# Automated Verification of Cyber-Physical Systems A.A. 2024/2025

# Simulation of Systems

Igor Melatti

Università degli Studi dell'Aquila

Dipartimento di Ingegneria e Scienze dell'Informazione e Matematica





#### Simulation vs Model Checking

- In "standard" Model Checking, we are given
  - a non-deterministic Kripke Structure (KS)
  - an LTL or CTL property to be verified
- The output is either PASS or FAIL
  - if PASS, then all evolutions (paths) of the given model fulfill the given property
  - if FAIL, we also have a counterexample
- In probabilistic model checking, we consider probabilities of sets of evolutions
- Simulation of a system only considers one path





#### Murphi Simulation

```
\mathbf{void} Make_a_run(NFSS \mathcal{N}, invariant \varphi)
 let \mathcal{N} = \langle S, I, \text{Post} \rangle;
 s_curr = pick_a_state(I);
 if (!\varphi(s\_curr))
  return with error message;
 while (1) { /* loop forever */
  if (Post(s_curr) = \emptyset)
    return with deadlock message;
  s_next = pick_a_state(Post(s_curr));
  if (!\varphi(s_next))
    return with error message;
  s_curr = s_next;
```



#### SPIN Simulation

```
void Make_a_run(NFSS \mathcal{N})
 let \mathcal{N} = \langle S, \{s_0\}, \text{Post} \rangle;
 s_curr = s_0;
 if (some assertion fail in s_curr))
  return with error message;
 while (1) { /* loop forever */
  if (Post(s_curr) = \emptyset)
   return with deadlock message;
  s_next = pick_a_state(Post(s_curr));
  if (some assertion fail in s_curr))
   return with error message;
  s_curr = s_next;
```



# Repeating a Simulation

- Simulations may be deterministic or probabilistic
  - both Murphi and SPIN simulations are probabilistic
  - at each step, a transition is chosen among the *n* possible ones with probability  $\frac{1}{n}$
  - of course, *n* may be different at each step
- Running multiple probabilistic simulations typically implies obtaining different paths
  - the longest the path, the more likely this is to happen





#### Repeating a Simulation

- For deterministic simulations, all runs are the same
  - choice between successor states is fixed (e.g., always the first)
  - multiple simulations all result in the same path
- Deterministic simulation are however important when *inputs* from the environment are present
  - this is actually true for many systems
  - some inputs must be given an system startup, others must be given during the system evolution
  - for a given input tuple, there is only one possible successor
- Running multiple deterministic simulation results in different paths if we vary the inputs to be received
  - this is actually true for many systems





#### Simulation

- Similar to testing
- If an error is found, the system is bugged
  - or the model is not faithful
  - actually, simulation in standard model checking is also used to understand if the model itself contains errors
- If an error is not found, we cannot conclude anything
- The error state may lurk somewhere, out of reach for the random choice in pick\_a\_state



#### Simulation vs Modeling

- However, for complex CPSs simulation is needed
- In fact, accurately modeling a complex CPS in a classical model checker is often too difficult or inconvenient
  - plant must be modeled by real variables: inherently infinite state systems
  - can be approximated, but accuracy may be low
- Simulators are often already available for testing, why can't we rely on them?
  - o not "real" model checking, but something close to it
  - far better than "simple" testing
- May be either build from scratch, or implemented with dedicated tools
  - C/Java/Python dedicated programs vs. Simulink/Modelica

# Simulation-based Model Checking

- Too many states, we cannot store them in an hash table
  - transition relation defined by a complex simulator, translation in OBDDs cannot be done
- Two main workhorses:
  - System Level Formal Verification chooses system inputs so as to cover as much as possible
    - mainly for safety, but also some sort of LTL may be used
  - Statistical Model Checking uses powerful statistical methods to perform model checking
    - something like Monte-Carlo sampling
    - i.e., we run the simulation several times, and we try to derive some guaranteed answer

#### Simulation

- A simulation is an experiment on a model
  - we focus on simulations performed by a computer
- Simulation is very easy to implement in the case of classical model checkers
  - no problems with RAM or execution time
- This stems from the fact that classical model checking deals with finite state systems
  - one step at a time, time passing typically not important
  - state space is finite and described by discrete-typed variables
  - computing a transition from a state to another is straightforward





#### Simulation

- What if we need to simulate a cyber-physical system?
  - e.g., simulate the Apollo mission
  - many subsystem interacting with each other via continuous signals
  - some subsystems are described by ODEs (Ordinary Differential Equations)
- In some cases, system developers also builds a simulator from scratch, e.g., in the C language
  - directly experimenting on the physical object may be dangerous, expensive, or simply impossible (e.g., it still does not exist)
- Many tools are available to easily describe complex models to be simulated
  - e.g., able to approximate solutions for ODEs
- Here we will deal with the open-source Modella
  - we will also have a look to Simulink



# Some Background: ODEs

With some simplification, an ODE is an equation

$$F(x, y, y', y'', \dots, y^{(n)}) = 0$$

- The unknown y is a function y = f(x), and  $y^{(n)}$  is the the *n*-th derivative of y w.r.t. x
- In our context, the independent variable is time, denoted by t
  - in simulations, we are interested in the system evolution over time...
- Thus, we have functions x = f(t) and n-th derivatives expressed in Newton's notation  $\dot{x}$
- Finally, our ODE is an equation

$$F(t, x, \dot{x}, \dots, \dot{\dot{x}}) = 0$$
UNIVERSITÀ DEGLI STUDI DELL'AQUILA







#### Some Background: ODEs

- We will consider explicit ODEs  $\dot{x} = F(t, x)$ 
  - x usually is in some n-dimensional space, e.g.,  $x \in \mathbb{R}^n$
  - thus, this is a system of equations
  - note that, with explicit ODEs, derivatives higher than 1 are not needed
  - simply put  $x_1 = x, x_2 = \dot{x}_1...$



# Example ODEs

- $\dot{x} = t + x$
- $(\dot{x}_1, \dot{x}_2) = (t + x_2 e^{x_1}, x_1 \log t)$
- Model of an infectious desease (HIV):

$$(\dot{x}_1, \dot{x}_2) = (\lambda - dx_1 - \eta \beta x_1 x_2, -\dot{x}_2(\dot{a} + I) + \eta \beta x_1 x_2)$$

- $x_1, x_2$  are uninfected and infected cells, I is an action by the immune system
- $a, d, \lambda, \beta, \eta$  are system parameters
- this latter example is time-invariant (see next slide)





#### **ODEs:** Euler Approximation

- Given the *time-invariant* ODE  $\dot{x} = F(x)$ , we may use the Euler approximation
  - for small au,  $\dot{x} pprox rac{x(t+ au)-x(t)}{ au}$
  - if we sample time with  $\tau$ , i.e., we only consider  $t \in \{0, \tau, 2\tau, \dots, k\tau, \dots\}\dots$
  - ... we have that  $F(x(k\tau)) = \frac{x(k\tau+\tau)-x(k\tau)}{\tau}$
  - thus by setting  $x_k = x(k\tau)$ , we have a discrete-time difference equation  $x_{k+1} = x_k + \tau F(x_k)$
- ullet This only works for small au and small k
  - it can be proved that  $||x_k x(k\tau)|| \le \tau \psi(k)$ , where  $\psi$  is *not* bounded
  - at least, in the general case



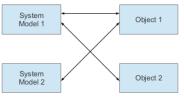




- A system is a mathematical concept used to study properties of physical objects
  - sometimes also called abstract system, or system model
- It is typically used to study evolutions as a function of time
- Virtually infinite examples:
  - population of rabbits
  - spread of diseases
  - physical objects: a fridge, an oven, a building, a car, ...
  - part of physical objects: a resistor, a brick, a wheel, ...
  - controllers for physical objects: ABS, autonomous driving, ...
- First distinction is among objects (what we want to model) and system (the mathematical model)
  - a system is defined through functions, sets, et



- Given an object, one may devise different systems
  - as we may have different programs to solve the same problem
  - not only because of different people doing it: different properties on the same object may be investigated
- Given a system, it may be applied to different objects
  - spreading of different diseases may have a common model
  - wheel of a car and of a motorcycle







- We start defining systems by looking at their inputs and outputs
  - keeping in mind that it is all as a function of time
- Deterministic systems: given an input sequence from some "start", the output is the same
  - probabilistic systems also exist, we do not consider them here
- Black-box system: at first, we perform experiments on the system
  - we provide sequences of inputs and observe the sequence of outputs





- We begin experiments at some time  $t_0 \in T$ , with  $T \subseteq \mathbb{R}$ 
  - ullet for some systems,  $\mathcal{T}\subseteq\mathbb{N}$
- We consider all input functions  $u: T \to U$  for our object
  - U is some set on which inputs may vary
  - ullet it may be multidimensional, e.g.  $U=\mathbb{N} imes\mathbb{Z} imes\mathbb{R}^2$
  - of course, such input functions are uncountably many, this is a conceptual experiment
- For each u, we have an output function  $y: T \to Y$  coming out of the object
  - U and Y may be different
  - again, Y may be multidimensional
- We define the system  $S = \{(u, y) \mid u \text{ is an input function and } y \text{ the corresponding output function}\}$ 
  - ullet thus,  $\mathcal{S} \subset \mathcal{U} \times \mathcal{Y}$

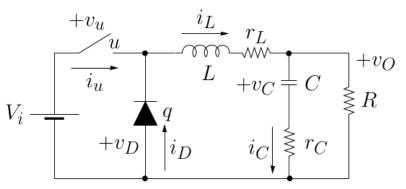


- Example: determine the number of students which graduate in same bachelor course
  - ullet assumption: student enrolls once every year, thus  $T\subset \mathbb{N}$
  - $U \subset \mathbb{N}$ : number of students enrolling "from outside"
  - $Y \subset \mathbb{N}$ : number of graduated students
  - $\mathcal{U} = \{f \mid f : T \to U\}$ , analogous for  $\mathcal{Y}$
  - example of input-output:
    - $u_1(2020) = 200$ ,  $u_1(2021) = 221$ ,  $u_1(2022) = 198$ , and  $u_1(x) = 0$  for  $x \notin \{2020, 2021, 2022\}...$
    - ... and we observe  $y_1(2020) = 51$ ,  $y_1(2021) = 51$ ,  $y_1(2022) = 60$
    - $u_1(2020) = 136$ ,  $u_2(2021) = 231$ ,  $u_2(2022) = 90$ , and  $u_2(x) = 0$  for  $x \notin \{2020, 2021, 2022\}...$
    - ... and we observe  $y_2(2020) = 42, y_2(2021) = 37, y_2(2022) = 98$
    - $u_1, u_2 \in \mathcal{U}, y_1, y_2 \in \mathcal{Y}, (u_1, y_1), (u_2, y_2) \in \mathcal{Y}$





Example: determine the output voltage of a buck DC-DC converter





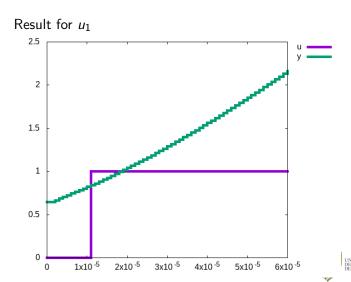


- Example: determine the output voltage of a buck DC-DC converter
  - $\bullet$   $T \subset \mathbb{R}$
  - $U \subset \{0,1\} \times \mathbb{R}$ 
    - u may be closed (0) or open (1) at any time
    - ullet  $V_i$  may be any real number
  - $Y \subset \mathbb{R}$ : observed output voltage  $v_O$
  - example of input-output (times are in microseconds):
    - $u_1(t) = (0,5)$  for all  $t \in [0,10]$ ,  $u_1(t) = (1,5)$  for all  $t \in [10,100]$
    - $u_2(t) = (0,15)$  for all  $t \in [0,9]$ ,  $u_2(t) = (1,10)$  for all  $t \in (9,15]$ ,  $u_2(t) = (0,7)$  for all t > 15

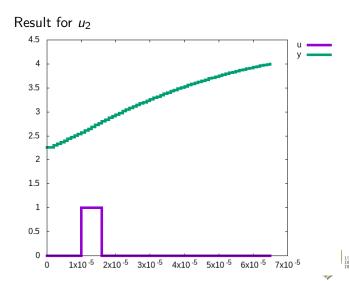








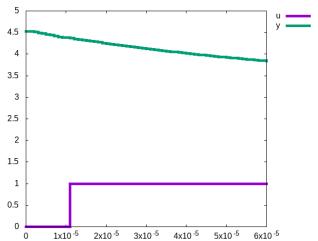






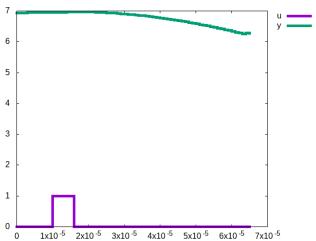


... but this is also a result for  $u_1$ 





... and this is also a result for  $u_2$ 





- Is this a non-deterministic system??? NO!
- The point is that output is not determined by input only
  - though for some systems this is the case: number of students above
- The missing element is the state
  - essentially, the input/output function has side effects...
- Thus, the output (for deterministic systems) is a function of both the input and the state
  - in the examples above, we made different choices for the starting state





- For our purposes, a system will be defined by a 6-tuple  $S = \langle T, U, Y, X, \eta, \phi \rangle$ :
  - *U*, *Y* are sets of possible input and output values, resp.
  - T is a set of times
    - ullet if  $T\subseteq\mathbb{R}$  then we have a *continuous-time* system
    - if  $T \subseteq \mathbb{N}$  then we have a *discrete-time* system
  - X is a set of states
    - may be either finite or infinite
    - if  $T\subset \mathbb{N}$  and  $|X|<\infty$  then we essentially have a Kripke structure
  - $\eta: T \times X \times U \rightarrow Y$  defines the output function
  - $\phi: T \times T \times X \times \mathcal{U} \to X$  defines the state transition function
    - recall that  $\mathcal{U} = \{f \mid f : T \to U\}$







- $\eta: T \times X \times U \rightarrow Y$  is as expected
  - given the current time, the current state, and the current input, we can compute the output
- $\phi: T \times T \times X \times \mathcal{U} \to X$  is somewhat more complicated than expected
  - one would expect  $\phi: T \times X \times U \to X$
  - actually, this is enough for most systems
- For some systems, the state transition function depend on some sequence of inputs, not only the last one
  - so we need a function, defined at least on an interval  $[t_0, t)$
  - $\bullet$  this is why  $\phi$  also takes two times instead of one
- 3 conditions must hold for  $\eta$  and  $\phi$ : causality, consistency and separation

# Some Background: Causal Systems

- $\forall t, t_0 \in T, x_0 \in X. (t \ge t_0 \land u|_{[t_0,t)]} = u'|_{[t_0,t)}) \Rightarrow \phi(t, t_0, x_0, u|_{[t_0,t)]}) = \phi(t, t_0, x_0, u'|_{[t_0,t)})$
- That it is, if we fix the first 3 arguments  $t, t_0, x_0$  of  $\phi$ ...
- ... and we provide, as a fourth argument, two possible different functions u, u'...
- ... which however output the same values in the interval  $[t_0, t)$ ...
- ullet ... then the final value of  $\phi$  does not change
- Thus, what happens in the interval  $[t_0, t)$  causes the system to go to one single state





# Some Background: Consistent Systems

- $\forall t \in T, x_0 \in X, u \in \mathcal{U}. \ \phi(t, t, x_0, u) = x_0$
- Recall that, for a call  $\phi(t, t_0, x, u)$ , u is considered in the interval  $[t_0, t)$
- Thus, in a call  $\phi(t, t, x_0, u)$ , we are considering the empty interval [t, t)
- Hence, we have no input at all!
- Of course, without inputs, the system cannot change its current state





# Some Background: Separation Property of Systems

- $\forall t, t_0, t_1 \in \mathcal{T}, x_0 \in \mathcal{X}, u \in \mathcal{U}. (t > t_1 > t_0) \Rightarrow \phi(t, t_0, x_0, u|_{[t_0, t_1)}) = \phi(t, t_1, \phi(t_1, t_0, x_0, u|_{[t_0, t_1)}), u|_{[t_1, t)})$
- In few words: the state you obtain if you go straight from t<sub>0</sub> to t, is the same state you would obtain if:
  - you first go from  $t_0$  to some intermediate  $t_1$ 
    - i.e.,  $x_1 = \phi(t_1, t_0, x_0, u|_{[t_0, t_1)})$
  - and then from  $t_1$  to t
    - i.e.,  $\phi(t, t_1, x_1, u|_{[t_1,t)})$





# Some Background: Hybrid Systems

- Note that the set of states X may be multi-dimensional
  - e.g.,  $X = \mathbb{R}^3$ , or  $X = \{1, 2, 3\} \times \mathbb{Z}$
- ullet Thus, also  $\phi$  may be multi-dimensional
- Informally, if X has dimension n, then we will have n state variables
  - recall that the same holds for U, Y: we may have multiple input and output variables
- Hybrid systems: those for which some variables are continuous and other are discrete
  - in some texts, a "hybrid system" have some variables depending on  $T=\mathbb{N}$  and some other on  $T=\mathbb{R}$
- This is exactly the case of cyber-physical systems!
  - plant + controller/monitor
  - plant is continuous, controller/monitor is discrete





#### Some Background: Special Systems

• With some semplification, a system is *time-invariant* iff  $\forall t, t_0, t_1 \in T, x \in X, u \in \mathcal{U}. \ \phi(t, t_0, x, u) = \phi(t - t_0, 0, x, u) \land \eta(t, x, u(t)) = \eta(t_1, x, u(t_1))$ 

- that is, the absolute time is not important
- the relative time is
  - given a state x, system evolution from 1 to 3 seconds and from 10 to 12 seconds is the same
- For time-invariant systems, we can always set  $t_0 = 0$
- For time-invariant systems, we can also write  $\dot{x}(t) = \phi(x(t), u(t)), y(t) = \eta(x(t), u(t))$





# Some Background: Special Systems

- With some semplification, a system is linear iff
  - U, Y, X are linear spaces
    - that is, any linear combination  $\sum_{i=1}^{n} a_i x_i$  is in X etc
  - $\mathcal{U}$  is a linear subspace of  $U^T = \{f \mid f : T \to U\}$ 
    - again, any linear combination  $\sum_{i=1}^n a_i u_i(t)$  is in  $\mathcal U$
  - fixed any 2 times  $t, t_0 \in T$  as first 2 arguments,  $\phi$  is linear in the remaining 2 arguments
    - $\phi(t, t_0, x, u) = A \cdot [x, u] + b$  for some A and b
    - A, b may depend on  $t, t_0$ , but not on x, u
  - fixed any time  $t \in T$  as first argument,  $\eta$  is linear in the remaining 2 arguments
- Linear systems are easy to model, simulate and verify
- With some semplification, a system is:
  - a finite-state system if U, X, Y are finite sets (Kripke structure)
  - a finite-dimensional system if U, X, Y are linear photostation spaces



# Some Background: Generating Functions

- For discrete-time systems, we have that  $x(t+1) = \phi(t+1,t,x(t),u|_{[t,t+1)}) = \phi(t+1,t,x(t),u(t)) = f(t,x(t),u(t))$ 
  - $\bullet$  first and second argument of  $\phi$  are not independent...
  - f has the same domain of  $\eta$
- For continuous-time systems, we focus on *regular* systems, i.e., those systems for which  $\phi$  is differentiable and there exists a function f s.t.
  - $\frac{d\phi(t,t_0,x,u)}{dt} = f(t,\phi(t,t_0,x,u),u(t))$
  - with the initial condition that exists an  $x_0 \in X$  s.t.
    - $x_0 = \phi(t_0, t_0, x_0, u)$
  - ullet often, it is easier to provide f than  $\phi$
- Using Newtonian notation, we have  $\dot{x}(t) = f(t, x(t), u(t))$
- For time-invariant systems, we have  $\dot{x}(t) = f(x(t), u(t)), y(t) = \eta(x(t))$







# Some Background: System For Students' Example

- $X \subseteq \mathbb{N}^3$ ,  $U, Y \subseteq \mathbb{N}$ ,  $T = \mathbb{N}$
- Parameters  $\alpha_i(t) \in [0,1]$  is ratio of students passing an year t
- $x_1(t+1) = (1 \alpha_1(t))x_1(t) + u(t)$
- $x_i(t+1) = (1 \alpha_i(t))x_i(t) + \alpha_{i-1}(t)x_{i-1}(t)$  for i = 2, 3
- $y(t) = \alpha_3(t)x_3(t)$
- Note that, if  $\alpha_i(t) = 1$  for all t, then states are not needed, as we have y(t) = u(t-3)
- Summing up:

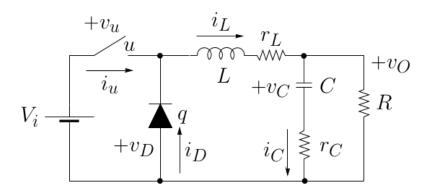
• 
$$f(t) = \begin{pmatrix} (1 - \alpha_1(t))x_1(t) + u(t) \\ (1 - \alpha_2(t))x_2(t) + \alpha_1(t)x_1(t) \\ (1 - \alpha_3(t))x_3(t) + \alpha_2(t)x_2(t) \end{pmatrix}$$







# Some Background: System For Buck DC/DC Converter







# Some Background: System For Buck DC/DC Converter

- $X \subseteq \mathbb{R}^2$ ,  $U \in \{0,1\} \times \mathbb{R}$ ,  $Y \subseteq \mathbb{R}$ ,  $T = \mathbb{R}$
- $L, C, R, r_L, r_C \in \mathbb{R}$  are system parameters
  - we will also use 6 real numbers  $a_{i,j}$  for  $i \in \{1,2\}, j \in \{1,2,3\}$  which are functions of such parameters
  - e.g.,  $a_{2,3} = -\frac{1}{L} \frac{r_C R}{r_C + R}$
- Variables for state are  $i_L, v_O, v_D, i_D, v_u, i_u$  (real) and q (boolean)
- Variables for input are u (boolean) and  $V_i$  (real)
- $y(t) = v_O(t)$ , thus  $\eta$  is easy
- $\bullet$   $\phi$  is defined by cases in the following slide
  - for the definition of  $\phi$ , some other (auxiliary) variables are useful:  $v_D$ ,  $i_D$ ,  $v_u$ ,  $i_u$  (real) and q (boolean)
  - $R_{on} \approx 0, R_{off} >> R_{on}$  are fixed parameters



# Some $\mathsf{Background}\colon\mathsf{System}$ For $\mathsf{Buck}$ $\mathsf{DC}/\mathsf{DC}$ Converter

We omit (t) for better readability There must exist a value for  $v_D, i_D, v_u, i_u \in \mathbb{R}, q \in \{0,1\}$  s.t.

$$i_L = a_{1,1}i_L + a_{1,2}v_O + a_{1,3}v_D$$
 (1)

$$\dot{v_O} = a_{2,1}i_L + a_{2,2}v_O + a_{2,3}v_D$$
 (2)

$$q \rightarrow v_D = R_{\rm on} i_D$$
 (3)  $\bar{q} \rightarrow v_D = R_{\rm off} i_D$  (7)

$$q \rightarrow i_D \geq 0$$
 (4)  $\bar{q} \rightarrow v_D \leq 0$  (8)

$$u \rightarrow v_u = R_{\rm on}i_u$$
 (5)  $\bar{u} \rightarrow v_u = R_{\rm off}i_u$  (9)

$$v_D = v_u - V_{in}$$
 (6)  $i_D = i_L - i_u$  (10)

Both ODEs and algebraic equations







#### Modeling in Modelica

- Modelica is an open-source language for specifying (complex) systems
  - developed by experts starting in late 1990s
- Many implementations exist
  - OpenModelica+simForge, Dymola, Simulation X, MapleSim, MathModelica
  - here we will stick to OpenModelica+simForge
- Also see Modelica slides





#### Modeling in Modelica

- Object-oriented language: classes and objects (i.e., class instances)
  - strongly typed
- Compositional modeling:
  - break up the system in subsystems (components)
  - connect the components
- Very useful for complex systems, with many components
  - some standard components already defined, e.g., resistors, flows etc
- May use equations, also with derivatives
- Generates a C program, thus it is very efficient





#### Modeling in Modelica

- Synchronous data flow principle: time is the same for all components
  - such as clocks for digital systems, but in Modelica it may be continuous
- May specify "algorithms" using assignments, ifs, whiles, etc
  - all variables must be instances of some class
  - this also includes integers and reals
- Acausal modeling: simply first provide the equations for each object, then connect the objects between them
  - other modeling languages, e.g., Simulink, requires to first design the full chain of connections...
  - ...and to make computation in sequence
  - Modelica allows both causal and acausal modeling
  - physical "reality" is lost
- Modelica easier for modelers, Simulink easier for computers



#### **Acausal Modeling**

#### The order of computations is not decided at modeling time

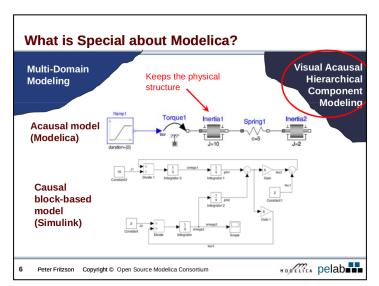
	Acausal	Causal
Visual Component Level	Toget bots Syrigt	
Equation Level	A resistor <i>equation</i> : R*i = v;	Causal possibilities: i := v/R; v := R*i; R := v/i;





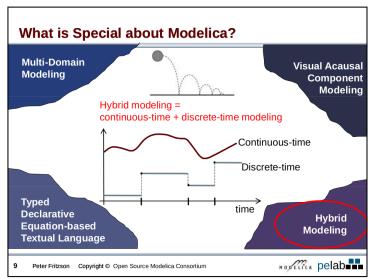
HODELICA pelab

#### Modelica: Acausal Modeling





#### Modelica: Hybrid Modeling







#### Modelica: Toy Example

• Text file with .mo extension, let us say model.mo

```
class Example
  output Real x, y, z;
  algorithm
    when initial() then //at 0, both this...
      x := 0; // Pascal-like assignments
    elsewhen sample (0, 1) then // \dots and this
      x := 1;
      y := pre(x); //0 till 1, then always 2
    elsewhen sample(0, 0.5) then
        //elsewhen order is important! from bottom to top
      x := 2;
      z := pre(x); //0 till 0.5, then always 1
    end when;
end Example;
```

#### Modelica: Toy Example

Text file with .mos extension, let us say run.mos



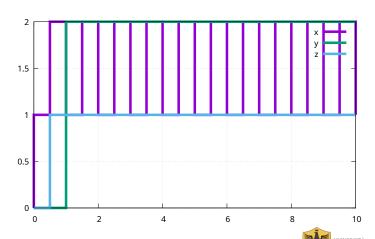


#### Modelica: Toy Example

- Run the command omc run.mos
  - of course, you must have installed omc for your OS
- This has the following effect:
  - generates a C program model.c
  - Occupiles model.c to obtain the executable file model
  - executes model
  - outputs both a file Example\_res.mat and a graphical window with the graph of variables x, y, z as function of time

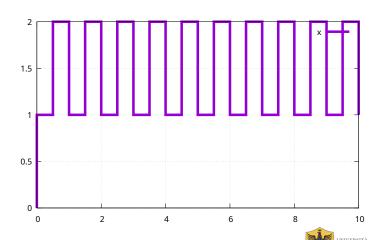




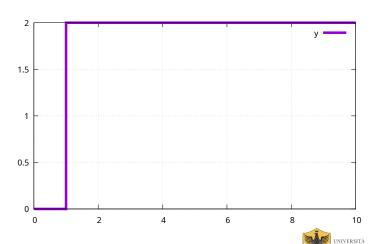






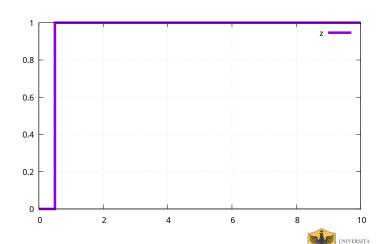






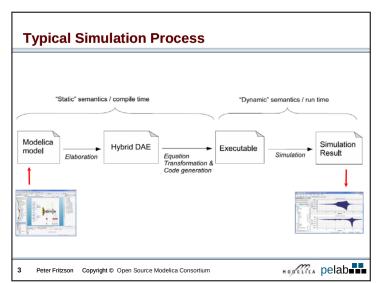














# Modelica: Toy Example Examined

```
class Example
  output Real x, y, z;
  algorithm
    when initial() then //at 0, both this...
      x := 0:
    elsewhen sample (0, 1) then // \dots and this
      x := 1:
      y := pre(x); //0 till 1, then always 2
    elsewhen sample(0, 0.5) then
      x := 2;
      z := pre(x); //0 till 0.5, then always 1
    end when;
end Example;
```

# Modelica: Toy Example Examined

- Example is a class defined by the modeler: Modelica is OO
- It has 3 real-valued variables, which may become the input for other blocks
- The dynamics is an algorithm based on the sample construct
  - when initial() C executes code in C at time 0
  - when sample (A, B) C executes code in C every A + Bx seconds, for  $x \in \mathbb{N}$
  - there may be multiple elsewhen sample(A, B) C triggered at a given time
  - they are all executed, starting from the bottom of the file
  - now explain the output of the previous example...
- In expressions, pre(var) holds the value of var before the current event



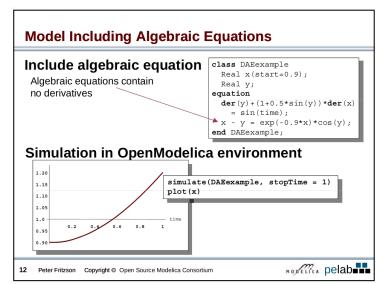
#### Modelica: Derivatives

#### Simplest Model – Hello World! A Modelica "Hello World" model Equation: x' = -xclass HelloWorld "A simple equation" Real x(start=1): Initial condition: x(0) = 1equation der(x) = -x: end HelloWorld; Simulation in OpenModelica environment simulate(HelloWorld, stopTime = 2) 0.8 plot(x) 0.6 0.4 0.2 0.5 1.5 HODELICA pelab Copyright @ Open Source Modelica Consortium 11





# Modelica: Derivatives and Algebraic Equations





# Modelica: Derivatives and Algebraic Equations

- time is a special variable, holding current simulation time
- System dynamics of previous example is defined as

$$\dot{y} + \left(1 + \frac{\sin y}{2}\right)\dot{x} = \sin t$$
$$x - y = e^{-0.9x}\cos y$$

 Can be transformed in "normal" form by adding state variables:

$$\dot{x} = \frac{\sin t - y_1}{1 + \frac{\sin y}{2}}$$

$$\dot{y} = y_1$$

$$z_1 = e^{-0.9x}$$

$$z_2 = \cos y$$

$$x = z_1 z_2 + y$$

$$y = x - z_1 z_2$$







# Modelica: Subsystems and Connections

- Till now, stand-alone systems with just one component
- Modelica allows compositional modeling of many components
- Each component is modeled autonomously, by simply looking at the interaction with the environment (input/output)
- Complex systems are made of *connected* components
- Connectors can be explicitly defined
  - causal: input/output relation is explicitly stated
  - acausal: input/output relation is left unspecified
  - Modelica will understand which is input and which is output





```
Mind the difference between = and :=
model ContinuousBehav

Boolean x;
Real i (start = 1);

equation
   (if x then 0.5*time else -0.1*time)*der(i) = time;
end ContinuousBehav;
```





```
model GenerateBoolInputs
  Boolean x;
  parameter Real sampling = 1.0;
  algorithm
    when initial() then
      x := false;
    elsewhen sample(sampling, sampling) then
      x := not(x);
    end when;
end GenerateBoolInputs;
```

```
model BoolCont

GenerateBoolInputs gbi;
ContinuousBehav cb;

equation
   gbi.x = cb.x;
end BoolCont;
```









#### Modelica: Passing of Time

- For all objects defined, the time passes in the same way
  - it is a kind of common clock, as in digital circuits
- This is of course consistent with physical reality
  - components are close enough...
- It is always continuous time, but using sample we can also have discrete time





#### Modelica: Algorithms and Equations

- Both may be used, the modeler has to choose
  - of course, x := x + 1 inside an algorithm is ok, x = x + 1 in an equation is not
  - using imperative vs. declarative style is left to the modeler
  - in some cases, algorithm is more natural, in some other equation has to be preferred
  - note that loops and ifs are available in both formats
  - e.g., a = (if b then 1 else 2); vs if (b) then a:=1; else a:=2; end if;
  - or simply a := (if b then 1 else 2);
- Algorithms, as well as equations solving, does not cause time to pass
  - number of computation steps required is not important



#### Modelica: Algorithms

- Generally speaking, when A B clauses triggers the corresponding block B when condition A is true
- A can be any boolean expression, not only sample
- Functions may also be defined and used
  - time does not pass during function calls
  - again, number of computation steps is not important
  - must have input and output
- External C or Fortran functions may be called
  external "C" result = myfun();
  annotation(Include = "#include \"myfile.c\"");





#### Modelica: Events

- Discrete events happens in a discrete number of time points
  - given that the simulation terminates somewhere, it is actually a finite number of points
- We saw initial and sample, there is also terminal
  - triggered at the end of the simulation
- Simulation ends either because of:
  - the stopTime attribute inside simulate command
  - a terminate statement





- Simulink is a graphical extension to MATLAB
  - MATLAB itself is proprietary, but UnivAQ provides it to students
- Main goal: modeling and simulation of systems
  - also non-linear ones
- Also see https://ctms.engin.umich.edu/CTMS/index. php?aux=Basics\_Simulink
- No way of simply writing a text file: you have to use the GUI and manipulate graphical objects
  - model files are saved in a binary proprietary SLX format



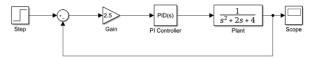


- Two major classes of objects: blocks and lines
  - blocks used to generate, modify, combine, output, and display signals
  - lines used to connect blocks, i.e., transfer signals from one block to another
  - again, a common clock for all objects in a model
- Suppose you create a new or open an existing Simulink model file
- How to add a new block:
  - click "Library Browser"
  - select the type of block you need
    - hundreds of types available, could also be searched by name
  - drag it to the model window
  - by double clicking, you can edit the properties





- How to add a new connecting line:
  - simply drag the mouse from the first object to the second object
- If you are connecting an object with a line:
  - first make a dangling line from the destination
  - connect the end of such line with the "source" line
  - this will make the source line bifurcated







#### Most notable types of blocks:

- Sources: used to generate various signals
- Sinks: used to output or display signals
- Continuous: continuous-time system elements
  - transfer functions, state-space models, PID controllers, etc.
- Discrete: linear, discrete-time system elements
  - discrete transfer functions, discrete state-space models, etc.
- Math Operations: contains many common math operations
  - gain, sum, product, absolute value, etc.
- Ports & Subsystems: contains useful blocks to build a system



