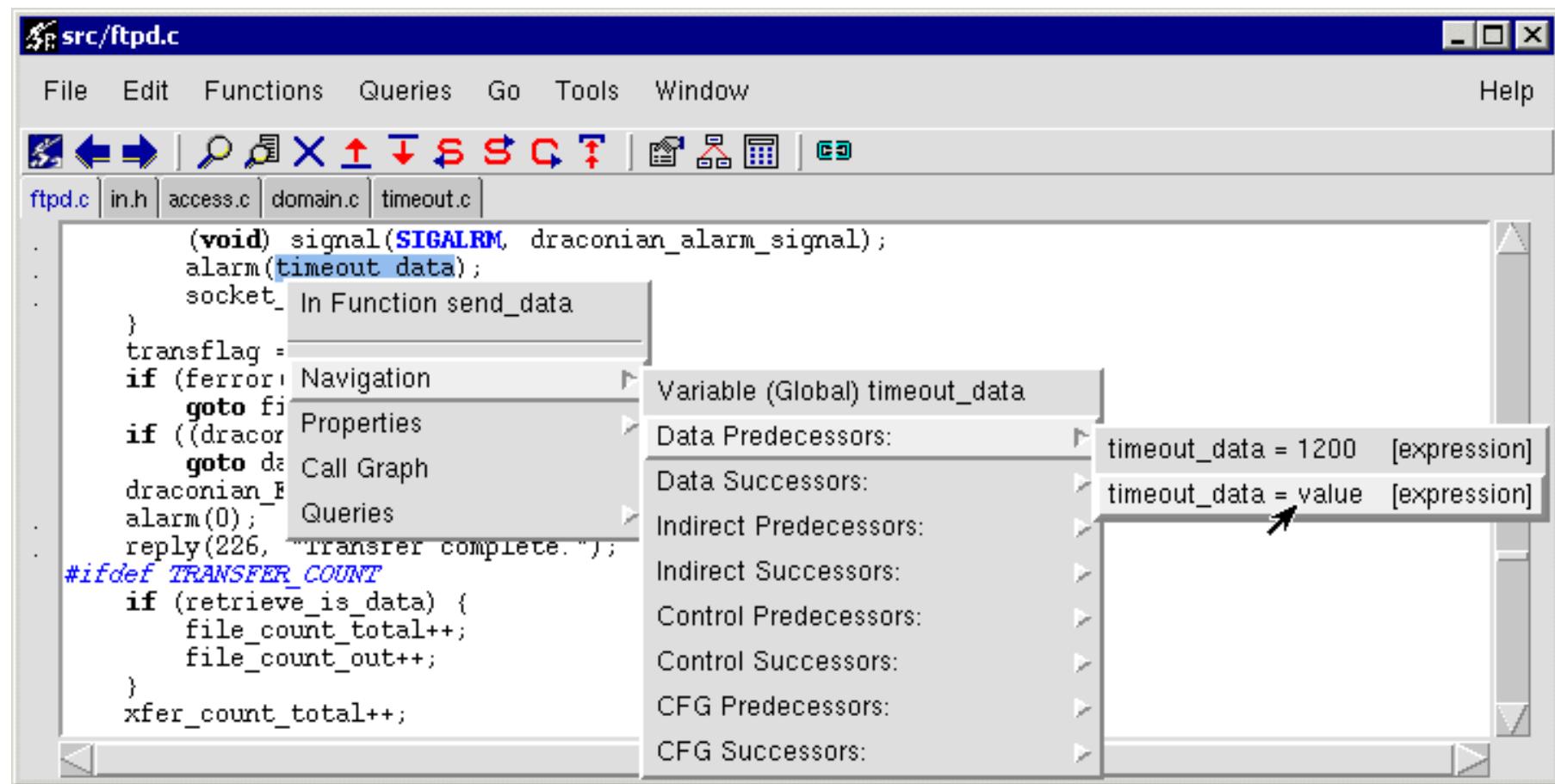


# Dependences

- Following the dependences, we can answer questions like
- Where does this value go to?
- Where does this value come from?



# Navigating along Dependencies



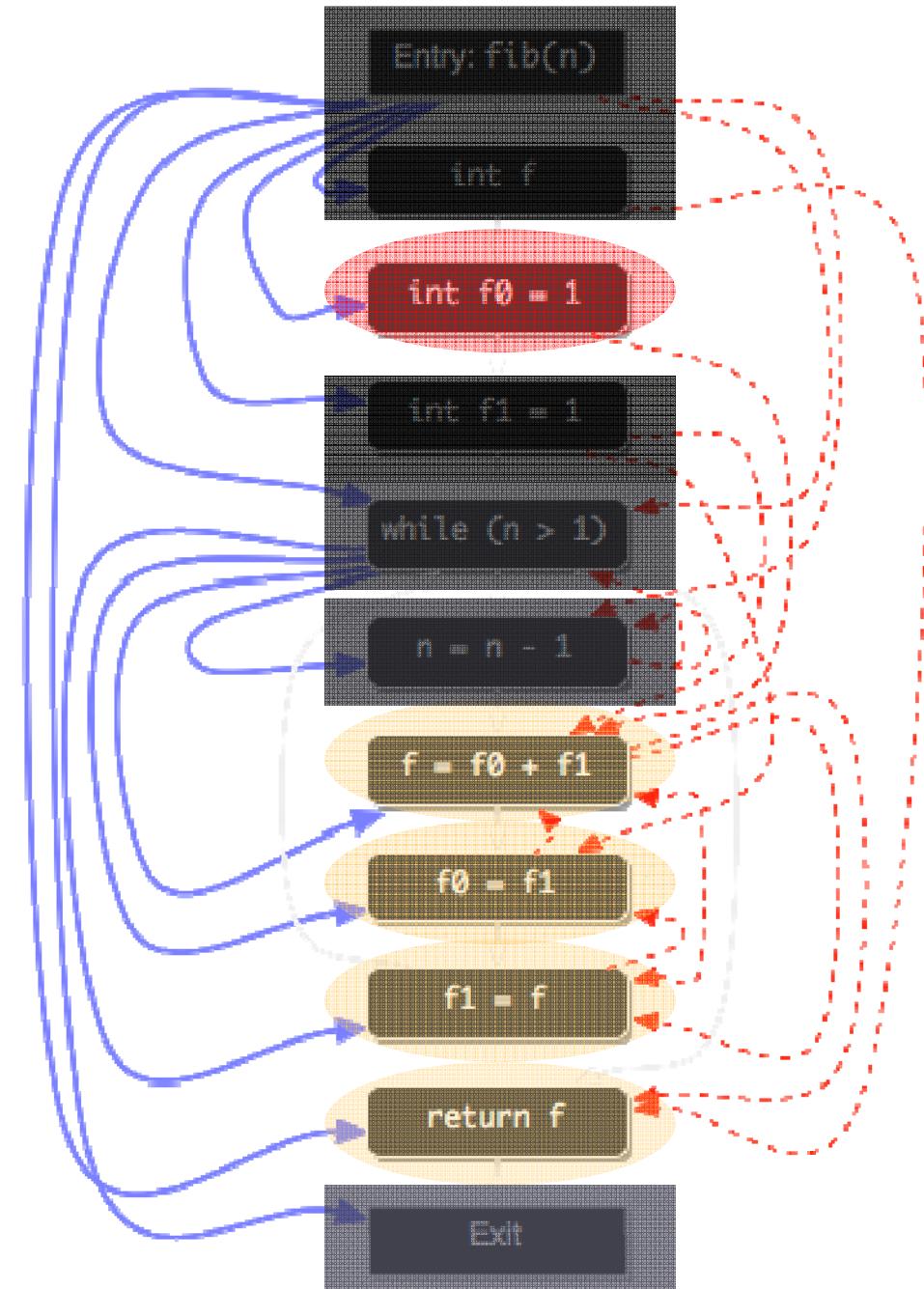
# Program Slicing

- A *slice* is a subset of the program
- Allows programmers to *focus on what's relevant* with respect to some statement S:
  - All statements influenced by S
  - All statements that influence S

# Forward Slice

- Given a statement A, the forward slice contains all statements whose read variables or execution could be influenced by A
- Formally:

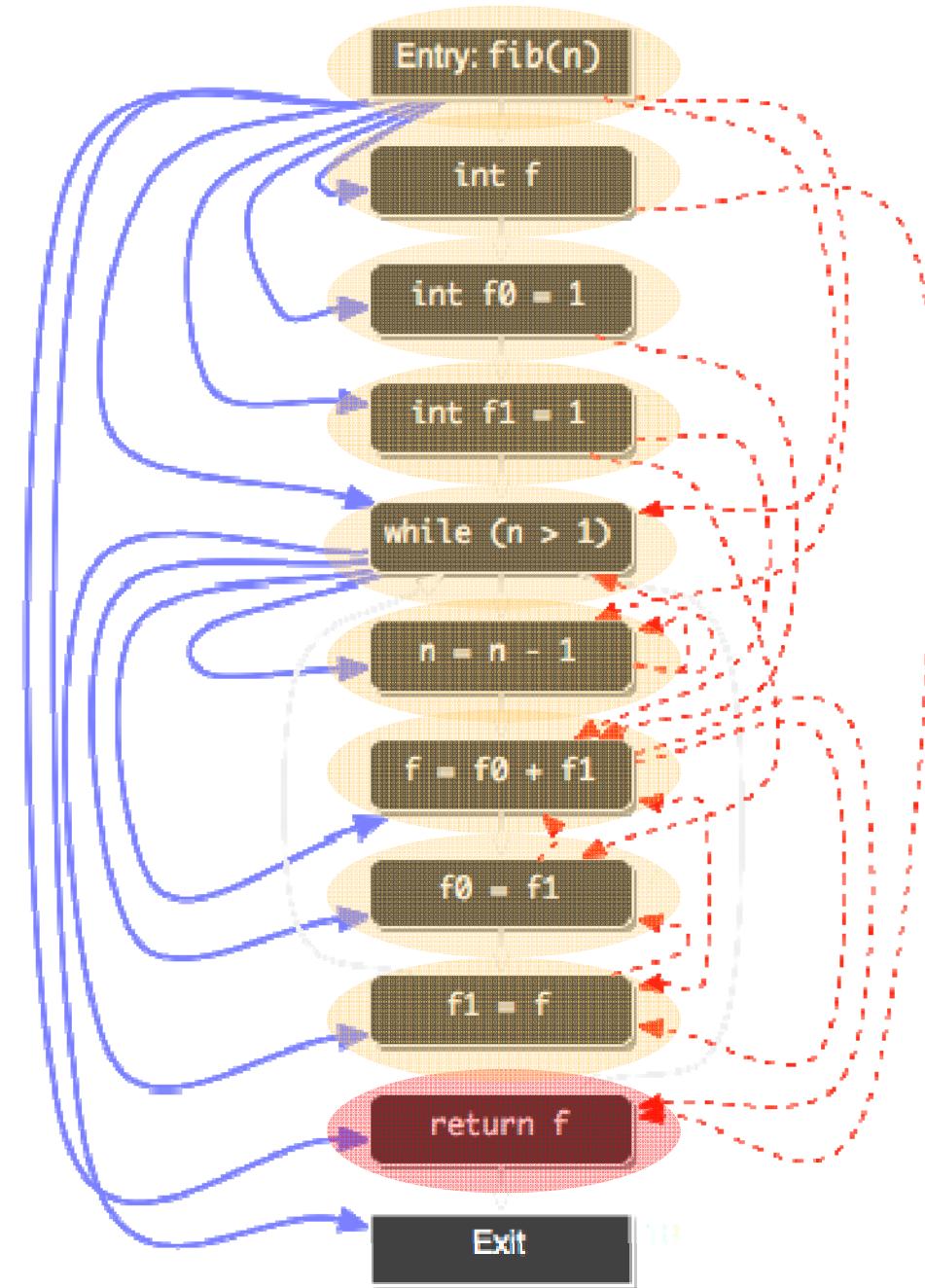
$$S^F(A) = \{B | A \rightarrow^* B\}$$



# Backward Slice

- Given a statement B, the backward slice contains all statements that could influence the read variables or execution of B
- Formally:

$Slice(B) = \{A | A \in P\}$



# Two Slices

```
int main() {
    int a, b, sum, mul;
    sum = 0;
    mul = 1;
    a = read();
    b = read();
    while (a <= b) {
        sum = sum + a;
        mul = mul * a;
        a = a + 1;
    }
    write(sum);
    write(mul);
}
```

- Slice Operations:
- Backbones
- Dices
- Chops

← **Backward slice of sum**  
← **Backward slice of mul**

# Backbone

```
a = read();  
b = read();  
while (a <= b) {
```

```
a = a + 1;
```

- Contains only those statement that occur in both slices
- Useful for focusing on common behavior

# Two Slices

```
int main() {  
    int a, b, sum, mul;  
    sum = 0;  
    mul = 1;  
    a = read();  
    b = read();  
    while (a <= b) {  
        sum = sum + a;  
        mul = mul * a;  
        a = a + 1;  
    }  
    write(sum);  
    write(mul);  
}
```

- Slice Operations:

- Backbones
- Dices
- Chops

← **Backward slice of sum**  
← **Backward slice of mul**

# Dice

```
sum = 0;
```

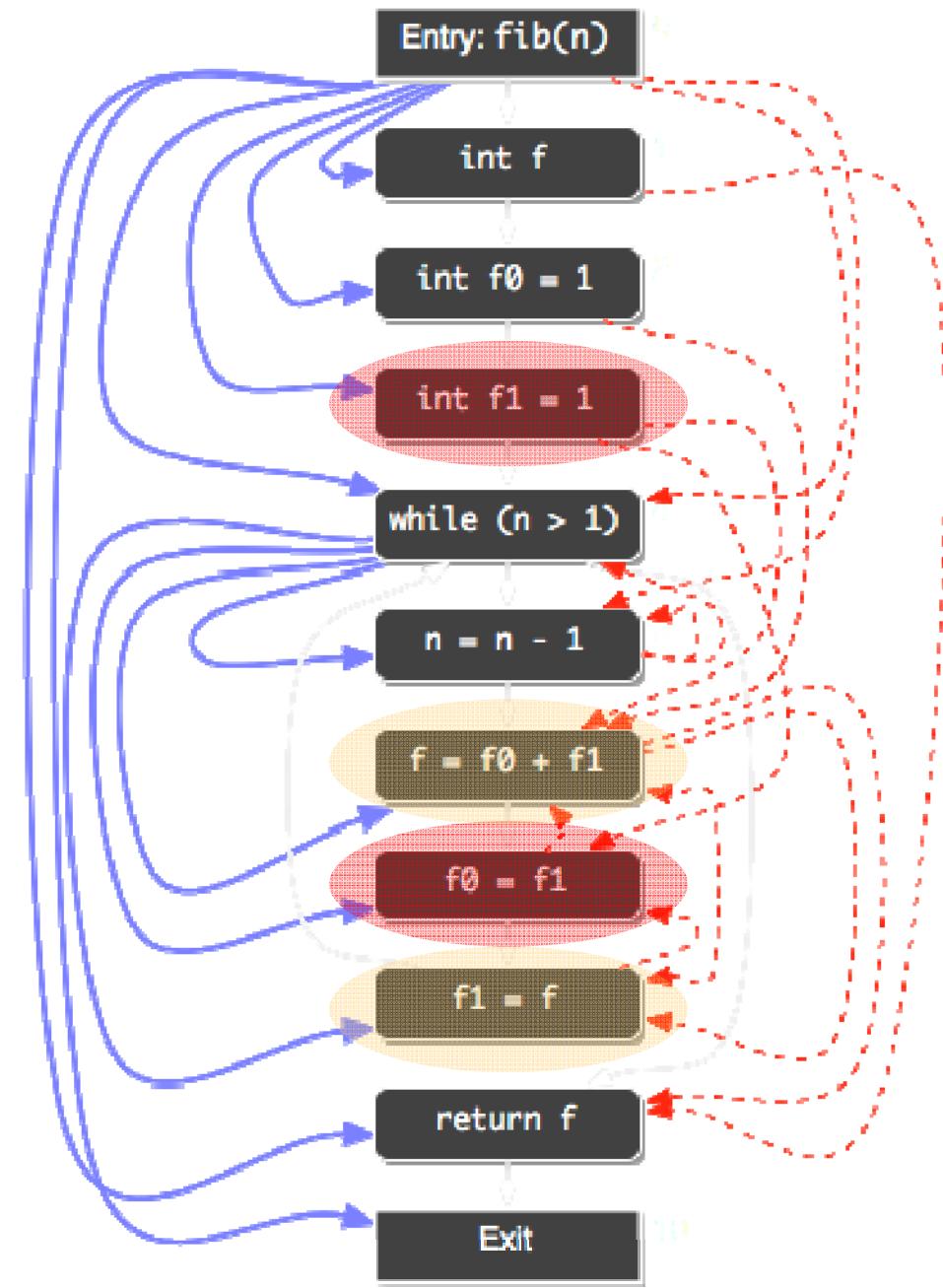
```
sum = sum + a;
```

```
write(sum);
```

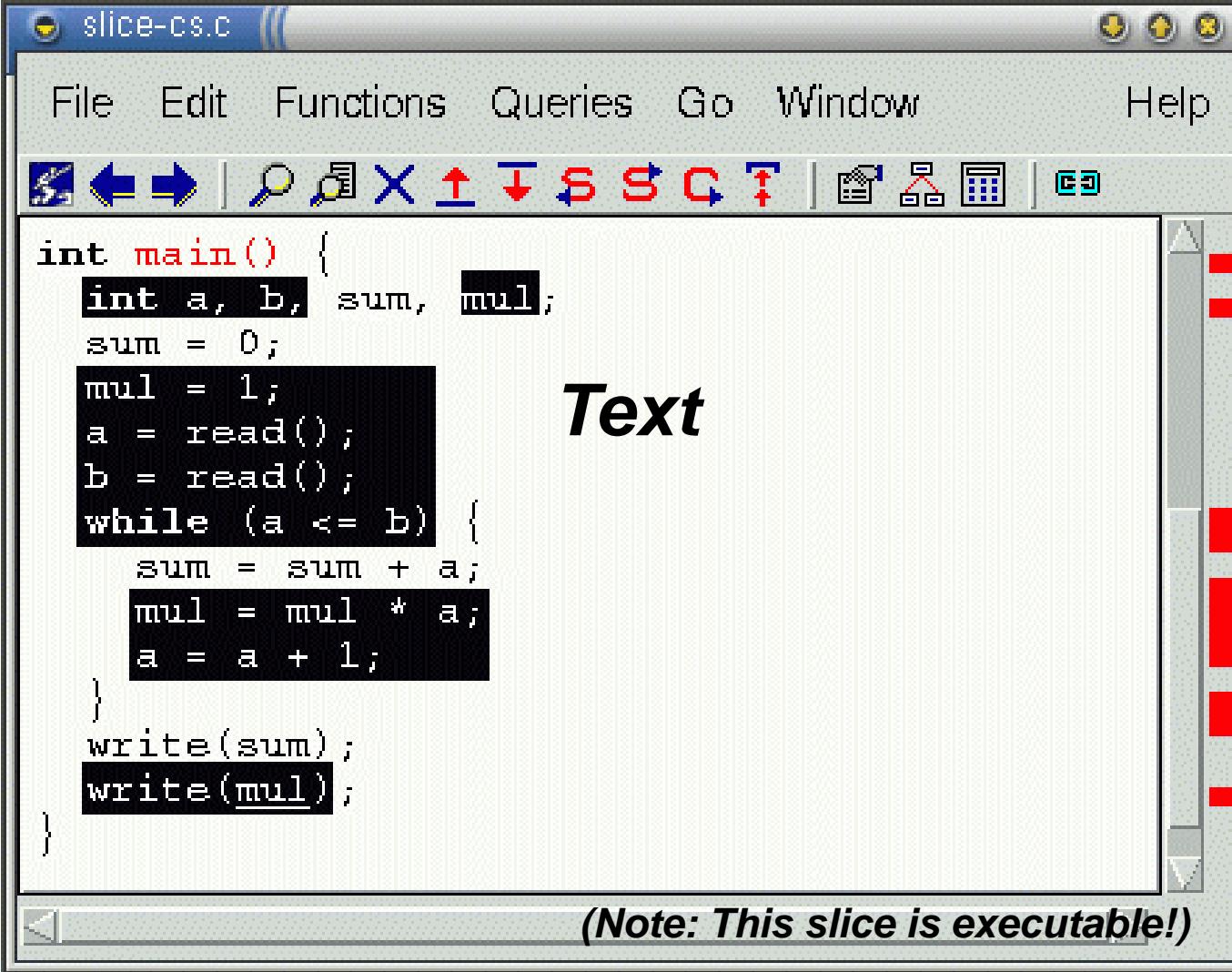
- Contains only the difference between two slices
- Useful for focusing on differing behavior

# Chop

- Intersection between a forward and a backward slice
- Useful for determining influence paths within the program



# Leveraging Slices



The screenshot shows a software window titled "slice-cs.c". The menu bar includes File, Edit, Functions, Queries, Go, Window, and Help. The toolbar contains various icons for file operations. The main text area displays the following C code:

```
int main() {
    int a, b, sum, mul;
    sum = 0;
    mul = 1;
    a = read();
    b = read();
    while (a <= b) {
        sum = sum + a;
        mul = mul * a;
        a = a + 1;
    }
    write(sum);
    write(mul);
}
```

A vertical red bar on the right side of the code editor highlights specific parts of the code. The word "Text" is overlaid in large black letters over the middle portion of the code. A note at the bottom of the code editor states: "(Note: This slice is executable!)"

# Concepts

- ★ **To reason about programs, use**
  - deduction (0 runs)
  - observation (1 run)
  - induction (multiple runs)
  - experimentation (controlled runs)
  
- ★ **To slice a program, follow dependences from a statement S to find all statements that**
  - could be influenced by S (forward slice)
  - could influence S (backward slice)

# Concepts

- ★ **Using deduction alone includes a number of risks, including code mismatch, abstracting away, and imprecision.**
- ★ **Any deduction is limited by the halting problem and must thus resort to conservative approximation.**
- ★ **For debugging, deduction is best combined with actual observation.**

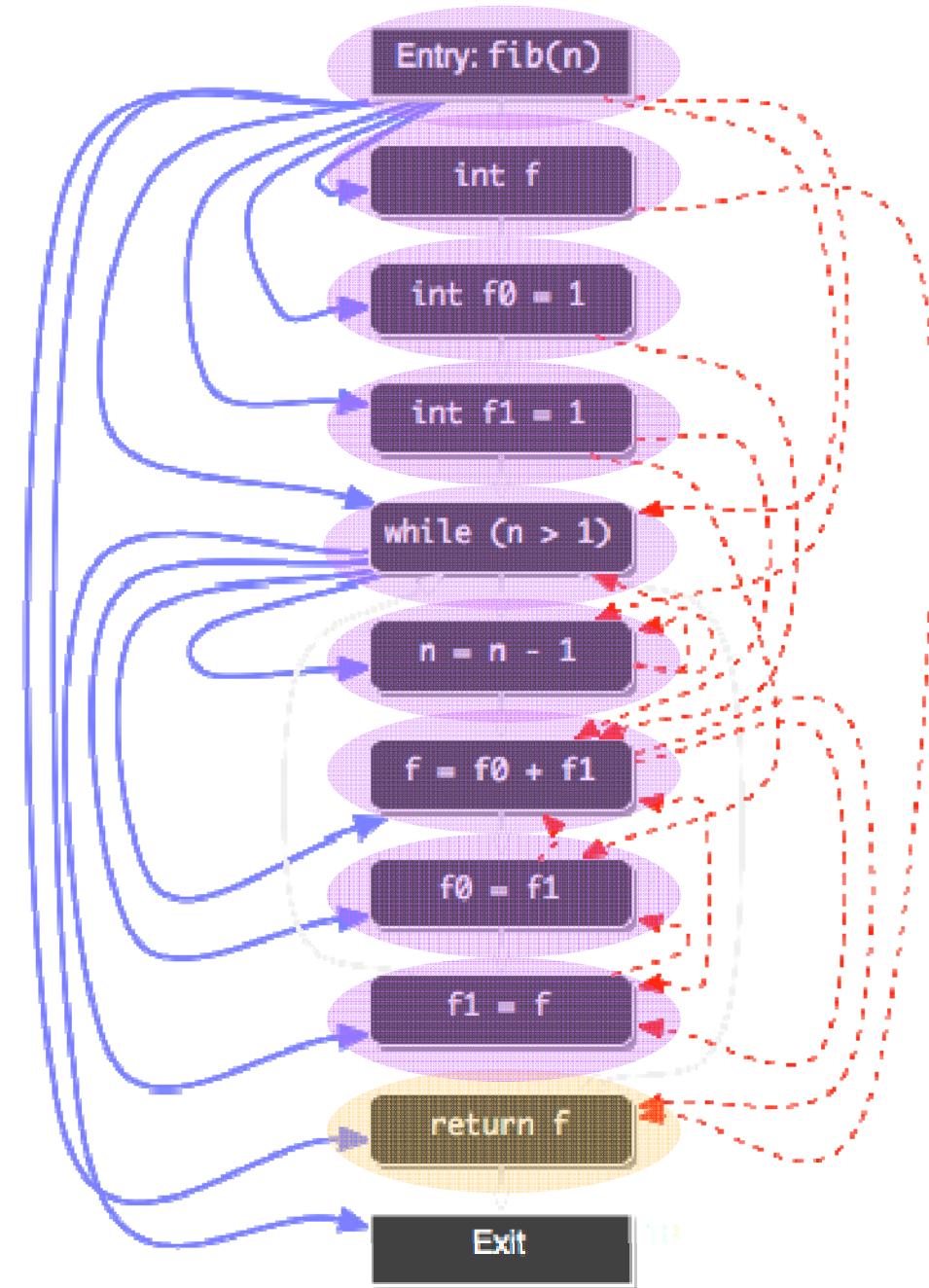
# Dynamic Slicing

- Static slices apply to *all* program runs:
  - General + reusable, but imprecise
- A *dynamic slice* applies to a *single run*:
  - Specific and precise

# Static Slicing

- Given a statement B, the backward slice contains all statements that could influence the read variables or execution of B
- Formally:

$Slice(B) = \{A | A \in P\}$



```
1 n = read();
2 a = read();
3 x = 1;
4 b = a + x;
5 a = a + 1;
6 i = 1;
7 s = 0;
8 while (i <= n) {
9   if (b > 0)
10    if (a > 1)
11      x = 2;
12    s = s + x;
13    i = i + 1;
14 }
15 write(s);
```

```
1 n = read(); // n = 2
2 a = read(); // a = 0
3 x = 1;
4 b = a + x;
5 a = a + 1;
6 i = 1;
7 s = 0;
8 while (i <= n) {
9   if (b > 0)
10    if (a > 1)
11      x = 2;
12    s = s + x;
13    i = i + 1;
14 }
15 write(s);
```

**Static slice for (s, 15) Dynamic slice for (s, 15)**

```

1 n = read();
2 a = read();
3 x = 1;
4 b = a + x;
5 a = a + 1;
6 i = 1;
7 s = 0;
8 while (i <= n) {
9   if (b > 0)
10     if (a > 1)
11       x = 2;
12   s = s + x;
13   i = i + 1;
14 }
15 write(s);

```

- 1. Obtain a *trace* of the execution**
- 2. Get the variables that are read and written**
- 3. Assign an empty slice to each written variable**
- 4. Compute the slices from start to end:**

$$\text{DynaSlice}(v_0) \leftarrow \bigcup_{v_1} (\text{DynaSlice}(v_1) \cup \{\text{line}(v_1)\})$$

Trace	Write	Read	Dynamic Slice
1 n = read();	n		
2 a = read();	a		$DynSlice(w) = \bigcup_i (DynSlice(r_i) \cup \{line(r_i)\})$
3 x = 1;	x		
4 b = a + x;	b	a, x	2, 3
5 a = a + 1;	a	a	2
6 i = 1;	i		
7 s = 0;	s		
8 while (i <= n) {	p8	i, n	6, 1
9 if (b > 0)	p9	b, p8	4, 2, 3, 8, 6, 1
10 if (a > 1)	p10	a, p9	5, 2, 9, 4, 2, 3, 8, 6, 1
12 s = s + x;	s	s, x, p8	7, 3, 8, 6, 1
13 i = i + 1;	i	i, p8	8, 6, 1
8 while (i <= n) {	p8	i, n	13, 8, 6, 1
9 if (b > 0)	p9	b, p8	4, 2, 3, 13, 8, 6, 1
10 if (a > 1)	p10	a, p9	5, 2, 9, 4, 2, 3, 13, 8, 6, 1
12 s = s + x;	s	s, x, p8	12, 7, 3, 6, 8, 1, 13
13 i = i + 1;	i	i, p8	13, 8, 6, 1
8 while (i <= n) {	p8	i, n	13, 8, 6, 1
15 write(s);	o15	s	12, 7, 3, 6, 8, 1, 13

Trace	Write	Read	Dynamic Slice
1 n = read();	n		
2 a = read();	a		
3 x = 1;	x		
4 b = a + x;	b	a, x	2, 3
5 a = a + 1;	a	a	2
6 i = 1;	i		
7 s = 0;	s		
8 while (i <= n) {	p8	i, n	6, 1
9     if (b > 0)	p9	b, p8	4, 2, 3, 8, 6, 1
10        if (a > 1)	p10	a, p9	5, 2, 9, 4, 2, 3, 8, 6, 1
12        s = s + x;	s	s, x, p8	7, 3, 8, 6, 1
13        i = i + 1;	i	i, p8	8, 6, 1
8 while (i <= n) {	p8	i, n	13, 8, 6, 1
9     if (b > 0)	p9	b, p8	4, 2, 3, 13, 8, 6, 1
10        if (a > 1)	p10	a, p9	5, 2, 9, 4, 2, 3, 13, 8, 6, 1
12        s = s + x;	s	s, x, p8	12, 7, 3, 6, 8, 1, 8, 13
13        i = i + 1;	i	i, p8	13, 8, 6, 1
8 while (i <= n) {	p8	i, n	13, 8, 6, 1
15 write(s);	o15	s	12, 7, 3, 6, 8, 1, 13

Trace	Write	Read	Dynamic Slice
1 n = read();		n	
2 a = read();		a	
3 x = 1;		x	
4 b = a + x;	b	a, x	2, 3
5 a = a + 1;	a	a	2
6 i = 1;	i		
7 s = 0;	s		
8 while (i <= n) {	p8	i, n	6, 1
9     if (b > 0)	p9	b, p8	4, 2, 3, 8, 6, 1
10        if (a > 1)	p10	a, p9	5, 2, 9, 4, 2, 3, 8, 6, 1
12        s = s + x;	s	s, x, p8	7, 3, 8, 6, 1
13        i = i + 1;	i	i, p8	8, 6, 1
8 while (i <= n) {	p8	i, n	13, 8, 6, 1
9     if (b > 0)	p9	b, p8	4, 2, 3, 13, 8, 6, 1
10        if (a > 1)	p10	a, p9	5, 2, 9, 4, 2, 3, 13, 8, 6, 1
12        s = s + x;	s	s, x, p8	12, 7, 3, 6, 8, 1, 8, 13
13        i = i + 1;	i	i, p8	13, 8, 6, 1
8 while (i <= n) {	p8	i, n	13, 8, 6, 1
15 write(s);	o15	s	12, 7, 3, 6, 8, 1, 13

```
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3 x = 1;
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8 while (i <= n) {
9   if (b > 0)
10    if (a > 1)
11      x = 2;
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14 }
15 write(s);
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```
1 n = read(); // n = 2
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3 x = 1;
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8 while (i <= n) {
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10    if (a > 1)
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12    s = s + x;
13    i = i + 1;
14 }
15 write(s);
```

**Static slice for (s, 15) Dynamic slice for (s, 15)**

# Discussion

- Dynamic slices are much more precise than static slices (applied to the one run, that is)
- From some variable, a backward slice encompasses on average
  - 30% of the *entire* program (static slice)
  - 5% of the *executed* program (dynamic slice)
- Overhead as in omniscient debugging

The Whyline

Your world is paused.

further questions can be asked

camera focuses on subject of question

tooltips show properties' current values

access to previous questions and answers

code related to the selection is highlighted

runtime actions

causality arrows

time cursor traverses execution history

**Pac's details**

- properties: current direction = forward
- methods: create new variable, capture pose, color = yellow, opacity = 1 (100%), vehicle = World

**World.move Pac**

No variables

Pac move Pac.current direction 3 meters duration = 1 second style = gentle true

If both Pac is within 2 meters of Ghost and not Big Dot.isEaten

Pac resize 0.5

**Question: Why didn't Pac resize 0.5?**

**Answer:**  
One or more of these actions prevented Pac resize 0.5 from happening.  
Try following the arrows and checking each action to find out what went wrong.

3.821010  
Big Dot.isEaten set to true → isEaten → Pac is within 2 of Ghost → and → Doing else X

3.854011  
false

This image shows a Scratch interface with several annotations explaining its features. At the top, there are buttons for Resume, Stop, Why..., Undo, Redo, and a trash bin. A message says 'Your world is paused.' Below the stage, a script is shown: 'When Pac is within 1 meter of Big Dot becomes true' which runs 'Do this once, when it becomes true'. Inside, it has two blocks: 'Do in order' and 'Big Dot set isShowing to false' followed by 'Big Dot.isEaten set value to true'. A callout 'further questions can be asked' points to the 'Why...' button. Another callout 'camera focuses on subject of question' points to the stage area where a yellow Pac sprite is located. On the left, 'Pac's details' panel shows properties like 'current direction = forward' and methods like 'create new variable'. A tooltip 'tooltips show properties' current values' points to the 'current direction' property. In the center, a script for 'World.move Pac' is shown with a tooltip 'access to previous questions and answers' pointing to the 'If' condition. A callout 'code related to the selection is highlighted' points to the 'Doing else' block. At the bottom, a 'Question: Why didn't Pac resize 0.5?' section has an 'Answer' with troubleshooting steps. A causality diagram shows 'Big Dot.isEaten set to true' leading to 'isEaten', which then leads to 'Pac is within 2 of Ghost', then 'and', and finally 'Doing else X'. A time cursor is shown traversing the execution history between frame numbers 3.821010 and 3.854011.

**Resume** **Stop** **Why...** **Undo** **Redo** **Your world is paused.**

further questions can be asked

camera focuses on subject of question

tooltips show properties' current values

access to previous questions and answers

code related to the selection is highlighted

runtime actions

causality arrows

time cursor traverses execution history

**Pac's details**

**properties** **methods** **questions**

**current direction** = forward

**create new variable**

**capture pose**

**color** = yellow

**opacity** = 1 (100%)

**vehicle** = World

**World.move Pac** No parameters

No variables

**create new parameter**

**create new variable**

**Pac** move Pac.current direction 3 meters duration = 1 second style = gentle true

**If both** Pac is within 2 meters of Ghost and not Big Dot.isEaten

**Pac** resize 0.5

**Do in order** **Do together** **If/Else** **Loop** **While** **For all in order** **print**

**Questions I've asked**

**Answer:**  
One or more of these actions prevented Pac resize 0.5 from happening.  
Try following the arrows and checking each action to find out what went wrong.

3.821010  
**Big Dot.isEaten set to true**

3.854011  
**Doing else**

**Your world is paused**

questions are chosen from a hierarchical menu

code related to the question is highlighted

selection is highlighted

runtime actions

causality arrows

time cursor traverses execution history

Answer:  
One or more of these actions prevented Pac resize 0.5 from happening.  
Try following the arrows and checking each action to find out what went wrong.

The image shows a Scratch interface with a paused world. At the top, there are buttons for Resume, Stop, Why..., Undo, and Redo. A message "Your world is paused" is displayed. On the left, there's a sprite list with Light, Ground, Pac, Ghost, Dot1, Dot2, and Dot3. Below it is a "Pac's details" panel with tabs for properties, methods, and questions. The properties tab shows current direction = forward, and the methods tab shows World.move Pac. The questions tab is active, showing a hierarchical menu with options like "Why did...", "Why didn't...", and "What happened while the world was running?". An arrow points from this menu to the main workspace where a blue Pac sprite and a yellow ghost are visible. The workspace contains a script for Pac with blocks like move forward 3 steps, resize 0.5, and pointOfView change to something else?. A comment "Do this once, when it becomes true" is present. To the right of the workspace is a script editor with blocks for Big Dot. A callout bubble highlights the "questions" tab with the text "questions are chosen from a hierarchical menu". In the center, a script for World.move Pac is shown with a conditional block "If both [Pac is within 2 meters of Ghost] and [not Big Dot.isEaten] then [Pac resize 0.5 more...]" and an "Else" block. A callout bubble highlights the conditional block with the text "code related to the question is highlighted". Another callout bubble points to the "Else" block with the text "leads to previous questions and answers". At the bottom, a "Questions I've asked" panel shows two entries: "Big Dot.isEaten set to true" at time 3.821010 and "Doing else" at time 3.854011. A causality graph shows arrows between these events and the execution history, with labels like "true", "false", and "and". A callout bubble highlights the "Questions I've asked" panel with the text "selection is highlighted". A callout bubble highlights the "time cursor traverses execution history" at the bottom with the text "time cursor traverses execution history".

# “Why did” questions

- Take the dynamic slice of the variable
- Follow at most two dependencies
- If programmer wants, follow dependencies transitively

# “Why did $s = 2$ in Line 15?”

```
1 n = read(); // n = 2
2 a = read(); // a = 0
3 x = 1;
4 b = a + x;
5 a = a + 1;
6 i = 1;
7 s = 0;
8 while (i <= n) {
9     if (b > 0)
10        if (a > 1)
11            x = 2;
12        s = s + x;
13        i = i + 1;
14    }
15 write(s);
```

“Because  $s = 1$  and  $i = 2$ ”

# **“Why didn’t” questions**

- Follow back control dependencies to closest controlling statement(s)
- Do a “why did” question on each
- Again, follow at most two dependencies

# “Why didn’t $x = 2$ in Line 11?”

```
1 n = read(); // n = 2
2 a = read(); // a = 0
3 x = 1;
4 b = a + x;
5 a = a + 1;
6 i = 1;
7 s = 0;
8 while (i <= n) {
9   if (b > 0)
10    if (a > 1)
11      x = 2;
12   s = s + x;
13   i = i + 1;
14 }
15 write(s);
```

**“Because  $a = 1$  and  $b = 1$ ”**

# Discussion

- The WHYLINE combines
  - omniscient debugging
  - static slicing
  - dynamic slicing
- in an attractive package, showcasing the state of the art in interactive debugging

# Tracking Infections

- 1. Start with the infected value as seen in the failure**
- 2. Follow back the dependencies**
- 3. Observe and judge origins – are they sane?**
- 4. If some origin is infected, repeat at Step 2**
- 5. All origins are sane? Here's the infection site!**

# Concepts

- ★ **Omniscient debugging allows for simple exploration of the entire execution history**
- ★ **Dynamic slicing tells the origin of a value**
- ★ **To track down an infection, follow dependencies and observe origins, repeating the process for infected origins**