

Semantic Analysis

Christian Cooper

Session Plan

Session 5: Semantic analysis

- Symbol tables
 - Environments
 - Hash tables
 - Symbol table for MiniJava
- Typechecking

5th March, 2007

IN2009 Language Processors - Session 5

2

Semantic analysis

- Connects variable (and type/class, and function) definitions to their uses
- Checks that each use matches an appropriate declaration
 - must also use scope rules of language...
- Checks that each expression/part of the program is of correct type
- Translates abstract syntax into simpler representation suitable for generating machine code

5th March, 2007

IN2009 Language Processors - Session 5

3

Semantic analysis

- Collect all identifiers in symbol table(s) with attributes such as: type; scope; for a function or method, parameter types & return type
 - by traversal of declaration ASTs
- Check that applied occurrences of identifiers have related declarations (are in table) at right scope
 - by traversal of rest of program ASTs

5th March, 2007

IN2009 Language Processors - Session 5

4

Semantic analysis

- Check that types of (sub-)expressions and parts of statements are of expected type
 - by traversal of rest of program ASTs and collecting types
 - eg actual param matches type of formal; in $A+B$, A, B are ints

5th March, 2007

IN2009 Language Processors - Session 5

5

Symbol tables (environments)

- Map identifiers to their types & locations
- Identifiers (variable names, function or method names, type or class names) are bound to "meanings" (bindings) in symbol tables
 - eg type of variable name
 - eg number and type of parameters of function
 - eg scope
- Perform lookup in symbol table when there is a use (non-defining occurrence) of identifier

Scope for each local variable in which it is visible:

```
class X { int y;... }
  local variable in a Java
  method?
```

Discard identifier bindings local to scope at end of scope analysis

Environment is a set of bindings → e.g.

$\sigma_0 = \{g \rightarrow \text{string}, a \rightarrow \text{int}\}$

5th March, 2007

IN2009 Language Processors - Session 5

6

Example

```

class C {
  int a; int b; int c;
  public void m() {
    System.out.println(a+c);
    int j = a+b;
    String a = "hello"
    System.out.println(a);
    System.out.println(j);
    System.out.println(b);
  }
}

```

- 1 σ_0 is starting environment
- 2 $\sigma_1 = \sigma_0 + \{a \rightarrow \text{int}, b \rightarrow \text{int}, c \rightarrow \text{int}\}$
- 3
- 4 lookup ids a, c in σ_1
- 5 $\sigma_2 = \sigma_1 + \{j \rightarrow \text{int}\}$
- 6 $\sigma_3 = \sigma_2 + \{a \rightarrow \text{string}\}$
- 7 lookup a in σ_3
- 8 lookup j in σ_3
- 9 lookup b in σ_1
- 10 discard σ_3 revert to σ_1
- 11 discard σ_1 revert to σ_0

Need to deal with clashes (different bindings for same symbol)

$\sigma_2: a \rightarrow \text{int}$, $\sigma_3: a \rightarrow \text{string}$

Prefer *most recent binding* so as to implement scope rules of language

5th March, 2007

IN2009 Language Processors - Session 5

7

Imperative implementation style

- Modify σ_1 to σ_2
- Undo modification to get back to σ_1
- Use *hash tables*
- $\sigma' = \sigma + \{a \rightarrow \tau\}$ Implement by inserting τ into hash table with key a
- Simple hash table with *external chaining*: i th bucket = linked list of all elements whose keys hash to $i \bmod \text{SIZE}$

5th March, 2007

IN2009 Language Processors - Session 5

8

Hashing

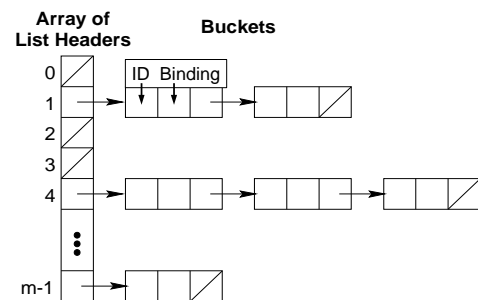
- `java.util.Hashtable` implements this for us, and a hashtable will be used for storing classes, and within them two hashtables will store the global variables (fields) and methods, and within methods a hashtable will store the local variables.

5th March, 2007

IN2009 Language Processors - Session 5

9

Hashing with external chaining

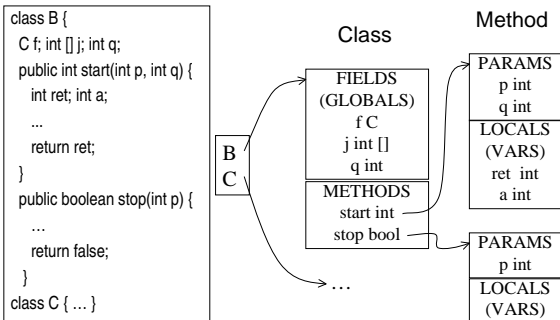


5th March, 2007

IN2009 Language Processors - Session 5

10

Symbol table for MiniJava



5th March, 2007

IN2009 Language Processors - Session 5

11

MiniJava symbols & typechecking

- Each variable name and formal parameter bound to its type
- Each method name bound to its formal parameters, result type, and local variables
- Each class name bound to its variable (field) and method declarations
- Typechecking:
 - First, build the symbol table
 - Then typecheck the statements and expressions in the AST, consulting the symbol table for each identifier found.

5th March, 2007

IN2009 Language Processors - Session 5

12

Symbol table implementation

```
class SymbolTable {
    public SymbolTable();
    public boolean addClass(String id, String parent);
    public Class getClass(String id);
    public boolean containsClass(String id);
    public Type getVarType(Method m, Class c, String id);
    public Method getMethod(String id, String classScope);
    public Type getMethodType(String id, String classScope);
    public boolean compareTypes(Type t1, Type t2);
}
```

- SymbolTable contains a hashtable of Class objects...

5th March, 2007

IN2009 Language Processors - Session 5

13

Symbol table implementation

- **getVarType(Method m, Class c, String id)**
 - in **c.m**, find variable **id**
 - may be...
 - local variable in method
 - parameter in formal parameters in method
 - variable in the class
 - variable in a parent class

5th March, 2007

IN2009 Language Processors - Session 5

14

Symbol table implementation

- **getMethod(), getMethodType()**
 - may be defined in the parent classes
- **compareTypes()**
 - primitive types **IntegerType**, **BooleanType**, **IntegerArrayType**
 - **IdentifierTypes** (class types) stored as strings, returns true if identical or if equals a parent

5th March, 2007

IN2009 Language Processors - Session 5

15

Symbol table class Class

```
class Class {
    public Class(String id, String parent);
    public String getId();
    public Type type();
    public boolean addMethod(String id, Type type);
    public Method getMethod(String id);
    public boolean containsMethod(String id);
    public boolean addVar(String id, Type type);
    public Variable getVar(String id);
    public boolean containsVar(String id);
    public String parent();
}
```

- Class contains a hashtable of global variables (fields) and a hashtable of methods

5th March, 2007

IN2009 Language Processors - Session 5

16

Symbol table variable representation

```
class Variable {
    public Variable(String id, Type type);
    public String id();
    public Type type();
}
```

5th March, 2007

IN2009 Language Processors - Session 5

17

Symbol table method representation

```
class Method {
    public Method(String id, Type type);
    public String getId();
    public Type type();
    public boolean addParam(String id, Type type);
    public Variable getParamAt(int i);
    public boolean getParam(String id);
    public boolean containsParam(String id);
    public boolean addVar(String id, Type type);
    public Variable getVar(String id);
    public boolean containsVar(String id);
}
```

- Method contains a vector of parameters (formals) and a hashtable of (local) variables, and a return type

5th March, 2007

IN2009 Language Processors - Session 5

18

Implementation

- By visitors, just like the pretty printer
 - Interfaces `Visitor` is as before
 - A new visitor interface `TypeVisitor` is just the same but its methods return a `Type`
 - General classes `DepthFirstVisitor` and `TypeDepthFirstVisitor` implement these interfaces and traverse the AST depth-first visiting each node, but taking no action

5th March, 2007

IN2009 Language Processors - Session 5

19

Implementation

- Building the symbol table
 - `BuildSymbolTableVisitor` extends `TypeDepthFirstVisitor`, overriding methods so as to add classes, methods, vars, etc
- Typechecking
 - `TypeCheckVisitor` extends `DepthFirstVisitor` overriding methods so as to check statements, `TypeCheckExpVisitor` extends `TypeDepthFirstVisitor` overriding methods so as to check expressions

5th March, 2007

IN2009 Language Processors - Session 5

20

BuildSymbolTableVisitor

- Note that some checks are done as the table is built, eg for variable declarations (see below) a check is made on whether the variable is already declared

```
public class BuildSymbolTableVisitor extends
    TypeDepthFirstVisitor {
    ...
    private Class currClass;
    private Method currMethod;
    ...
    // Type t;
    // Identifier i;
    public Type visit(VarDecl n) {
        Type t = n.t.accept(this);
        String id = n.i.toString();
```

5th March, 2007

IN2009 Language Processors - Session 5

21

BuildSymbolTableVisitor (2)

```
if (currMethod == null){
    if (!currClass.addVar(id,t)){
        System.out.println(id + "is already defined in "
            + currClass.getId());
    }
} else {
    if (!currMethod.addVar(id,t)){
        System.out.println(id + "is already defined in "
            + currClass.getId() + "." +
            currMethod.getId());
    }
}
return null;
}
```

5th March, 2007

IN2009 Language Processors - Session 5

22

TypeCheckVisitor

```
public class TypeCheckVisitor extends
    DepthFirstVisitor {
    static Class currClass;
    static Method currMethod;
    static SymbolTable symbolTable;

    public TypeCheckVisitor(SymbolTable s){
        symbolTable = s;
    }
    ...
}
```

5th March, 2007

IN2009 Language Processors - Session 5

23

TypeCheckVisitor (2)

```
...
// Identifier i;
// Exp e;
public void visit(Assign n) {
    Type t1 =
        symbolTable.getVarType(currMethod,currClass,n.i.toString());
    Type t2 = n.e.accept(new TypeCheckExpVisitor() );

    if (symbolTable.compareTypes(t1,t2)==false){
        System.out.println("Type error in assignment to
            "+n.i.toString());
    }
}
...
```

5th March, 2007

IN2009 Language Processors - Session 5

24

TypeCheckVisitor (3)

```
// Type t;
// Identifier i;
// FormalList fl;
// VarDeclList vl;
// StatementList sl;
// Exp e;
public void visit(MethodDecl n) {
    n.t.accept(this);
    String id = n.i.toString();
    currMethod = currClass.getMethod(id);
    Type retType = currMethod.type();
    for (int i = 0; i < n.fl.size(); i++)
        {n.fl.elementAt(i).accept(this); }
    for (int i = 0; i < n.vl.size(); i++)
        {n.vl.elementAt(i).accept(this); }
    for (int i = 0; i < n.sl.size(); i++)
        {n.sl.elementAt(i).accept(this); }
    if (symbolTable.compareTypes(retType, n.e.accept(new
        TypeCheckExpVisitor()))==false) {
        System.out.println("Wrong return type for method "+ id);
    }
}
```

5th March, 2007

IN2009 Language Processors - Session 5

25

TypeCheckExpVisitor

```
public class TypeCheckExpVisitor extends TypeDepthFirstVisitor {
    ...
    // Exp e1,e2;
    public Type visit(Plus n) {
        if (! (n.e1.accept(this) instanceof IntegerType) ) {
            System.out.println("Left side of Plus must be of type integer");
        }

        if (! (n.e2.accept(this) instanceof IntegerType) ) {
            System.out.println("Right side of Plus must be of type integer");
        }

        return new IntegerType();
    }
    ...
}
```

5th March, 2007

IN2009 Language Processors - Session 5

26

Method Calls $e.i(e1_1, e1_2, \dots)$

- Lookup method in the `symbolTable` to get parameter list and result type
- Find `i` in class `e`
- The parameter types in the parameter list for the method must be matched against the actual arguments $e1_1, e1_2, \dots$
- Result type becomes the type of the method call as a whole.

5th March, 2007

IN2009 Language Processors - Session 5

27

TypeCheckExpVisitor (2)

```
// Exp e;
// Identifier i;
// ExpList el;
public Type visit(Call n) {
    if (! (n.e.accept(this) instanceof IdentifierType)){
        System.out.println("method "+ n.i.toString()+" called on something
        that is not a class or Object.");
    }

    String mname = n.i.toString();
    String cname = ((IdentifierType) n.e.accept(this)).s;
    Method calledMethod = TypeCheckVisitor.symbolTable.getMethod(mname, cname);

    for (int i = 0; i < n.el.size(); i++) {
        Type t1 = null;
        Type t2 = null;

        if (calledMethod.getParamAt(i)!=null)
            t1 = calledMethod.getParamAt(i).type();
        t2 = n.el.elementAt(i).accept(this);
        if (!TypeCheckVisitor.symbolTable.compareTypes(t1,t2)){
            System.out.println("Type Error in arguments passed to " +cname+"."
            +mname);
        }
    }
    return TypeCheckVisitor.symbolTable.getMethodType(mname, cname);
}
```

5th March, 2007

IN2009 Language Processors - Session 5

28

What you should do now...

- Read and digest chapter 5
 - you don't need functional implementation styles or mutiple tables
- Get ready further to develop the MiniJava typechecker... assessment 2 awaits!

5th March, 2007

IN2009 Language Processors - Session 5

29

Schedule

- **Activation records (stack frames)**
- Monday 12th March, 2007
 - 12:00 - 13:50
 - CM383

5th March, 2007

IN2009 Language Processors - Session 5

31