# **2nd Lab** - CES 27 - Distributed Programming September 4th, 2019



## Prof. Celso Massaki HIRATA, hirata@ita.br Prof. JULIANA de Melo Bezerra, juliana@ita.br

### Igor Bragaia¹, igor.bragaia@gmail.com

<sup>1</sup>Computer Engineering undergraduate student at Aeronautics Institute of Technology (ITA)

This project aims to implement mutual exclusion through Ricart-Agrawala algorithm using scalar logical clocks from past project. It is the second assignment of Distributed Programming 2019.2 ITA's class.

There are two main codes, one for a process resource management and other to receive messages showing how the resource has been used on this distributed system.

```
type Message struct {
 Time int
 Text string
var myPortId int
var ResourceConn *net.UDPConn
var logicalClockFreeze int
var queue []int
func PrintError(err error) {
```

```
func doServerJob() {
 var receivedMessage Message
 err = json.Unmarshal(buf[:n], &receivedMessage)
 if logicalClock < receivedMessage.Time {</pre>
     logicalClock = receivedMessage.Time
 if receivedMessage.Text == "REPLY" {
    fmt.Printf("[logical clock %d] REPLY from %d\n", logicalClock, receivedMessage.Processor)
    if counter == qty {
   state = "HELD"
        fmt.Printf("[logical clock %d] Replying to: ", logicalClock)
 } else if receivedMessage.Text == "REQUEST" {
   if state == "HELD" || (state == "WANTED" && logicalClockFreeze < receivedMessage.Time) {</pre>
        queue = append(queue, receivedMessage.Processor)
        fmt.Printf("[logical clock %d] REPLYING TO %d\n", logicalClock, receivedMessage.Processor)
        doClientJob(receivedMessage.Processor-1, logicalClock, "REPLY")
func useResource() {
func doClientJob(otherProcess int, logicalClock int, text string) {
```

```
msg := Message{
   Processor: myPortId,
   Text: text,
id, err := strconv.Atoi(os.Args[1])
myPortId = id
myPort = os.Args[myPortId+1]
nServers = len(os.Args) - 2
qty = len(os.Args) - 3
ServConn = Conn
   ServerAddr,err := net.ResolveUDPAddr("udp","127.0.0.1" + os.Args[i+2])
   Conn, err := net.DialUDP("udp", LocalAddr, ServerAddr)
CheckError(err)
ResourceConn = Conn
reader := bufio.NewReader(os.Stdin)
   text, _, _ := reader.ReadLine()
ch <- string(text)</pre>
```

Code 1: process.go file

```
type Message struct {
var ServConn *net.UDPConn
 n, _, err := ServConn.ReadFromUDP(buf)
 err = json.Unmarshal(buf[:n], &receivedMessage)
 fmt.Printf("[logical clock %d] Process %d says %s\n", receivedMessage.Time, receivedMessage.Processor,
receivedMessage.Text)
```

```
Address, err := net.ResolveUDPAddr("udp", ":10001")
CheckError(err)
Connection, err := net.ListenUDP("udp", Address)
ServConn = Connection
CheckError(err)
defer ServConn.Close()
for {
    doServerJob()
}
```

Code 2: sharedResource.go file

In order to validate the algorithm developed, there are two test cases

#### **TEST CASE 1**

We aim to simulate three different processes, then we ran three different processes in three different terminals as suggested by the professor. Initially, process 1 tries to get into the Critical Section (CS), then send request message to both process 1 and 2. When process 1 gets all replies, it goes into CS and stay there for 5 seconds. Then process 2 tries to get into the CS, send requests as well and goes in when gets all responses, as seen on the following image. Note that on this test case, Ricart-Agrawala process queues are always empty and for this reason it have not been shown at the simulation.

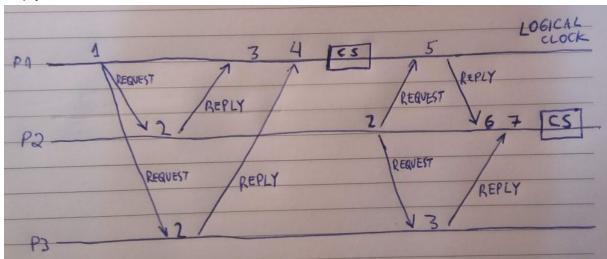


Image 1: expected output from simulation for processes 1, 2 and 3

Results are correct, as seen on the following screenshots that agree with simulation above.

#### STEP 1

Process 1 requests the shared resource and goes in and out of the CS

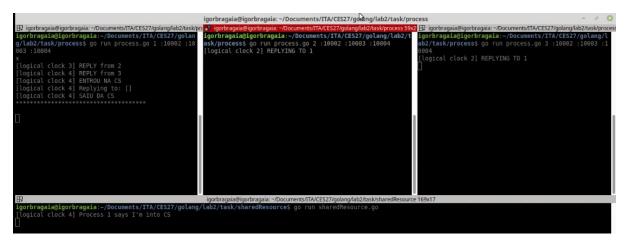


Image 2: screenshot from step 1

#### STEP 2

Process 2 requests the shared resource and goes in and out of the CS

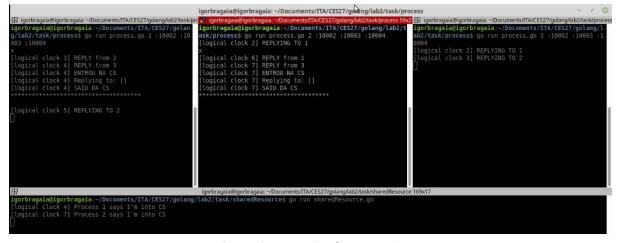


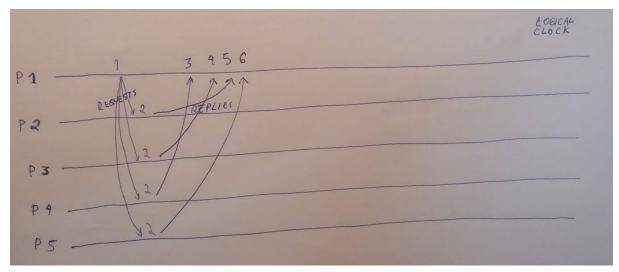
Image 3: screenshot from step 1

#### **TEST CASE 2**

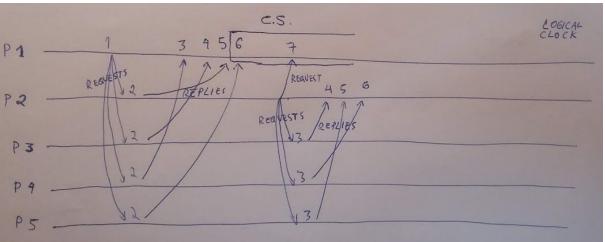
We aim to simulate five different processes, then we ran five different processes in five different terminals as suggested by the professor.

We have 6 steps:

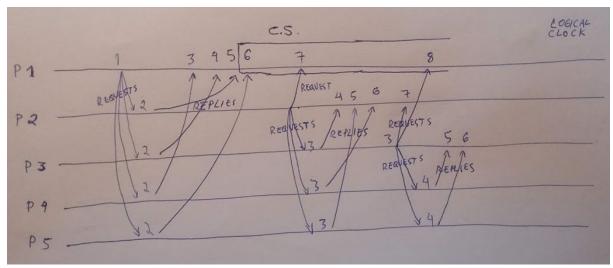
 Process 1 requests access to CS to all other processes, gets all replies because all of them are idle and then goes to CS. See simulation at the following image.



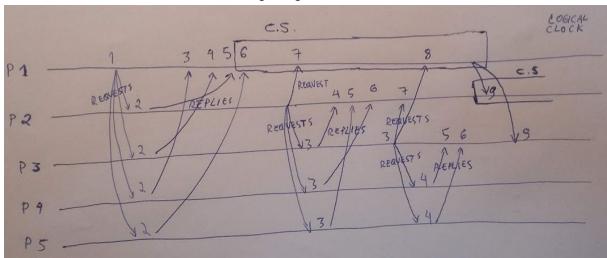
2. Before process 1 exit the CS, process 2 requests access to CS to all other processes, but does not get reply from process 1, which adds process 2 to its reply queue regarding its state HELD. See simulation at the following image.



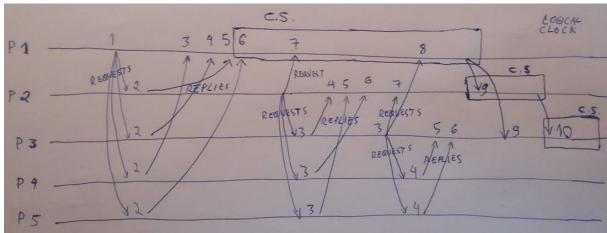
3. Before process 1 exit the CS, process 3 also requests access to CS to all other processes but does not get reply from process 1, which adds process 2 to its reply queue regarding its state HELD. Regarding state WANTED from process 3 and that logical clock for process 3 is larger than process 2, process 2 also adds process 3 to its reply queue. See simulation at the following image.



4. Process 1 exits the CS, then reply processes 2 and 3 that are at its queue. At this moment, process 2 got all necessary replies and then goes to CS. See simulation at the following image.



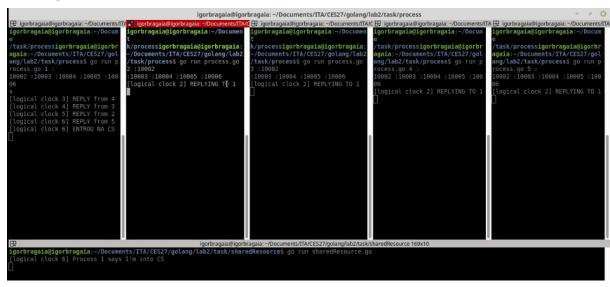
5. Process 2 exits the CS, then reply process 3 that is at its queue. At this moment, process 3 got all necessary replies and then goes to CS. See simulation at the following image.



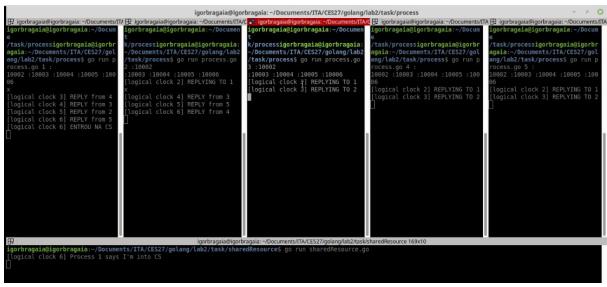
6. Process 3 exits the CS and does not reply any other process due to its empty reply queue.

Note that on this test case, Ricart-Agrawala process queues are not always empty. Finally, results are correct, as seen on the following screenshots that agree with simulation above.

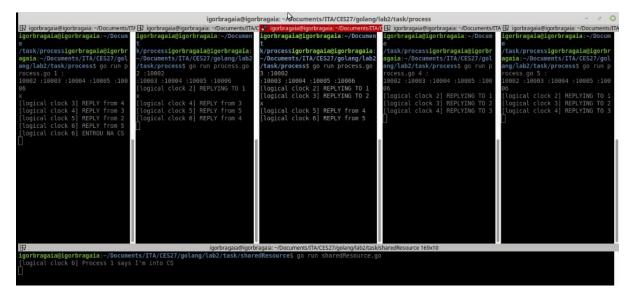
#### STEP 1



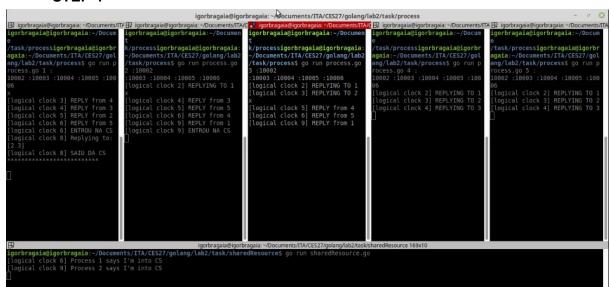
STEP 2



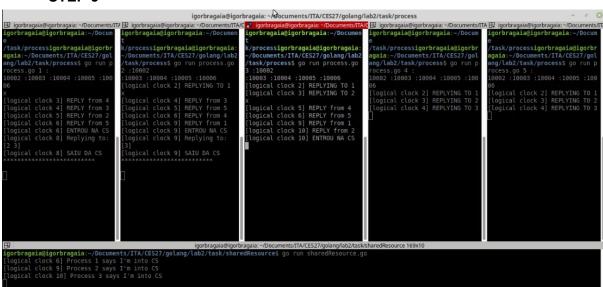
STEP 3



#### STEP 4



#### STEP 5



STEP 6

