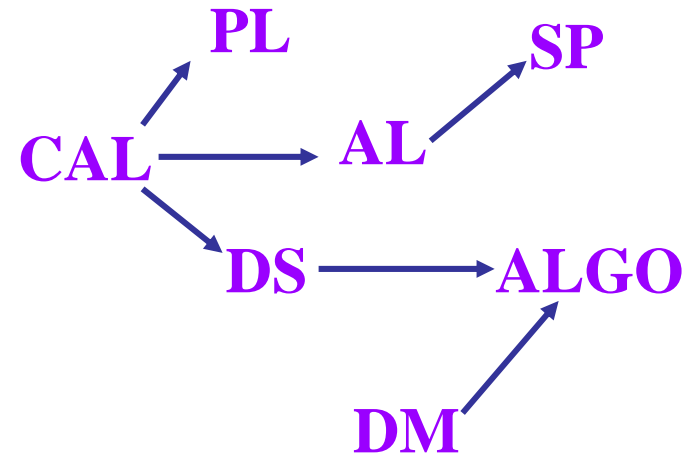
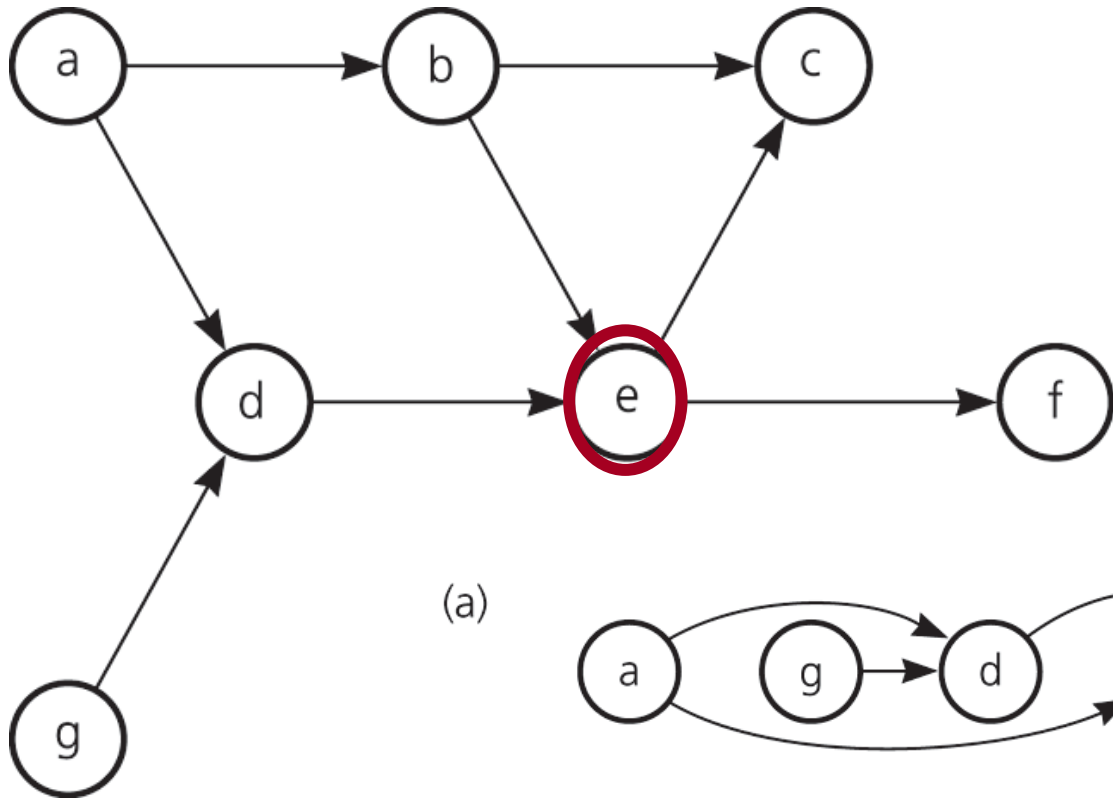


Graph Problems

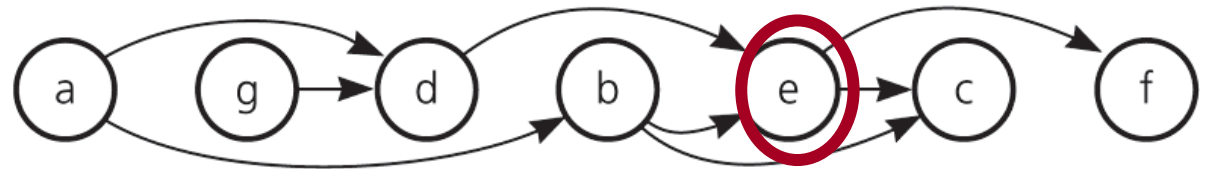
- ❑ **Critical Path Analysis**
- ❑ **Maximum Flow Problem**
- ❑ **Other Difficult Problems**

Activity-on-Arrow (AOA) Network

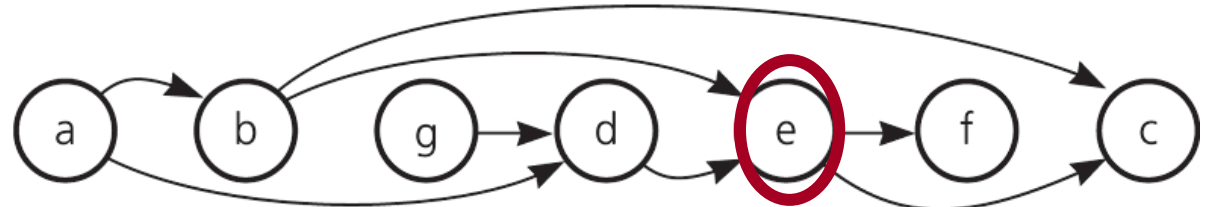
Data Structures



(a)

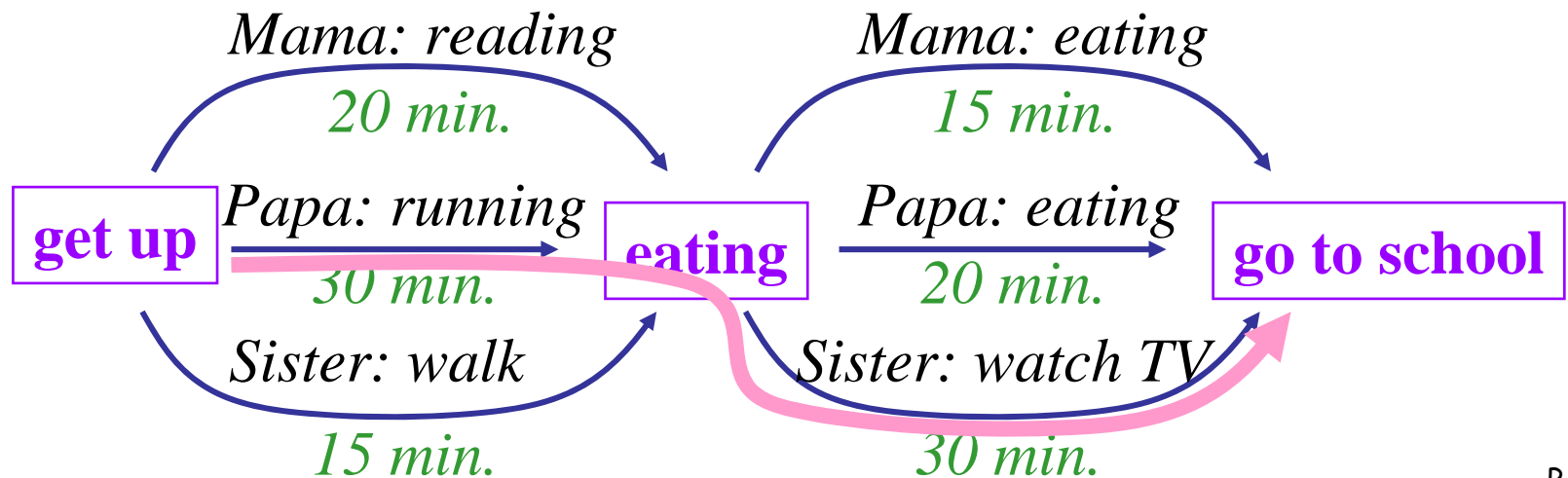


(b)



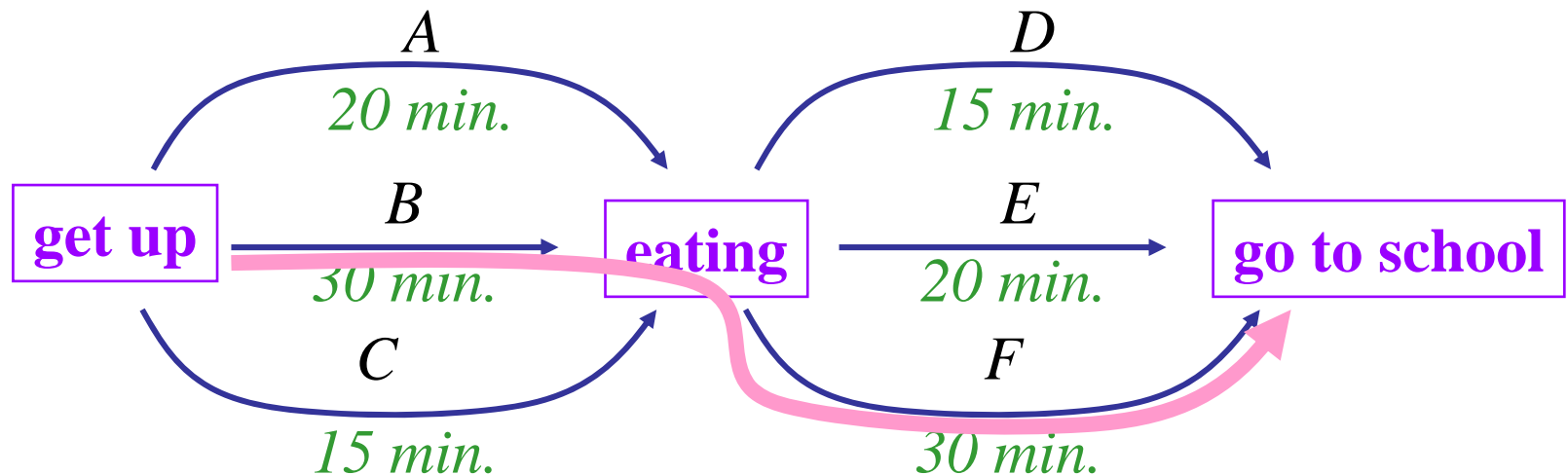
Activity-on-Edge (AOE) Network

- Directed edge: activity (task) to be performed
- **Vertex**: event to signal the completion of certain activities
- **Edge weight**: the time required to perform an activity
- Path length: the total time from the start to the last event
- **Critical Path**: a path with the longest length
 - *the minimum time required to complete the project*



Critical Path Analysis

- ❑ Earliest time of an activity/event: $\text{early}(E) = 30$
- ❑ Latest time of an activity/event: $\text{late}(E) = 60 - 20 = 40$
- ❑ **Critical activity:** $\text{late}(F) = \text{early}(F) = 30$



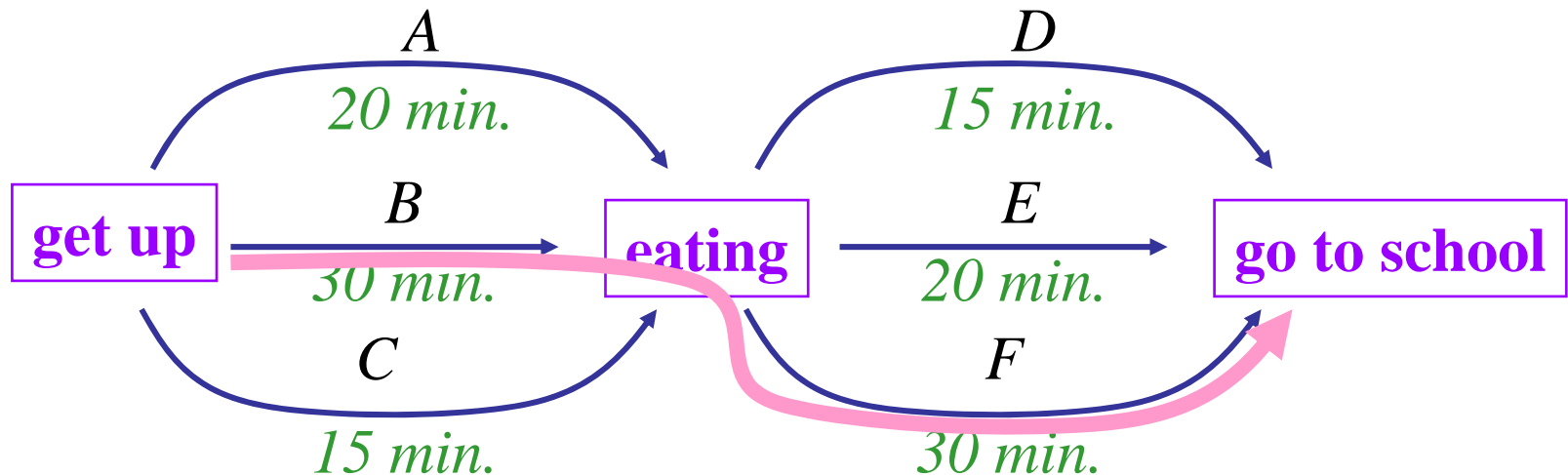
Critical Path Analysis

□ AOE network is very useful for evaluating the performance of many types of projects

– Project Evaluation and Review Techniques (PERT)

Q1. What is *the least amount of time* in which the project may be complete (assuming there is no cycle in the network)?

Q2. Which *activities* should be speeded to reduce project length?



Critical Path Analysis: *Background*

❑ Developed in the 1950s by the US Navy

- Originally, it considered only *logical dependencies* among *project activities*. Since then, it has been expanded to include the *resources* related to each activity, through the process called *resource leveling*.

❑ John Fondahl

- US Marine Corps Sergeant
- Stanford CE Professor Emeritus
- 1961 Paper for the US Navy – "*Non-Computer Approach to the Critical Path Method for the Construction Industry*"

Critical Path Analysis: *Model*

□ Input

- A list of all *activities* required to complete the project
- The time (*duration*) that each activity will take to completion
- The *dependencies* between the activities.

□ Output

- The *longest* path of planned activities to the end of the project
- The *earliest* and *latest* that each activity can start and **finish without making the project longer**
- Determines “*critical*” activities (on the longest path)
- Prioritize activities for the effective management and to shorten the critical path of a project

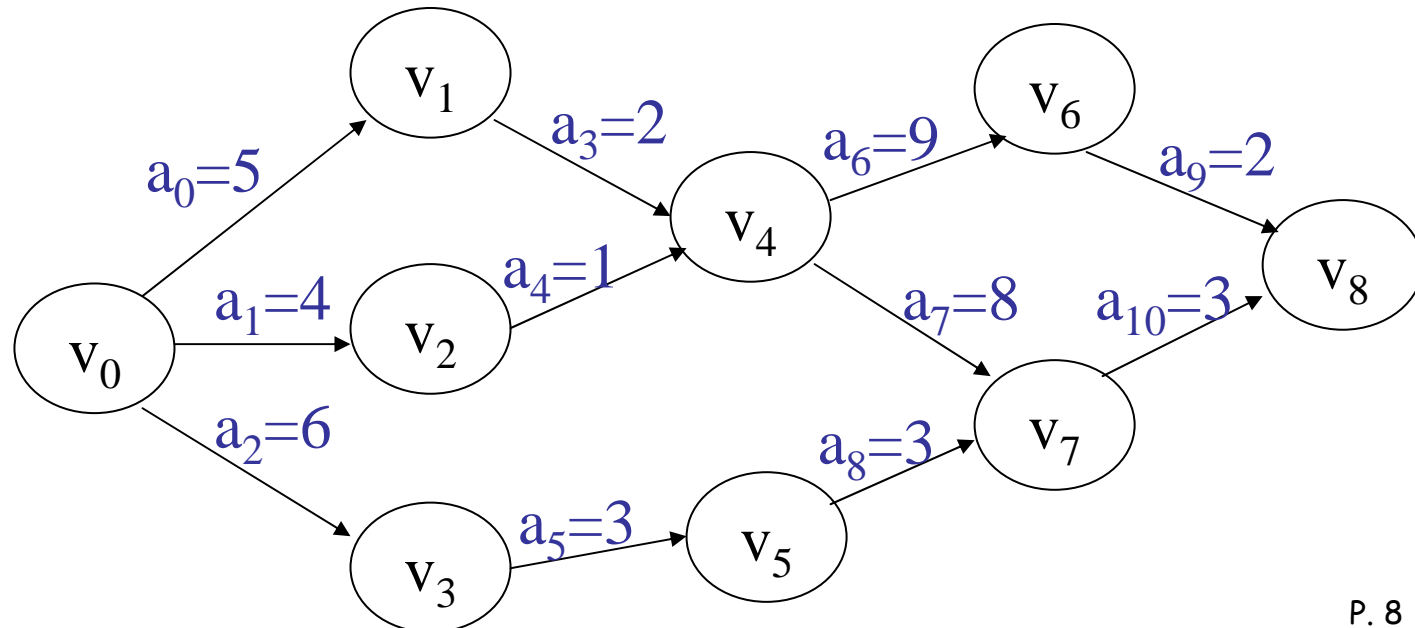
Critical Path Analysis: *Example*

□ Activities: a_0 a_1 a_2 a_3 a_4 a_5 a_6 a_7 a_8 a_9 a_{10}

□ Dependencies among Activities

$a_0 \rightarrow a_3$, $a_1 \rightarrow a_4$, $a_2 \rightarrow a_5$, $a_5 \rightarrow a_8$, $a_6 \rightarrow a_9$

a_3 $a_4 \rightarrow a_6$ a_7 , a_7 $a_8 \rightarrow a_{10}$



Critical Path Analysis: *Forward*

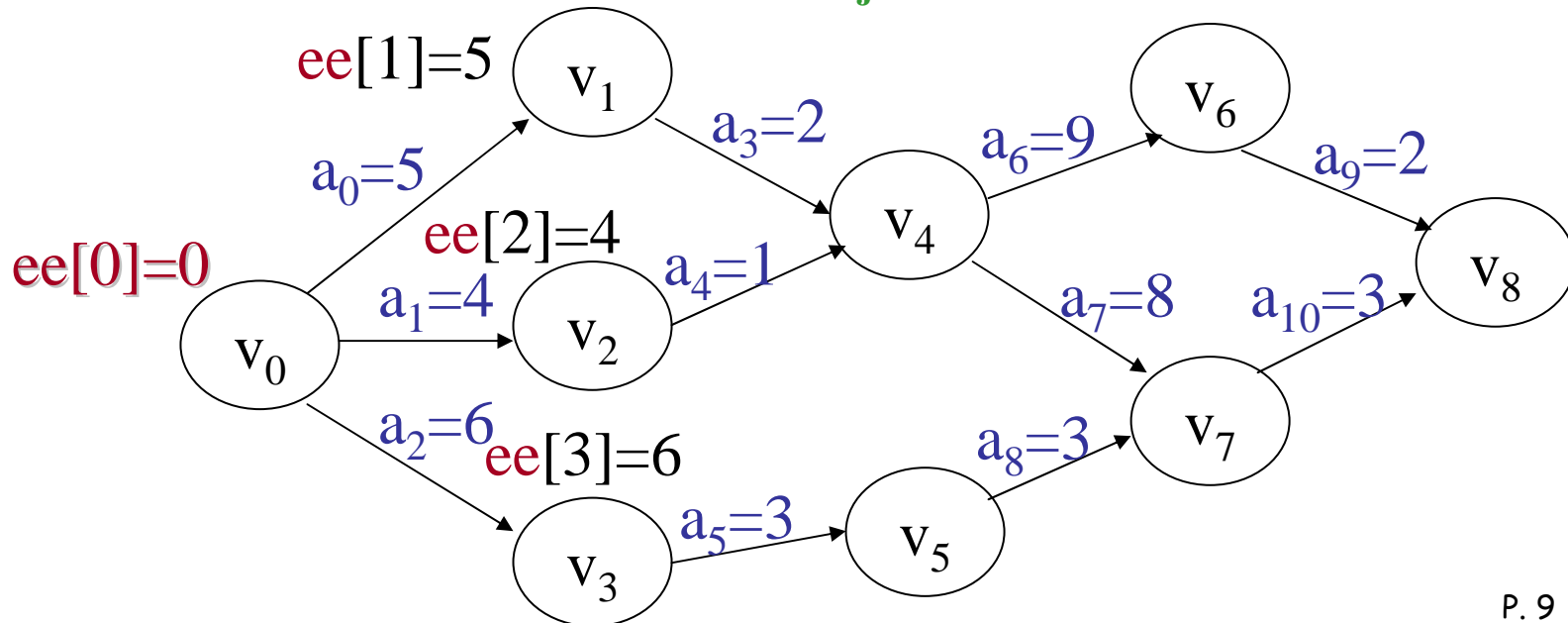
□ earliest time of an activity: $ea[0..10]$

– earliest time of an event: $ee[0..8]$

■ $ea[x] = ee[i]$ if a_x is on the edge $\langle v_i, v_j \rangle$

■ $ee[j] = \max\{ee[i] + \text{duration of } \langle v_i, v_j \rangle\}$ for every v_i that is an *immediate predecessor* of v_j

ea
[0]: 0
[1]: 0
[2]: 0



Critical Path Analysis: *Forward*

ea

[0]: 0

[1]: 0

[2]: 0

[3]: 5

[4]: 4

[5]: 6

[6]: 7

[7]: 7

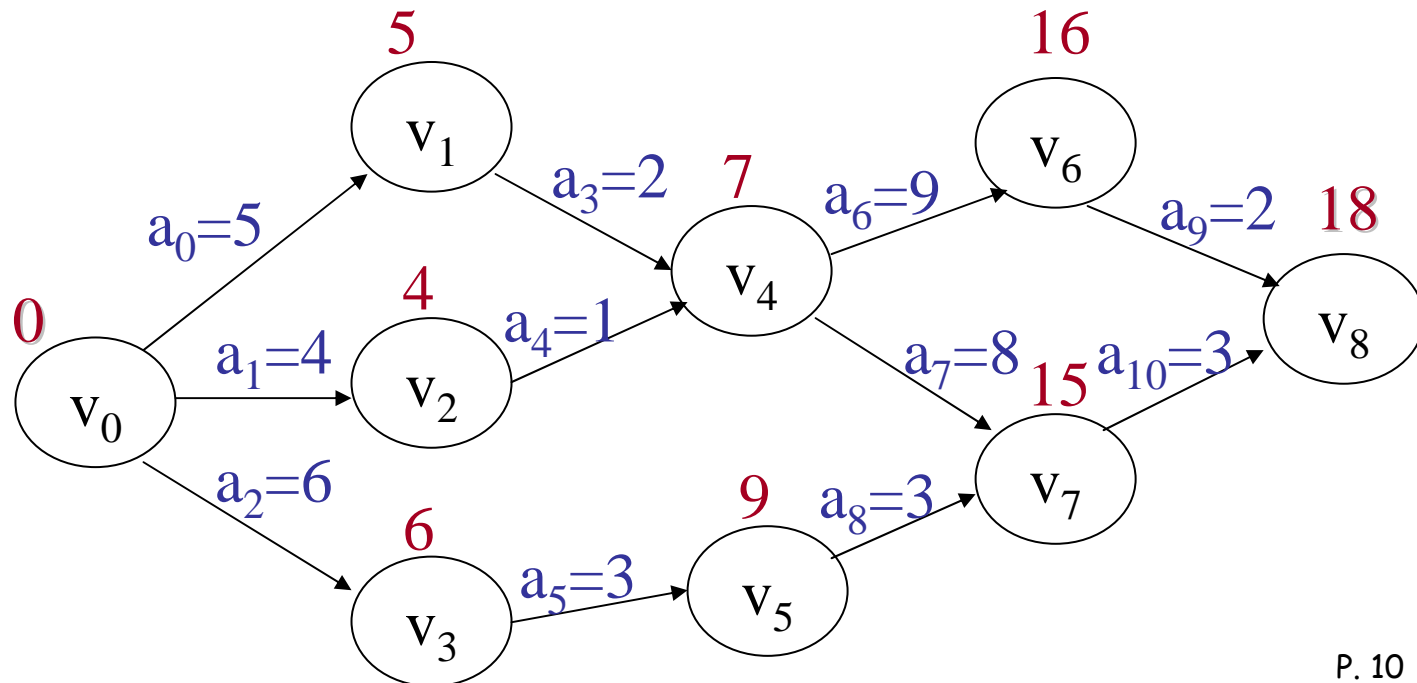
[8]: 9

[9]: 16

[10]: 15

■ $ea[x] = ee[i]$ if a_x is on the edge $\langle v_i, v_j \rangle$

■ $ee[j] = \max\{ee[i] + \text{duration of } \langle v_i, v_j \rangle\}$ for every v_i that is an *immediate predecessor* of v_j



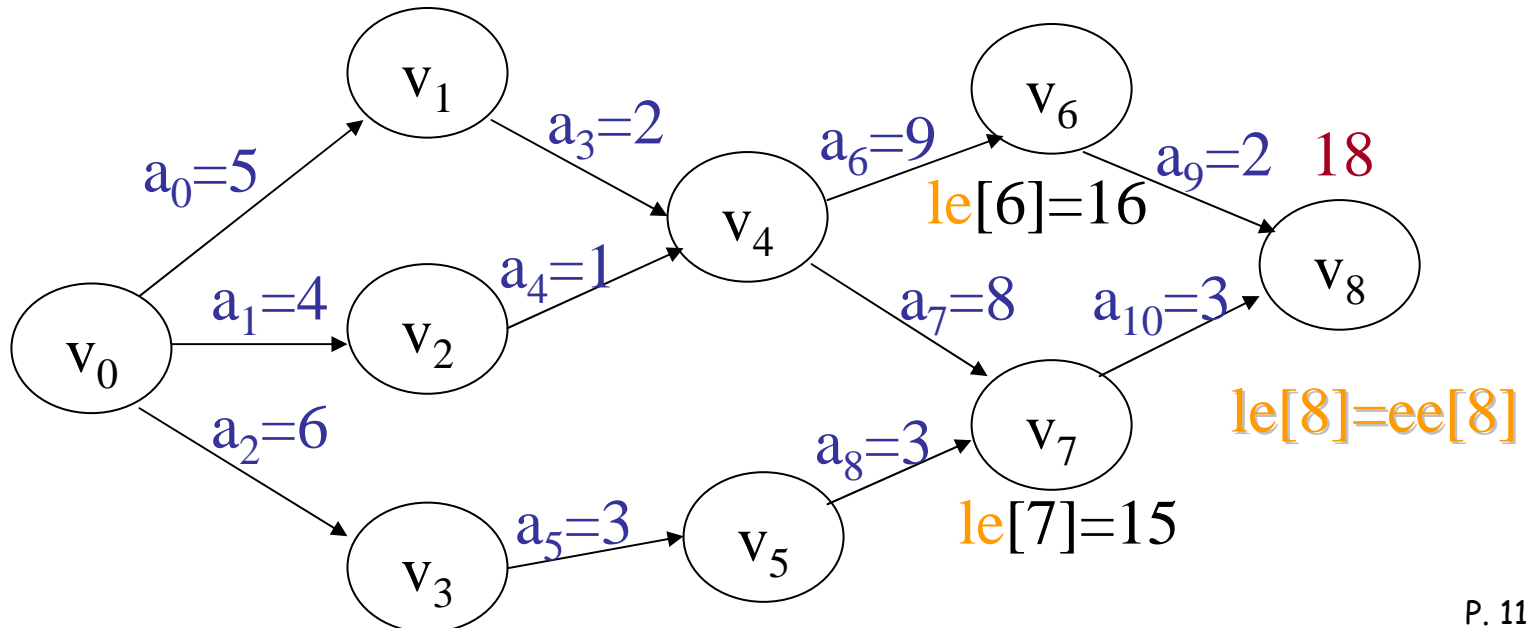
Critical Path Analysis: *Backward*

□ latest time of an activity: $la[0..10]$

– latest time of an event: $le[0..8]$

■ $la[x] = le[j] - \text{duration of } \langle v_i, v_j \rangle$, where a_x is on $\langle v_i, v_j \rangle$

■ $le[i] = \min\{le[j] - \text{duration of } \langle v_i, v_j \rangle\}$ for every v_j that is an immediate successor of v_i



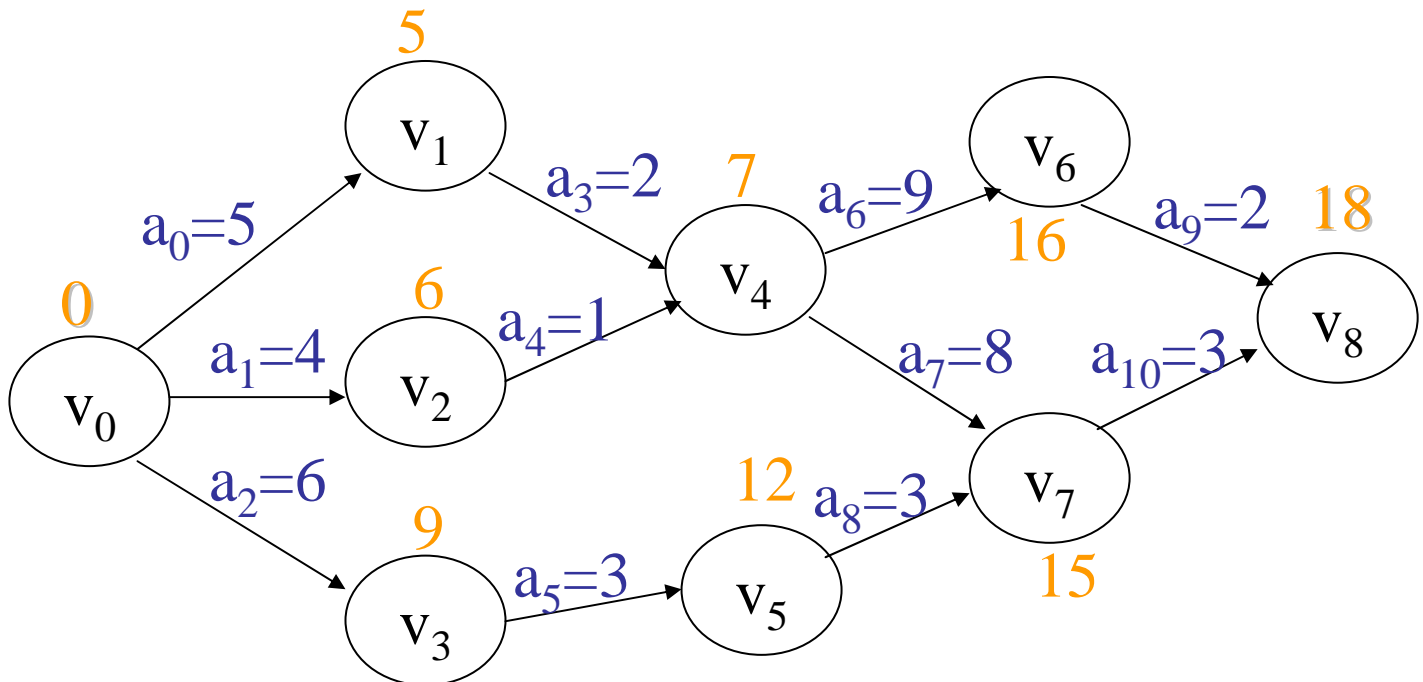
Critical Path Analysis: *Backward*

■ $la[x] = le[j] - \text{duration of } \langle v_i, v_j \rangle$, where a_x is on $\langle v_i, v_j \rangle$

■ $le[i] = \min\{le[j] - \text{duration of } \langle v_i, v_j \rangle\}$ for every v_j that is an *immediate successor* of v_i

[0]: 0
[1]: 2
[2]: 3
[3]: 5
[4]: 6
[5]: 9
[6]: 7
[7]: 7
[8]: 12
[9]: 16
[10]: 15

la



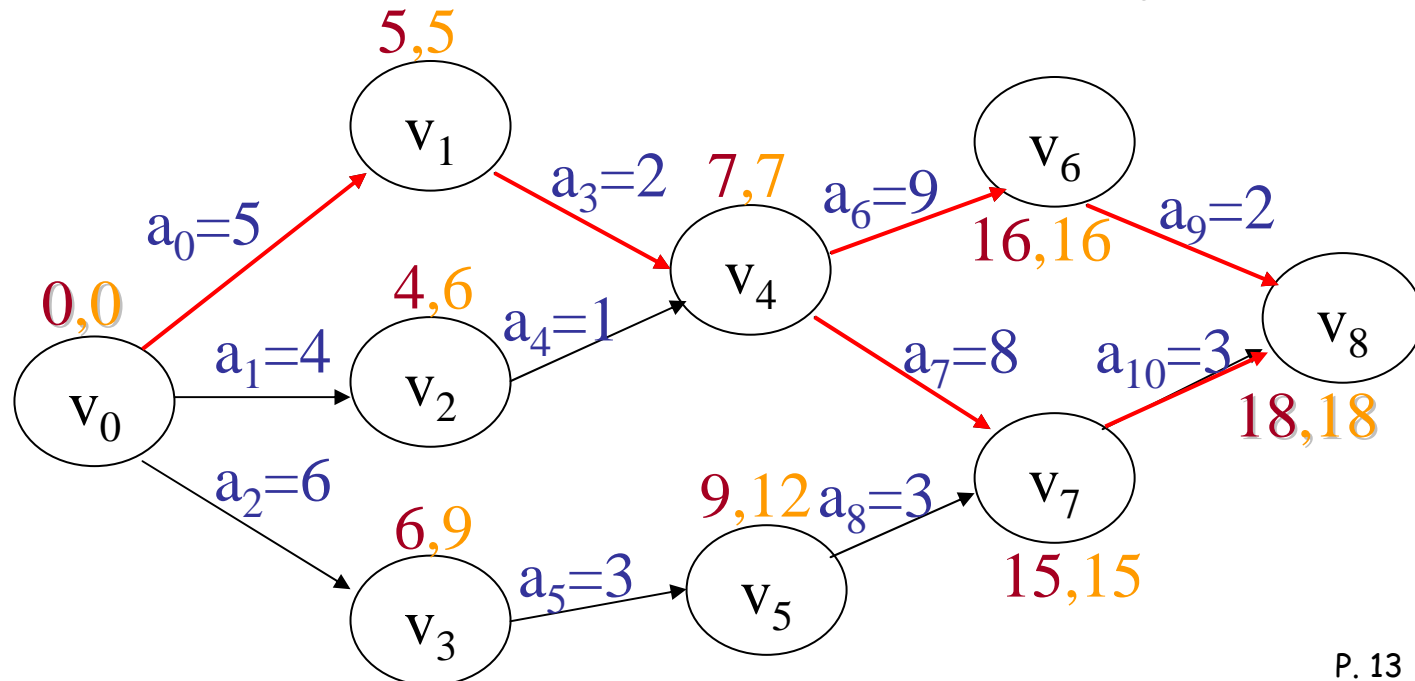
Critical Path Analysis: *Results*

<u>ea</u>	<u>la</u>	<u>la-ea</u>
[0]: 0	0	0
[1]: 0	2	2
[2]: 0	3	3
[3]: 5	5	0
[4]: 4	6	2
[5]: 6	9	3
[6]: 7	7	0
[7]: 7	7	0
[8]: 9	12	3
[9]: 16	16	0
[10]: 15	15	0

□ **la-ea** is called (total) *float* or *slack*

– amount of time that a task can be delayed without causing a delay to project completion time

la-ea=0 means a *critical activity*



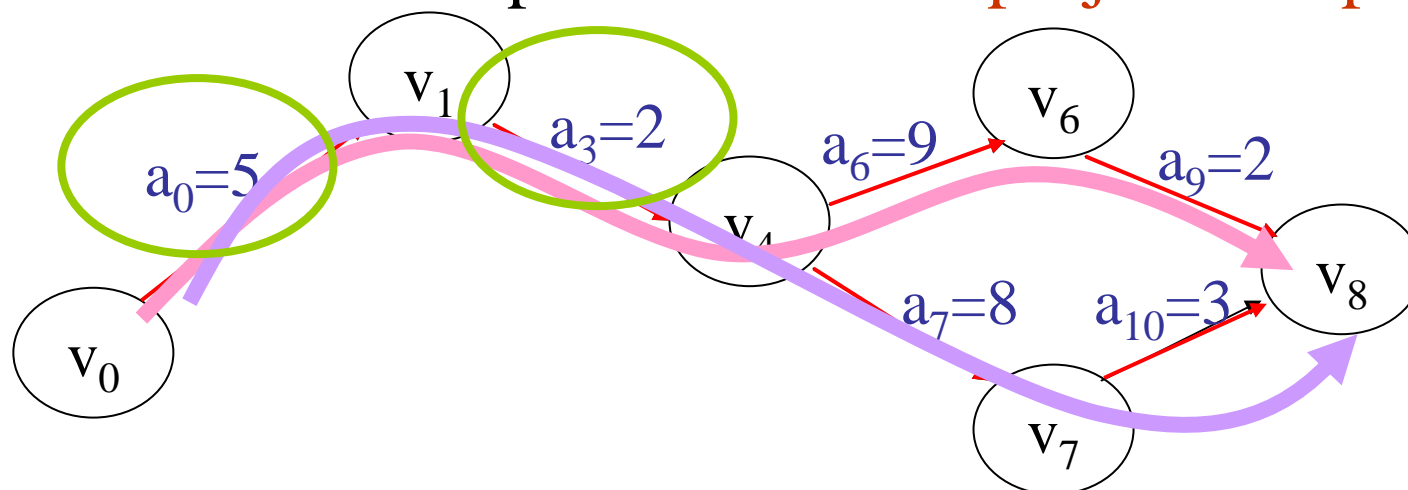
Critical Path Analysis: *Results*

□ Determine Critical Paths

- Delete all non-critical activities (*nonzero slack*)
- Generate all the paths from the start to the end

□ Speed up the activities on **all** critical paths

- *resource* can be concentrated on these activities in an attempt to **reduce the project completion time**



Critical Path Method: *Forward Phase*

□ Like *topSort1*

Initialize: $ee[v_0]=0$

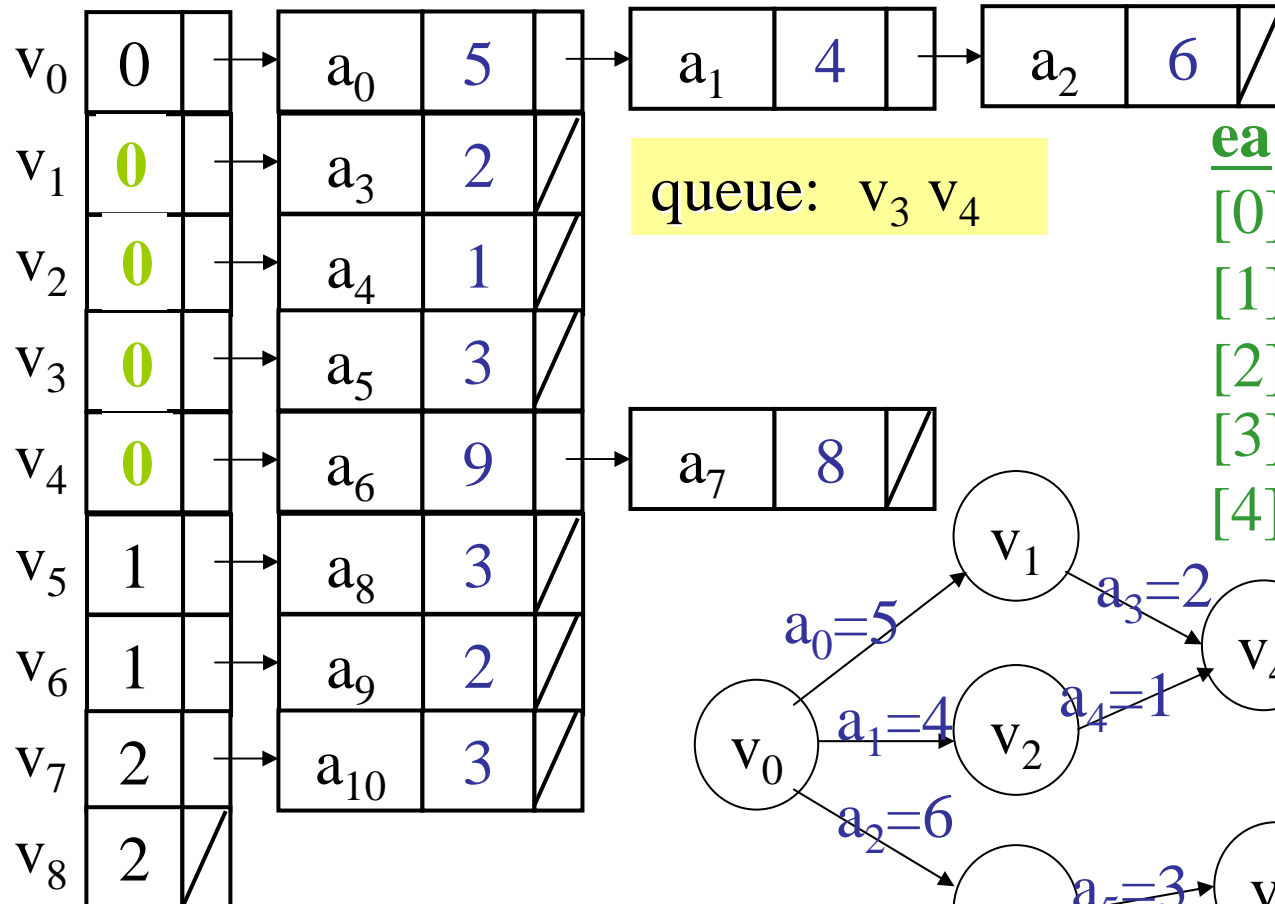
1. Find vertex v that has **no predecessor** (in-degree=0)
2. For each immediate successor u , do the following:
 - Set $ea[x] = ee[v]$, where x is the *activity* on $\langle v, u \rangle$
 - Set $ee[u] = \max\{ee[u], ee[v] + \text{duration of } \langle v, u \rangle\}$
 - Decrease the in-degree of u
3. Repeat the steps until *all vertices are visited*
 - For the vertex w that has no successor, $le[w]=ee[w]!$

Put u into queue or stack

Critical Path Method: *Forward Phase*

Data Structures

in-degree activity duration



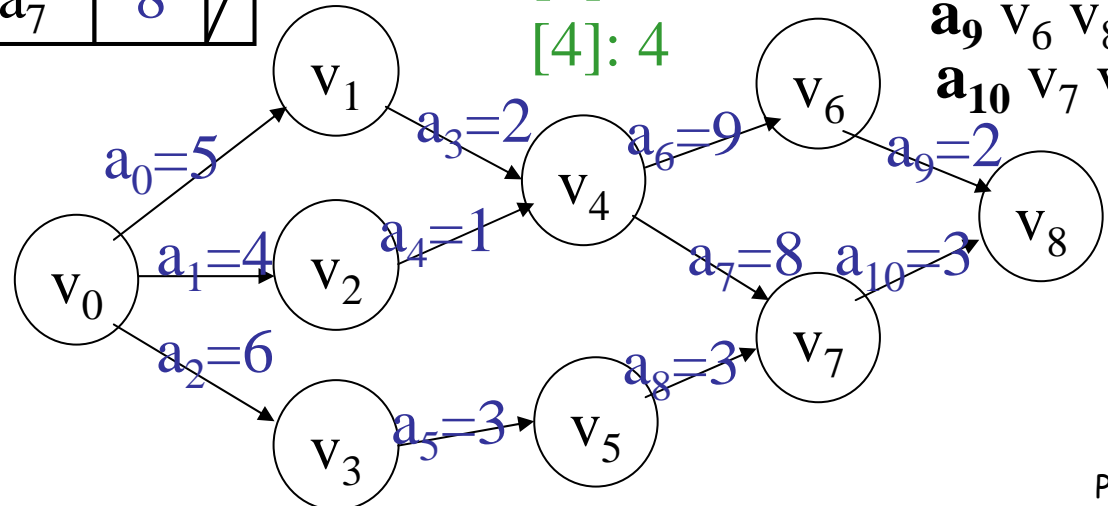
ea

[0]: 0
[1]: 0
[2]: 0
[3]: 5
[4]: 4

ee

[0]: 0
[1]: 5
[2]: 4
[3]: 6
[4]: 7

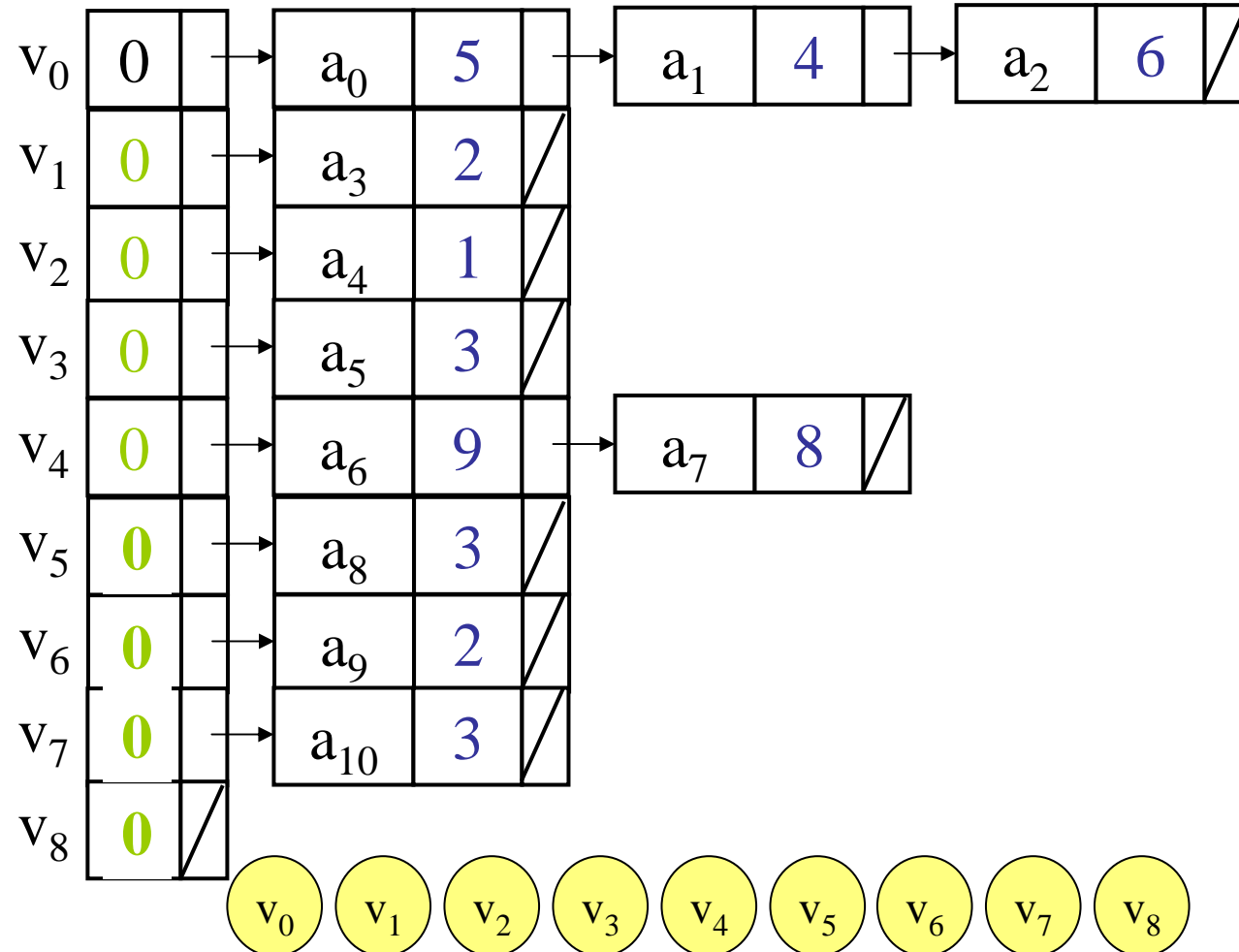
a₀ v₀ v₁
a₁ v₀ v₂
a₂ v₀ v₃
a₃ v₁ v₄
a₄ v₂ v₄
a₅ v₃ v₅
a₆ v₄ v₆
a₇ v₄ v₇
a₈ v₅ v₇
a₉ v₆ v₈
a₁₀ v₇ v₈



Critical Path Method: *Forward Phase*

Data Structures

in-degree activity duration



ea

[0]: 0
[1]: 0
[2]: 0
[3]: 5
[4]: 4
[5]: 6
[6]: 7
[7]: 7

ee

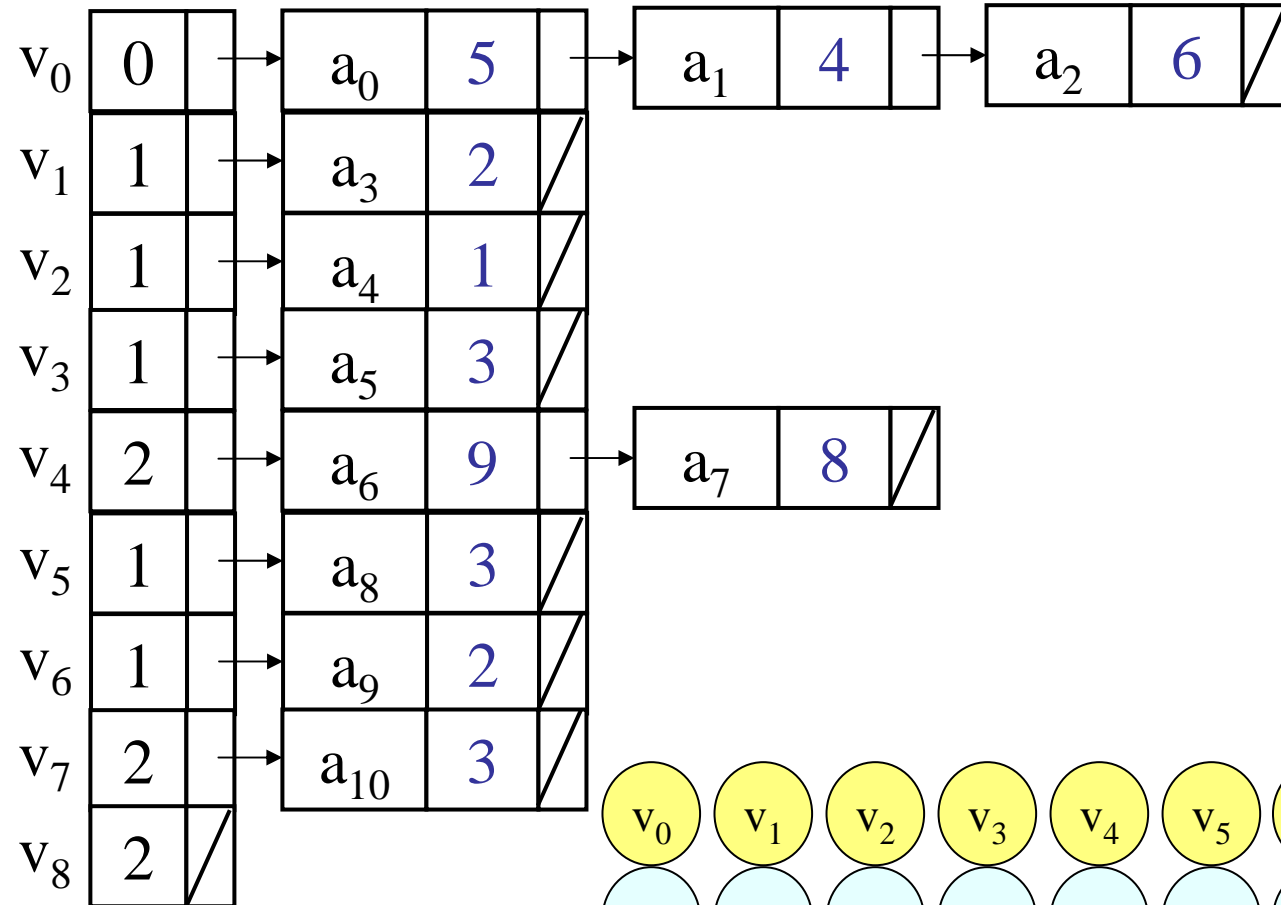
[0]: 0
[1]: 5
[2]: 4
[3]: 6
[4]: 7
[5]: 9
[6]: 16
[7]: 15

a₀ v₀ v₁
a₁ v₀ v₂
a₂ v₀ v₃
a₃ v₁ v₄
a₄ v₂ v₄
a₅ v₃ v₅
a₆ v₄ v₆
a₇ v₄ v₇
a₈ v₅ v₇
a₉ v₆ v₈
a₁₀ v₇ v₈

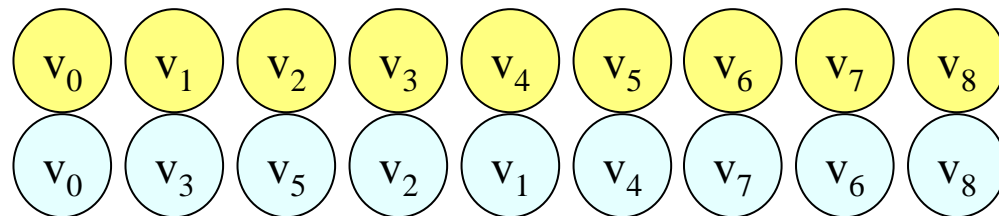
Critical Path Method: *Forward Phase*

Data Structures

in-degree activity duration



a_0 v_0 v_1
 a_1 v_0 v_2
 a_2 v_0 v_3
 a_3 v_1 v_4
 a_4 v_2 v_4
 a_5 v_3 v_5
 a_6 v_4 v_6
 a_7 v_4 v_7
 a_8 v_5 v_7
 a_9 v_6 v_8
 a_{10} v_7 v_8



Critical Path Method: *Backward Phase*

Initialize: $le[u] = ee[u]$

1. Find vertex u that has **no successor** (out-degree=0)
2. For each immediate predecessor v , do the following:
 - Set $la[x] = le[u] - \text{duration of } \langle v, u \rangle$, where x is the *activity* on $\langle v, u \rangle$
 - Set $le[v] = \min\{le[v], le[u] - \text{duration of } \langle v, u \rangle\}$
 - Decrease the out-degree of v
3. Repeat the steps until *all vertices are visited*
 - For the vertex w that has no predecessor, $le[w] = ee[w]!$

Put v into queue or stack

Critical Path Method: *Backward Phase*

Inverse adjacency list

out-degree activity duration

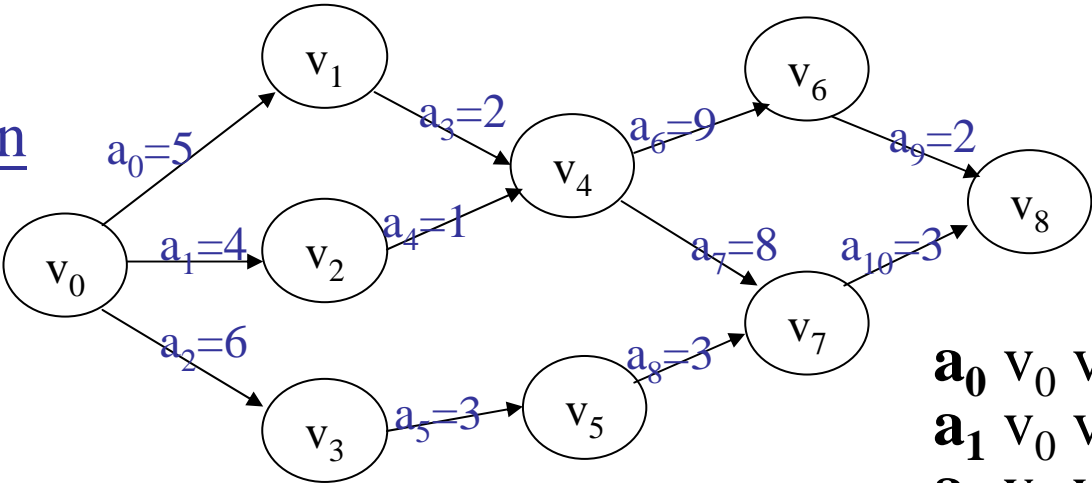
v ₀	2	/
v ₁	1	
v ₂	1	
v ₃	0	
v ₄	1	
v ₅	0	
v ₆	0	
v ₇	0	
v ₈	0	

a ₀	5	/
a ₁	4	/
a ₂	6	/
a ₃	2	
a ₄	1	/
a ₅	3	/
a ₆	9	/
a ₇	8	
a ₉	2	

a ₄	1	/
----------------	---	---

a ₈	3	/
a ₁₀	3	/

v₆
stack



[0]: 3
[2]: 3
[5]: 9
[7]: 7
[8]: 12
[9]: 16
[10]: 15
la

[0]: 3
[3]: 9
[4]: 7
[5]: 12
[6]: 16
[7]: 15
[8]: 18
le

- a₀ v₀ v₁
- a₁ v₀ v₂
- a₂ v₀ v₃
- a₃ v₁ v₄
- a₄ v₂ v₄
- a₅ v₃ v₅
- a₆ v₄ v₆
- a₇ v₄ v₇
- a₈ v₅ v₇
- a₉ v₆ v₈
- a₁₀ v₇ v₈

Critical Path Method: *Backward Phase*

Data Structures

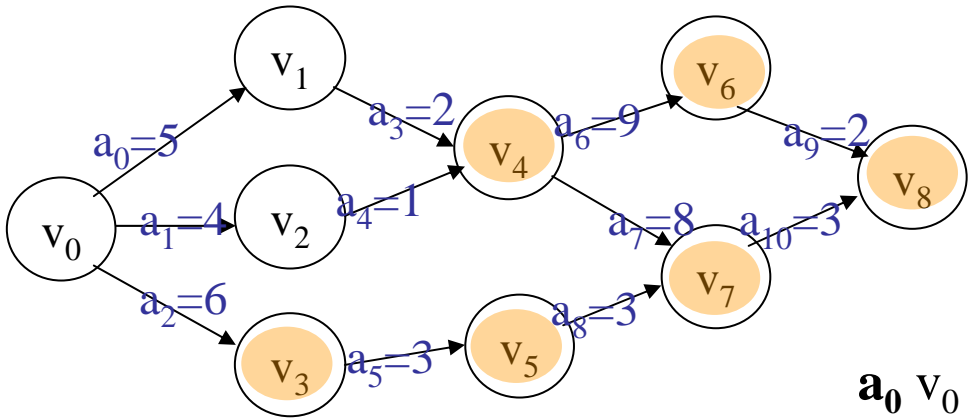
Inverse adjacency list

out-degree activity duration

v ₀	2	/		
v ₁	0		a ₀	5
v ₂	0		a ₁	4
v ₃	0		a ₂	6
v ₄	0		a ₃	2
v ₅	0		a ₅	3
v ₆	0		a ₆	9
v ₇	0		a ₇	8
v ₈	0		a ₉	2

a ₄	1	/
----------------	---	---

a ₈	3	/
a ₁₀	3	/



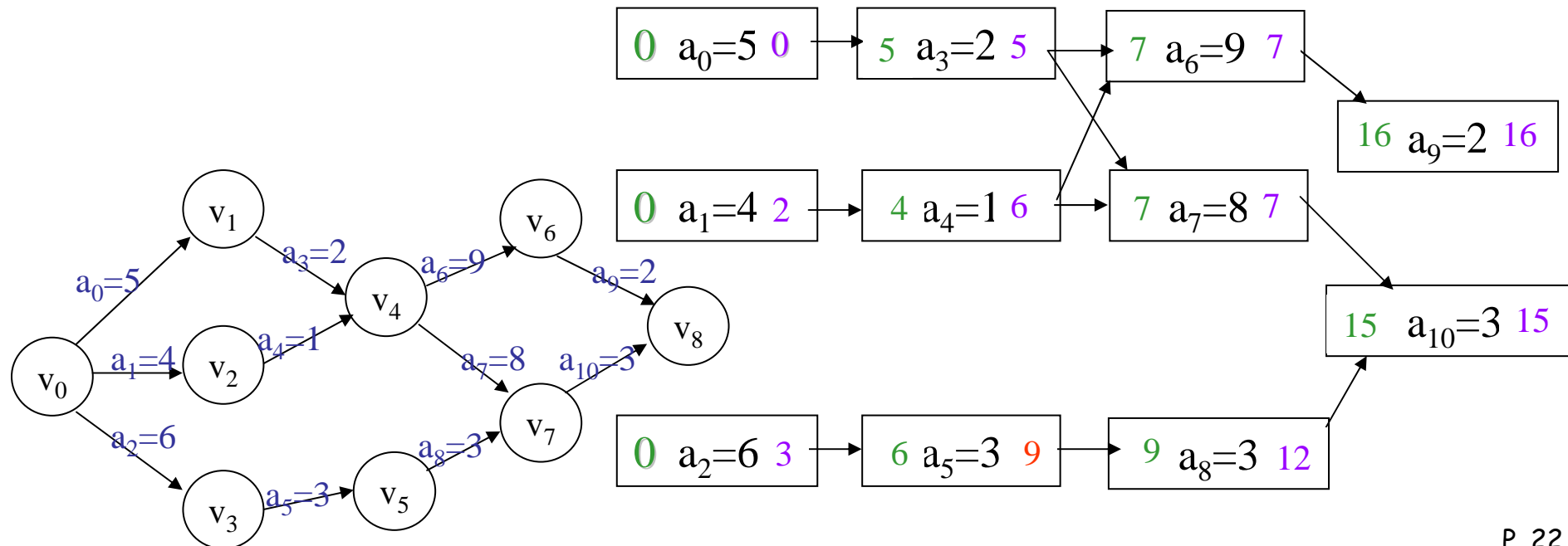
[2]: 3	[0]: 3
[3]: 7-2=5	[1]: 5
[4]: 7-1=6	[2]: 6
[5]: 9	[3]: 9
[6]: 16-9=7	[4]: 7
[7]: 7	[5]: 12
[8]: 12	[6]: 16
[9]: 16	[7]: 15
[10]: 15	[8]: 18
<u>la</u>	<u>le</u>

v₂
v₁
stack

a₀ v₀ v₁
a₁ v₀ v₂
a₂ v₀ v₃
a₃ v₁ v₄
a₄ v₂ v₄
a₅ v₃ v₅
a₆ v₄ v₆
a₇ v₄ v₇
a₈ v₅ v₇
a₉ v₆ v₈
a₁₀ v₇ v₈

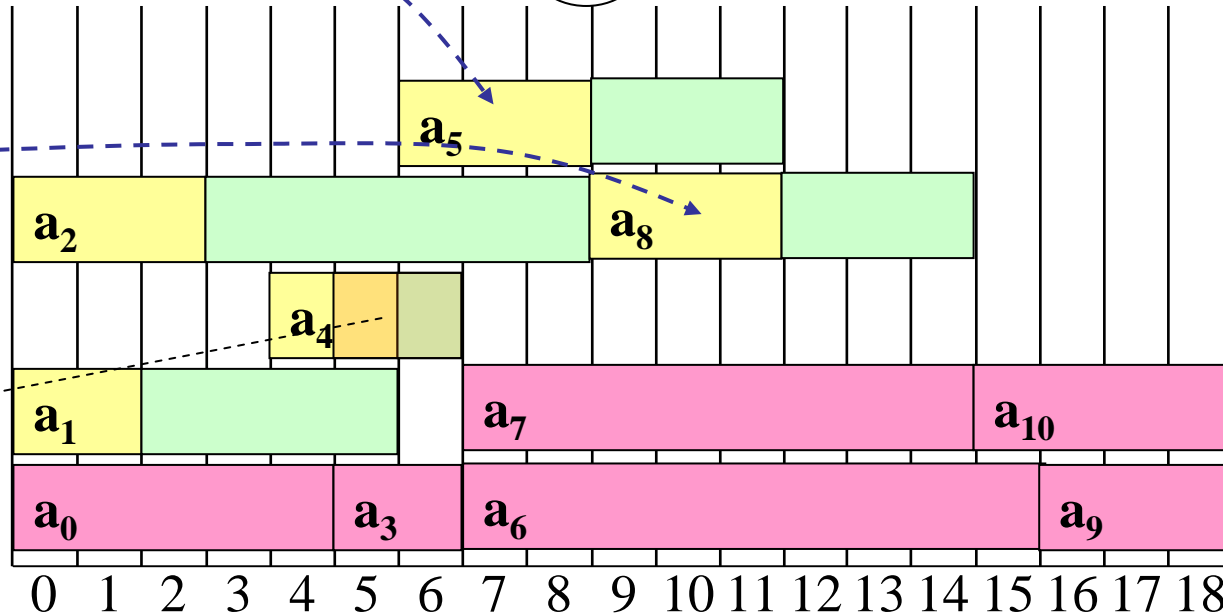
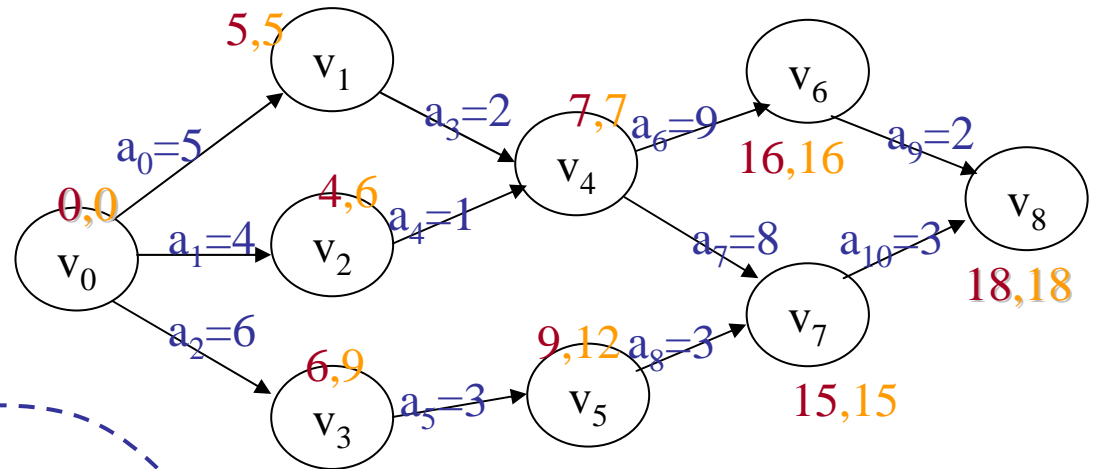
Critical Path Analysis: *Extensions*

- ❑ Critical-path analysis can be carried out with **AOV network**
- ❑ **Free float**: amount of time that a task can be delayed without causing a delay to the *earliest start time* of any *immediately following activities*.
 - earliest finish time & latest finish time for each activity



Critical Path Method: *Gantt Chart*

<u>ea</u>	<u>la</u>	<u>la-ea</u>
[0]: 0	0	0
[1]: 0	2	2
[2]: 0	3	3
[3]: 5	5	0
[4]: 4	6	2
[5]: 6	9	3
[6]: 7	7	0
[7]: 7	7	0
[8]: 9	12	3
[9]: 16	16	0
[10]: 15	15	0



free float =
 $7 - (4 + 1) = 2$

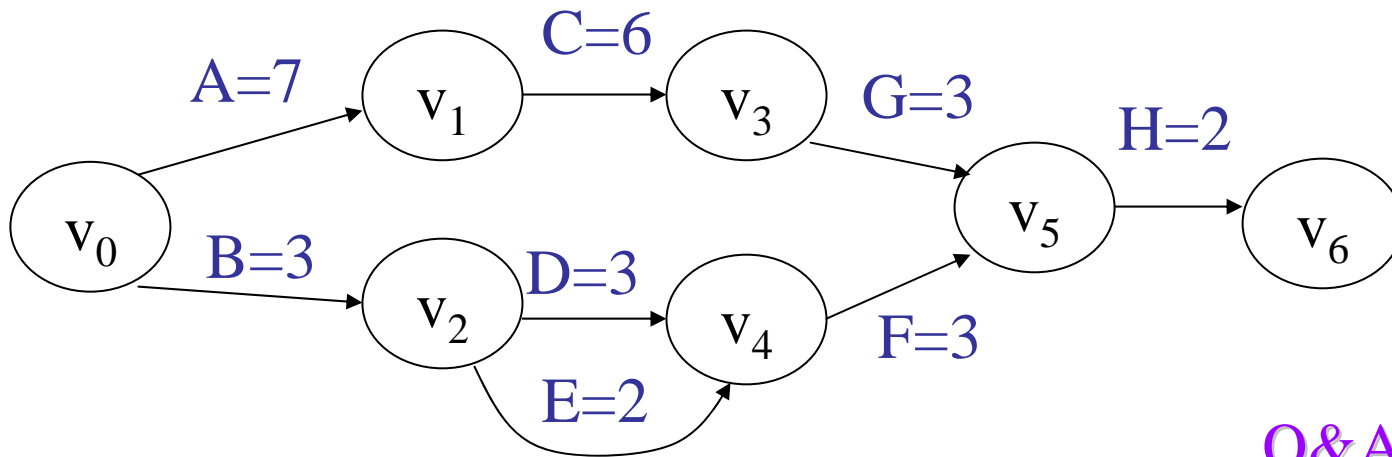
Practice 11: *Critical Path*

□ Find the critical path in the following network

<u>activity</u>	<u>duration</u>	<u>dependencies</u>
A	7	
B	3	
C	6	A
D	3	B
E	2	B
F	3	D, E
G	3	C
H	2	F, G

Practice 11: *Solution*

□ Find the critical path in the following network



Q&A: free float?

Self-exercise 7

1. Consider the activities and durations required to complete a project and dependencies among them.
 - (a) What is the latest time that the activity E can start?
 - (b) What is the minimum time required to complete the project?

- (c) What are the critical activities?

<u>activity</u>	<u>duration</u>	<u>predecessor</u>
A	2	
B	3	
C	6	
D	3	A
E	2	B
F	4	D
G	3	C, E