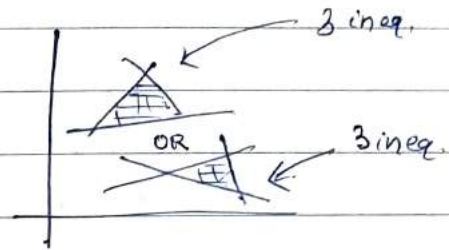


IP FORMULATIONS

cb

$$f_i(x) \leq b_i$$



$$f_1(x) \leq b_1 + yB$$

$$f_2(x) \leq b_2 + yB$$

$$f_3(x) \leq b_3 + yB$$

constraints on y .

$$f(x) + b_n + (1-y)B$$

assume constraint

$$N \text{ convex} = \text{convex}$$

$$x \leq By$$

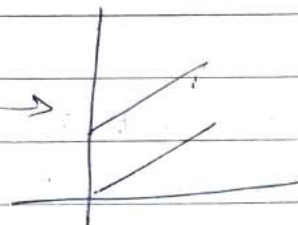
$$U \text{ convex} = x \text{ convex}$$

$$x \neq \text{nonzero when } y = 1.$$

$$\frac{K + Cx}{\downarrow}$$

$$Ky + Cx$$

$$y \in \{0, 1\}$$

When x is nonzero $y = 1$ 

$$0 \leq x_1 \leq 4K$$

$$C_1x_1 + C_2x_2 + C_3x_3$$

$$0 \leq x_2 \leq 5K$$

$$0 \leq x_3 \leq 6K$$



$$\text{optima} = c_1 x_1 + c_2 x_2 + c_3 x_3$$

$$0 \leq x_1 \leq 4k$$

$$0 \leq x_2 \leq 5k$$

$$0 \leq x_3 \leq 6k$$

$$4k\omega_1 \leq x_1 \leq 4k\omega_1$$

$$5k\omega_2 \leq x_2 \leq 5k\omega_1$$

$$0 \leq x_3 \leq 6k\omega_2$$

$$\omega_1, \omega_2 \in \{0, 1\}$$

① $\omega_1 = 0, \omega_2 = 0$

$$0 \leq x_1 \leq 4k$$

$$0 \leq x_2 \leq 0$$

$$0 \leq x_3 \leq 0$$

② $\omega_1 = 1, \omega_2 = 0$

$$4k \leq x_1 \leq 4k$$

$$0 \leq x_2 \leq 5k$$

$$0 \leq x_3 \leq 0$$

③ $\omega_1 = 1, \omega_2 = 1$

$$4k \leq x_1 \leq 4k$$

$$5k \leq x_2 \leq 5k$$

$$0 \leq x_3 \leq 6k$$

x_1	x_2	x_3
$4k$	$5k$	$6k$

~~$\omega_1 = 0, \omega_2 = 1$~~

$$0 \leq x_1 \leq 4k$$

$$5k \leq x_2 \leq 0$$

$$0 \leq x_3 \leq 6k$$

BALA'S ALGO

$$\text{min. } Z = 3x_1 + 5x_2 + 6x_3 + 9x_4 + 10x_5 + 10x_6$$

sf =

$$x_i \in \{0, 1\}, i = 1, 2, \dots, 6$$

$x_1 = 0$ (all)

$$C_1 = X$$

$$C_2 = \checkmark$$

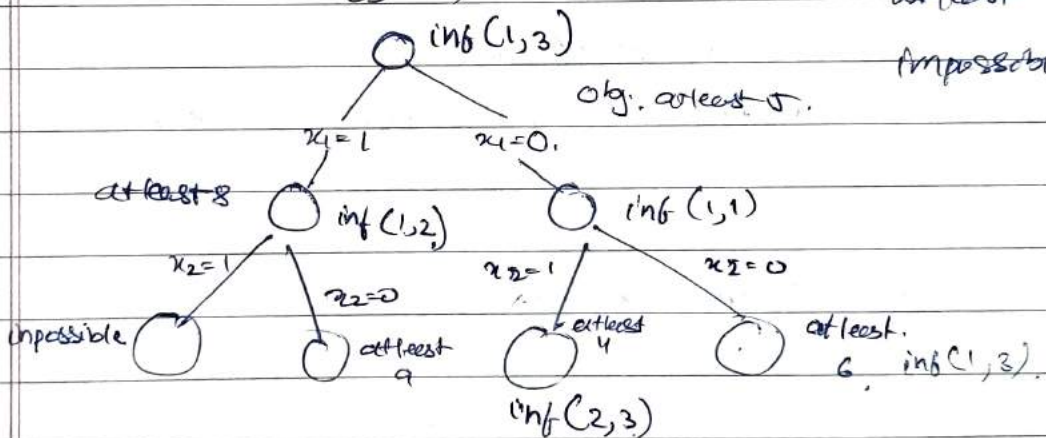
$$C_3 = X$$

branches

infeasible

at least

impossible,



$$\text{coz } 2:2 \mid a, a, (b, b) \in \infty$$

$$x \geq 3 \neq y \geq$$

or

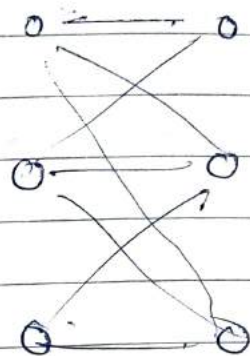
$$x \leq -3 - (1-y) \leq$$

bipartite matching

max. cut bipartite matching.

at every vertex only 1 edge.

get max. edge cut.



min. vertex cover.

select set or vertex st.

$x \in \{0, 1\}$

select least no. of vertex.

st. vertex of every edge is selected.

max independent set.

max. set. st. no 2 vertices are connected by an edge.

LP relaxation - rounding trick.

LOGICAL CONSTRAINTS

$$5x_1 + 2x_2 \leq 12 \quad \begin{cases} \text{true} \\ \text{false} \end{cases}$$

$$5x_1 + 2x_2 \leq 12 + \frac{1000}{\text{using binary variable.}} \rightarrow \text{not necessary.}$$

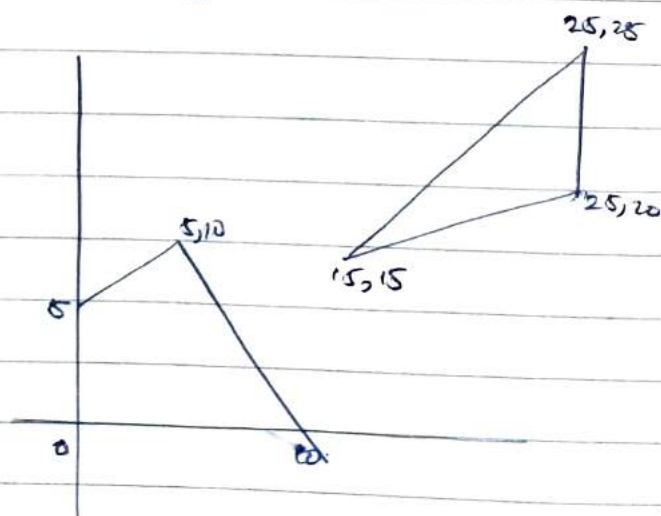
$$\left. \begin{aligned} b_1(\bar{x}) &\leq b_1 + y_1 B_1 \\ b_2(\bar{x}) &\leq b_2 + y_2 B_2 \end{aligned} \right\} \text{at least 1}$$

$$b_1(\bar{x}) \leq b_1 \Rightarrow b_2(\bar{x}) \leq b_2$$

$$A \rightarrow B \quad \approx \quad \bar{A} + B$$

K-fold alternatives.

f must satisfy at least K of the following constraints





$$x \leq 20$$

$$y \leq x$$

$$y = \frac{20-15}{25-15}x + \frac{20-15}{25-15}$$

$$y - y_1 = \frac{y_2 - y_1}{x_2 - x_1} (x - x_1)$$

$$y - 15 = \frac{5}{10} (x - 15)$$

2 →



$$2y - 30 = x - 15$$

$$-x + 2y \geq 15$$

$$\sum y_i = 1$$

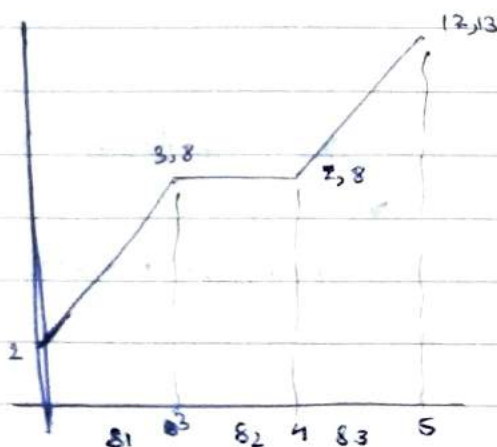
$$\sum y_i \leq 1$$

FIXED COSTS

$$C = \begin{cases} 0 & x = 0 \\ k + cx & x > 0 \end{cases}$$

$$\bar{c} = ky + cx \quad x \leq B_y$$

PIECEWISE LINEAR FUNCTIONS



$$C = 2 + 3\delta_1 + 0\delta_2 + 8\delta_3$$

$$3\omega_1 \leq \delta_1 \leq 3\omega_1$$

$$4\omega_2 \leq \delta_2 \leq 4\omega_2$$

$$0 \leq \delta_3 \leq 5\omega_2$$