

## 2.6 BINARY SEARCH FOR PEAK ELEMENT.

### Question:

A peak element is an element that is strictly greater than its neighbors. Given a 0-indexed integer array `nums`, find a peak element, and return its index. If the array contains multiple peaks, return the index to any of the peaks. You may imagine that `nums[-1] = nums[n] = -∞`. In other words, an element is always considered to be strictly greater than a neighbor that is outside the array. You must write an algorithm that runs in  $O(\log n)$  time.

### AIM

To design an algorithm that finds the index of a peak element in an array using binary search in  $O(\log n)$  time.

### ALGORITHM

1. Let `low = 0`, `high = n-1`.
2. While `low < high`:
  - Compute `mid = (low + high) // 2`.
  - If `nums[mid] > nums[mid + 1]`:
  - Then the peak lies in the left half (including mid), so set `high = mid`.

Else:

The peak lies in the right half, so set `low = mid + 1`.

3. When the loop ends, `low == high` and that index is a peak.

## PROGRAM

```
def find_peak_element(nums):
    left, right = 0, len(nums) - 1
    while left < right:
        mid = (left + right) // 2
        if nums[mid] > nums[mid + 1]:
            right = mid
        else:
            left = mid + 1
    return left

def run_peak_element():
    nums = list(map(int, input("Enter array elements: ").split()))
    index = find_peak_element(nums)
    print("Peak element index:", index)
run_peak_element()
```

Input:

[1, 2, 3, 1]

Output:

```
Enter array elements: 1 2 3 1
Peak element index: 2
>>> |
```

## RESULT:

Thus the program is successfully executed and the output is verified.

## PERFORMANCE ANALYSIS:

- Time Complexity:
  - Each step reduces the search space by half  $\rightarrow O(\log n)$ .
- Space Complexity:
  - $O(1)$  (binary search done in-place, no extra memory).
  - Best Case: Found in first comparison  $\rightarrow O(1)$ .
  - Worst Case:  $O(\log n)$  comparisons until narrowing down to one element.