Distributed Systems Direct Communication PART II

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October 3, 2023















This chapter: overview

- Part I
 - Data representation
 - Message passing
 - Request-reply protocols
 - Remote procedure calls
- Part II
 - Object request brokers



Object request brokers

• Basics:

Object request broker =

Objects +

RPC

+ ...



Object request brokers

• Examples:

- CORBA (Common Object Request Broker Architecture)
- DCOM (Distributed Component Object Model)
- Java RMI
- NET Remoting



Object request brokers

- Overview
 - distributed object systems
 - case study: Java RMI

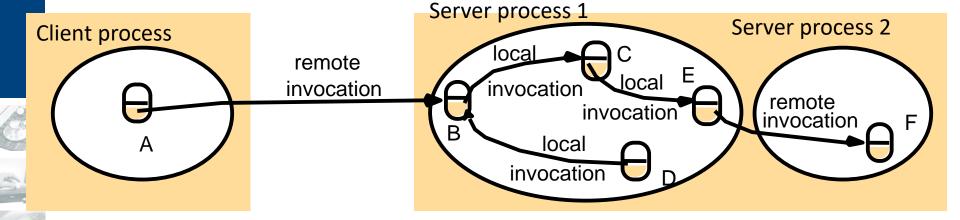


ORB: Distributed object systems

- Why objects?
 - **—** ...
- Why distributed objects?
 - Maps naturally to distributed services
 - Services are modeled as objects
 - Functionality emerges by cooperation between services and clients



Local ⇔ remote invocation





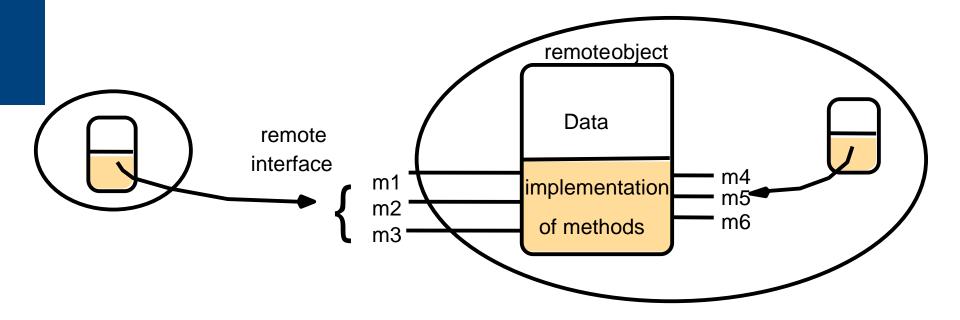
- Requirements
 - synchronous invocation semantics
 - location transparency: location transparent object references
 - access transparency: local & remote object invocations identical
 - inheritance
 - polymorphism



- A distributed object consists of
 - an implementation-independent service definition specified by an interface
 - method signature
 - multiple inheritance
 - a concrete service implementation
 - implements one or more interfaces
 - provides the concrete implementation



Remote object & its remote interface





- Distributed objects act like normal objects
 - can be used as parameters and return values
 - can be invoked (across address space boundaries)
 - Location transparency
 - Access transparency
 - can be subclassed (inheritance)
 - are invoked synchronously



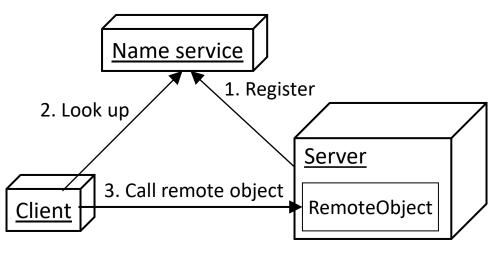
- But there is some complexity (non-transparent!)
 - Invocation semantics
 - Java RMI: at most once
 - CORBA: at most once or maybe (if no result)

⇔ exactly once

- Parameter passing
 - distributed objects are passed by reference (≠ pointer)
 - normal objects are passed by value
- Additional exceptions due to distribution
 - distribution is not transparent
- Concurrent access
- Latency (performance)



- Typical architecture of distributed object systems:
 - 2 separate processes: client & server
 - Server hosts remote object(s) and waits for incoming calls.
 - Client obtains references to remote objects and makes calls.
 - Name service
 - Helps clients to locate remote objects from their remote object references. (name ⇒ address)





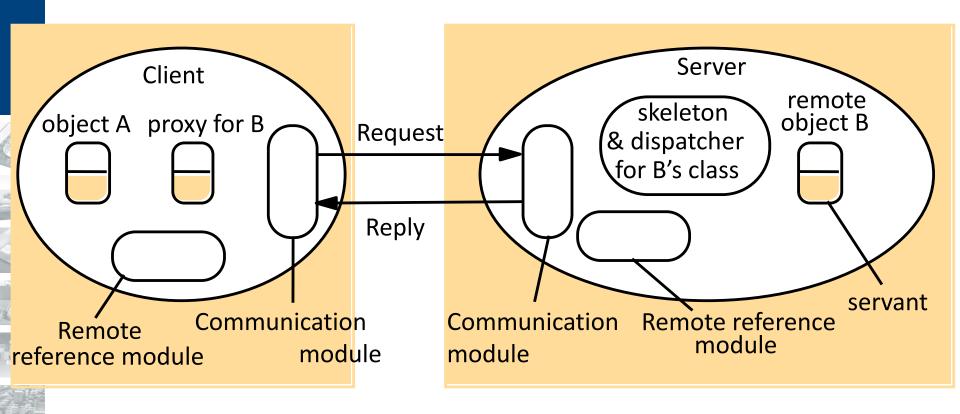
- Distributed garbage collection
 - Garbage collection of remote objects.
 - In cooperation with local garbage collector.
- Example: Java distributed garbage collection algorithm:
 - Reference counting:

```
if distributed reference count == 0
    garbage collect remote object
```

- Leasing-based: client obtains lease for a period of time.
 - → Tolerates failure of client processes
 - → Tolerates failure of communication

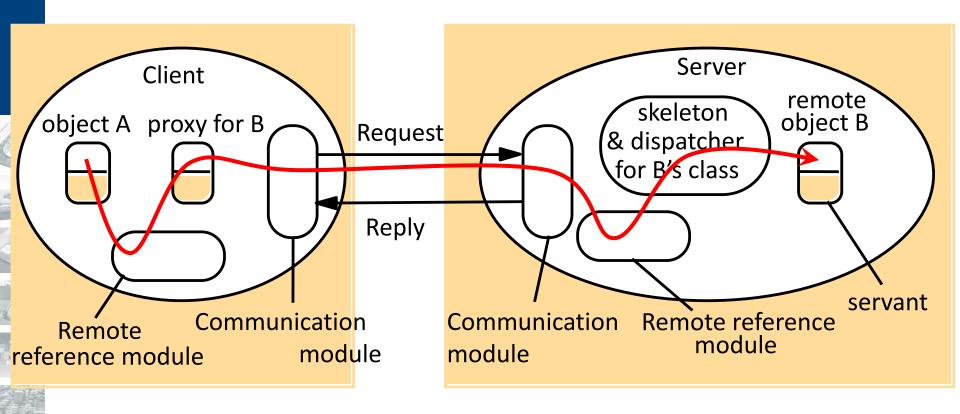


ORB: Object invocations

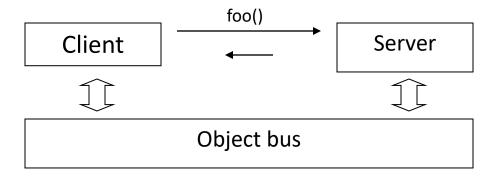




ORB: Object invocations

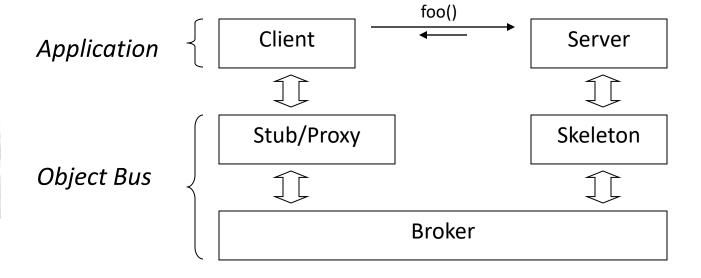






Client - the invoking application object Server - the invoked application object

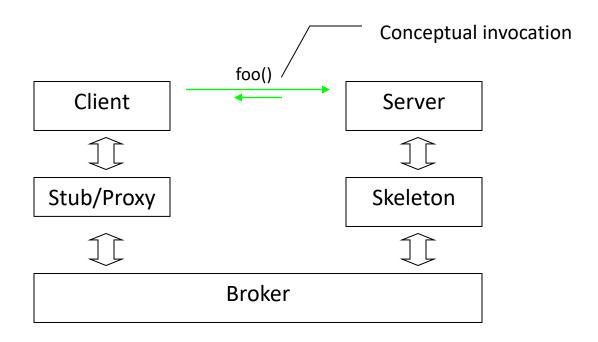






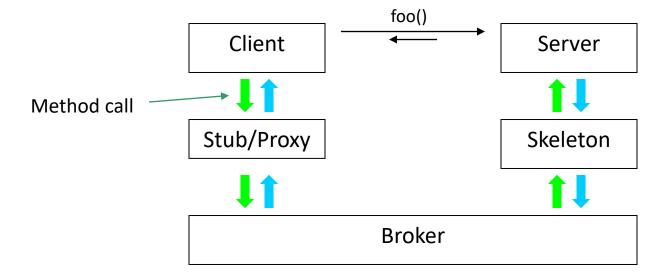
- Client the invoking application object
- Server the invoked application object
- Stub/Proxy reifies the invocation
- Skeleton dispatches the invocation to the actual object implementation
- Broker invocation distributor





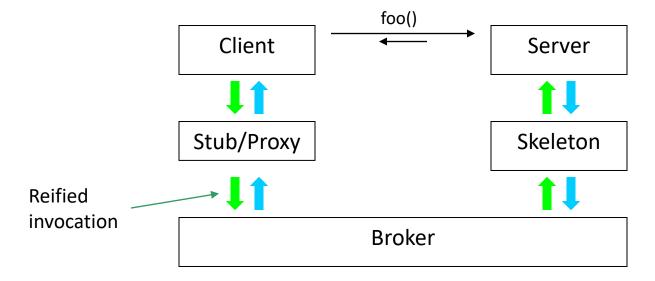


- Look up remote object (in <u>remote reference module</u>)
 - Not local ⇒ Create / return stub
 - Invoke stub



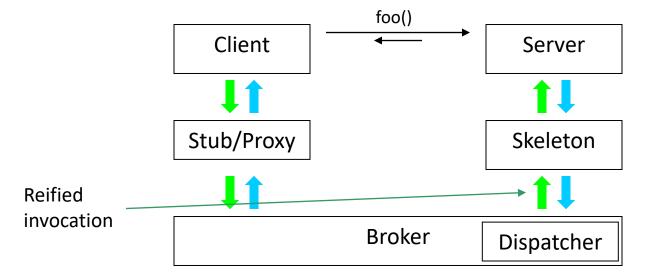


- Forward invocation in a message to the remote object
 - Marshalling of object reference, method name and arguments into request message



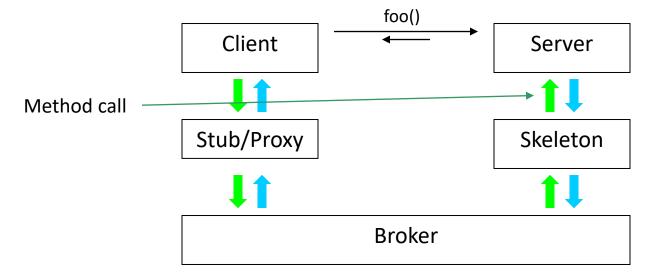


- Look up remote object (at server side)
 - Local ⇒ pass local reference of target object to dispatcher
 - Dispatcher sends request to appropriate method in skeleton



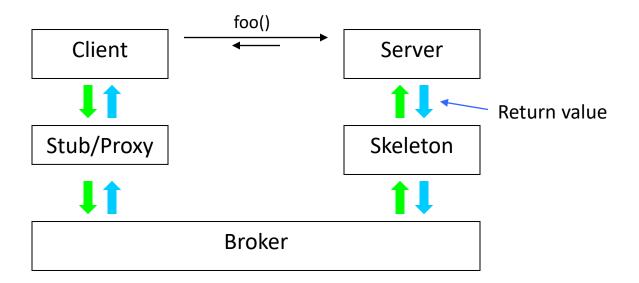


- Forward invocation to actual remote object (i.e. server)
 - Unmarshalling arguments in request message
 - Invoking corresponding method of target object



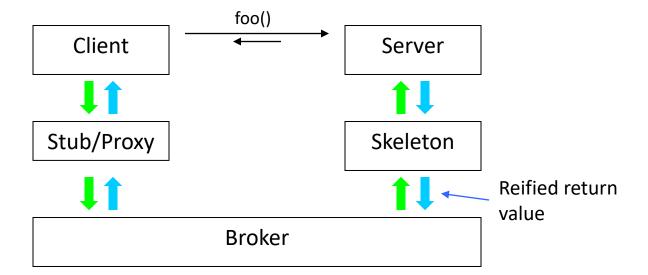


Return result to skeleton



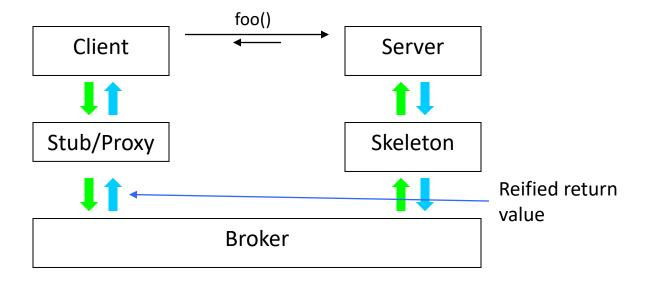


- Forward return in a message to the waiting stub
 - Marshalling of return value (and exceptions) into reply message



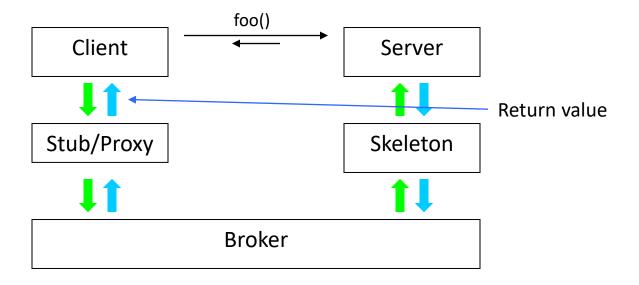


• Waiting for *reply* message





- Return result to invoking object (i.e. client)
 - Unmarshalling reply message





Questions?













