

Session Plan

Session 5: Semantic analysis

- · Symbol tables
 - Environments
 - Hash tables
 - Symbol table for MiniJava
- Typechecking

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Semantic analysis

Christian Cooper

- Connects variable (and type/class, and function) definitions to their uses
- Checks that each use matches an appropriate declaration
 - must also use scope rules of language...
- Checks that each expression/part of the program is of correct type
- Translates abstract syntax into simpler representation suitable for generating machine code

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Semantic analysis

- Collect all identifiers in symbol table(s) with attributes such as: type; scope; for a function or method, parameter types & return type
 - by traversal of declaration ASTs
- Check that applied occurrences of identifiers have related declarations (are in table) at right scope
 - by traversal of rest of program ASTs

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Semantic analysis

- Check that types of (sub-)expressions and parts of statements are of expected type
 - by traversal of rest of program ASTs and collecting types
 - eg actual param matches type of formal; in A+B, A,B are ints

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Symbol tables (environments)

- Map identifiers to their types & locations
- Identifiers (variable names, function or method names, type or class names) are bound to "meanings" (bindings) in symbol tables
 - eg type of variable name
 - eg number and type of parameters of function
- eg scope
- Perform lookup in symbol table when there is a *use* (nondefining occurrence) of identifier

Scope for each local variable in which it is visible:

class X { int y;... }
local variable in a Java
method?

Discard identifier bindings local to scope at end of scope analysis

Environment is a set of bindings \rightarrow e.g.

 $\sigma_0 = \{g \rightarrow \text{string,a} \rightarrow \text{int}\}$

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```
Example
class C {
                                                          σ<sub>0</sub> is starting environment
    int a; int b; int c;
                                                        2 \sigma_1 = \sigma_0 + \{a \rightarrow int, b \rightarrow int, c \rightarrow int\}
    public void m() {
       System.out.println(a+c);
                                                        4 lookup ids a, c in G,
       int j = a+b;
String a = "hello"
                                                        5 \sigma_2 = \sigma_1 + \{j \rightarrow int\}
                                                        6 \sigma_3 = \sigma_2 + \{a \rightarrow string\}
       System.out.println(a);
                                                        7 lookup \mathbf{a} in \mathbf{\sigma}_3
        System.out.println(j);
                                                        8 lookup j in σ<sub>3</sub>
        System.out.println(b);
                                                        9 lookup b in G
                                                        10 discard \sigma_3 revert to \sigma_1
                                                        11 discard \sigma_1 revert to \sigma_0
 Need to deal with clashes (different bindings for same symbol)
                                 \sigma_2: a \rightarrow int , \sigma_3: a \rightarrow string
  Prefer most recent binding so as to implement scope rules of language
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```

Imperative implementation style

- Modify σ₁ to σ₂
- Undo modification to get back to σ₁
- · Use hash tables
- $\sigma' = \sigma + \{a \rightarrow \tau\}$ Implement by inserting τ into hash table with key a
- Simple hash table with external chaining: i th bucket = linked list of all elements whose keys hash to i mod SIZE

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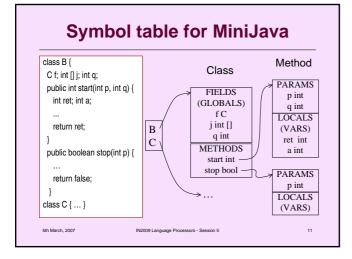
Hashing

 java.util.Hashtable implements this for us, and a hashtable will be used for storing classes, and within them two hashtables will store the global variables (fields) and methods, and within methods a hashtable will store the local variables.

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Hashing with external chaining Array of List Headers Buckets D Binding m-1 March, 2007 IND Binding



MiniJava symbols & typechecking

- Each variable name and formal parameter bound to its type
- Each method name bound to its formal parameters, result type, and local variables
- Each class name bound to its variable (field) and method declarations
- · Typechecking:
 - First, build the symbol table
 - Then typecheck the statements and expressions in the AST, consulting the symbol table for each identifier found.

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Symbol table implementation

```
class SymbolTable {
  public SymbolTable();
  public boolean addClass(String id, String parent);
  public Class getClass(String id);
  public boolean containsClass(String id);
  public Type getVarType(Method m, Class c, String id);
  public Method getMethod(String id, String classScope);
  public Type getMethodType(String id, String classScope);
  public boolean compareTypes(Type t1, Type t2);
}
```

• SymbolTable contains a hashtable of Class objects...

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Symbol table implementation

- getVarType(Method m, Class c, String id)
 - in c.m, find variable id
 - may be...
 - · local variable in method
 - parameter in formal parameters in method
 - variable in the class
 - · variable in a parent class

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Symbol table implementation

- getMethod(), getMethodType()
 - may be defined in the parent classes
- compareTypes()
 - primitive types IntegerType,
 BooleanType, IntegerArrayType
 - IdentifierTypes (class types) stored as strings, returns true if identical or if equals a parent

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Symbol table class Class

```
class Class {
  public Class(String id, String parent);
  public String getId();
  public Type type();
  public boolean addMethod(String id, Type type);
  public Method getMethod(String id);
  public boolean containsMethod(String id);
  public boolean addVar(String id, Type type);
  public Variable getVar(String id);
  public boolean containsVar(String id);
  public String parent();
}
```

• Class contains a hashtable of global variables (fields) and a hashtable of methods

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Symbol table variable representation

```
class Variable {
  public Variable(String id, Type type);
  public String id();
  public Type type();
}
```

Symbol table method representation

```
class Method {
  public Method(String id, Type type);
  public String getId();
  public Type type();
  public boolean addParam(String id, Type type);
  public Variable getParamAt(int i);
  public boolean getParam(String id);
  public boolean containsParam(String id);
  public boolean addVar(String id, Type type);
  public Variable getVar(String id);
  public boolean containsVar(String id);
}
```

 Method contains a vector of parameters (formals) and a hashtable of (local) variables, and a return type

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Implementation

- By visitors, just like the pretty printer
 - Interfaces Visitor is as before
 - A new visitor interface TypeVisitor is just the same but its methods return a Type
 - General classes DepthFirstVisitor and TypeDepthFirstVisitor implement these interfaces and traverse the AST depth-first visiting each node, but taking no action

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Implementation

- · Building the symbol table
 - BuildSymbolTableVisitor extends
 TypeDepthFirstVisitor, overriding methods
 so as to add classes, methods, vars, etc
- Typechecking
 - TypeCheckVisitor extends
 DepthFirstVisitor overriding methods so as
 to check statements, TypeCheckExpVisitor
 extends TypeDepthFirstVisitor overriding
 methods so as to check expressions

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BuildSymbolTableVisitor

 Note that some checks are done as the table is built, eg for variable declarations (see below) a check is made on whether the variable is already declared

```
public class BuildSymbolTableVisitor extends
   TypeDepthFirstVisitor {
...
   private Class currClass;
   private Method currMethod;
...
   // Type t;
   // Identifier i;
   public Type visit(VarDecl n) {
    Type t = n.t.accept(this);
    String id = n.i.toString();
}
```

BuildSymbolTableVisitor (2)

TypeCheckVisitor

```
public class TypeCheckVisitor extends
   DepthFirstVisitor {
    static Class currClass;
    static Method currMethod;
    static SymbolTable symbolTable;

   public TypeCheckVisitor(SymbolTable s){
        symbolTable = s;
    }
...
```

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TypeCheckVisitor (2)


```
TypeCheckExpVisitor

public class TypeCheckExpVisitor extends TypeDepthFirstVisitor {

"// Exp el.e2;
public Type visit(Plus n) {

if (! (n.el.accept(this) instanceof IntegerType) ) {

System.out.println("Left side of Plus must be of type integer");
}

if (! (n.e2.accept(this) instanceof IntegerType) ) {

System.out.println("Right side of Plus must be of type integer");
}

return new IntegerType();
}

"**

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```

Method Calls $e.i(el_1,el_2,...)$

- Lookup method in the symbolTable to get parameter list and result type
- Find i in class e
- The parameter types in the parameter list for the method must be matched against the actual arguments el₁, el₂,
- Result type becomes the type of the method call as a whole.

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TypeCheckExpVisitor (2)

```
|
|/ Exp e;
|/ Identifier i;
|/ ExpList el;
|/ public Type visit(Call n) {
| if (! (n.e.accept(this) instanceof IdentifierType)) {
| system.out.println("method "+ n.i.toString()+ "called on something that is not a class or Object.");
| String mname = n.i.toString();
| String cname = ((IdentifierType) n.e.accept(this)).s;
| Method calledMethod = TypeCheckVisitor.symbolTable.getMethod(mname,cname);
| for (int i = 0; i < n.el.size(); i++) {
| Type tl = null;
| Type tl = null;
| if (calledMethod.getParamAt(i)!=null)
| tl = calledMethod.getParamAt(i).type();
| t2 = n.el.elementAt(i).accept(this);
| if (!TypeCheckVisitor.symbolTable.compareTypes(tl,t2)) {
| System.out.println("Type Error in arguments passed to " +cname+"."
| +mname);
| }
| return TypeCheckVisitor.symbolTable.getMethodType(mname,cname);
| }
| String name = n.i.toString();
| System.out.println("Type Error in arguments passed to " +cname+"."
| +mname);
| }
| return TypeCheckVisitor.symbolTable.getMethodType(mname,cname);
| }
| String name = n.i.toString();
| System.out.println("Type Error in arguments passed to " +cname+"."
| +mname);
| }
| return TypeCheckVisitor.symbolTable.getMethodType(mname,cname);
| }
| String name = n.i.toString();
| System.out.println("Type Error in arguments passed to " +cname+"."
| +mname);
| }
| Return TypeCheckVisitor.symbolTable.getMethodType(mname,cname);
| }
| String name = n.i.toString();
```

What you should do now...

- Read and digest chapter 5
 - you don't need functional implementation styles or mutiple tables
- Get ready further to develop the MiniJava typechecker... assessment 2 awaits!

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Schedule

- Activation records (stack frames)
- Monday 12th March, 2007
 - 12:00 13:50
 - CM383

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