CS/MATH111 ASSIGNMENT 2

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Problem 1: Let $n = p_1 p_2 ... p_k$, where $p_1, p_2, ..., p_k$ are different primes. Prove that n has exactly 2^k different divisors. For example, if n = 105, then $n = 3 \cdot 5 \cdot 7$, so k = 3, and thus n has $2^3 = 8$ divisors. These divisors are: 1, 3, 5, 7, 15, 21, 35, 105. Hint. You can reduce the problem to counting other objects that we already know how to count. Alternatively, this can be proved by induction on k.

Solution 1: We will be using Induction to prove this problem

Base case:

Assume that k = 1. As a result, we should get 2 divisors if we have one prime p_1 .

Assume that n is any prime number. As a result we know that the two divisors are 1 and n itself Therefore, the base case holds

Inductive Step:

We want to prove that for $n = p_1 p_2 ... p_k$, that there are different prime numbers that there are 2^k different divisors. Assumption: We know that for $n = p_1 p_2 ... p_{k-1}$ that there k-1 different primes and there are 2^{k-1} different divisors.

$$n_1 = p_1 p_2 ... p_{k-1}$$
 Following our assumption with a certain number n_1 $n_2 = n_1 \cdot p_k$ Multiply every prime number in n_1 by p_k to get a new divisor for each old divisor $d_1 \cdot p^k = d_2$ d_1 is the set of all divisors from $n_1.d_2$ is the new set of divisors d_1 multiplied by p_k $... d_{1_{k-1}} \cdot p^k = ... d_{2_{k-1}}$

We need to add all the old and new divisors together. We know for both n_1 and n_2 , that there are 2^{k-1} divisors

$$2^k=d_1+d_2$$
 Following from before
$$2^k=2^{k-1}+2^{k-1}$$
 Counting current primes and new primes
$$2^k=\frac{1}{2}\cdot 2^k+\frac{1}{2}\cdot 2^{k-1}$$

Therefore, by induction we know that for $n = p_1 p_2 ... p_k$ that there are 2^k divisors. We accounted for all the divisors that either contained p_k or did not contain p_k . d_1 contained all the divisors that did not contain p_k and d_2 accounted for divisors that contained p_k

Problem 2: Alice's RSA public key is P = (e, n) = (13, 77). Bob sends Alice the message by encoding it as follows. First he assigns numbers to characters: A is 2, B is 3, ..., Z is 27, and blank is 28. Then he uses RSA to encode each number separately.

Bob's encoded message is:

10	7	58	30	23	62
7	64	62	23	62	61
7	41	62	21	7	49
75	7	69	53	58	37
37	41	10	64	50	7
10	64	21	62	61	35
62	61	62	7	52	10
21	58	7	49	75	7
62	26	22	53	30	21

10 37 64

Decode Bob's message. Notice that you don't have Bob's secrete key, so you need to "break" RSA to decrypt his message.

For the solution, you need to provide the following:

- Describe step by step how you arrived at the solution. In particular, explain how you determined $p, q, \phi(n)$, and d.
- Show the calculation that determines the first letter in the message from the first number in ciphertext.
- Give Bob's message in plaintext. The message is a quote. Who said it?
- If you wrote a program, attach your code to the hard copy. If you solved it by hand (not recommended), attach your scratch paper with calculations for at least 5 first letters.

Suggestion: this can be solved by hand, but it will probably be faster to write a short program.

Solution 2:

• Describe steps of solution:

We find p and q from n, we are given n = 77. The only prime factors of 77 are 7 and 11. So, p = 7 and q = 11.

Next, we find totient function $\phi(n)$. $\phi(77) = (p-1)(q-1) = 6*10 = 60$

Next, we solve for d using formula: $d = e^{-1} mod(\phi(n))$ Plugging in e = 13 and $\phi(n) = 60$: $d = 13^{-1} mod(60)$ Since 60 is not prime, we cannot use Fermat's Little Theorem

But we can use Euclid's Algorithm: gcd(60, 13) $60 = 13(4) + 8 \rightarrow 8 = 60 - 13(4)$

$$13 = 8(1) + 5 \to 5 = 13 - 8$$

$$8 = 5(1) + 3 \rightarrow 3 = 8 - 5$$

$$5 = 3(1) + 2 \rightarrow 2 = 5 - 3$$

$$3 = 2(1) + 1 \rightarrow 1 = 3 - 2$$

$$\begin{split} 1 &= 3 - [5 - 3] \rightarrow 1 = 2(3) - 5 \\ 1 &= 2[8 - 5] - 5 \rightarrow 1 = 2(8) - 3(5) \\ 1 &= 2[8 - 5] - 5 \rightarrow 1 = 5(8) - 3(13) \\ 1 &= 5[60 - 4(13)] - 3(13) \rightarrow 1 = 5(60) - 20(13) - 3(13) \\ 1 &= 5(60) - 23(13) \\ d &= -23 mod(60) \rightarrow d = 37 \end{split}$$

Next, we decrypt message using formula: $M = C^d mod(n)$ Our formula is: $M = C^{37} mod(77)$

• We can decrypt the letter for the first number in the ciphertext using C = 10. $M = C^{37} mod(77) \rightarrow M = 10^{37} mod(77)$

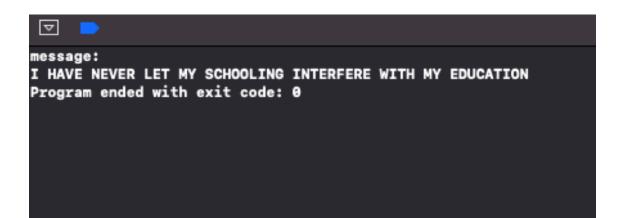
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\begin{array}{l} M=(10^2)^{18}*10 mod (77) \\ 10^2 mod (77)=23 \\ 10^4 mod (77)=(10^2)^2 mod (77)=(23)^2 mod (77)=67 \\ 10^8 mod (77)=(10^4)^2 mod (77)=(67)^2 mod (77)=23 \\ 10^{16} mod (77)=(10^8)^2 mod (77)=(23)^2 mod (77)=67 \\ 10^{32} mod (77)=(10^{16})^2 mod (77)=(67)^2 mod (77)=23 \\ 10^{37} mod (77)=(10^{32}*10^4*10) mod (77)=(23*67*10) mod (77)=10 \\ M=10 \end{array}
```

• Give Bob's message in plaintext. The message is a quote. Who said it?
"I HAVE NEVER LET MY SCHOOLING INTERFERE WITH MY EDUCATION" by Mark Twain

Code for RSA Decryption:

```
RSA > My Mac
         RSA > RSA > RSA > a main.cpp > f decrypt(vector<int> message, int n, int d)
 7 #include <sstream>
   #include <vector>
    #include <cmath>
11 using namespace std;
13 //find prime numbers p1 & p2 given n
14 int breakRSA();
15 void decrypt(vector<int> message, int n, int d);
   int main(int argc, const char * argv[]) {
             and message (each letter: C)
        //2. Find phi_n = (p - 1)(q - 1).
        // d = (1/e)%phi_n
// d = (1/13)%60 = 37
// Euclid's Algorithm: d = gcd(e, n)
        int e = 13, n = 77;
        int phi_n = 60;
        int d = 37;
```

```
vector<int> message =
    { 10, 7, 58, 30, 23, 62,
       7, 64, 62, 23, 62, 61,
       7, 41, 62, 21, 7, 49,
      75, 7, 69, 53, 58, 37,
      37, 41, 10, 64, 50, 7,
      10, 64, 21, 62, 61, 35,
      62, 61, 62, 7, 52, 10,
      21, 58, 7, 49, 75, 7,
      62, 26, 22, 53, 30, 21,
      10, 37, 64};
    decrypt(message, n, d);
    return 0;
}
void decrypt(vector<int> numText, int n, int d)
{
    char charText;
    long x;
    int m, C = 0;
    cout << "message: " << endl;
    for(int i = 0; i < numText.size(); i++)</pre>
    {
        C = numText.at(i);
        //cout << C << endl;
        string str = to_string(m);
        m = static_cast<int>(pow(C,2)) % n;
        m = static_cast<int>(pow(m,3)) % n;
        m = (m*C) \% n;
        if (m == 28) { cout << ' '; }
        else
        €
            m += 63;
            charText = m;
            cout << charText;</pre>
    cout << endl;
```



Problem 3:

(a) Compute $13^{-1} \pmod{19}$ by enumerating multiples of the number and the modulus. Show your work. Referring to the slides in number theory. We want to find an integer a and b that satisfies $13^{-1} \pmod{19}$

$$a=13^{-1}\pmod{19}$$
 Multiply both sides by 13
 $13\cdot a=1\pmod{19}$ turn into multiple form $13\cdot a=19\cdot b+1$

On the LHS, we have multiples 13,26,29. On the RHS, we have multiples 20,39,58. Since $39 = 19 \cdot 2 + 1$, we have a = 3 and b = 2. So as a result we have $13^{-1} = 3 \pmod{19}$. Finally, substituting what we just found $13^{-1} \pmod{19} = 3 \pmod{19}$

(b) Compute $13^{-1} \pmod{19}$ using Fermat's theorem. Show your work. Referring to the slides in number theory. We want to find use Fermat's little Theorem

$$a^{p-1}=1\pmod{p}$$
 If p is a prime number, and a is not divisible by p(FLT) $13^{19-1}=1\pmod{19}$ Multiplying, we get this $13^{18}=1\pmod{19}$

Next we apply FLT to our equation. We will use what we just found above

$$13^{18} \cdot 13^{-1} = 1 \pmod{19}$$
 If p is a prime number, and a is not divisible by p(FLT) $13^{17} = 1 \pmod{19}$ Multiplying, we get this

Listing out exponential and their remainders. This will help us get the final answer

$$13^2 \pmod{19} = 17$$
 $13^4 \pmod{19} = 4$
 $13^8 \pmod{19} = 16$
 $13^{16} \pmod{19} = 9$

Applying what we just found

$$13^{17}\pmod{19}=13^{16}\cdot 13^1\pmod{19}$$
 We know that $13^{16}\pmod{19}=9$
$$13^{16}\cdot 13^1\pmod{19}=9\pmod{19}\cdot 13\pmod{19}$$
 9 $\pmod{19}\cdot 13\pmod{19}=142\pmod{19}$ 117 $\pmod{19}=3$

This matches our final answer in (a). Therefore $13^{-1} \pmod{19} = 3 \pmod{19}$

(c) Compute $13^{-40} \pmod{19}$ using Fermat's theorem. Show your work. For this equation

$$a^{p-1}=1\pmod{p}$$
 If p is a prime number, and a is not divisible by p(FLT)
$$13^{19-1}=1\pmod{19}$$
 Multiplying, we get this
$$13^{18}=1\pmod{19}$$

Next we apply FLT to our equation. We will use what we just found above

$$13^{18} * 13^{-40} \pmod{19} \equiv 13^{-22} \pmod{19}$$

 $13^{18} * 13^{-22} \pmod{19} \equiv 13^{-4} \pmod{19}$
 $13^{18} * 13^{-4} \pmod{19} \equiv 13^{14} \pmod{19}$

Listing out exponential and their remainders, will help us find answer for $13^{14} \pmod{19}$

$$13^2 \pmod{19} = 17$$

 $13^4 \pmod{19} = 4$
 $13^8 \pmod{19} = 16$

$$13^{14} \pmod{19} = (13^8 * 13^4 * 13^2) \pmod{19}$$

 $13^{14} = (16 * 4 * 17) \pmod{19}$
 $= (1088) \pmod{19}$
 $= 5$ Final answer

(d) Find a number $x \in \{1, 2, ..., 36\}$ such that $8x \equiv 3 \pmod{37}$. Show your work. (You need to follow the method covered in class; brute-force checking all values of x will not be accepted.)

$$8x\equiv 3\pmod{37}$$
 Multiply both sides by the inverse of 8
$$8^{-1}\cdot 8x\equiv 3\pmod{37}\cdot 8^{-1}$$

$$x\equiv 3\pmod{37}\cdot 8^{-1}$$

Solving for $8^{-1} \pmod{37}$

$$a = 8^{-1} \pmod{37}$$
 Multiply both sides by 8
8 · $a = 37 \cdot b + 1$ turn into multiple form

Listing out multiples on the RHS: 37,75,112. We know that 112 is divisible by 8. Therefore we have a=14 and b=3. At the end we know that $8^{-1} \pmod{37} = 14 \pmod{37}$ Plugging it in and solving, we get

$$x = 3 \cdot 14 \pmod{37}$$
 Multiply $x = 42 \pmod{37}$ Get remainder $x = 5$ Final answer

Submission. To submit the homework, you need to upload the pdf file into gradescope by Friday, May 4 (noon).