CS1.305: Introduction to Algorithms Engineering t-Spanner Construction

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1 3-Spanner Construction

1.1 Algorithm

```
Algorithm 1: 3-Spanner Algorithm
 1 Initialization: E_S = R = \phi
 \mathbf{2} \ \mathbf{for} \ v \in V \ \mathbf{do}
       add v to \mathcal{R} with probability \frac{1}{\sqrt{n}}
 4 end
 5 for v \in V - \mathcal{R} do
        if v is not adjacent to x \in \mathcal{R} then
            add all edges incident on v to E_S
        else
             N(v, \mathcal{R}) \leftarrow \text{nearest neighbor in } \mathcal{R}
            add (v, N(v, \mathcal{R})) and all lighter edges (v, *) to E_S
10
11
12 end
13 for v belonging to a cluster do
        for each adjacent cluster c do
            add the least weight edge in E(v,c) to E_S
16
17 end
```

1.2 Results

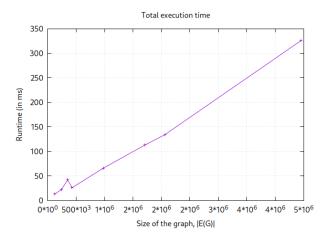


Figure 1: Total execution time $\,$

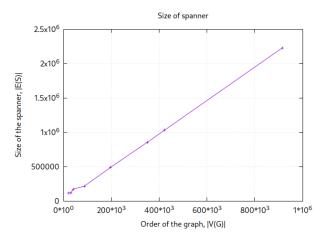


Figure 2: Spanner size

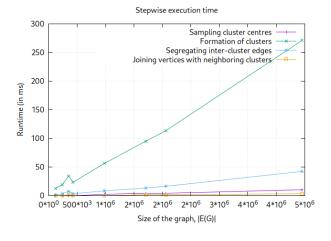


Figure 3: Stepwise execution time $\,$

2 (2k - 1)-Spanner Construction

2.1 Algorithm

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Algorithm 2: (2k - 1)-Spanner Algorithm: Cluster Formation
 1 Initialization: E' = E, V' := V(E' \cup \mathcal{E}_{i-1}), E_S = \mathcal{E}_0 = \phi, C_0 = \{\{v\} \mid v \in V\}
 2 for i \in \{1, 2, \dots, num\_iterations\} do
        sample clusters C_i from C_{i-1} independently with probability n^{-\frac{1}{k}}
       for v \in V' \setminus \cup C_i do
 4
           if v is not adjacent to c \in C_i then
 5
                for c' \in \mathcal{C}_{i-1} do
 6
                    add the least weight edge E'(v,c') to E_S
                    remove the edges E'(v,c') from E'
                end
           else
10
                c, e_v \leftarrow \text{nearest cluster in } \mathcal{C}_i \text{ and leading edge}
11
                add the edge e_v to E_s and \mathcal{E}_i
12
                remove the edges E'(v,c) from E'
13
                for c' \in \mathcal{C}_{i-1} do
14
                    if c' is reachable with edge lighter than e_v then
                        add the least weight edge E'(v,c') to E_S
16
                        remove the edges E'(v,c') from E'
17
                    end
                end
19
20
           end
21
       remove all intra-cluster edges of C_i from E'
22
23 end
24 for v \in V', c \in \mathcal{C}_{k-1} do
       add the least weight edge in E'(v,c) to E_S
       remove the edges E'(v,c) in E'
26
27 end
   Algorithm 3: (2k - 1)-Spanner Algorithm: Vertex Cluster Joining
 {f 1} execute cluster formation for k-1 iterations
 2 for v \in V', c \in \mathcal{C}_{k-1} do
       add the least weight edge in E'(v,c) to E_S
       remove the edges E'(v,c) in E'
 5 end
```

Algorithm 4: (2k – 1)-Spanner Algorithm: Cluster Cluster Joining

```
1 execute cluster formation for \lfloor \frac{k}{2} \rfloor iterations
2 if k is odd then
3 | for c, c' \in \mathcal{C}_{\lfloor \frac{k}{2} \rfloor} do
4 | add the least weight edge in E'(c, c') to E_S
5 | end
6 else
7 | for c \in \mathcal{C}_{\lfloor \frac{k}{2} \rfloor}, c' \in \mathcal{C}_{\lfloor \frac{k}{2} - 1 \rfloor} do
8 | add the least weight edge in E'(c, c') to E_S
9 | end
10 end
```

2.2 Results

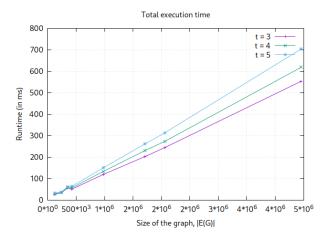


Figure 4: Total execution time

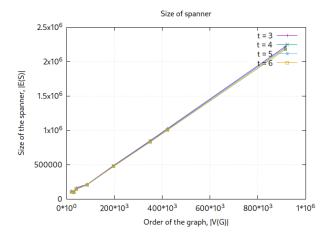


Figure 5: Spanner size

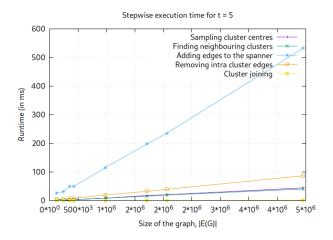


Figure 6: Stepwise execution time

Bibliography

- 1. Baswana, Surender Sen, Sandeep. (2003). A Simple and Linear Time Randomized Algorithm for Computing Sparse Spanners in Weighted Graphs
- 2. Reyan Ahmed, Greg Bodwin, Faryad Darabi Sahneh, Keaton Hamm, Mohammad Javad Latifi Jebelli, Stephen Kobourov, Richard Spence (2019). Graph Spanners: A Tutorial Review.