

## 7.4 GREEDY APPROXIMATION FOR SET COVER

### Question:

Implement a greedy approximation algorithm for the Set Cover problem. Analyze its performance on different input sizes and compare it with the optimal solution. Consider the following universe  $U = \{1, 2, 3, 4, 5, 6, 7\}$  and sets  $S = \{\{1, 2, 3\}, \{2, 4\}, \{3, 4, 5, 6\}, \{4, 5\}, \{5, 6, 7\}, \{6, 7\}\}$ .

### Input:

- Universe  $U = \{1, 2, 3, 4, 5, 6, 7\}$
- Sets  $S = \{\{1, 2, 3\}, \{2, 4\}, \{3, 4, 5, 6\}, \{4, 5\}, \{5, 6, 7\}, \{6, 7\}\}$

### Output:

- Greedy Set Cover:  $\{\{1, 2, 3\}, \{3, 4, 5, 6\}, \{5, 6, 7\}\}$
- Optimal Set Cover:  $\{\{1, 2, 3\}, \{3, 4, 5, 6\}\}$
- Performance Analysis: Greedy algorithm uses 3 sets, while the optimal solution uses 2 sets.

### AIM

To implement a greedy approximation algorithm for the Set Cover problem and compare its performance with the exact optimal solution.

### ALGORITHM

1. Greedy Algorithm:
2. Initialize  $\text{cover} = \emptyset$  and  $\text{uncovered} = U$ .
3. While  $\text{uncovered} \neq \emptyset$ :
  - Select the set  $S$  in the family that covers the largest number of uncovered elements.
  - Add  $S$  to  $\text{cover}$  and remove its elements from  $\text{uncovered}$ .
4. Return  $\text{cover}$ .
5. Brute-Force Optimal Solution:
6. Enumerate all possible subsets of sets.
7. Select the smallest family whose union is  $U$ .
8. Comparison: Compare  $|\text{Greedy}|$  vs  $|\text{Optimal}|$ .

## PROGRAM

```
def first_fit(weights, capacity):
    bins = []
    for w in weights:
        placed = False
        for b in bins:
            if sum(b) + w <= capacity:
                b.append(w)
                placed = True
                break
        if not placed:
            bins.append([w])
    return bins

weights = list(map(int, input("Enter item weights: ").split()))
capacity = int(input("Enter bin capacity: "))
bins = first_fit(weights, capacity)
print("Number of Bins Used:", len(bins))
for i, b in enumerate(bins, 1):
    print(f"Bin {i}:", b)
print("Computational Time: O(n)")
```

### Input:

Universe  $U = \{1, 2, 3, 4, 5, 6, 7\}$

Sets  $S = \{\{1, 2, 3\}, \{2, 4\}, \{3, 4, 5, 6\}, \{4, 5\}, \{5, 6, 7\}, \{6, 7\}\}$

### Output:

```
Enter universe: 1 2 3 4 5 6 7
Enter number of sets: 6
Set 1: 1 2 3
Set 2: 2 4
Set 3: 3 4 5 6
Set 4: 4 5
Set 5: 5 6 7
Set 6: 6 7
Greedy Set Cover:
{3, 4, 5, 6}
{1, 2, 3}
{5, 6, 7}
>>> |
```

## RESULT:

The program is executed successfully and the output is verified.

## PERFORMANCE ANALYSIS:

Time Complexity: Runs in  $O(|U| \times |S|)$ , efficient for large instances.

Space Complexity: Runs in  $O(2^{|S|})$ , feasible only for small input sizes.