Gramática sin problemas de factorización ni recursividad izquierda:

P→Block#

Block → ConstDecl VarDecl ProcDecl Statement

ConstDecl \rightarrow const ConstAssigList;| λ

ConstAssigList→ id = num ConstAssigList`

ConstAssigList \rightarrow , id = num ConstAssigList λ

 $VarDecl \rightarrow var IdList ; |\lambda|$

IdList \rightarrow id IdList IdList` \rightarrow ,id IdList`| λ

ProcDecl→procedure id; Block; ProcDecl|\(\lambda\)

Statement → id :=Expression|call id|begin StatementList end|if Condition then Statement|while

Condition do Statement|\(\lambda\)

StatementList→ Statement StatementList`

StatementList \rightarrow ;Statement StatementList $\mid \lambda$

Condition→Expression Relation Expression|odd Expression

Expression→ SumOperator Term Expression` |Term Expression`

Expression'→ SumOperator Term Expression'|\(\lambda\)

Term→ Factor Term`

Term'→MultOperator Factor Term'|\(\lambda\)

Relation $\rightarrow = |<>|<|>|<=|>=$

SumOperator→+|-

MultOperator→/|*

Factor→(Expression)|id|num

Tabla de análisis sintáctico correspondiente:

	Program	Block	ConstDec 1	ConstAssi gList	ConstAssi gList`	VarDecl	IdList
Const	P→Block #	Block → ConstDec l VarDecl ProcDecl Statement	ConstDec 1 → const ConstAssi gList;				
#	P→Block #	Block → ConstDec l VarDecl ProcDecl Statement	ConstDec $1 \rightarrow \lambda$			VarDecl → λ	
var	P→Block #	Block → ConstDec l VarDecl ProcDecl Statement	ConstDec $1 \rightarrow \lambda$			VarDecl → var IdList;	
procedure	P→Block #	Block → ConstDec l VarDecl ProcDecl Statement	ConstDec $1 \rightarrow \lambda$			VarDecl → λ	
id	P→Block	Block →	ConstDec	ConstAssi		VarDecl	IdList →

	#	ConstDec 1 VarDecl ProcDecl Statement	$1 \rightarrow \lambda$	gList→ id = num ConstAssi gList`		$\rightarrow \lambda$	id IdList`
call	P→Block #	Block → ConstDec l VarDecl ProcDecl Statement	ConstDec $1 \rightarrow \lambda$			$\begin{array}{c} \text{VarDecl} \\ \rightarrow \lambda \end{array}$	
begin	P→Block #	Block → ConstDec l VarDecl ProcDecl Statement	ConstDec $1 \rightarrow \lambda$			VarDecl → λ	
if	P→Block #	Block → ConstDec l VarDecl ProcDecl Statement	ConstDec $1 \rightarrow \lambda$			VarDecl → λ	
while	P→Block #	Block → ConstDec l VarDecl ProcDecl Statement	ConstDec $1 \rightarrow \lambda$			$\begin{array}{c} \text{VarDecl} \\ \rightarrow \lambda \end{array}$	
;		Block → ConstDec l VarDecl ProcDecl Statement	ConstDec $1 \rightarrow \lambda$		ConstAssi gList→ λ	$\begin{array}{c} \text{VarDecl} \\ \rightarrow \lambda \end{array}$	
,					ConstAssi gList→, id = num ConstAssi gList`		
	IdList`	ProDecl	Statement	Statement List	Statement List`	Condition	Expressio n
Const							
#		$\begin{array}{c} ProDecl \\ \rightarrow \lambda \end{array}$	Statement $\rightarrow \lambda$				
var							
procedure		ProcDecl →proced					

	T	I	<u> </u>	<u> </u>	T		<u> </u>
		ure id ; Block ; ProcDecl					
id		ProDecl $\rightarrow \lambda$	Statement → id :=Express ion	Statement List→ Statement Statement List`			
call		ProDecl $\rightarrow \lambda$	Statement →call id	Statement List→ Statement Statement List`			
begin		ProDecl $\rightarrow \lambda$	Statement →begin Statement List end	Statement List→ Statement Statement List`			
if		ProDecl $\rightarrow \lambda$	Statement →if Condition then Statement	Statement List→ Statement Statement List`			
while		ProDecl $\rightarrow \lambda$	Statement →while Condition do Statement	Statement List→ Statement Statement List`			
;	IdList → ,id IdList`	ProDecl → λ	Statement $\rightarrow \lambda$		Statement List`→ ;Statemen t Statement List`		
end					Statement List` $\rightarrow \lambda$		
,	IdList → ,id IdList`						
+						Condition →Expres sion Relation Expressio n	Expressio n→ SumOper ator Term Expressio n`
(Condition	Expressio

$\begin{array}{c ccccccccccccccccccccccccccccccccccc$				→Expres sion Relation Expressio n	n→ Term Expressio n`
→Expres sion Relation Expression Odd Condition → odd Expression Condition → odd Expression Condition → Expression Relation Expression The expression is considered and in the expression of the expression is considered at the expression of the	id			→Expres sion Relation Expressio	Expressio n→ Term Expressio n`
$\begin{array}{c} \longrightarrow \text{odd} \\ \text{Expressio} \\ \text{n} \end{array}$	num			→Expres sion Relation Expressio	Expressio n→ Term Expressio n`
→Expres n→ sion SumOp Relation ator Ter Expressio Express	odd			→ odd Expressio	
	-			→Expres sion Relation Expressio	SumOper ator Term Expressio

	Expressio n'	Term	Term`	Relation	SumOper ator	MultOper ator	Factor
#	Expressio n`→λ						
*			Term'→ MultOper ator Factor Term'			MultOper ator→*	
=	Expressio n`→λ			Relation →=			
\Leftrightarrow	Expressio n`→λ			Relation →<>			
<	Expressio n`→λ			Relation →<			
>	Expressio			Relation			

	n`→λ			→ >			
<=	Expressio n'→λ			Relation →<=			
>=	Expressio n'→λ			Relation →>=			
;	Expressio n`→λ						
num		Term→ Factor Term`					Factor→n um
/			Term'→ MultOper ator Factor Term'			MultOper ator→/	
+	Expressio n'→ SumOper ator Term Expressio n'		Term`→λ		SumOper ator→+		
(Term→ Factor Term`					Factor→(Expressio n)
id		Term→ Factor Term`					Factor→i d
do	Expressio n'→λ						
then	Expressio n'→λ						
-	Expressio n'→ SumOper ator Term Expressio n'		Term`→λ		SumOper ator→-		
)	Expressio n`→λ						