

verizon  
class

230 + 50

THE ADDITION LANGUAGE

Syntax

 $\text{exp} ::= \text{num} \mid \bar{+} \text{exp} \text{ exp}$ 

Semantics

 $\text{num} ::= 0 \mid 1 \mid \dots$ 

BNF grammar.

→ syntactic category.

ex  $\bar{2} + \bar{4}$

↑ literals

1) what are we talking about?  
(in the addition language).

SEMANTIC DOMAINS

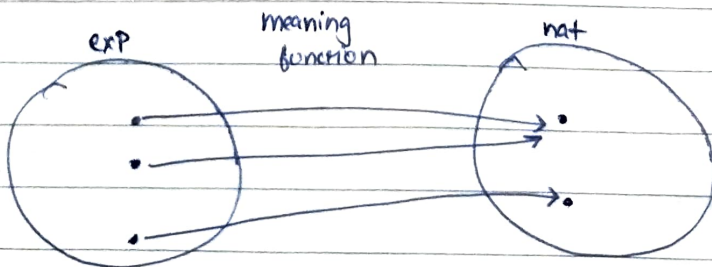
natural nos.

operations

 $+: \text{Nat} \times \text{Nat} \longrightarrow \text{Nat}$ 

notion of composition already built into semantic domain.

Now - we want a representation for that.

 $+(2, 3)$  $+(7, +(2, 3))$ 

## INDUCTIVE DEFNS.

antecedent  $\downarrow$   
 $\frac{n \text{ Nat}}{\bar{n} \text{ Num}}$  ~~Num~~ NAT  
 consequent  $\uparrow$   
 if  $n$  is a Nat,  
 then  $\bar{n}$  is a Num

$\frac{e \text{ Num} \quad \text{NUM}}{e \text{ Exp}}$   
 if  $e$  is a Num,  
 then  $e$  is an Exp

antecedent  $\downarrow \quad \downarrow$   
 $\frac{e_1 \text{ Exp} \quad e_2 \text{ Exp}}{+ e_1 e_2 \text{ Exp}}$  PLUS  
 if  $e_1$  is an Exp and  $e_2$  is an Exp,  
 then  $+ e_1 e_2$  is an Exp.

## NAT

rules give us a way to define a valid expression,  
 a valid expr. is one which allows me to derive an expr.

$+ \bar{3} \bar{4}$

- |                        |                       |                                   |
|------------------------|-----------------------|-----------------------------------|
| 1. 3 is a Nat          |                       | } proof / deduction / derivation. |
| 2. $\bar{3}$ is a Num  | from 1 using NAT      |                                   |
| 3. 4 is a Nat          |                       |                                   |
| 4. $\bar{4}$ is a Num  | from 3 using NAT      |                                   |
| 5. $+ \bar{3} \bar{4}$ | from 2, 4 using PLUS. |                                   |

$+ \bar{3} \text{ Exp?}$  X

$\hookrightarrow$  this is not even an expression  $\neq$  Exp.

$\vdash_{Exp}$  — deduce in the system Exp.  
that  $+ 3 4$  is an Exp.

$\vdash_{Exp} + \bar{3} \bar{4} \text{ Exp}$

$\nvdash_{Exp} \bar{+} \bar{3} \text{ Exp}$

→ used to prove if programs are syntactically correct.

helps to have a semantic domain in mind  
how to compose operations.

### VALUATION JUDGEMENTS (VAL)

$\frac{e \Rightarrow n}{\text{}} \quad \begin{array}{l} e: \text{Exp} \\ n: \text{Nat} \end{array}$

①  $\frac{\text{---}}{\bar{n} \Rightarrow n} \text{ NUM}$

②  $\frac{e_1 \Rightarrow n_1 \quad e_2 \Rightarrow n_2}{\bar{e}_1 \bar{e}_2 \Rightarrow n_1 + n_2} \text{ PLUS}$   
or  $\frac{\text{---}}{\text{or } \{ \bar{+}(\bar{e}_1, \bar{e}_2) \} }$

$$(+ \bar{7} (+ \bar{2} \bar{3}))$$

1. 7 is a Nat
2.  $\bar{7}$  is a Num from ①, NAT
3. 2 is a Nat
4.  $\bar{2}$  is a Num from ③, NAT
5. 3 is a Nat
6.  $\bar{3}$  is a Num from ④, NAT
7.  $\bar{2}$  is an Exp from ④, NUM
8.  $\bar{3}$  is an Exp from ⑤, NUM
9.  $\bar{7} \bar{2} \bar{3}$  is an Exp from ④, ⑤, PLUS
- ⑩.  $\bar{7}$  is an Exp from ②, NUM
11.  $(+ \bar{7} (+ \bar{2} \bar{3}))$  ⑩, ⑨, PLUS.

$$\vdash_{\text{Exp}} (+ \bar{7} (+ \bar{2} \bar{3})) \text{ Exp}$$

commutative, associative

$$11 : \text{Exp} \rightarrow \text{Nat}.$$

$$|11| = 1$$

$$|11 e_1 e_2| = 1 + |e_1| + |e_2|$$

CONJECTURE

thm

(uniqueness)

$$\forall e : \text{Exp} \exists ! n \in \text{Nat} : e \Rightarrow n$$

→ unique

$$+ \bar{3} \bar{4} \Rightarrow 7$$

to prove this claim.

1.  $\bar{3} \Rightarrow 3$  from NUM ( $\vdash_{\text{VAL}} \text{NUM}$ )
2.  $\bar{4} \Rightarrow 4$   $\vdash_{\text{VAL}} \text{NUM}$
3.  $+(3, 4) \Rightarrow 7$  ARITHMETIC.

$$\therefore + \bar{3} \bar{4} \Rightarrow 7 \quad 1, 2, 3 \text{ using PLUS } (\vdash_{\text{VAL}} \text{PLUS})$$

proof  $\forall e: \text{Exp} \exists ! n : e \Rightarrow n$

$:$  = element of

1.  $\forall e: \text{Exp} \exists n \in \text{Nat} : e \Rightarrow n$
2.  $e \Rightarrow n_1 \ \& \ e \Rightarrow n_2$   
IMPLIES  $n_1 = n_2$

proof of ① by induction on the structure of  $e$ .  
(by induction on  $\bullet \text{Exp}$ )

base case:

1.  $e \equiv \bar{n}$  assumption
2.  $\bar{n} \Rightarrow n$  NUM rule

DONE.

inductive case:

1.  $e \equiv \bar{+} e_1 e_2$  assumption.

2.  $e_1 \Rightarrow n_1$  by IH  
on  $e_1$   
for some  $n_1$

3.  $e_2 \Rightarrow n_2$  by IH  
on  $e_2$

4.  $\bar{+} e_1 e_2 \Rightarrow + (n_1, n_2)$

5.  $+ (n_1, n_2)$  is well defined

$+$  is total

$\text{NAT} \times \text{NAT}$

5. let  $n = + (n_1, n_2)$

6.  $e \Rightarrow n$  3, 5, 1