

# Approximate Solution

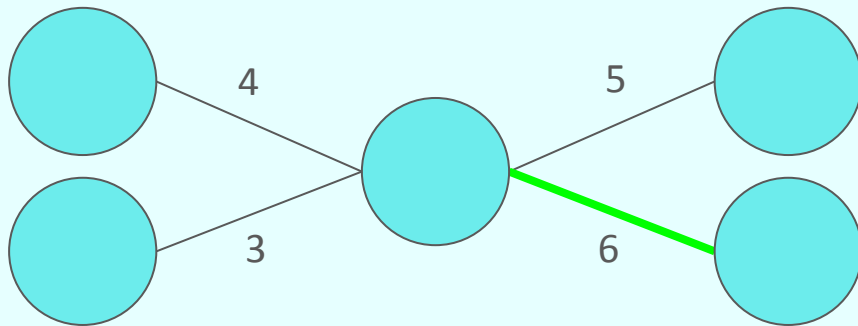
### Algorithm (Approximate Greedy Longest Path):

For each vertex in the graph: → runs  $n$  times  
start a new path at that vertex  
mark it as visited

While the current vertex has any unvisited neighbors: →  $\leq n$  iterations (each vertex visited once)  
choose the unvisited neighbor connected by the heaviest edge →  $\leq n$  (scans all neighbors)  
add that neighbor to the path  
mark it visited  
move to that neighbor

keep track of the longest path found across all starting vertices

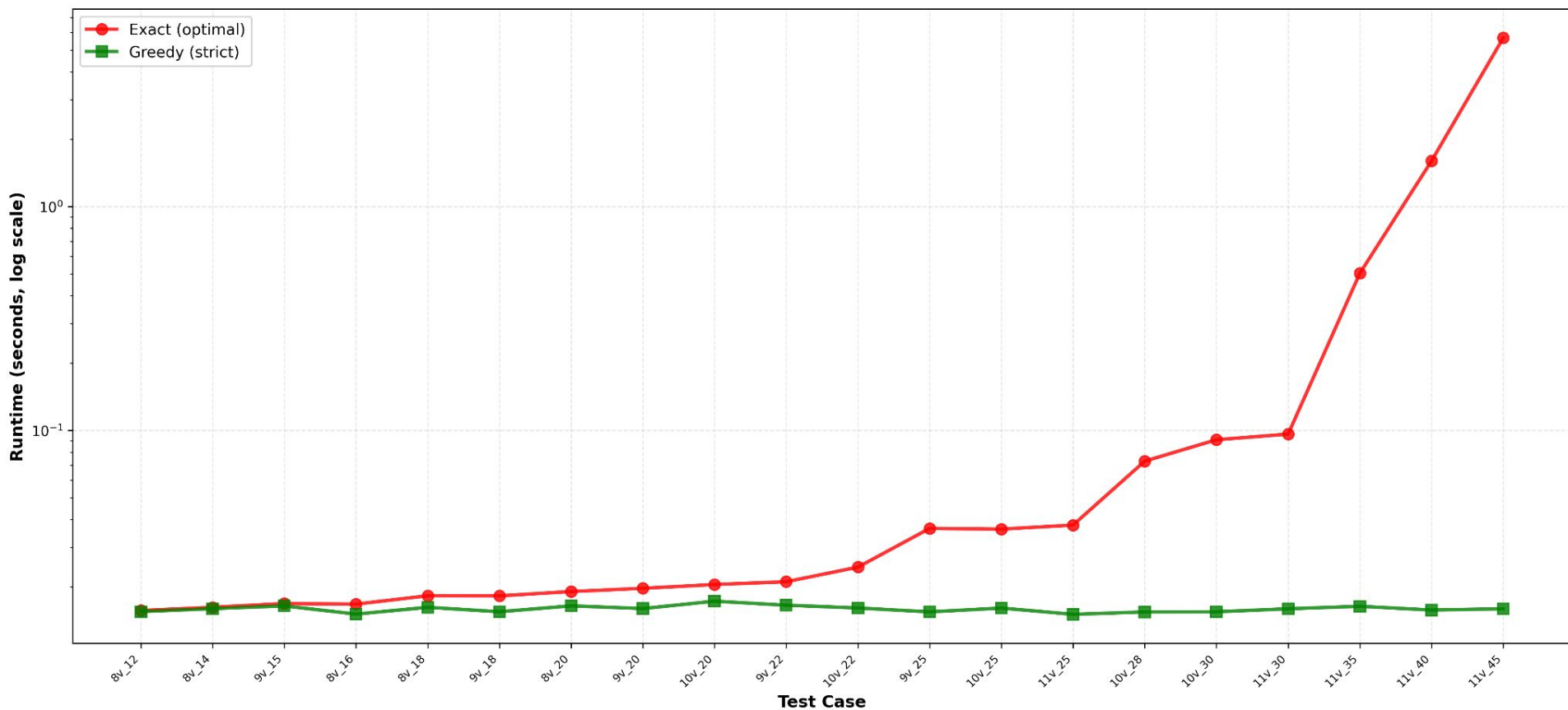
return the longest path discovered



Runtime Analysis:

$O(n^3)$

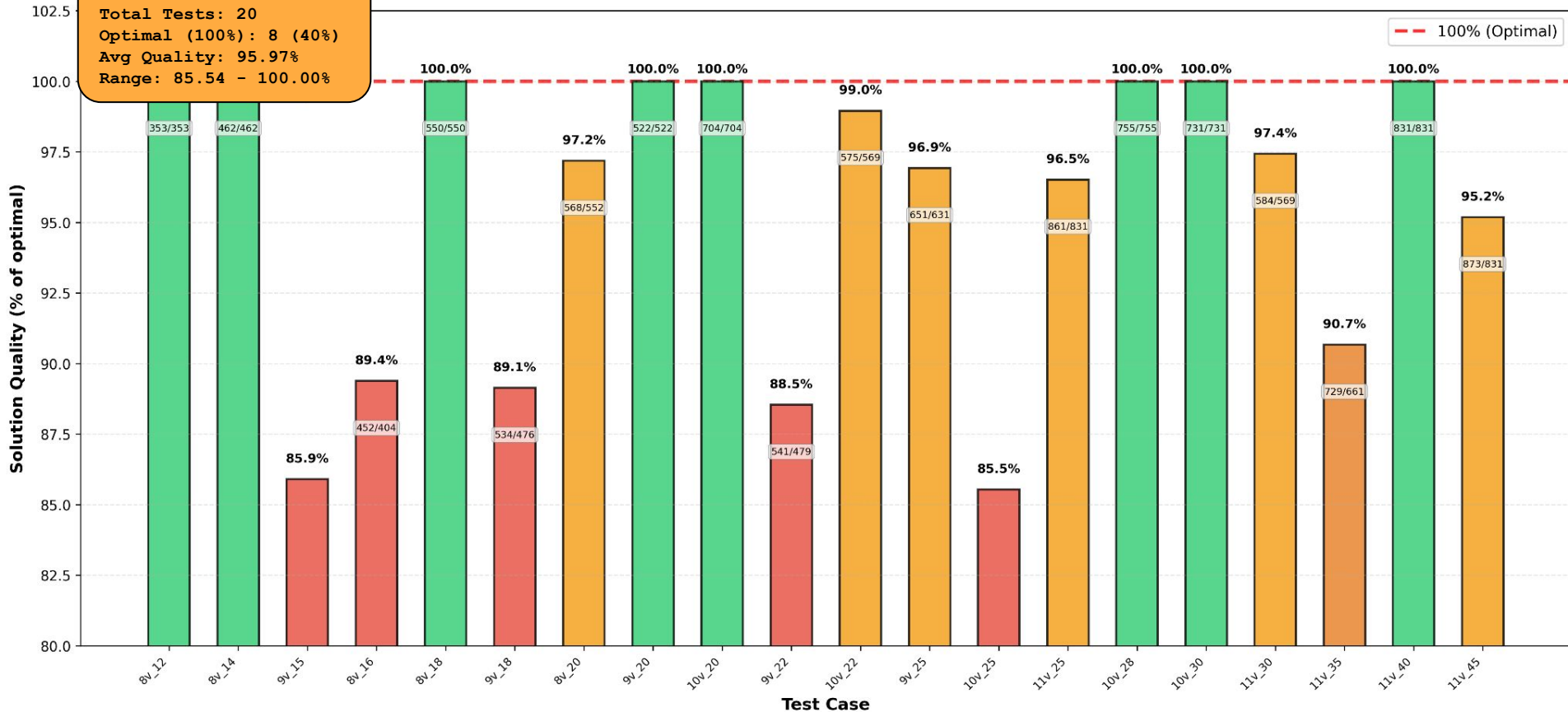
### Exact vs Greedy Runtime Comparison (Part D Additional Test Cases)



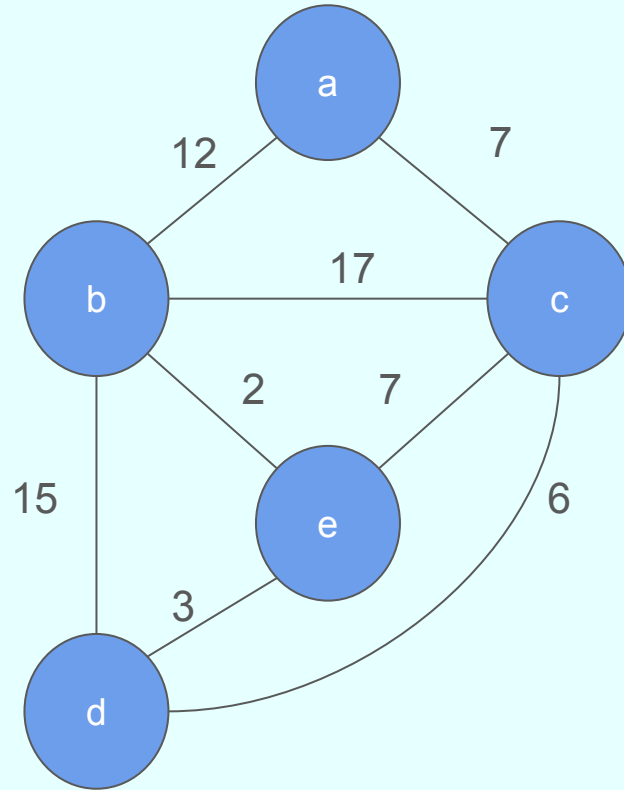
8v\_16  $\rightarrow$  8 vertices, 16 edges

## Greedy Solution Quality vs Optimal (Percentage = Greedy/Exact × 100%)

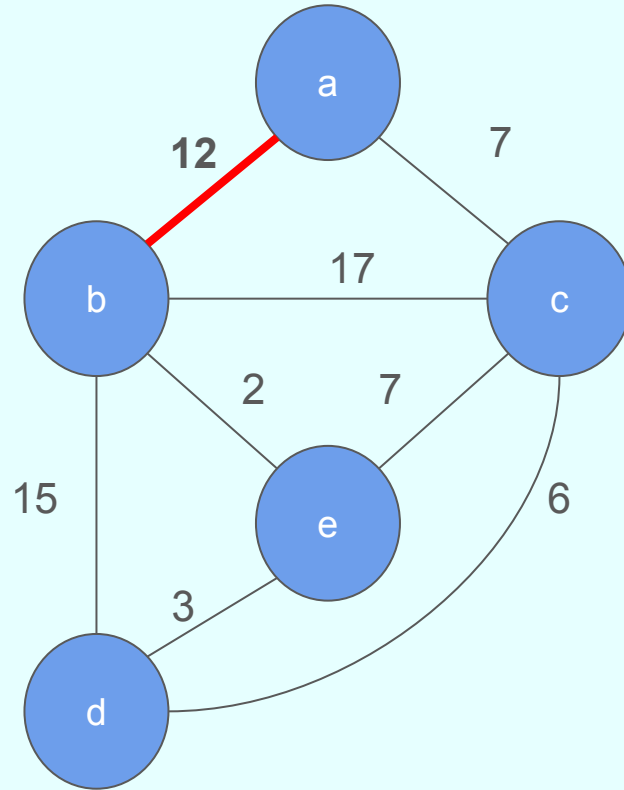
**Statistics:**  
Total Tests: 20  
Optimal (100%): 8 (40%)  
Avg Quality: 95.97%  
Range: 85.54 - 100.00%



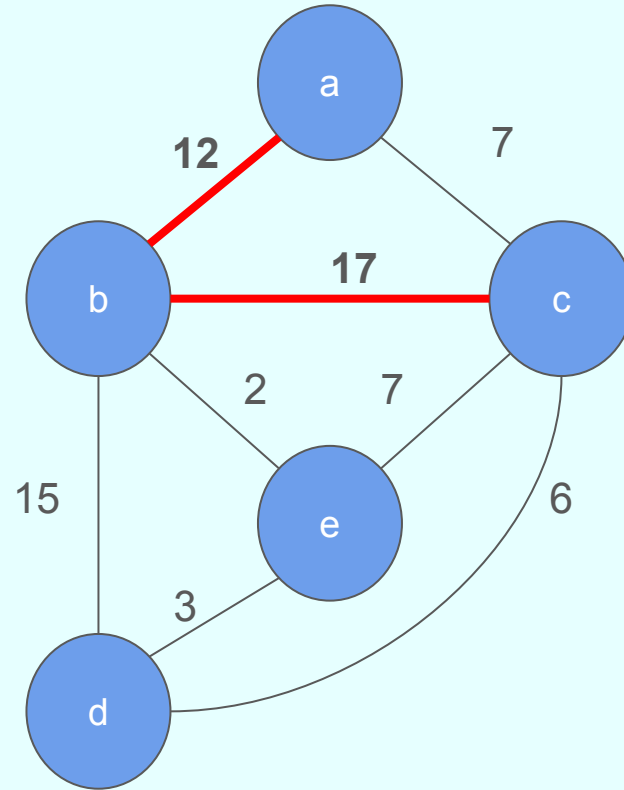
## Why Greedy Can Fail



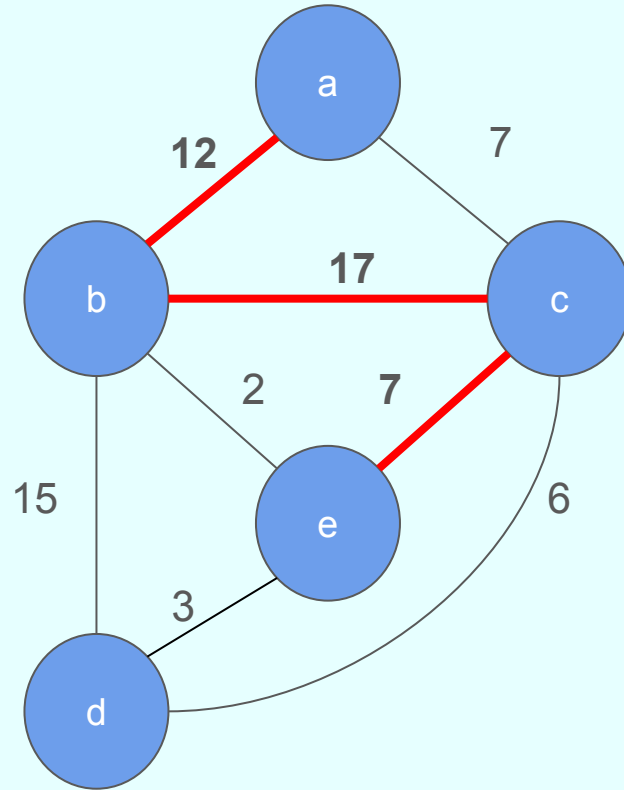
Greedy starting at a



Greedy starting at a

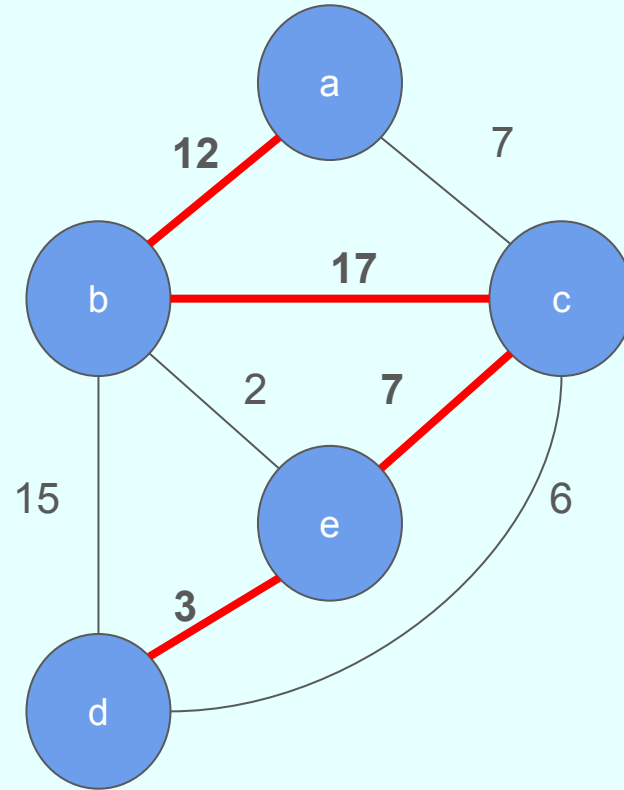


Greedy starting at a



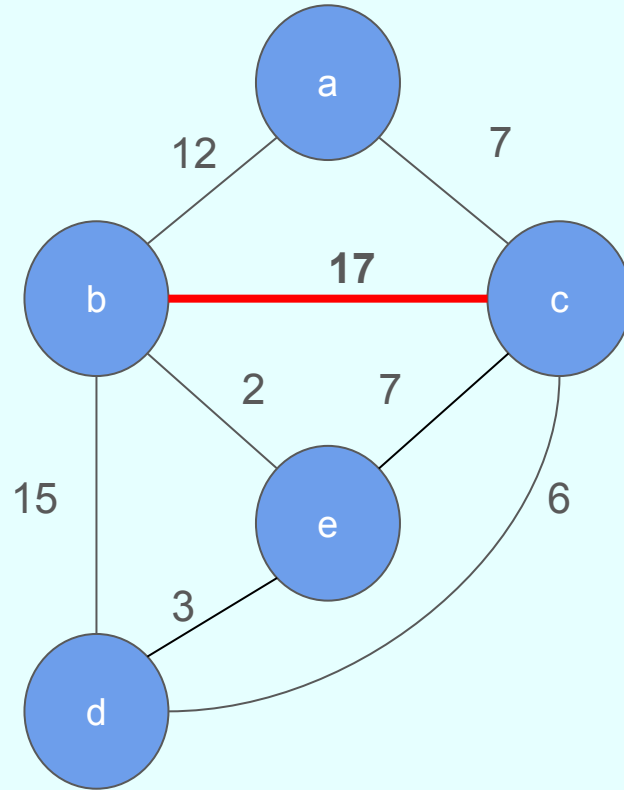


Greedy starting at a

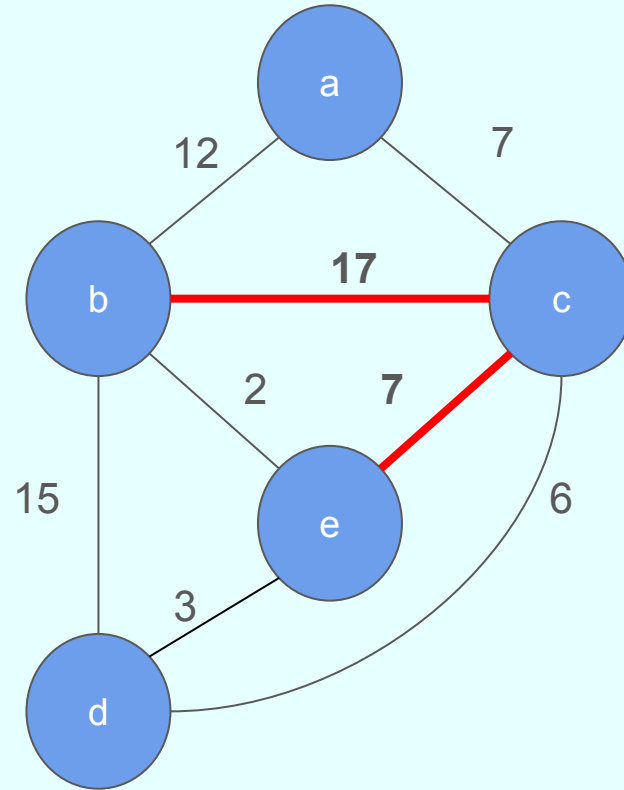


39 total

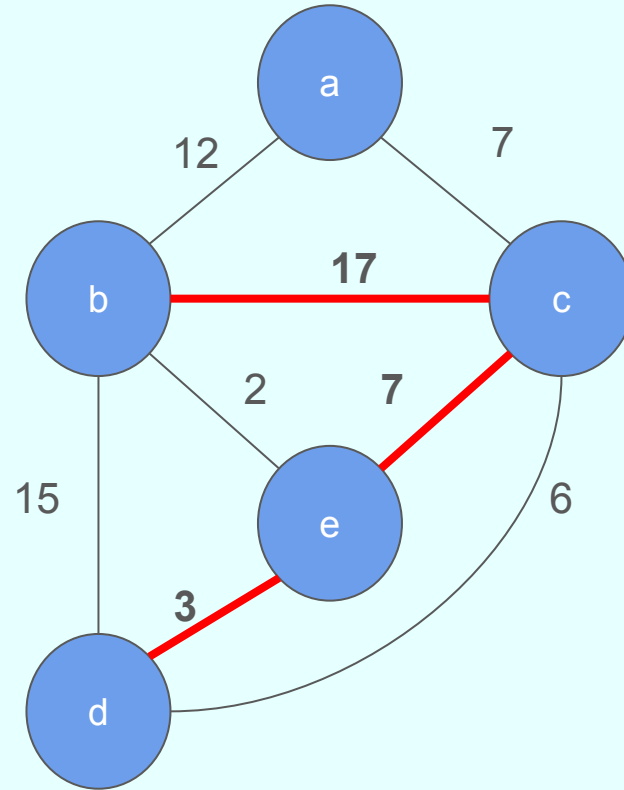
Greedy starting at b



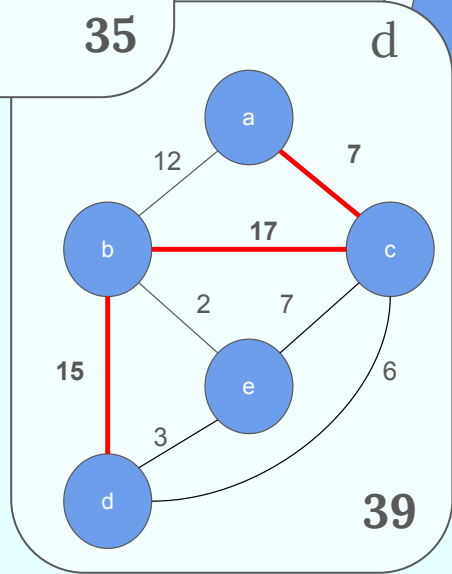
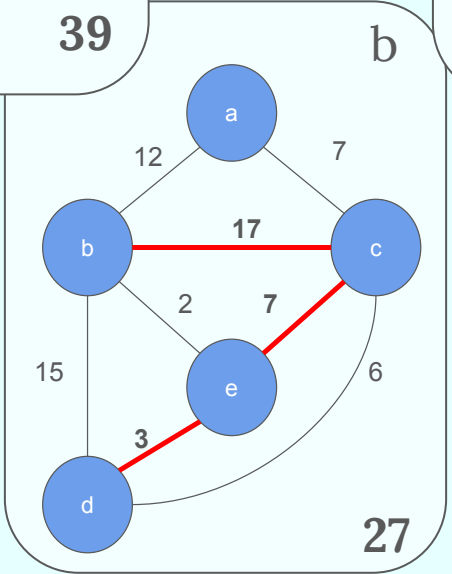
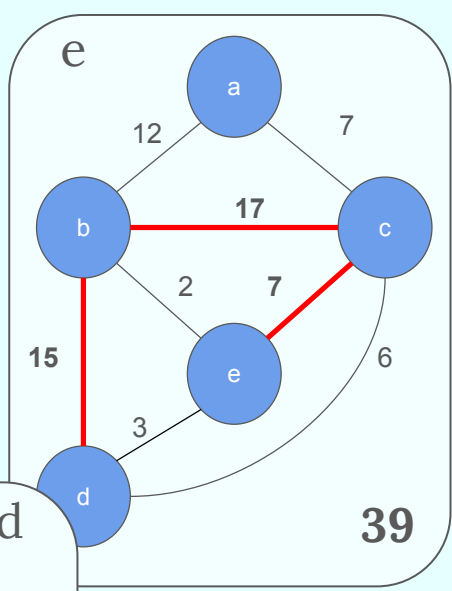
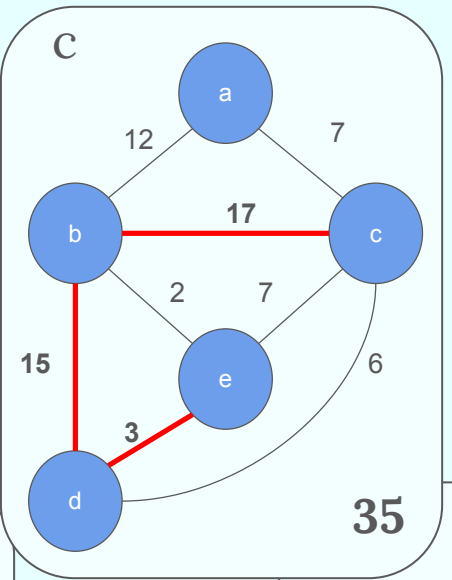
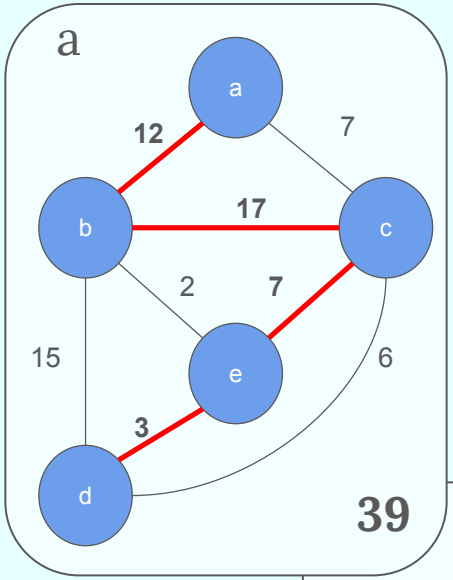
Greedy starting at b



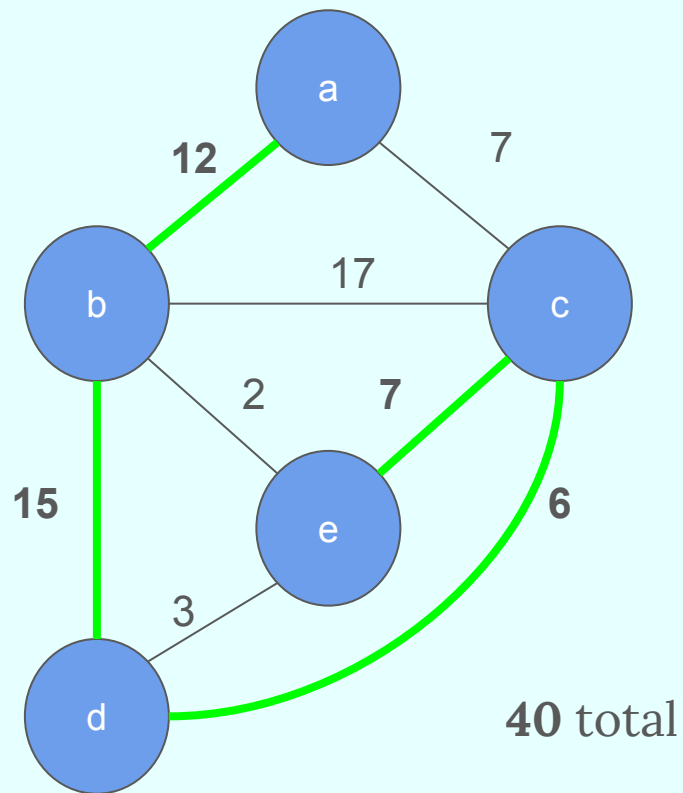
Greedy starting at b



27 total



## Optimal solution



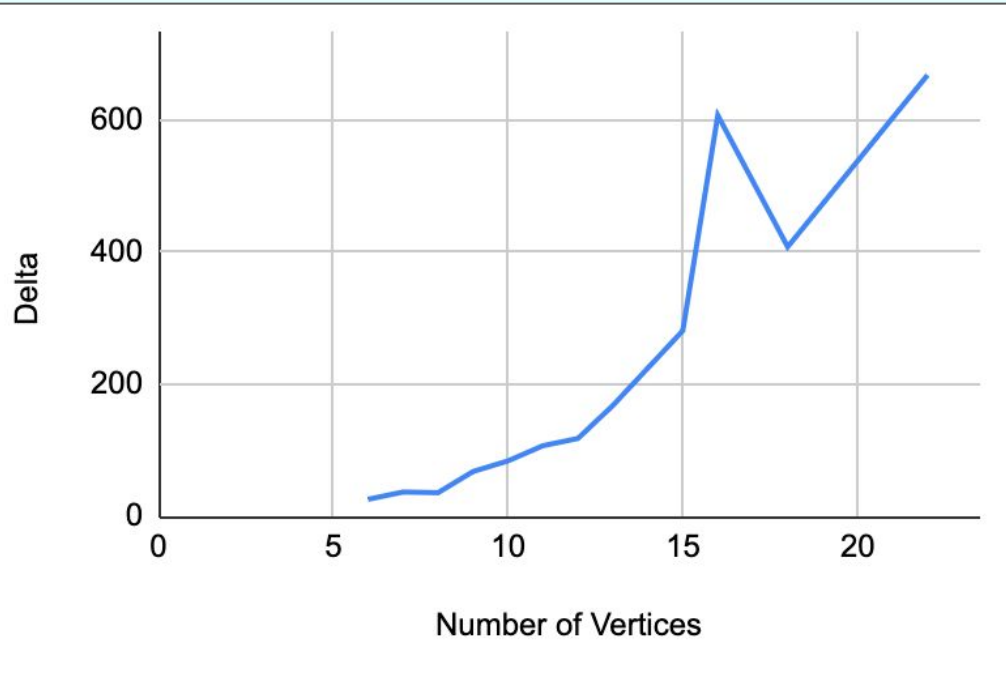
## Lower Bound and $\Delta$

- $\Delta$  = Greedy Path Weight – Lower Bound
- Maximum edge weight is valid lower bound
- $\Delta$  grows dramatically on large graphs  $\rightarrow$  greedy scales well.

Case	Vertices	LB	Greedy	$\Delta$
test_input_4	6	9	35	26
test_input_medium_1	7	12	49	37
test_input_5	8	10	46	36
test_input_medium_2	9	15	83	68
test_input_6	10	20	104	84
test_input_large_1	11	20	127	107
test_input_7	12	18	136	118
test_input_large_2	13	25	193	168
test_input_8	15	30	311	281
test_input_dense_large	16	60	667	607
test_input_large	18	34	442	408
test_input_extreme	22	45	713	668



## $\Delta$ vs Number of Vertices



$\Delta$  tells how much better this solution is than the lower bound

$\Delta$  increases with vertices because greedy paths grow with the graph while lower bound remains small