

1.設計

A. System call :

- 1.get_time : using getnstimeofday() to get the start time and using copy_to_user to write the time into user space buffer.
- 2.prinf_ans : similar to get_time but adding printk to write 2 time_string to dmesg.

B. main :

1.pcb structure

I use circular link list to maintain the schedule struct which store the process status of fork process.

2. priority setting :

- a. for scheduler(parent) :
 - except idle : set its priority as 10 (2nd order)
 - idle : set its priority as 1 (3rd order) , only happen with no job in queue
- b. for child process unselected : set its priority as 1 (3rd order)
- c. for child process selected (only one in same time): set its priority as 99 (1st order)

3. scheduler process : follow is a while loop and finish when all task finish

- a. if next job don't ready and no task in queue => idle until timer become ready time
 - b. check if there are some task ending and remove it from the queue
 - c. create tasks(fork) whose ready time < timer (using a while loop) , and insert it to the queue (pending in the tail) , at the same time I would set child process's priority as 1
 - d. base on the queue , select the task and set its priority as 99
 - FIFO : the head of the queue
 - RR : head->next of the queue , and let head = head->next
 - SJF / PSJF: the shortest exec_time in the queue
- child process which been selected run-----
- back to schedule when :
- FIFO / SJF : finish
 - RR : finish or running time = 500
 - PSJF : timer = next job's ready time
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- e. (only PSJF) renew the exec_time depend on the running time it run

2.核心版本

linux 4.14.25

3.比較實際結果與理論結果，並解釋造成差異的原因

在順序上，結果與理論相同，但在運行時間上，因為 context swicth 的 overhead，導致部分 TASK 時間並沒有明顯的節省 ex : PSJF 理論上時間要小於 SJF，但部分 TASK 無顯著差異