Introduction to Programming

Week – 7, Lecture – 3

Solving Problems Recursively

SAURABH SRIVASTAVA

DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

IIT KANPUR

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Using loops to solve problems

```
Statements before the loop
Initialisation - setting values for important variables
    Statements, using different values of one or more variables
   Update to the values of one or more variables
   condition over one or more variables (to go out of the loop)
Statements after the loop
```

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We used the example of Factorial Calculation as the example to show the process

We saw how it could be done using three different loops

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• We will see some examples in a minute... just grasp the concepts right now

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- ... which then calls the first function, then too, it is called recursion
- Actually, there can be any number of functions in the chain, not just two...

While the former type is called *direct* recursion, the latter is called *indirect* recursion

Direct and Indirect Recursion

Direct Recursion

Direct and Indirect Recursion

Direct Recursion

Indirect Recursion – with 2 functions

Direct and Indirect Recursion

```
function f()
                        function f()
                                                    function f()
     f();
                             g();
                                                        g();
                         function g()
                                                   function g()
                             f();
                                                        h();
                                                    function h()
                                                        f();
                        Indirect Recursion –
                                                                       Indirect Recursion –
Direct Recursion
                        with 2 functions
                                                                       with 3 functions
```

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You use it to prove that a particular premise is true for all Natural Numbers

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- and, how solutions to smaller problems or subproblems, can be combined to give the complete solution

In terms of programming, it is often easy to do this using recursion

• For instance, you write a function, and call it repeatedly, to construct a solution in a bottom-up fashion

```
#include<stdio.h>
long factorial(int number)
        if(number < 0)
                return -1:
        else if(number == 0 \mid \mid number == 1)
                return 1;
        else
                return number * factorial(number-1);
int main()
        int number;
        do
                printf("Give me a small positive integer: ");
                scanf("%d", &number);
        while(number < 0 || number > 20);
        printf("Calculated Factorial: %ld\n", factorial(number));
```

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We have a function here to calculate factorial

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We have a function here to calculate factorial

It takes as input, an integer, and returns its factorial

```
if(number < 0)
        return -1;
```

If it is a negative number, it returns -1 — a convention to show an error

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This is really some validation, not actual factorial calculation

```
else if(number == 0 \mid \mid number == 1)
         return 1;
```

If the number is 0 or 1, the answer is simply, 1

```
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         return 1;
```

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The base case here is actually when the number is 1, we have to include 0, because by convention, 0! = 1

```
else
        return number * factorial(number-1);
```

If the number is anything else, it can be calculated as:

number * factorial(number-1)

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Basically, we decide to solve the same problem (i.e., calculating factorial), but with a smaller sized input

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long factorial(int number)
        if(number < 0)
                return -1;
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                return 1;
        else
                return number * factorial(number-1);
```

Run this code in your head a few times, to convince yourself that it works !!

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Run this code in your head a few times, to convince yourself that it works !!

Then try to understand how we solved a large problem by solving "smaller problems recursively", and then "combining their solutions appropriately"

Example 2 – Binary Search (let's code it!)

```
int binary_search(int *sorted_array, int beg, int end, int key)
       int mid;
       if(beg > end)
                return -1;
       mid = beg + (end - beg)/2;
       if(sorted array[mid] == key)
                return mid;
       else if(sorted array[mid] < key)</pre>
               return binary search(sorted array, mid+1, end, key);
       else
               return binary search(sorted array, beg, mid-1, key);
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The function expects the following inputs:

- 1. The sorted array
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- 3. The ending index of the part of the array to search
- 4. The value to search

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The function expects the following inputs:

- 1. The sorted array
- 2. The beginning index of the part of the array to search
- 3. The ending index of the part of the array to search
- 4. The value to search

If the value is found in the part, it returns the index at which it was found, otherwise, it returns -1 (another convention)

```
if(beg > end)
       return -1;
```

If the beginning index is greater than the ending index, it means we have a part of size 0

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So, we return -1 irrespective of what the key is (we cannot find anything in a zero-sized array)

As discussed before, a good point to check the sorted array for the key, is the mid-point (or almost the mid-point) of the part of the array to search

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So, we find the mid-point of the array part

```
if(sorted_array[mid] == key)
        return mid;
```

Then we check if the key is found at the mid-point

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Then we check if the key is found at the mid-point

If so, we return the mid-point of the array part as the output

```
else if(sorted array[mid] < key)</pre>
        return binary search(sorted array, mid+1, end, key);
```

Otherwise, if the key greater than the element at the midpoint, we can be sure that "if at all" the key is present in this part of the array, it can only be on the right of the midpoint

```
else if(sorted array[mid] < key)</pre>
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So, we find the result of the binary search in the part of array which starts at mid+1 and ends at end, and pass on the result as the result of binary search in the whole part

```
else
       return binary search(sorted array, beg, mid-1, key);
```

Or else, i.e., the key is smaller than the element at the midpoint, the only possibility of finding it in the current part is to the left of the mid-point

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else
       return binary search(sorted array, beg, mid-1, key);
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Or else, i.e., the key is smaller than the element at the midpoint, the only possibility of finding it in the current part is to the left of the mid-point

So, we find the result of the binary search in the part of array which starts at beg and ends at mid-1, and pass on the result as the result of binary search in the whole part

```
search_index = binary_search(array, 0, 9, key);
```

To use the function, we call it with the beg value set to 0, and the end value set to capacity-1

```
int binary_search(int *sorted_array, int beg, int end, int key)
       int mid;
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```

Again, make sure you get the flow here, i.e., you understand why this works !!

Factorial

At each step, we reduce the problem size by 1

Binary Search

At each step, we reduce the problem size by roughly half

Factorial

- At each step, we reduce the problem size by 1
- At each step, the solution to the larger problem is either trivial (if the number is 0 or 1), or, given by combining the solution of a smaller problem with a trivial element (the number itself)

Binary Search

- At each step, we reduce the problem size by roughly half
- At each step, the solution to the larger problem is either trivial (if the element is found at the mid-point, or the part of array to search is of size 0), or, given by the solution to either of the smaller problems

Factorial

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- At each step, the solution to the larger problem is either trivial (if the number is 0 or 1), or, given by combining the solution of a smaller problem with a trivial element (the number itself)
- We know we have the final solution, when the base case is found (i.e. the number is 1 or 0)

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- At each step, the solution to the larger problem is either trivial (if the number is 0 or 1), or, given by combining the solution of a smaller problem with a trivial element (the number itself)
- We know we have the final solution, when the base case is found (i.e. the number is 1 or 0)
- It will take n steps to find the final solution, when the input size is n (i.e. number is n)

Binary Search

- At each step, we reduce the problem size by roughly half
- At each step, the solution to the larger problem is either trivial (if the element is found at the mid-point, or the part of array to search is of size 0), or, given by the solution to either of the smaller problems
- We know we have the solution, when either of the two base cases is found (element found or array exhausted)
- It will take, at max, [log₂n] steps to find the final solution, when the input size is n (i.e. the array size)

Homework!!

While Divide and Conquer and Recursion are orthogonal concepts, they are often used together

- You can, theory, write Divide and Conquer algorithms iteratively as well
- Write a recursive as well as iterative Divide and Conquer algorithm to find the smallest element in an array

Read more about the overall Divide and Conquer approach

 You may start here: https://www.geeksforgeeks.org/divide-and-conquer/