Exercises for the course Cryptography and Network Security

[Exercise II]

carl.bootland@esat.kuleuven.be

Public Key Cryptography

I. General Exercises

- 1. Compute 9⁹ mod 133 using repeated squaring.
 - Extra: Compute $9^3 \mod 133$, what can you deduce about $\lambda(133)$ without factoring 133? (Recall that for a positive integer n we have $a^{\lambda(n)} \equiv 1 \mod n$ for all a coprime to n.)
- 2. Compute efficiently $\varphi(100)$.
- 3. Compute $10^{82} \mod 33$ without using repeated squaring. Is 33 a strong pseudoprime to the base 10?
- 4. Compute gcd(1624, 6363).
- 5. Compute $101^{-1} \mod 195$
- 6. Compute $5^{-1} \mod 8$ using Euler's generalisation of Fermat's Theorem.
 - Extra: What are the inverses of 3 and 7 modulo 8? Deduce the value of $\lambda(8)$? (In general, the Carmichael function of n, denoted $\lambda(n)$, is defined to be the minimal natural number λ such that $a^{\lambda} \equiv 1 \mod n$ for all a coprime with n.)
- 7. Find the smallest non-negative solution to the following system of congruences:

 $x \equiv 4 \mod 19$

 $x \equiv 7 \mod 11$

 $x \equiv 1 \mod 7$.

II RSA encryption and decryption

- 1. (a) Consider an RSA encryption-system with modulus $n = 629 = 37 \cdot 17$. Choose the smallest possible public exponent. What is the corresponding secret exponent?
 - (b) Calculate the ciphertext for the message '591'. Is there a problem?
 - (c) Decrypt your result using the Chinese Remainder Theorem; verify whether you retrieve the plaintext.
 - (d) Is it a problem to have a common modulus for RSA? That is, suppose a message m is encrypted with both of the public keys (n, e) and (n, f), with gcd(e, f) = 1, then can you break RSA given these two ciphertexts $c_e = m^e \mod n$ and $c_f = m^f \mod n$. If so, how? (Hint: The Extended Euclidean Algorithm allows one to find integers x and y such that gcd(a, b) = xa + yb.)
 - (e) Is a small public exponent e a security problem? Why? (Hint: Suppose an entity wishes to send the same message m to three entities whose public moduli are n_1, n_2, n_3 , and whose encryption exponents are all e = 3.)
- 2. Prove that if n = pq is a product of two primes, then determining $\varphi(n)$ is equivalent to factoring n.

Extra Questions on RSA

- 3. Suppose you have somehow gained access to the decryption device of a competitor who is using the basic RSA encryption scheme. You notice that their decryption times are extremely fast and you suspect their decryption exponent d is very small. You know that their public key is (e, n) = (31, 247) and have intercepted the ciphertext c = 23. Decrypt c without factoring n.
- 4. You are told that the ACME RSA software builds an RSA modulus by choosing a random k-bit prime p and setting q to be the smallest prime larger than p. You manage to get hold of an RSA public key (e,n) = (1921, 141367) and a ciphertext c = 70918 of someone using the ACME software. Decrypt the ciphertext.

III ElGamal

- 1. Given p = 89, a = 3, y = 69. Verify for a message m = 77 whether the signature (r, s) = (66, 77) is a valid signature.
- 2. Show that a different random number k must be selected for each message signed; otherwise the private key x can be determined with high probability.