

## A) Binary search algorithm

→ Searching algorithm used to search in a sorted array by dividing the search interval in half. The time complexity is reduced to  $O(\log n)$ .

### Basic steps:

- Sort the array in ascending order.
- Set the low index to first element of the array and high index to the last element.
- Set the middle index to the average of the low and high indices.
- if element is at middle index, then return the middle index.
- if target element is greater than middle element, set low index to  $mid + 1$ .
- if target element is lesser than middle element, set high index to  $mid - 1$ .
- Repeat steps 3-6, until the element is found or it is clear that element is not present in the array.

### 1) Iterative method

```
binarysearch(arr, x, l, h) {  
    repeat till l = h {  
        mid = (l + h) / 2  
        if (x == arr[mid])  
            return mid  
        else if (x > arr[mid])  
            l = mid + 1  
        else h = mid - 1  
    }  
}
```

2) recursive method

```

binarySearch(arr, n, L, h) {
    if low > high
        return 0
    else
        mid = (l + h) / 2
        if (x == arr[mid])
            return mid
        else if (x > arr[mid])
            return binarySearch(arr, n, mid + 1, h)
        else
            return binarySearch(arr, n, l, mid - 1)
}

```

Time complexity:  $O(\log n)$

recurrence relation:  $T(n) = T(n/2) + 1$

where,  $T(n)$  is time required for binary search in an array of size  $n$ .

3) Merge sort algorithm

→ Merge sort is sorting algorithm that works by dividing an array into smaller subarrays, sorting each subarray and then merging the sub-sorted subarrays back together to form the final sorted array.

→ In simple terms, divide the array into two halves, sort each half and then merge sorted halves back together, until the <sup>whole</sup> array is sorted.



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Complexities

Time complexity :

Best case :  $O(n \log n)$ Average case :  $O(n \log n)$ Worst case :  $O(n \log n)$ 

Space complexity :

✓ Stable,  $O(n)$ Algorithm:

```

mergesort(arr, s, e) {
    if (s >= e)
        return

```

arr = array

s = start of array

e = End of array

```

    mid = (s + e) / 2

```

```

    mergesort(arr, s, mid)

```

```

    mergesort(arr, mid + 1, e)

```

```

    merge(arr, s, e)

```

```

}

```

```

merge(arr, s, e) {

```

```

    mid = (s + e) / 2

```

```

    l1 = mid - s + 1

```

```

    l2 = e - mid

```

```

    int a1 as array of length l1

```

```

    int a2 as array of length l2

```

```

    // copy elements from main array to left and
    right array

```

```

    k = s; // index of main array

```

```

    for (i = 0 to l1) {

```

Spiral

```

a1[i] ← arr[k++];
}
k = mid + 1;
for (i = 0 to l2) {
    a2[i] ← arr[k++];
}
// merge 2 sorted arrays
index1 = 0
index2 = 0
k = r
while (index1 < l1 && index2 < l2) {
    if (a1[index1] < a2[index2]) {
        arr[k++] = a1[index1++];
    } else {
        arr[k++] = a2[index2++];
    }
}
while (index1 < l1) { // for remaining of a1
    arr[k++] = a1[index1++];
}
while (index2 < l2) {
    arr[k++] = a2[index2++];
}
//end

```

recurrence relation:

$$T(n) = \begin{cases} O(1) & \text{if } n = 1, \\ 2T(n/2) + O(n) & \text{if } n \geq 2 \end{cases}$$

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## 1) Quicksort

```
int Partition (int Arr, L, h) {  
    Pivot = Arr[L]  
    i = L, j = h;  
    do {  
        do { i++; } while (Arr[i] <= Pivot);  
        do { j--; } while (Arr[j] > Pivot);  
        if (i < j) {  
            swap (Arr[i], Arr[j]);  
        } while (i < j);  
        swap (Arr[L], Arr[j]);  
        return j;  
    }  
}
```

time complexity:

Best case:  $O(n \log n)$

Worst case:  $O(n^2)$

Average case:  $O(n \log n)$

```
void quicksort (Arr, L, h) {  
    if (L < h) {  
        int j = Partition (Arr, L, h);  
        quicksort (Arr, L, j);  
        quicksort (Arr, j+1, h);  
    }  
}
```