Graph App.

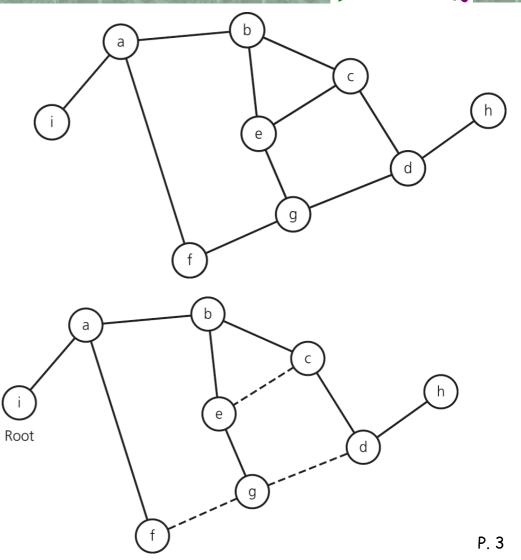
- **□**Topological Sort
- **□**Spanning Tree
- **□Shortest Paths**

Spanning Tree: Definition

- ☐ A tree is an undirected connected graph without cycles (acyclic)
- \square A spanning tree of a connected undirected graph G is
 - A subgraph of G that contains all of G's
 vertices and enough of its edges to form a tree
 - Application example: *communication network*

Spanning Tree: Definition

- ☐ To obtain a spanning tree from a connected undirected graph with cycles
 - Remove edges until there are no cycles



Spanning Tree: Properties

Data Structures

□ Detecting a cycle in an undirected connected graph

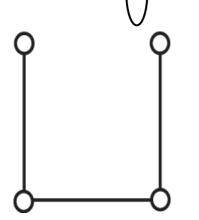
- DFS or BFS?
- A connected undirected graph that has n vertices must have at least n-1 edges
- A connected undirected graph that has n
 vertices and exactly n 1 edges cannot contain
 a cycle
- A connected undirected graph that has n vertices and more than n 1 edges must contain at least one cycle

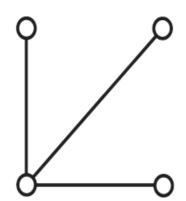
Practice 4: Number of Spanning Trees

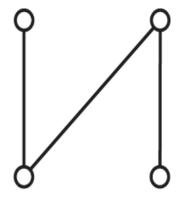
Data Structures

☐ How many different spanning trees?

- Two graphs G and H are isomorphic if and only if there is a bijection f between their vertex sets
 - For any vertices x and y in G, they are adjacent in G iff f(x) and f(y) are adjacent in H



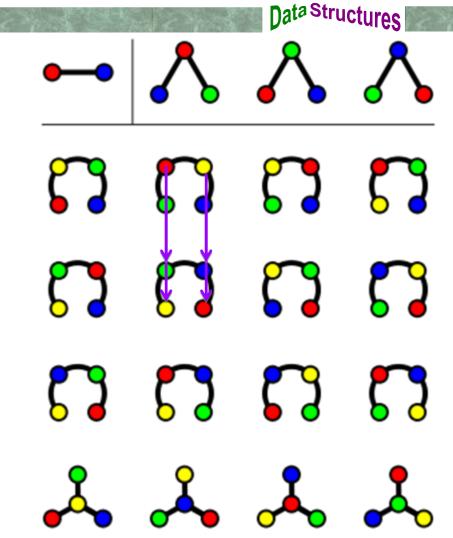




Counting Spanning Trees

Various vertex labeling: nⁿ⁻²

- \square 2 nodes \rightarrow 2²-2=1
- $\square 3 \text{ nodes} \rightarrow 3^{3-2}=3$
- $\Box 4 \text{ nodes} \Rightarrow 4^{4-2} = 16$
- \square Why n^{n-2} ?
 - One proof is based on Prüfer sequence
- ☐ Graph isomorphism
 - One of the NP problems



Counting Spanning Trees

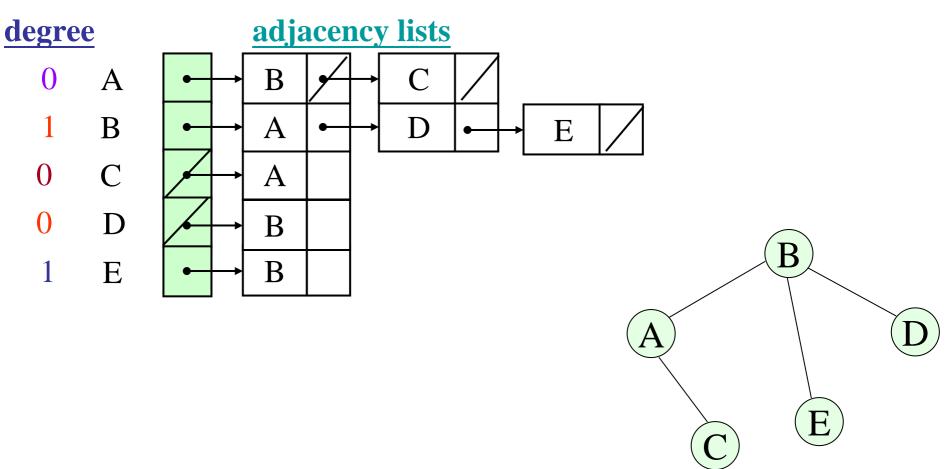
Prüfer sequence [Heinz Prüfer, 1918] Each labeled tree with n vertices has a unique Prüfer sequence of length n-2 Conversion algorithms Leaf with the smallest label Keep the label of its parent Each Prüfer sequence of length n-2 has a unique labeled tree with n vertices

Wikipedia ©

Prüfer Sequence

Data Structures | | |

Prüfer sequence: A B B



Prüfer Sequence

Data Structures

□ Rationale behind the proof

- For *n* labels, there are n^{n-2} sequences of length n-2
- Each Prüfer sequence can be converted to a unique labeled tree with n vertices (a spanning tree)
- For a *complete* graph with n vertices, there are n^{n-2} spanning trees

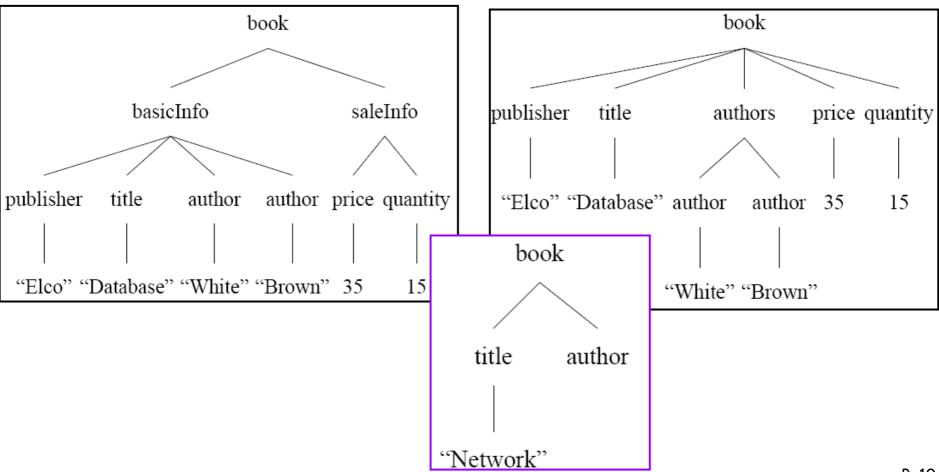
□ Other applications

- Compact tree representation → Prüfer code
- XML structural matching → twig query

Prüfer Sequence

Data Structures

☐ An example of twig query on XML documents



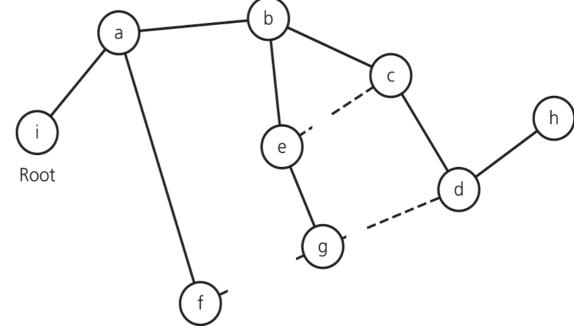
P. 10

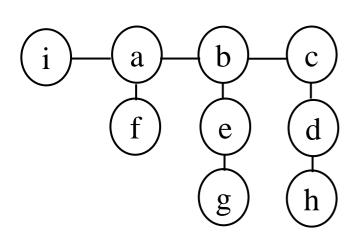
Practice 5: Tree -> Prüfer Sequence

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□ Prüfer sequence?

aebdcba





DFS / BFS for Spanning Trees

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□ To create a spanning tree

- Traverse the graph using either depth-first search (DFS) or breadth-first search (BFS) and mark the edges that you follow
- After the traversal is completed, the graph's vertices and marked edges form a spanning tree

DFS in iterative form (stack)

Nata Structures iterativeDFS(Vertex v) DFS traversal sequence: AB BC BD s.createStack(); count=0; s.push(v);Mark v as visited; while (!s.isEmpty() && count < |V|-1) B u = s.getTop();// at top of the stack **if** (unvisited vertex w is adjacent to u) s.**push**(w); count++; Mark w as visited; // (u,w) // backtrack else **s.pop()**;

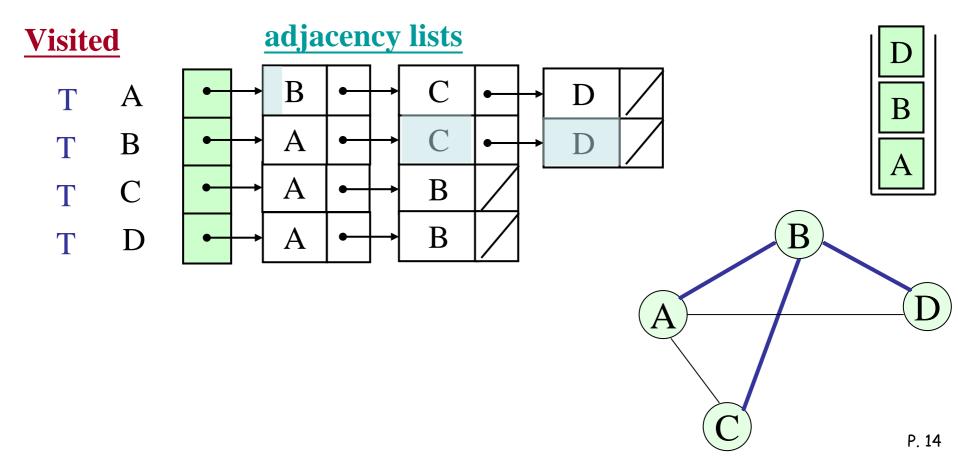
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DFS in iterative form (stack)

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DFS traversal sequence: **AB** BC BD

count = 3



BFS in iterative form (queue)

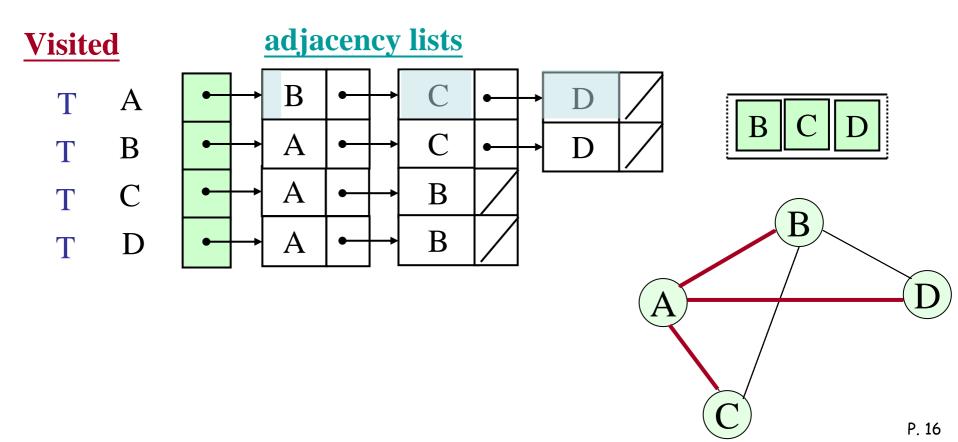
```
BFS traversal sequence: AB AC AD
iterativeBFS(Vertex v)
  q.createQueue(); count=0;
  q.enqueue(v);
  Mark v as visited:
  while (!q.isEmpty() && count < |V|-1)
       q.dequeue(u);
       for (each unvisited vertex w adjacent to u)
             Mark w as visited; // (u,w)
             q.enqueue(w); count++;
                                                                P. 15
```

BFS in iterative form (queue)

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BFS traversal sequence: AB AC AD

count = 3



Self-exercise 4

- 1. Consider the graph to answer the following:
- (a) How many **spanning trees** does it have?
- (b) Draw all the spanning trees
- (c) Compute Prüfer sequence of each spanning tree

