# Quantum Channels: Transition from Classical to Quantum Channels

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Accessible Information

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Schumacher's Quantum Noiseless Channel Coding Theorem

Classical Capacity of Quantum Channels

Private Capacity of Quantum Channels

Quantum Capacity of Quantum Channels

Super Additivity of Quantum Channels

**Examples of Quantum Channels** 

Author 1 3

Author 1

Author 2

Author 2

#### Private Capacity of a Wiretap Channel

The private capacity  $P(\mathcal{N})$  of a classical wiretap channel  $\mathcal{N} = p_{Y,Z|X}$  is defined as follows.

$$P(\mathcal{N}) = \max_{p_{U,X}(u,x)} [I(U;Y) - I(U;Z)]$$

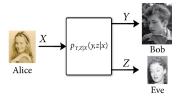


Figure: Classical wiretap channel.

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## Properties: Private Capacity of a Wiretap Channel

Non-negativity

$$P(\mathcal{N}) \geq 0$$

Additivity

$$P(\mathcal{N} \otimes \mathcal{M}) = P(\mathcal{N}) + P(\mathcal{M})$$

Equivalence of asymptotic and single use capacity

$$\lim_{n\to\infty}\frac{1}{n}P(\mathcal{N}^{\otimes n})=P(\mathcal{N})$$

#### Private Capacity of a Quantum Channel

Consider a classical-quantum state of the form given below.

$$\rho_{XA'} = \sum_{x} \rho_{X}(x) |x\rangle \langle x|_{X} \otimes \rho_{A'}^{x}$$

The private capacity  $P(\mathcal{N})$  of a quantum channel  $\mathcal{N}$  is defined as follows.

$$P(\mathcal{N}) = \max_{\rho_{XA'}} [I(X; B)_{\rho} - I(X; E)_{\rho}]$$

$$\begin{array}{c|c}
\rho_{A} & \sigma_{B} = \mathcal{N}(\rho_{A}) \\
|0\rangle & \sigma_{E} = E(\rho_{A})
\end{array}$$

Figure: Private communication through a quantum channel.

## Properties: Private Capacity of a Quantum Channel

Non-negativity

$$P(\mathcal{N}) \geq 0$$

Bounded by quantum capacity

$$P(\mathcal{N}) \leq Q(\mathcal{N})$$

Non-additivity

$$P(\mathcal{N} \otimes \mathcal{M}) \neq P(\mathcal{N}) + P(\mathcal{M})$$

## Quantum Capacity of a Quantum Channel

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