

Open Queueing Networks

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Queueing Networks

- ▶ A queueing network is a connected directed graph, whose nodes represent servers (each with their own queue), and edges between nodes represent channels through which jobs can be routed.

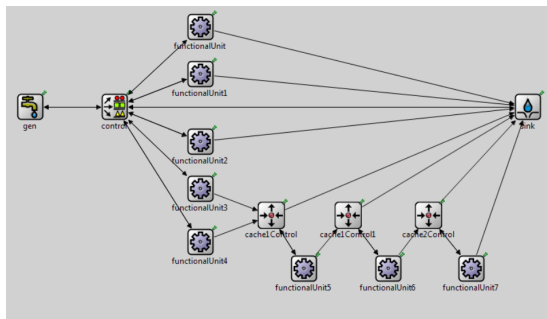


Figure: Intel Ivy Bridge queue model [2]

Open Networks

- ▶ A network is said to be open if there is at least one edge for an 'external' arrival, and from every node there is a path for an 'external' departure.

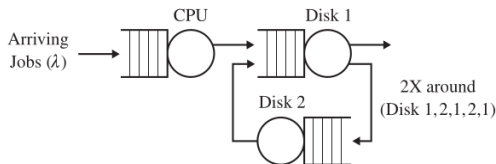


Figure: An open queueing network

Jackson Network

► A general form of open network

- Unbounded queues, with FCFS service order
- External arrivals $\sim \text{Poisson}(r_i)$
- Service rates $\sim \text{Exp}(\mu_i)$
- Probabilistic routing - P_{ij}
- States: (n_1, n_2, \dots, n_k)

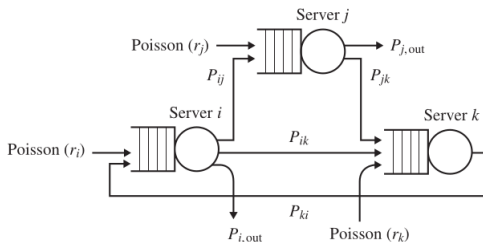
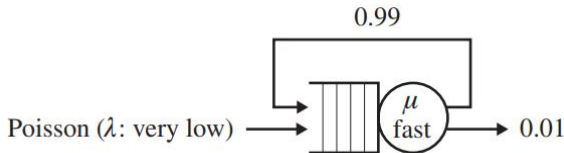


Figure: A Jackson network

Poisson Arrivals?

- ▶ $\lambda_i = r_i + \sum_j \lambda_j P_{ji}$
- ▶ Consider a single server network with properties:
 - Low arrival rate
 - High service rate
 - High P_{ii}



Violates independent increment property!

Balance Equations

- Rate of jobs leaving the state = Rate of jobs entering the state.

$$\begin{aligned}
 & \pi_{n_1, n_2, \dots, n_k} \cdot \left[\sum_{i=1}^k r_i + \sum_{i=1}^k \mu_i (1 - P_{ii}) \right] = \\
 & \underbrace{\sum_{i=1}^k \pi_{n_1, \dots, n_i-1, \dots, n_k} \cdot r_i}_{\text{outside arrival}} + \underbrace{\sum_{i=1}^k \pi_{n_1, \dots, n_i+1, \dots, n_k} \cdot \mu_i P_{i, \text{out}}}_{\text{departure to outside}} \\
 & + \underbrace{\sum_{i=1}^k \sum_{j \neq i} \pi_{n_1, \dots, n_i-1, \dots, n_j+1, \dots, n_k} \cdot \mu_j P_{ji}}_{\text{internal transition from server } j \text{ to server } i, j \neq i}
 \end{aligned}$$

Balance Equations (Contd.)

- ▶ In general, very hard to solve.
- ▶ Local Balance Approach
 - Break both sides of the equation into $k + 1$ matching components and equate each component.
 - Define A = Rate of leaving state due to an external arrival.
 - Define A' = Rate of entering state due to an external departure.
 - Define B_i = Rate of leaving state due to a departure from server i .
 - Define B'_i = Rate of entering state due to an arrival at server i .

Balance Equations (Contd.)

- ▶ Solving $A = A'$

$$\sum_{i=1}^k \pi_{n_1, \dots, n_i, \dots, n_k} \cdot r_i = \sum_{i=1}^k \pi_{n_1, \dots, n_i+1, \dots, n_k} \cdot \mu_i P_{i, \text{out}}$$

- ▶ Can we make a guess for π_{n_1, \dots, n_k} such that it satisfies the above equality?
- ▶ Observe that $\pi_{n_1, \dots, n_i, \dots, n_k}$ term in A and the $\pi_{n_1, \dots, n_i+1, \dots, n_k}$ term in A' only differ in n_i .
- ▶ Let us assume that

$$\pi_{n_1, \dots, n_i, \dots, n_k} \cdot c_i = \pi_{n_1, \dots, n_i+1, \dots, n_k}$$

- ▶ Substituting, we get

$$\sum_{i=1}^k r_i = \sum_{i=1}^k (c_i \cdot \mu_i) P_{i, \text{out}}$$

Balance Equations (Contd.)

- ▶ If we further assume $c_i \cdot \mu_i = \lambda_i$, we get

$$\sum_{i=1}^k r_i = \sum_{i=1}^k \lambda_i P_{i,\text{out}}$$

- ▶ By our assumption, $c_i = \frac{\lambda_i}{\mu_i} = \rho_i$. Substituting this we get

$$\pi_{n_1}, \dots, \pi_{n_i}, \dots, \pi_{n_k} \cdot \rho_i = \pi_{n_1}, \dots, \pi_{n_i+1}, \dots, \pi_{n_k}$$

- ▶ Repeating the argument yields the following equation

$$\pi_{n_1}, \dots, \pi_{n_k} = C \rho_1^{n_1} \dots \rho_k^{n_k}$$

Balance Equations (Contd.)

- We still need to check if our assumption yields a valid solution to $B_i = B'_i$

$$B_i = B'_i$$

$$C \rho_1^{n_1} \rho_2^{n_2} \cdots \rho_k^{n_k} \mu_i (1 - P_{ii}) = \sum_{j \neq i} C \rho_1^{n_1} \rho_2^{n_2} \cdots \rho_k^{n_k} \left(\frac{\rho_j}{\rho_i} \right) \mu_j P_{ji}$$

$$+ C \rho_1^{n_1} \rho_2^{n_2} \cdots \rho_k^{n_k} \left(\frac{1}{\rho_i} \right) r_i$$

$$\mu_i (1 - P_{ii}) = \sum_{j \neq i} \frac{\rho_j \mu_j P_{ji}}{\rho_i} + \frac{r_i}{\rho_i}$$

$$\rho_i \mu_i (1 - P_{ii}) = \sum_{j \neq i} \rho_j \mu_j P_{ji} + r_i$$

$$\lambda_i (1 - P_{ii}) = \sum_{j \neq i} \lambda_j P_{ji} + r_i$$

Balance Equations (Contd.)

- Lastly, we need to find the normalizing constant C .

$$\sum_{n_1, \dots, n_k} \pi_{n_1, \dots, n_k} = 1$$

$$C \sum_{n_1, \dots, n_k} \rho_1^{n_1} \cdots \rho_k^{n_k} = 1$$

$$C \left(\sum_{n_1} \rho_1^{n_1} \right) \left(\sum_{n_2} \rho_2^{n_2} \right) \cdots \left(\sum_{n_k} \rho_k^{n_k} \right) = 1$$

$$C \left(\frac{1}{1 - \rho_1} \right) \left(\frac{1}{1 - \rho_2} \right) \cdots \left(\frac{1}{1 - \rho_k} \right) = 1$$

$$C = (1 - \rho_1)(1 - \rho_2) \cdots (1 - \rho_k).$$

$$\pi_{n_1, \dots, n_k} = \rho_1^{n_1} (1 - \rho_1) \rho_2^{n_2} (1 - \rho_2) \cdots \rho_k^{n_k} (1 - \rho_k)$$

Balance Equations (Contd.)

- What about the jobs at a given server?

$$\begin{aligned}P\{n_1 \text{ jobs at server 1}\} &= \sum_{n_2, \dots, n_k} \pi_{n_1, \dots, n_k} \\&= \sum_{n_2, \dots, n_k} \rho_1^{n_1} (1 - \rho_1) \rho_2^{n_2} (1 - \rho_2) \cdots \rho_k^{n_k} (1 - \rho_k) \\&= \rho_1^{n_1} (1 - \rho_1)\end{aligned}$$

$$P\{n_i \text{ jobs at server } i\} = \rho_i^{n_i} (1 - \rho_i)$$

- Thus, all servers behave like M/M/1 queues in terms of their stationary queue length distributions, even though their arrival processes are not necessarily Poisson.

Summary

$$\lambda_i = r_i + \sum_j \lambda_j P_{ji}$$

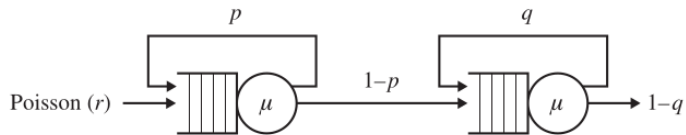
$$\rho_i = \frac{\lambda_i}{\mu_i}$$

$$\pi_{n_1, \dots, n_k} = \rho_1^{n_1} (1 - \rho_1) \rho_2^{n_2} (1 - \rho_2) \cdots \rho_k^{n_k} (1 - \rho_k)$$

$$P\{n_i \text{ jobs at server } i\} = \rho_i^{n_i} (1 - \rho_i)$$

$$E\{N_i\} = \frac{\rho_i}{1 - \rho_i}$$

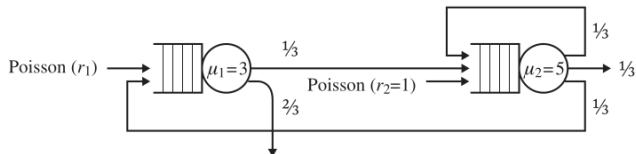
Case 1



Case 1

- ▶ $\lambda_1 = r + p\lambda_1$
 $\lambda_2 = (1 - p)\lambda_1 + q\lambda_2$
- ▶ $\lambda_1 = \frac{r}{1-p}, \lambda_2 = \frac{r}{1-q}$
- ▶ $\rho_1 = \frac{r}{\mu(1-p)}, \rho_2 = \frac{r}{\mu(1-q)}$
- ▶ $E[N_1] = \frac{r}{\mu(1-p)-r}, E[N_2] = \frac{r}{\mu(1-q)-r}$
- ▶ Let $r = 0.5, \mu = 3, p = 0.1, q = 0.2$.
- ▶ $\lambda_1 = \frac{5}{9}, \lambda_2 = \frac{5}{8}$
- ▶ $\rho_1 = \frac{5}{27}, \rho_2 = \frac{5}{24}$
- ▶ $E[N_1] = 0.23, E[N_2] = 0.26$
- ▶ $E[T_1] = 0.41, E[T_2] = 0.41$

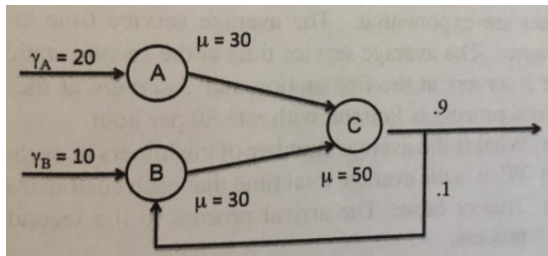
Case 2



Case 2

- ▶ $\lambda_1 = r_1 + \frac{1}{3}\lambda_2$
 $\lambda_2 = 1 + \frac{1}{3}\lambda_1 + \frac{1}{3}\lambda_2$
- ▶ $\lambda_1 = \frac{6r_1+3}{5}, \lambda_2 = \frac{3r_1+9}{5}$
- ▶ $\rho_1 = \frac{2r_1+1}{5}, \rho_2 = \frac{3r_1+9}{25}$
- ▶ For $\rho_1 < 1$, $r_1 < 2$. Let $r_1 = 1.5$.
- ▶ $\lambda_1 = 2.4, \lambda_2 = 2.7$
- ▶ $\rho_1 = 0.8, \rho_2 = 0.34$
- ▶ $E[N_1] = 4, E[N_2] = 1.17$
- ▶ $E[T_1] = 1.67, E[T_2] = 0.43$

Case 3



Case 3

- ▶ $\lambda_1 = 20$
 $\lambda_2 = 10 + 0.1\lambda_3$
 $\lambda_3 = \lambda_1 + \lambda_2$
- ▶ $\lambda_1 = 20, \lambda_2 = 13, \lambda_3 = 33.33$
- ▶ $\rho_1 = 0.67, \rho_2 = \frac{13}{30}, \rho_3 = 0.67$
- ▶ $E[N_1] = 2, E[N_2] = \frac{13}{17}, E[N_3] = 2$
- ▶ $E[T_1] = 0.1, E[T_2] = \frac{1}{17}, E[T_3] = 0.06$

References

- [1] M. Harchol-Balter. *Performance Modeling and Design of Computer Systems: Queueing Theory in Action*. Cambridge University Press, 2013.
- [2] Damián Roca Marí. *High Level Queueing Architecture Model for High-end Processors*. 2014.
- [3] Isi Mitrani. *Queueing networks*. Cambridge University Press, 1997.