

IN2009

Language Processors

Introduction to Language **Processors**

Igor Siveroni

IN2009

- · Instructor: Igor Siveroni
 - Research Associate/Visiting Lecturer
 - Room A306, College Building
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- · Pre-Requisite: Java
- · Lectures: 11.00-12.50 at C350
- · Tutorial/Labs
 - 13.00-13.50 at A217
 - 15.00-15.50 at A218

Dates for your diary...

- Lectures (10):
 - Weeks 1-5, 7-11.
 - No lecture weeks 6 and 12.
- Labs (10):
 - Start this week!
 - Weeks 1-5, 7-11.
 - No labs weeks 6 and week 12.

This Week

Week 1

- · Module details, aims, resources
- · What is language processing?
- · Introduction to syntax definition
 - Example: A straight-line programming language
- Abstract syntax trees
- · ...and some Java.

Module Details

Support

- Online support is through Cityspace: http://www.city.ac.uk/cityspace
 - · E.g. Digital copies of all lectures and labs.
- Use discussion boards to post questions.
- Questions (except personal issues) sent by email will be cut & pasted into the discussion boards on Cityspace.
- ...with up to a week's delay!

Core Text

· Appel, A., "Modern Compiler Implementation in Java", 2nd ed., 2002, Cambridge University Press, ISBN-13: 978-0521820608

Appel (2002)



Assessment

- · In-module coursework tasks
 - 3 practical assessments (weeks 4 11).
 - 30% of module marks.
 - Pair working is allowed, but you must declare if you do so! More on this at a later date
- End-of-module exam
 - Unseen, written 1.5-hour paper.
 - 70% of module marks (min threshold 30%).
 - In May/June 2009.

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Coursework

- The first assessment brief will be published on Cityspace next week.
- · Deadlines:
 - All deadlines are by 5pm on Friday of the indicated week.
- Lateness Penalties
 - Assignments handed in after the deadline will receive a 0.
 - Special circumstances will be reviewed by a committee.

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Coursework

- Coursework marks and feedback will be returned to you via Cityspace.
 - Usually within two weeks of the deadline.
- It is my intention to produce a guideline answer for all assessments - published around 3 weeks after each deadline.
 - Not a "model" solution, but one which is fit for purpose and would have achieved approximately 50-60%

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Module Details Language Processors - Why?

- Programming Languages is one of the cornerstones of Computer Science
- Programming Languages design and implementation covers a wide range of theoretical and practical aspects of Computer Science.

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Module Details - Why?

- "An understanding of how to write PL definitions, and how to turn a definition into a recogniser and translator for the language, will give students an understanding of programming language structure and implementation that will compliment programming skills and aid the learning of new programming languages."
- "Also, an understanding of the run-time environment for the translated program will give insight into how high-level programs execute at machine level."

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Module Aims

"To introduce students to the specification and implementation of programming languages. In particular the module aims to provide an introduction to the structure of programming languages, the algorithms and data structures used in compilers, the tools which may be used to automate compiler construction, and to the run time environments in which programs execute."

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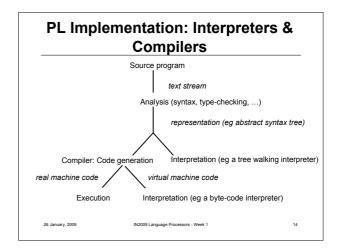
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Indicative Content

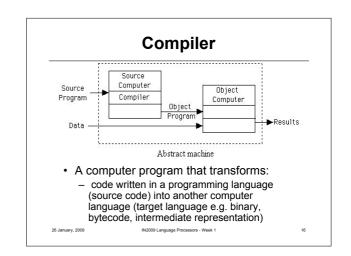
- Syntax definition using regular and context free grammars
- Abstract and concrete syntax, and abstract syntax tree representations
- · Introduction to type-checking
- Translation
- · Runtime environments

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Interpreter Source Program Data Abstract machine • A computer program that executes instructions written in a Programming Language. 15



Compilers & Interpreters

- But only rarely do we have a pure interpreter or compiler.
- Typically code is first compiled to intermediate form. Then ...
 - It is interpreted, or
 - code is generated for a virtual machine interpreter, a real machine or even another programming language.
- Interpreters may also perform internal transformations.

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Example: Java implementation Java Compiler (produces Java bytecode) ₹⋝ Class loader Java library classes Java Virtual Machine Bytecode verifier (Java bytecode) (JVM) Bytecode interpreter Optional Just In Time (JIT) Compiler JVM implemented on many platforms (eg in Unix PC Mac OS X

Language processing: Compiler source code lexical analysis lexical items (tokens) syntactic analysis abstract syntax tree semantic analysis symbol table (eg type-checking) other tables annotated tree or optimisation code generation object code from object code (for virtual or real (eg libraries) linking and loading executable object code

Concepts and Techniques

Concepts and Techniques used:

- Language Specification
 - Syntax: Regular Expressions and Grammars
 - · Semantics: Formal Methods
- Lexical and Syntactic Analysis: Finite Automata and Parsing. Parser generators.
- Semantics Analysis and Optimisation:
 - Type Checking, Intermediate representation.
 - Control and Data Flow Analysis (Program Analysis)
- Code Generation: Runtime environments, machine language, operating systems, garbage collection, etc.
- Programming!!!!

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Learning Outcomes

- On successful completion of this module, you will be able to:
 - Use formal languages to define input language syntax.
 - Explain techniques for syntactic and semantic analysis, and translation.
 - Explain the compiled code, and run-time environment requirements, for various common programming language structures.
 - Program data structures and algorithms for representation and analysis and translation of programming languages.
 - Use standard compiler generation tools.

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Lectures

- 1. Introduction to Language Processing (this week)
- 2. Language processing & lexical analysis
- 3. Parsing I (syntax analysis)
- 4. Parsing II (abstract syntax)
- 5. MiniJava abstract syntax trees
- 6. Semantic analysis
- 7. Activation records (stack frames)
- 8. Translation to intermediate representation

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What we will do in this module...

- Work through Appel (2002)
 - Leaving out much of the theory
- Define a simple programming language
- Implement some example abstract trees for the simple programming language
- · Introduce a more complicated language
- Implement a lexical analyser using JavaCC

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What we will do in this module...

- Implement a syntax analyser using JavaCC
- Implement an abstract syntax tree builder
- Look at semantic analysis (e.g. typechecking)
- Look at runtime environments and translation

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Intermission

- Since this is a 2 hour lecture, it is time for a short break...
 - Please be seated and ready to continue in 10 minutes time (that doesn't mean arrive back at the room in 10 minutes time!)

Lecture recommences at:

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Syntax definition

 Suppose we want to implement a compiler for a language that looks as follows:

```
a := 5+3;
b := (print(a, a-1), 10*a);
print (b)
```

- · How do we:
 - · Specify the language's syntax?
 - · Represent a program internally?

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Syntax definition

 Use context-free grammars (e.g. Backus–Naur Form or BNF) to define a grammar for the language

```
| Stm → id := Exp | Stm ; Stm | print(ExpList) | Exp → id | num | Exp Binop Exp | ...
```

| "or" - separates alternatives

→ "is defined as"

Exp, Stm non-terminals

print, id, num, := terminals (tokens)

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Syntax definition

- Each definition is called a production; many productions define a grammar
- · Repetition by recursive definition
- You may have seen BNF before...
 - ...any MySQL users? The manual and instruction list shows the statement syntax in BNF.
 - Java specification.

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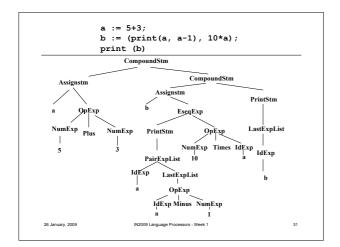
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Straight-line programming language

```
Stm; Stm
Stm
                                        (CompoundStm)
Stm
                                       (AssignStm)
             id := Exp
Stm
             print (ExpList)
                                        (PrintStm)
                                       (IdExp)
Exp
             id
                                       (NumExp)
Exp
             num
             Exp Binop Exp
                                       (OpExp)
Exp
Exp
             (Stm , Exp)
                                        (EseqExp)
ExpList→
             Exp , ExpList
                                        (PairExpList)
.
ExpList→
                                       (LastexpList)
             Exp
Binop →
                                        (Plus)
                                       (Minus)
Binop →
Binop →
                                        (Times)
Binop →
                                        (Div)
```

Straight-line programming language

```
Stm
              Stm ; Stm | id := Exp | print( ExpList )
              id | num | Exp Binop Exp | (Stm , Exp)
Exp
ExpList →
              \textbf{Exp , ExpList} \mid \textbf{Exp}
              + |- |x| /
Binop →
· A program in this language:
   - What does it do?
   - How do we represent it inside the computer?
    Answer: We use trees.
                a := 5+3;
                b := (print(a, a-1), 10*a);
                print (b)
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```



Abstract Syntax Tree (AST) Java Representation - Program 1.5

```
abstract class Stm {}

class CompoundStm extends Stm {
   Stm stm1, stm2;
   CompoundStm(Stm s1, Stm s2) {stm1=s1;
   stm2=s2;}
}

class AssignStm extends Stm {
   String id; Exp exp;
   AssignStm(String i, Exp e) {id=i; exp=e;}
}

class PrintStm extends Stm {
   ExpList exps;
   PrintStm(ExpList e) {exps=e;}

20 }

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```

Java representation AST

```
abstract class Exp {}

class IdExp extends Exp {
   String id;
   IdExp(String i) {id=i;}
}

class NumExp extends Exp {
   int num;
   NumExp(int n) {num=n;}
}

class OpExp extends Exp {
   Exp left, right; int oper;
   final static int Plus=1,Minus=2,Times=3,Div=4;
   OpExp(Exp l, int o, Exp r) {left=1; oper=o;
   right=r;}
   **Tight=r; **Tight=r;
```

Java representation AST

```
class EseqExp extends Exp {
   Stm stm; Exp exp;
   EseqExp(Stm s, Exp e) {stm=s; exp=e;}
}
abstract class ExpList {}

class PairExpList extends ExpList {
   Exp head; ExpList tail;
   public PairExpList(Exp h, ExpList t) {head=h;
   tail=t;}
}
class LastExpList extends ExpList {
   Exp head;
   public LastExpList (Exp h) {head=h;}
   public LastExpList(Exp h) {head=h;}
   Public LastExpList(Exp h) {head=h;}
}
```

Java representation AST Instantiation

Programming the AST maxargs and interp

- int maxargs (Stm s) returns the maximum number of arguments of any print statement within any subexpression of a given statement in a Straightline program.
 - maxargs (prog) returns 2
 - remember that print statements can contain expressions that contain other print statements
- void interp(Stm s) "interprets" a program written in the Straightline language

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What you should do now...

- · Read, digest and understand these slides!
 - in particular work out how you would write maxargs (Stm s) - next slides

interp(Stm s), which is much harder...

```
class interp {
  static void interp(Stm s) { /* you write this part */ }
  static int maxargs(Stm s) { /* you write this part */ }
  return 0;
}

public static void main(String args[])
  throws java.io.IOException {
  Stm prog = Parser.parse(args[]); // next lectures
  System.out.println(maxargs(prog));
  interp(prog);
  }
}

Summary.2008

ROOGG Language Processors-Week!
```

Method maxargs (Stm s)

method maxargs(ExpList e)

```
static int maxargs (ExpList e) {
  if (e instanceof PairExpList)
    return Math.max
        (maxargs (((PairExpList) e).head),
            maxargs (((PairExpList) e).tail));
  else if (e instanceof LastExpList)
    return maxargs (((LastExpList) e).head);
  else {
    System.out.println
        ("maxargs(ExpList): unrecognised ExpList");
    return 0;
  }
}
```

method maxargs (Exp e)

```
static int maxargs (Exp e) {
   if (e instanceof OpExp) {
     return Math.max
        (maxargs (((OpExp) e).left),
        maxargs (((OpExp) e).right));
   } else if (e instanceof EseqExp) }
   return Math.max
        (maxargs (((EseqExp) e).stm),
        maxargs (((EseqExp) e).exp));
   } else {
     /* it's an IdExp or a NumExp */
     return 0;
}
```

method maxargs

- But this solution is not very "object oriented".
- What about adding the method maxargs() to the abstract classes Stm and Exp?
- Exercise: Try this approach.
- Exercise: Implement interp

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Java - doing without BlueJ

- Lab and Coursework: We'll develop Java programs without BlueJ
- · This is easier than it might seem;
- Our main program needs an entry point which is the method main():
 - public static void main(String args[])
- This method needs to set up and construct your initial objects, and then that is it!

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Required reading for Week 2

- Before next week's labs and lecture you should:
 - Read Appel (2002), ch. 1-2.
 - Review Java.

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Next Lecture

- Language processing & lexical analysis
- Monday 2nd February, 2009
- 11:00 12:50
 - C.350

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