

Outline



- Worksharing Constructs
 - ✓ Loop worksharing construct
 - Ordered and schedule clauses
 - ✓ Reduction
- Data Sharing
 - ✓ Private
 - ✓ Firstprivate and lastptivate
 - √ Threadprivate
- Data Copying
- Other Constructs
 - ✓ Single worksharing construct
 - ✓ Sections worksharing construct
 - ✓ Master construct

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(Lecture 2) Loop Worksharing Construct



• The loop worksharing construct splits up loop iterations among the threads in a team.

```
#pragma omp parallel
{
    #pragma omp for
        for(i=0; i<N; i++){
            do_something(i);
        }
}</pre>
```

```
#pragma omp parallel
{
    #pragma omp for
        for(i=0; i<N; i++) {
            c[i] = a[i]+b[i];
        }
}</pre>
```

Sequential code

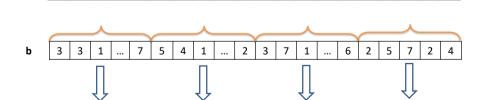
```
for(i=0; i<N; i++) {
   c[i] = a[i]+b[i];
}</pre>
```

Loop Worksharing Construct (Example)



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Vector Addition: c = a + b



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Loop Worksharing Construct



- The loop worksharing construct splits up loop iterations among the threads in a team.
- The following are 3 equivalent examples of vector addition.

```
//Example 1.
#pragma omp parallel
  int id, i, Nthrds, istart, iend;
   id = omp get thread num();
   Nthrds = omp_get_num_threads();
   istart = id * N / Nthrds;
   iend = (id+1) * N / Nthrds;
   if (id == Nthrds-1) iend = N;
   for(i=istart ; i<iend ;i++)</pre>
       c[i] = a[i] + b[i];
```

- Manually assign the loop to each thread, which is complicated.
- The same work can be done succinctly by using "parallel for".

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Loop Worksharing Construct (Cont'd) WESTERN AUSTRALIA



- //Example 2. #pragma omp parallel #pragma omp for for(i=0;i<N;i++)</pre> c[i] = a[i] + b[i];
- //Example 3. #pragma omp parallel for for(i=0;i<N; i++)</pre> c[i] = a[i] + b[i];
- OpenMP parallel region and a worksharing for construct
- OpenMP shortcut: Put the "parallel" and the worksharing directive on the same line

Note: loop index "i" is private by default.

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Loop Worksharing Construct (Cont'd) WESTERN AUSTRALIA



Working with loops

- Find compute intensive loops
- Make the loop iterations independent. So they can safely execute in any order without loop-carried dependencies.
- Place the appropriate OpenMP directive and test

```
int i, A[MAX];
int i, j, A[MAX];
                                                       #pragma omp parallel for
j = 5;
                                                         for (i=0;i< MAX; i++)
                        Remove loop carried dependence
for (i=0;i< MAX; i++)
                                                            int j = 5 + 2* (i+1);
   j +=2;
                                                            A[i] = big(j);
   A[i] = big(j);
```

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Worksharing Constructs

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Loop Worksharing Constructs: The ordered Clause



The ordered region executes in the sequential order.

```
void test(int first, int last)
#pragma omp parallel
#pragma omp for schedule(static) ordered
   for (int i = first; i <= last; ++i)</pre>
      // Do something here.
      if (i % 2)
         #pragma omp ordered
         printf s("test() iteration %d\n", i);
int main(int argc, char *argv[])
   test(1, 8);
```

Output:

test() iteration 1 test() iteration 3 test() iteration 5 test() iteration 7

Exercise:

- Delete "#pragma omp ordered", compile and run the program multiple times.
- Do you see any difference?

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Loop Worksharing Constructs: The ordered Clause (Cont'd)



The omp ordered directive must be used as follows:

- It must appear within the extent of an "omp for" or "omp parallel for" construct containing an ordered clause.
- It applies to the statement block immediately following it. Statements in that block are executed in the same order in which iterations are executed in a sequential loop.
- An iteration of a loop must not execute the same omp ordered directive more than once.
- An iteration of a loop must not execute more than one distinct omp ordered directive.

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Loop Worksharing Constructs: The schedule Clause



- The schedule clause is sometimes used in parallel for loops.
- When the amount of computation for each iteration in the loop is not equal, simply assigning the same number of iterations to each thread will result in unbalanced computing.
- This causes some lines to execute first and some to execute later, thus causing CPU idle and affecting program performance. For example, the following code:

```
int i, j;
int a[100][100] = \{0\};
for ( i =0; i < 100; i++)
   for(j = i; j < 100; j++)
      a[i][j] = i*j;
```

- If you parallelise the outer loop, say using 4 threads, and each thread is evenly allocated 25 iterations of the loop.
- There is a 100-fold difference in the amount of computation between i = 0 and i = 99.

Loop Worksharing Constructs: The schedule Clause (Cont'd)



The schedule clause affects how loop iterations are mapped onto threads to solve these problems.

schedule(type[,size])

type

The kind of scheduling can be: static, dynamic, guided, or runtime

(Optional) Specifies the size of iterations. size must be an integer. Not valid when type is runtime

Loop Worksharing Constructs: The schedule Clause (Cont'd)



- schedule(static [,chunk])
 - Deal-out blocks of iterations of size "chunk" to each thread.
- schedule(dynamic[,chunk])
 - Each thread grabs "chunk" iterations off a queue until all iterations have been handled.
- schedule(guided[,chunk])
 - Threads dynamically grab blocks of iterations. The size of the block starts large and shrinks down to size "chunk" as the calculation proceeds.
- schedule(runtime)
 - Schedule and chunk size taken from the OMP_SCHEDULE environment variable (or the runtime library ... for OpenMP 3.0).

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Loop Worksharing Constructs: The schedule Clause



Schedule Clause	When to Use	_	
static	Pre-determined and predictable by the programmer	←	least work at run- time: scheduling done at compile-time
dynamic	Unpredictable, highly variable work per iteration	←	most work at run- time: complex scheduling logic used at run-time
guided	Special case of dynamic to reduce scheduling overhead		
		_	15

Loop Worksharing Constructs: The schedule Clause (Cont'd)



<pre>#include<stdio.ho #include<comp.h=""> #include<comp.ho #define="" #include<comp.ho="" #pragma="" %="" (int="" ++i)="" 0)<="" 1)="" 20="" 4="" 5="" ;="" <="" dynamic_chunk="" for="" i="" if((i="" int="" nguided(num_loops];="" noynamicn[num_loops];="" nstatic1[num_loops],="" nstaticn[num_loops];="" num_loops="" num_loops],="" num_threads="" omp="" omp_set_num_threads(num_threads);="" parallel="" schedule(static,="" sleep_every_n)="=" static_chunk="" td="" {=""><td><pre>#pragma omp for schedule(dynamic, 1) for (int i = 0; i < NUM_LOOPS; ++i) { if((i % SLEEP_EVERY_N) == 0) sleep(0); nDynamic1[i] = omp_get_thread_num(); } #pragma omp for schedule(dynamic, DYNAMIC_CHUNK) for (int i = 0; i < NUM_LOOPS; ++i) { if((i % SLEEP_EVERY_N) == 0) sleep(0); nDynamicN[i] = omp_get_thread_num(); } #pragma omp for schedule(guided) for (int i = 0; i < NUM_LOOPS; ++i) { if((i % SLEEP_EVERY_N) == 0) sleep(0); nGuided[i] = omp_get_thread_num(); } }</pre></td></comp.ho></stdio.ho></pre>	<pre>#pragma omp for schedule(dynamic, 1) for (int i = 0; i < NUM_LOOPS; ++i) { if((i % SLEEP_EVERY_N) == 0) sleep(0); nDynamic1[i] = omp_get_thread_num(); } #pragma omp for schedule(dynamic, DYNAMIC_CHUNK) for (int i = 0; i < NUM_LOOPS; ++i) { if((i % SLEEP_EVERY_N) == 0) sleep(0); nDynamicN[i] = omp_get_thread_num(); } #pragma omp for schedule(guided) for (int i = 0; i < NUM_LOOPS; ++i) { if((i % SLEEP_EVERY_N) == 0) sleep(0); nGuided[i] = omp_get_thread_num(); } }</pre>
<pre>sleep(0); nStaticN[i] = omp_get_thread_num();</pre>	16

Loop Worksharing Constructs: The schedule Clause (Cont'd)

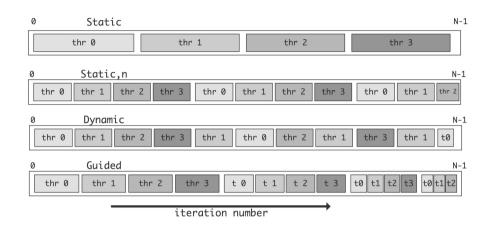


Possible Output:

static static dynamic dynamic guided 1					
0	static	static	dynamic	dynamic	guided
1	1	5	1	5	
1					
2	0	0	0	2	1
3	1	0	3	2	1
	2	0	3	2	1
1	3	0	3	2	1
2	0	0	2	2	1
3	1	1	2	3	3
0	2	1	2	3	3
1	3	1	0	3	3
2	0	1	0	3	3
3	1	1	0	3	2
0	2	2	1	0	2
1	3	2	1	0	2
2 2 2 0 0 0 1 3 3 2 1 0 0 1 1 1 1 1 1 1	0	2	1	0	3
3 3 2 1 0 0 3 3 1 1 1 3 3 1 1 2 3 3 1 1	1	2	2	0	3
0 3 3 1 1 1 1 1 1 1 1	2	2	2	0	0
1 1 3 3 1 1 1 1 2 3 3 1 1 1 1	3	3	2	1	0
2 3 3 1 1 1 1	0	3	3	1	1
	1	3	3	1	1
3 3 0 1 3	2	3	3	1	1
	3	3	0	1	3

The schedule Clause: Overview





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Outline



Worksharing Constructs

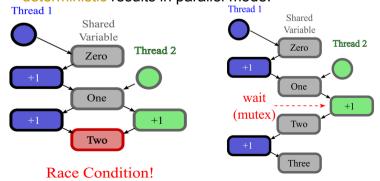
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Race Condition



- A data race occurs when two threads access the same memory without proper synchronisation.
- This can cause the program to produce nondeterministic results in parallel mode.



Race Condition

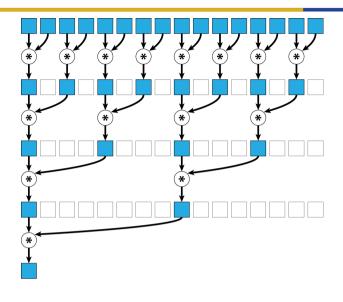


· What happen in this case?

- A shared variable (sum) is modified in every iteration.
- Resulting in race condition, because multiple threads could try
 to update the shared variable at the same time if we parallelise
 the "for" loops by simply adding #pragma omp parallel for.
- The problem can be solved by reduction clause in OpenMP.

Parallel Reduction

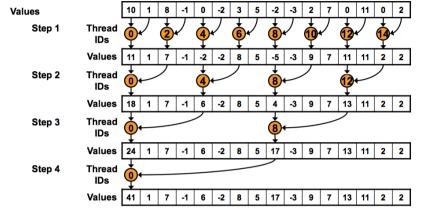




Parallel Reduction (Continued)



Many problems can be modelled as a reduction problem.
 ✓ Sum, Min, Max, ...



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Reduction in OpenMP



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- OpenMP has the special reduction clause which can express the reduction of a for loop.
- · OpenMP reduction clause:
 - ✓ reduction (op : list)
- · Implementation:
 - ✓ OpenMP creates a team of threads and then shares the iterations of the for loop between the threads.
 - ✓ Each thread has its own **local** copy of the reduction variable, and the thread modifies only the local copy of this variable.
 - ✓ Therefore, there is no data race.
 - ✓ When the threads join together, all the local copies of the reduction variable are combined to the global shared variable.

Reduction in OpenMP



Implementation:

Assume each thread has **sum**, which is a local copy of the reduction variable. The threads then perform the following computations:

- Thread 1 : sum1 = a[0] + a[1] + a[2]
- Thread 2 : sum2 = a[3] + a[4] + a[5]
- Thread 3: sum3 = a[6] + a[7] + a[8]
- sum = sum1 + sum2 + sum3

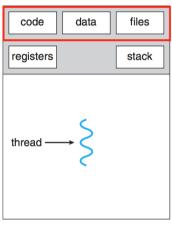
Outline

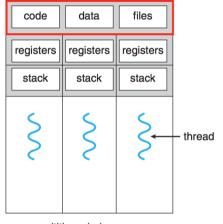


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(Lecture 1) Threads in a Process







single-threaded process

multithreaded process

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Data Sharing



Data Environment:

- An important consideration for OpenMP programming is the understanding and use of data scoping.
- As OpenMP is based upon the shared memory programming model, most variables are shared by default.
- · Global variables include:
 - ✓ File scope variables, static
- But not everything is shared...
- · Private variables include:
 - ✓ Loop index variables
 - ✓ Stack variables in subroutines called from parallel regions

Data Sharing: private Clause



Purpose:

✓ private(var) creates a new local copy of var for each thread.

Implementation:

```
int main(int argc, _TCHAR* argv[])
{
    int B;
#pragma omp parallel for private(B)
    for(int i=0; i<10;i++)
    {
        B = 100;
    }
    printf("%d\n",B);
    return 0;
}</pre>
```

Notes:

- A new object of the same type is declared once for each thread in the team
- All references to the original object are replaced with references to the new object
- Should be assumed to be uninitialized for each
 thread

May have problem here.

Data Sharing: shared Clause



Purpose:

• The **shared(list)** clause declares variables in its list to be shared among all threads in the team.

Implementation:

```
#pragma omp parallel for shared(n, a)
    for (int i = 0; i < n; i++) {
        int b = a + i;
        ...
}</pre>
```

Notes:

- n and a are shared variables.
- Shared variable exists in only one memory location and all threads can read or write to that address
- OpenMP does not put any restriction to prevent data races between shared variables. This is a responsibility of a programmer.

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Data Sharing: default Clause



Purpose:

- The default clause allows the user to specify a default scope for all variables in the lexical extent of any parallel region.
- **default(shared)** means that any variable in a parallel region will be treated as if it were specified with the **shared** clause.
- default(none) means that any variables used in a parallel region that aren't scoped with the private, shared, reduction, firstprivate, or lastprivate clause will cause a compiler error.

Notes:

• C/C++ only has default(shared) or default(none).

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Data Sharing: default(shared)



Implementation:

```
int a, b, c, n;
...
#pragma omp parallel for default(shared)
   for (int i = 0; i < n; i++) {
     // using a, b, c
}</pre>
```

- The default(shared) clause sets the data-sharing attributes of all the variables in the construct to "shared".
- a, b, c and n are shared variables.

Data Sharing: default(none)



Implementation:

```
int a = 0, b = 0;
#pragma omp parallel default(none) shared(b)
{
    b += a;
}
```

- The program above will cause compilation error.
- Because variable 'a' has not scope with any data-sharing clause.

Data Sharing: firstprivate Clause



• **firstprivate(list):** All variables in the list are initialised with the value the original object had before entering the parallel construct.

```
void test firstprivate()
    int tmp = 0;
#pragma omp for firstprivate(tmp)
                                               Each thread gets its
    for(int j = 0; j < 1000; ++j)
         tmp += j;
                                                own tmp with an
     printf("%d\n", tmp);
                                                 initial value of 0
                                               Exercise:
                tmp: 0 in OpenMP
                                                 change "firstprivate" to "private"
                3.0, unspecified in
                                                 print "tmp" within the for loop
                   OpenMP 2.5.
                                                                            34
```

Data Sharing: lastprivate Clause



• lastprivate(list): The thread that executes the sequentially last iteration or section updates the value of the objects in the list.

```
void test_firstprivate_lastprivate()
{
   int tmp = 0;
#pragma omp parallel for firstprivate(tmp) lastprivate(tmp)
   for (int j = 0; j < 1000; ++j)
        tmp += j;
   printf("%d\n", tmp);
}</pre>

Each thread gets its
   own tmp with an
   initial value of 0

tmp is defined as its value at
   the "last sequential" iteration
        (i.e., for j=999)
```

Data Sharing: threadprivate



Purpose:

- The threadprivate directive specifies that variables are replicated, with each thread having its own copy.
- Can be used to make global file scope variables and persistent to a thread through the execution of multiple parallel regions.

Notes:

- The directive must appear after the declaration of listed variables. Each thread then gets its own copy of the variable, so data written by one thread is not visible to other threads.
- On first entry to a parallel region, data in threadprivate variables should be assumed undefined, unless a COPYIN clause is specified in the PARALLEL directive
- threadprivate variables differ from private variables because:
 - with private global variables are masked.
 - threadprivate preserves global scope within each thread

Data Sharing: threadprivate Example



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• Use threadprivate to create a counter for each thread:

```
int counter = 0;
#pragma omp threadprivate(counter)
int increment_counter()
{
    counter++;
    Return counter;
}
```

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Data Copying: copyin Clause



Purpose:

 copyin(list) allows to copy master thread's threadprivate variables to corresponding threadprivate variables of other threads.

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Data Copying: copyin Clause (Cont'd)



Notes:

- List contains the names of variables to copy.
- The master thread variable is used as the copy source. The team threads are initialised with its value upon entry into the parallel construct.

Data Copying: copyprivate Clause



- Purpose:
 - The copyprivate clause can be used to broadcast values acquired by a single thread directly to all instances of the private variables in the other threads.

Usage:

- The typical usage is to have one thread read or initialise private data that is subsequently used by the other threads as well.
- After the "single" construct has ended, but before the threads have left the associated barrier, the values of variables specified in the associated list are copied to the other threads.
- Do not use copyprivate in combination with the nowait clause!

Data Copying: copyprivate Example



Implementation:

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Master Construct



- The **master** construct denotes a structured block that is only executed by the master thread.
- The other threads just skip it (no synchronisation is implied).

Single Worksharing Construct



- The single construct denotes a block of code that is executed by only one thread (not necessarily the master thread).
- A barrier is **implied** at the end of the single block (can remove the barrier with a **nowait** clause).

```
#pragma omp parallel
{
    do_many_things();

#pragma omp single
    {
       exchange_boundaries();
    }

    do_many_other_things();
}
```

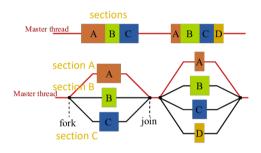
Master vs. Single Constructs





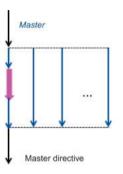


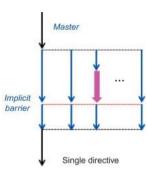
- The sections construct is a non-iterative worksharing construct that contains a set of structured blocks that are to be distributed among and executed by the threads in a team.
- Each structured block is executed once by one of the threads in the team in the context of its implicit task.
- There is an implicit barrier at the end of a sections construct unless a nowait clause is specified.



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Sections Worksharing Construct (Cont'd) WESTERN AUSTRALIA



• The syntax of the sections construct is as follows:

```
#pragma omp parallel
#pragma omp sections
         #pragma omp section
            X calculation();
         #pragma omp section
            Y_calculation();
         #pragma omp section
            Z_calculation();
```

Barriers and Worksharing Constructs



- If the parallel region is worksharing construct, barriers will be automatically added after it;
- Worksharing construct includes for, sections, and single; pay special attention to that the master construct is not;
- nowait can remove barriers.

References



- Readings
 - Ordered Clause
 - OpenMP Data Sharing
 - <u>Difference of Directive, Construct and Clause</u>