

Data Structures

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Arrays, Records and Pointers



- Data Structure are Classified as either linear or nonlinear.
 - A data structure is said to be linear if its elements form a sequence, or in other words, a linear list.
 - ☐ There are two basic ways of representing such linear structures in memory.
- One way is to have the linear relationship between the elements represented by means of sequential memory locations. These linear structures are called Arrays[1].





- The other way is to have the linear relationship between the elements represented by means of pointers or links.
- These linear structures are called linked lists.



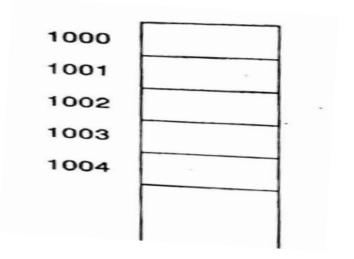
- A linear array is a list of finite number n of homogeneous data elements (i. e. data elements of the same type) such that:
 - The elements of the array are referenced respectively by an index set consisting of n consecutive numbers.
 - The elements of the array are stored respectively in successive memory locations [1].
- The number of n elements is called the length or size of the array.



Representation of Linear Array in Memory:

❖ Let LA be a linear array in the memory of computer (A memory of a computer is simply a sequence of addressed locations).

LA(LA[k]) = address of the element LA[k] of the array LA.





- **Operations on Linear Arrays:**
 - **Traversing:**
 - **Ex.** Count number of elements on the array.
 - **A** Inserting and Deleting:
 - **Sorting:**
 - **Searching:**



♦ Multidimensional Arrays:

- The linear array discussed so far are also called one-dimensional arrays, each elements in the array is referenced by a single subscript [1].
- Most of the programming language allow two-dimensional and three-dimensional arrays, i. e. arrays where elements are referenced, respectively, by two and three subscripts.



Two-dimensional Arrays:

- A two dimensional M×N array A is a collection of M.N data elements such that each element is specified by pair of integers called subscript with a property that: $1 \le j \le N$, $1 \le k \le M$ [1].
- lacktriangle Denoted as A[j, k] or $A_{j,k}$
- **♦** Example: A[3,4]

```
1 2 3 4
1 [A[1,1] A[1,2] A[1,3] A[1,4]
2 [A[2,1] A[2,2] A[2,3] A[2,4]
3 [A[3,1] A[3,2] A[3,3] A[3,4]
```



♦ Two-dimensional Arrays: 2-D A

 \Leftrightarrow Example: A[3,3]

j/k	0	1	2
0	44	55	88
1	22	12	87
2	123	76	41

2-D Array Rep. Col. Major
44
22
123
55
12
76
88
87
41

2-D Array Rep. Row Major
44
55
88
22
12
87
123
76
41



Two-dimensional Arrays:

Example: A[3,3], now delete A[2,2]

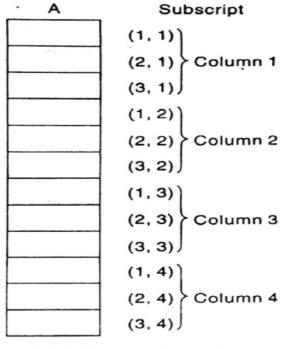
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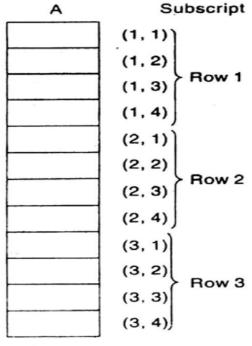
2-D Array
Rep. Row
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- Representation of two-dimensional Arrays in memory [1]:
 - ☐ Column major order:
 - ☐ Row Major order:



(a) Column-major order



(b) Row-major order



- **Representation of two-dimensional Arrays in memory [1]:**
 - ☐ Direct formula to find the memory address of an element:

LOC(LA[K]) =
$$Base(LA) + w(K - 1)$$

(Column-major order) $LOC(A[J, K]) = Base(A) + w[M(K - 1) + (J - 1)]$

(Row-major order)
$$LOC(A[J, K]) = Base(A) + w[N(J-1) + (K-1)]$$

❖ Here, w is the number of words per memory cell for the LA.

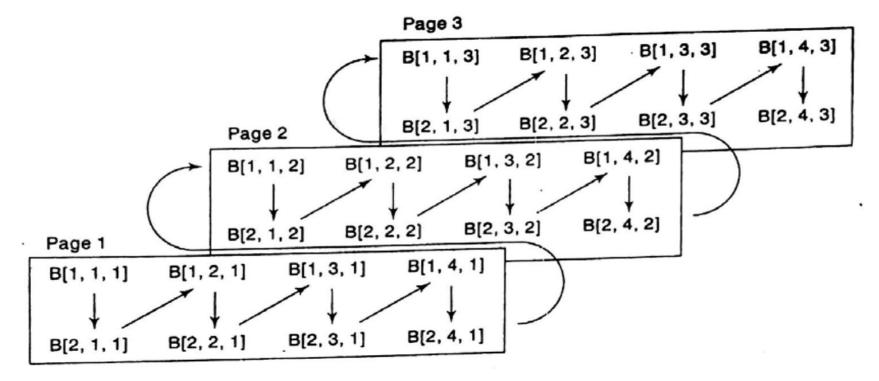


❖ General Multidimensional Arrays [1]:

- An n dimensional $m_1 \times m_2 \times m_3 \dots \times m_n$ array B is a collection of m_1 . $m_2 \cdot m_3 \cdot \dots \cdot m_n$ data elements in which each element is specified by a list of n integers such as $k_1, k_2, k_3, \dots k_n$ called subscripts. With the property that $1 \le k_1 \le m_1$, $1 \le k_2 \le m_2 \cdot \dots \cdot 1 \le k_n \le m_n$
- \triangleright These elements of B are denoted as: B[k_1 , k_2 , k_3 , ... k_n].
- \triangleright Ex.: **B**[2, 4, 3].

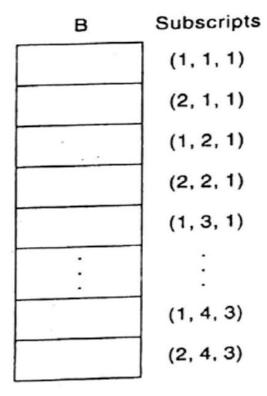


- **♦** General Multidimensional Arrays [1]:
 - \Box Ex.: **B**[2, 4, 3].

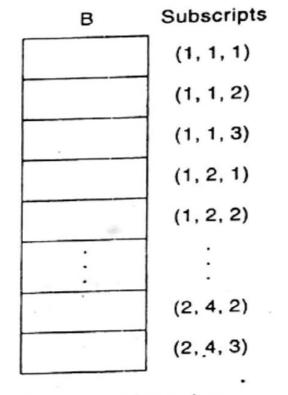




 \Box Ex.: **B**[2, 4, 3].



(a) Column-major order



(b) Row-major order



- Formula for finding the memory address of a General Multi-dimensional Array [1]:
- **Suppose** C is the multidimensional array:
 - * The length L_i of dimension i of C is the number of elements in the index set, and L_i can be calculated as:

 L_i = upper bound – lower bound + 1



- ❖ Formula for finding the memory address of a General Multi-dimensional Array:
 - For a given subscript k_i , the effective index E_i of L_i is the number of indices preceding k_i in the index set, and E_i can be calculated from:

$$E_i = K_i - lower bound$$



Formula for finding the memory address of a General Multi-dimensional Array:

$$Base(C) + w[(((... (E_N L_{N-1} + E_{N-1}) L_{N-2}) + ... + E_3) L_2 + E_2) L_1 + E_1]$$

OR

Base(C) +
$$w[(...(E_1L_2 + E_2)L_3 + E_3)L_4 + ... + E_{N-1})L_N + E_N]$$



- ❖ Suppose a three-dimensional array MAZE is declared using MAZE(2:8, -4:1, 6:10)
- ❖ Then the lengths of the three dimensions of the MAZE are, respectively, $L_1 = 8 2 + 1 = 7$ $L_2 = 1 (-4) + 1 = 6$ $L_3 = 10 6 + 1 = 5$
- * Accordingly, MAZE contains L_1 . L_2 . $L_3=7.6.5=210$ elements.
- ❖ Suppose the programming language stores MAZE in the memory in row-major order, and suppose Base(MAZE)=200 and there are w=4 words per memory cell.



- * Suppose the programming language stores MAZE in the memory in rowmajor order, and suppose Base(MAZE)=200 and there are w=4 words per memory cell.
- * The address of element of the MAZE for example, MAZE[5, -1, 8] is obtained as follows:
- * The address indices of the subscripts are, respectively,

$$E_1 = 5 - 2 = 3$$

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 $E_2 = -1 - (-4) = 3$ $E_3 = 8 - 6 = 2$

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Using row major order, we have:

$$E_1L_2 = 3.6 = 18$$

 $E_1L_2 + E_2 = 18 + 3 = 21$
 $(E_1L_2 + E_2) L_3 = 21.5 = 105$
 $(E_1L_2 + E_2) L_3 + E_3 = 105 + 2 = 107$

Therefore, LOC(MAZE(5, -1, 8))= 200+4(107)=200+428=628.

Base(C) +
$$w[(...(E_1L_2 + E_2)L_3 + E_3)L_4 + ... + E_{N-1})L_N + E_N]$$

References



1. Seymour Lipschutz, "Data Structures", Schaum's Series McGraw Hill edition 2013.