## **Assignment #6**

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**Problem 6.1:** completeness of  $\rightarrow$  and  $\neg$ 

Claim: implication and negation are universal

**Proof:** *Proof by Construction* 

It is enough to express an universal elementary boolean function(NOR/NAND gate) with negation and implication to guarantee that they are universal. Manually check that:

p	q	$p \rightarrow q$	$\neg p$	$\neg q$	$p \downarrow q$	$p \uparrow q$
F	F	T	T	T	T	T
F	T	T	T	F	F	T
T	F	F	F	T	F	T
T	T	T	F	F	F	F

Approach 1: From the table deduce that  $(p \downarrow q) \equiv (\neg((\neg p) \rightarrow q))$ .

Approach 2: From the table deduce that  $(p \uparrow q) \equiv (p \to (\neg q))$ .

**Problem 6.2:** conjunctive and disjunctive normal form

a)  $\varphi(P,Q,R,S) = (\neg P \lor Q) \land (\neg Q \lor R) \land (\neg R \lor S) \land (\neg S \lor P)$ 

P	Q	R	S	$\neg P \lor Q$	$\neg Q \lor R$	$\neg R \lor S$	$\neg S \lor P$	φ
F	F	F	F	T	T	T	T	T
F	F	F	T	T	T	T	F	F
F	F	T	F	T	T	F	T	F
F	F	T	T	T	T	T	F	F
F	T	F	F	T	F	T	T	F
F	T	F	T	T	F	T	F	F
F	T	T	F	T	T	F	T	F
F	T	T	T	T	T	T	F	F
T	F	F	F	F	T	T	T	F
T	F	F	T	F	T	T	T	F
T	F	T	F	F	T	F	T	F
T	F	T	T	F	T	T	T	F
T	T	F	F	T	F	T	T	F
T	T	F	T	T	F	T	T	F
T	T	T	F	T	T	F	T	F
T	T	T	T	T	T	T	T	T

 $\Rightarrow$  Only 2 interpretations satisfy  $\varphi$ .

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```
b) DNF = (P \land Q \land R \land S) \lor (\neg P \land \neg Q \land \neg R \land \neg S)
c) CNF = (\neg P \lor Q) \land (\neg Q \lor R) \land (\neg R \lor S) \land (\neg S \lor P)
       \xrightarrow{\text{distributive law}} ((\neg P \land \neg Q) \lor (\neg P \land R) \lor (\neg Q \land Q) \lor (Q \land \neg R)) \land
                                  \wedge ((\neg R \wedge \neg S) \vee (\neg R \wedge P) \vee (S \wedge \neg S) \vee (S \wedge \neg P))
          \xrightarrow{\text{contradiction}} ((\neg P \land \neg Q) \lor (\neg P \land R) \lor (Q \land \neg R)) \land
                                  \wedge ((\neg R \wedge \neg S) \vee (\neg R \wedge P) \vee (S \wedge \neg P))
       \xrightarrow{\text{distributive law}} ((\neg P \land \neg Q) \land (\neg R \land \neg S)) \lor ((\neg P \land R) \land (\neg R \land \neg S)) \lor
                                  \vee ((O \land R) \land (\neg R \land S)) \lor ((\neg P \land \neg O) \land (\neg R \land P)) \lor
                                  \vee ((\neg P \wedge R) \wedge (\neg R \wedge P)) \vee ((Q \wedge R) \wedge (\neg R \wedge P)) \vee
                                  \vee ((\neg P \wedge \neg Q) \wedge (S \wedge P)) \vee ((\neg P \wedge R) \wedge (S \wedge P)) \vee
                                  \vee ((O \wedge R) \wedge (S \wedge P))
            \xrightarrow{\text{associativity}} (\neg P \land \neg Q \land \neg R \land \neg S) \lor (\neg P \land R \land \neg R \land \neg S) \lor (Q \land R \land \neg R \land S) \lor
                                  \vee (\neg P \wedge \neg Q \wedge \neg R \wedge P) \vee (\neg P \wedge R \wedge \neg R \wedge P) \vee (Q \wedge R \wedge \neg R \wedge P) \vee
                                  \vee (\neg P \wedge \neg Q \wedge S \wedge P) \vee (\neg P \wedge R \wedge S \wedge P) \vee (Q \wedge R \wedge S \wedge P)
           \xrightarrow{\text{contradiction}} (\neg P \land \neg Q \land \neg R \land \neg S) \lor F \lor F \lor F \lor F \lor F \lor F \lor (Q \land R \land S \land P)
                   \xrightarrow{\text{identity}} (\neg P \land \neg Q \land \neg R \land \neg S) \lor (Q \land R \land S \land P)
         \xrightarrow{\text{commutativity}} (P \land Q \land R \land S) \lor (\neg P \land \neg Q \land \neg R \land \neg S)
```

## **Problem 6.3:** boolean expressions (haskell)

```
{-
Module: p6-boolexpr/boolexpr.hs
- }
module BoolExpr (Variable, BoolExpr(..), evaluate) where
import Data.List
type Variable = Char
data BoolExpr
= T
 1 F
 | Var Variable
 | Not BoolExpr
 | And BoolExpr BoolExpr
 | Or BoolExpr BoolExpr
 deriving (Eq, Ord, Show)
 -- evaluates an expression
evaluate :: BoolExpr -> [Variable] -> Bool
evaluate T ts = True
```

```
evaluate F ts = False
evaluate (Var v) ts = elem v ts
evaluate (Not e) ts = not (evaluate e ts)
evaluate (And e1 e2) ts = evaluate e1 ts && evaluate e2 ts
evaluate (Or e1 e2) ts = evaluate e1 ts || evaluate e2 ts
-- 6.3.a.
variables :: BoolExpr -> [Variable]
variables T = "" -- following the example on the sheet
variables F = "" --following the example on the sheet
variables (Var x) = [x]
variables (Not y) = variables y --declare the expression condition
variables (And a b) = sort(variables a `union` variables b)
variables (Or a b) = sort(variables a `union` variables b)
-- 6.3.b.
subsets :: [Variable] -> [[Variable]]
subsets [] = [[]] -- similiar to the prefixes and suffixes
subsets (head:tail) = subsets(tail) ++ map (head:)(subsets(tail))
truthtable :: BoolExpr -> [([Variable], Bool)]
truthtable n = zip (subsets(variables n))(map (evaluate(n)) (subsets(variables n)))
-- the zip relates tuples with elements in the same position for the 2 lists
```